

Adaptive Submodularity for Active Classification

Sergey Karayev

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- ▶ Have:
 - ▶ Hypothesis $H \in \mathcal{H}$, set of possible observations $\mathcal{X} = \{X_1, \dots, X_F\}$.
 - ▶ Each X_i has cost c_i and depends on H in some way.
- ▶ Goal: select subset \mathcal{X}_S to maximize reward $f(\mathcal{X}_S)$.
 - ▶ e.g. loss: $f(\mathcal{X}_S) = -\ell(h^*, \operatorname{argmax}_h P(h | \mathcal{X}_S))$.
 - ▶ e.g. entropy: $f(\mathcal{X}_S) = \sum_x P(h | \mathcal{X}_S) \log P(h | \mathcal{X}_S)$
- ▶ Variants:
 - ▶ *Budgeted maximization*: Maximize objective without exceeding cost budget.
 - ▶ *Minimum-cost cover*: Minimize cost while achieving some minimum objective value.

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- ▶ Furthermore, want to do this *adaptively*: selecting and observing values in sequence (motivation to follow).
- ▶ Lots of recent work in this setting: POMDPs, cascades, trees for spam detection, object recognition, sentence parsing.
- ▶ Different heuristic methods, but no performance guarantees.
- ▶ Here, we synthesize a few recent papers that may construe a promising avenue for theoretical analysis of the active classification problem.

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Offline problem formulation

Take the problem of selecting the whole subset before observing anything; assume uniform costs.

Budgeted maximization

$$\underset{\mathcal{X}_S \subseteq \mathcal{X}}{\operatorname{argmax}} f(\mathcal{X}_S)$$

$$s.t | \mathcal{X}_S | \leq N$$

Min-cost cover

$$\underset{\mathcal{X}_S \subseteq \mathcal{X}}{\operatorname{argmin}} |\mathcal{X}_S|$$

$$s.t f(\mathcal{X}_S) \geq Q$$

- ▶ For most interesting f , this is NP-hard.
- ▶ Nemhauser et al. (1978) showed that if f is *monotone submodular*, then a simple greedy algorithm is near-optimal with these bounds:

$$f(\mathcal{X}_{S_N}) \geq \left(1 - \frac{1}{e}\right) f(\mathcal{X}_{S_N}^*) \quad |\mathcal{X}_{S_Q}| \leq |\mathcal{X}_S^*| (1 + \log \max_X f(X))$$

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Monotone submodular functions

- ▶ Intuitively, submodularity is the property of diminishing returns. For proving near-optimality of the greedy method, we additionally require monotonicity.
- ▶ Define the *gain* of adding another item to a set as

$$\Delta_f(X \mid \mathcal{X}_A) = f(\mathcal{X}_A \cup \{X\}) - f(\mathcal{X}_A)$$

- ▶ For all X , the submodularity requirement must hold.
“Adding an element to a smaller set helps more than adding the same element to a larger set.”

$$\mathcal{X}_A \subseteq \mathcal{X}_B \implies \Delta_f(X \mid \mathcal{X}_A) \geq \Delta_f(X \mid \mathcal{X}_B)$$

- ▶ For all X , the monotonicity requirement must hold.
“Adding an element can never hurt.”

$$\Delta_f(X \mid \mathcal{X}_A) \geq 0$$

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Greedy algorithm

Input: Budget N ; observations \mathcal{X} ; objective f .

Output: Set $\mathcal{X}_S \subseteq \mathcal{X}$ of size N

begin

$\mathcal{X}_S \leftarrow \emptyset;$

for $i = 1$ **to** k **do**

foreach $X \in \mathcal{X} \setminus \mathcal{X}_S$ **do** compute $\Delta_f(X \mid \mathcal{X}_S)$;

Select $X^* \in \underset{X}{\operatorname{argmax}} \frac{\Delta_f(X \mid \mathcal{X}_S)}{c(X)}$;

Set $\mathcal{X}_S \leftarrow \mathcal{X}_S \cup \{X^*\}$;

end

end

- ▶ Costs have been omitted in the exposition, but are handled by the algorithm if the cost function is modular:

$$c(\mathcal{X}_S) = \sum_{X \in \mathcal{X}_S} c(X)$$

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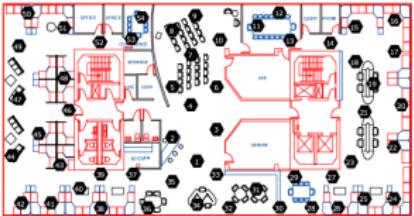
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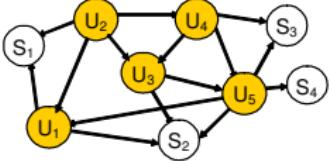
Future work

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Example



(a) Sensor network deployment



(b) Graphical model

- ▶ Sensor network: want minimal set of observations (Sensors) to maximize utility over hidden variables (Fire): $f(P(\text{Fire} \mid \text{Sensors}))$.
- ▶ If the observation model satisfies certain conditional independence requirements, then Information Gain (reduction in entropy) is submodular [Krause and Guestrin (2005)].
- ▶ Classification models like Naive Bayes are very similar.

