

**7.** Prove that the maximum number of nodes in a binary tree with height  $h$  is  $2^{h+1} - 1$ .

When the tree is a perfect binary tree every level is filled and we have the summation  $\sum_{i=0}^h 2^i = 2^{h+1} - 1$ . We conclude the maximum number of nodes in binary tree with height  $h$  is  $2^{h+1} - 1$ .

**8.** Prove that it takes  $\Omega(n \cdot \log(n))$  steps in the best case to build a binary search tree having  $n$  distinct keys.

In the best case the binary search tree insertion takes  $\log(k)$  where  $k$  is the current number of elements in the tree. we have the summation  $\sum_{i=1}^n \log(i) = \Omega(\int_1^n \log(x) \cdot dx) = \Omega(n \cdot \log(n))$ .

**9.** Prove that, when a binary tree with  $n$  nodes is implemented using links to the left and right child, then there will be a total of  $n + 1$  null links.

We prove by induction. When  $n = 1$  there is only the root which has a null left child and null right child hence it has 2 null links, hence  $P(1)$  holds. we assume it holds for all integers  $i \leq k$  for some positive integer  $k$  and show it must hold for the  $k + 1$  integer.

Case 1: the root only has child. WLOG we assume it is the left child. the left subtree has  $n$  elements and by the inductive hypothesis has  $n + 1$  null links. Thus we have  $(n + 1) + 1$ . Which proves the 1st case.

Case 2: the root has both children. Let the left subtree contain  $n - k$  elements and the right subtree contain  $k$  elements where  $k \leq n$ . by the inductive hypothesis the left subtree contains  $n - k + 1$  null links and the right subtree contains  $k + 1$  null links. we have in total  $(n + 1) + 1$  null links which proves the 2nd case.

Since both cases hold  $P(k + 1)$  is true. We conclude by the principle of mathematical induction the result holds for all integers  $n \geq 1$ .

**10.** A full node for a binary tree is one that has two children. Prove that the number of full nodes plus one equals the number of leaves on a binary tree.

**11.** Prove or disprove: deleting keys  $x$  and  $y$  from a BST is commutative. In other words, it does not matter which order the keys are deleted. the final trees will be identical. If true, provide a proof. If false, provide a counterexample.