

# COMP 3270 Homework Assignment 4

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Feel free to discuss this assignment with others; however, all work should reflect your own understanding and you should never copy another's work. One useful rule of thumb is, "would I be able to explain my answer to someone else?"

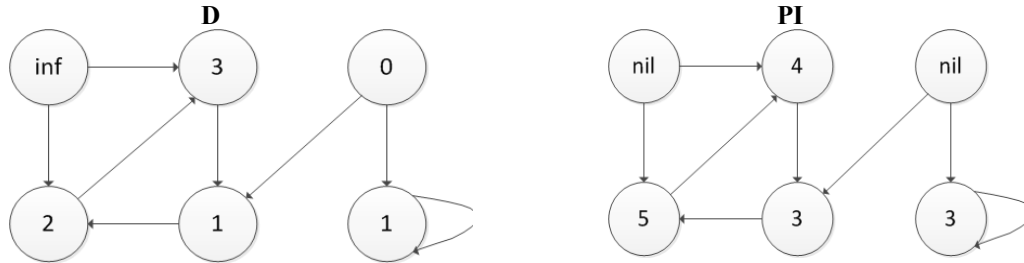
### Problems

#### Chap 22

Exercise 22.1-1 . Given an adjacency-list representation of a directed graph, how long does it take to compute the out-degree of every vertex? How long does it take to compute the in-degrees?

*The computation of the out-degree of a vertex of a directed graph in an adjacency list representation would consist of measuring the length of the linked list at that particular vertex's array index. This could be a linear search along the linked list, or, if there is a "size" field stored, it could be constant time for each index, resulting in  $O(V)$  time for all vertices. The time to compute the in-degree of each vertex will require a search of ALL linked lists in the adjacency list, which means that computing the in-degree of all vertices would take  $O(V * E)$  time.*

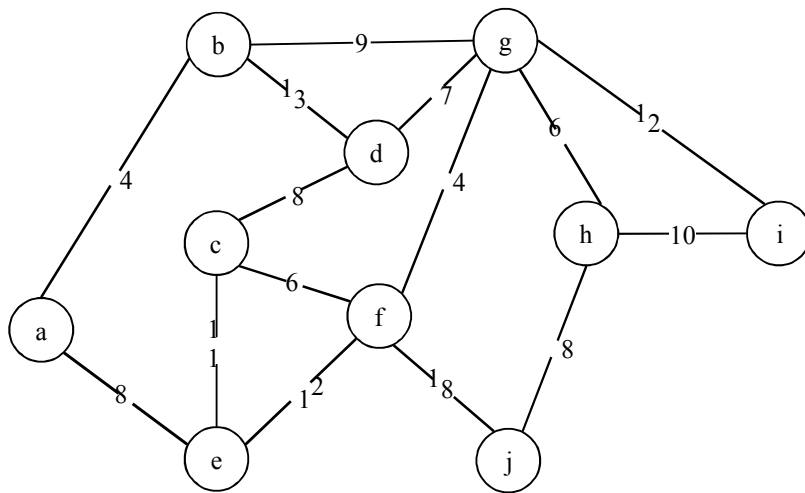
Exercise 22.2-1 Show the d and pi values that result from running breadth-first search on the directed graph of Fig 22.2(a) using vertex 3 as the source.



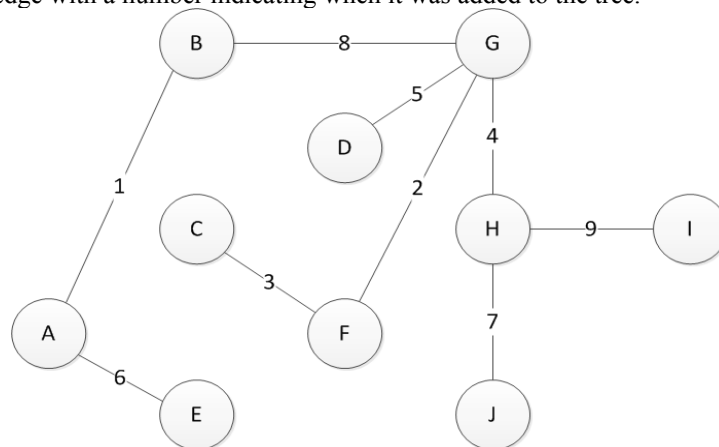
Exercise 22-4-1 Show the ordering of vertices produced by TOPOLOGICAL-SORT when it is run on the dag of Figure 22.8 under the assumption of Exercise 22,3,2, ie vertices are considered in alphabetical order.

