Beyond basics in parallelism

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Loop parallelism

As a reminder, most parallelism in loops

```
#pragma omp target teams distribute parallel for simd
for (int i=0; i<N; i++) {
    A[i] += 3*B[i]
}</pre>
```

Reductions

- When you compute in parallel, you should never write to the same variable/memory location from multiple iterations
 - That would lead to unpredictable results
- Common operations across iterations supported as reductions
 - The reduction variable is kept private during compute
 - All intermediate values then combined automatically at the end of the loop

Partitioned memory

- Reminder: CPU and GPU memories are typically separate
 - Buffers must be moved between them as needed
- OpenMP can try to infer what needs to be done
 - But doing it explicitly best

```
double* A = new double[N];
  double* B = new double[N];
  ...

#pragma omp target teams distribute parallel for simd \
    map(tofrom: A[0:N]) map(to: B[0:N])
  for (int i=0; i<N; i++) {
        A[i] += 3*B[i];
    }</pre>
```

Classes/Objects and partitioned memory

Works well for both numeric buffers and buffers of simple objects

```
double* A = new double[N];
   MyClass* B = new MyClass[N];
   ...

#pragma omp target teams distribute parallel for simd \
      map(tofrom: A[0:N]) map(to: B[0:N])
   for (int i=0; i<N; i++) {
        A[i] += B[i].compute();
   }</pre>
```

Classes/Objects and partitioned memory

- But will not automatically handle nested buffers!
 - Most compilers also will not automatically detect the need for copying internal buffers
 - Resulting in code crashing
- Using pointers is OK, but the pointed buffer must be handled independently

```
class MyClass {
   double *inarr;
 public:
   MyClass(): inarr(new double[34]) {}
   double compute() {return inarr[2] + inarr[7];}
 };
 double* A = new double[N];
 MyClass* B = new MyClass[N];
#pragma omp target teams distribute parallel for simd \
   map(tofrom: A[0:N]) map(to: B[0:N])
 for (int i=0; i<N; i++) {
   A[i] += B[i].compute();
```

Classes/Objects and managed/unified memory

- Nested buffers allowed when memory movement is managed by GPU driver
 - Makes life much easier!
 - But typically slower
 - And requires explicit management of static buffers
- Unified memory between
 CPU and GPU a step further
 - But only few platforms support it right now!
 E.g. NVIDIA Grace Hopper and AMD MI300A APU

```
#pragma omp requires unified_shared_memory
 class MyClass {
   double *inarr;
 public:
    MyClass(): inarr(new double[34]) {}
    double compute() {return inarr[2] + inarr[7];}
 double* A = new double[N];
 MyClass* B = new MyClass[N];
#pragma omp target teams distribute parallel for simd
 for (int i=0; i<N; i++) {
   A[i] += B[i].compute();
```

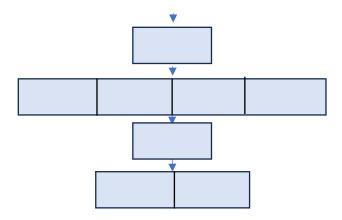
No dynamic memory in GPU compute

- Dynamic memory management in GPU code very limited (e.g. new and/or malloc)
 - Slow and typically limited to a KB of memory
 - Normally does not interoperate with CPU-based dynamic memory
- Makes using standard C++ containers virtually impossible (e.g., std::vector)

```
std::vector *A = new std::vector[N];
...
#pragma omp target teams distribute parallel for simd
for (int i=0; i<N; i++) {
    A[i].push_back(B[i].compute());
}</pre>
```

