

# Context-Free Grammars and CKY Algorithm

CS114 Lab 10

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# Review: Regular Languages

- ▶ A language is regular iff:

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  - ▶ It is recognized by a finite state automaton

# Review: Regular Languages

- ▶ A language is regular iff:
  - ▶ It is recognized by a finite state automaton
  - ▶ It is generated by a regular expression

# Natural languages are not regular

- ▶ Why not?

# Context-Free Languages

- ▶ More powerful than regular languages

# Context-Free Languages

- ▶ More powerful than regular languages
- ▶ Used in specification of natural and programming languages

# Context-Free Languages

- ▶ A language is context-free iff:



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- ▶ A language is context-free iff:
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# Context-Free Languages

- ▶ A language is context-free iff:
  - ▶ It is recognized by a pushdown automaton
  - ▶ It is generated by a context-free grammar

# Phrase structure grammars = context-free grammars

- $G = (T, N, S, R)$ 
  - $T$  is set of terminals
  - $N$  is set of nonterminals
    - For NLP, we usually distinguish out a set  $P \subset N$  of preterminals, which always rewrite as terminals
    - $S$  is the start symbol (one of the nonterminals)
    - $R$  is rules/productions of the form  $X \rightarrow \gamma$ , where  $X$  is a nonterminal and  $\gamma$  is a sequence of terminals and nonterminals (possibly an empty sequence)
- A grammar  $G$  generates a language  $L$ .

# Phrase structure grammars = context-free grammars

- ▶  $S \rightarrow aSb$
- ▶  $S \rightarrow \epsilon$

# Phrase structure grammars = context-free grammars

| Grammar                            | Lexicon   |
|------------------------------------|---|
| $S \rightarrow NP VP$              | $Det \rightarrow that \mid this \mid the \mid a$                      |
| $S \rightarrow Aux NP VP$          | $Noun \rightarrow book \mid flight \mid meal \mid money$              |
| $S \rightarrow VP$                 | $Verb \rightarrow book \mid include \mid prefer$                      |
| $NP \rightarrow Pronoun$           | $Pronoun \rightarrow I \mid she \mid me$                              |
| $NP \rightarrow Proper-Noun$       | $Proper-Noun \rightarrow Houston \mid NWA$                            |
| $NP \rightarrow Det Nominal$       | $Aux \rightarrow does$  |
| $Nominal \rightarrow Noun$         | $Preposition \rightarrow from \mid to \mid on \mid near \mid through$ |
| $Nominal \rightarrow Nominal Noun$ |   |
| $Nominal \rightarrow Nominal PP$   |   |
| $VP \rightarrow Verb$              |   |
| $VP \rightarrow Verb NP$           |   |
| $VP \rightarrow Verb NP PP$        |   |
| $VP \rightarrow Verb PP$           |   |
| $VP \rightarrow VP PP$             |   |
| $PP \rightarrow Preposition NP$    |   |

**Figure 13.1** The  $\mathcal{L}_1$  miniature English grammar and lexicon.

Source

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# Probabilistic CFGs

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- ▶ The probability of a parse  $T$  is the product of the probabilities of all of the  $n$  rules used to generate  $T$

- ▶ 
$$P(T) = \prod_{i=1}^n P(\text{RHS}_i|\text{LHS}_i)$$

# Probabilistic CFGs

| Grammar                            |       | Lexicon  |  |
|------------------------------------|-------|--|--|
| $S \rightarrow NP VP$              | [.80] | $Det \rightarrow that$ [.10]   $a$ [.30]   $the$ [.60] |  |
| $S \rightarrow Aux NP VP$          | [.15] | $Noun \rightarrow book$ [.10]   $flight$ [.30]         |  |
| $S \rightarrow VP$                 | [.05] | $meal$ [.05]   $money$ [.05]                           |  |
| $NP \rightarrow Pronoun$           | [.35] | $flight$ [.40]   $dinner$ [.10]                        |  |
| $NP \rightarrow Proper-Noun$       | [.30] | $Verb \rightarrow book$ [.30]   $include$ [.30]        |  |
| $NP \rightarrow Det Nominal$       | [.20] | $prefer$ [.40]   |  |
| $NP \rightarrow Nominal$           | [.15] | $Pronoun \rightarrow I$ [.40]   $she$ [.05]            |  |
| $Nominal \rightarrow Noun$         | [.75] | $me$ [.15]   $you$ [.40]                               |  |
| $Nominal \rightarrow Nominal Noun$ | [.20] | $Proper-Noun \rightarrow Houston$ [.60]                |  |
| $Nominal \rightarrow Nominal PP$   | [.05] | $NWA$ [.40]  |  |
| $VP \rightarrow Verb$              | [.35] | $Aux \rightarrow does$ [.60]   $can$ [.40]             |  |
| $VP \rightarrow Verb NP$           | [.20] | $Preposition \rightarrow from$ [.30]   $to$ [.30]      |  |
| $VP \rightarrow Verb NP PP$        | [.10] | $on$ [.20]   $near$ [.15]                              |  |
| $VP \rightarrow Verb PP$           | [.15] | $through$ [.05]  |  |
| $VP \rightarrow Verb NP NP$        | [.05] |  |  |
| $VP \rightarrow VP PP$             | [.15] |  |  |
| $PP \rightarrow Preposition NP$    | [1.0] |  |  |

