

## Scaling a Variant Calling Genomics Pipeline with FaaS

9th International Workshop on Serverless Computing

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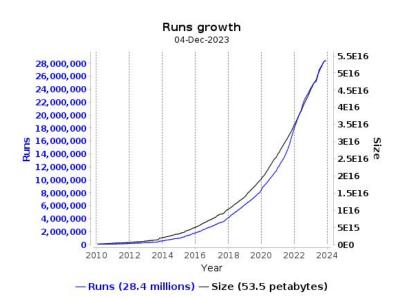


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### Genomics Workloads





- → Genomics is a computeand data intensive- task.
- → Exponential growth in data size and complexity.
- → Biomedical institutions with HPC struggle to keep up.

## Cloud IaaS for Genomics



The Cloud elasticity is key for scaling genomics workloads using short-term resources.









### Cloud IaaS for Genomics











The Cloud elasticity is key for scaling genomics workloads using short-term resources.

- **Complexity for bioinformatic users** 
  - Capacity/VM size for processing X GB of data?
  - Auto-scaling?
  - Hidden costs?

Configuring, deploying and scaling genomics workloads is challenging for bioinformatics users.

## Going serverless

#### Serverless (FaaS)

- Pay only for resources used at millisecond granularity, scale down to zero when not used
- 2. Instant scalability (~200 ms cold start, thousands of parallel functions)
- **3. Completely managed:** Scaling, security...

- → Why serverless for genomic pipelines?
  - No servers to manage!
  - Less friction to the Cloud for less experienced (bioinformatic) users
  - Allows to massively and effortlessly scale highly-parallel genomics workloads.

## Serverless Genomic Variant Calling

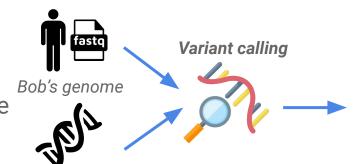
Objective: Adapt a single-node HPC variant calling genomics application to serverless in order to scale in parallelism, process larger datasets and decrease runtime.

## Serverless Genomic Variant Calling

Objective: Adapt a single-node HPC variant calling genomics application to serverless in order to scale in parallelism, process larger datasets and decrease runtime.

Reference human genome

 Variant Calling: detect differences (mutations, variants) in a sampled genome compared to a reference genome.



Personalized medicine, preventive disease detection, ... for Bob

## Distributing a Variant Calling pipeline

#### Sampled Genome (FASTQ)

```
@SEQ_ID1
CGGTAGCCAGCTGCGTTCAGTATGGAAGATTTGATTT
+
+&&-&%$%%$$$#)33&0$&%$''*''%$#%$%#+-5
@SEQ_ID2
TTCAGTTTATGGGTGCGGGTGTTATGTGACAAGAAAG+
"###""$$*#)%,+)+&'(,"###""&0$&%$''*&0
@SEQ_ID3
GCATGACCATACCGTGACAAGAAAGTCACCGCCCGTC
+
!''*((((***+))%%%++)(%%%%)'%%##%$('%#
@SEQ_ID4
CGGTAGCCAGCTGCGTTCAGTATGGAAGATTTGATTT
+
+&&-&%$%%$$$#)33&0$&%$''*'$$#%$%#+-5
@SEQ_ID5
TTCAGTTTATGGGTGCGGTGTTATGTGACAAGAAAG+
"###""$$$#)%,+)+&'(,"###""&0$&%$''*&0
```

#### Reference Genome (FASTQ)

```
>SEQUENCE 1
MTEITAAMVKELRESTGAGMMDCKNALSETNGDGLEKK
TEDFAAEVAAQLGLEKKTEDFAAEVAAQLFDKAVQLLR
EMGQFYVMDDKKTVEQVIAEKEKEFGGKIKIV
>SEQUENCE 2
SATVSEINSETDFVAKNDQFIALTKDTTAHIQSNSLQS
VEELHSSTINGVKFEEYLKSQIATIGENLVVRRFATLK
AGANGVVNGYIHTNGRVGVVIAAACDSAE
>SEQUENCE 3
MTEITAAMVKELRESTGAGMMDCKNALSETNGDGLEKK
TEDFAAEVAAQLGLEKKTEDFAAEVAAQLFDKAVQLLR
EMGQFYVMDDKKTVEQVIAEKEKEFGGKIKIV
>SEQUENCE 4
SATVSEINSETDFVAKNDQFIALTKDTTAHIQSNSLQS
VEELHSSTINGVKFEEYLKSQIATIGENLVVRRFATLK
AGANGVVNGYIHTNGRVGVVIAAACDSAE
>SEQUENCE 5
MTEITAAMVKELRESTGAGMMDCKNALSETNGDGLEKK
TEDFAAEVAAQLGLEKKTEDFAAEVAAQLFDKAVQLLR
EMGQFYVMDDKKTVEQVIAEKEKEFGGKIKIV
```

