

SMART CONTRACT AUDIT REPORT

for

SGNv2 Message

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Contents

| 1 | Intr | Introduction | | | | |
|----|--------|--|----|--|--|--|
| | 1.1 | About SGNv2 Message | 4 | | | |
| | 1.2 | About PeckShield | 5 | | | |
| | 1.3 | Methodology | 5 | | | |
| | 1.4 | Disclaimer | 7 | | | |
| 2 | Find | dings | 9 | | | |
| | 2.1 | Summary | 9 | | | |
| | 2.2 | Key Findings | 10 | | | |
| 3 | Det | ailed Results | 11 | | | |
| | 3.1 | Improper Logic Of MessageBusSender::withdrawFee() | 11 | | | |
| | 3.2 | Improper Logic Of MessageBusReceiver::verifyTransfer() | 12 | | | |
| | 3.3 | Meaningful Events For Important State Changes | 14 | | | |
| | 3.4 | Trust Issue Of Admin Keys | 15 | | | |
| 4 | Con | clusion | 17 | | | |
| Re | eferer | nces | 18 | | | |

1 Introduction

Given the opportunity to review the design document and related smart contract source code of the SGNv2 Message, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

1.1 About SGNv2 Message

Celer Network is an Internet-scale, trust-free, and privacy-preserving platform where everyone can quickly build, operate, and use highly scalable decentralized applications. Celer Network provides unprecedented performance and flexibility through innovation in off-chain scaling techniques and incentive-aligned crypto-economics. SGNv2 is an important part of Celer Network, which supports cross-chain transaction between different blockchains. In particular, the audited Message provides message cross-chain service, which brings more flexibility for cross-chain transactions.

Item Description
Target SGNv2 Message
Type Smart Contract
Language Solidity
Audit Method Whitebox
Latest Audit Report February 20, 2022

Table 1.1: Basic Information of SGNv2 Message

In the following, we show the Git repository of reviewed files and the commit hash value used in this audit. Please note that this audit only covers the contracts in the contracts/message directory and the contracts/message/apps sub-directory is out of the audit scope.

https://github.com/celer-network/sgn-v2-contracts.git (fbd4992)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

https://github.com/celer-network/sgn-v2-contracts.git (71a1955)

1.2 About PeckShield

PeckShield Inc. [9] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

Medium High High Impact Medium Medium Low High Medium Low Low Low High Medium Low Likelihood

Table 1.2: Vulnerability Severity Classification

1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [8]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild:
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

Table 1.3: The Full List of Check Items

| Category | Check Item |
|-----------------------------|---|
| | Constructor Mismatch |
| | Ownership Takeover |
| | Redundant Fallback Function |
| | Overflows & Underflows |
| | Reentrancy |
| | Money-Giving Bug |
| | Blackhole |
| | Unauthorized Self-Destruct |
| Basic Coding Bugs | Revert DoS |
| Dasic Couling Dugs | Unchecked External Call |
| | Gasless Send |
| | Send Instead Of Transfer |
| | Costly Loop |
| | (Unsafe) Use Of Untrusted Libraries |
| | (Unsafe) Use Of Predictable Variables |
| | Transaction Ordering Dependence |
| | Deprecated Uses |
| Semantic Consistency Checks | Semantic Consistency Checks |
| | Business Logics Review |
| | Functionality Checks |
| | Authentication Management |
| | Access Control & Authorization |
| | Oracle Security |
| Advanced DeFi Scrutiny | Digital Asset Escrow |
| Advanced Deri Scrutilly | Kill-Switch Mechanism |
| | Operation Trails & Event Generation |
| | ERC20 Idiosyncrasies Handling |
| | Frontend-Contract Integration |
| | Deployment Consistency |
| | Holistic Risk Management |
| | Avoiding Use of Variadic Byte Array |
| | Using Fixed Compiler Version |
| Additional Recommendations | Making Visibility Level Explicit |
| | Making Type Inference Explicit |
| | Adhering To Function Declaration Strictly |
| | Following Other Best Practices |

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- <u>Basic Coding Bugs</u>: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [7], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings.

