

On the feasibility for the system of quadratic equations.

Minimizing $z(c)$ when $|c_{\text{bad}}| = \infty$

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1 The framework

In this section, the framework for finding convex cuts is described. We start from the definition of a quadratic map, then the task is set formally and then is formulated as a constrained optimization problem.

Let $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$ be a quadratic map: $f_i(x) = x^T A_i x + 2b_i^T x$, $A_i^T = A_i$. We want to explore the image $F = f(\mathbb{R}^n)$ and find convex parts of F .

1. Consider a problem

$$g(c) = \min_{x \in \mathbb{R}^n} (c, f(x)) \quad (1)$$

It has a finite solution iff 3.2

$$\begin{cases} c \cdot A \geq 0 \\ (c \cdot b)^T \text{Ker}(c \cdot A) = 0 \end{cases}$$

The x , minimizing 1 are (Q^g is a pseudoinverse of Q)

$$x = \text{Ker}(c \cdot A) - (c \cdot A)^g (c \cdot b)$$

This way, we found pre-images $X(c)$ of boundary points of F lying on the supporting hyperplane with a normal vector c :

$$X(c) = \arg \min (c, f(x)) = \text{Ker}(c \cdot A) - (c \cdot A)^g (c \cdot b)$$

2. Let c_+ be a vector from \mathbb{R}^m s.t. $c_+ \cdot A > 0$. Then the supporting hyperplane with the normal vector c_+ is touching F at a single point $y_0 = f(x_0)$, where

$$x_0 = -(c_+ \cdot A)^{-1} (c_+ \cdot b)$$

We want to find the maximum value z_{\max} such that $Z \subseteq F$ is convex:

$$Z = \{y \mid (c_+, y - y_0) \in [0, z_{\max}]\}$$

To do that, we find c such that set $f(X(c))$ is nonconvex, project points of $f(X(c))$ to c_+ and find a minimum of the projection:

$$z_{\max} = \inf_c \tilde{z}(c), \quad \tilde{z}(c) = \inf (c_+, f(X(c)) - f(x_0)) \quad (2)$$

3. To find $\tilde{z}(c) = \inf_{x \in \text{Ker}(c \cdot A) - (c \cdot A)^g (c \cdot b)} (c_+, f(x) - f(x_0))$, we introduce the plus norm in the $q \in \mathbb{R}^n$ space

$$\|q\|_+^2 = q^T (c_+ \cdot A) q$$

Note that 3.4 $(c_+, f(x) - f(x_0)) = \|x - x_0\|_+^2$. Therefore, $\tilde{z}(c) = \rho_+^2(\text{Ker}(c \cdot A) - (c \cdot A)^g (c \cdot b), x_0)$

4. Next, we change variables to find $\tilde{z}(c)$ explicitly. We use the following change $\begin{cases} x = S\hat{x} + x^0 \\ y = \hat{y} + y^0 \end{cases}$

$y = f(x)$, $\hat{y} = \hat{f}(\hat{x})$ and demand $\begin{cases} c_+ \cdot A = I \\ c_+ \cdot b = 0 \end{cases}$. It leads 3.5 to

$$x^0 = x_0, y^0 = y_0, S = S_1 S_2, S_1^T (c_+ \cdot A) S_1 = \Lambda, S_1^T S_1 = I, S_2 = \Lambda^{-1/2}, \begin{cases} \hat{A}_i = S^T A_i S \\ \hat{b}_i = S^T (b_i + A_i x^0) \end{cases}$$

In new variables $\|\hat{q}\|_+^2 = \hat{q}^T \hat{q}$, i.e. the new basis is orthonormal w.r.t. $\|\cdot\|_+$, because $c_+ \cdot A = I$. Notice that $\tilde{z}(c) = \rho_+^2(x_0, \text{Ker}(c \cdot A) - (c \cdot A)^g(c \cdot b)) = \text{diam}(\text{Ker}(c \cdot \hat{A}) - (c \cdot \hat{A})^g(c \cdot \hat{b})) = \|(c \cdot A)^g(c \cdot b)\|^2$ since $\text{Ker}(c \cdot \hat{A})^T((c \cdot A)^g(c \cdot b)) = 0$ 3.2

5. We associate the condition $\text{Rg}(c \cdot A) \leq n - 1$, $c \cdot A \geq 0$ with nonconvexity of $f(X(c))$. The real sufficient condition for nonconvexity is given by the Theorem 3.4, which requires some additional constraints on c :

$$\begin{cases} \text{Rg}(c \cdot A) = n - 1 \\ f^1 \nparallel f^2 \end{cases}$$

But we assume that other cases, in particular, $\text{Rg}(c \cdot A) < n - 1$ or $f^1 \parallel f^2$ are rare.

6. Let us note that if $f(X(c))$ is nonconvex, then $\lambda_{\min}(c \cdot A) = 0$. Consider $\tilde{z}(c + \alpha c_+) = \|(c \cdot A + \alpha I)^g(c \cdot b)\|^2$, since $c_+ \cdot \hat{A} = I$, $c_+ \cdot \hat{b} = 0$. Therefore, adding αc_+ to c adjusts the spectrum of $c \cdot A$. If and only if we choose $\alpha = -\lambda_{\min}(c \cdot A)$, \tilde{z} will be associated with a $\tilde{c} = c + \alpha c_+$ for which $\lambda_{\min}(\tilde{c} \cdot A) = 0$. This way, for every given $c \in \mathbb{R}^m$ we can find one and only one $\tilde{c} = c + \alpha c_+$ such that $\lambda_{\min}(\tilde{c} \cdot A) = 0$

We define $z(c) = \|((c \cdot \hat{A}) - \lambda_{\min}(c \cdot \hat{A}))(c \cdot \hat{b})\|^2$, which is the same as $\tilde{z}(c - \lambda_{\min} c_+)$. This way, we automatically choose the adjustment to the spectrum to ensure that $\lambda_{\min}(c \cdot A - \lambda_{\min}(c \cdot A)) = 0$

Notice that $z(c + \mu c_+) = z(c)$.

7. Next, we define $Q(c) = c \cdot \hat{A} - \lambda_{\min}(c \cdot \hat{A})$, $\hat{b}_c = c \cdot \hat{b}$, $v(c) = Q(c)^g b_c$. Then

$$z(c) = \|v(c)\|^2 = v^T v$$

Note that $z(c)$ has its geometrical meaning 2 only if $c \in c_{\text{bad}} = \{c \mid \text{Ker } Q(c) \perp b_c\}$.