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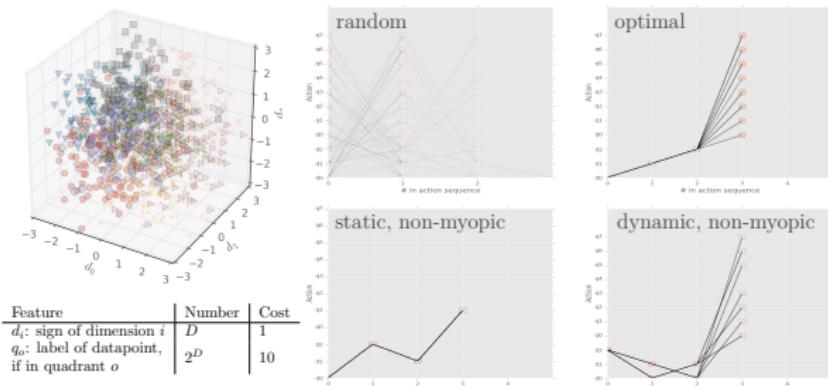
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Adaptivity motivation

- With the greedy algorithm and monotone submodular objective, we can find a near-optimal subset.
- But what if we can observe the results of our actions upon execution?
- In this case, offline submodularity also can be exponentially far from optimal [Hollinger et al. (2011)].



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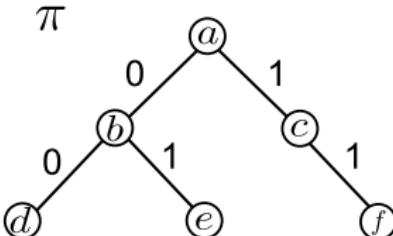
Future work

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Adaptive submodularity

- ▶ Introduced by Golovin and Krause (2011)¹.
- ▶ All F potential observations X_i are now random variables taking on values in \mathcal{O} , and have *realization* \mathbf{x} .
- ▶ Selected subset $\mathcal{X}_S \subseteq \mathcal{X}$ forms a *partial realization* \mathbf{x}_S .
- ▶ The objective function $f : 2^F \times |\mathcal{O}|^F \rightarrow \mathcal{R}$ takes a realization as well as the selected subset.
- ▶ The *policy* π maps partial realizations to new observations, and can be thought of as a decision tree.
- ▶ $F(\pi) = \mathbb{E}[f(\pi(\mathbf{x}), \mathbf{x})] = \sum_{\mathbf{x}} P(\mathbf{x}) f(\pi(\mathbf{x}), \mathbf{x})$

$$\begin{aligned}\pi(\emptyset) &= a \\ \pi(\{(a, 0)\}) &= b \\ \pi(\{(a, 1)\}) &= c \\ \pi(\{(a, 0), (b, 0)\}) &= d \\ \pi(\{(a, 0), (b, 1)\}) &= e \\ \pi(\{(a, 1), (c, 1)\}) &= f\end{aligned}$$



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¹ JAIR Best Paper Prize 2013

Adaptive submodular functions

- ▶ We transfer the guarantees of greedy optimization with submodular functions to the adaptive setting by taking expectation over realizations, conditioned on the current partial realization.
- ▶ Define the *conditional expected gain* of adding another item to a set as

$$\Delta_f(X \mid \mathbf{x}_A) = \mathbb{E}[f(\mathcal{X}_A \cup \{X\}, \mathbf{x}) - f(\mathcal{X}_A, \mathbf{x}) \mid \mathbf{x}_A]$$

- ▶ For all X , the function must be adaptive submodular.

$$\mathbf{x}_A \preceq \mathbf{x}_B \implies \Delta_f(X \mid \mathbf{x}_A) \geq \Delta_f(X \mid \mathbf{x}_B)$$

- ▶ For all X , the function must be adaptive monotone.

$$\Delta_f(X \mid \mathbf{x}_A) \geq 0$$

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Input: Budget N ; observation set \mathcal{X} ; $p(\mathbf{x})$; objective f .

Output: Set $\mathcal{X}_S \subseteq \mathcal{X}$ of size N

begin

$\mathcal{X}_S \leftarrow \emptyset$; $\mathbf{x}_S \leftarrow \emptyset$;

for $i = 1$ **to** k **do**

foreach $X \in \mathcal{X} \setminus \mathcal{X}_S$ **do** compute $\Delta_f(X | \mathbf{x}_S)$;

 Select $X^* \in \operatorname{argmax}_X \frac{\Delta_f(X | \mathbf{x}_S)}{c(X)}$;

 Set $\mathcal{X}_S \leftarrow \mathcal{X}_S \cup \{X^*\}$;

 Observe value \mathbf{x}^* and update $\mathbf{x}_S \leftarrow \mathbf{x}_S \cup \mathbf{x}^*$;

end

end

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Submodularity → Adaptive submodularity

Remarkably, optimality bounds stay the same!

