

# **ADS Project-4**

## **The Best Peak Shape**

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## I. Introduction

### 1.1 Problem Description

The Best Peak is a very classic max-increasing-substring which can be easily solved by DP method.

After input a series of string, we find the longest sub-sequence that forms a peak shape. If there is a tie (a tie refers to the same length of sub-sequence), the most symmetric one will be chosen. The peak shape means the whole string can be divided into three parts --- the rising edge, peak and falling edge. In mathematical formula, there exists an index  $i$  ( $1 < i < L$ ) such that  $D_1 < \dots < D_{i-1} < D_i > D_{i+1} > \dots > D_L$ .

The Best Peak Shape problem is very common in science, such as physics, data analysis, image processing and so on. Today, we will introduce a quite simple method towards this problem.

## II. Algorithm Specification

The whole program is rather simple to be designed. First, we insert the size of input and a series of numbers. We can interpret the whole string as the string  $[a, i]$  and  $string[i, b]$ . Where the  $string[a, i]$  is the left increasing substring of the whole peak shape. And  $string[i, b]$  is the decreasing substring of the input. However, we can change its sequence and see it as the  $string[b, i]$ , then it becomes the increase substring as well. Thus, the whole problem can be simplified into solving two longest increasing substring problems.

During longest-increasing-substring problems, a state transition equation can easily be obtained that  $length[i] = \max(length[i], length[i - 1] + 1)$ . And we can just use two loops to traverse through all the substring of the whole string to get the correct answer which can be vividly seen in the pseudo-code.

### 2.1 pseudo-code

The pseudo-code of solving `max_increasing_substring` function is as follows.

```

1.  int main()
2.  {
3.      input L and string;
4.      for(i from 0 to L) // Loop the end of substring
5.      {
6.          Initialize length_longest_substring[i]
7.          for(j from 0 to i) // Loop the start of substring
8.          {
9.              if(string(i) > string (j))
10.                 // state transition equation
11.                 length = MAX(length,length_longest_substring[j] + 1)
12.
13.          }
14.      }
15.      return length;
16. }

```

## 2.2 Postscripts

There may be some differences between the pseudo-code given above and the actual program. But the whole algorithm is more or less the same.

## III. Testing Results

### 3.1.1 comprehensive cases

```

20
1 3 0 8 5 -2 29 20 20 4 10 4 7 25 18 6 17 16 2 -1
Results: 10 14 25

```

```

20
527 -620 -9455 -1272 -2674 -8815 -5328 -9025 -5597 5177 -6431 -3044 1533
5908 -6572 -4645 3864 -351 -2693 3928 7118 -9671 -9221 -5904 52 8374 -9
853 1261 1638 -6103 -932 8156 -4477 -1517 5219 -6783 -330 1432 -7788 12
61 2229 1325 -2043 -6003 -2482 498 1916 -6169 -3010 1125 2237 -8318 989
3 2008 -6276 6309 -1259 3704 -36 -6912 1746 1305 2488 -434 9432 -2833 1
566 -2403 374 -4098 2141 8136 1154 -3953 -8994 -861 6512 290 -1706 1570
-1990 -884 -3396 -718 -3675 6327 6201 -6905 -8569 -8915 -3427 4636 2402
4540 -8454 -3596 -4235 -7983 -1973 -6332 7763 -55 436 -6441 7955 1674 -
1615 -7064 2280 -9526 -4824 3380 2387 -6278 -7871 1459 1049 9279 5655 -
9692 464 -4484 4238 -1463 -7170 -1414 9479 -7949 -6364 8352 -7630 7392
1077 5730 7715 -1663 5007 -2430 -9815 -8832 9506 1362 -4645 2915 5651

```

```
1017 -5678 -4597 -6554 1978 -5825 5408 -4520 -7569 1013 -112 -8529 -3684
-6904 7669 -6709 226 -2150 7190 -3992 -9456 7537 -7247 -5251 -927 -5877
-4521 386 8385 -6507 -8392 1694 273 -8581 770 -6187 -6579 -1510 8101 -
8280 4302 7678 -2056 -1253 23 -18 789 -7589 -1345 4587 432 1 3 15 96
Results: 34 65 9432
```

### 3.1.2 no heap shape

```
5
-1 3 8 10 20
Results: No peak shape
```

```
15
1211 189 75 23 20 11 4 -15 -40 -45 -46 -77 -89 -91 -152
Results: No peak shape
```

### 3.1.3 Minimum scale

```
3
-1 3 1
Results: 3 2 3
```

```
3
-1 3 4
Results: No peak shape
```

### 3.1.4 Maximum scale

```
10000
The test data is too long, so they are not shown here. Please refer to "case3.t
xt"
Results: 275 4185 9991
```

After verification, we can know that all the test results are correct.