1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

| Category | Summary |
|------------------------------|---|
| Configuration | Weaknesses in this category are typically introduced during |
| | the configuration of the software. |
| Data Processing Issues | Weaknesses in this category are typically found in functional- |
| | ity that processes data. |
| Numeric Errors | Weaknesses in this category are related to improper calcula- |
| | tion or conversion of numbers. |
| Security Features | Weaknesses in this category are concerned with topics like |
| | authentication, access control, confidentiality, cryptography, |
| | and privilege management. (Software security is not security |
| | software.) |
| Time and State | Weaknesses in this category are related to the improper man- |
| | agement of time and state in an environment that supports |
| | simultaneous or near-simultaneous computation by multiple |
| 5 C IV | systems, processes, or threads. |
| Error Conditions, | Weaknesses in this category include weaknesses that occur if |
| Return Values, | a function does not generate the correct return/status code, |
| Status Codes | or if the application does not handle all possible return/status |
| Describe Management | codes that could be generated by a function. |
| Resource Management | Weaknesses in this category are related to improper manage- |
| Behavioral Issues | ment of system resources. |
| Denavioral issues | Weaknesses in this category are related to unexpected behaviors from code that an application uses. |
| Business Logics | Weaknesses in this category identify some of the underlying |
| Dusilless Logics | problems that commonly allow attackers to manipulate the |
| | business logic of an application. Errors in business logic can |
| | be devastating to an entire application. |
| Initialization and Cleanup | Weaknesses in this category occur in behaviors that are used |
| mitialization and Cicanap | for initialization and breakdown. |
| Arguments and Parameters | Weaknesses in this category are related to improper use of |
| / inguinents and i diameters | arguments or parameters within function calls. |
| Expression Issues | Weaknesses in this category are related to incorrectly written |
| | expressions within code. |
| Coding Practices | Weaknesses in this category are related to coding practices |
| 3 | that are deemed unsafe and increase the chances that an ex- |
| | ploitable vulnerability will be present in the application. They |
| | may not directly introduce a vulnerability, but indicate the |
| | product has not been carefully developed or maintained. |

2 | Findings

2.1 Summary

Here is a summary of our findings after analyzing the SGNv2 Message implementation. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logic, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

| Severity | # of Findings |
|---------------|---------------|
| Critical | 0 |
| High | 2 |
| Medium | 0 |
| Low | 1 |
| Informational | 1 |
| Total | 4 |

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities that need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 2 high-severity vulnerabilities, 1 low-severity vulnerability, and 1 informational recommendation.

ID Title Severity Category **Status** PVE-001 High **Improper** Logic Of Message-**Business Logic** Fixed BusSender::withdrawFee() **PVE-002** Improper Logic Of MessageBusRe-High **Business Logic** Fixed ceiver::verifyTransfer() **PVE-003** Informational Meaningful Events For Important Coding Practices Fixed State Changes **PVE-004** Low Trust Issue Of Admin Keys Security Features Confirmed

Table 2.1: Key SGNv2 Message Audit Findings

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

3 Detailed Results

3.1 Improper Logic Of MessageBusSender::withdrawFee()

ID: PVE-001Severity: High

• Likelihood: High

• Impact: High

• Target: MessageBusSender

Category: Business Logic [6]

• CWE subcategory: CWE-841 [3]

Description

By design, the MessageBusSender contract is one of the core entries of the SGNv2 Message, which supports arbitrary message transfer along with carrying out cross-chain assets transactions. Additionally, the transaction fee is related to the length of the message. In particular, one entry routine, i.e., withdrawFee(), is designed to withdraw the fee to the specified account. While examining its logic, we notice there is an improper implementation that needs to be improved.

To elaborate, we show below the related code snippet of the MessageBusSender contract. In the withdrawFee() routine, the input _cumulativeFee parameter stores the cumulative fees of the account (specified by the input _account parameter), and the withdrawnFees[_account] storage variable stores the fees that the _account has withdrawn. The statement (i.e., uint256 amount = _cumulativeFee - withdrawnFees[_account]) (line 92) is executed to calculate the fees that the _account can withdraw this time. However, we notice the withdrawnFees[_account] storage variable is not updated anywhere in the routine, which results in that the _account may withdraw unexpected fees. Given this, we suggest to update the withdrawnFees[_account] variable as below: withdrawnFees[_account] += amount (line 93).

```
function withdrawFee(

address _account,

uint256 _cumulativeFee,

bytes[] calldata _sigs,

address[] calldata _signers,

uint256[] calldata _powers
```

```
89
       ) external {
90
            bytes32 domain = keccak256(abi.encodePacked(block.chainid, address(this), "
                withdrawFee"));
91
            sigsVerifier.verifySigs(abi.encodePacked(domain, _account, _cumulativeFee),
                _sigs, _signers, _powers);
92
            uint256 amount = _cumulativeFee - withdrawnFees[_account];
93
            require(amount > 0, "No new amount to withdraw");
94
            (bool sent, ) = _account.call{value: amount, gas: 50000}("");
95
            require(sent, "failed to withdraw fee");
96
```

Listing 3.1: MessageBusSender::withdrawFee()

Recommendation Correct the implementation of the withdrawFee() routine by properly updating the withdrawnFees[_account] storage variable.

