

# AVR Microcontroller

Microprocessor Course

Chapter 5

ARITHMATIC, LOGIC INSTRUCTIONS AND PROGRAMS

Aban 1401

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

Unsigned numbers are defined as data in which all the bits are used to represent data and no bits are set aside for the positive or negative sign. This means that the operand can be between 00 and FFH (0 to 255 decimal) for 8-bit data.

### Addition of unsigned numbers

In the AVR, the add operation has two general purpose registers as inputs and the result will be stored in the first (left) register. One form of the ADD instruction in the AVR is:

**ADD Rd,Rr            ;Rd = Rd + Rr**

It could change any of the Z, C, N, V, H or S bits of the status register, depending on the operands involved.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Example 5-1

Show how the flag register is affected by the following instructions.

```
LDI    R21,0xF5      ;R21 = F5H
LDI    R22,0x0B      ;R22 = 0x0BH
ADD    R21,R22        ;R21 = R21+R22 = F5+0B = 00 and C = 1
```

### Solution:

|       |             |
|-------|-------------|
| F5H   | 1111 0101   |
| + 0BH | + 0000 1011 |
| 100H  | 0000 0000   |

After the addition, register R21 contains 00 and the flags are as follows:

C = 1 because there is a carry out from D7.

Z = 1 because the result in destination register (R21) is zero.

H = 1 because there is a carry from D3 to D4.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Example 5-2

Assume that RAM location 400H has the value of 33H. Write a program to find the sum of location 400H of RAM and 55H. At the end of the program, R21 should contain the sum.

#### Solution:

```
LDS    R2,0x400    ;R2 = 33H (location 0x400 of RAM)
LDI    R21,0x55    ;R21 = 55
ADD     R21,R2      ;R21 = R21 + R2 = 55H + 33H = 88H, C = 0
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### ADC and addition of 16-bit numbers

When adding two 16-bit data operands, we need to be concerned with the propagation of a carry from the lower byte to the higher byte. This is called multibyte addition to distinguish it from the addition of individual bytes. The instruction ADC (ADD with carry) is used on such occasions.

For example, look at the addition of  $3CE7H + 3B8DH$ , as shown next.

$$\begin{array}{r} 1 \\ 3C \ E7 \\ + \quad 3B \ 8D \\ \hline 78 \ 74 \end{array}$$

When the first byte is added, there is a carry ( $E7 + 8D = 74$ ,  $C = 1$ ). The carry is propagated to the higher byte, which results in  $3C + 3B + 1 = 78$  (all in hex).

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Example 5-3

Write a program to add two 16-bit numbers. The numbers are 3CE7H and 3B8DH. Assume that R1 = 8D, R2 = 3B, R3 = E7, and R4 = 3C. Place the sum in R3 and R4; R3 should have the lower byte.

#### Solution:

```
;R1 = 8D
;R2 = 3B
;R3 = E7
;R4 = 3C

ADD    R3,R1      ;R3 = R3 + R1 = E7 + 8D = 74 and C = 1
ADC    R4,R2      ;R4 = R4 + R2 + carry, adding the upper byte
                        ;with carry from lower byte
                        ;R4 = 3C + 3B + 1 = 78H (all in hex)
```

Notice the use of ADD for the lower byte and ADC for the higher byte.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Subtraction of unsigned numbers

In many microprocessors, there are two different instructions for subtraction: SUB and SUBB (subtract with borrow). In the AVR we have five instructions for subtraction: SUB, SBC, SUBI, SBCI, and SBIW.

|      |            |                     |
|------|------------|---------------------|
| SUB  | Rd, Rr     | ; Rd=Rd-Rr          |
| SBC  | Rd, Rr     | ; Rd=Rd-Rr-c        |
| SUBI | Rd, K      | ; Rd=Rd-K           |
| SBCI | Rd, K      | ; Rd=Rd-K-c         |
| SBIW | Rd:Rd+1, K | ; Rd+1:Rd=Rd+1:Rd-K |

**Figure 5-1.**

The SBC and SBCI instructions are subtract with borrow. In the AVR, we use the C (carry) flag for the borrow and that is why they are called SBC (SUB with Carry).

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

***SUB Rd,Rr (Rd = Rd - Rr)***

In subtraction, the AVR microcontrollers use the 2's complement method. Although every CPU contains adder circuitry, it would be too cumbersome (and take too many transistors) to design separate subtractor circuitry. For this reason, the AVR uses adder circuitry to perform the subtraction command. Assuming that the AVR is executing a simple subtract instruction and that  $C = 0$  prior to the execution of the instruction, one can summarize the steps of the hardware of the CPU in executing the SUB instruction for unsigned numbers as follows

1. Take the 2's complement of the subtrahend (right-hand operand).
2. Add it to the minuend (left-hand operand).
3. Invert the carry.



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Example 5-4

Show the steps involved in the following.

```
LDI    R20, 0x23      ;load 23H into R20
LDI    R21, 0x3F      ;load 3FH into R21
SUB     R21, R20       ;R21 <- R21-R20
```

### Solution:

|                   |           |                   |                      |
|-------------------|-----------|-------------------|----------------------|
| R21 = 3F          | 0011 1111 | 0011 1111         |                      |
| - R20 = <u>23</u> | 0010 0011 | + 1101 1101       | (2's complement)     |
| 1C                |           | 1 0001 1100       |                      |
|                   |           | C = 0, D7 = N = 0 | (result is positive) |

The flags would be set as follows: N = 0, C = 0. (Notice that there is a carry but C = 0. We will discuss this more in the next section.) The programmer must look at the N (or C) flag to determine if the result is positive or negative.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Example 5-5

Write a program to subtract 18H from 29H and store the result in R21 (a) without using the SUBI instruction, and (b) using the SUBI instruction.

#### Solution:

(a)

```
LDI    R21,0x29    ;R21 = 29H
LDI    R22,0x18    ;R22 = 18H
SUB     R21,R22     ;R21 = R21 - R22 = 29 - 18 = 11 H
```

(b)

```
LDI    R21,0x29    ;R21 = 29H
SUBI    R21,0x18    ;R21 = R21 - 18 = 29 - 18 = 11 H
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Example 5-6

Write a program to subtract 18H from 2917H and store the result in R25 and R24.

#### Solution:

```
LDI    R25,0x29          ;load the high byte (R25 = 29H)
LDI    R24,0x17          ;load the low byte (R24 = 17H)
SBIW   R25:R24,0x18      ;R25:R24 <- R25:R24 - 0x18
                        ;28FF = 2917 - 18
```

Notice that you should use `SBIW Rd+1:Rd,K` format. If `SBIW Rd:Rd+1,K` format is used, the assembler will assemble your code as if you had typed `SBIW Rd+1:Rd,K`. Change the third line of the code from `SBIW R25:R24,0x18` to `SBIW R24:R25,0x18` and examine the result.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

**SBC (Rd  $\leftarrow$  Rd - Rr - C) subtract with borrow (denoted by C)**

This instruction is used for multibyte numbers and will take care of the borrow of the lower byte. If the borrow flag is set to one ( $C = 1$ ) prior to executing the SBC instruction, this operation also subtracts 1 from the result.

