CAQM: Convexity Analysis of Quadratic Maps

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CAQM is a MATLAB library designed to analyze geometric properties of images of quadratic maps (which are manifolds of real-valued solutions of systems of quadratic equations).

This file is aimed at both guiding users with the installation steps and enumerating the methods and techniques which constitute the library. The main functionality described below is contained in several MATLAB functions located in the main folder.

Additional procedures which reside in the library/ folder are described in the accompanying file library.pdf.

Installation and Testing

The library requires MATLAB and CVX as prerequisites. Having obtained those, the following steps are required in order to use CAQM.

1. **Obtaining source code** can be done in two ways. First one requires basic knowledge of the GIT version control system. Simply clone the repository:

git clone git@github.com:sergeivolodin/CertificateCutting.git

The second option is to download the latest snapshot of the repository in form of a .zip archive. To do so, navigate to github.com/sergeivolodin/CertificateCutting, press the green "Clone or download" button and then choose "Download ZIP" in the drop-down menu. Then unpack the obtained archive. Both cases should give a folder called CertificateCutting (or CertificateCutting-master in the latter case). This folder must contain a file README.md.

- 2. Configuring MATLAB involves adding the folder to PATH. The following directories (assuming that the root folder is called CertificateCutting) should be added: CertificateCutting, CertificateCutting/library, CertificateCutting/tests. MATLAB gives two ways of adding a folder to MATLAB PATH. The first one involves console commands and the second one uses GUI.
- 3. Testing the library is an essential step which ensures that the library and all of its dependencies were installed correctly. To run the test, open MATLAB, navigate to CertificateCutting/tests and run the CAQMtest.m file. This file will go through essential parts of the functionality. During that period, a parallel cluster might be started (depending on the MATLAB version) and several GUI dialog boxes might be displayed. If the test succeeded, a message TEST PASSED should be displayed in the console output. Otherwise, if the test has failed, a YouTube video demonstrating normal test operations might prove helpful in diagnosing the issue and overcoming installation mistakes. In other cases, feel free to contact the authors.

After installation and testing, the library is ready to use in your applications. The sections below describe the functionality in details.

Notations

1. Real case, the map $f: \mathbb{R}^n \to \mathbb{R}^m$

$$f_k(x) = x^T A_k x + 2b_k^T x, \quad A_k = A_k^T, \quad x, b_k \in \mathbb{R}^n, \quad i = k \dots m.$$
 (1)

2. Complex case, the map $f: \mathbb{C}^n \to \mathbb{R}^m$

$$f_k(x) = x^* A_k x + b_i^* x + x^* b_k, \quad A_k = A_k^*, \quad x, b_k \in \mathbb{C}^n, \quad i = k \dots m,$$
 (2)

where \cdot^* is Hermitian conjugate.

We will use \mathbb{V} to denote \mathbb{R}^n in the real case and \mathbb{C}^n in the complex case. We also use the following notations:

Notation 1. For a vector $c = (c_1, ..., c_m)$ and a tuple of vectors $b = (b_1, ..., b_m)$, $b_k \in \mathbb{V}$, or a tuple of $n \times n$ matrices $A = (A_1, ..., A_m)$, $A_k \in \mathbb{V}^2$, the dot product is defined as follows,

$$c \cdot b = \sum_{k=1}^{m} c_k b_k, \qquad c \cdot A = \sum_{k=1}^{m} c_k A_k.$$

Notation 2. The image of f is denoted as F,

$$F = f(\mathbb{V}) \subset \mathbb{R}^m$$
.

Notation 3. The convex hull of F is denoted as G:

$$G = \operatorname{conv}(F) \subset \mathbb{R}^{\mathrm{m}}$$
.

Notation 4. The boundary points of F touched by a supporting hyperplane with the normal vector $c \in \mathbb{R}^m$,

$$\partial F_c = \operatorname*{arg\,min}_{y \in F} (c \cdot y)$$

Notation 5. The boundary points of G touched by a supporting hyperplane with the normal vector $c \in \mathbb{R}^m$,

$$\partial G_c = \operatorname*{arg\,min}_{y \in G} (c \cdot y)$$

Notation 6. Set of normal vectors c, such that ∂F_c is nonconvex is denoted as C_- ,

$$C_{-} = \{c \in \mathbb{R}^m \mid \text{Set } \partial F_c \text{ is nonconvex}\}$$

Functionality

The library consists of several functions, each of them is defined in a separate .m file. Input format for specify the quadratic map is as follows.

