

# chapter 9

## Sensitivity Analysis

### 9.1 INTRODUCTION

Suppose we take a practical problem and formulate it as the LP

$$\begin{cases} \text{Minimize } z(x) = cx \\ \text{Subject to } Ax = b \\ x \geq 0 \end{cases} \quad (9.1)$$

where  $A$  is a matrix of order  $m \times n$  and rank  $m$ , and obtain an optimum feasible solution for it. In most real-life problems, the coefficients in  $A$ ,  $c$ , and  $b$  are estimated from practical considerations. After the final optimal feasible solution has been obtained we may discover that some of the entries in  $b$ ,  $c$ , or  $A$  have to be changed or that extra constraints or variables have to be introduced into the model. Solving the modified problem from scratch will be wasteful. *Sensitivity analysis* (also called *post-optimality analysis*) deals with the problem of obtaining an optimum feasible solution of the modified problem starting with the optimum feasible solution of the old problem. We consider the problem of introducing only one change at a time (e.g., how of only one  $c_j$  has to be changed or if only one new constraint has to be added, etc.) If several changes have to be made, make them one at a time, or extend the methods discussed here in an obvious manner to take care of several simultaneous changes. In Chapter 8 we discussed some types of postoptimality analysis where the right-hand-side vector or the cost coefficient vector vary linearly as a function of a parameter as it ranges over the real line. Here we discuss various other types of postoptimality analyses. Some of the postoptimality analyses discussed here, for example, changes in cost coefficients or the right-hand-side constants, can be viewed as specializations of the parametric analysis applied to those cases. For an idea of the kind of questions

that can be answered using the methods of sensitivity analysis, see Exercise 9.16. In most practical applications using a linear programming model, marginal and sensitivity analysis provide economic information that is very useful in planning.

Let  $K$  denote the set of feasible solutions of (9.1). Suppose the optimal basis obtained for (9.1) is  $\bar{B}$ , associated with the basic vector  $x_{\bar{B}} = (x_1, \dots, x_m)$ ,  $c_{\bar{B}} = (c_1, \dots, c_m)$ . Let  $\bar{\pi} = c_{\bar{B}}\bar{B}^{-1}$  and let  $\bar{x}$  be the BFS of (9.1) corresponding to the basis  $\bar{B}$ . For illustration we will use the following problem, which has  $x_{\bar{B}} = (x_1, x_2, x_3)$  as an optimum basic vector. Therefore,  $\bar{x} = (3, 4, 2, 0, 0)^T$  is an optimum feasible solution of this problem and the minimum objective value is  $z(\bar{x}) = 11$ .

Tableau 9.1 Original Problem

$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$-z$	$b$
1	2	0	1	0	-6	0	11
0	1	1	3	-2	-1	0	6
1	2	1	3	-1	-5	0	13
3	2	-3	-6	10	-5	1	0

$x_j \geq 0$  for all  $j$ ;  $z$  to be minimized.

Tableau 9.2 Optimum Inverse Tableau for the Problem in Tableau 9.1

Problem in Tableau 9.1						
Basic Variables	Inverse Tableau					Basic Values
$x_1$	-1	-2	2	0	0	3
$x_2$	1	1	-1	0	4	2
$x_3$	-1	0	1	0	1	-11
$-z$	-2	4	-1	1	1	

### 9.2 INTRODUCING A NEW ACTIVITY

Suppose a new variable called  $x_{n+1}$  has to be introduced into the model (9.1). Let  $A_{n+1}$ ,  $c_{n+1}$  be the



input-output coefficient vector and the cost coefficient, respectively, of this new variable. Let  $X^1 = (x^1, x_{n+1})$ . The augmented problem is

$$\begin{cases} \text{Minimize } Z(X) = cx + c_{n+1}x_{n+1} \\ \text{Subject to } (A : A_{n+1})X = b \\ X \geq 0 \end{cases} \quad (9.2)$$

$\bar{B}$  remains a feasible basis to the augmented problem, and the BFS of (9.2) corresponding to it is  $\bar{X}^1 = (\bar{X}^1, 0)$ . From the optimality criterion,  $\bar{X}$  remains an optimum feasible solution of (9.2) if the relative cost coefficient of the new variable  $x_{n+1}$ , with respect to the basis  $\bar{B}$ , is nonnegative; that is, if  $\bar{c}_{n+1} = c_{n+1} - \bar{\pi}A_{n+1} \geq 0$ . On the other hand, if  $\bar{c}_{n+1} < 0$ , solve (9.2) by using the inverse tableau corresponding to the basis  $\bar{B}$  as an initial tableau. Bring  $x_{n+1}$  into the basic vector and complete the solution of (9.2) according to the revised simplex algorithm.