## IV. Analysis and Comments

As the test results have shown, our program has accomplished the goal to find the best peak shape. Using DP, our program divides the problem of best peak shape into

the sub problems of the best peak shape of the sub sequences. Processing with the nodes one by one, before the node, the best peak and its uphill and downhill is preserved for comparison.

In our DP function, there are two n-times iterations is done to find and construct the peak's uphill and downhill. The time complexity is  $O(N)$ . In the main function, the data is input first, so the time complexity is  $O(N)$ . Because each number must run the DP function and the scale of data is  $N$ , the time complexity is  $O(N^2)$ . At last, compare each peak to find out which of them has best peak shape. This is done through an iteration by comparing each pair of the peak. As a result, the time complexity is  $O(N)$ . To conclude, the overall time complexity of the program is  $O(N^2)$ .

As for the space complexity, we must allocate some space for several fixed-sized arrays to store the data, such as the array named "seq" to store the data and "up" to store information about the uphill, to name just a few. The size of the space is related to the scale of the data. As a result, the space complexity of our program is  $O(N)$ .

## V. Appendix

```
1. #include<iostream>
2. #include<cstdlib>
3.
4. using namespace std;
5.
6. int seq[11000]; //record sequence (sequence)
7. int up[11000]; //record the longest sequence of increments with i as the last
   node
8. int down[11000]; //records with i as the first node to decrement the longest
   sequence
9.
10.
11. //Dynamic programming algorithm
12. void DP(int index, int N)
13. {
14.     int temp;
15.     temp = 0;
16.
17.     //Calculating incremental sequences
```

```

18.     for (int i = 0; i < index; i++)
19.     {
20.         if (seq[index] > seq[i] && temp < up[i] + 1)//if the sequence[index]
            > sequence[i] and temp < up[i]+1, means that this incremental sequence is l
            onger
21.         {
22.             temp = up[i] + 1;
23.             up[index] = temp;                //let down[index] be the
            longer incremental sequence
24.         }
25.     }
26.
27.     temp = 0;
28.
29.     //Calculating descending sequence(from the last number to the first numb
        er, contrary to the incremental sequence)
30.     for (int i = N - 1; i > N - 1 - index; i--)
31.     {
32.         if (seq[N - 1 - index] > seq[i] && temp < down[i] + 1)//if the seque
            nce[index] > sequence[i] and temp < down[i]+1, means that this incremental s
            equence is longer
33.         {
34.             temp = down[i] + 1;
35.             down[N - 1 - index] = temp;        //let down[N-1-
            index] be the longer incremental sequence
36.         }
37.     }
38.
39.     return;
40. }
41.
42.
43. int main()
44. {
45.
46.     //N is the number of the sequence,
        max records the maximum value subscript and determine if there is a peak
47.     int N, max = 0, maxindex, difference = 1000000;
48.     printf("please input the number of the sequence(3-10000):\n");
49.     //initialize the data
50.     cin >> N;
51.     printf("please input %d numbers(split with spaces or carriage returns):\n", N);
52.     for (int i = 0; i < N; i++)

```

```

53.     {
54.         scanf("%d", &seq[i]);
55.         up[i] = 1;
56.         down[i] = 1;
57.     }
58.
59.     //excute the dynamic programming algorithm
60.     for (int i = 0; i < N; i++)
61.         DP(i, N);
62.
63.
64.     //Find the longest "peak" and compare the symmetry of the longest "peak"
65.     for (int i = 0; i < N; i++)
66.     {
67.         if (up[i] > 1 && down[i] > 1)    //First, find the longest "peak"
68.         {
69.             if (max < up[i] + down[i] - 1 )
70.             {
71.                 difference = abs(up[i] - down[i]);
72.                 max = up[i] + down[i] - 1;
73.                 maxindex = i;
74.             }
75.             //finally, if it is equal, find the most symmetrical "peak"(the
              absolute value of (up[i] - down[i]) is the smallest)
76.             else if (max == up[i] + down[i] - 1 && difference > abs(up[i] -
              down[i]))
77.             {
78.                 difference = abs(up[i] - down[i]);
79.                 max = up[i] + down[i] - 1;
80.                 maxindex = i;
81.             }
82.         }
83.     }
84.
85.
86.     //Output the corresponding result if it exists
87.     if (max)
88.         printf("%d %d %d\n", max, maxindex + 1, seq[maxindex]);
89.     else printf("No peak shape\n");
90.
91.     system("pause");
92.     return 0;
93. }

```