

## 1. Intro

Public key cryptography - is a method of encrypting, which allows individuals to securely communicate without sharing the secret key.

## 2. Rsa

Asymmetric encryption.

Trapdoor one way function - function that can be easily computed in one way, and hard in the inverse without special(secret) information(trapdoor). In this case forward -> encrypt. Inverse -> decrypt.

Private key - is key to decrypt message. Public key - is key to encrypt message.

For signatures private key enables to sign and public to verify this sign validity.

Arithmetic functions, that considered as a one way(not proofed)

1. Multiplication and factorization
2.  $x^y \bmod n = z \Rightarrow \text{find } y$
3.  $x^y \bmod n = z \Rightarrow \text{find } x$
4.  $x^2 \bmod n$ .  $n$  - not prime,  $\text{jacobi}(z/n)=1 \Rightarrow \text{find } x$
5.  $g^{ab} \bmod p \Rightarrow \text{find } a$

## 3. Group

Group is set of elements, that are related to each other according to certain well-defined rules.

$Z_p^*$  - is a group with nonzero integers between 1 and  $p-1$  modulo some prime number  $p$ .

### 3.1. Axioms

Operation for example is a multiplication.

1. Closure -  $a, b \in G, a * b = c \rightarrow c \in G$
2. Associativity  $a * (b * c) = (a * b) * c$
3. Identity existence  $a * 1 = a$
4. Inverse existence  $a * b = 1$

A group is commutative if  $a * b = b * a$

A group is cyclic if

$$g \in G, \forall x \in G, \exists n : g^n = x$$

## 4. Diffie-hellman

### 4.1. Discrete logarithm problem

Problem - find the  $y$ , where  $g$  and  $x$  are provided,  $g^y = x$ .

### 4.2. Core Equation

$$y_1 = (a^{x_1} \bmod p)^{x_2} \bmod p = a^{x_1 x_2} \bmod p$$

Some modular arithmetics to proof:

$$y_1 = (a^{x_1} \bmod p)^{x_2} \bmod p = ((a \bmod p)^{x_1} \bmod p)^{x_2} \bmod p =$$

Reduce extra modulo due to modulo properties:

$$(a \bmod p)^{x_1} \bmod p = a^{x_1} \bmod p$$

Continue:

$$\begin{aligned} &= (a \bmod p)^{x_1 x_2} \bmod p \\ &= a^{x_1 x_2} \bmod p \end{aligned}$$

### 4.3. Modular arithmetics

$$(ab) \bmod m = [(a \bmod m)(b \bmod m)] \bmod m$$