

Introduction to DNS and its vulnerabilities

Olaf M. Kolkman
olaf@nlnetlabs.nl

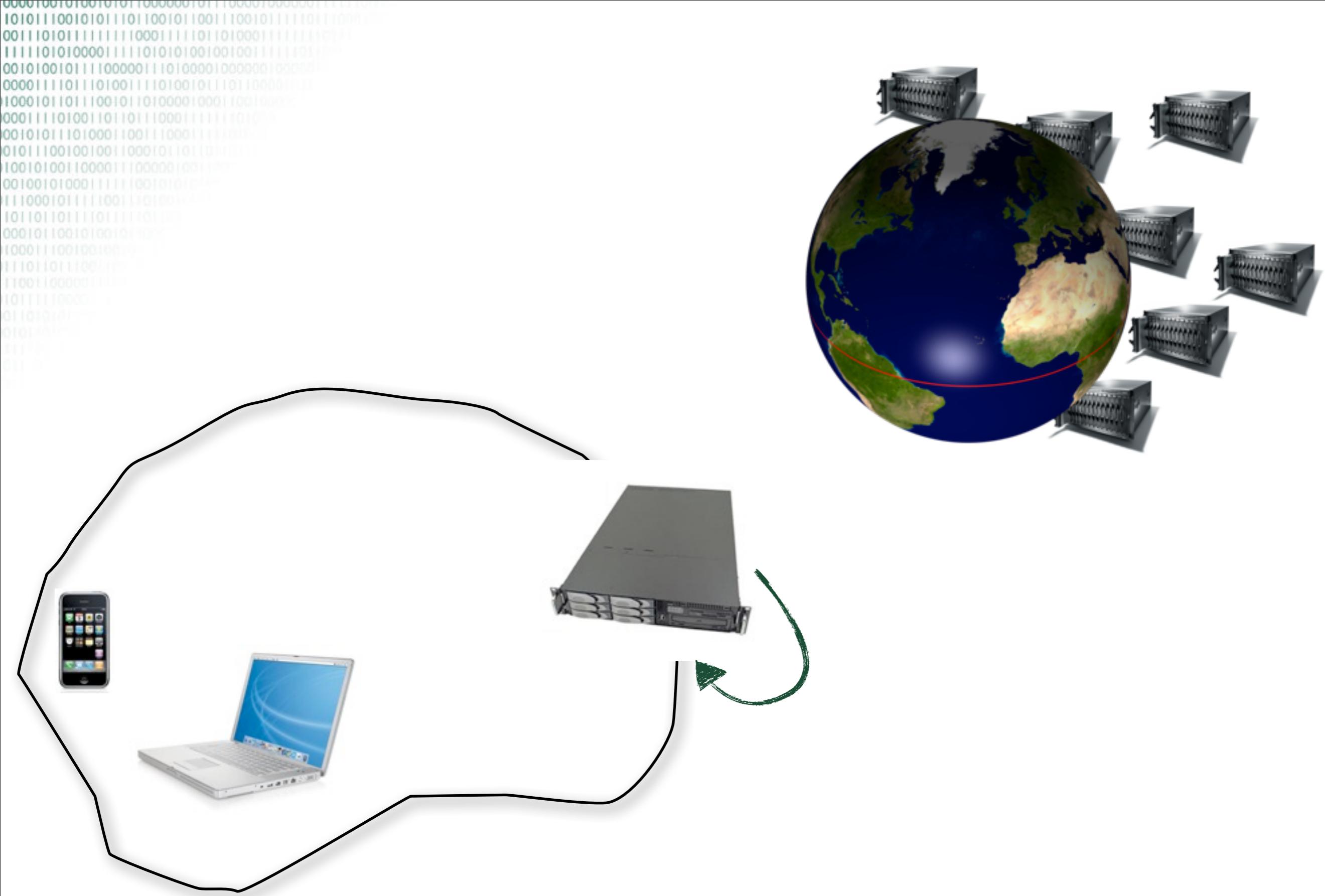


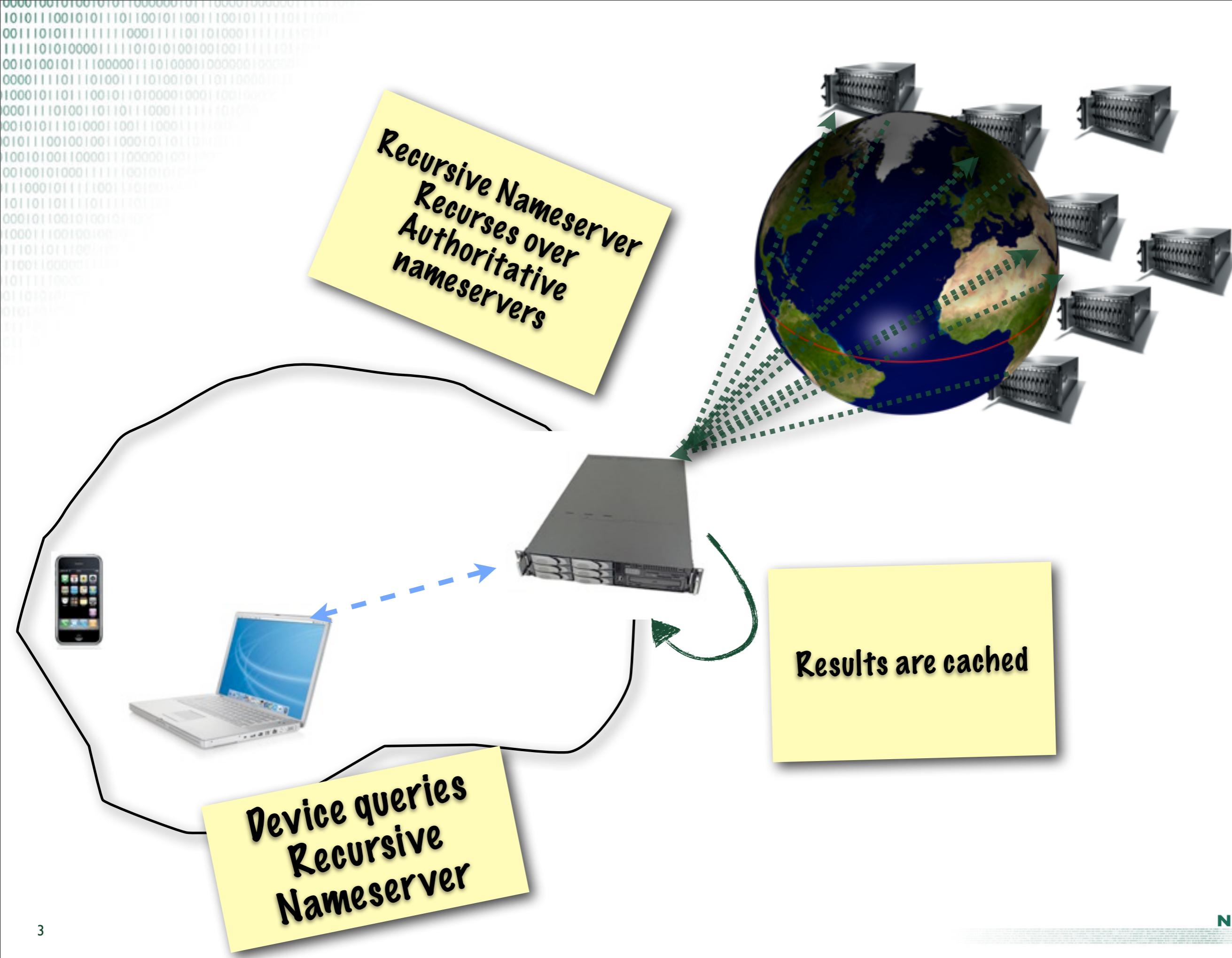
© 2006-2012 NLnet Labs, Licensed under a [Creative Commons Attribution 3.0 Unported License](#).