- The array A(i,j,k) denotes i'th row and j'th column of a $n \times n$ matrix $A_k \in \mathbb{V}^2$
- The array b(i,k) denotes i'th element of a vector $b_k \in \mathbb{V}$
- 1. Feasibility membership oracle, infeasibility_oracle.m

```
is_infeasible = infeasibility_oracle(A, b, y)
```

Input:

- the map f specified by matrices A_k and vectors b_k ,
- point $y \in \mathbb{R}^m$.

Output: determines if $y \in G$, returns is_infeasible=1 if $y \notin G$, is_infeasible=0 if $y \in G$.

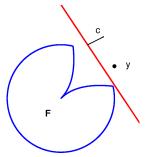


Figure 1: Infeasibility oracle: hyperplane orthogonal to c separates the point y from the convex hull G of F

This function attempts to certify that the point y does not belong to the convex hull G of F by separating y and G by an appropriate hyperplane. This is illustrated in Fig. 1, see Theorem 3.2 in the accompanying article for details. The function returns $is_infeasible=1$ if the desired hyperplane was found. In this case $y \notin G$ and consequently $y \notin F$, implying there is no $x \in \mathbb{V}$ such that y = f(x), i.e. this point is infeasible. If the hyperplane was not found the function returns $is_infeasible=0$, which means the feasibility of y with respect to F is uncertain. We note that feasibility of y is exactly the same as solvability of the system of quadratic equations y = f(x) with real-valued x when the map f is real, or with x and x^* being complex conjugate to each other if f is complex as defined in eq. (2).

2. Boundary oracle boundary_oracle.m

```
[t, is_in_F] = boundary_oracle(A, b, y, d)
```

Input:

- the map f specified by matrices A_k and vectors b_k ,
- a point $y \in G$,
- a direction $d \in \mathbb{R}^m$.

Output: finds and returns distance t to the boundary of G from the point y inside G in the direction d; verifies if boundary point belongs to F.

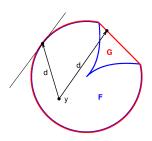


Figure 2: Boundary oracle: distance from y to the boundary ∂G in the direction d. The boundary point of G may or may not belong to ∂F (cases of $d = d_1$ and $d = d_2$ respectively).

This function finds point y+td on the boundary ∂G with the largest $t=\sup\{\tau\big|y+\tau d\in G\}$ and checks if this point belongs to F. This is illustrated in Fig. 2, see Theorem ??!! in the accompanying article for details. The function returns t with the value of t, variable $is_in_F=1$ if the boundary point y+td belongs to F, and variable $is_in_F=0$ if feasibility of y+td with respect to F is uncertain.

Exception: if on input $y \notin G$ or ∂G is not smooth at the boundary point $y + t d \in \partial G$ the function produces an exception.

3. "Non-convex direction" get_c_minus.m

$$c = get_c_minus(A, b, [y], [k])$$

Input:

- the map f specified by matrices A_k and vectors b_k ,
- [optional] a point $y \in G$,
- [optional] number of iterations k.

Output: finds and returns vector c such that ∂F_c is nonconvex through a stochastic algorithm using up to k iterations; returns empty vector if such c was not found.

Exception: really??!!

This function consequently generates up to k random directions d and for each one finds vector c normal to ∂G at the boundary point $y + t d \in \partial G$. Next it finds $\partial_c F$, the intersection of F with the supporting hyperplane orthogonal to c and checks if it is nonconvex. We note that non-convexity of $\partial_c F$ implies non-convexity of F. This function stops and returns c if non-convexity of $\partial_c F$ was established during one of the iterations, and empty vector $\mathbf{c} = [\]$ otherwise. If g and g are not specified, the function uses default values g and g are not specified.