	START	FINISH
M	1	20
N	21	26
O	22	25
P	27	28
Q	18	19
R	2	17
S	23	24
T	4	5
U	3	6
V	8	15
W	9	12
X	13	14
Y	7	16
Z	10	11

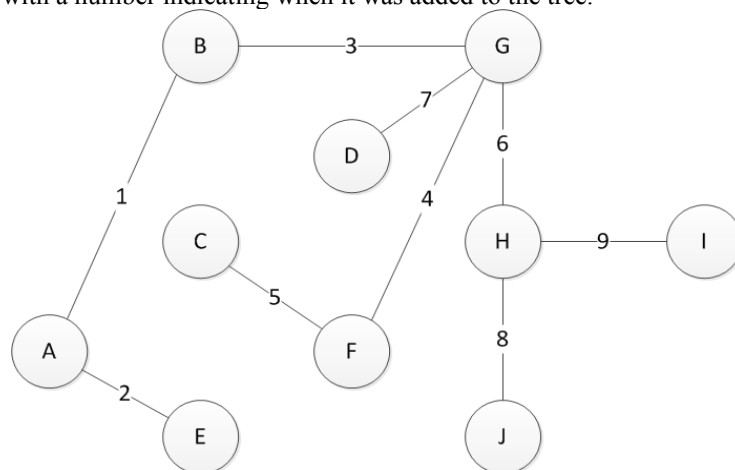
28 26 25 24 20 19 17 16 15 14 12 11 6 5  
 P N O S M Q R Y V X W Z T U



1. Show the minimum spanning tree generated by Kruskal's algorithm on the graph above. Label each tree edge with a number indicating when it was added to the tree.