Status The issue has been addressed by the following commit: ce081e8.

3.2 Improper Logic Of MessageBusReceiver::verifyTransfer()

ID: PVE-002Severity: High

• Likelihood: High

• Impact: High

Target: MessageBusReceiver

• Category: Business Logic [6]

CWE subcategory: CWE-841 [3]

Description

As mentioned in Section 3.1, SGNv2 Message supports arbitrary message transfer along with carrying out cross-chain assets transactions. The MessageBusReceiver contract is designed to deal with the cross-chain message on the destination chain. By design, only when the cross-chain assets transfer has finished on the destination chain, the corresponding message can be dealt with. In particular, one internal routine, i.e., MessageBusReceiver::verifyTransfer(), is called inside the executeMessageWithTransfer() routine to verify whether the cross-chain assets transfer associated with the message has finished. While examining its logic, we notice there is an improper implementation that needs to be improved.

To elaborate, we show below the related code snippet of the MessageBusReceiver contract. We notice the verifyTransfer() routine is compatible with three kinds of cross-chain bridges (The bridge-related contracts are stored in liquidityBridge, pegBridge/pegVault, and pegBridgeV2/pegVaultV2.) that the SGNv2 protocol supports. However, it comes to our attention that the pegBridge/pegVault rather than pegBridgeV2/pegVaultV2 are incorrectly used to support the last kind of cross-chain bridge, which directly undermines the assumption of the design.

```
263
         function verifyTransfer(TransferInfo calldata _transfer) private view returns (
             bytes32) {
264
             bytes32 transferId;
265
             address bridgeAddr;
266
             if (_transfer.t == TransferType.LqSend) {
267
268
             } else if (_transfer.t == TransferType.LqWithdraw) {
269
270
             } else if (_transfer.t == TransferType.PegMint _transfer.t == TransferType.
                 PegWithdraw) {
271
                 . . .
272
             } else if (_transfer.t == TransferType.PegMintV2 _transfer.t == TransferType.
                 PegWithdrawV2) {
273
                 if (_transfer.t == TransferType.PegMintV2) {
274
                     bridgeAddr = pegBridge;
275
                 } else {
276
                     // TransferType.PegWithdrawV2
277
                     bridgeAddr = pegVault;
                 }
278
279
                 transferId = keccak256(
280
                     abi.encodePacked(
281
                          _transfer.receiver,
282
                         _transfer.token,
283
                         _transfer.amount,
284
                         _transfer.sender,
285
                          _transfer.srcChainId,
286
                          _transfer.refId,
287
                         bridgeAddr
288
                     )
289
                 );
290
                 if (_transfer.t == TransferType.PegMintV2) {
291
                     require(IPeggedTokenBridge(bridgeAddr).records(transferId) == true, "
                         mint record not exist");
292
                 } else {
293
                     // \ {\tt TransferType.PegWithdrawV2}
                     require(IOriginalTokenVault(bridgeAddr).records(transferId) == true, "
294
                         withdraw record not exist");
295
                 }
             }
296
297
             return keccak256 (abi.encodePacked (MsgType.MessageWithTransfer, bridgeAddr,
                 transferId));
298
```

Listing 3.2: MessageBusReceiver::verifyTransfer()

Recommendation Correct the implementation of the verifyTransfer() routine.

Status The issue has been addressed by the following commit: 32410b8.

3.3 Meaningful Events For Important State Changes

• ID: PVE-003

• Severity: Informational

Likelihood: N/A

• Impact: N/A

• Target: Multiple Contracts

• Category: Coding Practices [5]

• CWE subcategory: CWE-563 [2]

Description

In Ethereum, the event is an indispensable part of a contract and is mainly used to record a variety of runtime dynamics. In particular, when an event is emitted, it stores the arguments passed in transaction logs and these logs are made accessible to external analytics and reporting tools. Events can be emitted in a number of scenarios. One particular case is when system-wide parameters or settings are being changed. Another case is when tokens are being minted, transferred, or burned.