At this point, we have a function $z(c)$ defined on c_{bad} and we want to minimize it. We assume that it is continuous and differentiable, and also that c_{bad} contains connected parts (manifolds)

2 Minimization

In this section we address the problem of finding minimum of $z(c)$ stated in the previous section. Again, we assume that c_{bad} contains a manifold and we have a point c in it already. We use well-known technique, Gradient projection method to solve it. We use not the Euclidean projection, but a very special one based on the geometry of c_{bad} . The method's two steps, gradient step and projection, are discussed in details below.

Formally, we want to find

$$z_{\max} = \inf_{c \in c_{\text{bad}}} z(c) \quad (3)$$

Let $\pi_{c_{\text{bad}}}$ be an operator which projects a point onto c_{bad} . Then the method is the following:

$$c^{(k+1)} = \underbrace{\pi_{c_{\text{bad}}}}_{\text{proj.}} \underbrace{(c^{(k)} - \beta^k \nabla z(c^{(k)}))}_{\text{gradient step}}$$

Nonconvexity certificate gives a single point $c \in c_{\text{bad}}$, which is used as a start point $c^{(1)}$. We assume $\text{Rg } Q(c) = n - 1$

Define $k(c) \in \mathbb{R}^n$ s.t. $k(c) \in \text{Ker } Q(c)$, $\|k(c)\| = 1$. We assume that for any c' the dot product $k(c')^T k(c) > 0$ (change sign if false). Each iteration has two key elements: gradient step and projection (see image below)

Assume that we have done $k - 1$ steps already and obtained a point $c = c^{(k)}$. The next two sections describe how to obtain the next point $\tilde{c} = c^{(k+1)}$.

2.1 Gradient step

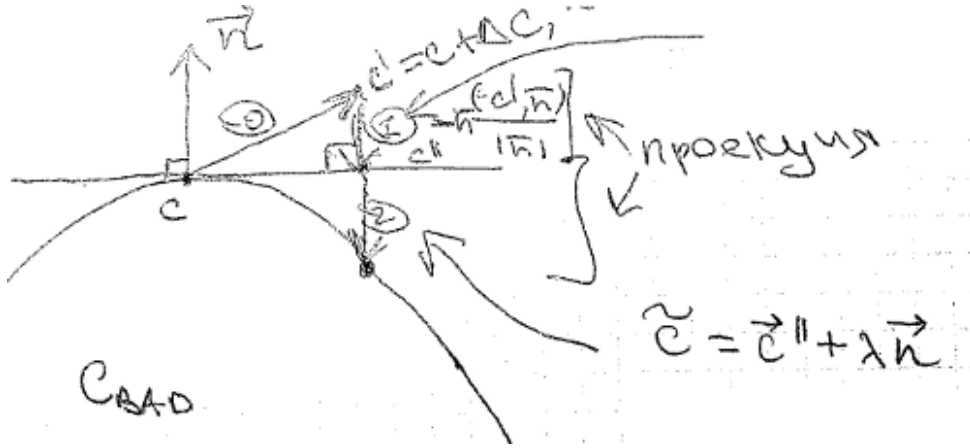
Define $R_i = \hat{A}_i - k^T \hat{A}_i k$, $q_i = \hat{b}_i - R_i v$. Then 3.9

$$\nabla_i z \equiv \frac{\partial z}{\partial c_i} = 2v^T Q^g q_i$$

Define $n_i = k^T q_i$. Then n is a normal vector to c_{bad} at a point c (see picture). We eliminate normal part of the gradient to decrease the distance to c_{bad} (see picture):

$$c' = c - \beta \left(I - n \frac{(\cdot, n)}{|n|^2} \right) \nabla z$$

2.2 Projection



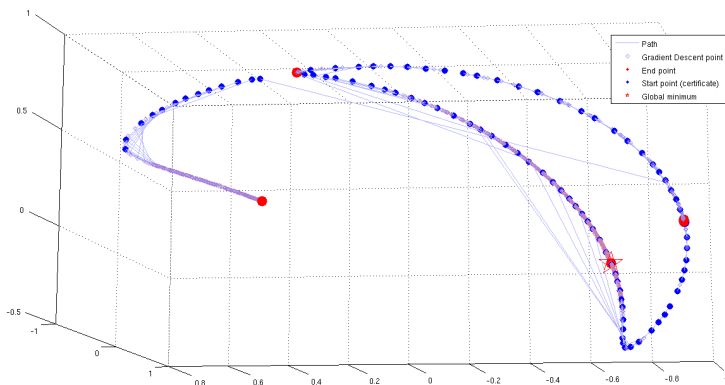
We next assume that for some $\lambda \in \mathbb{R}$ vector $\tilde{c}(\lambda) = c' + \lambda n \in c_{\text{bad}}$. To find λ , we write the condition $\tilde{c} \in c_{\text{bad}} \Leftrightarrow (\tilde{c} \cdot \tilde{b})^T k(\tilde{c}(\lambda)) = 0$. We define a function

$$m(\lambda) = (\tilde{c}(\lambda) \cdot \tilde{b})^T k(\tilde{c}(\lambda))$$

This function is continuous if $\text{Rg } Q(\lambda) = n - 1$ in the neighbourhood of λ . Next, we find its root using bisection method on $[-\lambda_0, \lambda_0]$, $\lambda_0 = \|c - c'\|$.

For some λ , $m(\lambda) = 0$ means that $\tilde{c} = c' + \lambda n \in c_{\text{bad}}$, and that the projection step was a success.

If the method does not converge on the interval given, or rank problem ($\text{Rg } Q \neq n - 1$) occurs, we reduce the gradient step $\beta \rightarrow \theta \beta$, $\theta < 1$, recalculate c' and try the projection again.



This way, we can construct a new point \tilde{c} from the previous one c , and \tilde{c} has lower value of $z(\cdot)$. We continue until the Gradient projection method condition holds:

$$\nabla z \|n$$

If this condition holds, then the resulting \tilde{c} is the same as c on previous iteration, and the iterations stop.

3 Explanations

This section contains explanations and proofs for the main article and draft.pdf.