### Example 5-7

Write a program to subtract two 16-bit numbers:  $2762H - 1296H$ . Assume  $R26 = (62)$  and  $R27 = (27)$ . Place the difference in  $R26$  and  $R27$ ;  $R26$  should have the lower byte.

;R26 = (62)

;R27 = (27)

```
LDI    R28,0x96    ;load the low byte (R28 = 96H)
LDI    R29,0x12    ;load the high byte (R29 = 12H)
SUB     R26,R28     ;R26 = R26 - R28 = 62 - 96 = CCH
                     ;C = borrow = 1, N = 1
SBC     R27,R29     ;R27 = R27 - R29 - C
                     ;R27 = 27 - 12 - 1 = 14H
```

After the SUB,  $R26$  has  $= 62H - 96H = CCH$  and the carry flag is set to 1, indicating there is a borrow (notice,  $N = 1$ ). Because  $C = 1$ , when SBC is executed  $R27$  has  $27H - 12H - 1 = 14H$ . Therefore, we have  $2762H - 1296H = 14CCH$ .

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Multiplication of unsigned numbers

The AVR has several instructions dedicated to multiplication. Here we will discuss the MUL instruction. Other instructions are similar to MUL but are used for signed numbers.

**Table 5-1: Multiplication Summary**

| Multiplication | Application                          | Byte1 | Byte2 | High byte of result | Low byte of result |
|----------------|--------------------------------------|-------|-------|---------------------|--------------------|
| MUL Rd, Rr     | Unsigned numbers                     | Rd    | Rr    | R1                  | R0                 |
| MULS Rd, Rr    | Signed numbers                       | Rd    | Rr    | R1                  | R0                 |
| MULSU Rd, Rr   | Unsigned numbers with signed numbers | Rd    | Rr    | R1                  | R0                 |

MUL is a byte-by-byte multiply instruction. In byte-by-byte multiplication, operands must be in registers. After multiplication, the 16-bit unsigned product is placed in R1 (high byte) and R0 (low byte).

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

The following example multiplies 25H by 65H.

```
LDI    R23,0x25    ;load 25H to R23
LDI    R24,0x65    ;load 65H to R24
MUL     R23,R24     ;25H * 65H = E99 where
                    ;R1 = 0EH and R0 = 99H
```

$$0x25 = 37$$

$$0x65 = 101$$

$$37 * 101 = 3737 = 0x0E99$$

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Division of unsigned numbers

**AVR has no instruction for divide operation.** We can write a program to perform division by repeated subtraction.

```
.DEF  NUM = R20
.DEF  DENOMINATOR = R21
.DEF  QUOTIENT = R22

      LDI    NUM,95           ;NUM = 95
      LDI    DENOMINATOR,10   ;DENOMINATOR = 10
      CLR    QUOTIENT         ;QUOTIENT = 0

L1:    INC    QUOTIENT
      SUB    NUM, DENOMINATOR
      BRCC   L1               ;branch if C is zero

      DEC    QUOTIENT         ;once too many
      ADD    NUM, DENOMINATOR ;add back to it

HERE:  JMP   HERE             ;stay here forever
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### An application for division

Sometimes a sensor is connected to an ADC (analog-to-digital converter) and the ADC represents some quantity such as temperature or pressure. The 8-bit ADC provides data in hex in the range of 00-FFH. This hex data must be converted to decimal. We do that by dividing it by 10 repeatedly, saving the remainders,

#### Example 5-8

Assume that the data memory location 0x315 has value FD (hex). Write a program to convert it to decimal. Save the digits in locations 0x322, 0x323, and 0x324, where the least-significant digit is in location 0x322.

```
1  .EQU HEX_NUM = 0x315
2  .EQU RMND_L = 0x322
3  .EQU RMND_M = 0x323
4  .EQU RMND_H = 0x324
5  .DEF NUM = R20
6  .DEF DENOMINATOR = R21
7  .DEF QUOTIENT = R22
```



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

```

 9      LDI      R16, 0xFD          ;$FD = 253 in decimal
10      STS      HEX_NUM, R16      ;store $FD in location 0x315
11
12      LDS      NUM, HEX_NUM
13      LDI      DENOMINATOR, 10    ;DENOMINATOR = 10
14
15      L1:      INC      QUOTIENT
16              SUB      NUM, DENOMINATOR
17              BRCC     L1          ;if C = 0 go back
18
19              DEC      QUOTIENT    ;once too many
20              ADD      NUM, DENOMINATOR ;add back to it
21              STS      RMND_L, NUM ;store remainder as the 1st digit
22
23              MOV      NUM, QUOTIENT
24              LDI      QUOTIENT, 0
25
26      L2:      INC      QUOTIENT
27              SUB      NUM, DENOMINATOR
28              BRCC     L2
29
30              DEC      QUOTIENT    ;once too many
31              ADD      NUM, DENOMINATOR ;add back to it
32              STS      RMNDM, NUM  ;store remainder as the 2nd digit
33
34              STS      RMND_H, QUOTIENT ;store quotient as the 3rd digit
35
36      HERE:    JMP      HERE      ;stay here forever
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.1 ARITHMETIC INSTRUCTIONS

### Example 5-9

Analyze the program in Example 5-8 for a numerator of 253.

#### Solution:

To convert a binary (hex) value to decimal, we divide it by 10 repeatedly until the quotient is less than 10. After each division the remainder is saved. In the case of an 8-bit binary, such as FDH, we have 253 decimal, as shown below.

|          | <b>Quotient</b> | <b>Remainder</b> |
|----------|-----------------|------------------|
| 253/10 = | 25              | 3 (low digit)    |
| 25/10 =  | 2               | 5 (middle digit) |
|          |                 | 2 (high digit)   |

Therefore, we have FDH = 253.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### 5.2 SIGNED NUMBER CONCEPTS AND ARITHMETIC OPERATIONS

All data items used so far have been unsigned numbers. Many applications require signed data. In this section the concept of signed numbers is discussed along with related instructions.

#### Concept of signed numbers in computers

In everyday life, numbers are used that could be positive or negative. To do that, computer scientists have devised the following arrangement for the representation of signed positive and negative numbers: The most significant bit (MSB) is set aside for the sign (+ or -), while the rest of the bits are used for the magnitude. The sign is represented by 0 for positive (+) numbers and 1 for negative (-) numbers.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### Signed bit operands

### Positive numbers

The range of positive numbers that can be represented by the format shown is 0 to +127. If a positive number is larger than +127, a 16-bit operand must be used.

