1.	Suppose there is a penalty for 1 for each gap and a penalty of	2
for	matching two different symbols in a column. What is the N	W
scc	ore of the strings AGTACG and ACATAG?	

- 1. 3
- 2. 4
- 3. 5
- 4. 6

## A - GTCGACATA - G

(2) is the correct answer. We have an upperbound of 8 if we compute the NW score of the original strings. The strings are of the same length so they must have an even NW score because inserting a gap in one of the strings will require we insert a gap in the other. We will have at least one mismatch will we give us a lowerbound of 4.

- **2.** Let  $X = x_1, x_2, \ldots, x_m$  and  $Y = y_1, y_2, \ldots, y_n$  be two input strings, with each symbol  $x_i$  or  $y_j$  in  $\{A, C, G, T\}$ . How many relevant possibilities are there for the contents of the final column of an optimal alignment?
  - 1. 2
  - 2. 3
  - 3. 4
  - 4.  $m \cdot n$

(2) is the correct answer. We choose for the  $x_m$  character to be a gap or the  $x_m$  character. We again can choose for the  $y_n$  character to be a gap or the  $y_n$  character. Making both  $x_m$  and  $y_n$  gaps would be redundant and we could have a better solution by having either  $x_m$  or  $y_n$  as a gap. Thus we have three options.  $x_m$  and  $y_m$  staying the same.  $x_m$  staying the same and matched with a gap.  $y_n$  staying the same and matched with a gap.

**3.** Suppose one of the two input strings is empty. What is the NW score of X and Y?

- 1. 0
- 2.  $\alpha_{gap} \cdot (length \ of \ X)$

- $3. +\infty$
- 4. undefined
- (2) is the correct answer. We make Y the "gap" string. which would result in a penalty of the length of X.
- **4.** Consider the following two search trees that store objects with key 1, 2 and 3: and the search frequency 1:.8, 2:.1, 3:.1. What are the average search times in the two trees, respectively?
  - 1. 1.9 and 1.2
  - 2. 1.9 and 1.3
  - 3. 2 and 1
  - 4. 2 and 3

The average search time for the first tree is  $\sum_{n\in N}p(n)\cdot[depth\:o\:f\:n\:in\:tree+1]=.8*2+.1*1+.1*3=1.9$ . The average search time for the second tree is  $\sum_{n\in N}p(n)\cdot[depth\:o\:f\:n\:in\:tree+1]=.8*1+.1*2+.1*3=1.2$ . Hence (2) is the correct answer.