

Estimating the Impact of BGP Prefix Hijacking

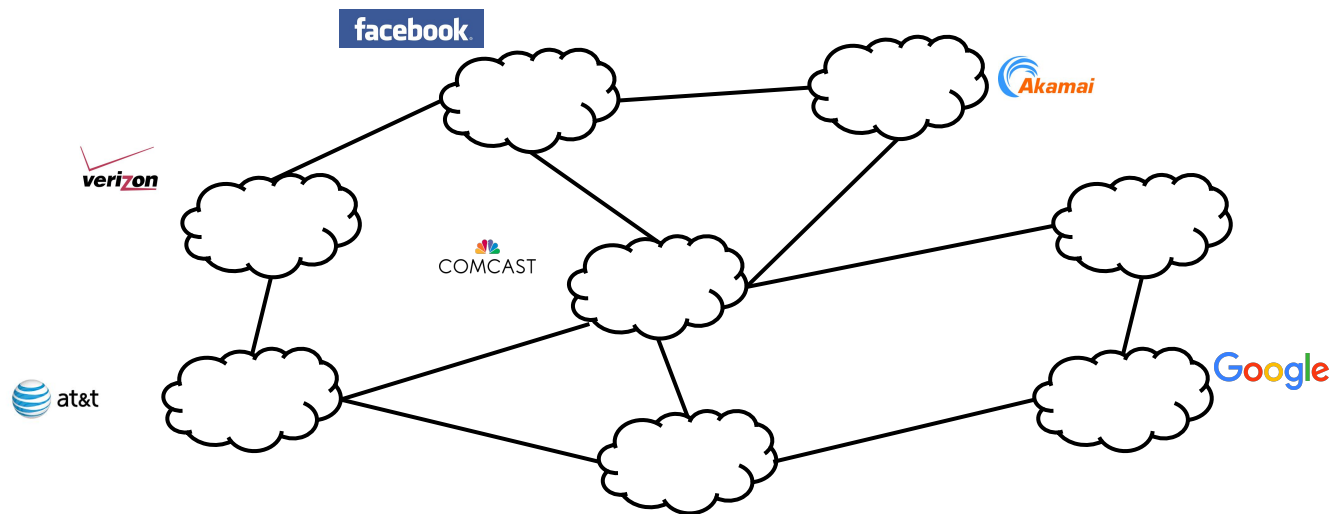
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² Institute of Computer Science, FORTH, Greece

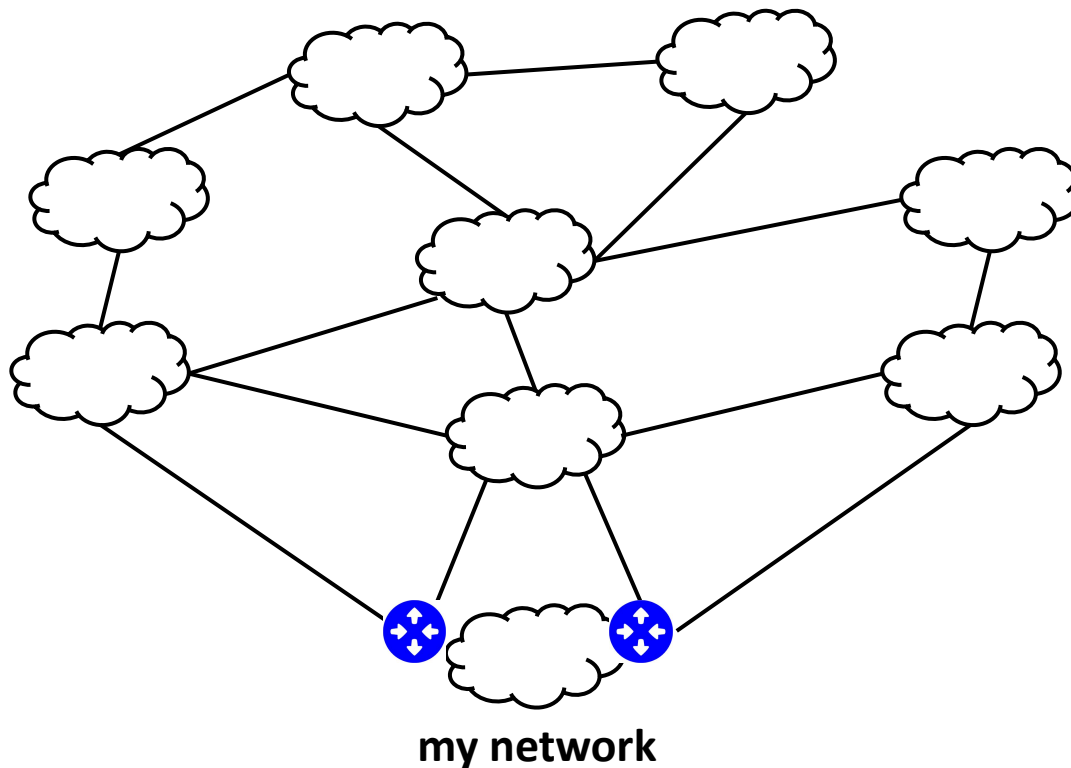
³ Computer Science Dept., University of Crete, Greece

The Internet

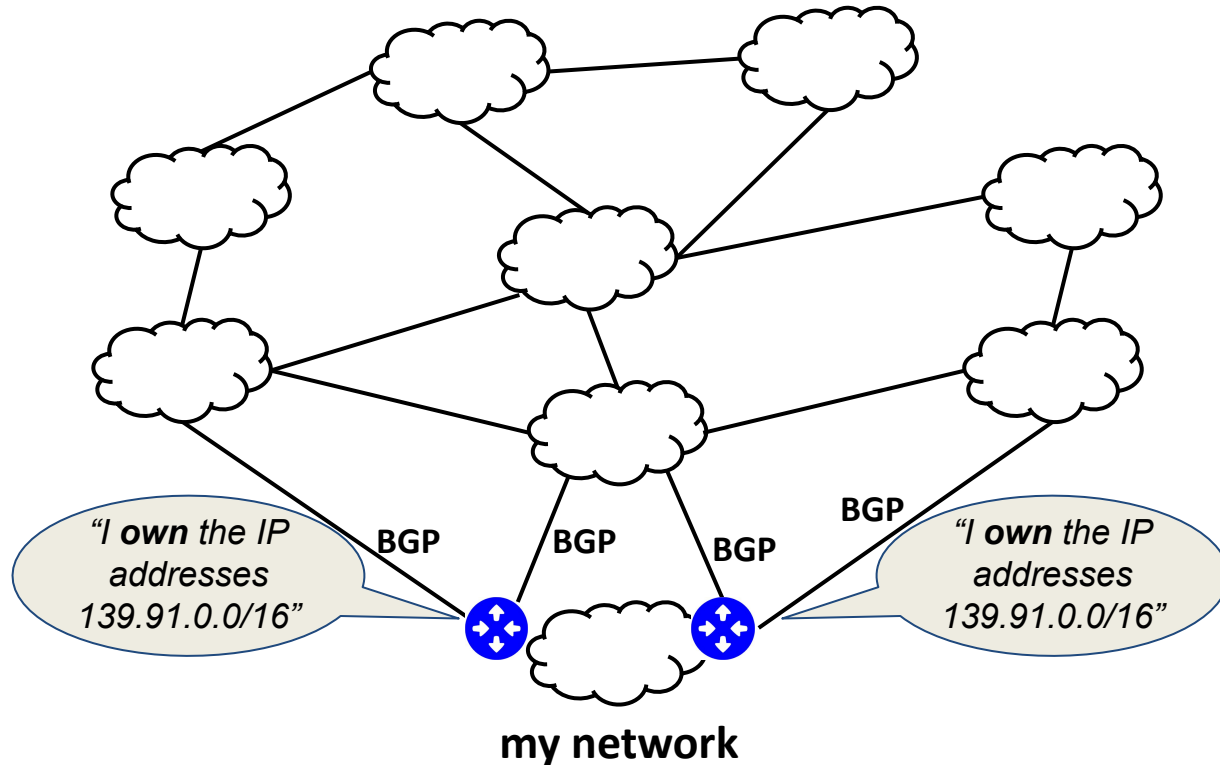


- The Internet is a network of networks or “**Autonomous Systems (AS)**”
- today **~70k** ASes