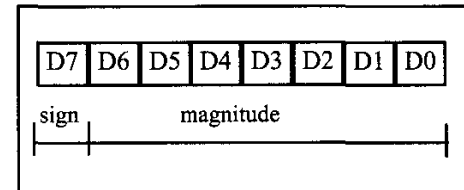


Figure 5-2. 8-Bit Signed Operand

### Negative numbers

For negative numbers, D7 is 1; however. Although the assembler does the conversion, it is still important to understand how the conversion works. To convert to negative number representation (2's complement), follow these steps:

1. Write the magnitude of the number in 8-bit binary (no sign).
2. Invert each bit.
3. Add 1 to it.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### Example 5-10

Show how the AVR would represent -5.

#### Solution:

Observe the following steps.

- |    |           |                                 |
|----|-----------|---------------------------------|
| 1. | 0000 0101 | 5 in 8-bit binary               |
| 2. | 1111 1010 | invert each bit                 |
| 3  | 1111 1011 | add 1 (which becomes FB in hex) |

Therefore, -5 = FBH, the signed number representation in 2's complement for -5. The D7 = N = 1 indicates that the number is negative.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

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### Example 5-11

---

Show how the AVR would represent -34H.

#### Solution:

Observe the following steps.

- |    |           |                            |
|----|-----------|----------------------------|
| 1. | 0011 0100 | 34H given in binary        |
| 2. | 1100 1011 | invert each bit            |
| 3  | 1100 1100 | add 1 (which is CC in hex) |

Therefore, -34 = CCH, the signed number representation in 2's complement for 34H.  
The D7 = N = 1 indicates that the number is negative.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### Example 5-12

Show how the AVR would represent -128.

#### Solution:

Observe the following steps.

- |    |           |                                 |
|----|-----------|---------------------------------|
| 1. | 1000 0000 | 128 in 8-bit binary             |
| 2. | 0111 1111 | invert each bit                 |
| 3  | 1000 0000 | add 1 (which becomes 80 in hex) |

Therefore,  $-128 = 80H$ , the signed number representation in 2's complement for -128. The  $D7 = N = 1$  indicates that the number is negative. Notice that 128 (binary 10000000) in unsigned representation is the same as signed -128 (binary 10000000).

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

From the examples above, it is clear that the range of byte-sized negative numbers is -1 to -128. The following lists byte-sized signed number ranges:

| <b>Decimal</b> | <b>Binary</b> | <b>Hex</b> |
|----------------|---------------|------------|
| -128           | 1000 0000     | 80         |
| -127           | 1000 0001     | 81         |
| -126           | 1000 0010     | 82         |
| ...            | .....         | ..         |
| -2             | 1111 1110     | FE         |
| -1             | 1111 1111     | FF         |
| 0              | 0000 0000     | 00         |
| +1             | 0000 0001     | 01         |
| +2             | 0000 0010     | 02         |
| ..             | .....         | ..         |
| +127           | 0111 1111     | 7F         |



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### Overflow problem in signed number operations

What is an overflow? If the result of an operation on signed numbers is too large for the register, an overflow has occurred and the programmer must be notified.

#### Example 5-13

Examine the following code and analyze the result, including the N and V flags.

```
LDI    R20,0x60    ;R20 = 0110 0000 (+70)
LDI    R21,0x46    ;R21 = 0100 0110 (+96)
ADD     R20,R21     ;R20 = (+96) + (+70) = 1010 0110
                        ;R20 = A6H = -90 decimal, INVALID!!
```

#### Solution:

```
    +96    0110 0000
+   +70    0100 0110
+  166    1010 0110  N = 1 (negative) and V = 1 Sum = -90
```

According to the CPU, the result is negative ( $N = 1$ ), which is wrong. The CPU sets  $V = 1$  to indicate the overflow error. Remember that the N flag is the D7 bit. If  $N = 0$ , the sum is positive, but if  $N = 1$ , the sum is negative.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### When is the V flag set?

In 8-bit signed number operations, V is set to 1 if either of the following two conditions occurs:

1. There is a carry from D6 to D7 but no carry out of D7 (C = 0).
2. There is a carry from D7 out (C = 1) but no carry from D6 to D7.

#### Example 5-14

Examine the following code, noting the role of the V and N flags:

```
LDI    R20,0x80    ;R20 = 1000 0000 (80H = -128)
LDI    R21,0xFE    ;R21 = 1111 1110 (FEH = -2)
ADD    R20,R21     ;R20 = (-128) + (-2)
                     ;R20 = 10000000 + 11111110 = 0111 1110,
                     ;N = 0, R0 = 7EH = +126, invalid
```

#### Solution:

|       |                  |                            |
|-------|------------------|----------------------------|
| -128  | 1000 0000        |                            |
| + - 2 | <u>1111 1110</u> |                            |
| - 130 | 0111 1110        | N = 0 (positive) and V = 1 |

According to the CPU, the result is +126, which is wrong, and V = 1 indicates that. Notice that the N flag indicates the sign of the corrupted result, not the sign that the real result should have.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### Further considerations on the V flag

In the ADD instruction, there are two different conditions. Either the operands have the same sign or the signs of the operands are different.

- When we ADD two numbers with different signs, the absolute value of the result is smaller than the operands before executing the ADD instruction. So overflow definitely cannot happen after two operands with different signs are added.
- **Overflow is possible only when we ADD two operands with the same sign.** In this case the absolute value of the result is larger than the operands before executing the ADD instruction. So it is possible that the result will be too large for the register and cause overflow. If we add two numbers with the same sign and the result sign is different, we know that overflow has occurred. That is exactly the way that the CPU knows when to set the V flag. In the AVR the equation of the V flag is as follows:

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ....

### **Further considerations on the V flag**

In the AVR the equation of the V flag is as follows:

$$V = Rd7 \cdot Rr7 \cdot \overline{R7} + \overline{Rd7} \cdot \overline{Rr7} \cdot R7$$

where Rd7 and Rr7 are the 7th bit of the operands and R7 is the 7th bit of the result.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### Example 5-15

Examine the following code, noting the role of the V and N flags:

```
LDI    R20,-2           ;R20 = 1111 1110 (R20 = FEH)
LDI    R21,-5           ;R21 = 1111 1110 (R21 = FBH)
ADD    R20,R21           ;R20 = (-2) + (-5) = -7 or F9H
                        ;correct, since V = 0
```

### Solution:

|      |                  |                                      |
|------|------------------|--------------------------------------|
| -2   | 1111 1110        |                                      |
| + -5 | <u>1111 1011</u> |                                      |
| - 7  | 1111 1001        | and V = 0 and N = 1. Sum is negative |

According to the CPU, the result is -7, which is correct, and the V indicates that (V = 0).