DNS and DNSSEC
in a Nutshell

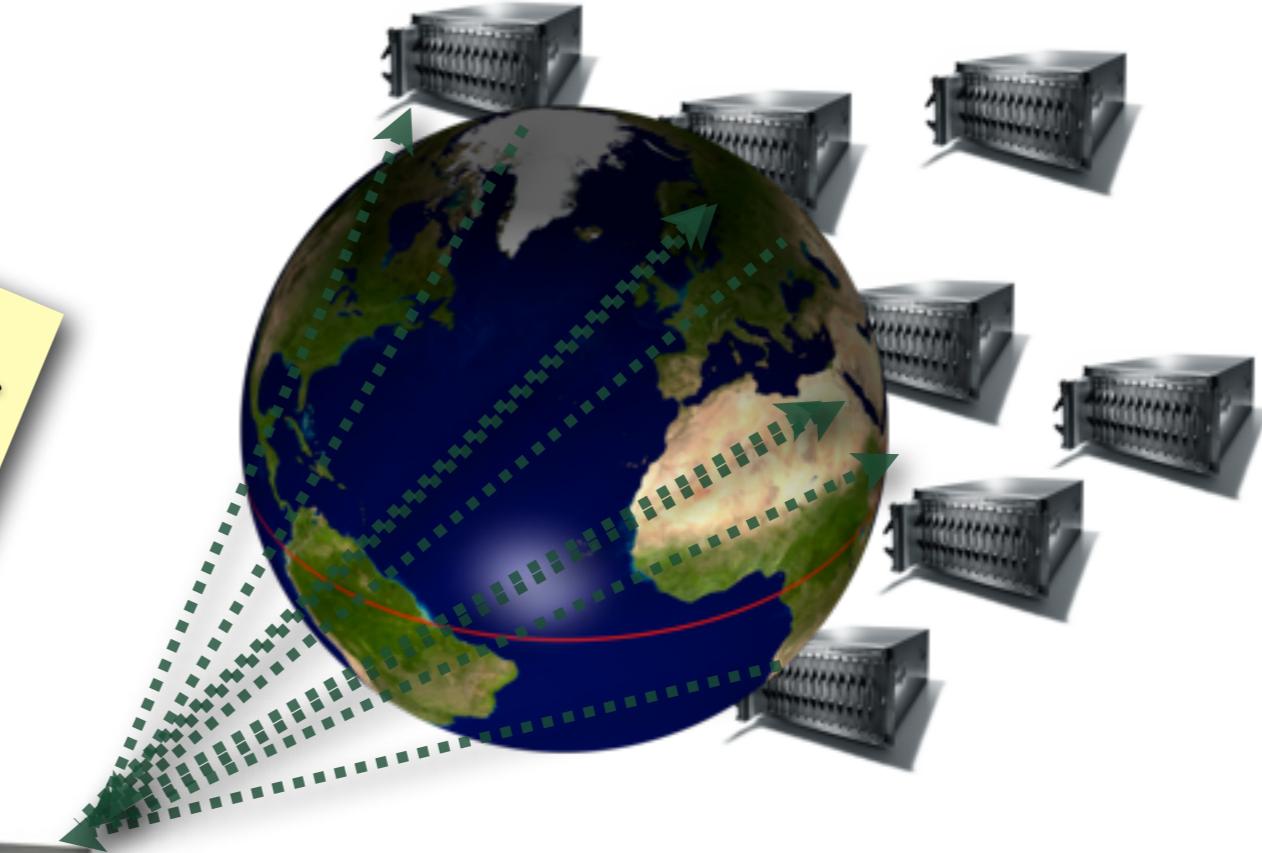






The DNS is highly distributive

Recursive Nameserver
queries over
Distributive
Nameservers

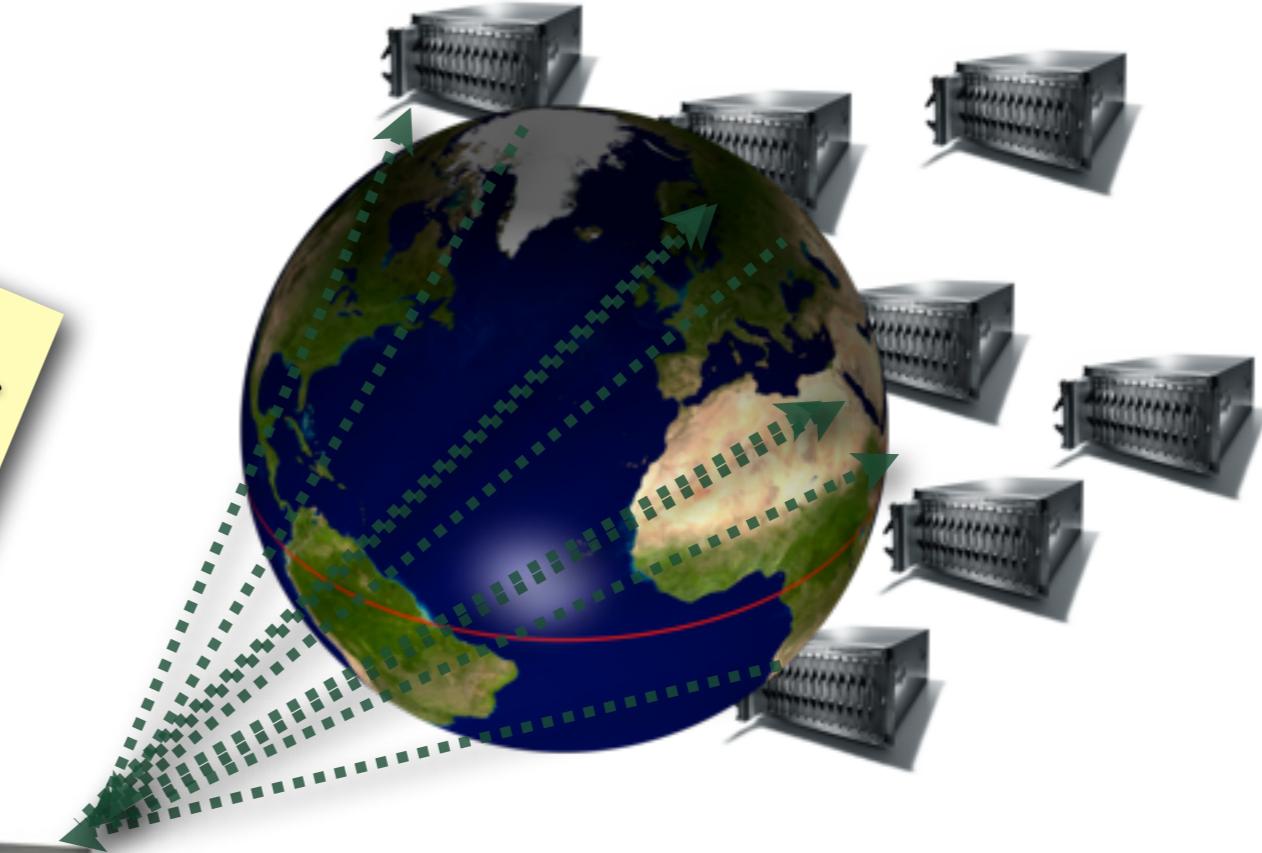


Device queries
Recursive
Nameserver

Results are cached

The DNS is highly distributive

Recursive Nameserver
queries over
Distributive
Nameservers



Device queries
Recursive
Nameserver

DNS is implemented
through 100s of
thousands of
machines

Authoritative Nameservers **ROOT**



Stub Resolver



Recursive Nameserver

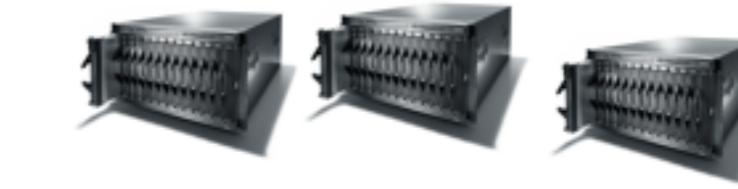


NL



NLnetLabs.NL

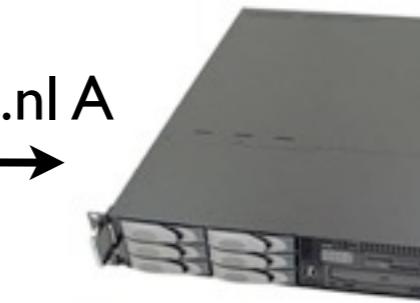
Authoritative Nameservers **ROOT**



Stub Resolver



www.nlnetlabs.nl A

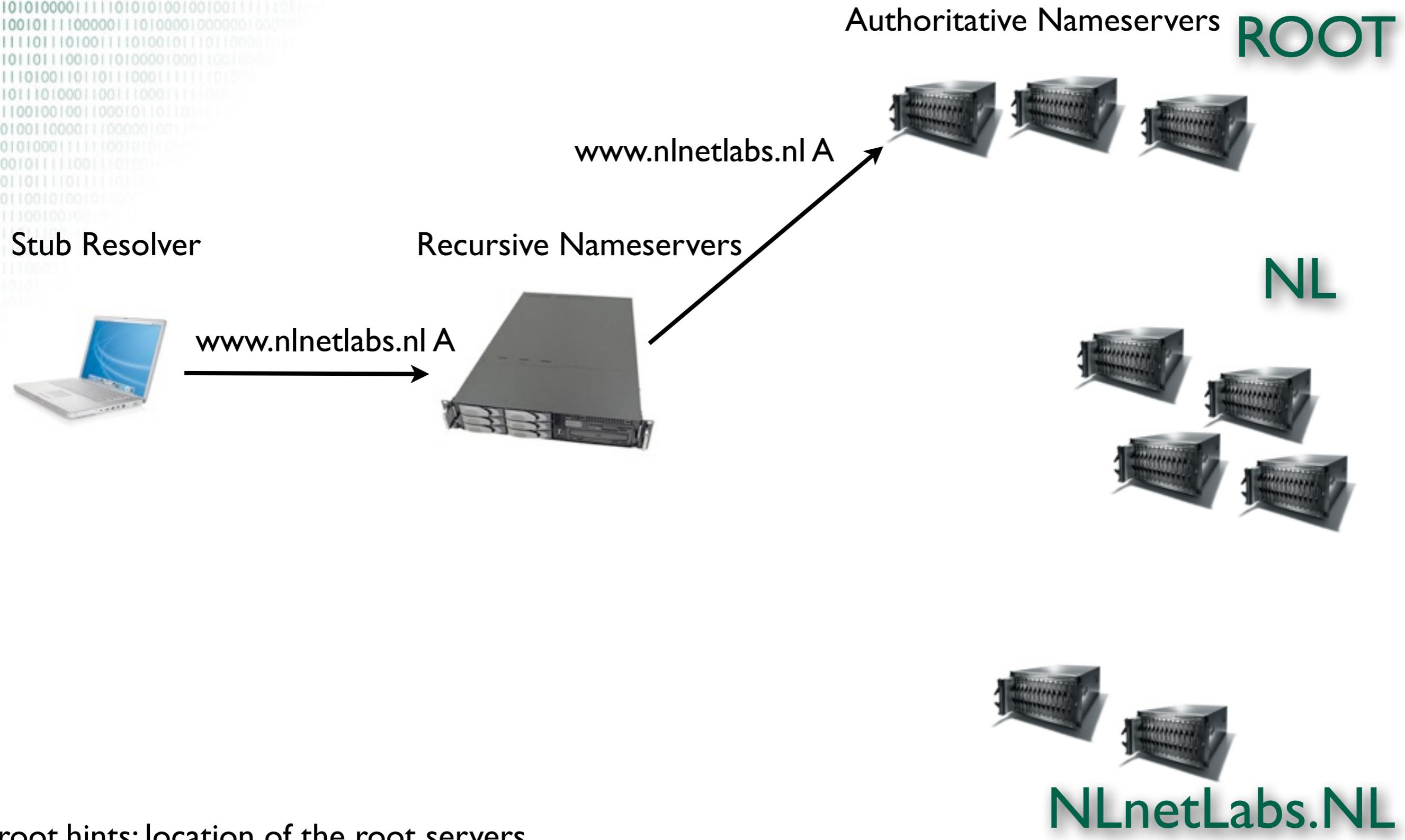


Recursive Nameservers

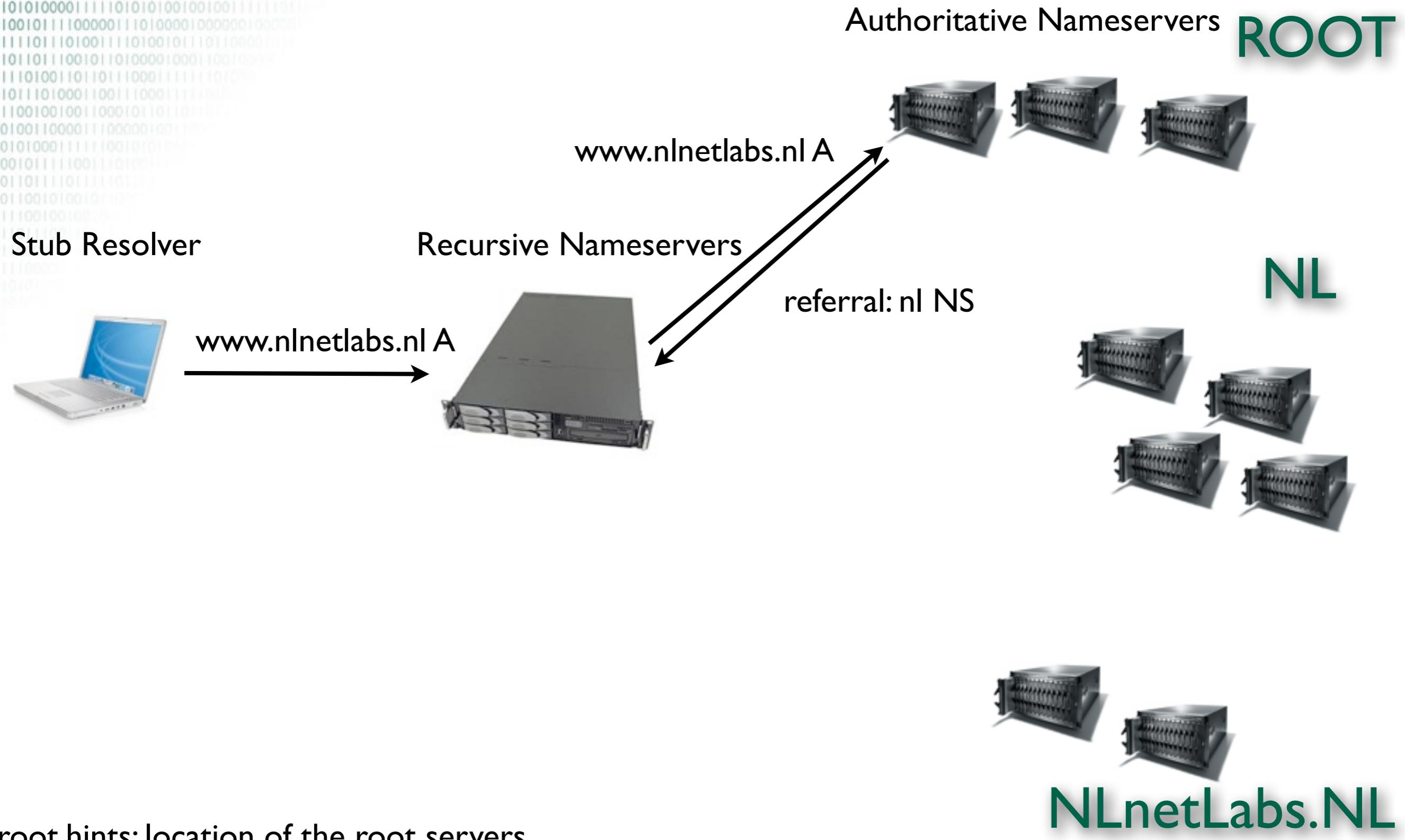
NL



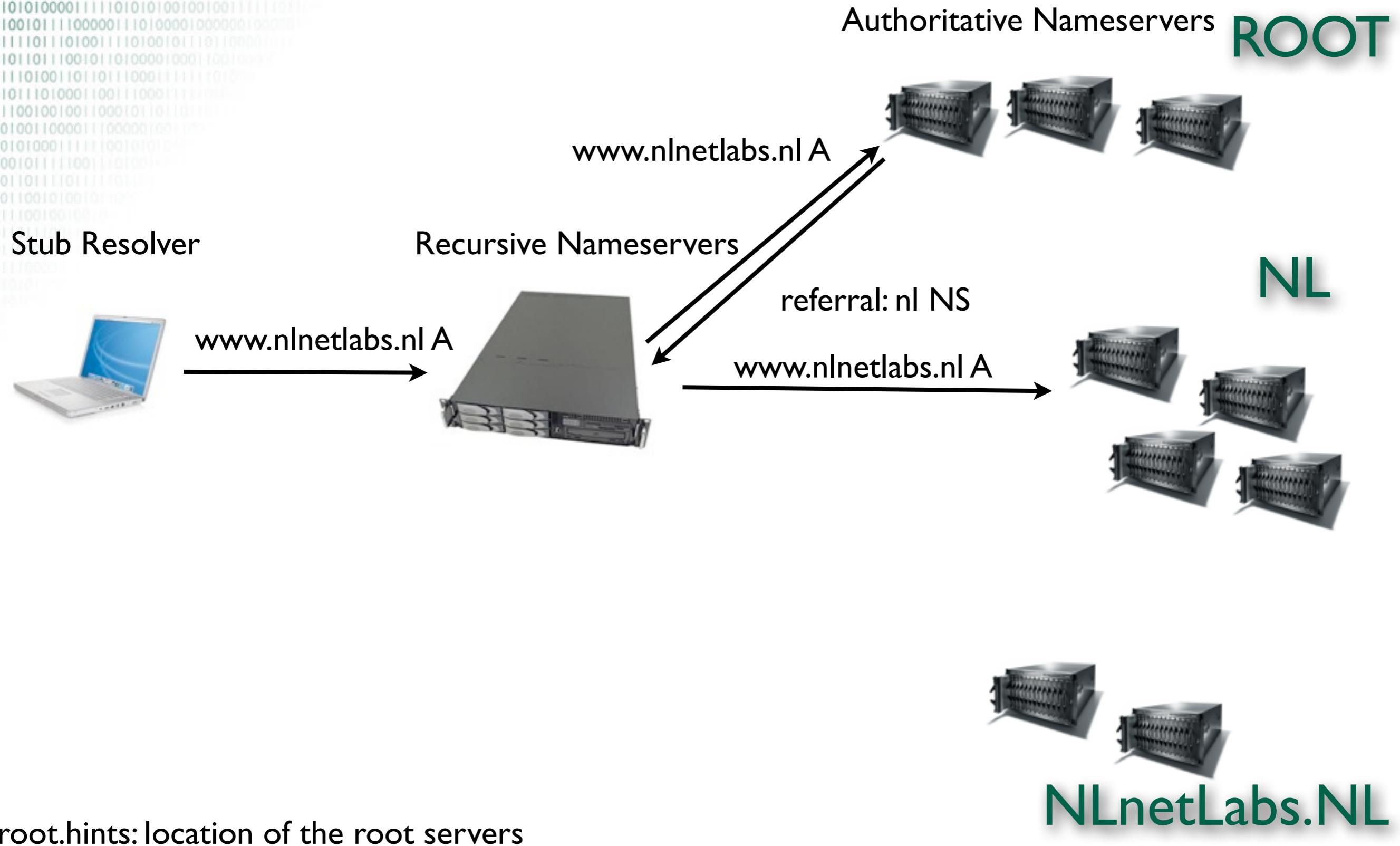
