1. Euler function

$$\varphi(n)=\operatorname{len}(\{1,2,3..n\},\operatorname{gcd}(k,n)=1)$$

Phi is multiplicative function

$$\varphi(ab) = \varphi(a)\varphi(b)$$

 $\varphi(n)=p_1^{k_1-1}(p_1-1)p_2^{k_2-1}(p_2-1)...$ Where p is prime number from factorization n.

Example

$$\varphi(54) = \varphi(2*3^3) = \varphi(2)*\varphi(3^3)$$

Euler's theorem If a and p is coprime, than $a^{\varphi(n)} \equiv 1 \mod n$.

2. Modular arithmetics

2.1. Congruence

We say that 3 is congruent to 15 by modulo 12, written $15 \equiv 3 \pmod{12}$

Coprime two integers GCD is 1.

2.2. Fermat's little theorem

Special case of euler theorem.

Theorem If p is a prime number, then for any integer a, the number $a^p - a$ is an integer multiple of p

$$a^p \equiv a \operatorname{mod} p$$

If a is coprime to p.

$$a^{p-1} \equiv 1 \bmod p$$

2.3. Primitive root modulo

power of g modulo p.

a|b a divides b=> b/a = 0

 $\forall a \in Z : \gcd(a, p) = 1, \exists n : g^n = a \longrightarrow g$ is primitive root modulo

Primitive root modulo n g is called primitive root modulo p if every a coprime number to n is congruent to a

N is not required to be prime. G is a *primitive root modulo*
$$n$$
 if and only if g is a generator of the multiplicative

group of integers modulo n.

2.4. P Group

2.4.1. How to check that group is cyclic 2.4.2. Theorem to check generator in p group

$\alpha \in \mathbb{Z}_p^*$ is a generator of \mathbb{Z}_p^* if and only if

For all primes q such that q|(p-1)

$$\alpha^{\frac{p-1}{q}} \neg \equiv 1 \bmod p$$

Task

Task find the all generators of Z_{11}^* Let's begin with p - 1 = 10, 10 = 2*5.

Generator check condition for each divider of (p-1):

• $\alpha^5 \neg \equiv 1 \mod 11$

- $\alpha^2 \neg \equiv 1 \mod 11$
- Solution is to check each element in group to match conditions.

2.4.3. How to count generators in group

Theorem let p be prime, that Z_p^* contains exactly $\varphi(p-1)$ generators.

2.4.4. How to find generator

2.4.5. Discrete logarithm in p

If $B \in \mathbb{Z}_p^*$, then $B = g^x$ for some unique $0 \le x \le p-2$. X is called the discrete logarithm of B to base g.

 $g^x = \mathbf{B} \Rightarrow \log_g \mathbf{B} = x \text{ in } Z$ **Problem** find the integer x, such that $\log_q \mathbf{B}$ in Z_p^*

The naive approach is exhaustive search: compute $g^x, g^2x, ...$ until B is obtained.

$$x \in \mathcal{G}^{-1}, \mathcal{G}^{-1}$$

 $Z_n^* = \{k \in \{1, ..., n\}/\gcd(k, n) = 1\}$

 $\operatorname{len}(Z_n^*) = \varphi(n)$

2.5.1. Theorem to check generator in n group For $n \ge 1$, we consider Z_n^*

$$Z_n^*$$
 is cyclic(has at least one generator) when:

2. $n = p^x, x \in \{1, 2...\}$

3. $n = 2p^x, x \in \{1, 2...\}$

1. n=2 or 4

2.5. N Group

Theorem to check generator Assume Z_n^* is cyclic. $\alpha \in Z_n^*$ is a generator if and only if

For each prime p divisor of $\varphi(n)$

2.5.2. How to count generators in group

Theorem if
$$Z_n^*$$
 is cyclic, then it has $\pi(\pi(n))$ generators.

If $B \in \mathbb{Z}_n^*$, then $B = g^x$ for some unique $0 \le x \le p-2$. X is called the discrete logarithm of B to base g.

 $\alpha^{\frac{\varphi(n)}{p}} \neg \equiv 1 \bmod n$

2.5.3. Discrete logarithm in n

2.5.2. How to count generators in group

Find $\log_{13} 47$ in Z_{50}^*

2. Check g 13 is generator. (requires find $\varphi(n)$)

3. Algorithms for computing discrete algorithms 1. Brute force

1. Check Z_{50}^* is cyclic(e.g. has generators)

3.1. Some equations

$$y_1 = (a^{x_1} \mod p)^{x_2} \mod p = a^{x_1 x_2} \mod p$$

 $3B \mod 13 = 1 \Rightarrow 3B \equiv 1 \mod 13$

 $(ab) \bmod m = [(a \bmod m)(b \bmod m)] \bmod m$

Some modular arithmetics to proof:

Continue:

3.2. Core Equation for DH

$$y_1 = (a^{x_1} \bmod p)^{x_2} \bmod p = \big((a \bmod p)^{x_1} \bmod p\big)^{x_2} \bmod p =$$
 Reduce extra modulo due to modulo properties:

 $(a \bmod p)^{x_1} \bmod p = a^{x_1} \bmod p$

$$= \left(a \operatorname{mod} p\right)^{x_1 x_2} \operatorname{mod} p$$

 $=a^{x_1x_2} \bmod p$

3.3. Factorization problem For example RSA relies on difficulty of factoring the product of two large prime numbers.

But for this need to determine or find large prime number.

2. Probabilistic tests(all primes + some non primes - never FN, but sometimes FP)

1. Simple methods(advanced brute force, without 2,3 and maybe some memoization,etc)

Determine if number is prime:

Example

- 2. Shank's baby-step giant-step method