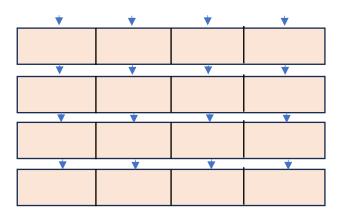
CPU

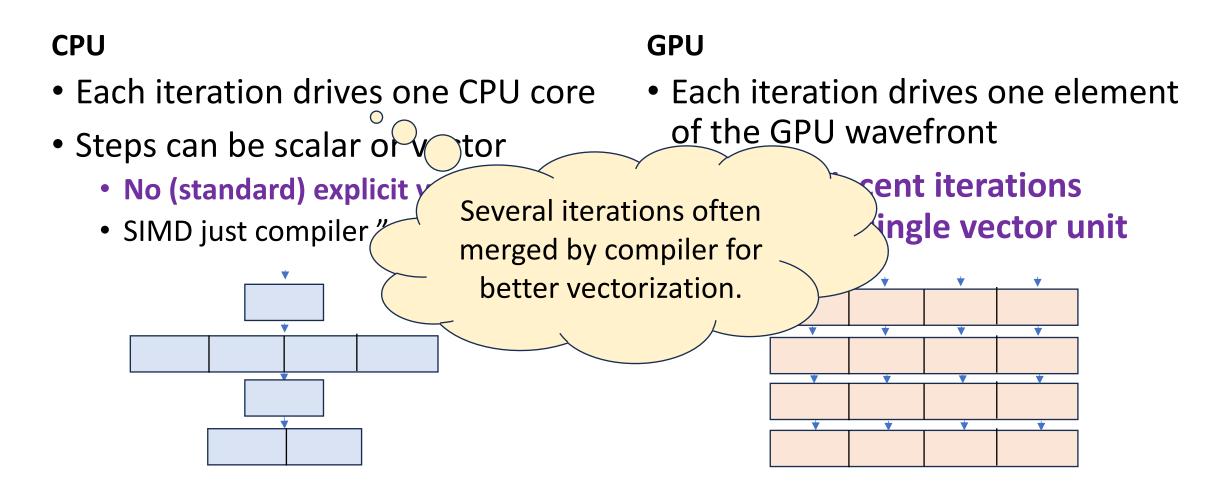
- Each iteration drives one CPU core
- Steps can be scalar or vector
 - No (standard) explicit vector syntax
 - SIMD just compiler "magic"



GPU

- Each iteration drives one element of the GPU wavefront
- Several adjacent iterations executed as single vector unit





CPU

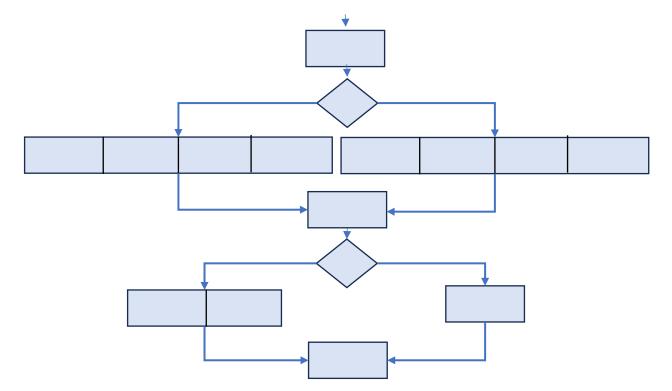
- Each iteration drives one core
- Steps can be scalar or vector
 - No (standard) explicit vector syntax
 - SIMD just compiler "magic"
- Each iteration independent
 - No coordination between iterations (unless compiler merges them)
- Logic (if) just another instruction
 - Typically predicted and/or speculated

GPU

- Each iteration drives one element in the wavefront vector
 - All steps are always scalar
 - HW groups them in vectors (SIMT)
- All instructions in a vector progress in lockstep
 - But no coordination between vectors (wavefronts)
- Logic (if) may split vector
 - Reducing parallelism!
 - And likely to yield until decision