NLnetLabs.NL



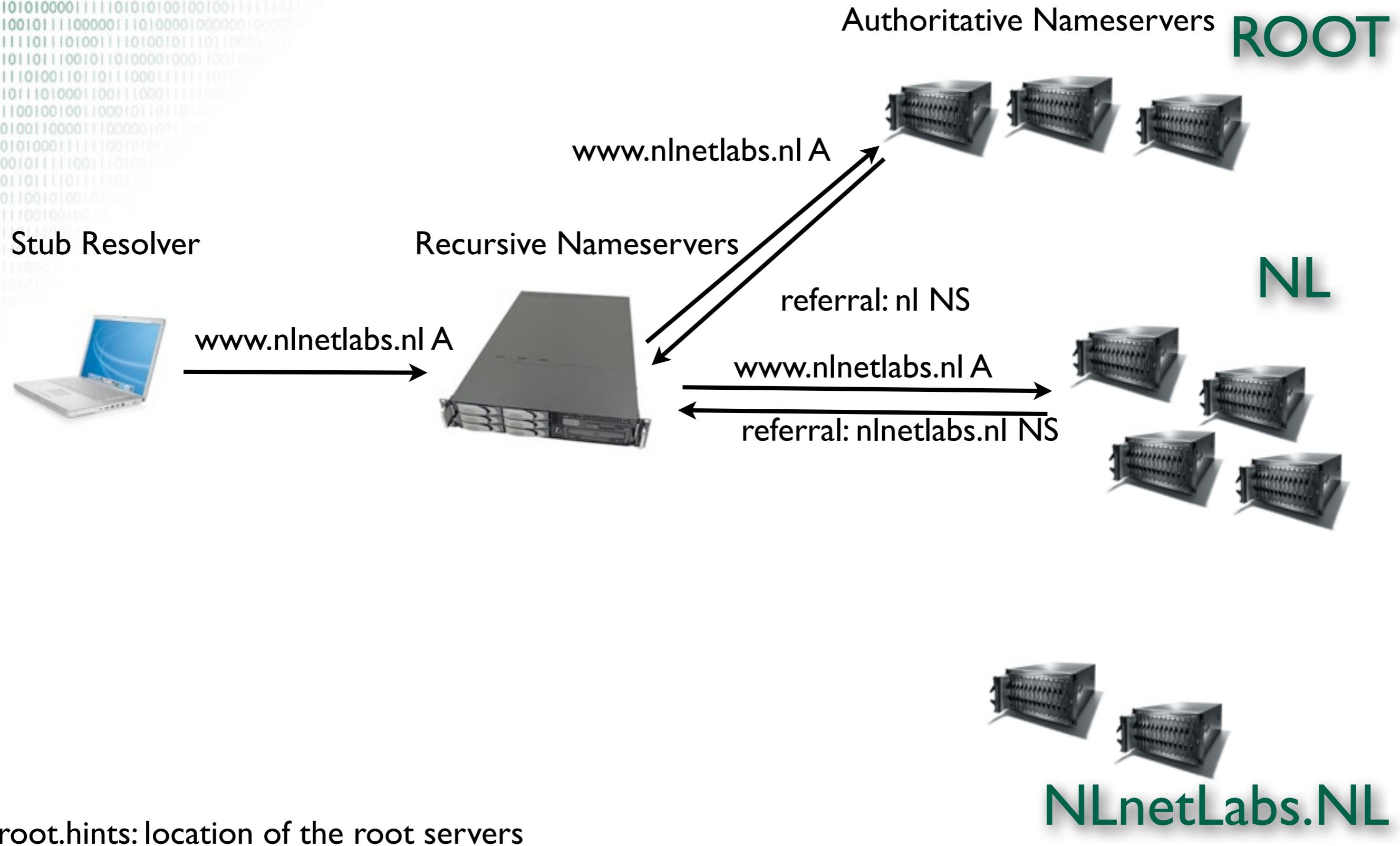
root.hints: location of the root server



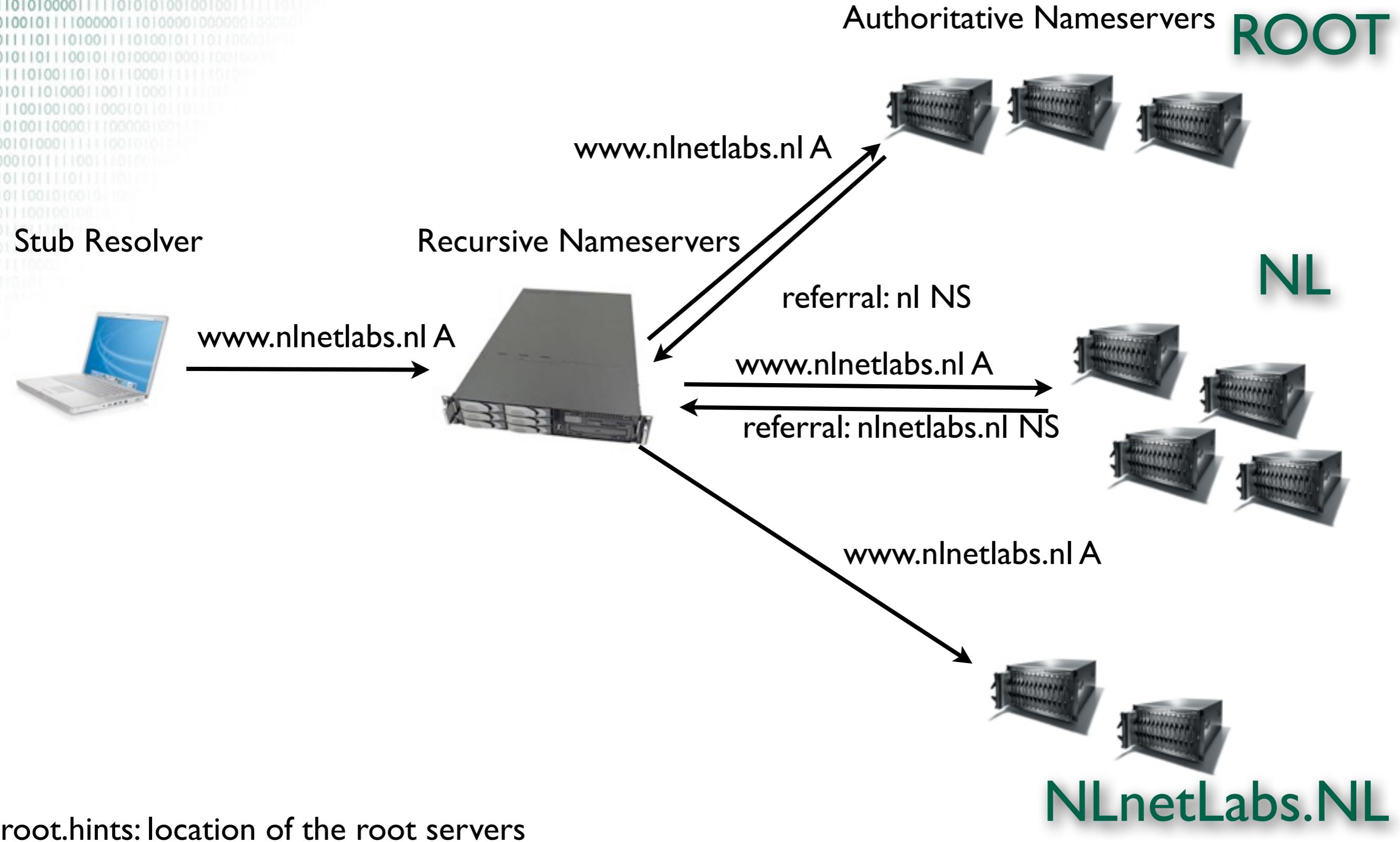
root.hints: location of the root server



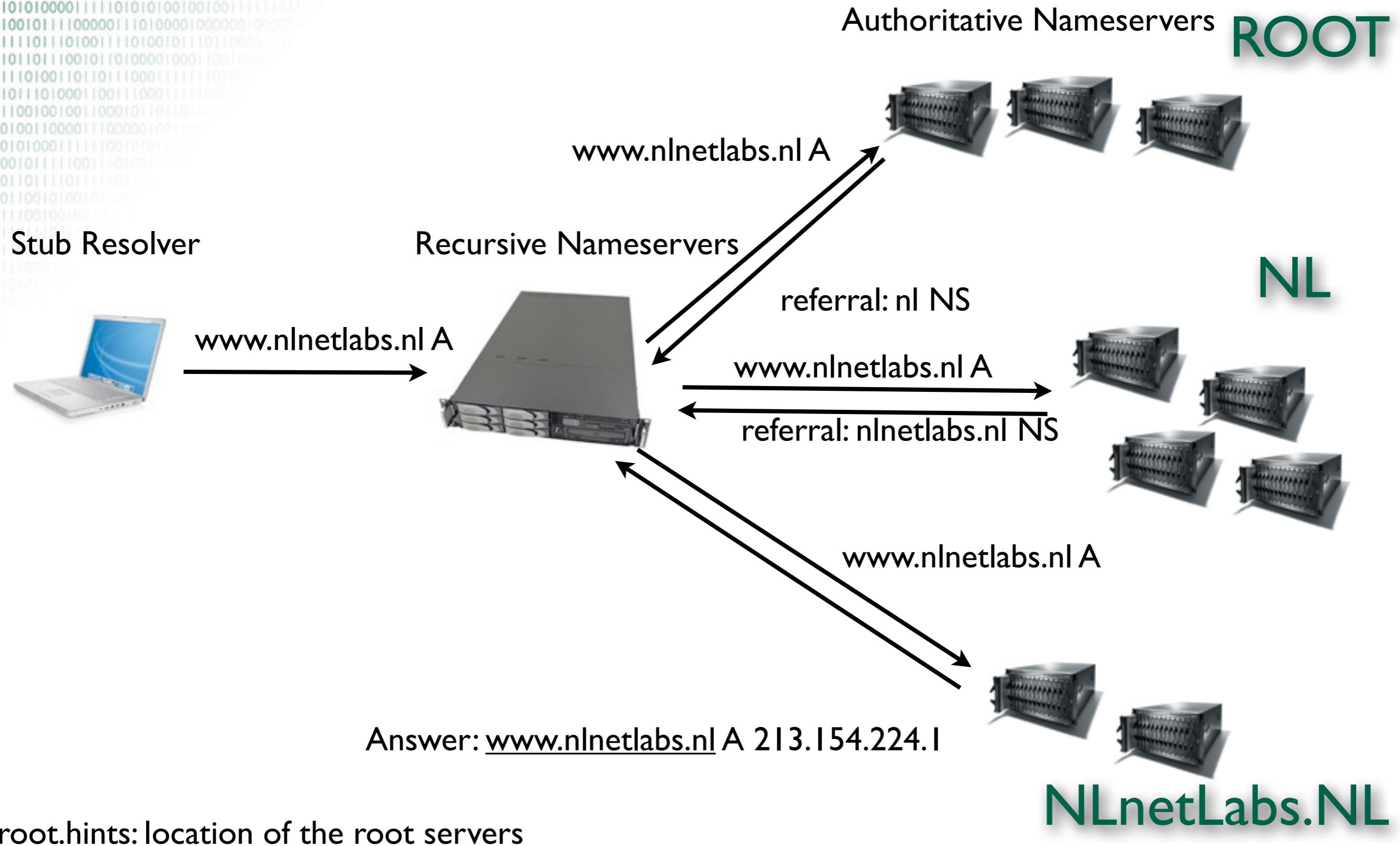
root.hints: location of the root server



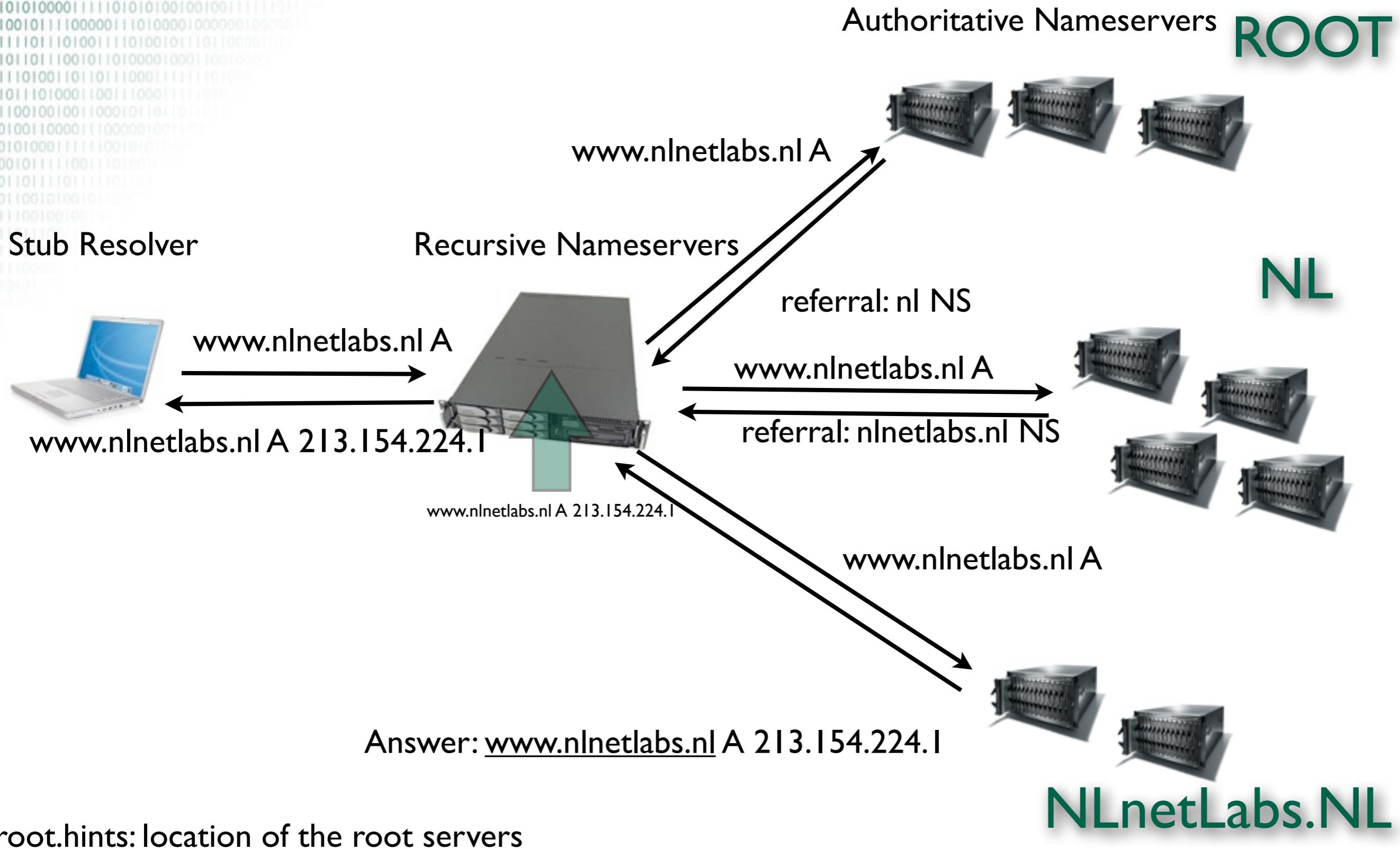
root.hints: location of the root server



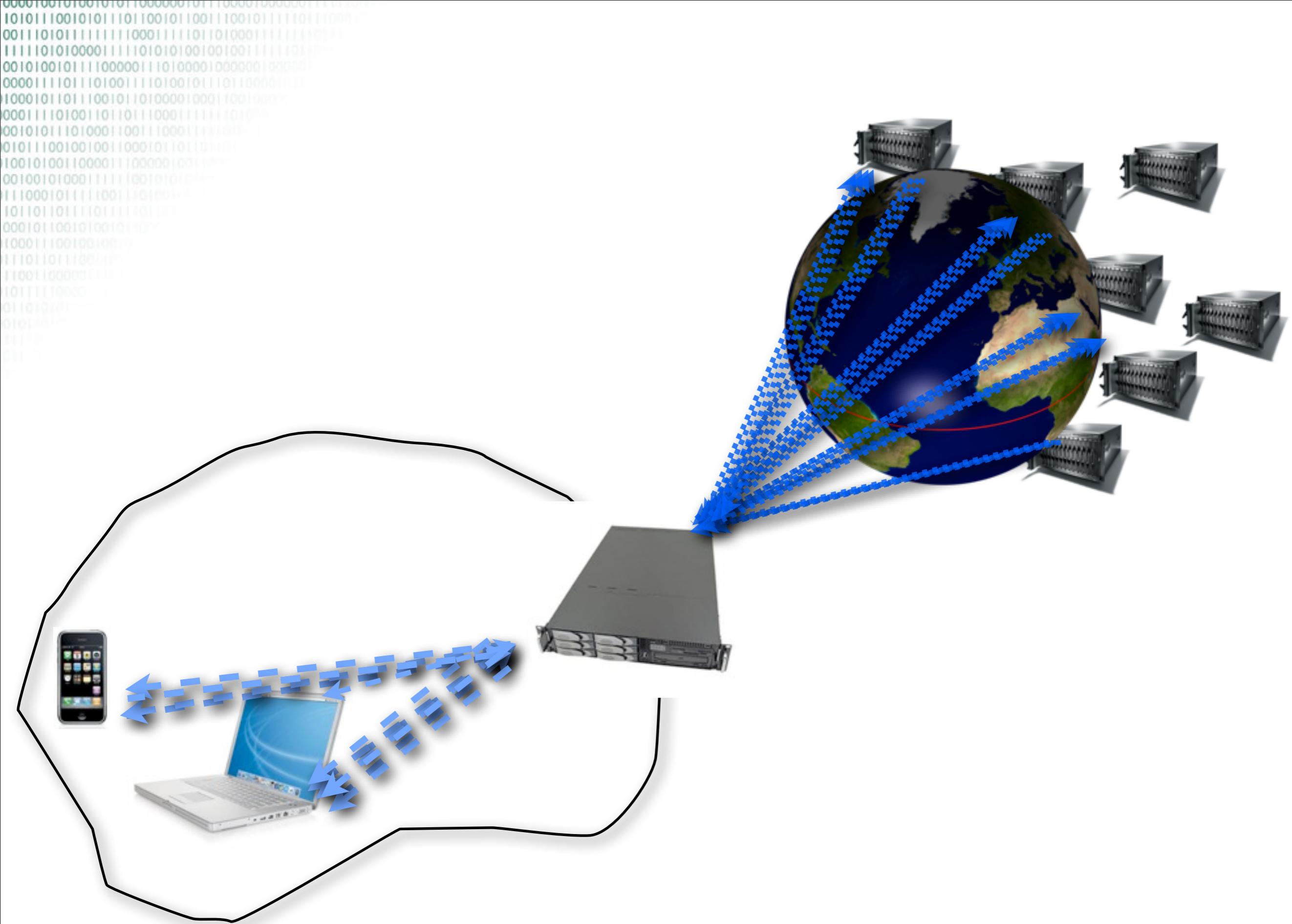
root.hints: location of the root server



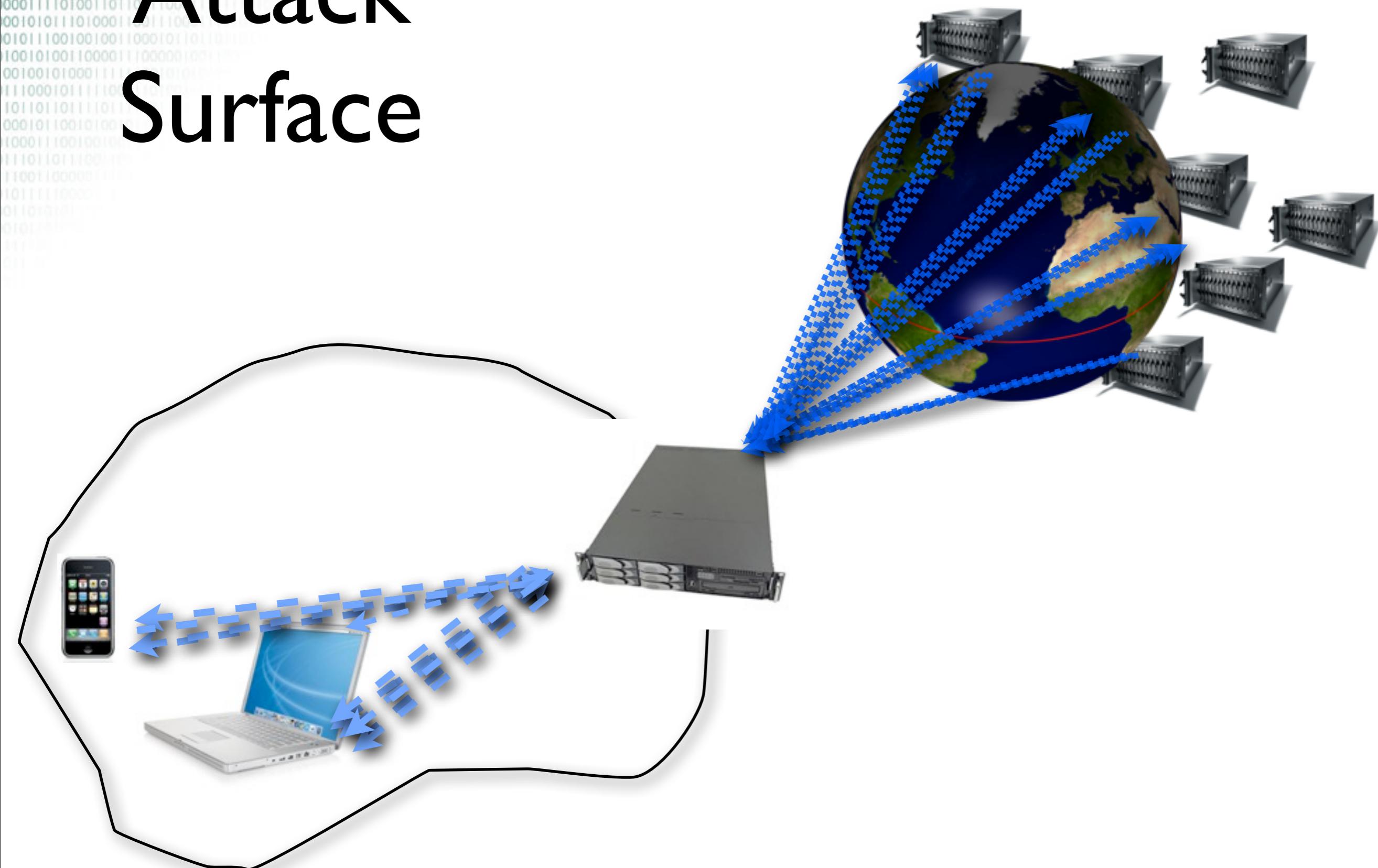
root.hints: location of the root server



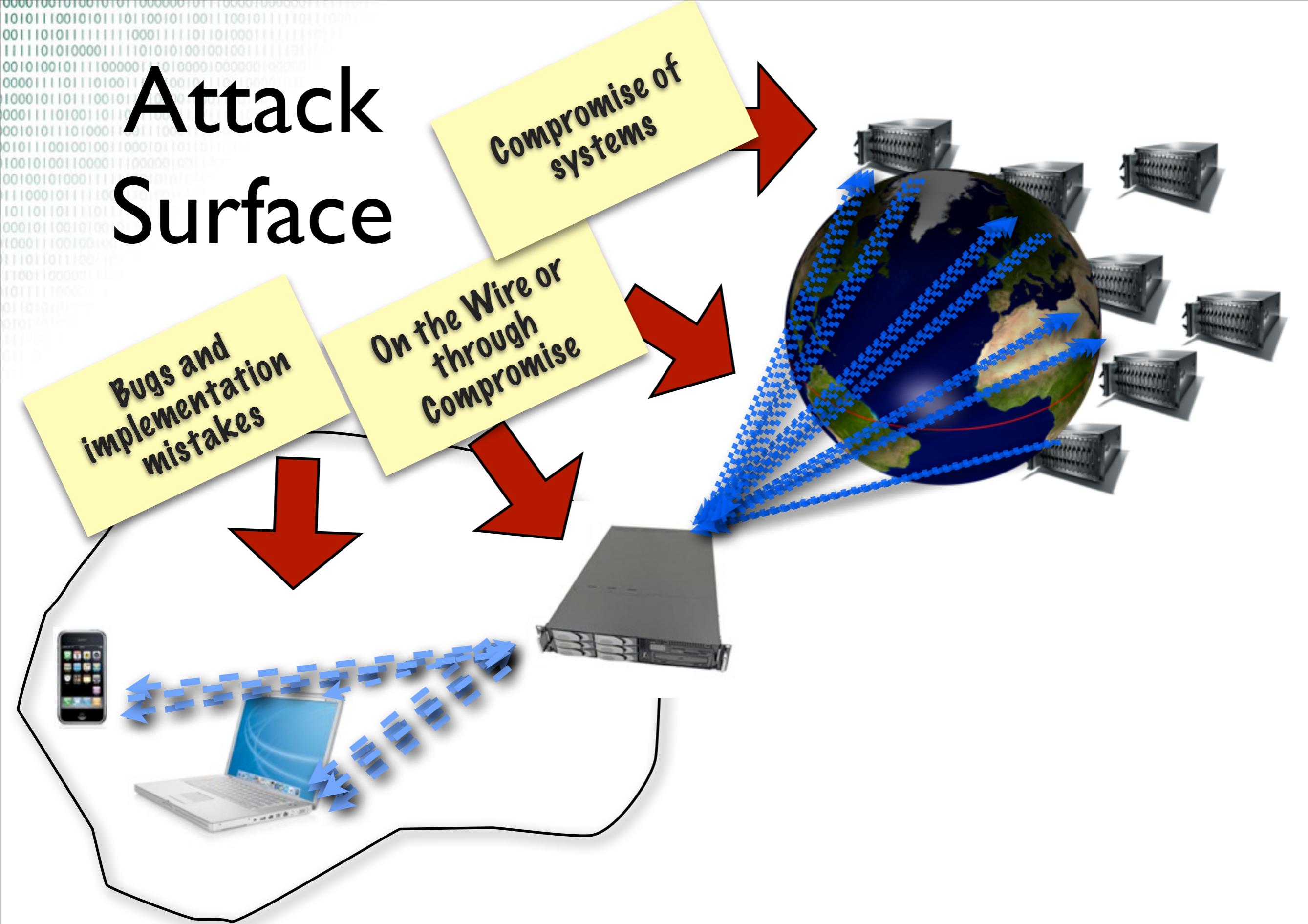
root.hints: location of the root server



Attack Surface



Attack Surface



Attack Surface

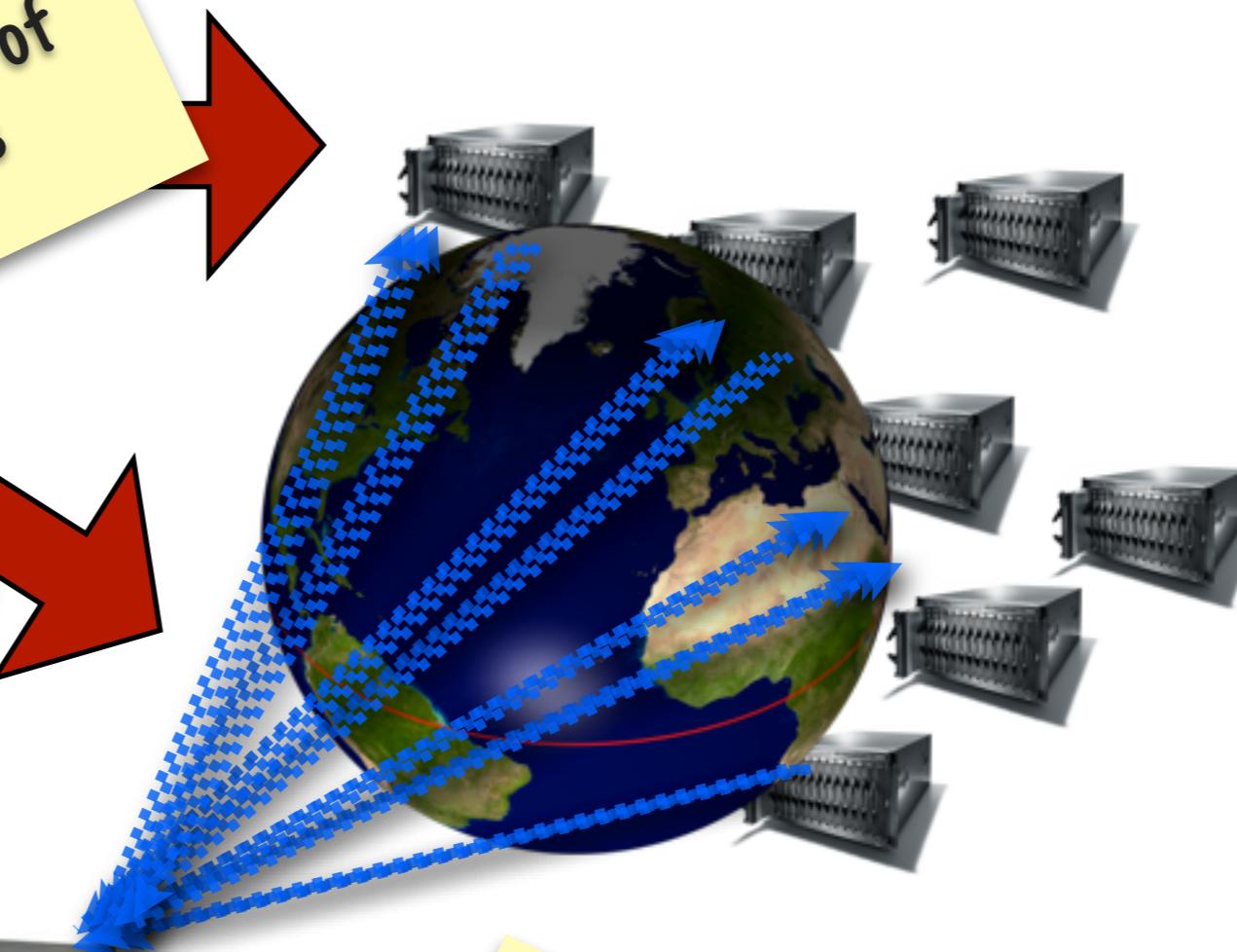
Bugs and implementation mistakes

On the Wire or through Compromise

Compromise of systems



Whoa, that looks bad!!! Who Uses This System?