The Internet

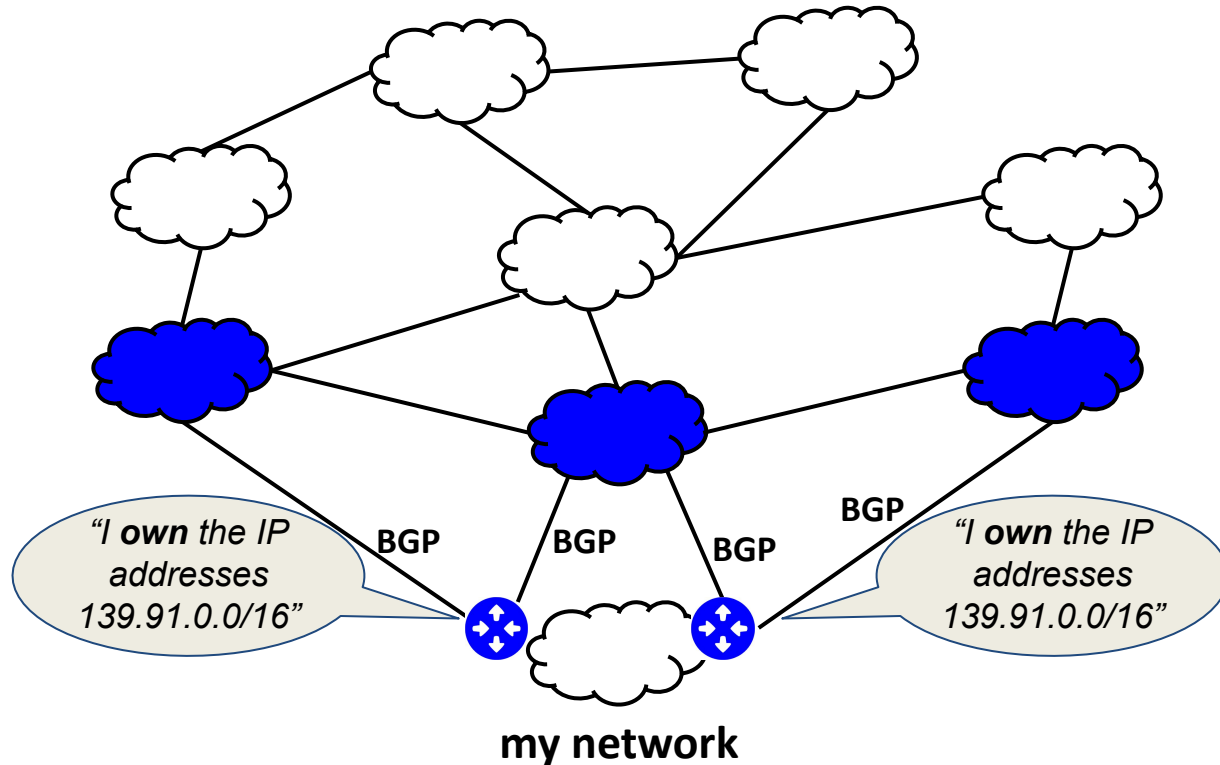


Internet routing



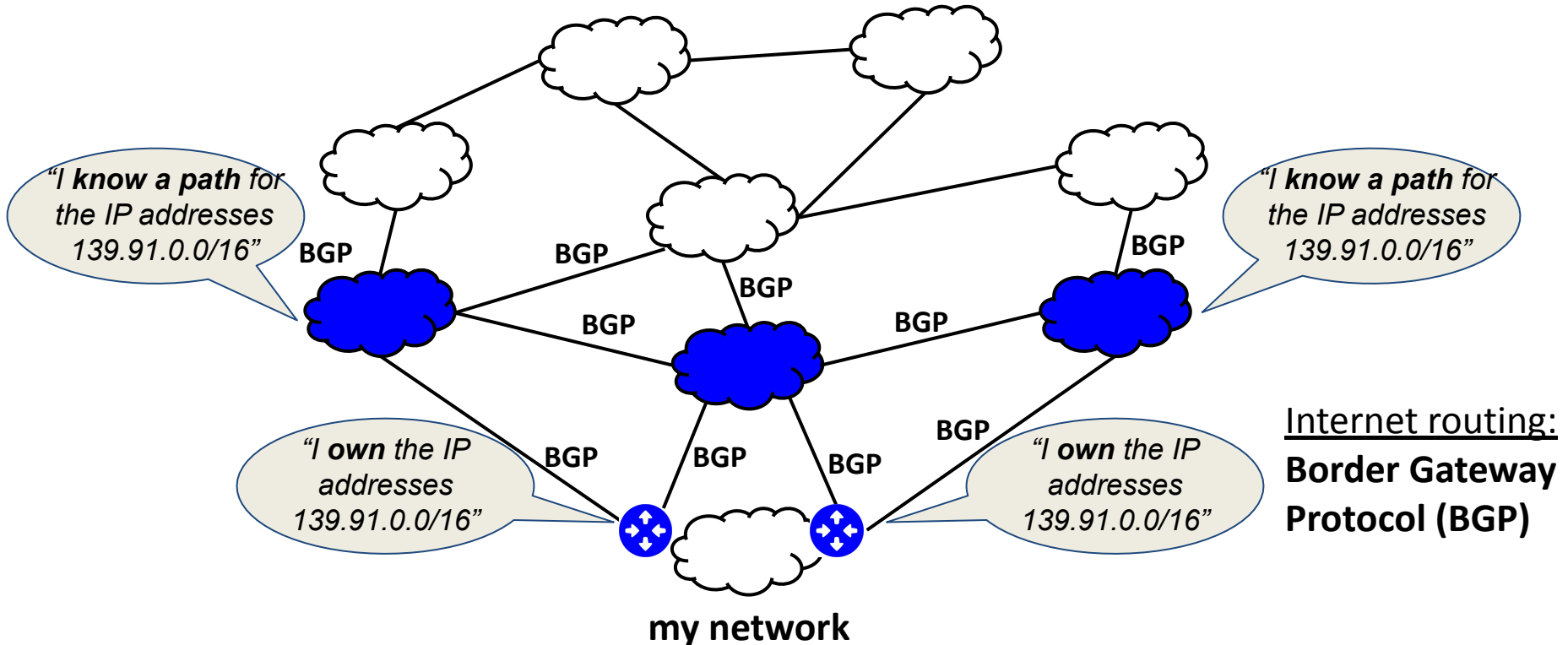
Internet routing:
**Border Gateway
Protocol (BGP)**