4. Nonconvexity certificate nonconvexity_certificate.m

```
is_nonconvex = nonconvexity_certificate(A, b, [y], [k])
```

Input:

- the map f specified by matrices A_k and vectors b_k ,
- [optional] a point $y \in G$,
- [optional] number of iterations k.

Output: attempts to establish if F is convex, returns $is_nonconvex=1$ if F is non-convex, $is_nonconvex=0$ if uncertain.

This function calls get_c_minus and returns is_nonconvex=1 if the latter returns a non-trivial c.

5. Positive-definite $c \cdot A$ get_c_plus.m

Input:

- matrices A_k
- \bullet [optional] number of iterations k

Output: finds and returns vector c_+ such that $c_+ \cdot A > 0$.

Exception: if c_+ was not found during k random iterations.

This function generates a random "seed" which is used to find c_+ such that $c_+ \cdot A > 0$. If successful, the function terminates and returns c_+ on the exit, otherwise the search attempt is repeated up to k times. If not specified explicitly, the default value of k = 10. If c_+ is not found during k iterations the function produces an exception.

6. Convex subpart get_z_max.m

z_max = get_z_max(A, b, c_plus, [z_max_guess], [k])

Input:

- the map f specified by matrices A_k and vectors b_k ,
- vector c_+ such that $c_+ \cdot A \succeq 0$,
- [optional] the guess value $z_{\text{max}}^{\text{guess}}$
- [optional] number of iterations k.

Output: finds and returns maximal value z_{max} such that the compact part of F "cut" by the hyperplane $c_+ \cdot (y - y_0) = z_{\text{max}}$, where $y_0 \in \partial_{c_+} F$, is still convex.

Exception: produces an exception if non-convexity of F confined within the half-plane $c_+ \cdot (y - y_0) \le z_{\text{max}}^{\text{guess}}$ has not been established. ??!!

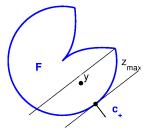


Figure 3: Maximal value of z_{max} such that the subpart of F, $F_{z_{\text{max}}} \equiv \{y | y \in F, c_+ \cdot (y - y_0) \le z_{\text{max}}\}$, which is "cut" from F by a hyperplane orthogonal to c_+ is still convex.

This function returns maximal value z_{\max} such that the hyperplane perpendicular to c_+ and located distance z_{\max} away from the boundary of F is still convex. More precisely, the compact part of F confined by the half-space $\{y \mid c_+ \cdot (y-y_0) \leq z_{\max}^{\text{guess}}\}$ is convex. Here $y_0 \in \partial_{c_+} F$, and since $c_+ \cdot A \succ 0$, the set $\partial_{c_+} F$ consists of only one point y_0 . The geometric meaning of z_{\max} is illustrated in Fig. fig:four. The function first tries to identify "non-convex directions" c_- using get_c_minus and then "follow" each non-convexity to the smallest value of z. This is described in detail in the section ??!! of the accompanying paper. If no "non-convex directions" found, the function returns z_{\max}^{guess} assuming the part of F confined within the half-space $\{y \mid c_+ \cdot (y-y_0) \leq z_{\max}^{\text{guess}}\}$ is convex. In case the input value of c_+ does not satisfy $c_+ \cdot A \succ 0$, the function produced as exception ??? If the maximal number of iterations k (to be used with get_c_minus) is not specified on input, the default value k=10 (or k=100?) is used by default. The guess value z_{\max}^{guess} is supposed to be substantially large to detect non-convexity of F. If it is not specified on input, a default value of $z_{\max}^{\text{guess}} = 10\text{Tr}(c_+ \cdot A)$ is used. It is important to keep in mind that the algorithm is heuristic. A non-trivial return value $z_{\max} \neq z_{\max}^{\text{guess}}$ does not guarantee convexity of $F_{z_{\max}} \equiv \{y \mid y \in F, c_+ \cdot (y-y_0) \leq z_{\max}\}$, but only that $\{y \mid y \in F, c_+ \cdot (y-y_0) \leq z\}$ for any $z > z_{\max}$ is non-convex. Nevertheless, increasing k would increase certainty (in the probabilistic sense) that $F_{z_{\max}}$ is indeed convex.