enterprise

Recursive DNS



Mail server

NLnet
Labs

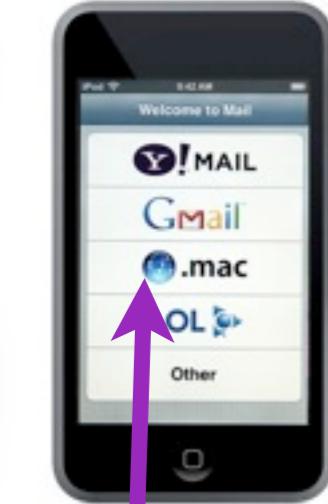
enterprise

Recursive DNS



Mail serve

Mail server



enterprise

Recursive DNS



Mail server

NLnet
Labs

enterprise

Recursive DNS



Mail server



enterprise

Recursive DNS



Mail server

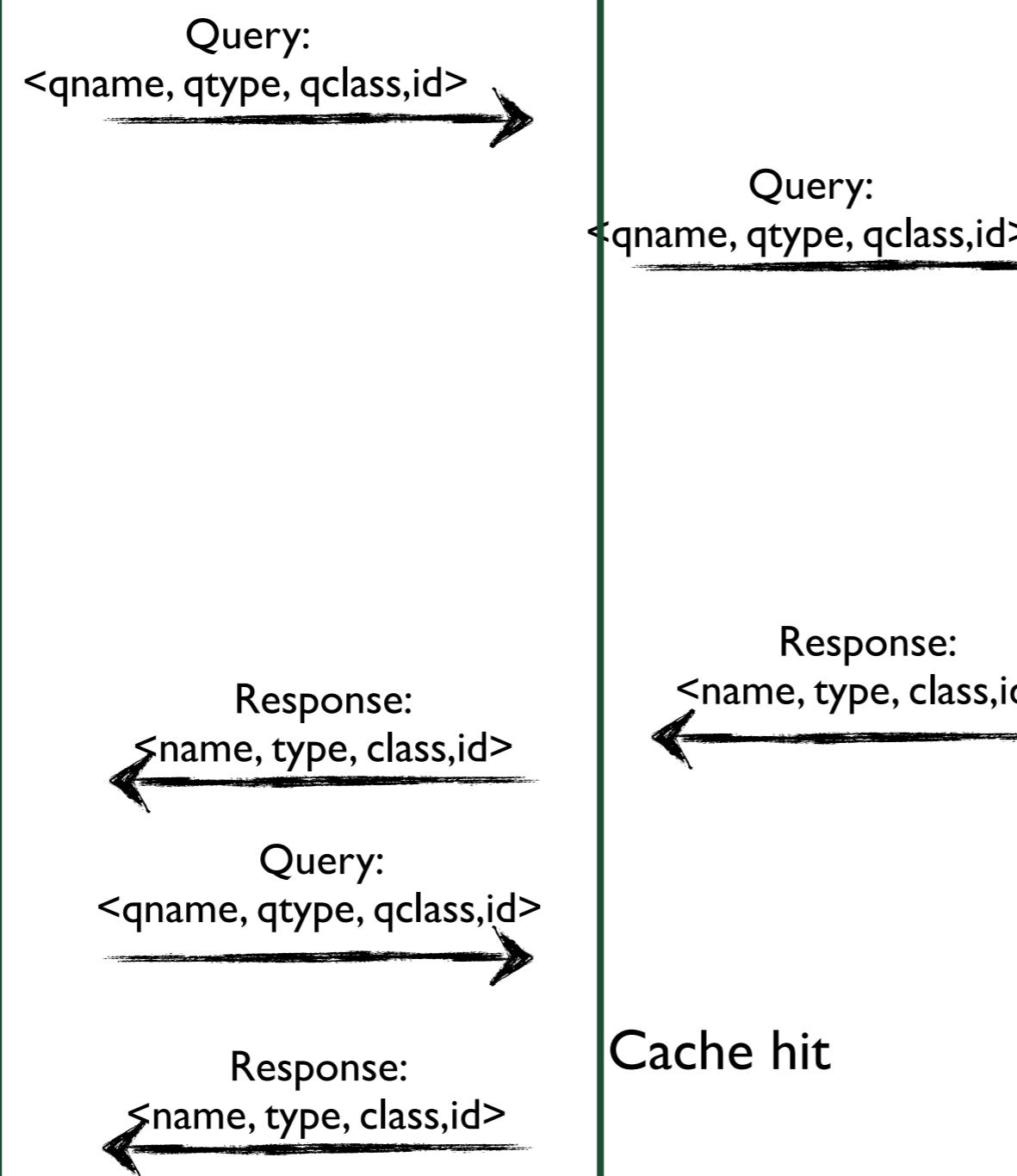
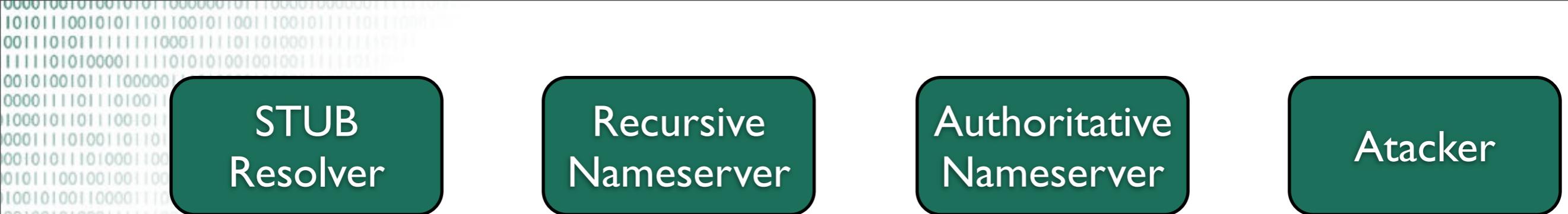


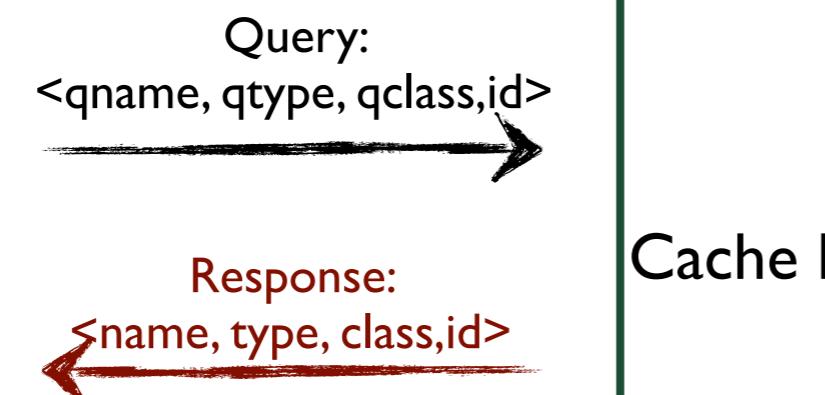
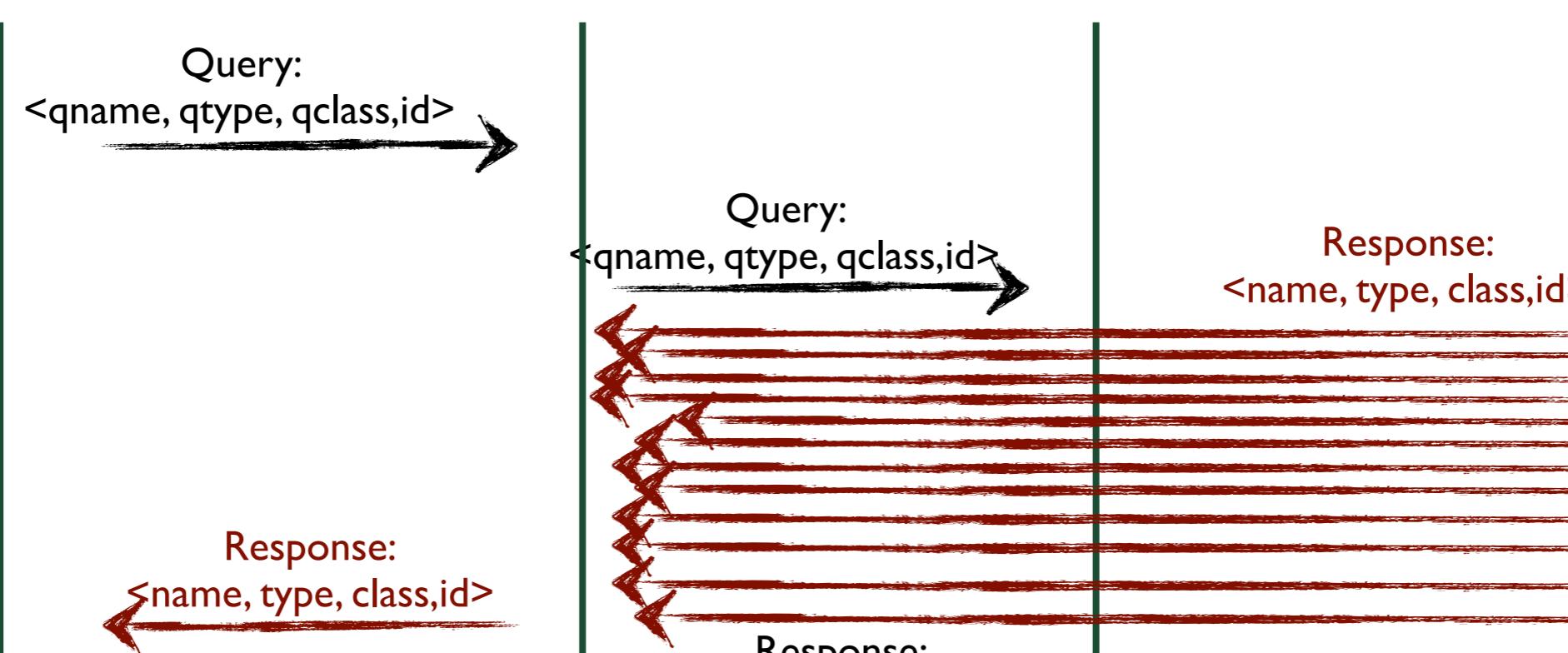
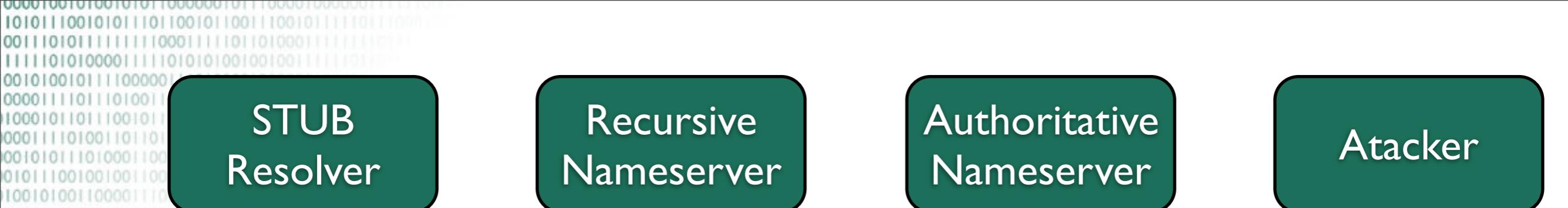
A black and white photograph of a dark fedora hat with a wide, slightly curved brim. A thin, light-colored band with a decorative edge runs around the base of the crown. The hat is shown from a three-quarter front angle, resting on a plain, light-colored surface.

Internet



Mail server





Success depends on legacy and speed of network.



And on various properties that the attacker needs to match

Response:
<name, type, class,id>

Query:
<qtype, qclass,id>

Response:
<name, type, class,id>

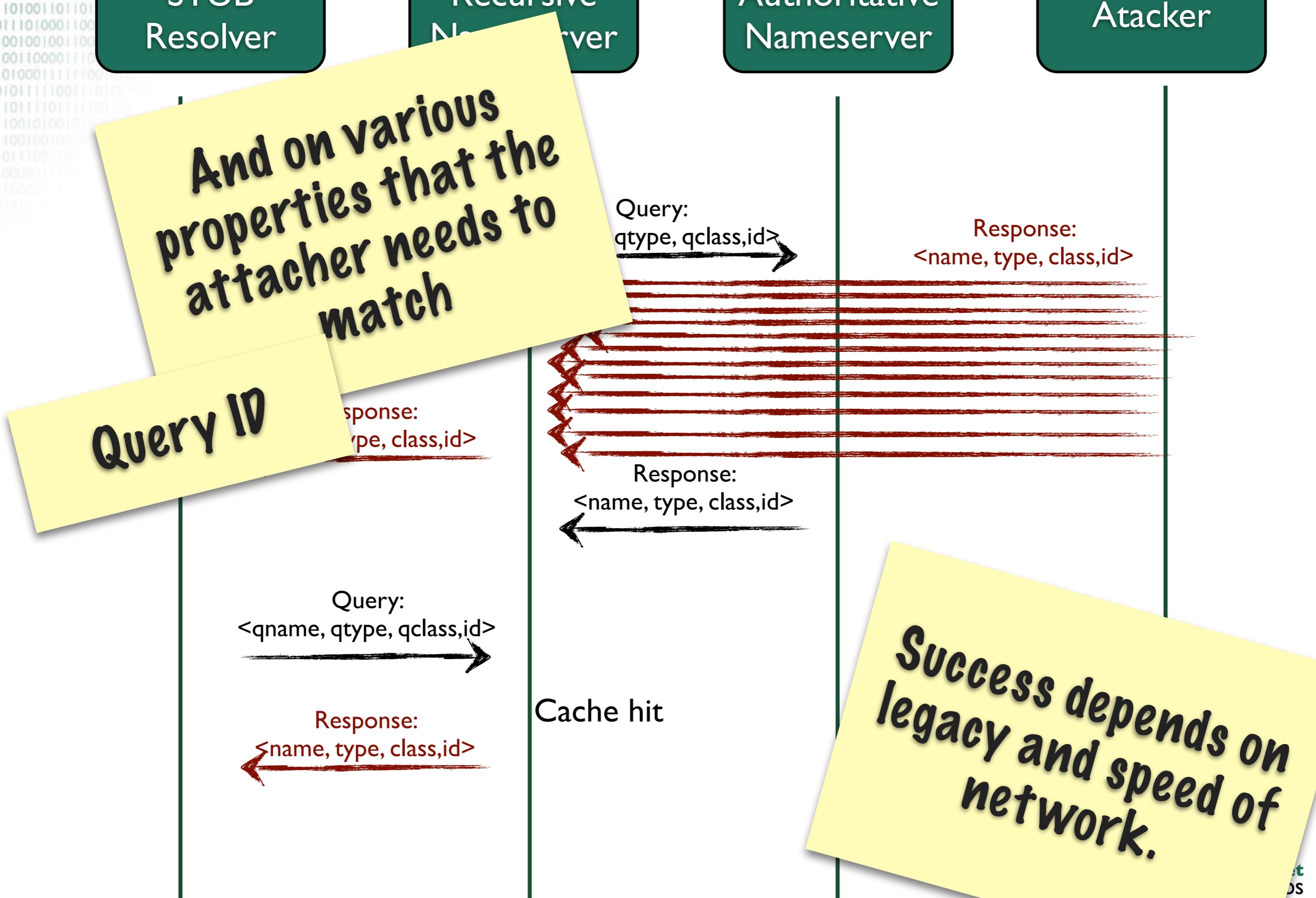
Query:
<qname, qtype, qclass,id>

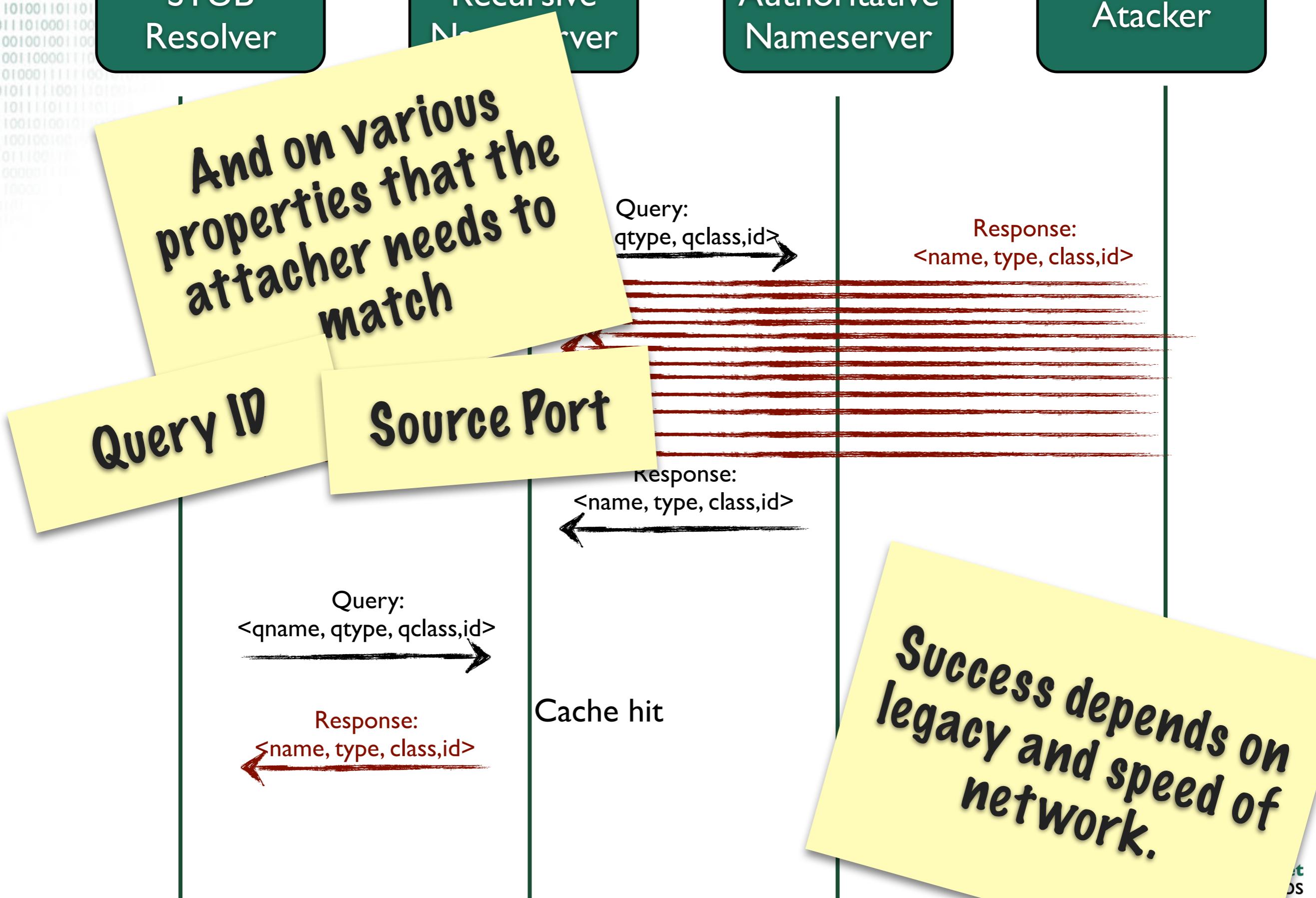
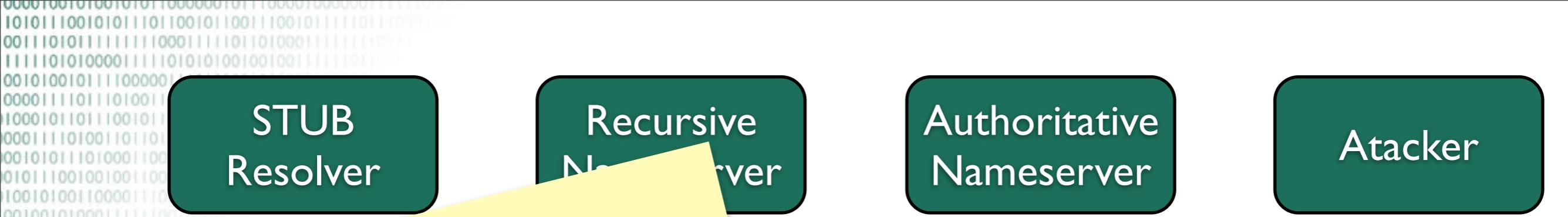
Response:
<name, type, class,id>