3.1 F and $\text{conv } F$

Consider $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, s.t. $f_i(x) = x^T A_i x + 2b_i^T x$, $A_i = A_i^T$. Define $F = f(\mathbb{R}^n)$. Then why $A = \inf_{y \in F} (c, y) = \inf_{y \in \text{conv } F} (c, y) = B$?

1. First, $F \subseteq \text{conv } F$, therefore, $B \leq A$.

2. Secondly, let $y_k \in \text{conv } F$ be a sequence s.t. $g_k = (c, y_k) \rightarrow B$. $y_k = \sum_{i=1}^{n_k} \alpha_i^k y_i^k$.

$g_k(c, y_k) = \sum_{i=1}^{n_k} \alpha_i^k (c, y_i^k) = \sum_{i=1}^{n_k} \alpha_i^k g_i^k$. Define $g_0^k = \min_{i \in 1, n_k} g_i^k$. Then $B \leq g_0^k \leq g^k$. Therefore, $g_0^k \rightarrow B$ also. This way, we have constructed a sequence $y_0^k \in F$ s.t. $(c, y_0^k) \rightarrow B$, therefore, $A \leq B$.

3.2 Minimum of $f(x)$

Consider $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$. $f_i(x) = x^T A_i x + 2b_i^T x$. $A_i^T = A_i$. Let $c \in \mathbb{R}^m$

We want to find $g(c) = \inf_{x \in \mathbb{R}^n} (c, f(x))$.

Define $A_c \equiv c \cdot A = \sum_{i=1}^m c_i A_i$, $b_c = c \cdot b = \sum_{i=1}^m c_i b_i$.

$(c, f(x)) = \sum_{i=1}^m c_i f_i(x) = \sum_{i=1}^m c_i (x^T A_i x + 2b_i^T x) = x^T A_c x + 2b_c^T x$.

If $\exists v: -\alpha = v^T A_c v < 0$ then $g(c) = -\infty$: $g(\beta v) = -\beta^2 \alpha + \beta 2b_c^T v \rightarrow -\infty$, $\beta \rightarrow +\infty$.

From this point on, we assume $A_c \geq 0$. Let R_0 be a zero eigenspace (=kernel) of A_c : $R_0 = \{v: A_c v = 0\} = \ker A_c$

If $\exists v \in R_0: v^T b_c \neq 0$ then $g(c) = -\infty$: Consider $f(\beta v) = \beta^2 v^T \overset{0}{A_c v} + 2\beta \underbrace{b_c^T v}_{\neq 0} \rightarrow -\infty$, $\beta \rightarrow \infty$

Then $R_0 \subseteq \{b_c\}^\perp$

Consider $A = \sum_{i=1}^n \lambda_i s_i s_i^T = S \Lambda S^T$, $S = \|s_1 \dots s_n\|$, $S^T S = E$, $s_i^T s_j = \delta_{ij}$.

f is differentiable, then for finding $g(c)$ the gradiend $\nabla(c, f(x)) = 2A_c x + 2b_c = 0$.

$$S \Lambda S^T x = -b_c \Leftrightarrow \Lambda S^T x = -S^T b_c (*)$$

Let x be $x = x^\parallel + x^\perp$, $x^\parallel \in R_0$, $x^\perp \perp R_0$.

Then neither $f(x)$ nor $\Lambda S^T x$ depend on x^\parallel . This means that the x minimizing $g(c)$ is defined in terms of x^\perp and x^\parallel is arbitrary.

Define $\lambda_i^g = \begin{cases} 0, & \lambda_i = 0 \\ 1/\lambda_i, & \lambda_i \neq 0 \end{cases}$. Define $\Lambda^g = \text{diag}(\lambda_1^g, \dots, \lambda_n^g)$. Then $\Lambda \Lambda^g = \delta_{ij} [\lambda_i \neq 0]$. Then $S \Lambda^g \Lambda S^T$ is

a projector on R_0^\perp .

Consider $(*) \Leftrightarrow \Lambda^g \Lambda S^T (x^\parallel + x^\perp) = -\Lambda^g S^T b_c$. But $\Lambda S^T x^\parallel = 0$, therefore, $\underbrace{S \Lambda^g \Lambda S^T}_{\text{projector}} x^\perp = -S \Lambda^g S^T b_c$.

But x^\perp is already in R_0^\perp , therefore, $x^\perp = -\underbrace{S \Lambda^g S^T}_{A_c^g} b_c$. Here A_c^g is a pseudoinverse of A_c .

Therefore, $\boxed{x = -A_c^g b_c + x^\parallel}$, where $x^\parallel \in R_0$.

Consider $(c, f(x)) = (c, f(x^\perp)) = b_c^T A_c^g A_c A_c^g b_c - 2b_c^T A_c^g b_c$. Consider $A_c A_c^g b_c = S \Lambda S^T \overset{E}{S^T} \Lambda^g S^T b_c$. Because $R_0 \subseteq \{b_c\}^\perp$, $\Lambda \Lambda^g S^T b_c = S^T b_c$. Therefore, $A_c A_c^g b_c = b_c$. Then $(c, f(x)) = b_c^T A_c^g b_c - 2b_c^T A_c^g b_c = \boxed{-b_c^T A_c^g b_c}$

3.3 Finding c provided d

Let $H: \mathbb{R}^{n+1, n+1} \rightarrow \mathbb{R}^n$ be a map s.t. $H_i(X) = \text{Tr}(H_i X)$,

$$H_i = \left\| \begin{bmatrix} A_i & b_i \\ b_i^T & 0 \end{bmatrix} \right\|^2$$

Consider a boundary point X , which is a solution of (main article, (4)):

$$\boxed{\begin{array}{l} \sup t \\ \left\{ \begin{array}{l} H(X) = y^0 + td \\ X \geq 0 \\ X_{n+1, n+1} = 1 \end{array} \right. \end{array}}$$

Define $f(t, X) = t$, $D_0 = \{(t, X) | X \geq 0, X_{n+1, n+1} = 1\}$, $D_1 = \{(t, X) | H(X) = y^0 + td\}$. Then supremum is equivalent to

$$\sup_{(t, X) \in D_0 \cap D_1} f(t, X)$$

Define a Lagrange function $L(c, t, X) = \underbrace{t}_{f(t, X)} + \sum_{i=1}^m c_i(y_i^0 + td_i - H_i(X))$.