Submodularity

- ▶ Works with sets.
- ▶ $\Delta_f(X \mid \mathcal{X}_A) = f(\mathcal{X}_A \cup \{X\}) - f(\mathcal{X}_A)$
- ▶ $\Delta_f(X \mid \mathcal{X}_A) \geq \Delta_f(X \mid \mathcal{X}_B)$ if $\mathcal{X}_A \subseteq \mathcal{X}_B$
- ▶ $\Delta_f(X \mid \mathcal{X}_A) \geq 0$
- ▶ $\max_{\mathcal{X}_S} f(\mathcal{X}_S)$ s.t. $|\mathcal{X}_S| \leq N$
- ▶ Budgeted maximization: greedy is within $(1 - 1/e)$ of OPT.
- ▶ Min-cost cover: greedy subset size is less than $1 + \log \max_X f(X)$ of OPT subset size.

Adaptive submodularity

- ▶ Works with policies.
- ▶ $\Delta(X \mid \mathbf{x}_A) = \mathbb{E}[f(\mathcal{X}_A \cup \{X\}, \mathbf{x}) - f(\mathcal{X}_A, \mathbf{x}) \mid \mathbf{x}_A]$
- ▶ $\Delta_f(X \mid \mathbf{x}_A) \geq \Delta_f(X \mid \mathbf{x}_B)$ if $\mathbf{x}_A \preceq \mathbf{x}_B$
- ▶ $\Delta_f(X \mid \mathbf{x}_A) \geq 0$
- ▶ $\max_{\pi} F(\pi)$ s.t. $|\pi| \leq N$
- ▶ Budgeted maximization: greedy policy is within $(1 - 1/e)$ of OPT policy.
- ▶ Min-cost cover: greedy policy size is less than $1 + \log \max_X f(X, \mathbf{x})$ of OPT policy size.

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Benefit: Lazy evaluation

- ▶ A neat algorithmic speed-up follows from the submodularity of the objective: **the expected utility of an observation cannot increase**. This means not all observations need to be evaluated on every iteration.

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	Expected Utility 1	Expected Utility 2		Expected Utility
action 1	0.85		action 1	0.85
action 2	0.72		action 2	0.72
action 3	0.52		action 3	0.52
action 4	0.63		action 4	0.63

Formulating the objective

- ▶ So far, things look promising!
- ▶ Unfortunately, infogain is only adaptive submodular in the noiseless setting (in Naive Bayes, only one outcome x_i given h).
- ▶ If there is noise, need to find an objective that both optimizes what we want and is adaptive submodular.
- ▶ Have two previously explored options:
 - ▶ *Hypothesis Pruning* [Golovin et al. (2010)]
 - ▶ *Equivalence Class Edge Cutting* [Javdani et al. (2012)]
- ▶ Both have proofs only for the Min-Cost Cover setting.

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Hypothesis pruning

- ▶ Motivation: efficiently figure out where an object is in 3D space so that a robotic arm can grasp it.
- ▶ Idea: downweight hypotheses with a non-normalized Gaussian, which effectively removes a portion of the hypothesis space. By construction, the probability mass of the posterior distribution can only become smaller with each step.
- ▶ Proof: lengthy (in report). Tricky to prove that the *expected utility* does not increase.

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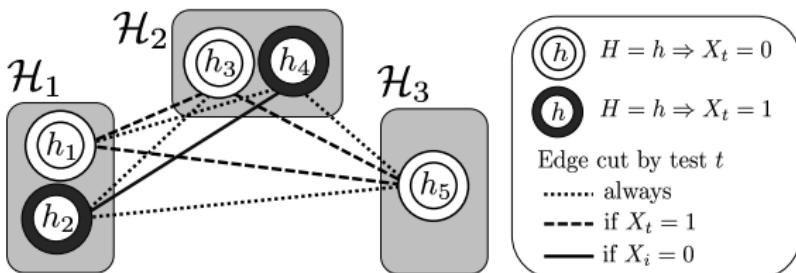
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Equivalence Class Edge Cutting

- ▶ Motivation: request most informative instance to be labeled by an oracle.
- ▶ Idea: expand the space of realizations to account for all possible noise settings. A single hypothesis can now be true in multiple states. Find the test that removes the most hypotheses.
- ▶ Proof: lengthy (in report).



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Future Work

- ▶ Empirically evaluate the two adaptive submodular objectives against my current POMDP approach and on my datasets.
 - ▶ Adaptive submodularity of a function can depend on the distribution $p(\mathbf{x})$.
- ▶ Modular² costs are fine. How about submodular costs?
- ▶ Analyze the adaptivity gap of these methods via data-dependent bounds.

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$$^2 c(\mathcal{X}_S) = \sum_{X \in \mathcal{X}_S} c(X)$$

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