Internet routing

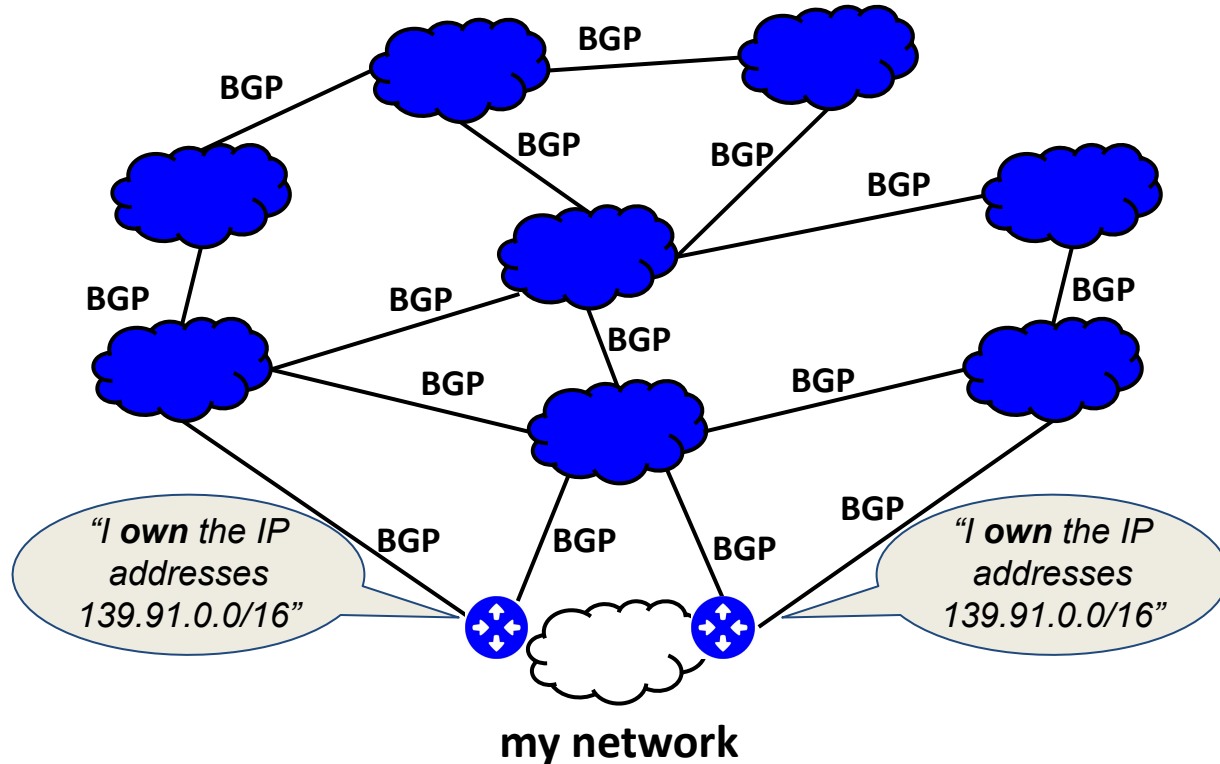


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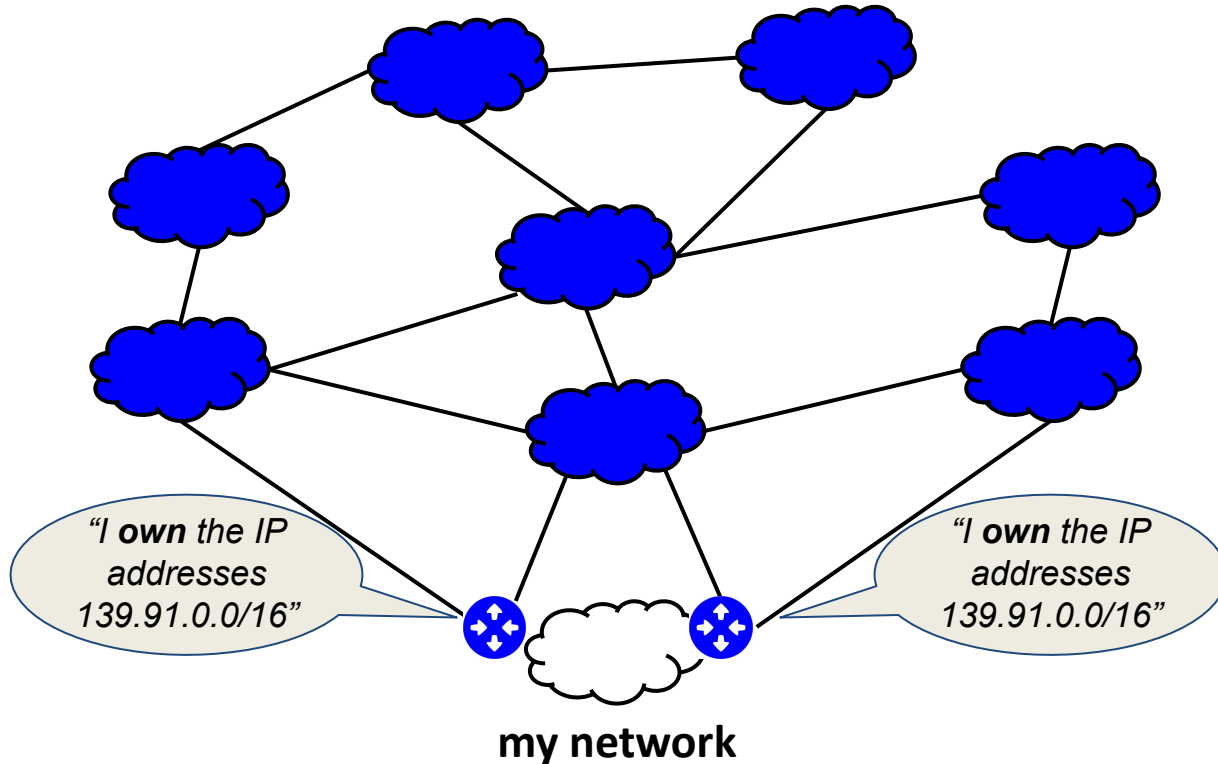