Here we divided the constraints into two parts: D_1 goes to the Lagrange function, D_0 goes to the inner supremum. *Stephen Boyd, Lieven Vandenberghe. Convex Optimization. Page ??? . Cambridge University Press*

Then the dual function is $g(c) = \sup_{(t, X) \in D_0} L(c, t, X)$.

Because $L = t(1 + \sum_{i=1}^m c_i d_i) + \sum_{i=1}^n c_i(y_i^0 - H_i(X))$, $g = +\infty$ when $(c, d) \neq -1$. From this point we assume that $\boxed{(c, d) = -1}$.

Now, $g(c) = \sup_{X_{n+1, n+1}=1, X \geq 0} (c, y^0 - H(X)) = (c, y^0) + \sup_{y \in \text{conv } F} -(c, y) = (c, y^0) - \inf_{y \in \text{conv } F} (c, y)$.

Then the dual problem is

$$g(c) \rightarrow \inf_{(c, d) = -1}$$

Let us prove that $\inf_{y \in \text{conv } F} (c, y) = \inf_{H = \left\| \begin{bmatrix} A_c & b_c \\ b_c^T & \gamma \end{bmatrix} \geq 0} (-\gamma)$

Via Schur complement $H \geq 0 \Leftrightarrow \begin{cases} A_c \geq 0 \\ \gamma - b_c^T A_c^{-1} b_c \geq 0 \\ (E - A_c A_c^{-1}) b_c = 0 \end{cases}$.

$A_c \geq 0$ is a necessary condition for $\exists g(c) \in \mathbb{R}$ (see part 2).

$(E - A_c A_c^{-1}) b_c = 0$ is another necessary condition for $\exists g(c) \in \mathbb{R}$.

Statement $\gamma \geq b_c^T A_c^{-1} b_c$ means $-\gamma \leq -b_c^T A_c^{-1} b_c = \inf_{y \in \text{conv } F} (c, y)$, which means that $-\gamma$ is a lower bound

for $\inf_{y \in \text{conv } F} (c, y)$.

Then $H \geq 0 \Leftrightarrow -\gamma \leq \inf_{y \in \text{conv } F} (c, y)$.

Then $g(c) = (c, y^0) - \inf_{H \geq 0} -\gamma$.

Then the dual problem is:

$$\inf_{(c, d) = -1} g(c) \Leftrightarrow \inf_{(c, d) = -1} \left[(c, y^0) - \inf_{H \geq 0} (-\gamma) \right] \Leftrightarrow \inf_{(c, d) = -1} \inf_{H \geq 0} (c, y^0) + \gamma = \boxed{\begin{array}{l} \inf \gamma + (c, y^0) \\ H \geq 0 \\ (c, d) = -1 \end{array}}$$

This problem is exactly (5) from main article ■.

3.4 What is z_{\max} ?

Consider $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, $f_i(x) = x^T A_i x + 2b_i^T x$, $A_i = A_i^T$.

Let $c_+ \in \mathbb{R}^m$ s.t. $A_+ = \sum c_i A_i > 0$. Then the minimum $\inf_x (c, f(x))$ is obtained at a single point

$$x_0 = -A_+^{-1} b_+, \quad b_+ = \sum c_i b_i.$$

Consider $S_\varepsilon^+ = \{x \in \mathbb{R}^n \mid (x - x_0)^T A_+ (x - x_0) = \varepsilon^2\}$. Then $f(S_\varepsilon^+) = \{y \in \mathbb{R}^m \mid (c_+, y) = (c_+, f(x_0)) + \varepsilon^2\}$.

Indeed, if $\begin{cases} P = (c_+, f(x) - f(x_0)) \\ Q = (x - x_0)^T A_+ (x - x_0) \end{cases}$ then $P - Q = \underline{x^T A_+ x} - x_0^T A_+ x_0 + 2b_+^T x - 2b_+^T x_0 - \underline{x^T A_+ x} - 2x_0^T A_+ x_0 + 2x^T A_+ x_0 = /x_0 = -A_+^{-1} b_+ / = 2x_0^T b_+ + 2b_+^T x - 2b_+^T x_0 - 2x^T b_+ = 0$.

Therefore, the image of B_ε^+ is a *convex cut* $\{y \mid (c_+, y) \in (c_+, f(x_0)) + [0, z_{\max}]\}$

3.5 Variables s.t. $c \cdot A = I$, $c \cdot b = 0$

Given: the map $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, $f_i(x) = x^T A_i x + 2b_i^T x$, $A_i = A_i^T$, vector $c \in \mathbb{R}^m : c \cdot A > 0$

We need to find a new pair of bases s.t. $\begin{cases} \tilde{c} \cdot \hat{A} = I & (1) \\ \tilde{c} \cdot \hat{b} = 0 & (2) \end{cases}$

We choose $x = S\hat{x} + x^0$, $y = \hat{y} + y^0$, $f(x) = y$, $\hat{f}(\hat{x}) = \hat{y}$. Consider $f_i(x) = x^{0T} A_i x^0 + \hat{x}^T S^T A_i S \hat{x} + 2x^{0T} A_i S \hat{x} + 2b_i^T S \hat{x} + 2b_i^T x^0 = \hat{x}^T \underbrace{S^T A_i S}_{\hat{A}_i} \hat{x} + 2\hat{x}^T \underbrace{S^T (b_i + A_i x^0)}_{\hat{b}_i} + \underbrace{x^{0T} A_i x^0 + 2b_i^T x^0}_{y_i^0}$

$c \cdot \hat{A} = S^T c \cdot A S = I$, then $S = S_1 S_2$, $S_1^T A_i S_1 = \Lambda = \text{diag}$, $S_1^T S_1 = I$, $S_2 = \Lambda^{-1/2}$.

$c \cdot \hat{b} = S^T (c \cdot b + (c \cdot A) x^0)$. Then $x^0 = -(c \cdot A)^{-1} (c \cdot b)$.