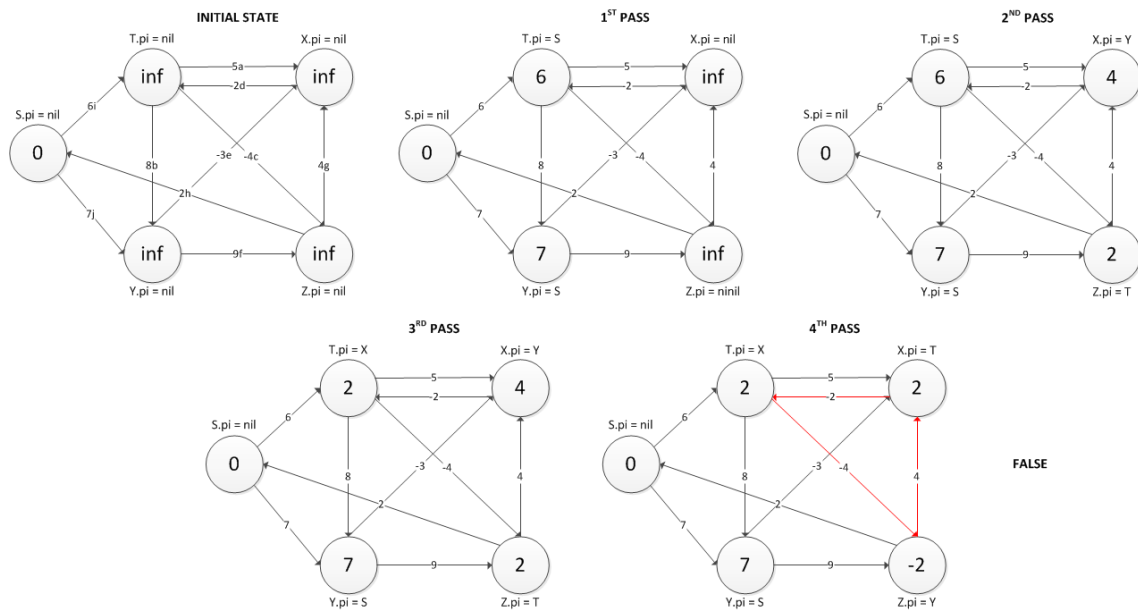
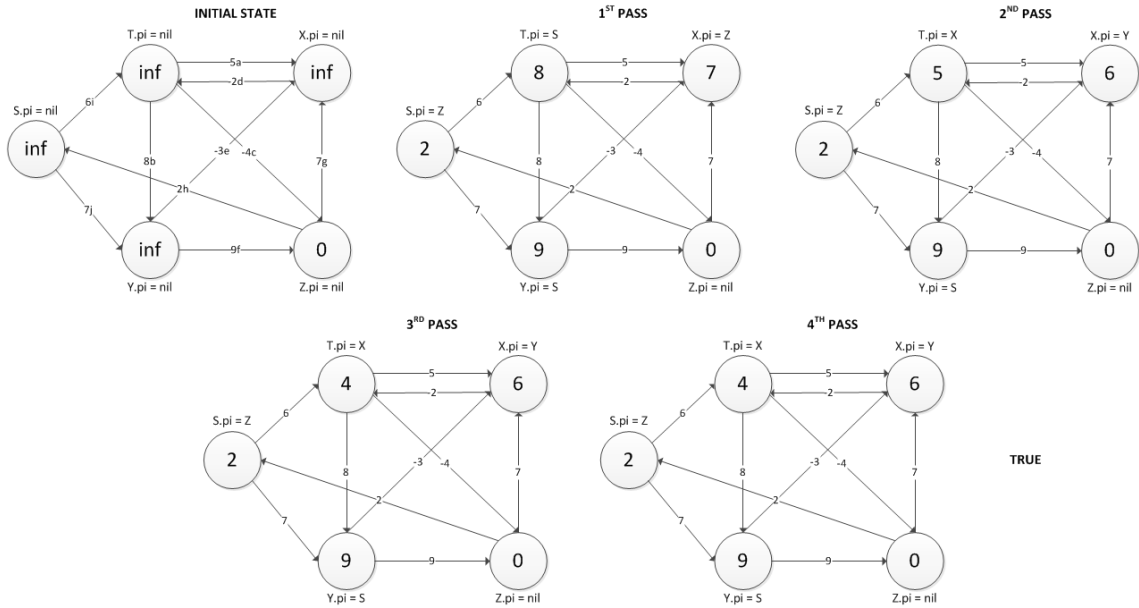


2. Show the minimum spanning tree generated by Prim's algorithm on the graph above. Label each tree edge with a number indicating when it was added to the tree.