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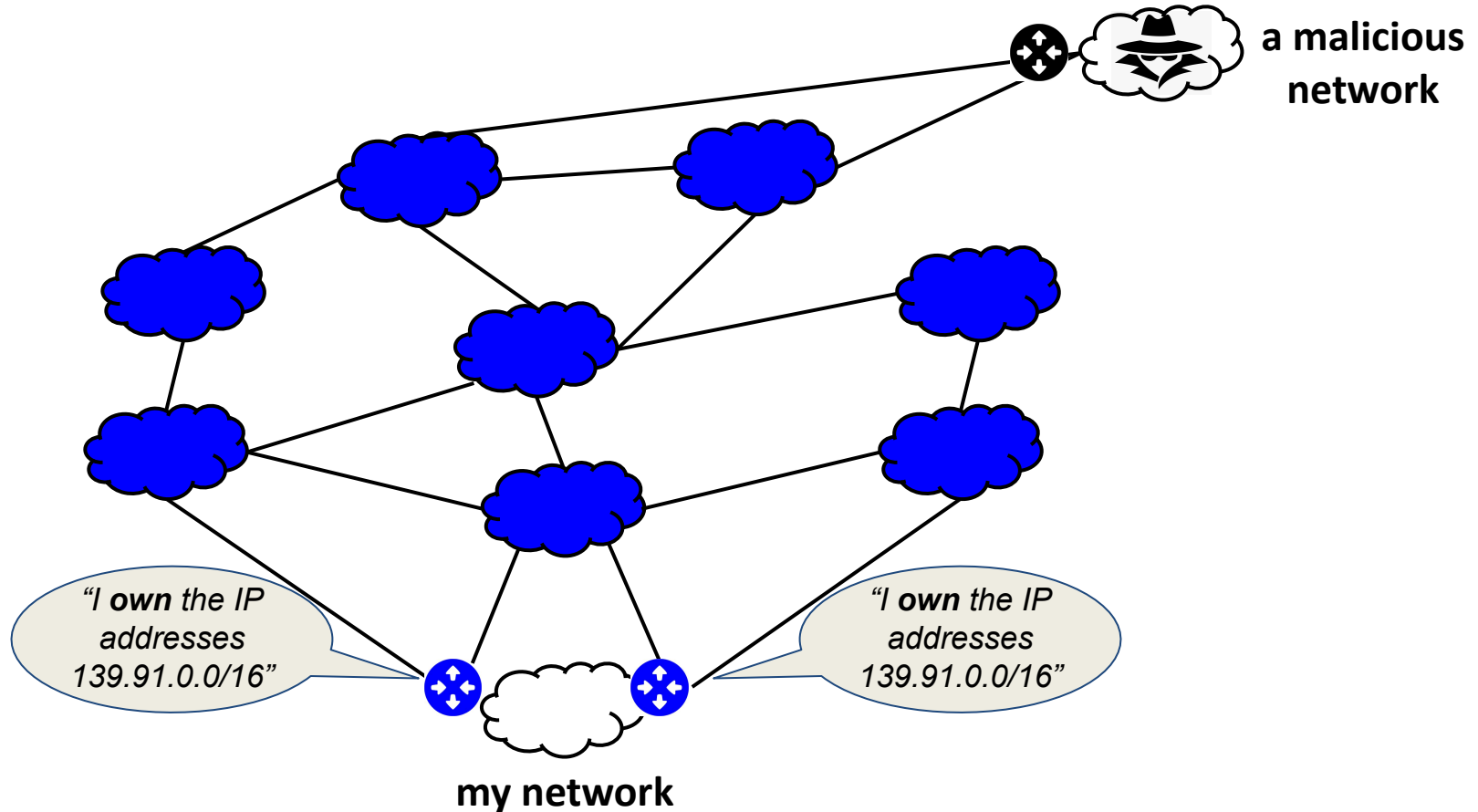


Internet routing:
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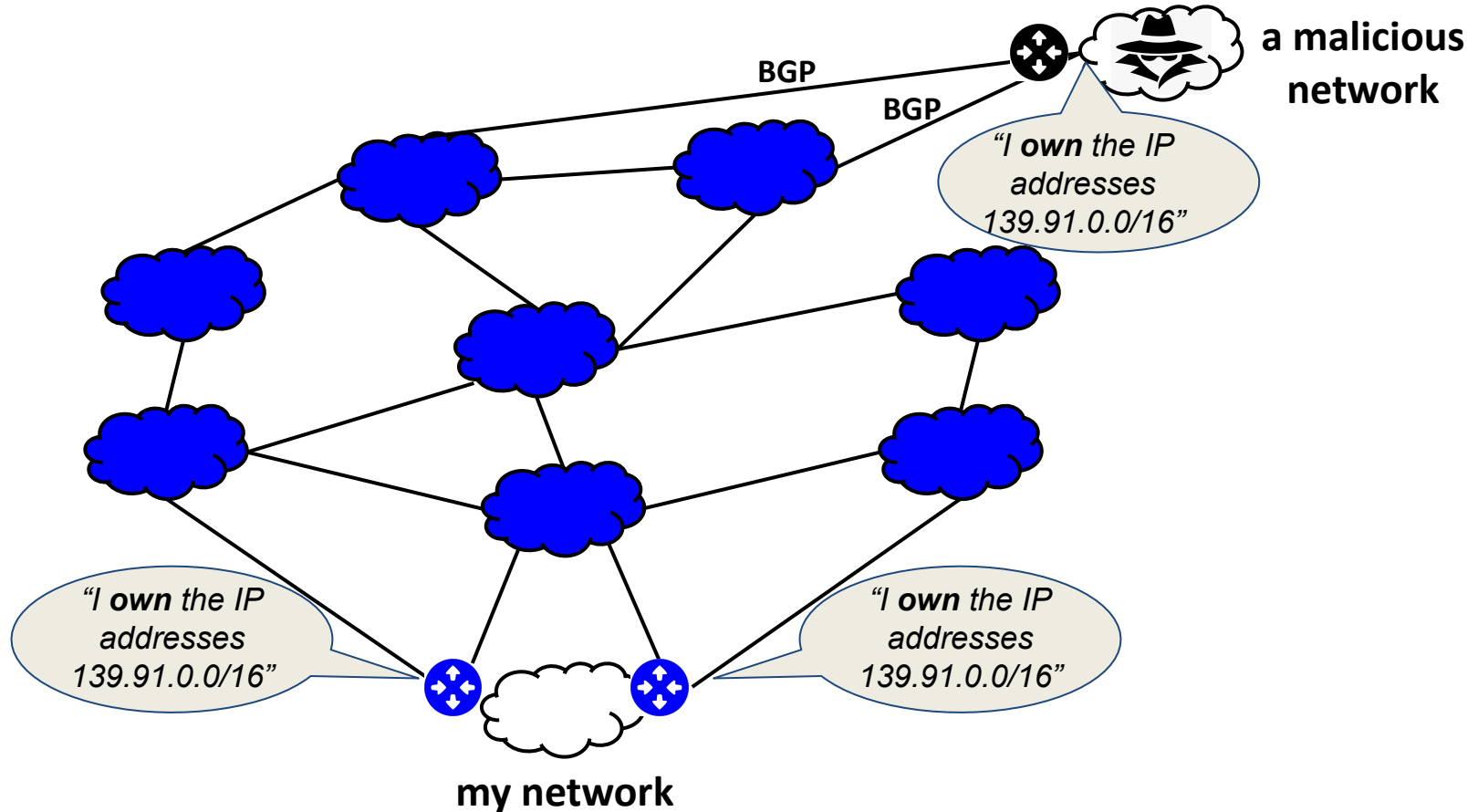
BGP prefix hijacking