**Figure 14.1** A PCFG that is a probabilistic augmentation of the  $\mathcal{L}_1$  miniature English CFG grammar and lexicon of Fig. ???. These probabilities were made up for pedagogical purposes and are not based on a corpus (since any real corpus would have many more rules, so the true probabilities of each rule would be much smaller).

Source

# Chomsky Normal Form

- ▶ Every rule is of the form

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  - ▶  $S \rightarrow \epsilon$
  - ▶  $A \rightarrow BC$
  - ▶  $A \rightarrow a$

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  - ▶  $S \rightarrow \epsilon$
  - ▶  $A \rightarrow BC$
  - ▶  $A \rightarrow a$
- ▶ Where  $S$  is the start symbol,  $A$  is a nonterminal,  $B$  and  $C$  are nonterminals (except for  $S$ ), and  $a$  is a terminal

# Converting a CFG into Chomsky Normal Form

- ▶ Add a new start symbol

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- ▶ Break up rules with more than 3 things on the right hand side

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# Converting a CFG into Chomsky Normal Form

- ▶ Add a new start symbol
- ▶ Remove  $\epsilon$ -rules
- ▶ Remove unary rules
- ▶ Break up rules with more than 3 things on the right hand side
- ▶ Replace terminals with nonterminals and add new rules as needed
  - ▶ We can modify CKY algorithm to handle unary rules

# Parsing Context-Free Languages

- ▶ Bottom-up

# Parsing Context-Free Languages

- ▶ Bottom-up
  - ▶ CKY algorithm

# Parsing Context-Free Languages

- ▶ Bottom-up
  - ▶ CKY algorithm
- ▶ Top-down

# Parsing Context-Free Languages

- ▶ Bottom-up
  - ▶ CKY algorithm
- ▶ Top-down
  - ▶ Earley algorithm

# CKY Algorithm

- ▶ Given tables *table* and *back*:



# CKY Algorithm

- ▶ Given tables *table* and *back*:
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  - ▶ Recursive case

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  - ▶ Base case
    - ▶ Fill in *table* cells  $[i, i + 1]$  with all possible nonterminals that can generate that word
  - ▶ Recursive case
    - ▶ In *table*, if  $A \rightarrow BC$  and  $B$  is in cell  $[i, j]$  and  $C$  is in cell  $[j, k]$ , fill in cell  $[i, k]$  with  $A$

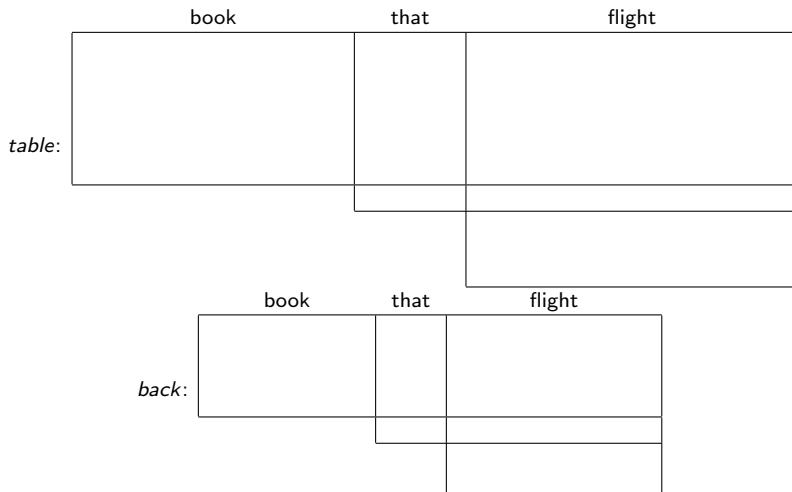
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- ▶ Given tables *table* and *back*:
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    - ▶ Fill in *table* cells  $[i, i + 1]$  with all possible nonterminals that can generate that word
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    - ▶ In *back*, fill in cell  $[i, k]$  with backpointers (e.g.  $A: j, B, C$ )