 Process: String similarity search -Compare all sample genome to all reference genome

## Distributing a Variant Calling pipeline

#### Sampled Genome (FASTQ)

```
@SEQ_ID1
CGGTAGCCAGCTGCGTTCAGTATGGAAGATTTGATTT
+
+&&-&%%%%%$$$#)33&0%&&$''*''%$#%$$#+-5
@SEQ_ID2
TTCAGTTTATGGGTGCGGGTGTTATGTGACAAGAAAG+
"###""$$%#)%,+)+&'(,"###""&0$&%$''*&0
```

```
@SEQ_ID3
GCATGACCATACCGTGACAAGAAAGTCACCGCCCGTC
+
!''*((((***+))%%%++)(%%%%)'%%##%$('%#
@SEQ_ID4
CGGTAGCCAGCTGCGTTCAGTATGGAAGATTTGATTT
+
+&&-&%%%%$$$#)33&0$&%$''*''%$#%$$#+-5
```

#### Reference Genome (FASTQ)

```
>SEQUENCE_1
MTEITAAMVKELRESTGAGMMDCKNALSETNGDGLEK
>SEQUENCE_2
SATVSEINSETDFVAKNDQFIALTKDTTAHIQSNSLQS
VEELHSSTINGVKFEEY
```

>SEQUENCE\_3
MTEITAAMVKELRESTGAGMMDCKNALSETNGDGLEKK
TEDFAAEVAAQLGLEK
>SEQUENCE\_4
MTEITAAMVKELRESTGAGMMDCKNALSETNGDGLEK

>SEQUENCE\_5
MTEITAAMVKELRESTGAGMMDCKNALSETNGDGLEKK
TEDFAAEVAAQLGLEKKTEDFAAEVAAQLFDKAVQLLR
EMGQFYVMDDKKTVEQVIAEKEKEFGGKIKIV

- Process: String similarity search -Compare all sample genome to all reference genome
- How to distribute the workload?
  - Partition the dataset and perform an all to all comparison
  - Cartesian product → Lots of parallel tasks
- Partitioning the dataset implies:
  - Partial correction process
  - Merging all partial results to produce the final output

## Serverless Variant Calling Pipeline Architecture

#### **Pre-Processing**

Partition input data to be distributed between lambdas.

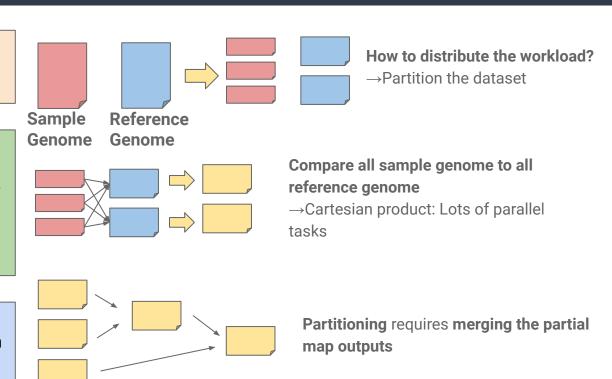
#### Map

Genome alignment (String similarity search) between sample and reference sets.

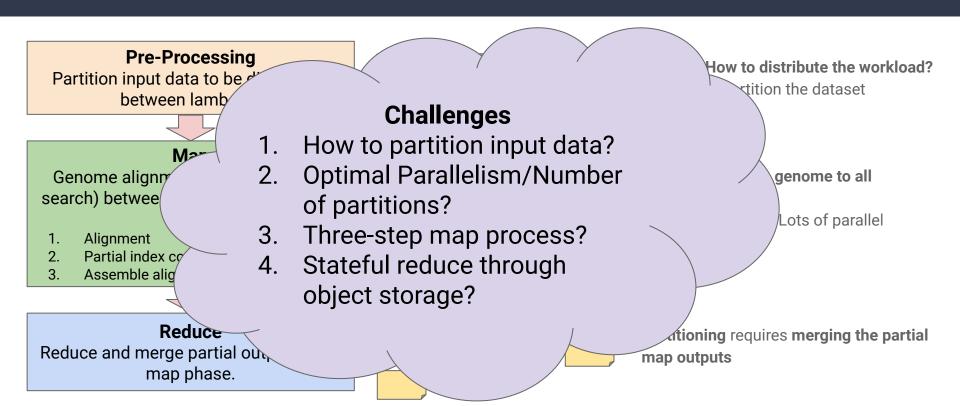
- Alignment
- Partial index correction
- 3. Assemble alignment