The algorithm:

1. Compute S_1 via the eigenbasis of $c \cdot A$, $S_1^T (c \cdot A) S_1 = \Lambda$
2. Compute $S_2 = \Lambda^{-1/2}$, $S = S_1 S_2$.
3. Compute $x^0 = -(c \cdot A)^{-1} (c \cdot b)$
4. Compute $\tilde{A}_i = S^T A_i S$, $\tilde{b}_i = S^T (b_i + A_i x^0)$

3.6 $z(c) = ?$

Given: the map $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, $f_i(x) = x^T A_i x + 2b_i^T x$, $A_i = A_i^T$, two vectors $c, c_+ \in \mathbb{R}^m$. $c_+ A = I$, $c_+ b = 0$. Find: $z(c) = \inf_{y \in Y} (c_+, y)$ where Y is an intersection of $f(\mathbb{R}^n)$ with a tangent hyperplane defined by its normal vector c .

Define $\sigma(Q) = \{\lambda \mid \dim \text{Ker}(Q - \lambda E) > 0\}$. Define $\lambda_{\min}(Q) = \min \sigma(Q)$ — minimal eigenvalue of Q .

1. If $\lambda_{\min}(c \cdot A) < 0$ then the tangent hyperplane does not exist, and $z(c) = +\infty$
2. Then $\lambda_{\min}(c \cdot A) > 0$, then there is no nonconvexity, and $z(c) = \inf$
3. For $\lambda_{\min} = 0$ in part 2 Y was found explicitly: $Y = \{f(x) \mid x = x^\parallel - A_c^g b_c, x^\parallel \in \text{Ker } A_c\}$. Then for $y \in Y$ $(c_+, y) = x^T (c_+ \xrightarrow{I} x) + 2(c_+ \xrightarrow{b} x)^0 = \boxed{x^T x}$
4. $x^T x = \|x\|^2 = \|x^\parallel\|^2 + \|A_c^g b_c\|^2$. We want to minimize (c_+, y) , therefore, we choose $x^\parallel = 0$. Then $z(c) = \|A_c^g b_c\|^2 = \boxed{|(c \cdot A)^g (c \cdot b)|^2}$

5. Consider $\inf_{x \in \mathbb{R}^n} (c_+, f(x)) = -(c_+ \xrightarrow{b} x^0) (A_{c_+} \xrightarrow{I})^{-1} (c_+ \cdot b) = 0$. Therefore, $z(c) = \inf_{y \in Y} (c_+, y) - \inf_{y \in F} (c_+, y)$

Now consider $z(c) = \begin{cases} |(c \cdot A)^g (c \cdot b)|^2, & \lambda_{\min}(c \cdot A) = 0 \\ +\infty, & \text{otherwise} \end{cases}$

Consider $z(c + \alpha c_+) = |(c \cdot A + \alpha c_+ \xrightarrow{I})^g b_c|^2 = |(c \cdot A + \alpha I)^g b_c|^2$. We need α s.t. $\lambda_{\min}(c \cdot A + \alpha I) = 0$, therefore, $\alpha = -\lambda_{\min}(c \cdot A)$.

Define $\hat{z}(c) = |(c \cdot A - \lambda_{\min}(A))^g (c \cdot b)|^2$. Consider $\mathbb{R}^m \ni c = c^\parallel + c^\perp$, $c^\parallel \parallel c_+$, $(c_+, c^\perp) = 0$.

Then $\hat{z}(c) = \hat{z}(c^\perp)$, i.e. \hat{z} does not depend on c^\parallel . It depends only on c^\perp , and c^\parallel is chosen in a way that $\lambda_{\min}(c \cdot A) = 0$.

3.7 Theorem 3.4 (Nonconvexity certificate)

Given.

1. The map $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, $f_i(x) = x^T A_i x + 2b_i^T x$, $A_i = A_i^T$, $m, n \geq 3$. The vector $c \in \mathbb{R}^m$.
2. $A_c \geq 0$, $\dim \text{Ker } A_c = 1$ (=simple zero eigenvalue), $\text{Ker } A_c = \text{Lin}\{e\}$
3. $b_c \perp \text{Ker } A_c$
4. $b \perp \text{Ker } A_c$, $e^0 = -A_c^g b_c$
5. $f(\alpha e + e^0) = f^0 + 2\alpha f^1 + \alpha^2 f^2$, $f^1 \nparallel f^2$

Then $F = \text{Im } f$ is nonconvex.

Consider $\inf_{y \in F} (c, y)$ (part 2). $x = \underbrace{x}_{\alpha e} + \underbrace{(-A_c^g b_c)}_{e^0}$. Then

$$f_i(\alpha e + e^0) = \alpha^2 \underbrace{e^T A_i e}_{f_i^2} + 2\alpha \underbrace{(b_i^T e + e^T A_i e^0)}_{f_i^1} + \underbrace{2b_i^T e^0 + e^{0T} A_i e^0}_{f_i^0}$$

If $f^1 \nparallel f^2$, then $\{f(\alpha e + e^0) | \alpha \in \mathbb{R}\} = F \cap \{y | (c, y) = \inf_{y \in F} (c, y)\}$ is nonconvex. Then F is nonconvex.

3.8 Equations (0.18)-(0.21)

Consider $A(t): n \times n$, $\exists \dot{A}$, $A^T = A$, $A \geq 0$, A has a simple zero eigenvalue: $\forall t A(t)x_0(t) = 0$, $x_0^T x_0 = 0$.

Then $A = S\Lambda S^T$, $S^T S = E$, $A^g = S\Lambda^g S^T$. Define $\lambda_i = \Lambda_{ii}$. $\Lambda_{ii}^g = \begin{cases} \frac{1}{\lambda_i}, & \lambda_i \neq 0 \\ 0, & \lambda_i = 0 \end{cases}$.

Then $AA^g = A^g A = S\Lambda S^T S\Lambda^g S^T = S\Lambda\Lambda^g S^T = \sum_{\lambda_i \neq 0} s_i s_i^T = 1 - x_0 x_0^T$ (0.20)

Consider $\frac{d}{dt} A x_0 = \dot{A} x_0 + A \dot{x}_0$. Multiplying by A^g from the left: $A^g \dot{A} x_0 + A^g A \dot{x}_0 = 0$. Consider $A^g A \dot{x}_0 = (1 - x_0 x_0^T) \dot{x}_0 = \dot{x}_0 - x_0 x_0^T \dot{x}_0$. Since $\|x_0\|^2 = x_0^T x_0 = 1$, $x_0^T \dot{x}_0 = 0$. Then $-A^g \dot{A} x_0 = \dot{x}_0$ (0.19).