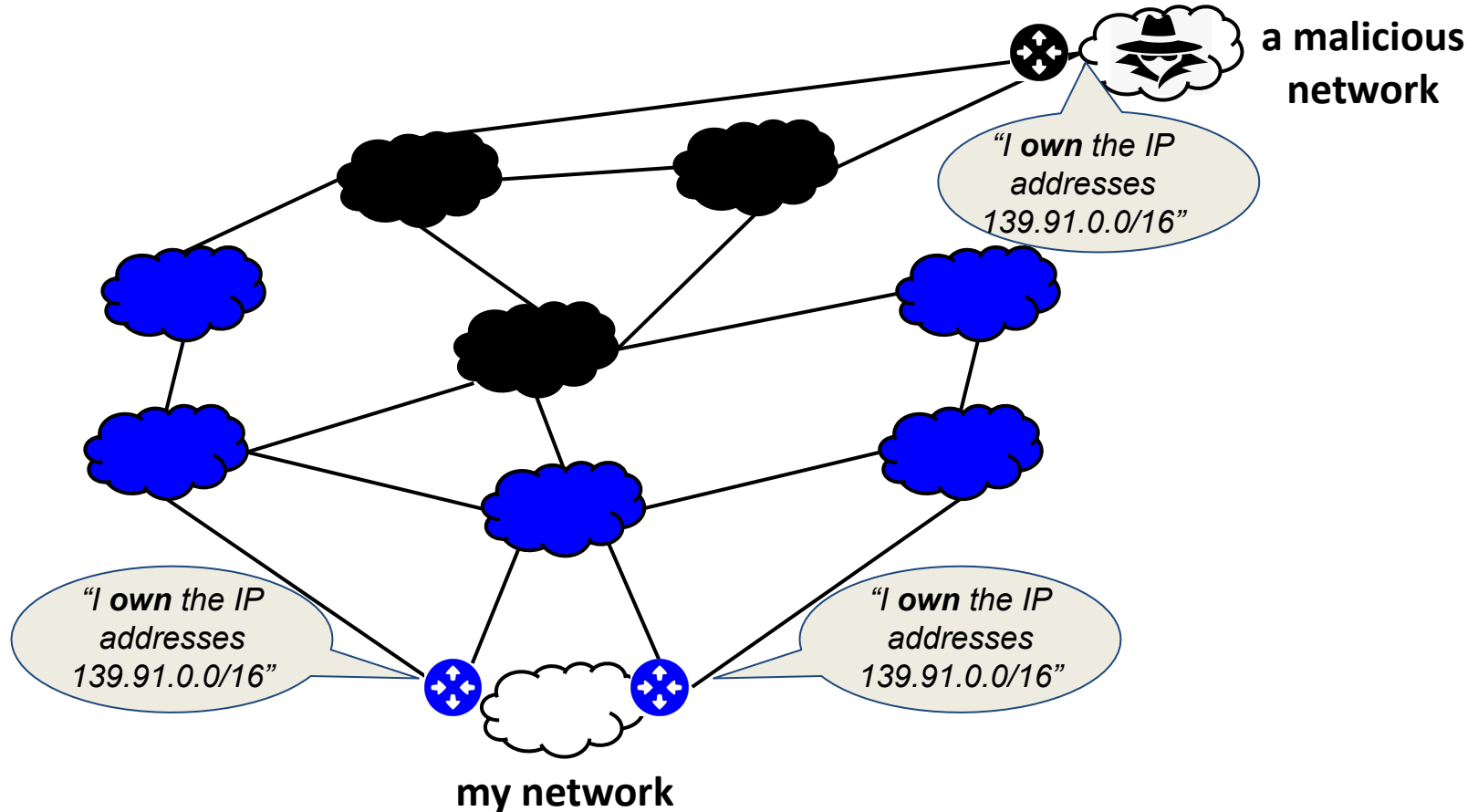
BGP prefix hijacking



BGP prefix hijacking



BGP prefix hijacking



BGP prefix hijacking

- **It's an important problem today! a few facts:**
 - ~2500 (reported) prefix hijacking events in 2020 ¹
 - examples of affected networks: Google, Amazon, Akamai, Visa, etc.
- **service outages & traffic interception**
 - can last for hours ²
 - can cost \$100k (or more) per minute!
- **no effective (proactive) defence**
 - RPKI: limited adoption & efficiency ^{2,3}
 - defences based upon detection & countermeasures ²

¹ APNIC, "BGP, RPKI, and MANRS: 2020 in review", Feb 2021, <https://blog.apnic.net/2021/02/05/bgp-rpki-and-manrs-2020-in-review/>

² P. Sermpezis, et. al., "A survey among Network Operators on BGP Prefix Hijacking", in ACM SIGCOMM CCR, Jan 2018.

³ NIST RPKI Monitor, <https://rpki-monitorantd.nist.gov/>

How do we defend?

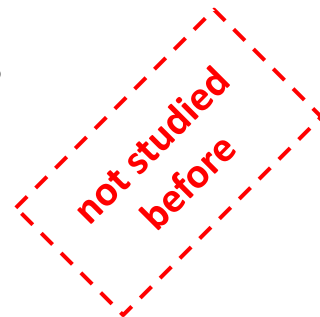
- **How do we defend against hijacks? → (mostly) reactively:**
 - *Step 1: detect the hijack*
 - *Step 2: proceed to mitigation action*
 - filtering, deaggregation, outsourcing (e.g., blackholing, anycast from large ISPs)
- **Detection**
 - ✓ a lot of research the last ~10 years
 - ✓ public monitoring infrastructure (RIPE RIS, RouteViews, etc.)
 - ✓ state-of-the-art: near real-time detection (in a few seconds) ¹
- **Mitigation**
 - ? different actions → different costs... *which one to choose?*
 - ? ok, I took an action... *was it effective? is the problem solved?*

→ **we need to know the impact of the hijack (before/after its mitigation) !!!**

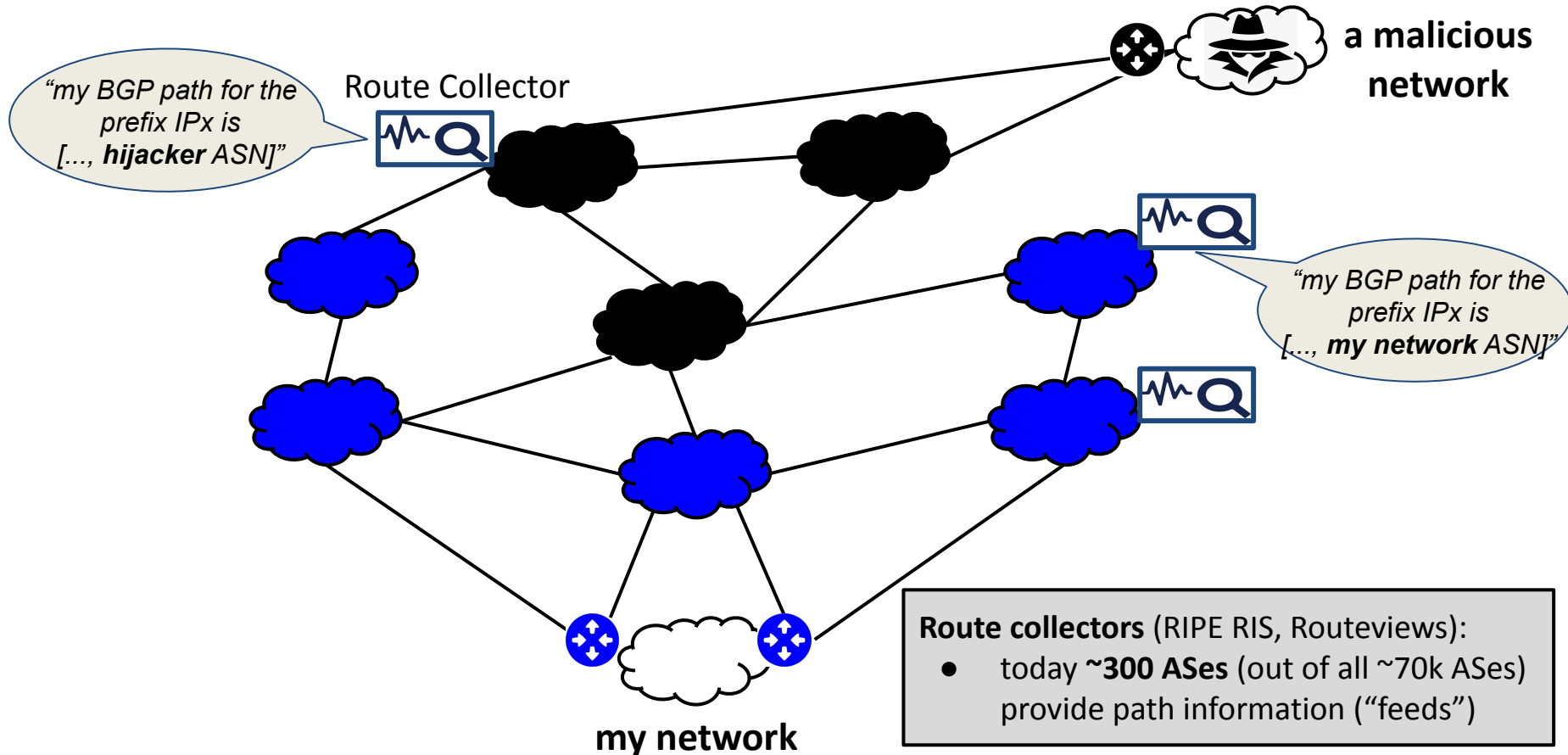
¹ ARTEMIS, open-source software, <https://bgpartemis.org/>

In this paper...