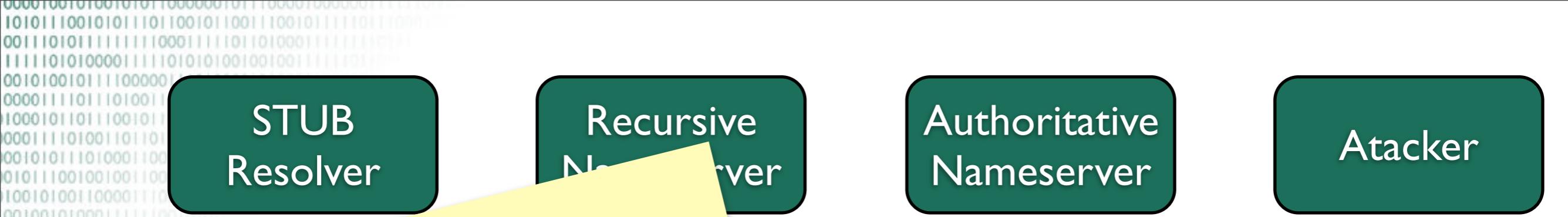
Response:
<name, type, class,id>

Cache hit

Success depends on legacy and speed of network.







And on various properties that the attacker needs to match

Query ID

Source Port

0X20

<qname, qtype, qclass, id>

Response:
<name, type, class, id>

Response:

<name, type, class, id>

Cache hit

Query:
<qtype, qclass, id>

Response:
<name, type, class, id>

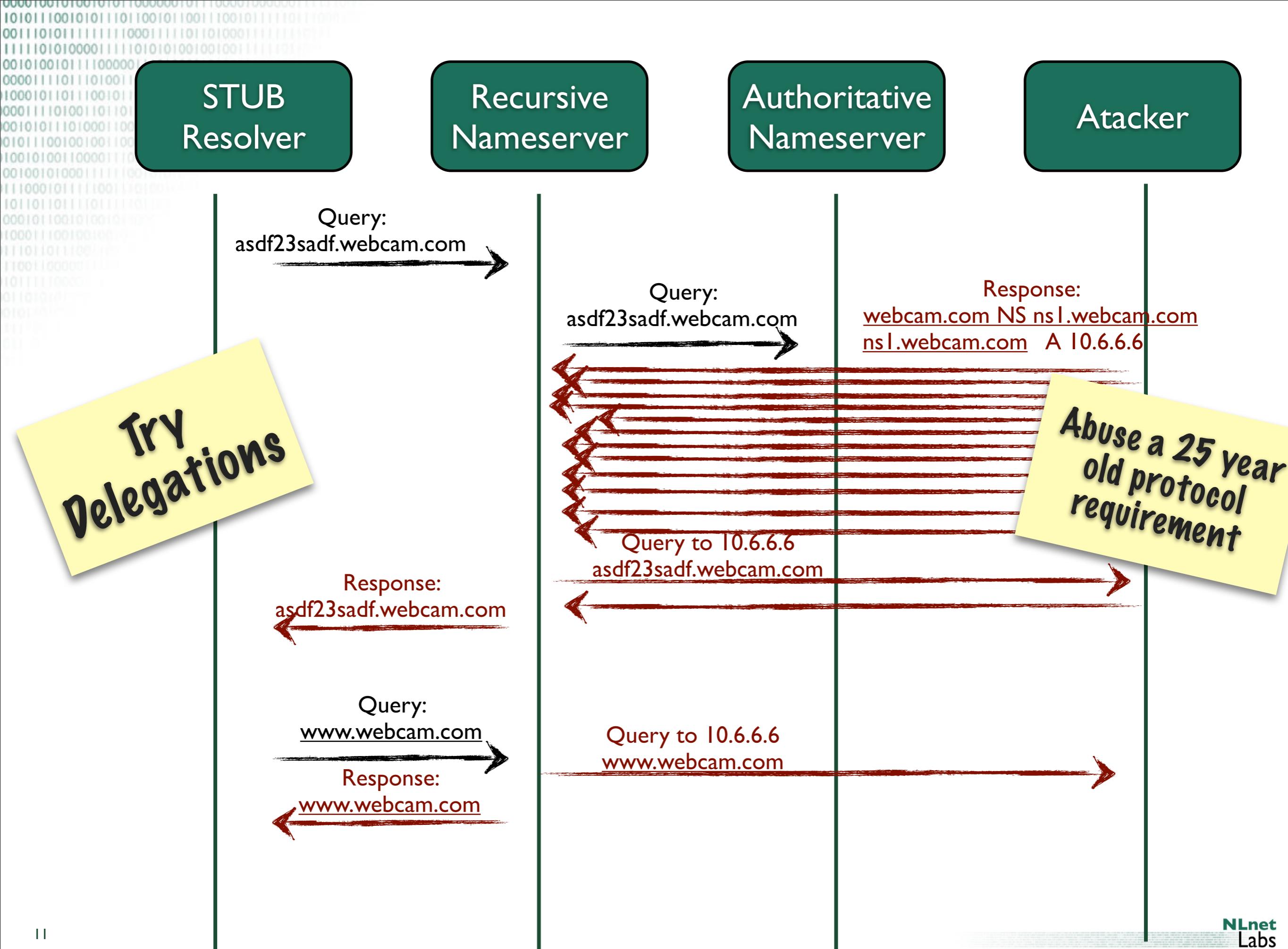
Success depends on legacy and speed of network.

TTL saves you?!?
I don't think so....



Dan Kaminsky's image from zdnet.com

Security
Popstar



Do attacks
happen in
practice?

Would you notice?



Do attacks
happen in
practice?

Why would one
attack the DNS?

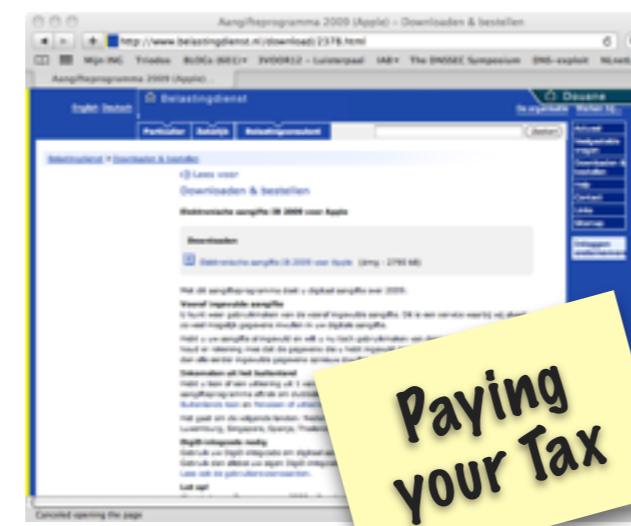
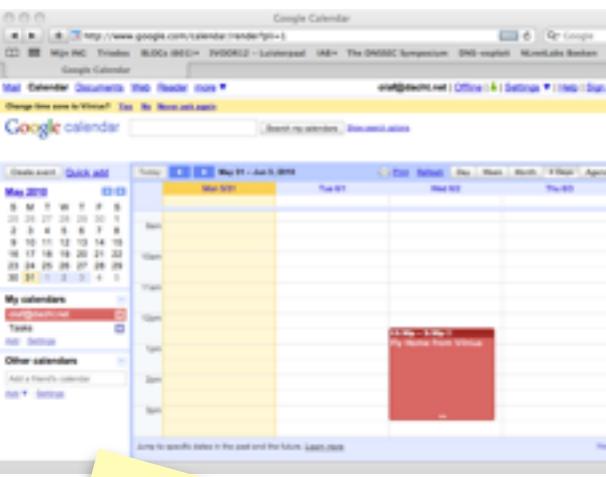
While one could be
doing other things



How to
Protect?

Why would one
attack the DNS?

Follow the Money



Money

Don't all these
transactions use
SSL and
Certificates?

The role of a
CA

3rd party
trust broker



The role of a
CA

3rd party
trust broker

Subject
Requests



The role of a
CA

3rd party
trust broker

Subject
Requests

RA performs
checks



The role of a CA

3rd party trust broker

Subject Requests

RA performs checks

RA tells CA to sign



The role of a CA

3rd party trust broker



Subject Requests

RA performs checks

RA tells CA to sign

Browser trusts CA signed certificates

The role of a CA

3rd party trust broker



EV
Extended validation

Subject Requests

RA performs checks

RA tells CA to sign

Browser trusts CA signed certificates

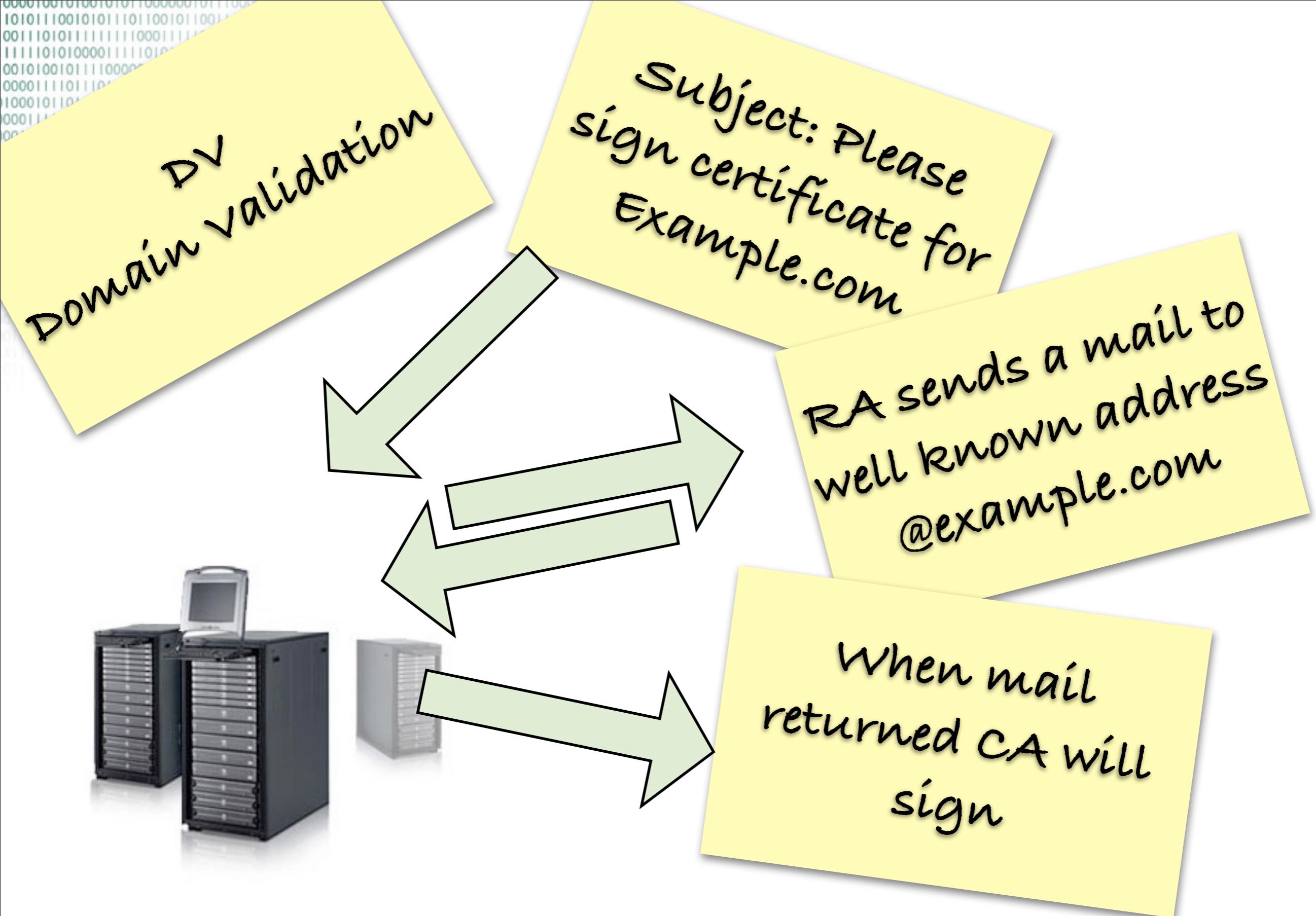


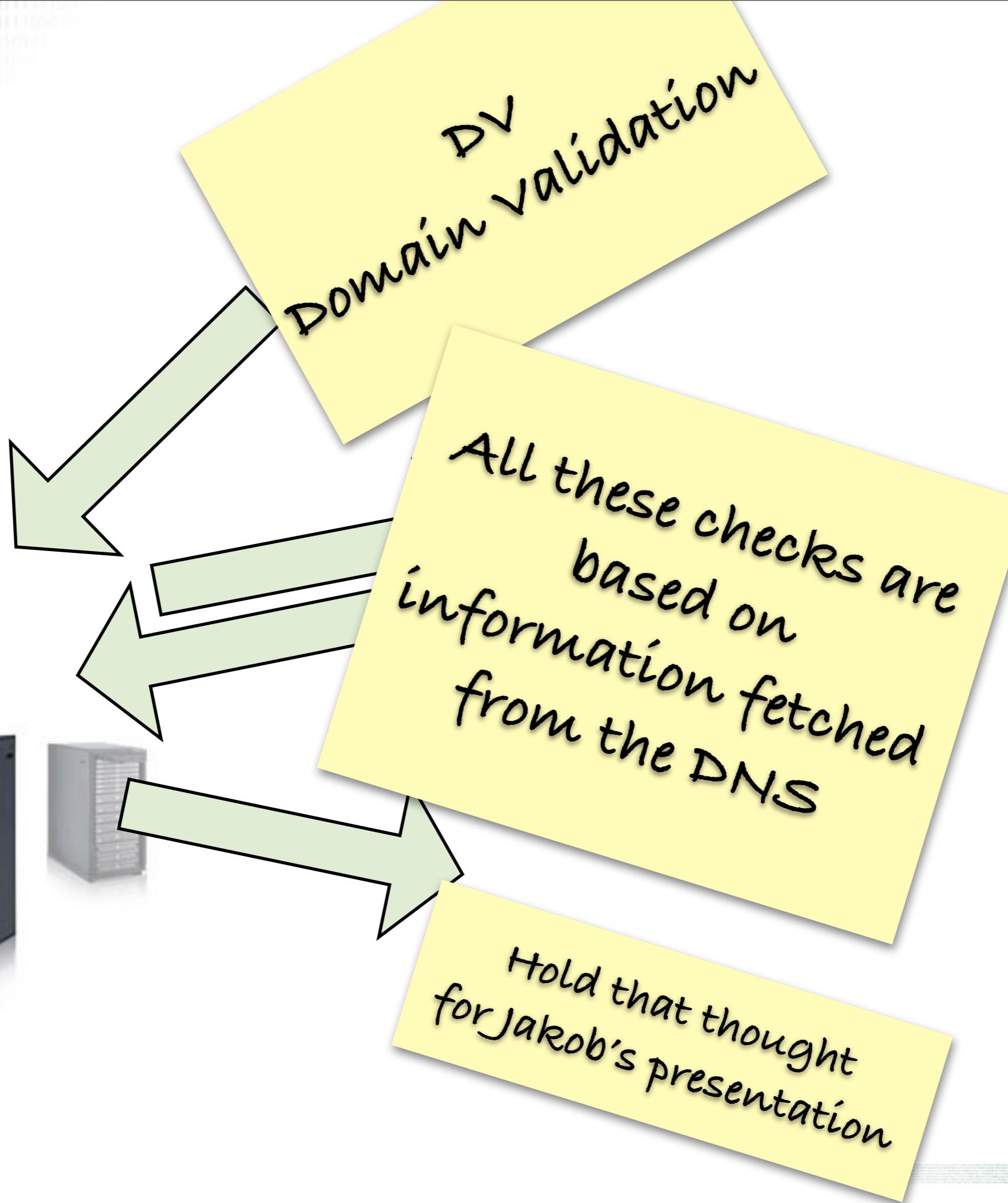
However all these
little men are a wee
bit expensive



However all these little men are a wee bit expensive

AUTOMATE THE LOT





Provisioning Vulnerabilities

Registrars
& Registrants



Man in the Middle



Registry

primary
DNS



Secondary
DNS



Secondary
DNS

Server vulnerability

spoofing
&
Man in the
Middle



What can one do to
protect...
(skipping DNSSEC)

Taking Unbound as example



Other servers might make other choices, but any modern resolver takes similar approaches



Security Choices in Unbound

- In general, a modern paranoid resolver
- DNSSEC support.
- RFC 2181 support completely
 - Fine grained. Keeps track of where RRSets came from and won't upgrade them into answers.
- Does not allow RRSets to be overridden by lower level rrsets

Filtering

- Scrubber:
- Only in-bailiwick data is accepted in the answer
 - The answer section must contain only answer
 - CNAME, DNAME checked that chain is correct
 - CNAME cut off and only the first CNAME kept
 - Lookup rest yourself do not trust other server
 - DNAME synthesize CNAME by unbound do not trust other server. Also cut off like above.
 - DNAME from cache only used if DNSSEC-secure.