It often happens that one pivot (that of bringing the new variable  $x_{n+1}$  into the basic vector  $x_B$ ) is all that is necessary to solve (9.2). However, there is no theoretical guarantee that this will be the general case. In some problems, it may be necessary to perform several pivots before the algorithm reaches a terminal basis for (9.2).

#### Example 9.1

Consider the LP in Tableau 9.1. Suppose a new variable  $x_7$  corresponding to the column vector  $A_7 = (1, 2, -3)^T$  and cost coefficient  $c_7 = -7$  is introduced. The relative cost coefficient of  $x_7$  with respect to the basis in Tableau 9.2 is  $\bar{c}_7 = c_7 + (-\bar{\pi})A_7 = -7 + (-2, 4, -1)(1, 2, -3)^T = 2$ . Hence, the previous optimum solution with  $x_7 = 0$  is still optimal. The optimum feasible solution of the new problem is  $\bar{X} = (3, 4, 2, 0, 0, 0)^T$ , with the optimal objective value  $\bar{Z} = 11$ . Thus, if  $x_7$  is the level of some new activity that has become available, we conclude that it is optimal to perform the new activity at zero level.

#### Example 9.2

Consider the LP in Tableau 9.1 again. Suppose we have to include a new variable  $x_7$  with  $A_7 = (3, -1, 1)^T$  and  $c_7 = 4$ . The relative cost coefficient of  $x_7$  with respect to the basis in Tableau 9.2 is  $\bar{c}_7 = 4 + (-2, 4, -1)(3, -1, 1)^T = -7 < 0$ . We conclude that it is optimal to include the new activity

in the basic vector. The updated column vector of  $x_7$  with respect to the basis in Tableau 9.2 is

$$\begin{pmatrix} \bar{A}_7 \\ \bar{c}_7 \end{pmatrix} = \begin{pmatrix} -1 & -2 & 2 & 0 \\ 1 & 1 & -1 & 0 \\ -1 & 0 & 1 & 0 \\ -2 & 4 & -1 & 1 \end{pmatrix} \begin{pmatrix} 3 \\ 4 \\ 2 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ 1 \\ 4 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ -2 \\ -7 \end{pmatrix}$$

This is the pivot column. The minimum ratio test indicates that  $x_1$  drops out of the basic vector when  $x_7$  enters.

Basic Variables	Inverse Tableau	Basic Values
$x_7$	-1   -2   2   0	3
$x_2$	1   1   -1   0	1
$x_3$	-1   0   1   0	8
$-z$	-2   4   -1   1	10

Since the relative cost coefficient of  $x_6$  here is  $-6$ , this basis is again not optimal to the augmented problem. Bring  $x_6$  into the basic vector and continue the algorithm until a terminal basis for the augmented problem is obtained.

#### Exercises

9.1 Complete the solution of this numerical problem.

9.2 Prove that the optimum objective value in (9.2) is less than or equal to the optimum objective value in (9.1). Construct a numerical example to show that the objective function in (9.2) may be unbounded below even though (9.1) has an optimum feasible solution.

9.3 How is the set of feasible solutions of the augmented problem related to  $K$ ?

#### 9.2.1 What Useful Planning Information Can Be Derived from This Type of Sensitivity Analysis?

The sensitivity analysis discussed above is very simple, and yet when a company has a linear programming model of its production operations, it can yield extremely useful planning information. For an example, consider the fertilizer problem (1.1) modeled in Section 1.1.3. From the discussion in Section 4.6.3, we know that the optimum dual solution associated with (1.1) is  $\bar{\pi} = (5, 5, 0)$ . Suppose a research chemist working for this manufacturer has