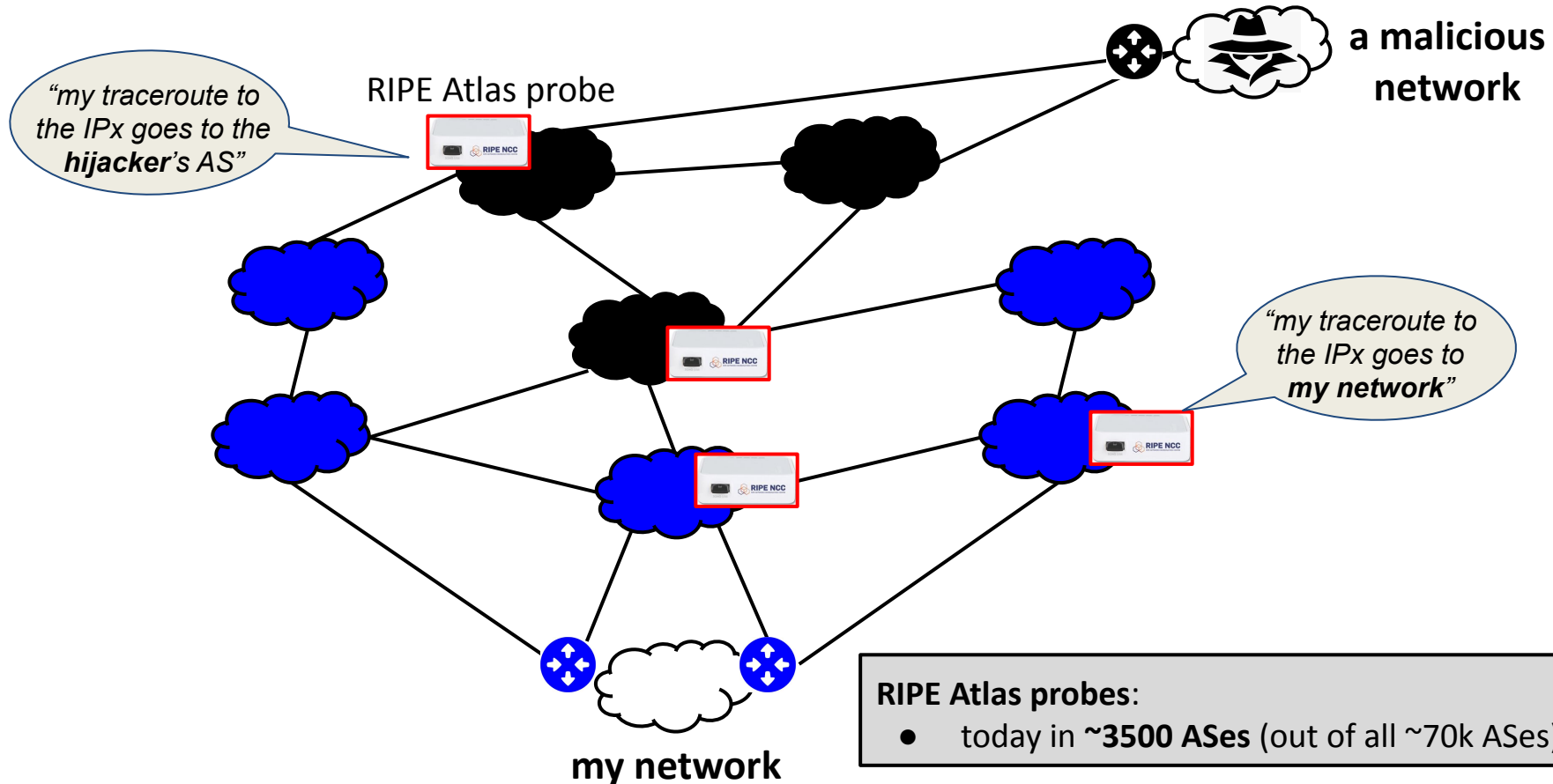
- **Goal** → Estimate the impact of an ongoing hijack through measurements
 - *sub-goal 1*: understand/characterise the estimation problem
 - types of measurements & public infrastructure
 - challenges & limitations
 - accuracy
 - *sub-goal 2*: design efficient estimation methodologies
 - with public infrastructure
 - without public infrastructure