Filtering II

- No address records in authority, additional section unless relevant – i.e. mentioned in a NS record in the authority section.
- Irrelevant data is removed
 - When the message only had preliminary parsing and has not yet been copied to the working region of memory

Entropy

- Randomness protects against spoof
 - Arc4random() (OpenBSD): crypto strong.
May not be perfectly random, but predicting it is a cryptographical breakin.
 - Real entropy from OS as seed
 - Query id – all 16 bits used.
 - Port randomisation – uses all 16bits there, goes out of its way to make sure every query gets a fresh port number

Entropy II

- Destination address, and ipv4/ipv6. RTT band of 400msec (=everything).
 - Its not the timewindow but the randomness
- Query aggregation – same queries are not sent out – unless by different threads
- QName strict match checked in reply
- 0x20 option
- Harden-referral-path (my draft) option
- Can use multiple source interfaces!
 - 4 outgoing IP address add +2 bits

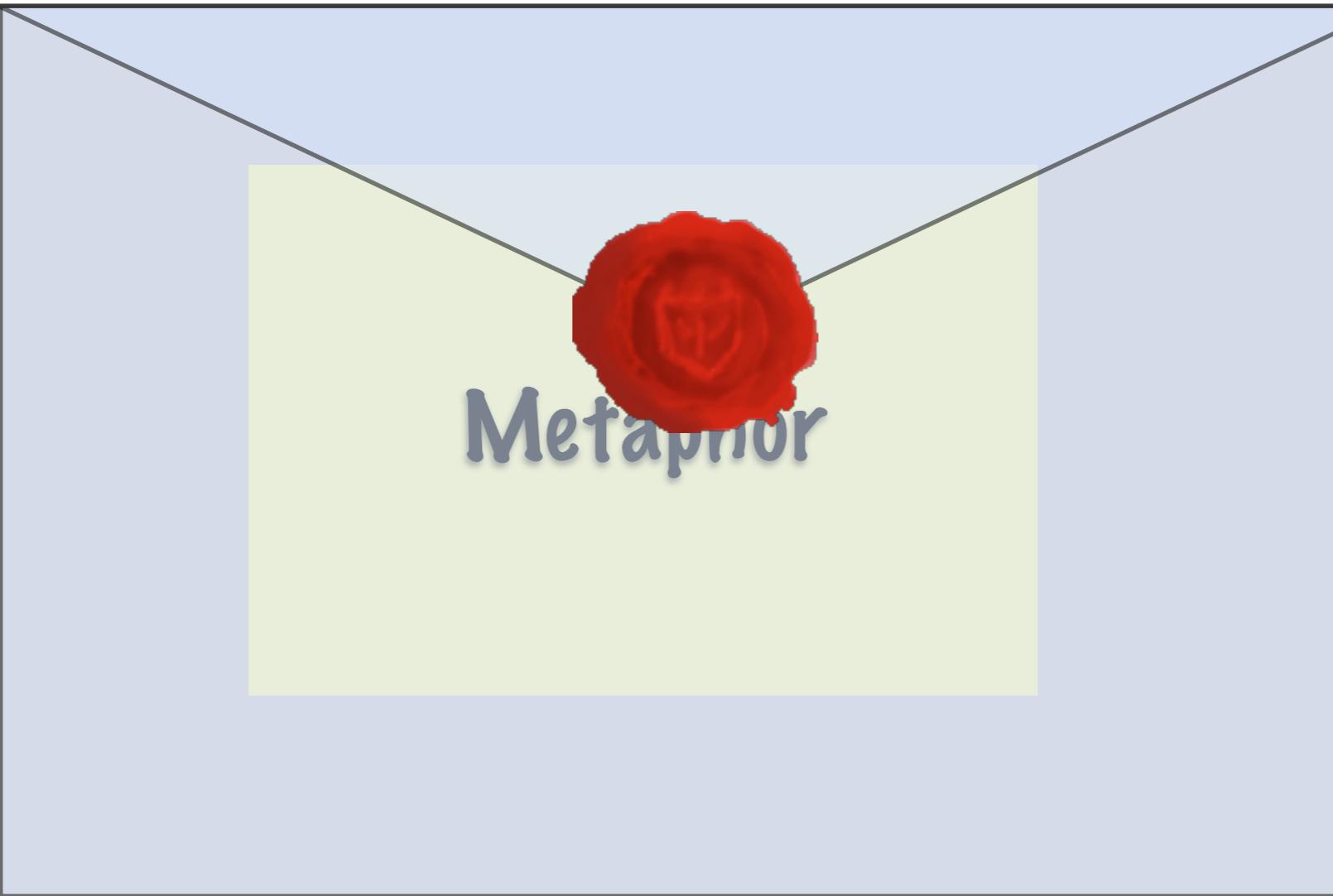
Other measures

- Not for the wire itself
 - Heap function pointer protection (whitelisted)
 - Chroot() by default
 - User privileges are dropped (lots of code!)
 - ACL for recursion
 - No detection of attacks – assume always under attack
 - version.bind hostname.bind can be blocked or configured what to return (version hiding)
 - Disprefer recursion lame servers – they have a cache that can be poisoned



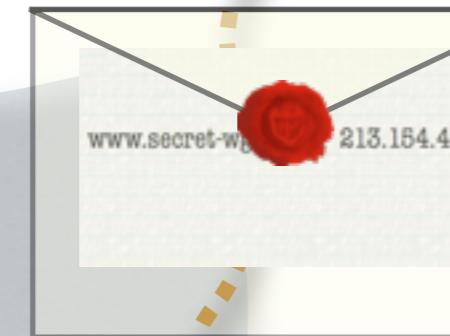
Arms Race...

Introducing
DNSSEC



End to End Security

**Registrars
& Registrants**

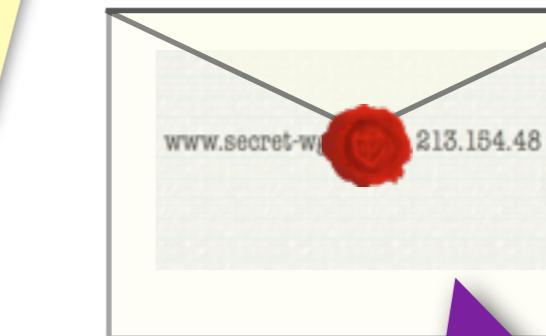


Registry

**primary
DNS**



**Secondary
DNS**



**Secondary
DNS**



DNSKEY:
public key from
the keypair

RRSIG: Signatures
made with a private
key from the keypair

All done using
Public Key crypto

NSEC and NSEC3
For pre-calculated
Denial of Existence

DS
For delegating
Security

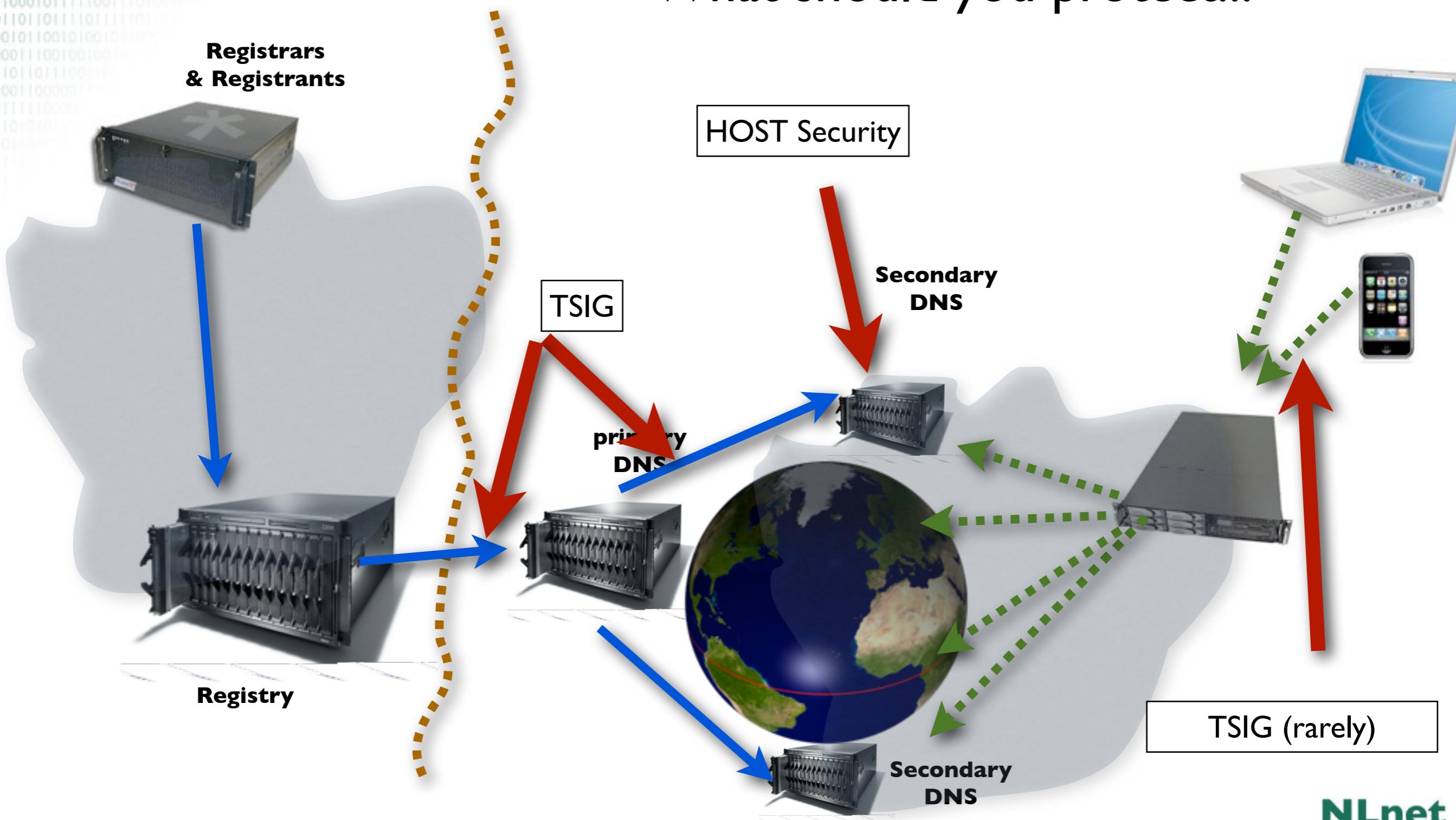
But more on that later

Let us have a look at
another cryptographic
DNS protection
mechanism

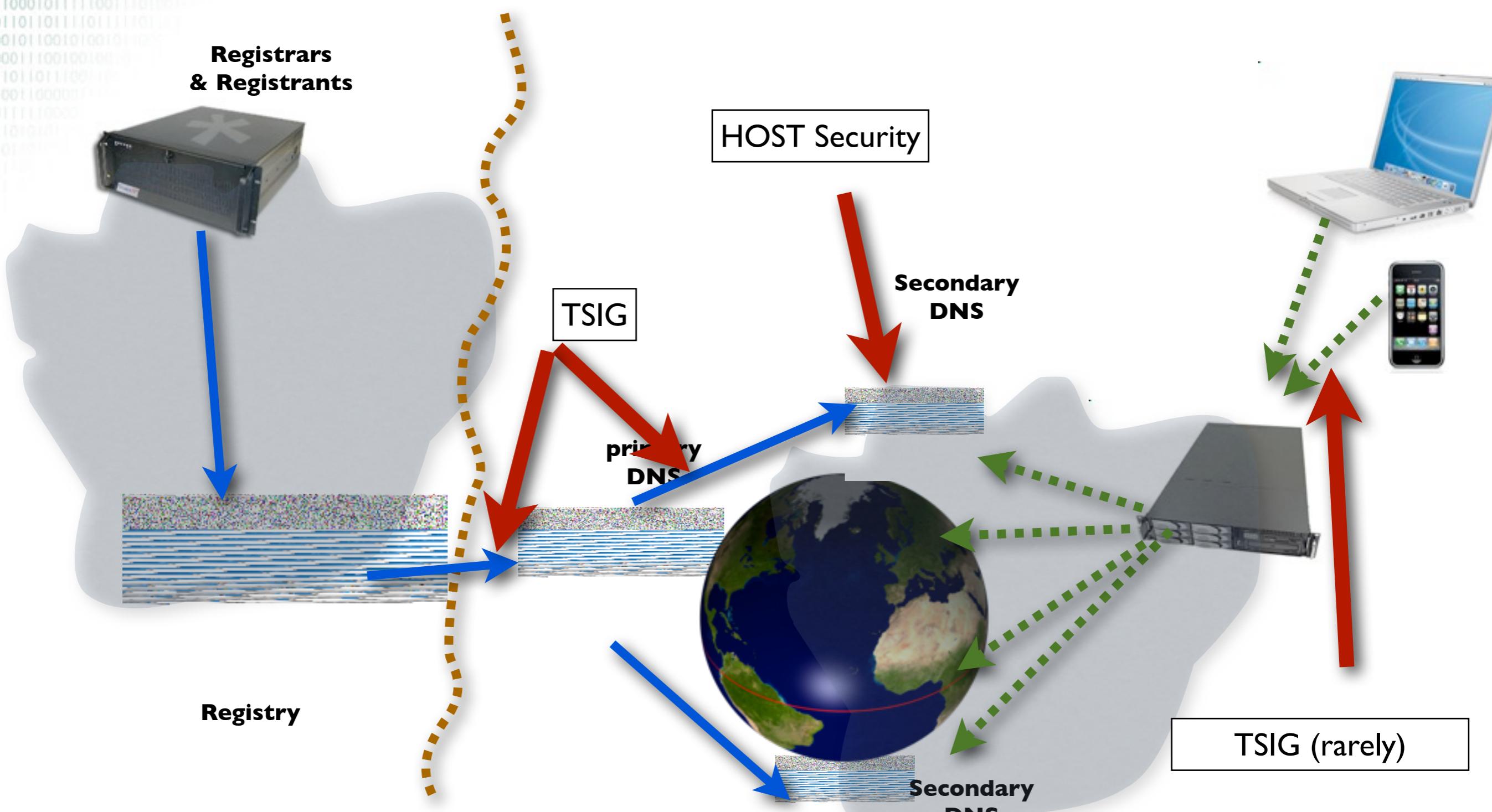
Securing Host-Host Communication

Data flow through the DNS

What should you protect...



Data flow through the DNS

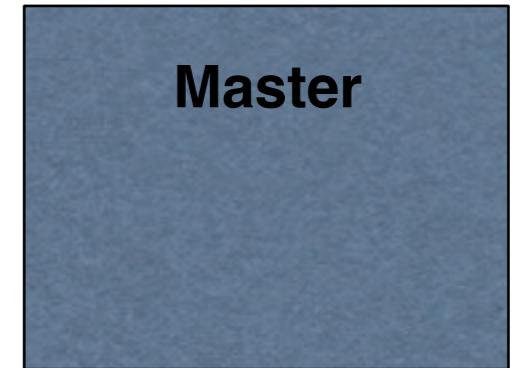


Transaction Signature: TSIG

- TSIG (RFC 2845)
 - Authorising dynamic updates and zone transfers
 - Authentication of caching forwarders
 - Independent from other features of DNSSEC
- One-way hash function
 - DNS question or answer and timestamp
- Traffic signed with “shared secret” key
- Used in configuration, **NOT** in zone file

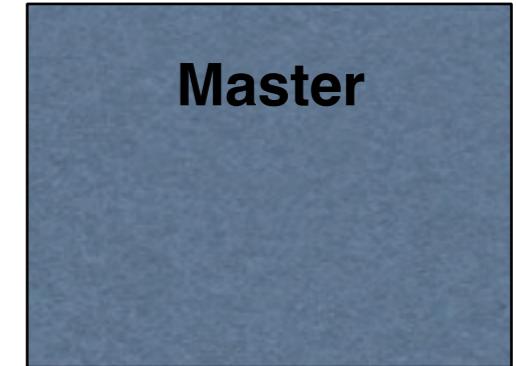
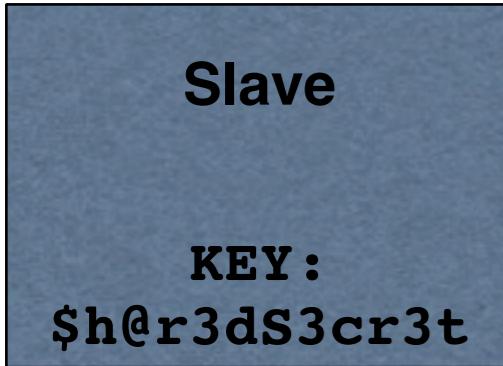
TSIG Example