While examining the events that reflect the protocol dynamics, we notice there are several privileged routines that lack meaningful events to reflect their changes. In the following, we show several representative routines.

```
function setMessageBus(address _messageBus) public onlyOwner {
    messageBus = _messageBus;
}
```

Listing 3.3: MessageBusAddress::setMessageBus()

```
364
         function setLiquidityBridge(address _addr) public onlyOwner {
365
             liquidityBridge = _addr;
366
367
368
         function setPegBridge(address _addr) public onlyOwner {
369
             pegBridge = _addr;
370
371
372
         function setPegVault(address _addr) public onlyOwner {
373
             pegVault = _addr;
374
        }
375
376
         function setPegBridgeV2(address _addr) public onlyOwner {
377
             pegBridgeV2 = _addr;
378
379
380
         function setPegVaultV2(address _addr) public onlyOwner {
381
             pegVaultV2 = _addr;
382
```

Listing 3.4: MessageBusReceiver

```
function setFeePerByte(uint256 _fee) external onlyOwner {
    feePerByte = _fee;
}

function setFeeBase(uint256 _fee) external onlyOwner {
    feeBase = _fee;
}
```

Listing 3.5: MessageBusSender::setFeePerByte()&&setFeeBase()

With that, we suggest to emit meaningful events in these privileged routines. Also, the key event information is better indexed. Note each emitted event is represented as a topic that usually consists of the signature (from a keccak256 hash) of the event name and the types (uint256, string, etc.) of its parameters. Each indexed type will be treated like an additional topic. If an argument is not indexed, it will be attached as data (instead of a separate topic). Considering that the key information is typically queried, it is better treated as a topic, hence the need of being indexed.

Recommendation Properly emit the above-mentioned events with accurate information to timely reflect state changes. This is very helpful for external analytics and reporting tools.

Status The issue has been addressed by the following commit: ce081e8.

3.4 Trust Issue Of Admin Keys

• ID: PVE-004

Severity: Low

Likelihood: Low

• Impact: Low

• Target: Multiple Contracts

• Category: Security Features [4]

CWE subcategory: CWE-287 [1]

Description

In the SGNv2 Message implementation, there is a privileged account that plays a critical role in governing and regulating the protocol-wide operations (e.g., configuring various system parameters). In the following, we show the representative functions potentially affected by the privilege of the account.

```
function setLiquidityBridge(address _addr) public onlyOwner {
    liquidityBridge = _addr;
}

function setPegBridge(address _addr) public onlyOwner {
    pegBridge = _addr;
}

function setPegVault(address _addr) public onlyOwner {
```

```
373
             pegVault = _addr;
374
        }
375
376
         function setPegBridgeV2(address _addr) public onlyOwner {
377
             pegBridgeV2 = _addr;
378
379
380
         function setPegVaultV2(address _addr) public onlyOwner {
381
             pegVaultV2 = _addr;
382
```

Listing 3.6: MessageBusReceiver

We emphasize that the privilege assignment may be necessary and consistent with the protocol design. However, it is worrisome if the privileged account is not governed by a DAO-like structure. Note that a compromised account would allow the attacker to modify a number of sensitive system parameters, which directly undermines the assumption of the protocol design.

Recommendation Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

Status

4 Conclusion

In this audit, we have analyzed the SGNv2 Message design and implementation. Celer Network is an Internet-scale, trust-free, and privacy-preserving platform where everyone can quickly build, operate, and use highly scalable decentralized applications. Celer Network provides unprecedented performance and flexibility through innovation in off-chain scaling techniques and incentive-aligned cryptoeconomics. SGNv2 is an important part of Celer Network, which supports cross-chain transaction between different blockchains. In particular, the audited Message provides message cross-chain service, which brings more flexibility for cross-chain transaction. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and addressed.

Meanwhile, we need to emphasize that smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.

References

- [1] MITRE. CWE-287: Improper Authentication. https://cwe.mitre.org/data/definitions/287.html.
- [2] MITRE. CWE-563: Assignment to Variable without Use. https://cwe.mitre.org/data/definitions/563.html.
- [3] MITRE. CWE-841: Improper Enforcement of Behavioral Workflow. https://cwe.mitre.org/data/definitions/841.html.
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