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### Example 5-16

Examine the following code, noting the role of the V and N flags:

```
LDI    R20,7           ;R20 = 0000 0111
LDI    R20,18          ;R20 = 0001 0010
ADD     R20,R21         ;R20 = (+7) + (+18)
                        ;R20 = 00000111 + 00010010 = 0001 1001
                        ;R20 = (+7) + (+18) = +25, N = 0, positive
                        ;and correct, V = 0
```

**Solution:**

```
  + 7 0000 0111
+ +18 0001 0010
-----
+25 0001 1001  N = 0 (positive 25) and V = 0
```

According to the CPU, this is +25, which is correct, and  $V = 0$  indicates that.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

From Examples 5-14 to 5-16, we conclude that in any signed number addition, **V indicates whether the result is valid or not**. If  $V = 1$ , the result is **erroneous**; if  $V = 0$ , the result is **valid**.

**We can state emphatically that in unsigned number addition, the programmer must monitor the status of C (carry flag), and in signed number addition, the V (overflow) flag must be monitored.**

In the AVR, instructions such as BRCS and BRCC allow the program to branch right after the addition of unsigned numbers according to the value of C flag.

There are also the BRVC and the BRVS instructions for the V flag that allow us to correct the signed number error. We also have two branch instructions for the N flag (negative), BRPL and BRMI.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.2: SIGNED NUMBER CONCEPTS AND ...

### What is the difference between the N and S flags?

In signed numbers the N flag represents the D7 bit of the result. It is called the Negative flag.

In operations on signed numbers, overflow is possible. **Overflow** corrupts the result and **negates the sign bit**. So if you ADD two positive numbers, in case of overflow, the N flag would be 1 showing that the result is negative! The S flag helps you to know the sign of the real result. It checks the V flag in addition to the D7 bit.

- If  $V = 0$ , it shows that overflow has not occurred and the S flag will be the same as D7 to show the sign of the result.
- If  $V = 1$ , it shows that overflow has occurred and the S flag will be opposite to the D7 to show the sign of the real (not the corrupted) result.



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

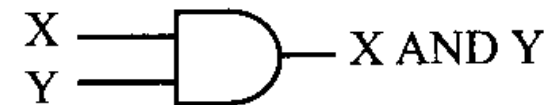
### AND

**AND Rd,Rr ;Rd = Rd AND Rr**

This instruction will perform a logical AND on the two operands and place the result in the left-hand operand. There is also the “**ANDI Rd,k**” instruction in which the right-hand operand can be a constant value. The AND instruction will affect the Z, S, and N flags.

#### Logical AND Function

| Inputs |   | Output  |
|--------|---|---------|
| X      | Y | X AND Y |
| 0      | 0 | 0       |
| 0      | 1 | 0       |
| 1      | 0 | 0       |
| 1      | 1 | 1       |



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### Example 5-18

Show the results of the following.

```
LDI    R20,0x35    ;R20 = 35H
ANDI    R20,0x0F    ;R20 = R20 AND 0FH (now R20 = 05)
```

**Solution:**

|     |       |      |      |                                  |
|-----|-------|------|------|----------------------------------|
|     | 35H   | 0011 | 0101 |                                  |
| AND | 0FH   | 0000 | 1111 |                                  |
|     | ----- |      |      |                                  |
|     | 05H   | 0000 | 0101 | ;35H AND 0FH = 05H, Z = 0, N = 0 |

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### OR

**OR Rd,Rr ;Rd = Rd OR Rr**

This instruction will perform a logical OR on the two operands and place the result in the left-hand operand. There is also the “**ORI Rd,k**” instruction in which the right-hand operand can be a constant value. The OR instruction will affect the Z, S, and N flags.

#### Logical OR Function

| <u>Inputs</u> |          | <u>Output</u> |
|---------------|----------|---------------|
| <u>X</u>      | <u>Y</u> | <u>X OR Y</u> |
| 0             | 0        | 0             |
| 0             | 1        | 1             |
| 1             | 0        | 1             |
| 1             | 1        | 1             |



# ARITHMETIC. LOGIC INSTRUCTIONS AND PROGRAMS

5.

## Example 5-19

(a) Show the results of the following:

```
LDI    R20, 0x04           ;R20 = 04
ORI     R20, 0x30           ;now R20 = 34H
```

(b) Assume that PB2 is used to control an outdoor light, and PB5 to control a light inside a building. Show how to turn “on” the outdoor light and turn “off” the inside one.

### Solution:

```
(a)      04H      0000 0100
        OR      30H      0011 0000
        -----
        34H      0011 0100      04 OR 30 = 34H, Z = 0 and N = 0

(b)
SBI      DDRB, 2           ;bit 2 of Port B is output
SBI      DDRB, 5           ;bit 5 of Port B is output
IN       R20, PORTB        ;move PORTB to R20. (Notice that we read
                           ;the value of PORTB instead of PINB
                           ;because we want to know the last value
                           ;of PORTB, not the value of the AVR
                           ;chip pins.)
ORI      R20, 0b00000100   ;set bit 2 of R20 to one
ANDI     R20, 0b11011111   ;clear bit 5 of R20 to zero
OUT      PORTB, R20        ;out R20 to PORTB

HERE:    JMP HERE          ;stop here
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

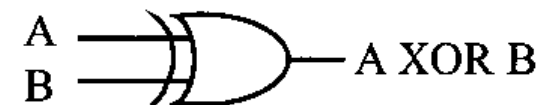
## 5.3: LOGIC AND COMPARE INSTRUCTIONS

**eOR Rd,Rr            ;Rd = Rd XOR Rr**

This instruction will perform a logical EX-OR on the two operands and place the result in the left-hand operand. There is also the “**eOR**” instruction will affect the Z, S, and N flags.

### Logical XOR Function

| Inputs |   | Output  |
|--------|---|---------|
| A      | B | A XOR B |
| 0      | 0 | 0       |
| 0      | 1 | 1       |
| 1      | 0 | 1       |
| 1      | 1 | 0       |



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### Example 5-20

Show the results of the following:

```
LDI    R20, 0x54
LDI    R21, 0x78
EOR     R20, R21
```

**Solution:**

```
      54H    0101 0100
XOR   78H    0111 1000
-----
      2CH    0010 1100
```

54H XOR 78H = 2CH, Z = 0, N = 0

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

EX-OR can also be used to see if two registers have the same value. The "EOR R0,R1" instruction will EX-OR the R0 register and R1, and put the result in R0.

### Example 5-21

The EX-OR instruction can be used to test the contents of a register by EX-ORing it with a known value. In the following code, we show how EX-ORing the value 45H with itself will raise the Z flag:

```
OVER:      IN      R20,PINB
           LDI      R21,0x45
           EOR      R20,R21
           BRNE     OVER
```

#### Solution:

|            |                 |
|------------|-----------------|
| 45H        | 01000101        |
| <u>45H</u> | <u>01000101</u> |
| 00         | 00000000        |

EX-ORing a number with itself sets it to zero with  $Z = 1$ . We can use the BREQ instruction to make the decision. EX-ORing with any other number will result in a nonzero value.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### Example 5-22

Read and test PORTB to see whether it has the value 45H. If it does, send 99H to PORTC; otherwise, it is cleared.

#### Solution:

```
LDI    R20,0xFF      ;R20 = 0xFF
OUT     DDRC,R20      ;Port C is output
LDI     R20,0x00      ;R20 = 0
OUT     DDRB,R20      ;Port B is input
OUT     PORTC,R20     ;PORTC = 00
LDI     R21,0x45      ;R21 = 45
HERE:
IN      R20,PINB      ;get a byte
EOR     R20,R21       ;EX-OR with 0x45

BRNE    HERE          ;branch if PORTB has value other than 45
LDI     R20,0x99      ;R20 = 0x99
OUT     PORTC,R20     ;PORTC = 99h
EXIT:   JMP     EXIT   ;stop here
```



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

Another widely used application of EX-OR is to toggle the bits of an operand. The following code demonstrates how to use EX-OR to toggle the bits of an operand.

```
LDI    R20, 0xFF
EOR     R0, R20      ;EX-OR R0 with 1111 1111 will
                    ;change all the bits of R0 to
                    ;opposite
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

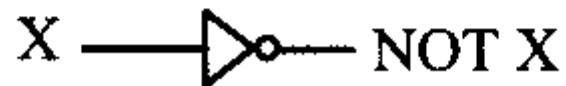
## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### COM (complement)

This instruction complements the contents of a register. The complement action changes the 0s to 1s, and the 1s to 0s. This is also called 1's complement.

#### Logical Inverter

| <u>Input</u> | <u>Output</u> |
|--------------|---------------|
| <u>X</u>     | <u>NOT X</u>  |
| 0            | 1             |
| 1            | 0             |



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### NEG (negate)

This instruction takes the 2's complement of a register.

#### Example 5-23

Find the 2's complement of the value 85H. Notice that 85H is -123.

**Solution:**

|     |           |            |                      |
|-----|-----------|------------|----------------------|
| LDI | R21, 0x85 |            | ; 85H = 1000 0101    |
|     |           |            | ; 1's = 0111 1010    |
|     |           |            | <u>          + 1</u> |
| NEG | R21       | ; 2's comp | 0111 1011 = 7BH      |

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### Compare instructions

**CP      Rd, Rr**

The AVR has the CP instruction for the compare operation. The compare instruction is really a subtraction, except that the values of the operands do not change. There is also the “**CPI Rd, k**” instruction in which the right-hand operand can be a constant value.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### Conditional branch instructions

As we studied in Chapter 3, conditional branches alter the flow of control if a condition is true. In the AVR there are at least two conditional jumps for each flag of the status register.

**Table 5-2: AVR Compare Instructions**

|             |   |                                     |
|-------------|---|-------------------------------------|
| <b>BREQ</b> | <b>Branch if equal</b>                          | <b>Branch if <math>Z = 1</math></b> |
| <b>BRNE</b> | <b>Branch if not equal</b>                      | <b>Branch if <math>Z = 0</math></b> |
| <b>BRSH</b> | <b>Branch if same or higher</b>                 | <b>Branch if <math>C = 0</math></b> |
| <b>BRLO</b> | <b>Branch if lower</b>                          | <b>Branch if <math>C = 1</math></b> |
| <b>BRLT</b> | <b>Branch if less than (signed)</b>             | <b>Branch if <math>S = 1</math></b> |
| <b>BRGE</b> | <b>Branch if greater than or equal (signed)</b> | <b>Branch if <math>S = 0</math></b> |
| <b>BRVS</b> | <b>Branch if Overflow flag set</b>              | <b>Branch if <math>V = 1</math></b> |
| <b>BRVC</b> | <b>Branch if Overflow flag clear</b>            | <b>Branch if <math>V = 0</math></b> |

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### **BREQ and BRNE instructions**

```
BREQ    k        ;if (Z = 1) then branch  
                ;else continue
```

The BREQ makes decisions based on the Z flag. If  $Z = 1$  the BREQ instruction branches

The BRNE instruction, like the BREQ, makes decisions based on the Z flag, but it branches when  $Z = 0$ .

Notice that the BREQ and BRNE instructions can be used for both signed, and unsigned numbers.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### Example 5-24

Write a program to monitor PORTB continuously for the value 63H. It should stop monitoring only if PORTB = 63H.

#### Solution:

```
    LDI    R20,0x00
    OUT    DDRB,R20      ;PORT B is input
    LDI    R21,0x63
AGAIN:
    IN     R20,PINB
    CP     R20,R21        ;compare with 0x63, Z = 1 if yes
    BRNE   AGAIN         ;go to AGAIN if PORTB is not equal to 0x63
    ....
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### **BRSH and BRLO instructions**

```
BRSH k          ;if (C = 0) then branch  
                ;else continue
```

The BRSH makes decisions based on the C flag. If  $C = 0$  the CPU will jump.

The BRLO instruction, like the BRSH, makes decisions based on the C flag, but it branches when  $C = 1$ .



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### Example 5-25

Write a program to find the greater of the two values 27 and 54, and place it in R20.

#### Solution:

```
.EQU  VAL_1=27
.EQU  VAL_2=54

    LDI    R20,VAL_1    ;R20 = VAL_1
    LDI    R21,VAL_2    ;R21 = VAL_2
    CP     R21,R20      ;compare R21 and R20
    BRLO   NEXT        ;if R21<R20 (branch if lower) go to NEXT
    LDI    R20,VAL_2    ;R20 = VAL_2
NEXT:
```

**Example 5-26**

Assume that Port B is an input port connected to a temperature sensor. Write a program to read the temperature and test it for the value 75. According to the test results, place the temperature value into the registers indicated by the following.

|           |              |                     |
|-----------|--------------|---------------------|
| If T = 75 | then R16 = T | ; R17 = 0 ; R18 = 0 |
| If T > 75 | then R16 = 0 | ; R17 = T ; R18 = 0 |
| If T < 75 | then R16 = 0 | ; R17 = 0 ; R18 = T |

**Solution:**

```
LDI    R20,0x00           ;R20 = 0
OUT     DDRB,R20          ;Port B = input

CLR     R16               ;R16 = 0
CLR     R17               ;R17 = 0
CLR     R18               ;R18 = 0

IN      R20,PINB
CPI     R20,75             ;compare R20 (PORTB) and 75
BRSH    SAME_HI           ;executes when R20 < 75

MOV     R18,R20
RJMP    CNTNU             ;executes when R20 >= 75
SAME_HI:

BRNE    HI
MOV     R16,R20           ;executes when R20 = 75
RJMP    CNTNU
HI:
MOV     R17,R20           ;executes when R20 > 75
CNTNU:  ....
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### **BRGE and BRLT instructions**

The BRGE makes decisions based on the S flag. If  $S = 0$  (which, after the CP instruction for signed numbers, means that the left-hand operand of the CP instruction was greater than or equal to the right-hand operand) the BRGE instruction branches in a forward or backward direction relative to program counter.

