

1. Compare the growth of $f(n) = 2^{\log_2(n^2)} + 4^{\log_2(\sqrt{n})} \cdot \log_2(n)$ and $g(n) = 6 \cdot n^2 + \log_2(n) + n^2$

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} &= \lim_{n \rightarrow \infty} \frac{2^{\log_2(n^2)} + 4^{\log_2(\sqrt{n})} \cdot \log_2(n)}{6 \cdot n^2 + \log_2(n) \cdot n^2} \\ &= \lim_{n \rightarrow \infty} \frac{n^2 \cdot \log_2(n)}{6 \cdot n^2 + \log_2(n) \cdot n^2} \\ &= c \end{aligned}$$

where $c \in \mathbb{R}$, hence $f(n) = \Theta(g(n))$.

2. Compare the growth of $f(n) = n^{1.06} + 6 \cdot n \cdot \log^5(n)$ and $g(n) = 2 \cdot n \cdot \log^{10}(n)$

$$f(n) = \omega(g(n))$$

3. Compare the growth of $f(n) = n^{2.01}$ and $g(n) = n^2 \cdot \log^{50}(n)$

$$f(n) = \omega(g(n))$$

4. Compare the growth of

$$\begin{aligned} f(n) &= n^{\log_2^2(n)} \\ g(n) &= \log_2^n(n) \end{aligned}$$

$$f(n) = o(g(n))$$

5. Compare the growth of

$$\begin{aligned} f(n) &= 8 \cdot n^{2+\sin(n)} \\ g(n) &= n \end{aligned}$$

$$f(n) = \Omega(g(n))$$

6. Prove $\forall k > 0, \epsilon > 0 \implies \log^k(n) = o(n^\epsilon)$.

We prove this by induction on k , let ϵ be a real number greater than 0. When $k = 1$ we have

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{\log(n)}{n^\epsilon} &= \lim_{n \rightarrow \infty} \frac{\frac{1}{n \cdot \ln(2)}}{\epsilon \cdot n^{\epsilon-1}} \\ &= \lim_{n \rightarrow \infty} \frac{1}{\epsilon \cdot \ln(2) \cdot n^\epsilon} \\ &= 0 \end{aligned}$$

Assume that the result holds for all integers n less than or equal to some integer k , we show it must also hold for the $k + 1$ integer.

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{\log^{k+1}(n)}{n^\epsilon} &= \lim_{n \rightarrow \infty} \frac{(k+1) \cdot \log^k(n) \cdot \frac{1}{n \cdot \ln(2)}}{\epsilon \cdot n^{\epsilon-1}} \\ &= \lim_{n \rightarrow \infty} \frac{(k+1) \cdot \log^k(n)}{\epsilon \cdot \ln(2) \cdot n^\epsilon} \\ &= 0 \end{aligned}$$

which is what we wanted to show. Hence by the principle of mathematical induction the result holds for all $n > 0$ and $\epsilon > 0$.

7. Prove $f(n) + g(n) = \Theta(\max(f(n), g(n)))$

$$\max(f(n), g(n)) \leq f(n) + g(n) \leq 2 \cdot \max(f(n), g(n))$$

We conclude $f(n) + g(n) = \Theta(\max(f(n), g(n)))$.