#### Reduce

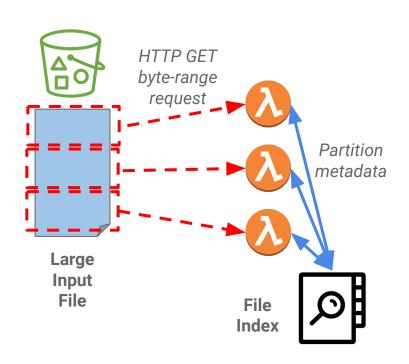
Reduce and merge partial outputs from map phase.



## Serverless Variant Calling Pipeline Architecture



# **Challenge 1**Input Data Partitioning

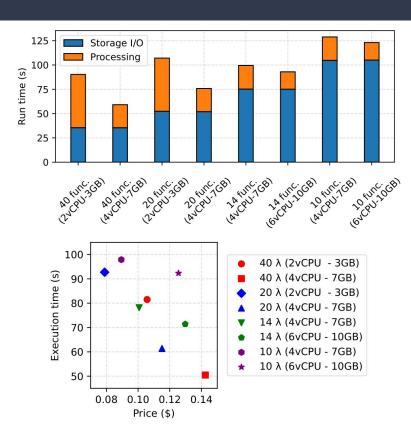


- Objective: partition input data
- Partial reads with HTTP GET byte-range requests
- Arbitrary byte-ranges breaks the genome file
- We require more metadata for each partition (sequence identifier and offset)
- Indexing to locate and identify each sequence for any arbitrary byte-range
- Lookup index for partition metadata

## Challenge 2 Data dependencies from synchronous HPC code

- Data dependencies→Functions calling functions: Stop the process, call another task, synchronize (wait), get result, then resume.
- No preemption in serverless →Blocked tasks occupy concurrency slot
- Can provoke deadlocks and limit scalability
- Blocking code must be split into many asynchronous tasks that can be scheduled independently
- Data dependencies must be passed through object storage

# **Challenge 3**Optimal degree of parallelism



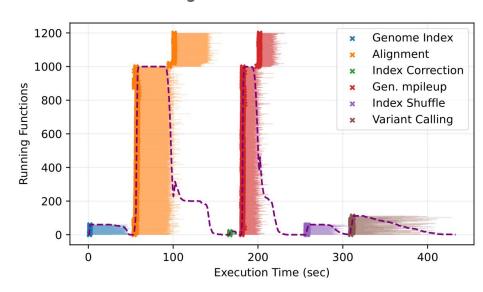
- How many tasks to launch?
- More tasks → More parallelism → Less scalable (concurrency limit)
- More tasks → Larger data chunks → Less efficient
- Leverage intra-function parallelism to launch less tasks without sacrificing performance

## Evaluation

#### **HPC vs Serverless**

Stage	HPC	Serverless
Genome Indexing Alignment Index correction Generate mpileup	0 min 14.20 s 0 min 14.20 s - 51 min 15.79 s	0 min 9.81 s 0 min 48.10 s 0 min 7.63 s 1 min 6.55 s
Index shuffle Variant Calling  Total	54 min 5.04s s 106 min 8.21 s	0 min 10.73 s 0 min 27.82 s 2 min 50.64 s

#### Large scale execution



## Insights and conclusions

- 1. Optimizing performance-cost with multi-CPU functions requires balancing intraand inter-function parallelism
- 2. Effective **data partitioning** is vital to exploit the massive parallelism offered by FaaS.
- 3. **Asynchronous and non-blocking** code is strictly necessary to guarantee scalability and to **avoid the concurrency limits** of serverless platforms.
- 4. Although **object storage** offers scalability for intermediate workflow data storage, it suffers from **slow performance and elevated costs**.
- Serverless data processing services for object storage, like S3 SELECT, provide a nice synergy with serverless FaaS workloads.

# Thank you! Any questions?

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# Challenge 4 Stateful data movements

 Reduce stage → Stateful data movements through object storage

- I/O time from Lambda to S3 is expensive
- We want to simplify the pipeline

- Delegate shuffle logic to S3 SELECT
- S3 SELECT allows to define simple SQL queries over semi-structured data
- Cheaper, more simple, and less error-prone than doing ad-hoc shuffling