

Fast Decomposable Submodular Function Minimization using Constrained Total Variation

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Neural Information Processing Systems, 2019.

Submodular functions

- ▶ **Definition:** $F : 2^V \rightarrow \mathbb{R}$ is **submodular** if and only if

$$\forall A, B \subseteq V, F(A) + F(B) \geq F(A \cup B) + F(A \cap B)$$

- ▶ Equality for *modular* functions
- ▶ Always assume $F(\emptyset) = 0$.

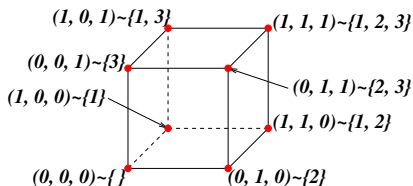
- ▶ **Equivalent definition :**

$$\forall k \in V, A \rightarrow F(A \cup \{k\}) - F(A) \text{ is non increasing}$$

- ▶ Diminishing returns property

Subsets as pseudo boolean functions

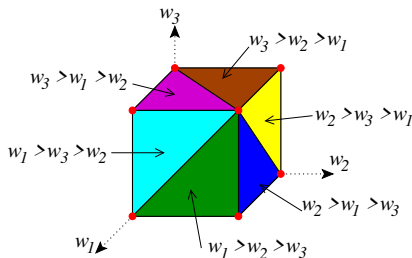
- Subsets may be identified with elements of $\{0, 1\}^P$.



Lovász extension

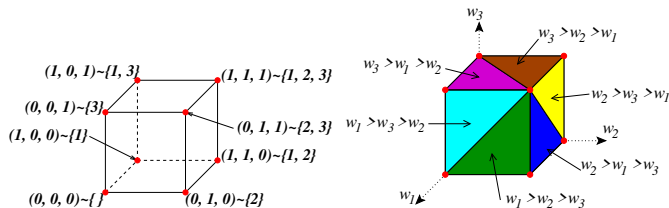
- ▶ Given **any** set-function F and w such that $w_{j_1} \geq \dots \geq w_{j_p}$,

$$f(w) = \sum_{k=1}^{p-1} (w_{j_k} - w_{j_{k+1}}) F(\{j_1, \dots, j_k\}) + w_{j_p} F(\{j_1, \dots, j_p\})$$



- ▶ f is piecewise-linear and positively homogeneous
- ▶ if $w = 1_A$, $f(w) = F(A)$
- ▶ Extension from $\{0, 1\}^p$ to \mathbb{R}^p
- ▶ **Theorem:** F is submodular if and only if f is convex

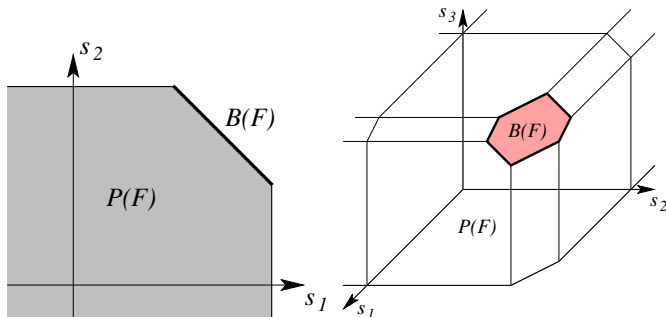
Cut functions and Lovász extension



- ▶ **Cut function.** $F : 2^V \rightarrow \mathbb{R}$ such that
$$F(A) = \sum_{i \in V} \sum_{j \in \mathcal{N}(i)} a_{ij} |1_{i \in A} - 1_{j \in A}|.$$
- ▶ **Lovász extension.** $f : [0, 1]^p \rightarrow \mathbb{R}$ such that
$$f(w) = \sum_{i \in V} \sum_{j \in \mathcal{N}(i)} a_{ij} |w_i - w_j|.$$
- ▶ If F is a cut function then f is the corresponding total variation problem.

Submodular Functions and base polyhedra

- ▶ Submodular polyhedron:
$$P(F) = \{s \in \mathbb{R}^p, \forall A \subset V, s(A) \leq F(A)\}$$
- ▶ Base polyhedron: $B(F) = P(F) \cap \{s(V) = F(V)\}$
- ▶ Many facets (up to 2^p), many extreme points (up to $p!$)



Lovász extension and base polyhedra

- ▶ **Fundamental property:** If F is submodular, maximizing linear functions may be done by a “greedy algorithm”
 - ▶ Let $w \in \mathbb{R}_+^p$ such that $w_{j_1} \geq \dots \geq w_{j_p}$
 - ▶ Let $s_{j_k} = F(\{j_1, \dots, j_k\}) - F(\{j_1, \dots, j_{k-1}\})$ for $k \in \{1, \dots, p\}$
 - ▶ Then $f(w) = \max_{s \in P(F)} w^\top s = \max_{s \in B(F)} w^\top s$
- ▶ **Representation of $f(w)$ as a support function:**

$$f(w) = \max_{s \in B(F)} w^\top s$$

From SFM to continuous minimization problems

► Related optimization problems

(SFMD) Discrete $\min_{A \subseteq V} F(A) - u(A)^1$

(SFMC) Continuous $\min_{w \in \mathbb{R}^p} f(w) - u^\top w + \sum_{j=1}^n \psi(w_j)$

► Solving (SFMC) is equivalent to

$$\min_{A \subseteq V} F(A) - u(A) + |A|\psi'(\alpha), \forall \alpha \in \mathbb{R}$$

► $\arg \min_{A \subseteq V} F(A) - u(A) + |A|\psi'(\alpha) = \{k, w_k^* \geq \alpha\}$

► Consequence of F being submodular.

► $\alpha = 0$ solves (SFMD).

► If (SFMD) is graph cut \Rightarrow (SFMC) is parametric maxflow.

¹ $u(A) = u^\top 1_A$

Overview

► Goal

(SFMD) Discrete $\min_{A \subset V} \sum_{i=1}^r F_i(A) - u(A)$

► Assuming the oracles

(SFMD_i) Discrete $\min_{A \subset V} F_i(A) - t(A)$

► Define

► (SFMC) $\min_{w \in \mathbb{R}^p} f(w) - u^\top w + \sum_{j=1}^n \psi(w_j)$

► (SFMC_i) $\min_{w \in \mathbb{R}^p} f_i(w) - t^\top w + \sum_{j=1}^n \psi(w_j)$

► Define ψ

$$\psi(w) = \begin{cases} \frac{1}{2}w^2 & \text{if } |w| \leq \varepsilon, \\ +\infty & \text{otherwise,} \end{cases} \quad (1)$$

► Approach

► Solve (SFMC_i) using (SFMD_i)

► Solve (SFMD) by solving (SFMC) using (SFMC_i)