Consider $\dot{A} x_0 + A \dot{x}_0 = 0$. Multiplying by x_0^T from the left: $x_0^T \dot{A} x_0 + \cancel{x_0^T A}^0 \dot{x}_0 = 0$. Then $x_0^T \dot{A} x_0 = 0$ (0.18)

Consider $AA^g = 1 - x_0 x_0^T$. Then $\dot{A} A^g + A \dot{A}^g = -\dot{x}_0 x_0^T - x_0 \dot{x}_0^T = A^g \dot{A} x_0 x_0^T + x_0 x_0^T \dot{A} A^g$ (a)

Consider $A^g x_0 = S\Lambda^g S^T x_0 = 0$.

Multiplying (a) by x_0 from the right: $A \dot{A}^g x_0 = A^g \dot{A} x_0$. Multiplying by A^g from the left: $AA^g \dot{A}^g x_0 = A^g A \dot{A}^g x_0$. Then $(1 - x_0 x_0^T) \dot{A}^g x_0 = A^g A \dot{A}^g x_0$. Then $\dot{A}^g x_0 = A^{-2} \dot{A} x_0 + x_0 x_0^T \dot{A}^g x_0$.

Let's multiply (a) by A^g from the left: $A^g \dot{A} A^g + A^g A \dot{A}^g = A^{-2} \dot{A} x_0 x_0^T + \cancel{A^g x_0}^0 x_0^T \dot{A} A^g$.

Consider $A^g A \dot{A}^g = \dot{A}^g - x_0 x_0^T \dot{A}^g = \dot{A}^g - x_0 x_0^T \dot{A} A^{-2} - x_0 x_0^T \dot{A}^g x_0 x_0^T$.

Consider $x_0^T \dot{A}^g x_0 = x_0^T (\dot{S} \Lambda^g S^T + S \dot{\Lambda}^g S^T + S \Lambda^g \dot{S}^T) x_0 = x_0^T \dot{S} \Lambda^g S^T x_0 = 0$.

Then $\frac{d}{dt} A^{-1} = -A^{-1} \dot{A} A^{-1} + x_0 x_0^T \dot{A} A^{-2} + A^{-2} \dot{A} x_0 x_0^T$ (0.21).

Case $\text{Rg } A < n - 1$ is not considered since probability of such event is small.

3.9 Gradient descent

1. (0.5): $\underbrace{(c \cdot A - \lambda_{\min}(c \cdot A))}_{Q} x_0 = 0$, then $\dot{x}_0 = -Q^{-1} \dot{Q} x_0$. $\frac{d}{dt} \lambda_{\min} = \frac{d}{dt} x_0^T (c \cdot A) x_0 = 2 \dot{x}_0 \underbrace{(c \cdot A) x_0}_{\lambda_{\min} x_0} + x_0^T (\dot{c} \cdot A - \lambda_{\min}(c \cdot A)) x_0$

$$A) x_0 = 2\lambda_{\min} \cancel{\dot{x}_0^T x_0}^0 + x_0^T (\dot{c} \cdot A) x_0.$$

Then $\dot{x}_0 = -(A_c - \lambda_{\min}(A_c))^{-1} (\dot{c} \cdot A - x_0^T (\dot{c} \cdot A) x_0) x_0$ (correct).

2. (0.6). $x_0^T (c \cdot b) = 0$, use (0.5) (correct).

3. (0.7). $\frac{\partial}{\partial t} \|v(c)\|^2 = \frac{\partial}{\partial t} \sum_j v_j^2(c) = 2 \sum_j v_j \frac{\partial}{\partial t} v_j = 2v^T(c) \frac{d}{dt} v(c(t))$ (correct). $v(c) = \underbrace{(c \cdot A - \lambda_{\min}(c \cdot A))}_{Q}^{-1} (c \cdot b)$

$$b) = Q^{-1}(c \cdot b).$$

4. (0.8) Define $Q = c \cdot A - \lambda_{\min}(c \cdot A)$. Define $v = Q^{-1}(c \cdot b)$. Then $z(c) = \|v\|^2$ and $\dot{z} = 2v^T \dot{v}$.

$$\dot{v} = \dot{Q}^{-1}(c \cdot b) + Q^{-1}(\dot{c} \cdot b).$$

$$\text{Consider (0.21)} \quad \dot{Q}^{-1} = -Q^{-1}\dot{Q}Q^{-1} + x_0 x_0^T \dot{Q}Q^{-2} + Q^{-2}\dot{Q}x_0 x_0^T.$$

$$\text{Then } \dot{z} = 2v^T \left(Q^{-1}(\dot{c} \cdot b) + (-Q^{-1}\dot{Q}Q^{-1} + x_0 x_0^T \dot{Q}Q^{-2} + Q^{-2}\dot{Q}x_0 x_0^T)(c \cdot b) \right) \equiv$$

$$\equiv 2v^T Q^{-1}(\dot{c} \cdot b) - 2v^T Q^{-1}\dot{Q}Q^{-1}(c \cdot b) + 2v^T x_0^T \dot{Q}Q^{-2}(c \cdot b) + 2v^T Q^{-2}\dot{Q}x_0 x_0^T(c \cdot b) \equiv$$

Since $x_0 \in \text{Ker } Q \perp (c \cdot b)$, we have $x_0^T(c \cdot b) = 0$.

Since $Qx_0 = 0$, $Q^{-1}x_0 = 0$: $Q^{-1}x_0 = S\Lambda^{-1}S^T x_0 = S * 0 = 0$. Since $v^T = (c \cdot b)^T Q^{-1}$. Then $v^T x_0 = 0$

$$\equiv 2v^T Q^{-1}(\dot{c} \cdot b) - 2v^T Q^{-1}\dot{Q} \underbrace{Q^{-1}(c \cdot b)}_v = 2v^T Q^{-1}(\dot{c} \cdot b) - 2v^T Q^{-1}\dot{Q}v = \boxed{\dot{z} = 2v^T Q^{-1}(\dot{c} \cdot b - \dot{Q}v)},$$

$$\dot{Q} = \dot{c} \cdot A - x_0^T(\dot{c} \cdot A)x_0$$

Since $\dot{z} = \sum \frac{\partial z}{\partial c_i} \frac{\partial c_i}{\partial t}$, $\frac{\partial z}{\partial c_i}$ can be found as a coefficient at \dot{c}_i in \dot{z}

