

# Kernelization

Serge Gaspers

UNSW

- 1 Vertex Cover
  - Simplification rules
  - Preprocessing algorithm
- 2 Kernelization algorithms
- 3 Kernel for HAMILTONIAN CYCLE
- 4 Kernel for EDGE CLIQUE COVER
- 5 Kernels and Fixed-parameter tractability
- 6 Further Reading

# Outline

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# Vertex cover

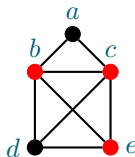
A **vertex cover** of a graph  $G = (V, E)$  is a subset of vertices  $S \subseteq V$  such that for each edge  $\{u, v\} \in E$ , we have  $u \in S$  or  $v \in S$ .

## VERTEX COVER

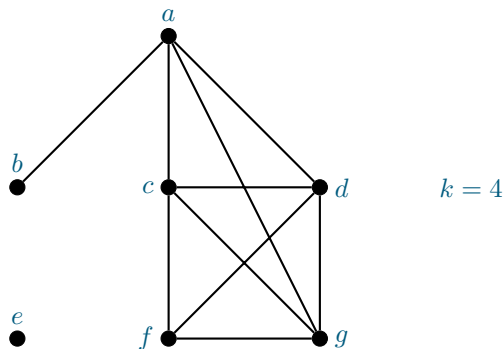
Input: A graph  $G = (V, E)$  and an integer  $k$

Parameter:  $k$

Question: Does  $G$  have a vertex cover of size at most  $k$ ?



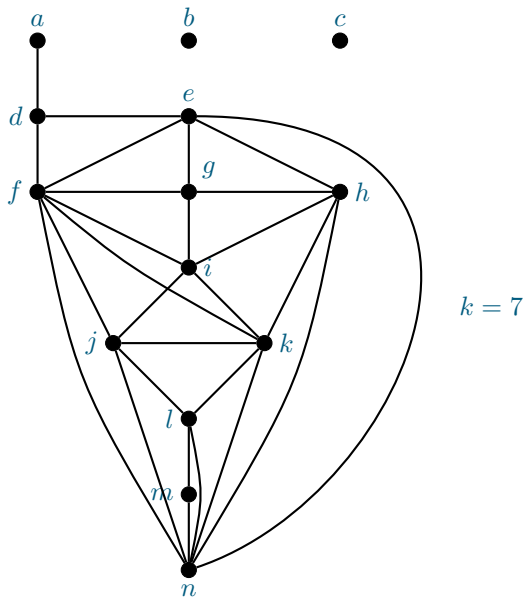
# Exercise 1



Is this a **YES**-instance for VERTEX COVER?

(Is there  $S \subseteq V$  with  $|S| \leq 4$ , such that  $\forall uv \in E, u \in S$  or  $v \in S$ ?)

## Exercise 2



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# Simplification rules for VERTEX COVER

## (Degree-0)

If  $\exists v \in V$  such that  $d_G(v) = 0$ , then set  $G \leftarrow G - v$ .



# Simplification rules for VERTEX COVER

## (Degree-0)

If  $\exists v \in V$  such that  $d_G(v) = 0$ , then set  $G \leftarrow G - v$ .

**Proving correctness.** A simplification rule is **sound** if for every instance, it produces an equivalent instance. Two instances  $I, I'$  are **equivalent** if they are both **YES**-instances or they are both **NO**-instances.

## Lemma 1

*(Degree-0) is sound.*

# Simplification rules for VERTEX COVER

## (Degree-0)

If  $\exists v \in V$  such that  $d_G(v) = 0$ , then set  $G \leftarrow G - v$ .

**Proving correctness.** A simplification rule is **sound** if for every instance, it produces an equivalent instance. Two instances  $I, I'$  are **equivalent** if they are both **YES**-instances or they are both **NO**-instances.