The BRLT is like the BRGE, but it branches when  $S = 1$ .  
**Notice that the BRGE, and the BRLT are used with signed numbers.**

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.3: LOGIC AND COMPARE INSTRUCTIONS

### BRVS and BRVC instructions

As we mentioned before, the V (overflow) flag must be monitored by the programmer to detect overflow and handle the error. The BRVC and BRVS instructions let you check the value of the V flag and change the flow of the program if overflow has occurred.

#### Example 5-27

Write a program to add two signed numbers. The numbers are in R21 and R22. The program should store the result in R21. If the result is not correct, the program should put 0xAA on PORTA and clear R21.

#### Solution:

```
LDI    R21,0xFA           ;R21 = 0xFA
LDI    R22,0x05           ;R22 = 0x05
LDI    R23,0xFF           ;R23 = 0xFF
OUT    DDRA,R23           ;Port A is output
ADD    R21,R22             ;R21 = R21 + R22
BRVC   NEXT              ;if V = 0 ( no error) then go to next
LDI    R23,0xAA           ;R23 = 0xAA
OUT    PORTA,R23          ;send 0xAA to PORTA
LDI    R21,0x00           ;clear R21
```

NEXT: ...

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

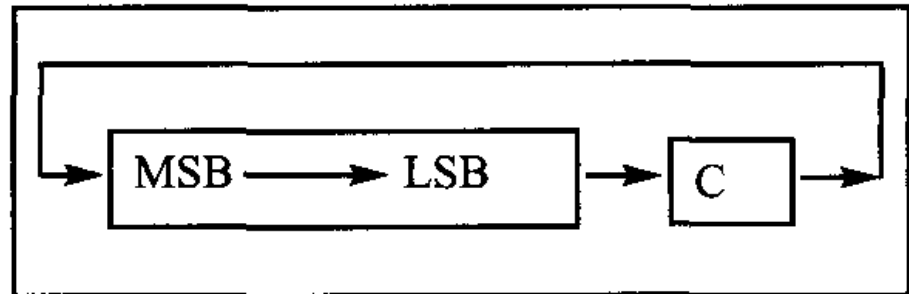
### Rotating through the carry

There are two rotate instructions in the AVR. They involve the carry flag.

#### *ROR instruction*

`ROR Rd ;rotate Rd right through carry`

In the ROR, as bits are rotated from left to right, the carry flag enters the MSB, and the LSB exits to the carry flag. In other words, in ROR the C is moved to the MSB, and the LSB is moved to the C. In reality, the carry flag acts as if it is part of the register, making it a 9-bit register.

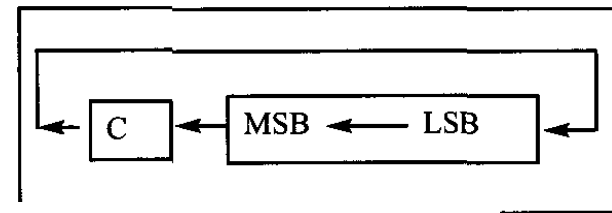


# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### ROL instruction

The other rotating instruction is ROL. In ROL, as bits are shifted from right to left, the carry flag enters the LSB, and the MSB exits to the carry flag. In other words, in ROL the C is moved to the LSB, and the MSB is moved to the C. Again, the carry flag acts as if it is part of the register, making it a 9-bit register.



Examine the following code.

|     |           |                        |
|-----|-----------|------------------------|
| SEC |           | ;make C = 1            |
| LDI | R20, 0x15 | ;R20 = 0001 0101       |
| ROL | R20       | ;R20 = 0010 1011 C = 0 |
| ROL | R20       | ;R20 = 0101 0110 C = 0 |
| ROL | R20       | ;R20 = 1010 1100 C = 0 |
| ROL | R20       | ;R20 = 0101 1000 C = 1 |

# ARITHMATIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### Serializing data

Serializing data is a way of sending a byte of data one bit at a time through a single pin of the microcontroller. There are two ways to transfer a byte of data serially:

1. Using the serial port. In using the serial port, programmers have very limited control over the sequence of data transfer.
2. The second method of serializing data is to transfer data one bit at a time and control the sequence of data and spaces between them. In many new generations of devices such as LCD, ADC, and ROM, the serial versions are becoming popular because they take up less space on a printed circuit board.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### Serializing a byte of data

Serializing data is one of the most widely used applications of the rotate instruction. We can use the rotate instruction to transfer a byte of data serially (one bit at a time). Shift instructions can be used for the same job.

#### Example 5-28

Write a program to transfer the value 41H serially (one bit at a time) via pin PB1. Put one high at the start and end of the data. Send the LSB first.

#### Solution:

```
.INCLUDE "M32DEF.INC"

        SBI    DDRB, 1           ;bit 1 of Port B is output
        LDI    R20, 0x41         ;R20 = the value to be sent

        CLC                     ;clear carry flag
        LDI    R16, 8            ;R16 = 8
        SBI    PORTB, 1          ;bit 1 of PORTB is 1

AGAIN:   ROR    R20               ;rotate right R20 (send LSB to C flag)
        BRCS   ONE              ;if C = 1 then go to ONE
        CBI    PORTB, 1         ;bit 1 of PORTB is cleared to zero
        JMP    NEXT             ;go to NEXT
ONE:     SBI    PORTB, 1         ;bit 1 of PORTB is set to one
NEXT:    DEC    R16              ;decrement R16
        BRNE   AGAIN           ;if R16 is not zero then go to AGAIN
        SBI    PORTB, 1         ;bit 1 of PORTB is set to one

HERE:    JMP    HERE            ;RB1 = high
```



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

Example 5-29 also shows how to bring in a byte of data serially.

### Example 5-29

Write a program to bring in a byte of data serially via pin RC7 and save it in R20 register. The byte comes in with the LSB first.