$$\dot{z} = 2v^T Q^{-1} \sum_i (\dot{c}_i b_i - \dot{c}_i A_i v + x_0^T \dot{c}_i A_i x_0 v) = \sum_i \dot{c}_i [2v^T Q^{-1}(b_i - (A_i - x_0^T A_i x_0)v)]$$

$$\text{Thus, } \boxed{\frac{\partial z}{\partial c_i} = 2v^T Q^{-1}(b_i - (A_i - x_0^T A_i x_0)v)}, \quad Q = c \cdot A - \lambda_{\min}(c \cdot A), \quad v = Q^{-1}(c \cdot b), \quad x_0 \in \text{Ker } Q, \quad \|x_0\| = 1$$

not the same as (0.8) in draft.pdf (but numerically same):

$$\dot{z}_{(0.8)} = 2 \underbrace{(c \cdot b)^T Q^{-2}(\dot{c} \cdot b)}_{v^T Q^{-1}} - v^T (Q^{-1}\dot{Q} + \dot{Q}Q^{-1})v = 2v^T Q^{-1}(\dot{c} \cdot b) - v^T Q^{-1}\dot{Q}v - v^T \dot{Q}Q^{-1}v$$

5. (0.10). If $\dot{c} = \beta c_+$, then $\dot{z} = 2v^T Q^{-1}(\dot{c} \cdot b - \dot{Q}v) = \boxed{\equiv}$. Since $c_+ \cdot b = 0$, $c_+ \cdot A = I$, $\dot{Q} = c_+ \cdot A - x_0^T(c_+ \cdot A)x_0 = I - x_0^T x_0 = I - 1 = 0$. And $\boxed{\equiv} 0$ (correct with new \dot{z}).

$$6. (0.14) \quad n_i = (b_i^T - v^T(A_i - x_0^T A_i x_0))x_0$$

$$7. (0.16) \quad P(\lambda) = Q^{-1}Q = S\Lambda^g S^T S\Lambda S^T = S\Lambda\Lambda^g S^T = 1 - x_0 x_0^T - \text{projector on } (\text{Ker } Q)^\perp \text{ (correct)}$$

$$8. (0.15) \quad P(\lambda)(c(\lambda) \cdot b) = c(\lambda) \cdot b \Leftrightarrow \underbrace{(c(\lambda) \cdot b) \perp \text{Ker } Q}_{\Leftrightarrow c(\lambda) \in c_{\text{bad}}} \Leftrightarrow c(\lambda) \cdot b \in \text{Im } Q \Leftrightarrow \exists \hat{x}: Q\hat{x} = c(\lambda) \cdot b \quad (0.17)$$

(correct)

3.10 Gradient descent. Projection

We have $c \in \mathbb{R}^m$, $c \in c_{\text{bad}} = \{c \mid \|c\| = 1, \text{Ker } Q(c) \perp (c \cdot b)\}$. $Q = c \cdot A - \lambda_{\min}(c \cdot A)$, $v = Q^{-1}(c \cdot b)$, $x_0 \in \text{Ker } Q$, $\|x_0\| = 1$

$$1. \text{ Calculate } \frac{\partial z}{\partial c_i} = 2v^T Q^{-1}(b_i - (A_i - x_0^T A_i x_0)v). \text{ Define } \Delta c = -\nabla z(c)$$

$$2. \text{ Calculate } \hat{n}_i = (b_i^T - v^T(A_i - x_0^T A_i x_0))x_0, \text{ define } n_i = \frac{\hat{n}_i}{\|\hat{n}_i\|}$$

$$3. \text{ Define } c' = c + \Delta c - n(\Delta c, n)$$

$$4. \text{ Define } c(\lambda) = c' + \lambda n. \text{ Define } x_0(\lambda) \text{ s.t. } x_0(\lambda) \in \text{Ker } Q(\lambda), x_0(\lambda)^T x_0 > 0, \|x_0(\lambda)\| = 1. \text{ Define } m(\lambda) = (c(\lambda) \cdot b)^T x_0(\lambda). \text{ Beware of } \text{Rg } Q < n - 1.$$

$$5. \text{ Find root of } m(\lambda) \text{ using binary search on } [-\lambda^0, \lambda^0], \lambda^0 = \|c - c'\|.$$

$$6. \text{ Next } c: c(\lambda)$$

$$m(\lambda) = 0 \Leftrightarrow (c(\lambda) \cdot b) \perp \{x_0\} = \text{Ker } Q \text{ if } \text{Rg } Q = n - 1.$$

3.11 Equation 0.19

Consider $A = S\Lambda S^T$, $S^T S = E$, x_0 is a simple zero eigenvector of A : $Ax_0 = 0$, $\|x_0\| = 1$.
 $A^{-1} = S\Lambda_1 S^T$, $\Lambda = (\lambda_1, \dots, \lambda_{n-1}, 0)$, $\Lambda_1 = (\lambda_1^{-1}, \dots, \lambda_{n-1}^{-1}, 0)$.
Equation (0.19):

$$\dot{x}_0 = -A^{-1}\dot{A}x_0$$

Consider $A^{-1}\dot{A}x_0 = S\Lambda_1 S^T(\dot{S}\Lambda S^T + S\dot{\Lambda}S^T + S\Lambda\dot{S}^T)x_0 \equiv$.

Consider $x_0 = Sy_0$, where $y_0 = (0, 0, \dots, 1)$. Therefore, $0 = \dot{y}_0 = \dot{S}^T x_0 + S^T \dot{x}_0$

Going back to 0.19, the part $\Lambda S^T x_0 = \Lambda y_0 = 0$, another part $\dot{\Lambda} S^T x_0 = 0$. Consequently,

$$\equiv S\Lambda_1 S^T S\Lambda\dot{S}^T x_0 = S\Lambda_1 \Lambda\dot{S}^T x_0 = -S\Lambda_1 \Lambda S^T \dot{x}_0 = -\sum_{i=1}^{n-1} s_i s_i^T \dot{x}_0 = -(E - x_0 x_0^T) \dot{x}_0 = -\dot{x}_0 + x_0 x_0^T \dot{x}_0 \equiv.$$

Taking a derivative $\|x_0\| = 1$, we get $x_0^T \dot{x}_0 = 0$, therefore,

$$\equiv -\dot{x}_0$$

3.12 Nonconvexity certificate in $b_i = 0$ case

Let $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$ be a quadratic map: $f_i(x) = x^T A_i x$, $A_i = A_i^T$.