## Lemma 1

*(Degree-0) is sound.*

## Proof.

First, suppose  $(G - v, k)$  is a **YES**-instance. Let  $S$  be a vertex cover for  $G - v$  of size at most  $k$ . Then,  $S$  is also a vertex cover for  $G$  since no edge of  $G$  is incident to  $v$ . Thus,  $(G, k)$  is a **YES**-instance.

Now, suppose  $(G - v, k)$  is a **NO**-instance. For the sake of contradiction, assume  $(G, k)$  is a **YES**-instance. Let  $S$  be a vertex cover for  $G$  of size at most  $k$ . But then,  $S \setminus \{v\}$  is a vertex cover of size at most  $k$  for  $G - v$ ; a contradiction.  $\square$

# Simplification rules for VERTEX COVER

(Degree-1)

If  $\exists v \in V$  such that  $d_G(v) = 1$ , then set  $G \leftarrow G - N_G[v]$  and  $k \leftarrow k - 1$ .

# Simplification rules for VERTEX COVER

## (Degree-1)

If  $\exists v \in V$  such that  $d_G(v) = 1$ , then set  $G \leftarrow G - N_G[v]$  and  $k \leftarrow k - 1$ .

## Lemma 1

*(Degree-1) is sound.*

## Proof.

Let  $u$  be the neighbor of  $v$  in  $G$ . Thus,  $N_G[v] = \{u, v\}$ .

If  $S$  is a vertex cover of  $G$  of size at most  $k$ , then  $S \setminus \{u, v\}$  is a vertex cover of  $G - N_G[v]$  of size at most  $k - 1$ , because  $u \in S$  or  $v \in S$ .

If  $S'$  is a vertex cover of  $G - N_G[v]$  of size at most  $k - 1$ , then  $S' \cup \{u\}$  is a vertex cover of  $G$  of size at most  $k$ , since all edges that are in  $G$  but not in  $G - N_G[v]$  are incident to  $u$ . □

# Simplification rules for VERTEX COVER

(Large Degree)

If  $\exists v \in V$  such that  $d_G(v) > k$ , then set  $G \leftarrow G - v$  and  $k \leftarrow k - 1$ .

# Simplification rules for VERTEX COVER

## (Large Degree)

If  $\exists v \in V$  such that  $d_G(v) > k$ , then set  $G \leftarrow G - v$  and  $k \leftarrow k - 1$ .

## Lemma 1

*(Large Degree) is sound.*

## Proof.

Let  $S$  be a vertex cover of  $G$  of size at most  $k$ . If  $v \notin S$ , then  $N_G(v) \subseteq S$ , contradicting that  $|S| \leq k$ . □

# Simplification rules for VERTEX COVER

## (Number of Edges)

If  $d_G(v) \leq k$  for each  $v \in V$  and  $|E| > k^2$  then return No

# Simplification rules for VERTEX COVER

## (Number of Edges)

If  $d_G(v) \leq k$  for each  $v \in V$  and  $|E| > k^2$  then return **No**

## Lemma 1

*(Number of Edges) is sound.*

## Proof.

Assume  $d_G(v) \leq k$  for each  $v \in V$  and  $|E| > k^2$ .

Suppose  $S \subseteq V$ ,  $|S| \leq k$ , is a vertex cover of  $G$ .

We have that  $S$  covers at most  $k^2$  edges.

However,  $|E| \geq k^2 + 1$ .

Thus,  $S$  is not a vertex cover of  $G$ . □



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# Preprocessing algorithm for VERTEX COVER

VC-preprocess

**Input:** A graph  $G$  and an integer  $k$ .

**Output:** A graph  $G'$  and an integer  $k'$  such that  $G$  has a vertex cover of size at most  $k$  if and only if  $G'$  has a vertex cover of size at most  $k'$ .

$G' \leftarrow G$

$k' \leftarrow k$

**repeat**

    Execute simplification rules (Degree-0), (Degree-1), (Large Degree), and  
    (Number of Edges) for  $(G', k')$

**until** *no simplification rule applies*

**return**  $(G', k')$

# Effectiveness of preprocessing algorithms

- How effective is VC-preprocess?
- We would like to study preprocessing algorithms mathematically and quantify their effectiveness.