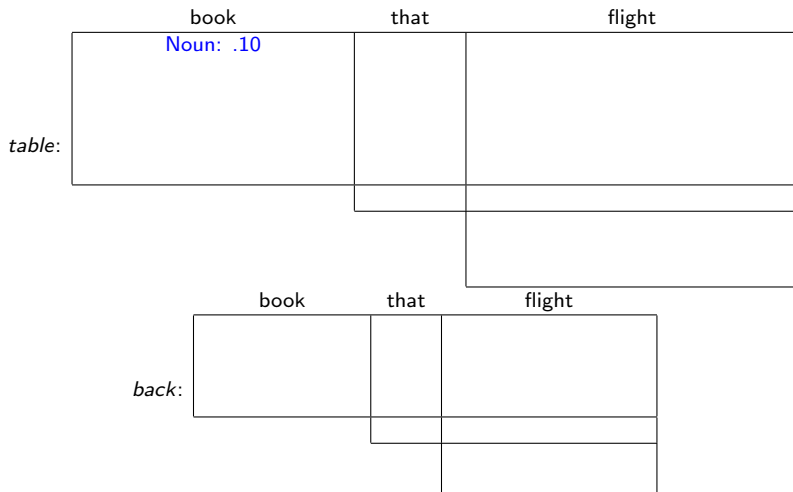
# Probabilistic CKY Algorithm

- ▶ Given tables *table* and *back*:
  - ▶ Base case
    - ▶ Fill in *table* cells  $[i, i + 1]$  with all possible nonterminals that can generate that word, and their probabilities
  - ▶ Recursive case
    - ▶ In *table*, if  $A \rightarrow BC$  and  $B$  is in cell  $[i, j]$  and  $C$  is in cell  $[j, k]$ , and  $table[i, j, A] < P(A \rightarrow BC) \times table[i, j, B] \times table[j, k, C]$ , fill in cell  $[i, k]$  with  $A$ :  
 $P(A \rightarrow BC) \times table[i, j, B] \times table[j, k, C]$
    - ▶ In *back*, fill in cell  $[i, k]$  with backpointers (e.g.  $A: j, B, C$ )

# Probabilistic CKY Algorithm



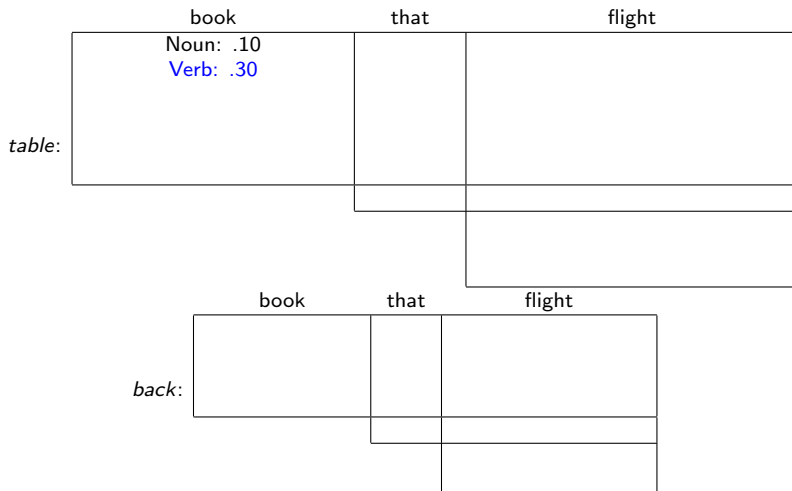
# Probabilistic CKY Algorithm



► Noun → book [.10]