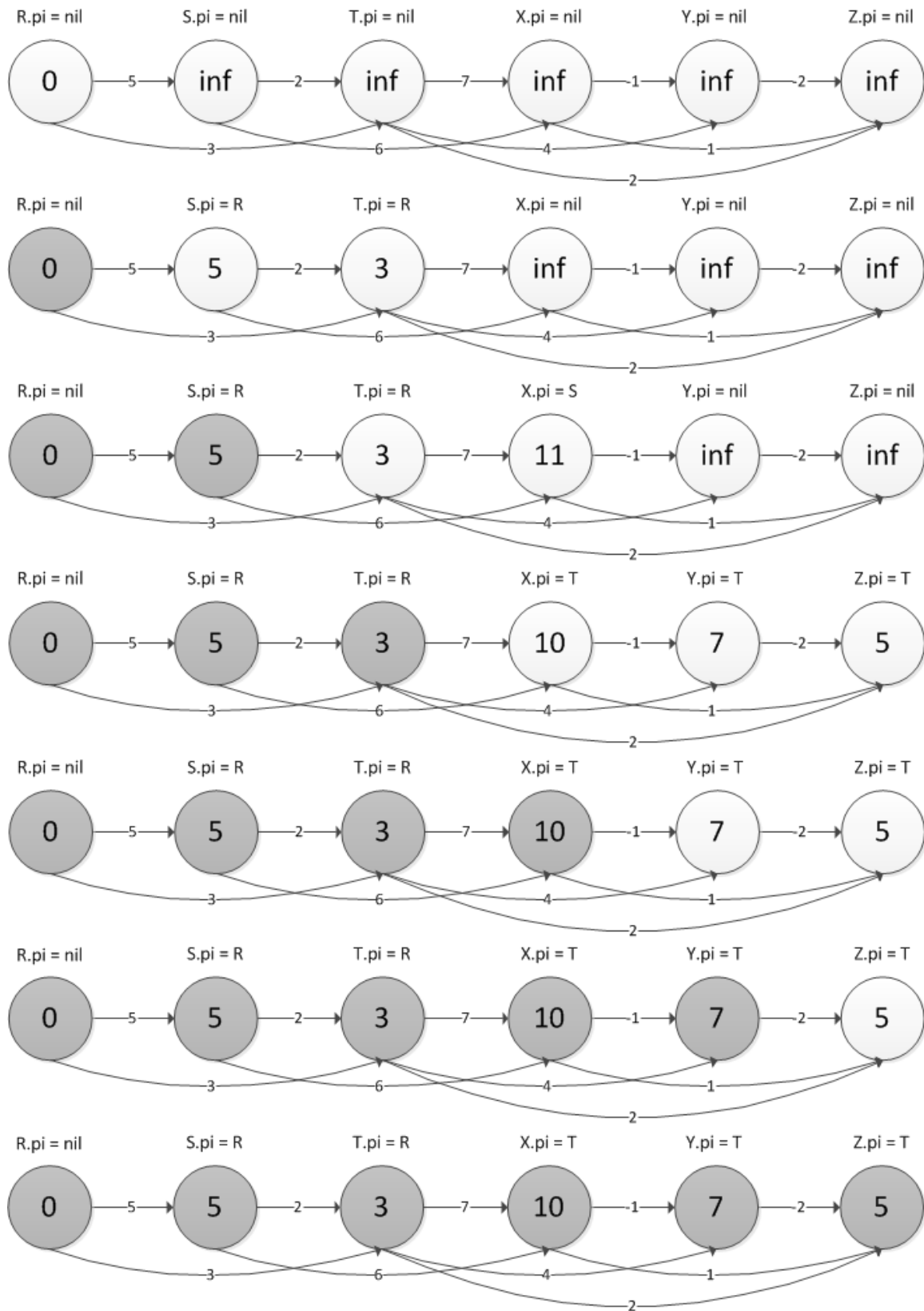


## Chap 24

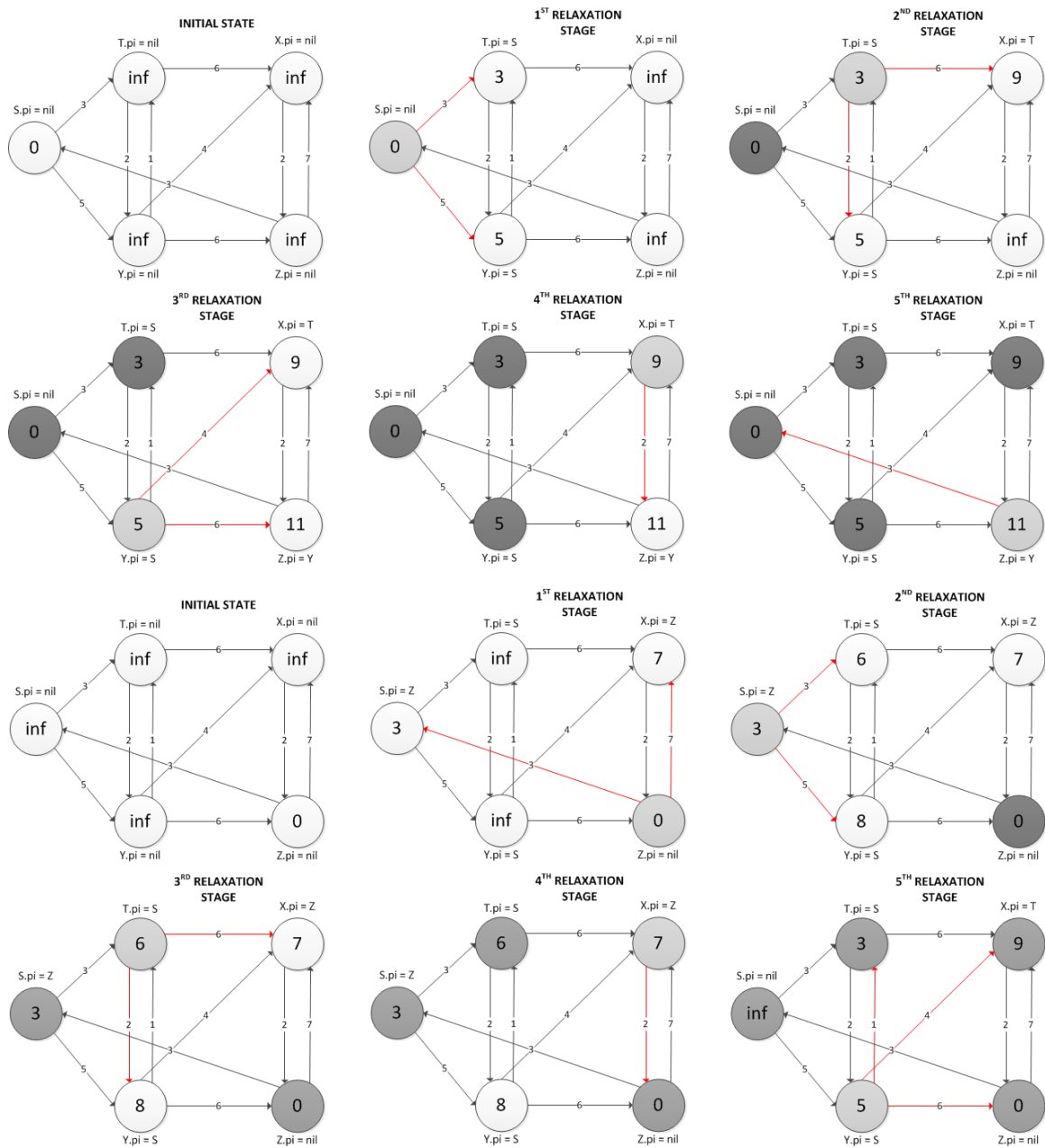
Exercise 24.1-1 Run the Bellman-Ford algorithm on the directed graph of Figure 24.2, using vertex  $z$  as the source. In each pass, relax edges in the same order as in the figure and show the  $d$  and  $pi$  values after each pass. Now, change the weight of edge  $(z,x)$  to 4 and run the algorithm again using  $s$  as the source.



Exercise 24.2-1 Run DAG\_SHORTEST\_PATHS on the directed graph of Figure 24.5 using vertex  $r$  as the source.

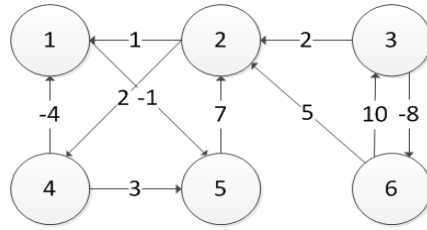


Exercise 24.3-1 Run Dijkstra's algorithm on the digraph of Figure 24.2, using vertex  $s$  as the source and then using vertex  $z$  as the source. In the style of Figure 24.6, show the  $d$  and  $\pi$  values on the vertices in the Set  $S$  after each iteration of the while loop.