- Say that a preprocessing algorithm for a problem  $\Pi$  is **nice** if it runs in polynomial time and for each instance for  $\Pi$ , it returns an instance for  $\Pi$  that is strictly smaller.

# First try

- Say that a preprocessing algorithm for a problem  $\Pi$  is **nice** if it runs in polynomial time and for each instance for  $\Pi$ , it returns an instance for  $\Pi$  that is strictly smaller.
- $\rightarrow$  executing it a linear number of times reduces the instance to a single bit
- $\rightarrow$  such an algorithm would solve  $\Pi$  in polynomial time
- For **NP**-hard problems this is not possible unless  $P = NP$
- We need a different measure of effectiveness

# Measuring the effectiveness of preprocessing algorithms

- We will measure the effectiveness in terms of the **parameter**
- How large is the resulting instance in terms of the parameter?

# Effectiveness of VC-preprocess

## Lemma 2

*For any instance  $(G, k)$  for VERTEX COVER, VC-preprocess produces an equivalent instance  $(G', k')$  of size  $O(k^2)$ .*

## Proof.

Since all simplification rules are sound,  $(G = (V, E), k)$  and  $(G' = (V', E'), k')$  are equivalent.

By (Number of Edges),  $|E'| \leq (k')^2 \leq k^2$ .

By (Degree-0) and (Degree-1), each vertex in  $V'$  has degree at least 2 in  $G'$ .

Since  $\sum_{v \in V'} d_{G'}(v) = 2|E'| \leq 2k^2$ , this implies that  $|V'| \leq k^2$ .

Thus,  $|V'| + |E'| \leq O(k^2)$ . □

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## Definition 3

A **kernelization** for a parameterized problem  $\Pi$  is a **polynomial time** algorithm, which, for any instance  $I$  of  $\Pi$  with parameter  $k$ , produces an **equivalent** instance  $I'$  of  $\Pi$  with parameter  $k'$  such that  $|I'| \leq f(k)$  and  $k' \leq f(k)$  for a computable function  $f$ .

We refer to the function  $f$  as the **size** of the kernel.

**Note:** We do not formally require that  $k' \leq k$ , but this will be the case for many kernelizations.

# VC-preprocess is a quadratic kernelization

## Theorem 4

*VC-preprocess is a  $O(k^2)$  kernelization for VERTEX COVER.*

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# HAMILTONIAN CYCLE I

A **Hamiltonian cycle** of  $G$  is a subgraph of  $G$  that is a cycle on  $|V(G)|$  vertices.

## vc-HAMILTONIAN CYCLE

Input: A graph  $G = (V, E)$ .

Parameter:  $k = vc(G)$ , the size of a smallest vertex cover of  $G$ .

Question: Does  $G$  have a Hamiltonian cycle?

**Thought experiment:** Imagine a very large instance where the parameter is tiny. How can you simplify such an instance?

**Issue:** We do not actually know a vertex cover of size  $k$ .  
We do not even know the value of  $k$  (it is not part of the input).

# HAMILTONIAN CYCLE III

- Obtain a vertex cover using an approximation algorithm. We will use a 2-approximation algorithm, producing a vertex cover of size  $\leq 2k$  in polynomial time.
- If  $C$  is a vertex cover of size  $\leq 2k$ , then  $I = V \setminus C$  is an independent set of size  $\geq |V| - 2k$ .
- No two consecutive vertices in the Hamiltonian Cycle can be in  $I$ .
- A kernel with  $\leq 4k$  vertices can now be obtained with the following simplification rule.

## (Too-large)

Compute a vertex cover  $C$  of size  $\leq 2k$  in polynomial time.

If  $2|C| < |V|$ , then return **No**

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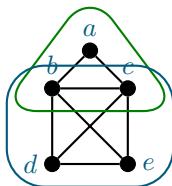
# EDGE CLIQUE COVER

## Definition 5

An **edge clique cover** of a graph  $G = (V, E)$  is a set of cliques in  $G$  covering all its edges.