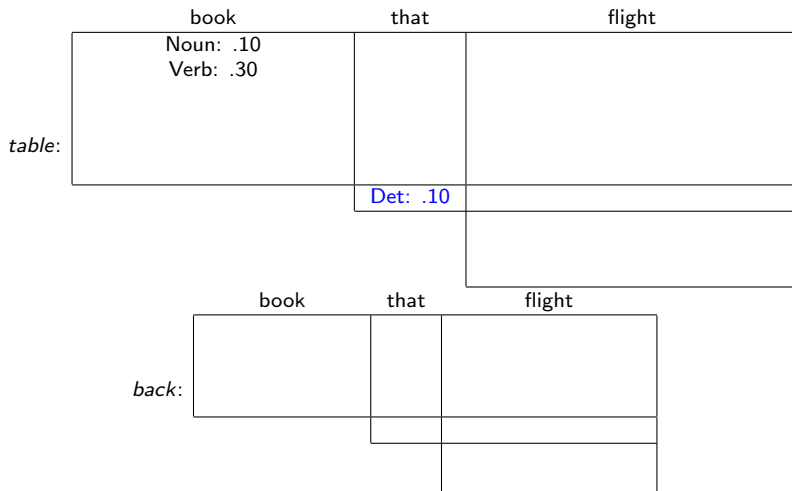


# Probabilistic CKY Algorithm



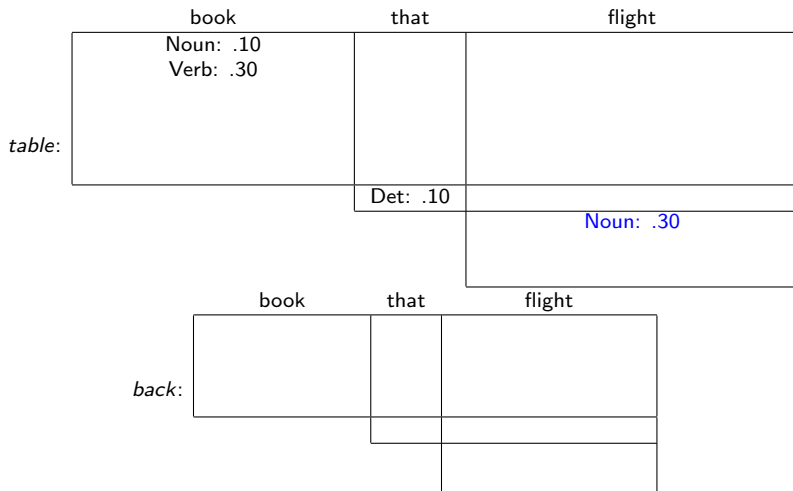
► Verb → book [.30]

# Probabilistic CKY Algorithm



► Det → that [.10]

# Probabilistic CKY Algorithm



► Noun → flight [.30]

# Probabilistic CKY Algorithm

- ▶ Unary rules

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  - ▶ In *table*, if  $A \rightarrow B$  and  $B$  is in cell  $[i, i + 1]$ , fill in cell  $[i, i + 1]$  with  $A$ :  $P(A \rightarrow B) \times \text{table}[i, i + 1, B]$

# Probabilistic CKY Algorithm

- ▶ Unary rules
  - ▶ In *table*, if  $A \rightarrow B$  and  $B$  is in cell  $[i, i + 1]$ , fill in cell  $[i, i + 1]$  with  $A$ :  $P(A \rightarrow B) \times \text{table}[i, i + 1, B]$
  - ▶ In *back*, fill in cell  $[i, i + 1]$  with backpointers (e.g.  $A: B$ )

# Probabilistic CKY Algorithm

|        |  |          |           |
|--------|--|----------|-----------|
|        | book   | that     | flight    |
| table: | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$ |          |           |
|        |  | Det: .10 |           |
|        |  |          | Noun: .30 |

|       |               |      |        |
|-------|---------------|------|--------|
|       | book          | that | flight |
| back: | Nominal: Noun |      |        |
|       |               |      |        |
|       |               |      |        |

► Nominal  $\rightarrow$  Noun [.75]

# Probabilistic CKY Algorithm

|        |  |          |           |
|--------|--|----------|-----------|
|        | book   | that     | flight    |
| table: | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$ |          |           |
|        |  | Det: .10 |           |
|        |  |          | Noun: .30 |
|        | book   | that     | flight    |
| back:  | Nominal: Noun<br>NP: Nominal   |          |           |
|        |  |          |           |
|        |  |          |           |

- NP  $\rightarrow$  Nominal [.15]



# Probabilistic CKY Algorithm

|        |   |          |           |
|--------|---|----------|-----------|
|        | book  | that     | flight    |
| table: | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$<br>VP: $.35 \times .30 = .105$ |          |           |
|        |   | Det: .10 |           |
|        |   |          | Noun: .30 |
|        | book  | that     | flight    |
| back:  | Nominal: Noun<br>NP: Nominal<br>VP: Verb  |          |           |
|        |   |          |           |
|        |   |          |           |

► VP → Verb [.35]