#### Solution:

```
.INCLUDE      "M32DEF.INC"

        CBI    DDRC, 7      ;bit 7 of Port C is input
        LDI     R16, 8       ;R16 = 8
        LDI     R20, 0       ;R20 = 0
AGAIN:
        SBIC    PINC, 7      ;skip the next line if bit 7 of Port C is 0
        SEC                      ;set carry flag to one
        SBIS    PINC, 7      ;skip the next line if bit 7 of Port C is 1
        CLC                      ;clear carry flag to zero
        ROR     R20          ;rotate right R20. move C flag to MSB of R21
        DEC     R16          ;decrement R16
        BRNE    AGAIN        ;if R16 is not zero go to AGAIN

HERE:    JMP     HERE        ;stop here
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### Example 5-30

Write a program that finds the number of 1s in a given byte.

#### Solution:

```
.INCLUDE      "M32DEF.INC"

        LDI    R20, 0x97
        LDI    R30, 0      ;number of 1s
        LDI    R16, 8      ;number of bits in a byte

AGAIN:
        ROR    R20          ;rotate right R20 and move LSB to C flag
        BRCC   NEXT        ;if C = 0 then go to NEXT
        INC    R30          ;increment R30
NEXT:
        DEC    R16          ;decrement R16
        BRNE   AGAIN        ;if R16 is not zero then go to AGAIN

        ROR    R20          ;one more time to leave R20 unchanged

HERE:   JMP    HERE        ;stop here
```

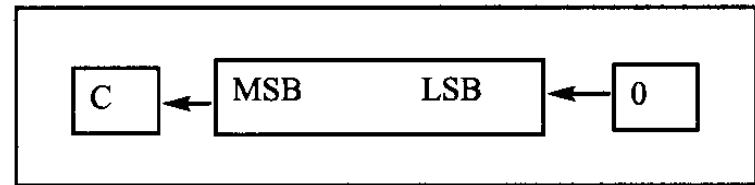
# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### Shift instructions

There are three shift instructions in the AVR. All of them involve the carry flag.

#### *LSL instruction*



`LSL Rd ;logical shift left`

In LSL, as bits are shifted from right to left, enters the LSB, and the MSB exits to the carry flag. In other words, in LSL, is moved to the LSB, and the MSB is moved to the C flag. Notice that this instruction multiplies the content of the register by 2 assuming that after LSL the carry flag is not set.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

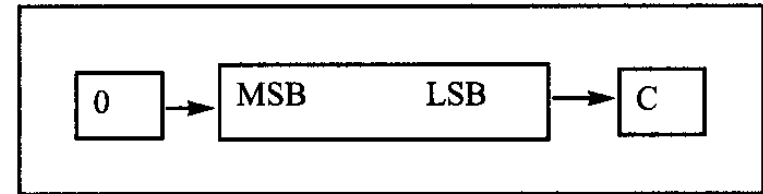
## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

Examine the following code.

```
CLC                ;make C = 0 (carry is 0 )
LDI                R20,0x26    ;R20 = 0010 0110(38)    c = 0
LSL                R20        ;R20 = 0100 1100(76)    C = 0
LSL                R20        ;R20 = 1001 1000(152)   C = 0
LSL                R20        ;R20 = 0011 0000(48)    C = 1
                    ;as C = 1 and content of R20
                    ;is not multiplied by 2
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION



### LSR instruction

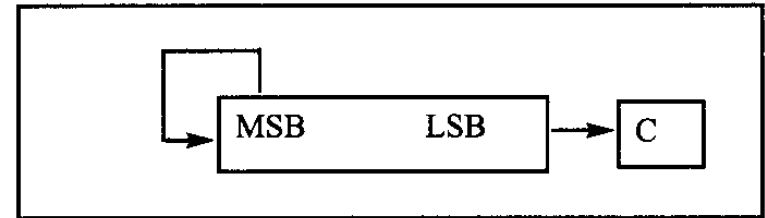
In LSR, as bits are shifted from left to right, 0 enters the MSB, and the LSB exits to the carry flag. In other words, in LSR, 0 is moved to the MSB, and the LSB is moved to the C flag. Notice that this instruction divides the content of the register by 2 and the carry flag contains the remainder of the division. Examine the following code.

|     |           |                        |       |
|-----|-----------|------------------------|-------|
| LDI | R20, 0x26 | ; R20 = 0010 0110 (38) |       |
| LSR | R20       | ; R20 = 0001 0011 (19) | C = 0 |
| LSR | R20       | ; R20 = 0000 1001 (9)  | C = 1 |
| LSR | R20       | ; R20 = 0000 0100 (4)  | C = 1 |

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### *ASR instruction*



This instruction is arithmetic shift right. The ASR instruction can divide signed numbers by two. In ASR, as bits are shifted from left to right, the MSB is held constant and the LSB exits to the carry flag. Examine the following code.

|     |           |                         |       |
|-----|-----------|-------------------------|-------|
| LDI | R20, 0D60 | ; R20 = 1101 0000 (-48) | C = 0 |
| LSR | R20       | ; R20 = 1110 1000 (-24) | C = 0 |
| LSR | R20       | ; R20 = 1111 0100 (-12) | C = 0 |
| LSR | R20       | ; R20 = 1111 1010 (-6)  | C = 0 |
| LSR | R20       | ; R20 = 1111 1101 (-3)  | C = 0 |
| LSR | R20       | ; R20 = 1111 1110 (-1)  | C = 1 |

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### Example 5-31

Assume that R20 has the number -6. Show that LSR cannot be used to divide the content of R20 by 2. Why?

#### Solution:

```
LDI    R20, 0xFA      ;R20 = 1111 1010 (-6)
LSR     R20            ;R20 = 0111 1101 (+125)
                        ;-6 divided by 2 is not +125 and
                        ;the answer is not correct
```

Because LSR shifts the sign bit it changes the sign of the number and therefore cannot be used for signed numbers.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### Example 5-32

Assume that R20 has the number 48. Show how we can use ROR to divide R20 by 8.

#### Solution:

```
                                ;to divide a number by 8 we can
                                ;shift it 3 bits to the right. without
                                ;LSR we have to ROR 3 times and
                                ;clear carry flag before
                                ;each rotation

LDI    R20,0x30                ;R20 = 0011 0000 (48)
CLC                                          ;clear carry flag
ROR     R20                     ;R20 = 0001 1000 (24)
CLC                                          ;clear carry flag
ROR     R20                     ;R20 = 0000 1100 (12)
CLC                                          ;clear carry flag
ROR     R20                     ;R20 = 0000 0110 (6)
                                ;48 divided by 8 is 6 and
                                ;the answer is correct
```



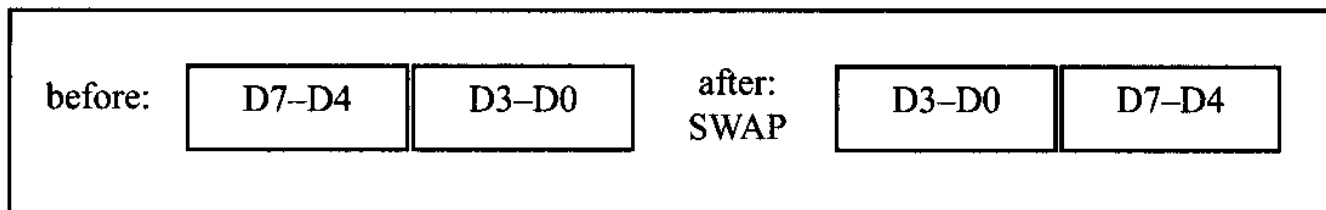
# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### SWAP instruction

```
SWAP Rd      ;swap nibbles
```

Another useful instruction is the SWAP instruction. It works on R0-R31. It swaps the lower nibble and the higher nibble. In other words, the lower 4 bits are put into the higher 4 bits, and the higher 4 bits are put into the lower 4 bits.



# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.4: ROTATE AND SHIFT INSTRUCTIONS AND DATA SERIALIZATION

### Example 5-33

(a) Find the contents of the R20 register in the following code .

```
LDI    R20, 0x72
SWAP   R20
```

(b) In the absence of a SWAP instruction, how would you exchange the nibbles?  
Write a simple program to show the process.

### Solution:

(a)

```
LDI    R20, 0x72      ;R20 = 0x72
SWAP   R20             ;R20 = 0x27
```

(b)

```
LDI    R20, 0x72
LDI    R16, 4
LDI    R21, 0
BEGIN:
    CLC
    ROL    R20
    ROL    R21
    DEC    R16
    BRNE   BEGIN
    OR     R20, R21
HERE:   JMP    HERE
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.5: BCD AND ASCII CONVERSION

### BCD (binary coded decimal) number system

BCD stands for binary coded decimal. BCD is needed because in everyday life we use the digits 0 to 9 for numbers, not binary or hex numbers. binary representation of 0 to 9 is called BCD. In computer literature, one encounters two terms for BCD numbers: (1) unpacked BCD, and (2) packed BCD.

| <b><i>Digit</i></b> | <b><i>BCD</i></b> |
|---------------------|-------------------|
| 0                   | 0000              |
| 1                   | 0001              |
| 2                   | 0010              |
| 3                   | 0011              |
| 4                   | 0100              |
| 5                   | 0101              |
| 6                   | 0110              |
| 7                   | 0111              |
| 8                   | 1000              |
| 9                   | 1001              |

**Figure 5-3. BCD Code**

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.5: BCD AND ASCII CONVERSION

### Unpacked BCD

In unpacked BCD, the lower 4 bits of the number represent the BCD number, and the rest of the bits are 0. For example, "0000 1001" and "0000 0101" are unpacked BCD for 9 and 5, respectively. Unpacked BCD requires 1 byte of memory, or an 8-bit register, to contain it.

### Packed BCD

In packed BCD, a single byte has two BCD numbers in it: one in the lower 4 bits, and one in the upper 4 bits. For example, "0101 1001" is packed BCD for 59H. Only 1 byte of memory is needed to store the packed BCD operands. Thus, one reason to use packed BCD is that it is twice as efficient in storing data.

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.5: BCD AND ASCII CONVERSION

### ASCII numbers

On ASCII keyboards, when the key “0” is activated, “011 0000” (30H) is provided to the computer. Similarly, 31H (011 0001) is provided for key “1”, and so on.

**Table 5-3: ASCII and BCD Codes for Digits 0–9**

| Key | ASCII (hex) | Binary   | BCD (unpacked) |
|-----|-------------|----------|----------------|
| 0   | 30          | 011 0000 | 0000 0000      |
| 1   | 31          | 011 0001 | 0000 0001      |
| 2   | 32          | 011 0010 | 0000 0010      |
| 3   | 33          | 011 0011 | 0000 0011      |
| 4   | 34          | 011 0100 | 0000 0100      |
| 5   | 35          | 011 0101 | 0000 0101      |
| 6   | 36          | 011 0110 | 0000 0110      |
| 7   | 37          | 011 0111 | 0000 0111      |
| 8   | 38          | 011 1000 | 0000 1000      |
| 9   | 39          | 011 1001 | 0000 1001      |

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.5: BCD AND ASCII CONVERSION

### Packed BCD to ASCII conversion

In many systems we have what is called a real-time clock (RTC). The RTC provides the time of day (hour, minute, second) and the date (year, month, day) continuously, regardless of whether the power is on or off. This data, however, is provided in packed BCD. For this data to be displayed on a device such as an LCD, or to be printed by the printer, it must be in ASCII format.

To convert packed BCD to ASCII, you must first convert it to unpacked BCD. Then the unpacked BCD is tagged with 011 0000 (30H). The following demonstrates converting packed BCD to ASCII.

#### ***Packed BCD***

29H

0010 1001

#### ***Unpacked BCD***

02H & 09H

0000 0010 &

0000 1001

#### ***ASCII***

32H & 39H

0011 0010 &

0011 1001

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.5: BCD AND ASCII CONVERSION

### Example 5-34

Assume that R20 has packed BCD. Write a program to convert the packed BCD to two ASCII numbers and place them in R21 and R22.

#### Solution:

```
.INCLUDE      "M32DEF.INC"

    LDI       R20,0x29      ;the packed BCD to be converted is 29

    MOV       R21,R20       ;R21 = R20 = 29H
    ANDI      R21,0x0F      ;mask the upper nibble (R21 = 09H)
    ORI       R21,0x30      ;make it ASCII (R21 = 39H)

    MOV       R22,R20       ;R22 = R20 = 29H
    SWAP      R22           ;swap nibbles (R22 = 92H)
    ANDI      R22,0x0F      ;mask the upper nibble (R22 = 02)
    ORI       R22,0x30      ;make it ASCII (R22 = 32H)

    HERE:     JMP          HERE
```

# ARITHMETIC, LOGIC INSTRUCTIONS AND PROGRAMS

## 5.5: BCD AND ASCII CONVERSION

### ASCII to packed BCD conversion

To convert ASCII to packed BCD, you first convert it to unpacked BCD (to get rid of the 3), and then combine it to make packed BCD. For example, for 4 and 7 the keyboard gives 34 and 37, respectively. The goal is to produce 47H or "0100 0111", which is packed BCD.

| <b>Key</b> | <b>ASCII</b> | <b>Unpacked BCD</b>                  | <b>Packed BCD</b>     |
|------------|--------------|--------------------------------------|-----------------------|
| 4          | 34           | 00000100                             |                       |
| 7          | 37           | 00000111                             | 01000111 which is 47H |
| LDI        | R21, '4'     | ;load character 4 to R21             |                       |
| LDI        | R22, '7'     | ;load character 7 to R22             |                       |
| ANDI       | R21, 0x0F    | ;mask upper nibble of R21            |                       |
| SWAP       | R21          | ;swap nibbles of R21                 |                       |
|            |              | ;to make upper nibble of packed BCD  |                       |
| ANDI       | R22, 0x0F    | ;mask upper nibble of R22            |                       |
| OR         | R22, R21     | ;join R22 and R21 to make packed BCD |                       |
| MOV        | R20, R22     | ;move the result to R20              |                       |