In other words, if  $\mathcal{C} \subseteq 2^V$  is an edge clique cover then each  $S \in \mathcal{C}$  is a clique in  $G$  and for each  $\{u, v\} \in E$  there exists an  $S \in \mathcal{C}$  such that  $u, v \in S$ .

Example:  $\{\{a, b, c\}, \{b, c, d, e\}\}$  is an edge clique cover for this graph.





# EDGE CLIQUE COVER

## EDGE CLIQUE COVER

Input: A graph  $G = (V, E)$  and an integer  $k$

Parameter:  $k$

Question: Does  $G$  have an edge clique cover of size at most  $k$ ?

The **size** of an edge clique cover  $\mathcal{C}$  is the number of cliques contained in  $\mathcal{C}$  and is denoted  $|\mathcal{C}|$ .

# Helpful properties

## Definition 5

A clique  $S$  in a graph  $G$  is a **maximal** clique if there is no other clique  $S'$  in  $G$  with  $S \subset S'$ .

## Lemma 6

*A graph  $G$  has an edge clique cover  $\mathcal{C}$  of size at most  $k$  if and only if  $G$  has an edge clique cover  $\mathcal{C}'$  of size at most  $k$  such that each  $S \in \mathcal{C}'$  is a maximal clique.*

## Proof sketch.

( $\Rightarrow$ ): Replace each clique  $S \in \mathcal{C}$  by a maximal clique  $S'$  with  $S \subseteq S'$ .

( $\Leftarrow$ ): Trivial, since  $\mathcal{C}'$  is an edge clique cover of size at most  $k$ . □

# Simplification rules for EDGE CLIQUE COVER

**Thought experiment:** Imagine a very large instance where the parameter is tiny. How can you simplify such an instance?

# Simplification rules for EDGE CLIQUE COVER II

The instance could have many degree-0 vertices.

(Isolated)

If there exists a vertex  $v \in V$  with  $d_G(v) = 0$ , then set  $G \leftarrow G - v$ .

Lemma 7

*(Isolated) is sound.*

Proof sketch.

Since no edge is incident to  $v$ , a smallest edge clique cover for  $G - v$  is a smallest edge clique cover for  $G$ , and vice-versa.  $\square$

# Simplification rules for EDGE CLIQUE COVER II

The instance could have many degree-0 vertices.

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If there exists a vertex  $v \in V$  with  $d_G(v) = 0$ , then set  $G \leftarrow G - v$ .

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## Proof sketch.

Since no edge is incident to  $v$ , a smallest edge clique cover for  $G - v$  is a smallest edge clique cover for  $G$ , and vice-versa.  $\square$

## (Isolated-Edge)

If  $\exists uv \in E$  such that  $d_G(u) = d_G(v) = 1$ , then set  $G \leftarrow G - \{u, v\}$  and  $k \leftarrow k - 1$ .

# Simplification rules for EDGE CLIQUE COVER III

(Twins)

If  $\exists u, v \in V$ ,  $u \neq v$ , such that  $N_G[u] = N_G[v]$ , then set  $G \leftarrow G - v$ .

Lemma 8

*(Twins) is sound.*

# Simplification rules for EDGE CLIQUE COVER III

## (Twins)

If  $\exists u, v \in V$ ,  $u \neq v$ , such that  $N_G[u] = N_G[v]$ , then set  $G \leftarrow G - v$ .

## Lemma 8

*(Twins) is sound.*

## Proof.

We need to show that  $G$  has an edge clique cover of size at most  $k$  if and only if  $G - v$  has an edge clique cover of size at most  $k$ .

( $\Rightarrow$ ): If  $\mathcal{C}$  is an edge clique cover of  $G$  of size at most  $k$ , then  $\{S \setminus \{v\} : S \in \mathcal{C}\}$  is an edge clique cover of  $G - v$  of size at most  $k$ .