Consider $c \in \mathbb{R}^m$ and boundary points ∂F_c :

$$c \cdot f(x) \rightarrow \min_c$$

Where $A = \sum c_i A_i$. Assuming $A \geq 0$.

Minimization leads to $Ax = 0$. The following cases hold:

1. $RgA = n$. $x = 0$ is a unique solution
2. $RgA = n - 1$. $x = \alpha e$, $f(x) = \alpha^2 f(e)$
3. $RgA = n - 2$. In this case $x = \alpha_1 e^1 + \alpha_2 e^2$. Consider $f(x) = \alpha_1^2 f_{11} + 2\alpha_1 \alpha_2 f_{12} + \alpha_2^2 f_{22}$.
 - (a) f_{11}, f_{22}, f_{12} are linearly independent. In this case ∂F_c is nonconvex
 - (b) $Rg\|f_{11} f_{12} f_{22}\| = 1$. ∂F_c convex.

Result If exist $c \in \mathbb{R}^m$:

1. $RgA = n - 2$
2. $A \geq 0$
3. f_{ij} are linearly independent

Then $F = \text{Im} f$ is nonconvex.

3.13 $\text{conv } F$

Пусть $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, $f_i(x) = x^T A_i x + 2b_i^T x$, $A_i = A_i^T$.

Обозначим $F = f(\mathbb{R}^n)$, $G = \text{conv} F$

Обозначим $H_i = \begin{vmatrix} A_i & b_i \\ b_i^T & 0 \end{vmatrix}$

Обозначим $X = \begin{vmatrix} x \\ 1 \end{vmatrix} \begin{vmatrix} x^T & 1 \end{vmatrix} = \begin{vmatrix} xx^T & x \\ x^T & 1 \end{vmatrix}$

Тогда $f_i(x) = \text{tr} H_i X$, $f(x) = H(X)$

Обозначим $V = \{X \in \mathbb{R}^{(n+1) \times (n+1)} | X = X^T, X \geq 0, X_{n+1, n+1} = 1\}$

Обозначим $G_1 = H(V)$.

Доказать: $G_1 = G$ (On the feasibility for the system of quadratic equations, Theorem 3.1. (Convex hull))

1. $G \subseteq G_1$. Пусть $y \in G$. Тогда $\exists \{y_i\}_{i=1}^l \subset F$, $\{\lambda_i\}_{i=1}^l: y = \sum_{i=1}^l \lambda_i y_i$, где $\lambda_i \geq 0$, $\sum \lambda_i = 1$.

Поскольку $y_i \in F$, $\exists \{X_i\}: y_i = H(X_i)$, причем $X_i = \begin{pmatrix} x_i x_i^T & x_i \\ x_i^T & 1 \end{pmatrix} \in V$. Рассмотрим j -ю компоненту $y^j = \sum \lambda_i y_i^j = \sum \lambda_i \text{tr} H_j X_i = \text{tr} H_j \underbrace{\sum \lambda_i X_i}_X$. То есть, найден $X \in V: y = H(X)$.

Значит, $y \in G_1$

2. $G_1 \subseteq G$. Пусть $y \in G_1$. Тогда $y = H(X)$, $X \in V$. Доказать: $y \in G = \text{conv} F$. Представим X в виде выпуклой комбинации $X = \sum \lambda_i X_i$, где $X_i \in V$, причем $X_i = \begin{pmatrix} x_i x_i^T & x_i \\ x_i^T & 1 \end{pmatrix} \in V$ для некоторого x_i . Это докажет $y \in G$.

Рассмотрим $X = \sum_{k=1}^n \lambda_k s_k s_k^T$ — спектральное разложение. Поскольку $X \in V$, $X \geq 0$, значит, $\lambda_k \geq 0$. Обозначим $\Lambda = \sum \lambda_k$. Пусть $s_{k,n+1} \neq \frac{1}{\Lambda}$. Тогда переопределим $s_{k,n+1} = \frac{1}{\Lambda}$. Это можно сделать, т.к. $H_i(s_k s_k^T)$ не изменится (прямая проверка, использовать $H_{i,n+1,n+1} = 0$).

Рассмотрим $X = \sum \underbrace{\frac{\lambda_k}{\Lambda}}_{\alpha_k} \underbrace{\Lambda s_k s_k^T}_{X_k}$. Получили представление $X = \alpha_k X_k$, X_k имеет нужный вид,

α_k — выпуклая комбинация. Получаем $y \in G$

3.14 Small ball

Пусть $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, $f_i(x) = x^T A_i x - 2b_i^T x$, $A_i = A_i^T$. Пусть $c \in \mathbb{R}^m$.

Обозначим $c \cdot A = \sum_{i=1}^n c_i A_i$, $c \cdot b = \sum_{i=1}^n c_i b_i$, $F_c(x) = c^T f(x)$

Хотим найти:

$$\min_{\|x\|^2=1} F_c(x)$$

Функция Лагранжа: $L(x, \lambda) = x^T (c \cdot A) x - 2(c \cdot b)^T x - \lambda(\|x\|^2 - 1)$.

Находим $L_x = 2(c \cdot A)x - 2c \cdot b - 2\lambda x = 0$, $L_\lambda = \|x\|^2 - 1 = 0$,

Получаем систему $\begin{cases} \|x\| = 1 \\ (c \cdot A - \lambda)x = c \cdot b \end{cases}$ Это совпадает с (2.3).

Далее переходим в базис из собственных векторов $\{x_i\}$ симметричной матрицы $c \cdot A$

$$S = \|x_1 \dots x_n\|, S^T S = E$$

$$x = Sy, \Lambda = S^T (c \cdot A) S, c \cdot b = S\alpha$$

Получаем

$$\begin{cases} \|y\| = 1 & (1) \\ (\Lambda - \lambda)y = \alpha & (2) \end{cases}$$

Выражаем из (2)

$$y_k = \frac{\alpha_k}{\lambda_k - \lambda}$$

Подставляем в (1) — это совпадает с (2.7)

$$\sum_{i=1}^n \left(\frac{\alpha_k}{\lambda_k - \lambda} \right)^2 = 1$$

Это соотношение определяет λ , по λ находим y_k , затем находим x в старом базисе.