# Probabilistic CKY Algorithm

|        |  |          |           |
|--------|--|----------|-----------|
|        | book   | that     | flight    |
| table: | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$<br>VP: $.35 \times .30 = .105$<br>S: $.05 \times .105 = .00525$ |          |           |
|        |  | Det: .10 |           |
|        |  |          | Noun: .30 |
|        | book   | that     | flight    |
| back:  | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP  |          |           |
|        |  |          |           |
|        |  |          |           |

►  $S \rightarrow VP [.05]$

# Probabilistic CKY Algorithm

|        |  |          |   |
|--------|--|----------|---|
|        | book   | that     | flight  |
| table: | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$<br>VP: $.35 \times .30 = .105$<br>S: $.05 \times .105 = .00525$ |          |   |
|        |  | Det: .10 |   |
|        |  |          | Noun: .30<br>Nominal: $.75 \times .30 = .225$ |
|        | book   | that     | flight  |
| back:  | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP  |          |   |
|        |  |          |   |
|        |  |          | Nominal: Noun                                 |

- Nominal  $\rightarrow$  Noun [.75]

# Probabilistic CKY Algorithm

|        |  |          |   |
|--------|--|----------|---|
|        | book   | that     | flight  |
| table: | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$<br>VP: $.35 \times .30 = .105$<br>S: $.05 \times .105 = .00525$ |          |   |
|        |  | Det: .10 |   |
|        |  |          | Noun: .30<br>Nominal: $.75 \times .30 = .225$<br>NP: $.15 \times .225 = .03375$ |
|        | book   | that     | flight  |
| back:  | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP  |          |   |
|        |  |          |   |
|        |  |          | Nominal: Noun<br>NP: Nominal  |

► NP → Nominal [.15]

# Probabilistic CKY Algorithm

|        |  |          |   |
|--------|--|----------|---|
|        | book   | that     | flight  |
| table: | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$<br>VP: $.35 \times .30 = .105$<br>S: $.05 \times .105 = .00525$ |          |   |
|        |  | Det: .10 | NP: $.20 \times .10 \times .225 = .0045$  |
|        |  |          | Noun: .30<br>Nominal: $.75 \times .30 = .225$<br>NP: $.15 \times .225 = .03375$ |
|        | book   | that     | flight  |
| back:  | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP  |          |   |
|        |  |          | NP: 2 Det Nominal   |
|        |  |          | Nominal: Noun<br>NP: Nominal  |

► NP → Det Nominal [.20]

# Probabilistic CKY Algorithm

|               |  |          |   |
|---------------|--|----------|---|
| <i>table:</i> | book   | that     | flight  |
|               | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$<br>VP: $.35 \times .30 = .105$<br>S: $.05 \times .105 = .00525$ |          | VP: $.20 \times .30 \times .0045 = .00027$                                      |
|               |  | Det: .10 | NP: $.20 \times .10 \times .225 = .0045$  |
|               |  |          | Noun: .30<br>Nominal: $.75 \times .30 = .225$<br>NP: $.15 \times .225 = .03375$ |
| <i>back:</i>  | book   | that     | flight  |
|               | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP  |          | VP: 1 Verb NP   |
|               |  |          | NP: 2 Det Nominal   |
|               |  |          | Nominal: Noun<br>NP: Nominal  |

► VP → Verb NP [.20]

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|               |  |          |   |
|---------------|--|----------|---|
| <i>table:</i> | book   | that     | flight  |
|               | Noun: .10<br>Verb: .30<br>Nominal: $.75 \times .10 = .075$<br>NP: $.15 \times .075 = .01125$<br>VP: $.35 \times .30 = .105$<br>S: $.05 \times .105 = .00525$ |          | VP: $.20 \times .30 \times .0045 = .00027$<br>S: $.05 \times .00027 = .0000135$ |
|               |  | Det: .10 | NP: $.20 \times .10 \times .225 = .0045$  |
|               |  |          | Noun: .30<br>Nominal: $.75 \times .30 = .225$<br>NP: $.15 \times .225 = .03375$ |
| <i>back:</i>  | book   | that     | flight  |
|               | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP  |          | VP: 1 Verb NP<br>S: VP  |
|               |  |          | NP: 2 Det Nominal   |
|               |  |          | Nominal: Noun<br>NP: Nominal  |