( $\Leftarrow$ ): Let  $\mathcal{C}'$  be an edge clique cover of  $G - v$  of size at most  $k$ . Partition  $\mathcal{C}'$  into  $\mathcal{C}'_u = \{S \in \mathcal{C}' : u \in S\}$  and  $\mathcal{C}'_{\neg u} = \mathcal{C}' \setminus \mathcal{C}'_u$ . Note that each set in  $\mathcal{C}_u = \{S \cup \{v\} : S \in \mathcal{C}'_u\}$  is a clique in  $G$  since  $N_G[u] = N_G[v]$  and that each edge incident to  $v$  is contained in at least one of these cliques. Now,  $\mathcal{C}_u \cup \mathcal{C}'_{\neg u}$  is an edge clique cover of  $G$  of size at most  $k$ .  $\square$

# Simplification rules for EDGE CLIQUE COVER IV

## (Size- $V$ )

If the previous simplification rules do not apply and  $|V| > 2^k$ , then return **No**.

## Lemma 9

*(Size- $V$ ) is sound.*



# Simplification rules for EDGE CLIQUE COVER IV

## (Size-V)

If the previous simplification rules do not apply and  $|V| > 2^k$ , then return **No**.

## Lemma 9

*(Size-V) is sound.*

## Proof.

For the sake of contradiction, assume neither (Isolated) nor (Twins) are applicable,  $|V| > 2^k$ , and  $G$  has an edge clique cover  $\mathcal{C}$  of size at most  $k$ . Since  $2^{\mathcal{C}}$  (the set of all subsets of  $\mathcal{C}$ ) has size at most  $2^k$ , and every vertex belongs to at least one clique in  $\mathcal{C}$  by (Isolated), we have that there exists two vertices  $u, v \in V$  such that  $\{S \in \mathcal{C} : u \in S\} = \{S \in \mathcal{C} : v \in S\}$ . But then,  $N_G[u] = \bigcup_{S \in \mathcal{C}: u \in S} S = \bigcup_{S \in \mathcal{C}: v \in S} S = N_G[v]$ , contradicting that (Twin) is not applicable.  $\square$

# Kernel for EDGE CLIQUE COVER

Theorem 10 ((Gramm et al., 2008))

EDGE CLIQUE COVER *has a kernel with  $O(2^k)$  vertices and  $O(4^k)$  edges.*

Corollary 11

EDGE CLIQUE COVER *is FPT.*

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## Theorem 12

*Let  $\Pi$  be a decidable parameterized problem.*

*$\Pi$  has a kernelization algorithm  $\Leftrightarrow \Pi$  is FPT.*

# Kernels and Fixed-parameter tractability

## Theorem 12

Let  $\Pi$  be a decidable parameterized problem.

$\Pi$  has a kernelization algorithm  $\Leftrightarrow \Pi$  is FPT.

## Proof.

( $\Rightarrow$ ): An FPT algorithm is obtained by first running the kernelization, and then any brute-force algorithm on the resulting instance.

( $\Leftarrow$ ): Let  $A$  be an FPT algorithm for  $\Pi$  with running time  $O(f(k)n^c)$ .

If  $f(k) < n$ , then  $A$  has running time  $O(n^{c+1})$ . In this case, the kernelization algorithm runs  $A$  and returns a trivial YES- or NO-instance depending on the answer of  $A$ .

Otherwise,  $f(k) \geq n$ . In this case, the kernelization algorithm outputs the input instance. □

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# Further Reading

- Chapter 2, *Kernelization* in (Cygan et al., 2015)
- Chapter 4, *Kernelization* in (Downey and Fellows, 2013)
- Chapter 7, *Data Reduction and Problem Kernels* in (Niedermeier, 2006)
- Chapter 9, *Kernelization and Linear Programming Techniques* in (Flum and Grohe, 2006)
- the kernelization book (Fomin et al., 2019)

# References

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