



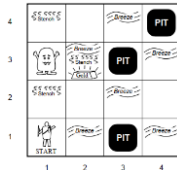
Logical Agents

ARTIFICIAL INTELLIGENCE
JUCHEOL MOON

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Wumpus World description

- Performance measure
 - gold +1000, death -1000
 - 1 per step, -10 for using the arrow
- Actuators
 - Up, Down, Left, Right move
 - No diagonal movement
 - Grab, Release, Shoot
- Sensors
 - Breeze, Glitter, Smell

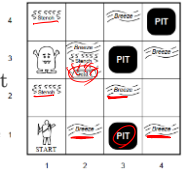


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Wumpus World description

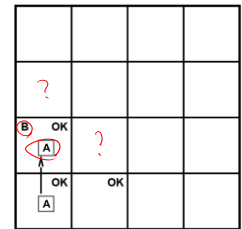
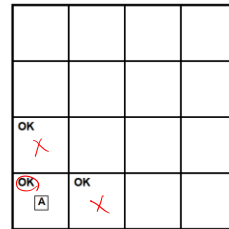
- Environment
 - Squares adjacent to wumpus are smelly
 - Squares adjacent to pit are breezy
 - Glitter iff gold is in the same square
 - Shooting kills wumpus if you are facing it
 - Shooting uses up the only arrow
 - Grabbing picks up gold if in same square
 - Releasing drops the gold in same square



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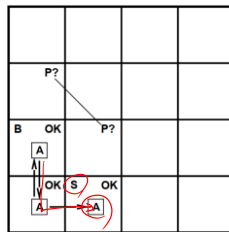
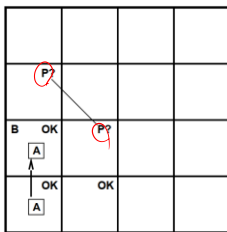
Exploring a wumpus world



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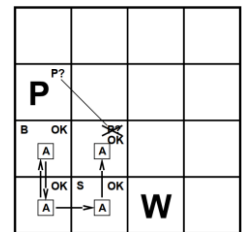
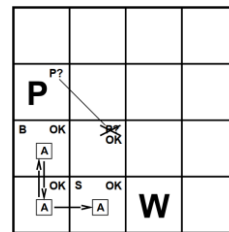
Exploring a wumpus world



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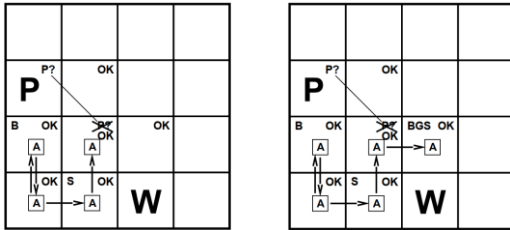
Exploring a wumpus world



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Exploring a wumpus world

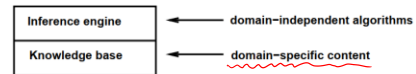


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Knowledge bases

- Knowledge base
 - set of sentences in a formal language
- Declarative approach to building an agent
 - Tell it what it needs to know
- Then it can Ask itself what to do
 - answers should follow from the KB



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The language of arithmetic

- Logics are formal languages for representing information such that conclusions can be drawn
- $x + 2 \geq y$ is a sentence
- $x2 + y >$ is not a sentence
- $x + 2 \geq y$ is true / false
 - in a world where $x = 7; y = 1$
- $x + 2 \geq y$ is true / false
 - in a world where $x = 0; y = 6$

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Entailment

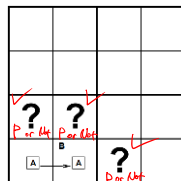
- Entailment means that one thing follows from another:
 - $KB \models \alpha$
 - Knowledge base KB entails sentence α if and only if α is true in all worlds where KB is true
 - E.g., $x + y = 4$ entails $4 = x + y$

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Entailment in the wumpus world

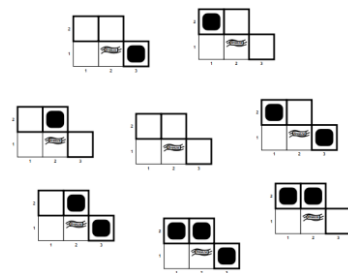
- Situation after detecting nothing in [1,1],
 - moving right, breeze in [2,1]
- Consider possible models for ?s
 - In terms of Pit
 - Ignore observations
 - They do not need be true
- 3 Boolean choices
 - \rightarrow 2 possible models



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Wumpus models

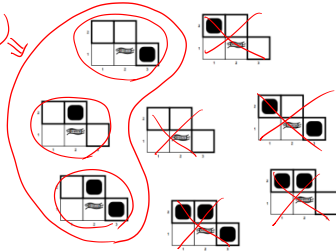


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Wumpus models

- KB = wumpus-world rules + observations
- nothing in [1,1]
- breeze in [2,1]



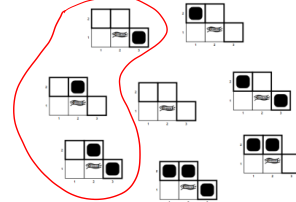
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Wumpus models