Chap 25

Exercise 25.1-1 Run SLOW-ALL-PAIRS-SHORTEST-PATH on the weighted, directed graph of Figure 25.2, showing the matrices that result for each iteration of the loop. Then do the same for FASTER-ALL-PAIRS-SHORTEST-PATH.



**L<sup>1</sup> SLOW-ALL-PAIRS-SHORTEST-PATH**

0	inf	inf	inf	-1	inf
1	0	inf	2	inf	inf
inf	2	0	inf	inf	-8
-4	inf	inf	0	3	inf
inf	7	inf	inf	0	inf
inf	5	10	inf	inf	0

**L<sup>2</sup>**

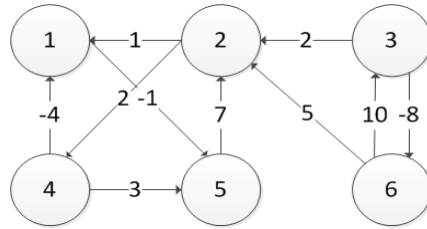
0	6	inf	inf	-1	inf
-2	0	inf	2	0	inf
3	-3	0	4	inf	-8
-4	10	inf	0	-5	inf
8	7	inf	9	0	inf
6	5	10	7	inf	0

**L<sup>3</sup>**

0	6	inf	8	-1	inf
-2	0	inf	2	-3	inf
-2	-3	0	-2	2	-8
-4	2	inf	0	-5	inf
5	7	inf	9	0	inf
3	5	10	7	10	0

**L<sup>4</sup> = L<sup>5</sup>**

0	6	inf	8	-1	inf
-2	0	inf	2	-3	inf
-5	-3	0	-1	-3	-8
-4	2	inf	0	-5	inf
5	7	inf	9	0	inf
3	5	10	7	2	0



**L<sup>1</sup> FASTER-ALL-PAIRS-SHORTEST-PATH**

0	inf	inf	inf	-1	inf
1	0	inf	2	inf	inf
inf	2	0	inf	inf	-8
-4	inf	inf	0	3	inf
inf	7	inf	inf	0	inf
inf	5	10	inf	inf	0

**L<sup>2</sup>**

0	6	inf	inf	-1	inf
-2	0	inf	2	0	inf
3	-3	0	4	inf	-8
-4	10	inf	0	-5	inf
8	7	inf	9	0	inf
6	5	10	7	inf	0

**L<sup>4</sup> = L<sup>5</sup>**

0	6	inf	8	-1	inf
-2	0	inf	2	-3	inf
-5	-3	0	-1	-3	-8
-4	2	inf	0	-5	inf
5	7	inf	9	0	inf
3	5	10	7	2	0

**Chap 26**

Exercise 26.2-2 In Figure 26.1(b) what is the flow across the cut  $(\{s, v_2, v_4\}, \{v_1, v_3, t\})$ ? What is the capacity of this cut?

**The flow across the cut,  $f = 11 + 1 + 7 + 4 - 4 = 19$ .**

**The capacity of the cut,  $c = 16 + 4 + 7 + 4 = 31$ .**

**Chap 34**

Exercise 34.4-6 Suppose someone gives you a polynomial-time algorithm to decide formula satisfiability. Describe how to use this algorithm to find satisfying assignments in polynomial time.

**A formula is satisfiable if it can be shown that there exists some combination of assignments to the boolean arguments  $x_1, \dots, x_m$  for which the boolean formula evaluates to *true*. If a polynomial-time algorithm to decide formula satisfiability is given, then said algorithm must provide a mechanism for finding such a combination of assignments in polynomial time. If the output of the algorithm is modified so that the algorithm no longer acts as a decision algorithm (simply returning *true* or *false*), and instead returns the satisfying assignment that was found for the given boolean formula (if any), then the algorithm can be used to find satisfying assignments in polynomial time.**