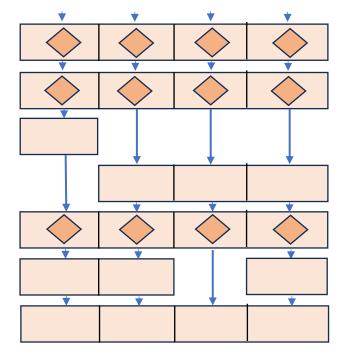
CPU

- Logic (if) just another instruction
 - Typically predicted and/or speculated



GPU

- Logic (if) may split vector (at runtime)
 - Reducing parallelism!
 - And likely to yield until decision



Implications on high-level parallelism

- Simple loops are a no-brainer
 - The compiler will always do the right thing
 - Both for CPUs and GPUs

```
#pragma omp parallel for for (int i=0; i<N; i++) {
    A(i) += 3*B(i)
}
```

All operations are vectorized

```
#pragma omp target teams distribute \
    parallel for simd
for (int i=0; i<N; i++) {
    A(i) += 3*B(i)
}</pre>
```

No logic -> no wavefront splitting

Implications on high-level parallelism

- Complex loops can be tricky
 - On CPUs, logic and nesting may make it hard to vectorize

```
#pragma omp parallel for
for (int i=0; i<N; i++) {
   double x = B[i] * C[i];
   double y = 0;
   for (int j=0; j<D[i/32]; j++) {
      if (E[i/4+j]>x) y+=A[i]*E[i/4+j];
      }
      A[i] = y
   }
```

GPU equivalent would have no problems "vectorizing" at all.

Implications on high-level parallelism

- Complex loops can be tricky
 - On GPUs, logic can easily fragment the vectors

```
#pragma omp target teams distribute \
      parallel for simd
 for (int i=0; i<N; i++) {
   // some complex code
   if (B[i]>1.0) {
     // some complex code
   } else {
     // some complex code
```

Effectively doubling the cost on GPUs (in many cases)

And will get worse if there is more logic inside the branches

CPUs would have no problems at all. Each branch would be vectorized separately. (wherever possible)

GPU compute and external functions

- Functions that can be inlined are fine
 - Compiler will treat them like if you cut-and-paste the code
- Separate compilation needs additional pragmas
 - I.e., must explicitly declare as GPU-enabled

```
double compute(double a, double b) {
   return a + 3*b;
#pragma omp declare target to(compute)
 double* A = new double[N];
 double* B = new double[N];
#pragma omp target teams distribute parallel for simd
 for (int i=0; i<N; i++) {
   A[i] = compute(A[i], B[i]);
```

GPU compute and vectorized functions

- Some functions are complex enough to warrant in-function vectorization
- Must use the alternative omp target teams loop syntax
 - Gives compiler more freedom
- Only supported for inlined functions
 - No good way to tell that an external function is vectorized

```
double compute(double a, double b) {
   double s = 0.0;
#pragma omp loop bind(parallel) reduction(+:s)
   for (int j=0; j<N; j++) {
     s += C[j]*a + D[j]*b;
   return s;
 double* A = new double[N];
 double* B = new double[N];
#pragma omp target teams loop
 for (int i=0; i<N; i++) {
   A[i] = compute(A[i], B[i]);
```

Asynchronous execution

- Reminder: CPU and GPU compute proceed independently
- However, OpenMP is by default blocking
 - The compiler will insert an implicit wait/barrier
 - Makes writing code much easier

```
#pragma omp target teams distribute parallel for simd
for (int i=0; i<N; i++) {
    A[i] += 3*B[i]
}
// implicit GPU wait
... // more CPU code here</pre>
```

Asynchronous execution

- Reminder: CPU and GPU compute proceed independently
- OpenMP does support async execution
 - But must be explicitly requested
 - Based on the notion of data dependencies

```
#pragma omp target teams distribute parallel for simd map(tofrom: A) nowait
for (int i=0; i<N; i++) {
    A[i] += 3*B[i]
}
... // more independent CPU code here
#pragma omp ... depend(A)
... // GPU finished, can safely use A here</pre>
```

Ready for some more hands-on?