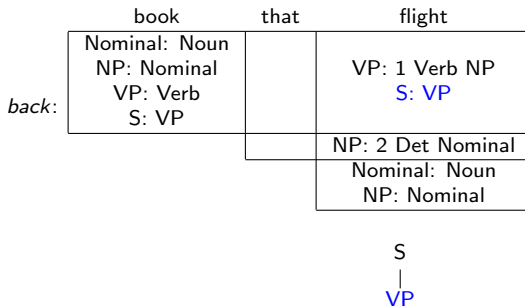
►  $S \rightarrow VP$  [.05]

# Probabilistic CKY Algorithm

|              | book  | that | flight                       |
|--------------|---|------|------------------------------|
| <i>back:</i> | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP |      | VP: 1 Verb NP<br>S: VP       |
|              |   |      | NP: 2 Det Nominal            |
|              |   |      | Nominal: Noun<br>NP: Nominal |

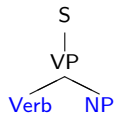


# Probabilistic CKY Algorithm



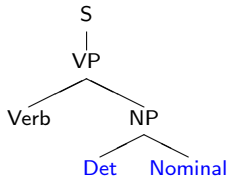
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|              |   |      |                              |
|--------------|---|------|------------------------------|
|              | book  | that | flight                       |
| <i>back:</i> | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP |      | VP: 1 Verb NP<br>S: VP       |
|              |   |      | NP: 2 Det Nominal            |
|              |   |      | Nominal: Noun<br>NP: Nominal |



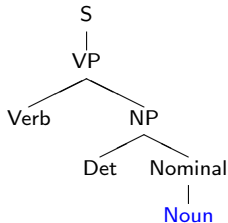
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|              | book  | that | flight                       |
|--------------|---|------|------------------------------|
| <i>back:</i> | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP |      | VP: 1 Verb NP<br>S: VP       |
|              |   |      | NP: 2 Det Nominal            |
|              |   |      | Nominal: Noun<br>NP: Nominal |



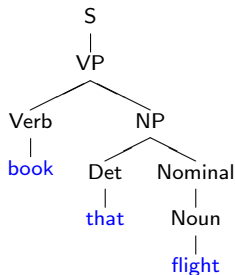
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|       |   |      |                              |
|-------|---|------|------------------------------|
|       | book  | that | flight                       |
| back: | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP |      | VP: 1 Verb NP<br>S: VP       |
|       |   |      | NP: 2 Det Nominal            |
|       |   |      | Nominal: Noun<br>NP: Nominal |



# Probabilistic CKY Algorithm

|       |   |      |                              |
|-------|---|------|------------------------------|
|       | book  | that | flight                       |
| back: | Nominal: Noun<br>NP: Nominal<br>VP: Verb<br>S: VP |      | VP: 1 Verb NP<br>S: VP       |
|       |   |      | NP: 2 Det Nominal            |
|       |   |      | Nominal: Noun<br>NP: Nominal |



# Earley Algorithm

- ▶ Fill in a chart with dotted rules

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  - ▶ Predictor
    - ▶ If the nonterminal after the dot is not a preterminal (POS tag), add a state for each possible expansion of the nonterminal

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- ▶ Fill in a chart with dotted rules
  - ▶ Predictor
    - ▶ If the nonterminal after the dot is not a preterminal (POS tag), add a state for each possible expansion of the nonterminal
  - ▶ Scanner
    - ▶ If the nonterminal after the dot is a preterminal (POS tag), read the next word in the input and, if the word can be tagged as that tag, add a state to the next position in the chart



# Earley Algorithm

- ▶ Fill in a chart with dotted rules
  - ▶ Predictor
    - ▶ If the nonterminal after the dot is not a preterminal (POS tag), add a state for each possible expansion of the nonterminal
  - ▶ Scanner
    - ▶ If the nonterminal after the dot is a preterminal (POS tag), read the next word in the input and, if the word can be tagged as that tag, add a state to the next position in the chart
  - ▶ Completer
    - ▶ If the dot is at the end of the rule, add states corresponding to previous states that were “looking” for the rule to be completed, and advancing the dot

# Pushdown Automata

- ▶ Finite state automaton with stack

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- ▶ Finite state automaton with stack
  - ▶ Infinite memory that can only be used last in, first out

# Pushdown Automata

- ▶ Finite state automaton with stack
  - ▶ Infinite memory that can only be used last in, first out
- ▶ Equivalent to CFGs

# Pushdown Automata

- ▶ Finite state automaton with stack
  - ▶ Infinite memory that can only be used last in, first out
- ▶ Equivalent to CFGs
- ▶ For more information, see theory of computation book

# Natural languages are not context-free

- ▶ Why not?