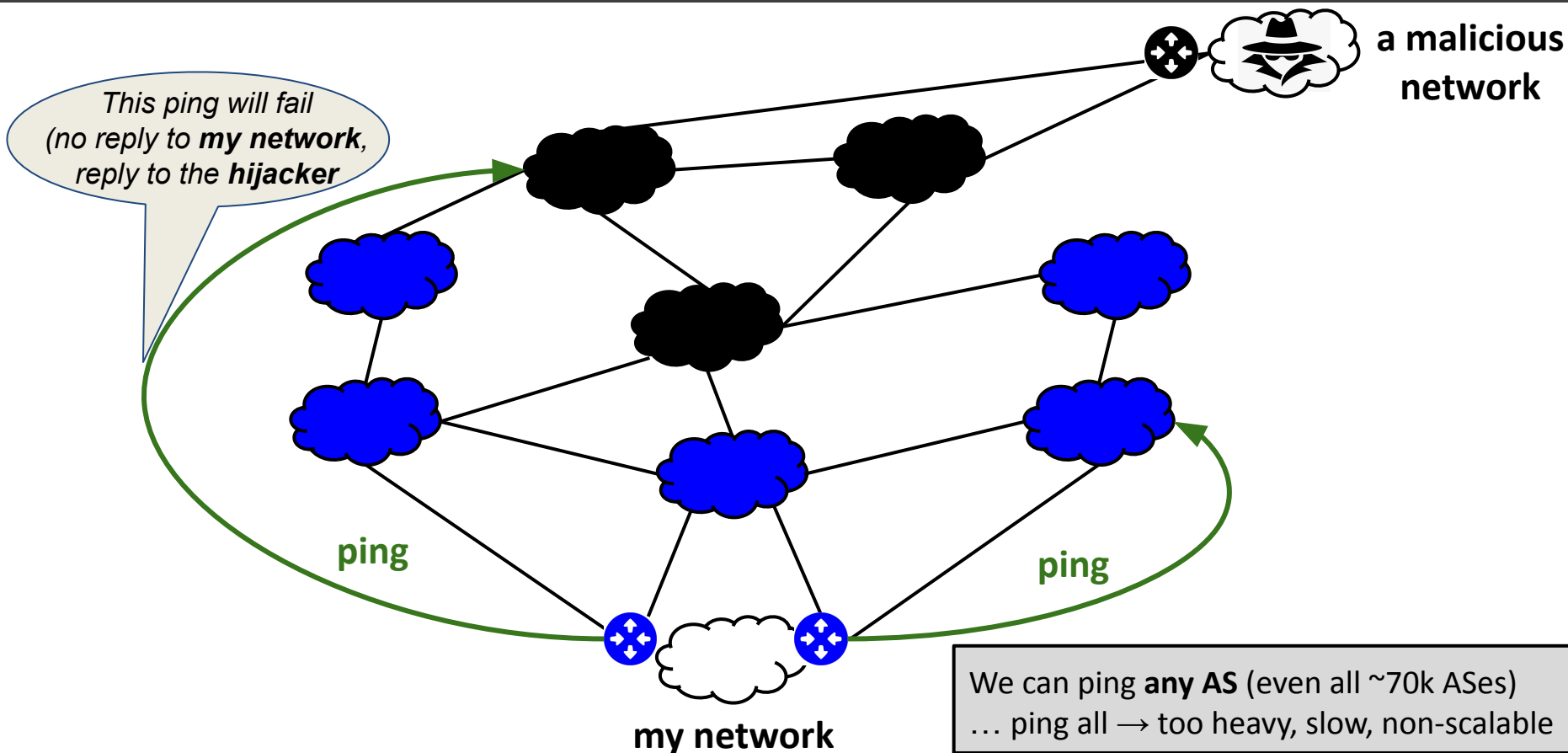
Measurements: BGP paths



Measurements: traceroutes



Measurements: pings



Hijack impact estimation with measurements

- Hijack impact == number of infected ASes

- **“infected AS”** == an AS that routes its traffic to the hijacker AS

$$\text{actual impact} = \frac{\text{\# infected ASes}}{\text{\# total ASes}}$$

- Measurements for hijack impact estimation

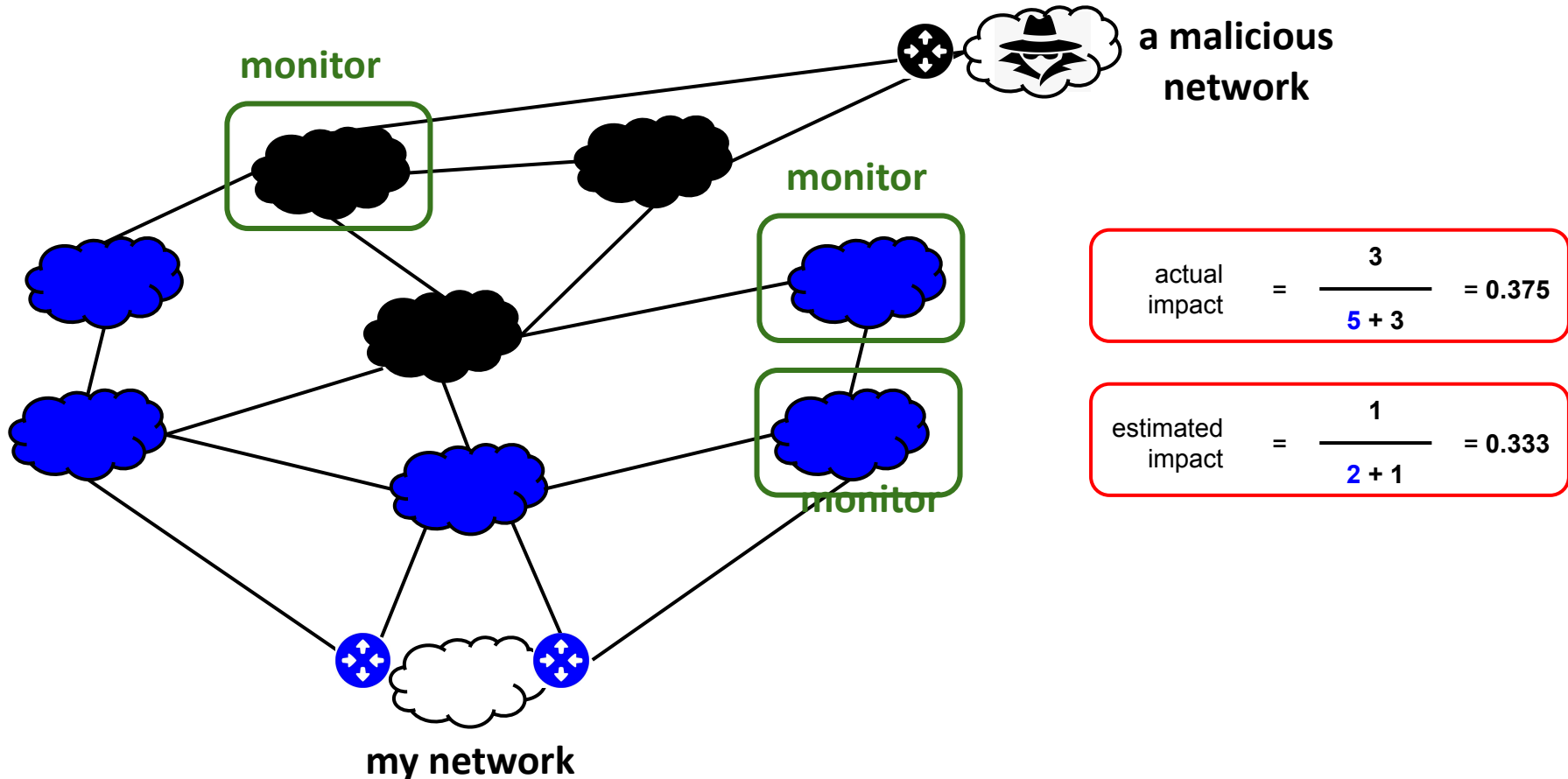
- measure some ASes
- measured AS == **“monitor”**
- any measurement type: BGP path (route collector), traceroute (RIPE Atlas probe), ping

- Estimate hijack impact

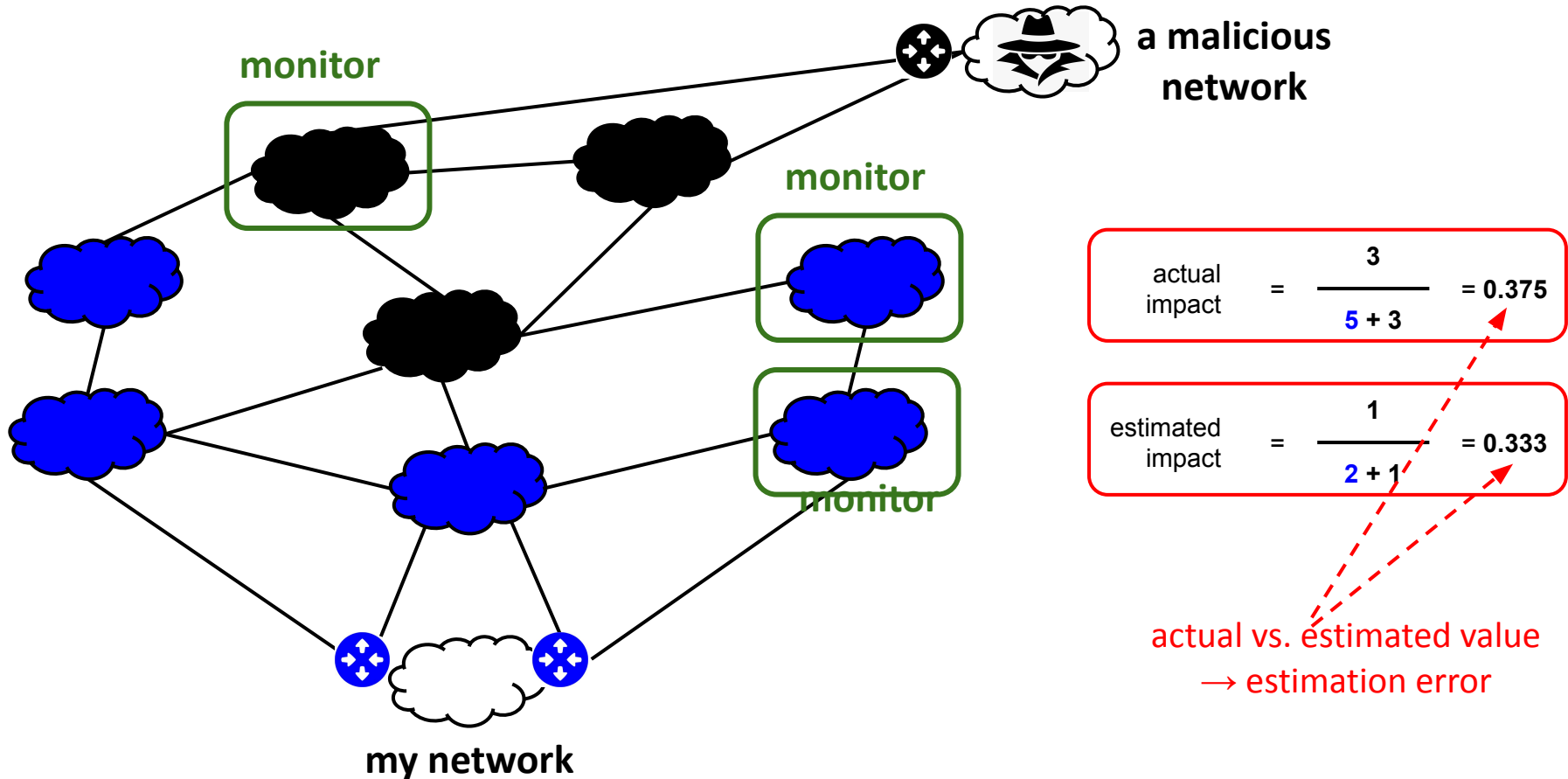
- from the number/percentage of “infected monitors”