Query: AXFR



TSIG Example

Query: AXFR



TSIG Example

Query: AXFR

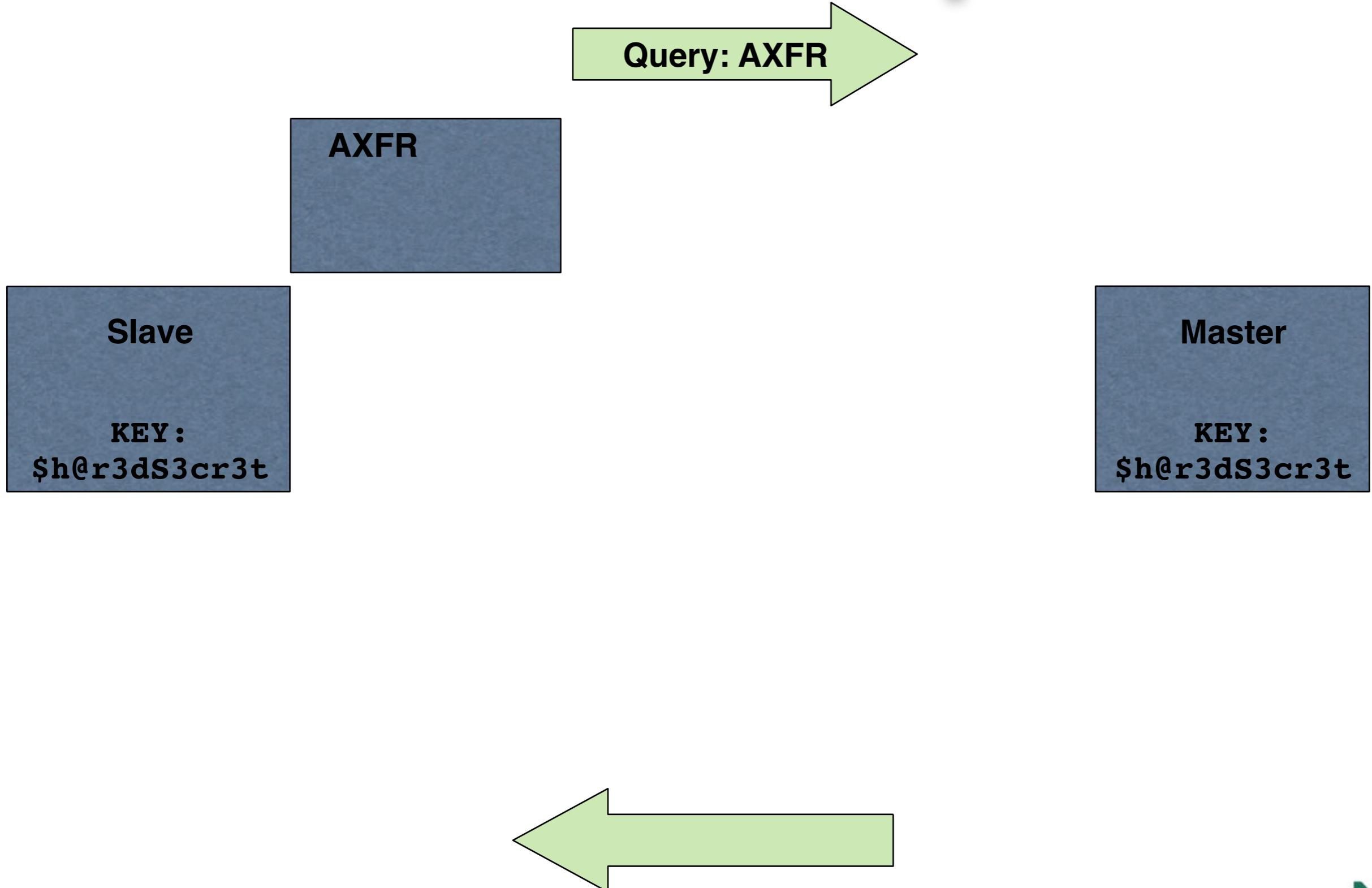
Slave

KEY:
\$h@r3dS3cr3t

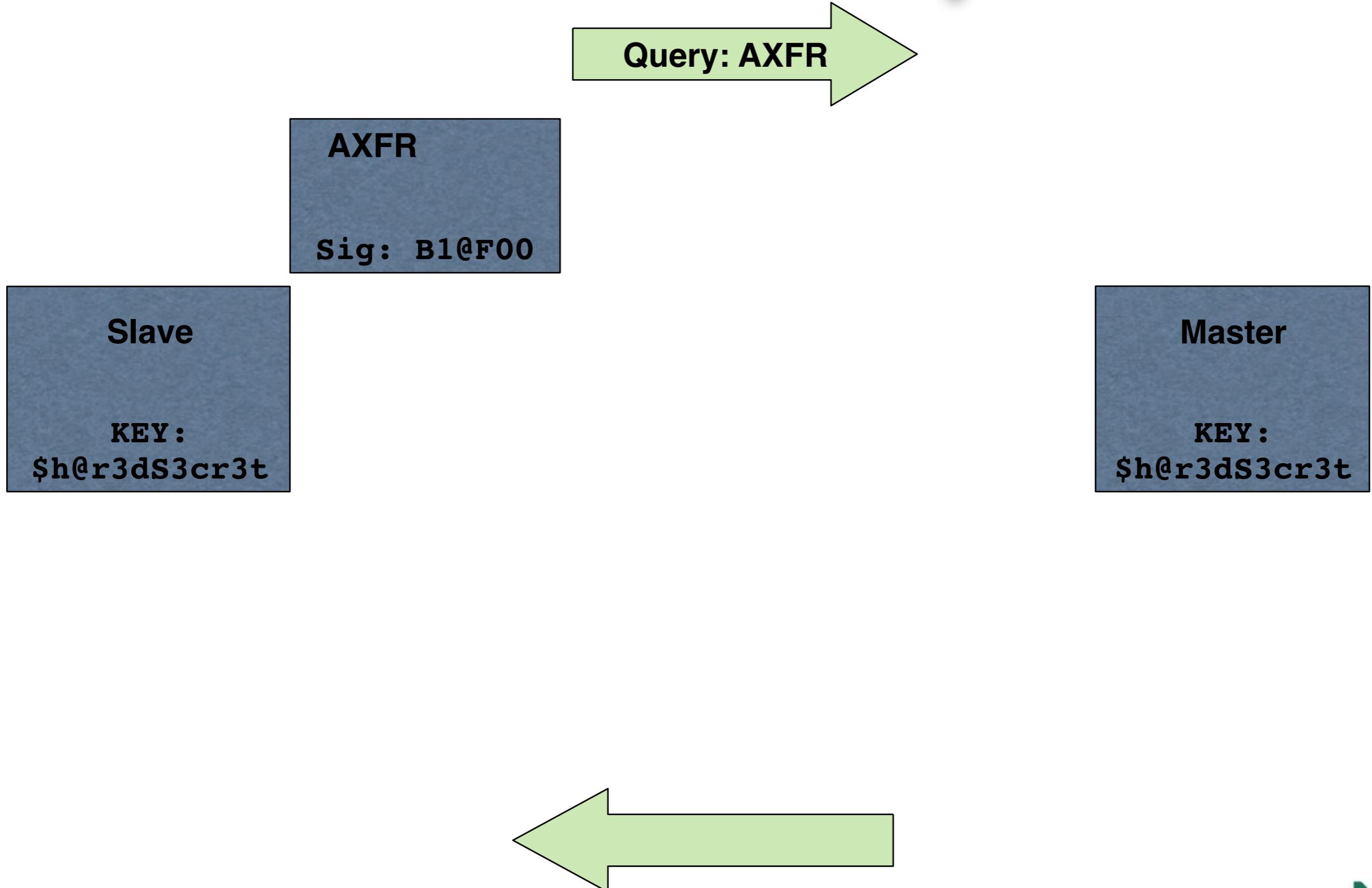
Master

KEY:
\$h@r3dS3cr3t

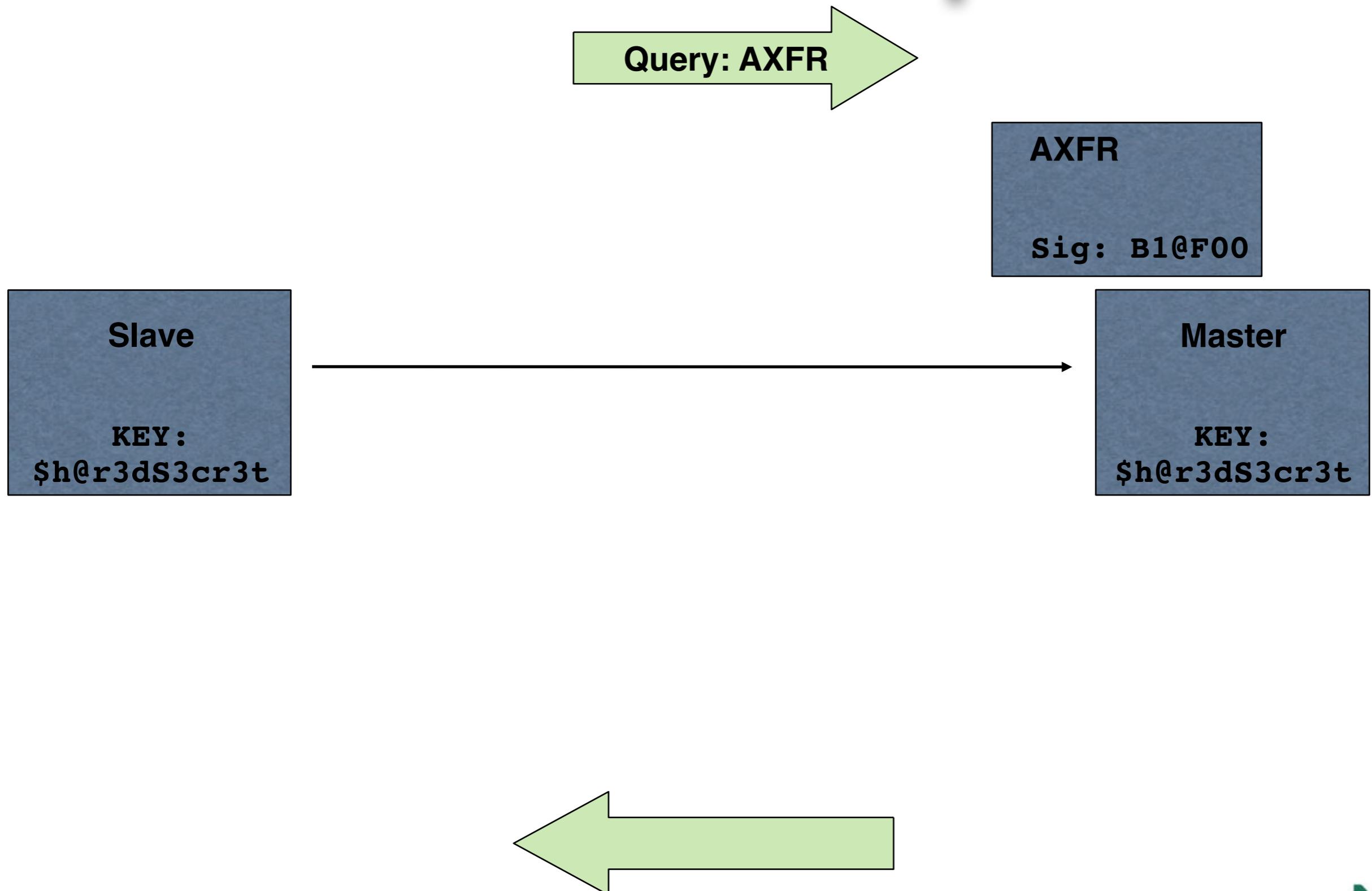
TSIG Example



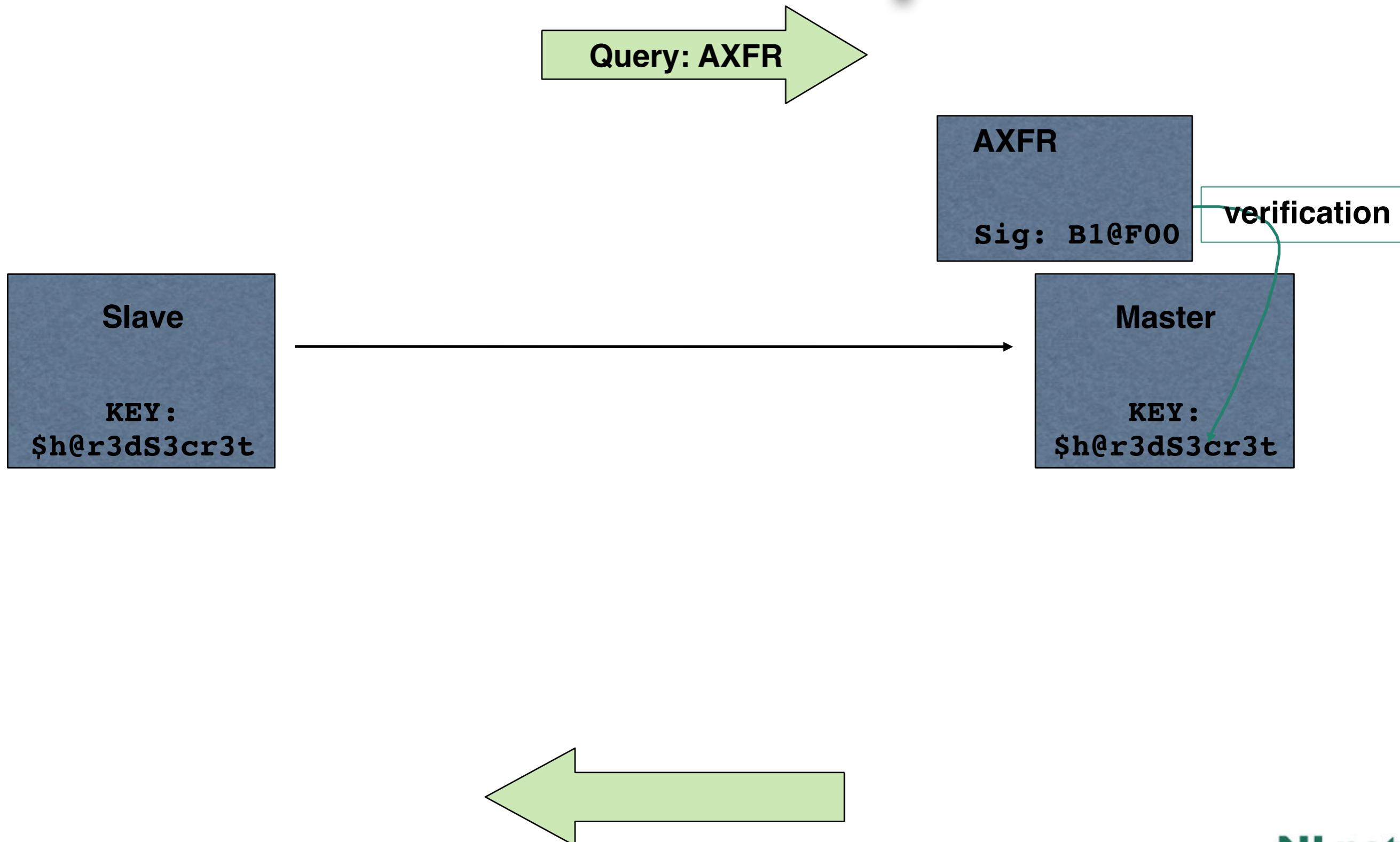
TSIG Example



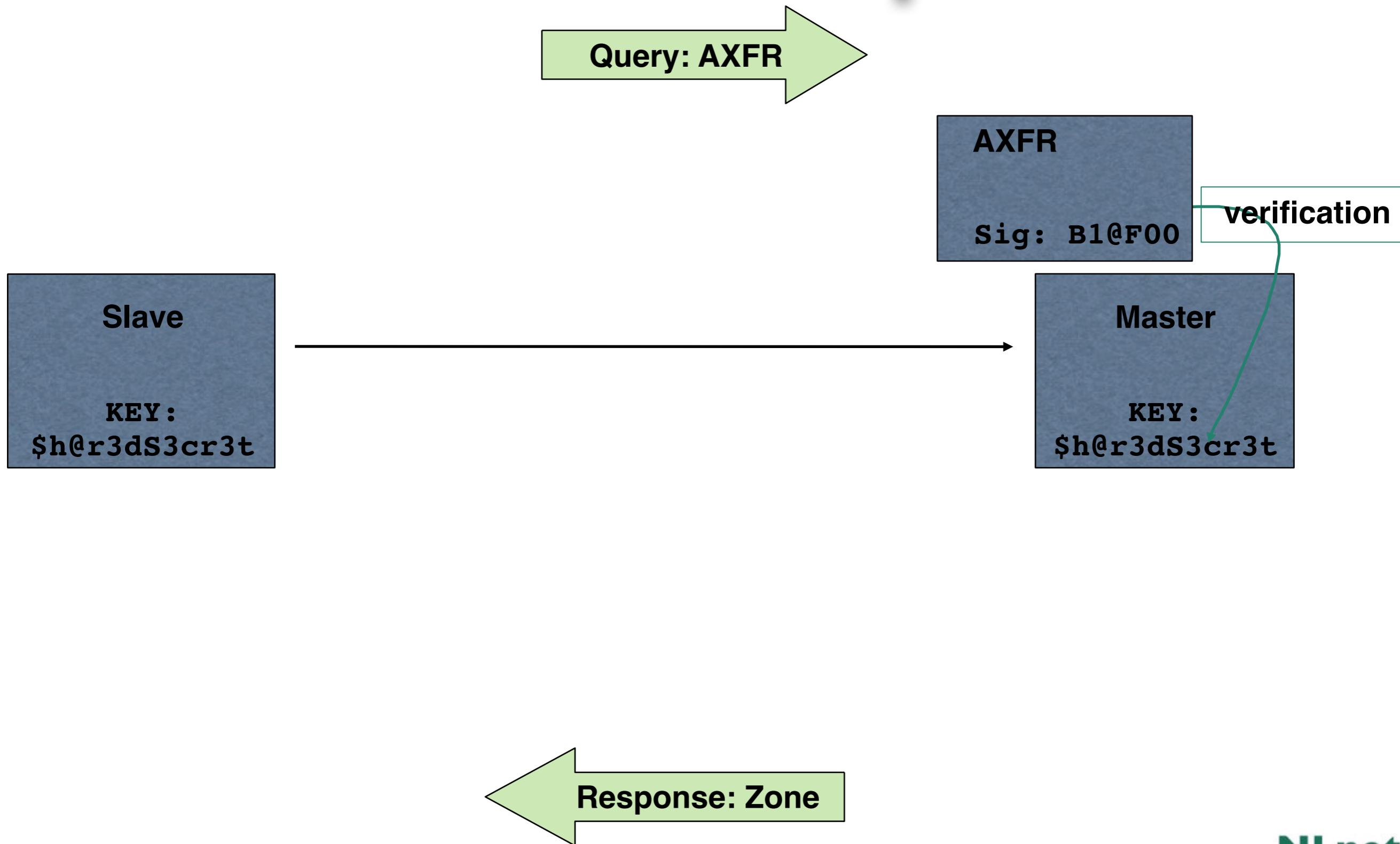
TSIG Example



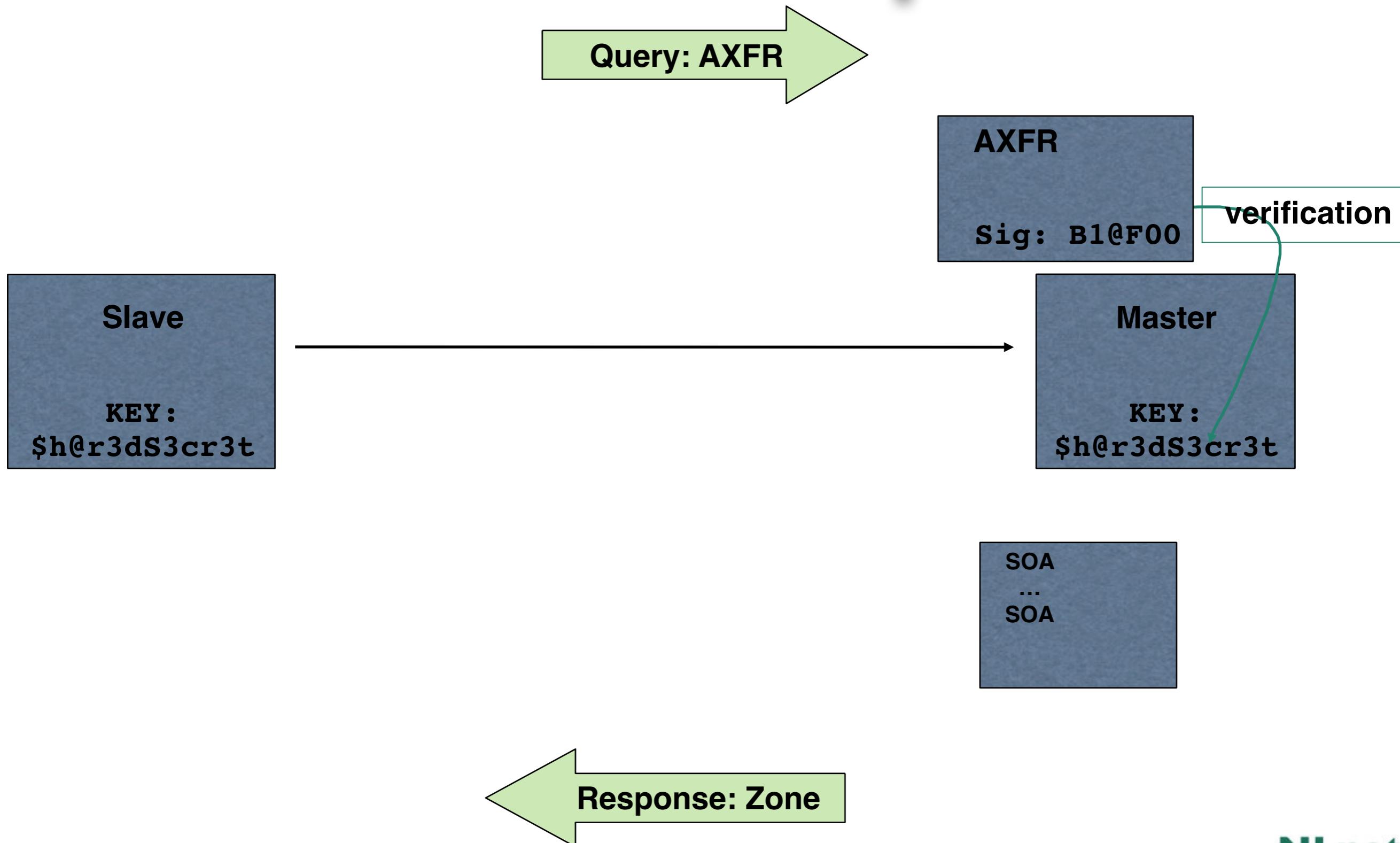
TSIG Example



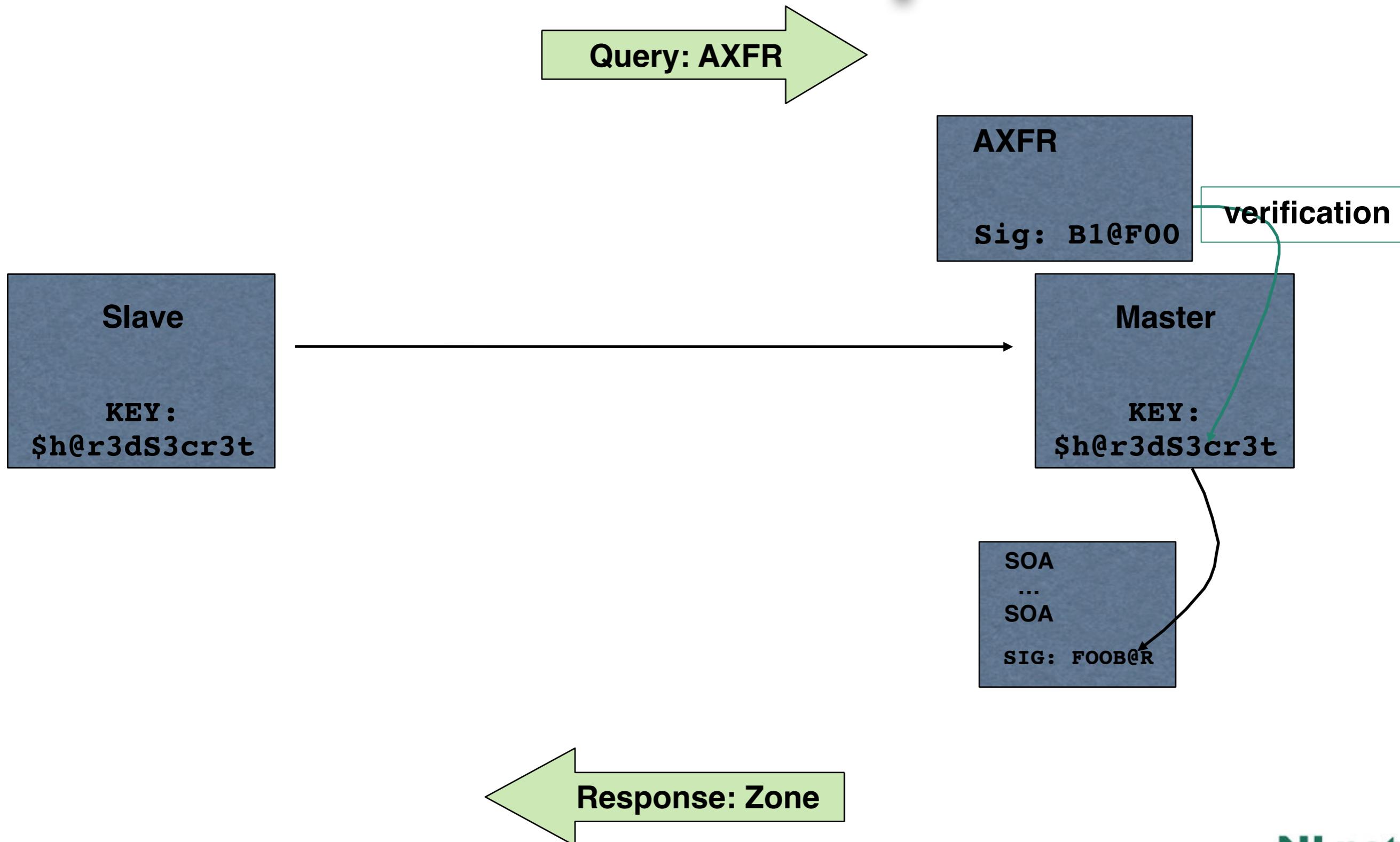
TSIG Example



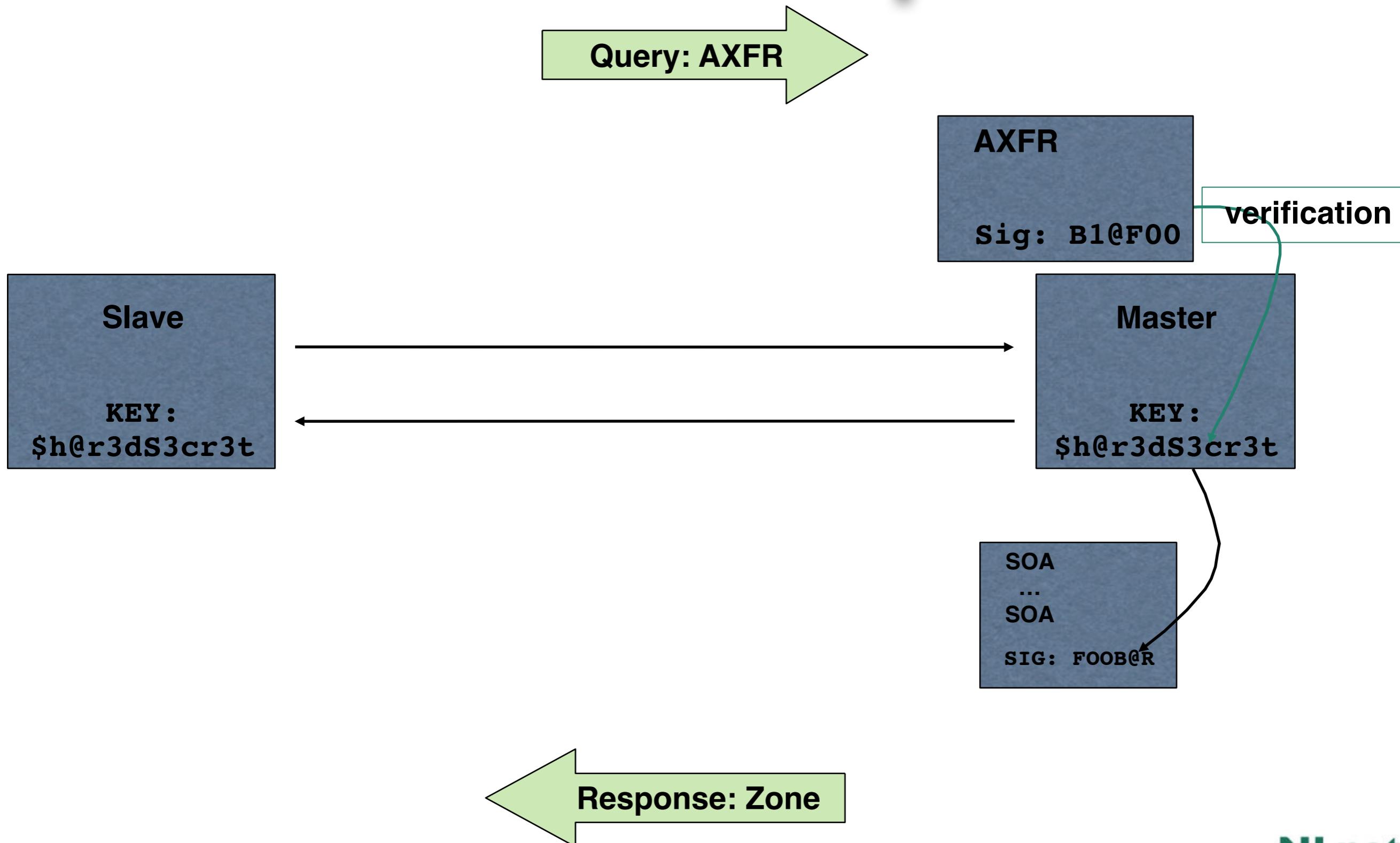
TSIG Example



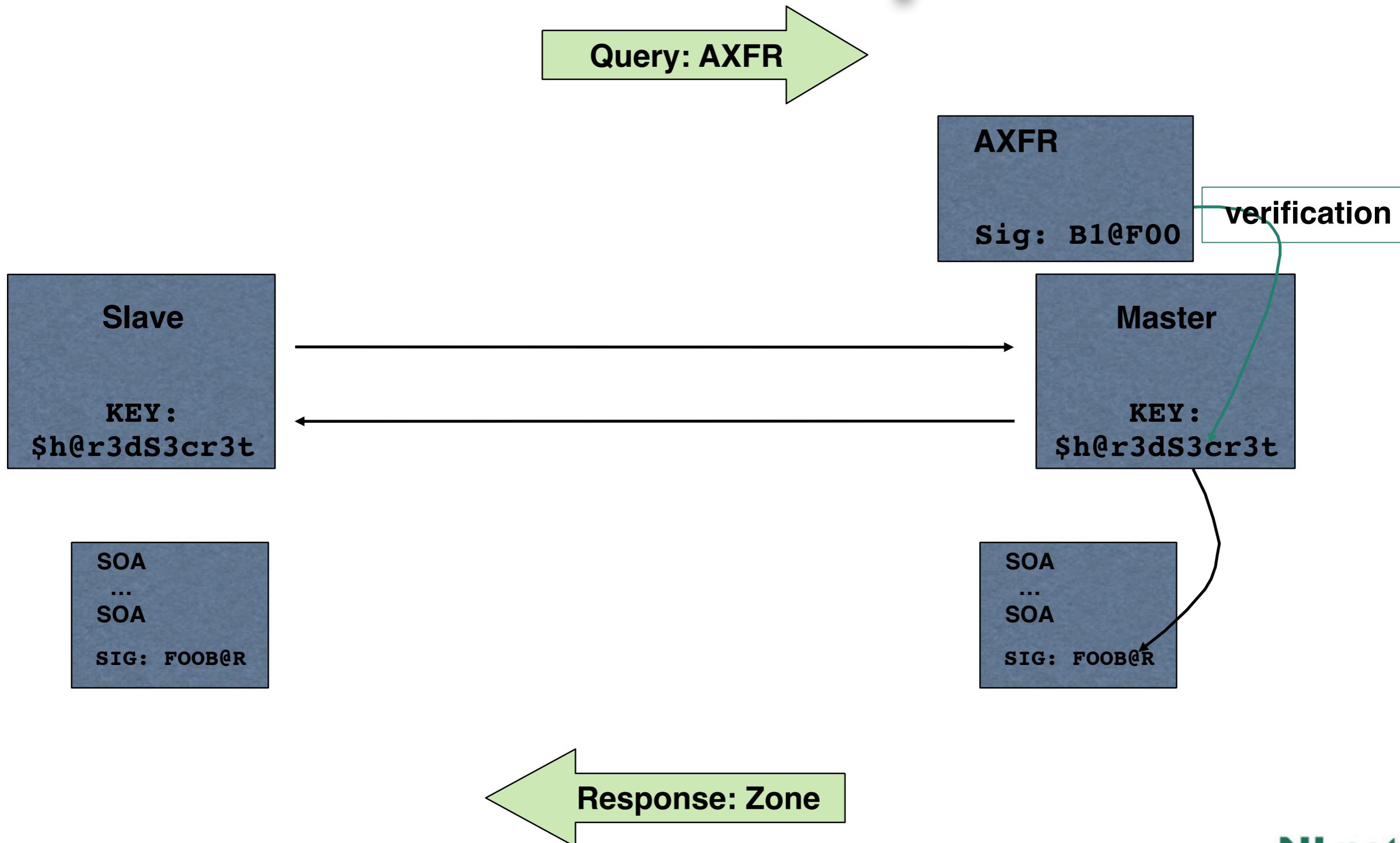
TSIG Example



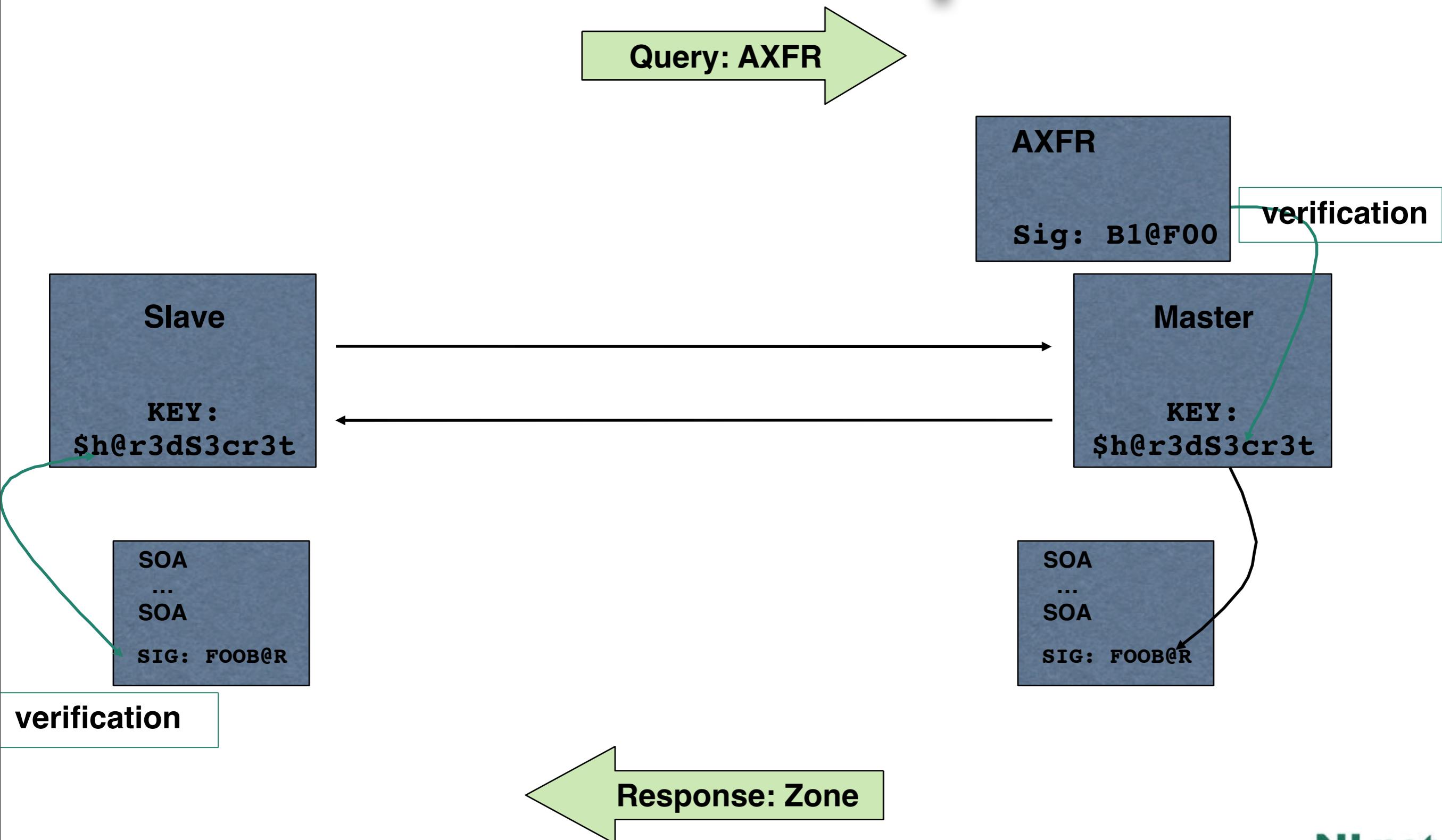
TSIG Example



TSIG Example



TSIG Example



TSIG for Zone Transfers

1. Generate secret
2. Communicate secret
3. Configure servers
4. Test

Importance of the Time Stamp

- TSIG/SIG(0) signs a complete DNS request / response with time stamp
 - To prevent replay attacks
 - Currently hardcoded at five minutes
- Operational problems when comparing times
 - Make sure your local time zone is properly defined
 - `date -u` will give UTC time, easy to compare between the two systems
 - Use NTP synchronisation!

Authenticating Servers Using SIG(0)

- Alternatively, it is possible to use SIG(0)
 - Not yet widely used
 - Works well in dynamic update environment
- Public key algorithm
 - Authentication against a public key published in the DNS
- SIG(0) specified in RFC 2931

Cool Application

- Use TSIG-ed dynamic updates to configure your laptop's name
- My laptop is known by the name of aagje.secret-wg.org
 - <http://ops.ietf.org/dns/dynupd/secure-ddns-howto.html>
 - Mac OS users: there is a bonjour based tool.
- www.dns-sd.org