$KB \models \alpha$ if and only if α is true in all worlds where KB is true

- KB = wumpus-world rules + observations
- $\alpha_1 = "[1,2] \text{ is safe}"$, $KB \models \alpha_1$? Yes



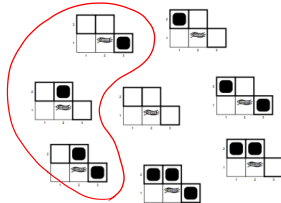
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Wumpus models

$KB \models \alpha$ if and only if α is true in all worlds where KB is true

- KB = wumpus-world rules + observations
- $\alpha_2 = "[2,2] \text{ is safe}"$, $KB \models \alpha_2$? No

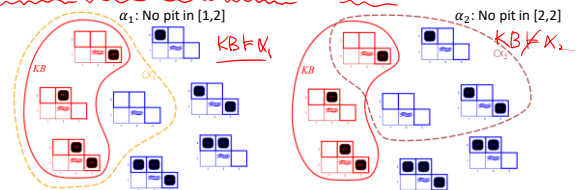


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Models

- We say M is a model of a sentence α if α is true in M
- $M(\alpha)$ is the set of all models of α
- $KB \models \alpha$ if and only if $M(KB) \subseteq M(\alpha)$



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Propositional logic: Syntax

- Propositional logic is the simplest logic
 - illustrates basic ideas
- The proposition symbols P_1, P_2 are sentences
- If P is a sentence, $\neg P$ is a sentence
 - negation

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Propositional logic: Syntax

- If P and Q are sentences, $P \wedge Q$ is a sentence
 - conjunction
- If P and Q are sentences, $P \vee Q$ is a sentence
 - disjunction
- If P and Q are sentences, $P \Rightarrow Q$ is a sentence
 - implication
- If P and Q are sentences, $P \Leftrightarrow Q$ is a sentence
 - biconditional

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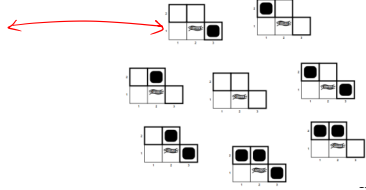
Wumpus world sentences

- Let $P_{i,j}$ be true if there is a pit in $[i,j]$.
- Let $B_{i,j}$ be true if there is a breeze in $[i,j]$.

$$\neg P_{1,1} = \text{true}$$

$$B_{2,1} = \text{true}$$

$$P_{3,1} = \text{true}$$



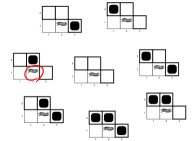
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Wumpus world sentences

- Let $P_{i,j}$ be true if there is a pit in $[i,j]$.
- Let $B_{i,j}$ be true if there is a breeze in $[i,j]$.
- Pits cause breezes in adjacent squares
 - A square is breezy if and only if there is an adjacent pit.

$$B_{2,1} \Leftrightarrow P_{3,1} \vee P_{2,2} \vee P_{1,1}$$

$$B_{1,1} \Leftrightarrow P_{1,2} \vee P_{2,1}$$



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Satisfiability

- A sentence is **satisfiable** if it is true in **some** model

$$(x \vee y)$$

$$(x \vee y) \wedge z$$

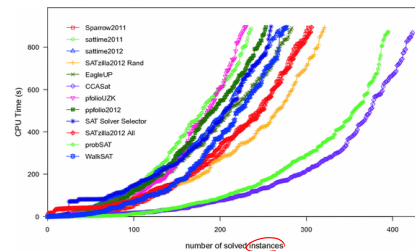
- A sentence is **unsatisfiable** if it is true in **no** models

$$x \wedge \neg x$$

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Satisfiability

- Is $B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$ satisfiable?



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Conversion to CNF

- Is $B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$ satisfiable?

- Replacing $\alpha \Leftrightarrow \beta$ with $(\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha)$

$$[B_{1,1} \Rightarrow (P_{1,2} \vee P_{2,1})] \wedge [(P_{1,2} \vee P_{2,1}) \Rightarrow B_{1,1}]$$

- Replacing $\alpha \Rightarrow \beta$ with $\neg \alpha \vee \beta$

$$[\neg B_{1,1} \vee P_{1,2} \vee P_{2,1}] \wedge [\neg (P_{1,2} \vee P_{2,1}) \vee B_{1,1}]$$

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Conversion to CNF

- Is $B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$ satisfiable?

- Move \neg inwards using de Morgan's rules and double negation

$$(\neg B_{1,1} \vee P_{1,2} \vee P_{2,1}) \wedge [(\neg P_{1,2} \wedge \neg P_{2,1}) \vee B_{1,1}]$$

- Apply distributivity law (\vee over \wedge) and flatten

$$(\neg B_{1,1} \vee P_{1,2} \vee P_{2,1}) \wedge (\neg P_{1,2} \vee B_{1,1}) \wedge (\neg P_{2,1} \vee B_{1,1})$$

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