$$\text{estimated impact} = \frac{\text{\# infected monitors}}{\text{\# total monitors}}$$

Impact estimation: an example



Impact estimation: an example



Goal 1: study the accuracy of estimation

- **Sampling in theory...**
 - The estimation error (RMSE) decreases with the number of samples/monitors (M)

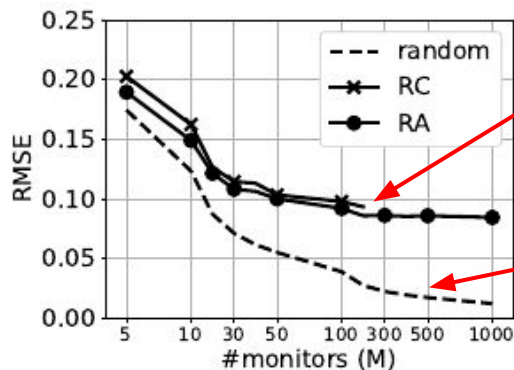
Theorem 1. Under a randomly selected set of monitors \mathcal{M} , the bias and root mean square error of NIE are given by

$$\text{Bias}_{NIE} = 0 \quad \text{RMSE}_{NIE} = \frac{1}{\sqrt{M}} \cdot c_I$$

where $c_I = \int_0^1 \sqrt{I \cdot (1 - I)} \cdot f(I) \cdot dI$, is a constant that depends on the impact distribution $f(I)$.

Goal 1: study the accuracy of estimation

- **Sampling in theory...**
 - The estimation error (RMSE) decreases with the number of samples/monitors (M)
- **Sampling in practice...**
 - with public infrastructure (Route Collectors, RIPE Atlas probes)



public
infrastructure

theory
(random
sampling)

Key findings:

- ▶ The error of public infrastructure does not decrease less than 9%
→ due to location bias; public infrastructure is not deployed uniformly around the world
- ▶ 50 samples from public monitors are enough
→ insights for lightweight measurements

Goal 1: study the accuracy of estimation

- **Sampling in theory...**
 - The estimation error (RMSE) decreases with the number of samples/monitors (M)
- **Sampling in practice...**
 - with public infrastructure (Route Collectors, RIPE Atlas probes)
 - what about ping measurements?

Goal 1: study the accuracy of estimation

- **Sampling in theory...**
 - The estimation error (RMSE) decreases with the number of samples/monitors (M)
- **Sampling in practice...**
 - with public infrastructure (Route Collectors, RIPE Atlas probes)
 - what about ping measurements?
 - ✓ we can have random sampling!
 - ✗ but... high measurement failures ($> 90\%$ non pingable IP addresses)

Theorem 2.
RMSE vs. failure probability p .

Key findings:

- ping measurements end-up being less accurate than public infrastructure (for $p > 20\%$)
- we would need at least $p < 10\%$

Goal 2: design accurate estimators

- Approach 1: based on **ping measurements**...
 - *Goal*: we need to decrease the failure probability p
- Approach 2: based on **public-infrastructure**...
 - *Goal*: we need to remove the measurement bias

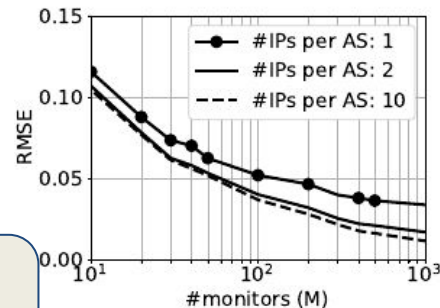
Ping-based impact estimator

- Approach 1: based on **ping measurements**...
 - *Goal*: we need to decrease the failure probability p

Ping-based impact estimator

1. Find “pingable” IP addresses for every AS [*ANT Lab’s IP hitlist*]
2. Ping multiple (N_{IP}) IP addresses per AS
3. If at least one ping reply from an AS → the AS is not affected by the hijack

N_{IP} (nb of pinged IPs per AS)	1	2	3	...	10
p (failure probability per AS)	12.8%	4.2%	2.1%	...	0%
$RMSE$ (estimation error; $M=100$)	7.9%	4.7%	4.1%	...	3.9%



Key findings:

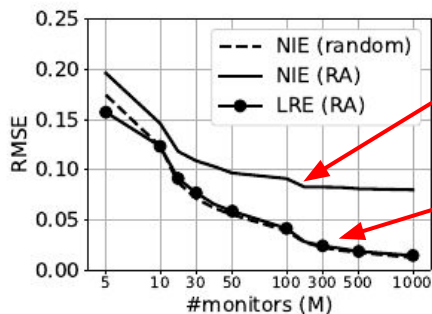
- $N_{IP} \geq 2$ for low error
- no need for $N_{IP} > 3$

Public infrastructure based estimator

- Approach 2: based on **public-infrastructure**...
 - *Goal*: we need to remove the measurement bias

Linear-regression estimator (LRE)

1. Collect past measurements of (public infrastructure) monitors
2. Fit a least-square estimator → give to each monitor i a weight w_i
3. Collect measurements m_i for the ongoing hijack
4. Estimate the impact as: $\sum_i m_i * w_i$



public
infrastructure

public infrastructure
with LRE

Key findings:

- ▶ LRE eliminates the bias in public infrastructure measurements & achieves close-to-theory efficiency
- ▶ Only a few past measurements are needed for fitting the LRE (e.g., it worked quite well even with 20 past events in our experiments)

Summarizing...

Estimating the impact of BGP prefix hijacking

- Important for network operations (e.g., mitigation actions)
- Not studied before
- We studied fundamental (limits, trade-offs, etc.) and practical aspects (use of public infrastructure, measurement failures, etc.)
 - theory (insights) & simulations (generality) & experiments (realism/verification)

Future research directions

- ML-based estimators (but... lack of labelled datasets)
- Generality of results - beyond BGP prefix hijacking
 - de-bias public infrastructure measurements
 - identify key locations for expanding public infrastructure