come up with the formula for a new fertilizer called Lushlawn. The manufacture of Lushlawn requires as inputs  $A_n = (3, 2, 2)^T$  tons of raw materials 1, 2, and 3, respectively, per ton. Should the manufacturer produce this new fertilizer? How much profit (in dollars per ton) should 1 ton of Lushlawn fetch in the marketplace for the manufacturer to consider it worth producing? To answer these questions (1.1) can be transformed into standard form, and the analysis discussed above is applied. It can be verified that the breakeven profit for Lushlawn is  $\bar{\pi}A_n = (5, 5, 0)(3, 2, 2)^T = 25$ /ton. That is, Lushlawn is worth producing if it can be sold in the market at a price that leads to a profit  $\geq 25$ /ton manufactured. The fertilizer manufacturer can conduct a survey and determine whether the market would accept Lushlawn at a price greater than or equal to this breakeven level. Based on the results of the survey, the decision whether to produce Lushlawn or not can be taken very easily. Also, once the profit level for Lushlawn is set, the new optimum solution for the manufacturer can be determined easily, as discussed above, beginning with an optimum feasible basic vector for the standard form of (1.1).

Many companies use this type of sensitivity analysis to determine whether new products or processes would turn out to be profitable, to set prices on new products, and to estimate how sales volumes of existing products will be affected by the introduction of new products.

### 9.3 INTRODUCING AN ADDITIONAL INEQUALITY CONSTRAINT

Consider the LP (9.1) and the optimum basis  $\bar{B}$  for it again. Suppose the additional constraint

$$A_{m+1}x \leq b_{m+1} \quad (9.3)$$

has to be introduced. Let  $K_1$  denote the set of feasible solutions of the augmented problem.  $K_1 \subset K$  (see Figure 9.1). If the original problem has an optimum feasible solution, the augmented problem is either infeasible or it has an optimum feasible solution. Also  $z(\bar{x}) = \text{minimum } \{z(x) : x \in K\} \leq \text{minimum } \{z(x) : x \in K_1\}$ . Hence, if  $\bar{x} \in K_1$ , that is,  $\bar{x}$  satisfies (9.3),  $\bar{x}$  remains optimal to the augmented problem. On the other hand, if  $\bar{x}$  does not satisfy (9.3), then

$$\bar{x}_{n+1} = -A_{m+1}\bar{x} + b_{m+1} < 0 \quad (9.4)$$

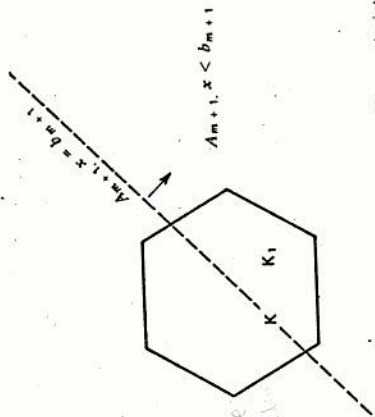


Figure 9.1 Introducing a new inequality constraint.

where  $x_{n+1}$  is the slack variable corresponding to (9.3). The augmented problem is

$$\text{Minimize } Z(x) = \sum_{j=1}^n c_j x_j + 0x_{n+1}$$

$$\text{Subject to } \begin{pmatrix} A & 0 \\ A_{m+1} & 1 \end{pmatrix} \begin{pmatrix} x \\ x_{n+1} \end{pmatrix} = \begin{pmatrix} b \\ b_{m+1} \end{pmatrix} \quad (9.5)$$

$$x_j \geq 0 \quad \text{for all } j$$

Let  $X^T = (x^T, x_{n+1})$ . By including  $x_{n+1}$  as an additional basic variable, we can enlarge  $\bar{B}$  into a basis for (9.5), denoted by  $\bar{B}$ .  $X_B = (x_1, \dots, x_m, x_{n+1})^T$ . The basic solution of (9.5) corresponding to the basis  $\bar{B}$  is  $\bar{x}^T = (\bar{x}^T, \bar{x}_{n+1})$ . By (9.4),  $\bar{B}$  is an infeasible basis for (9.5). Since  $c_{n+1} = 0$ , the relative cost coefficient of  $x_j$  with respect to the basis  $\bar{B}$  in (9.5) is equal to the relative cost coefficient of  $x_j$  with respect to the basis  $\bar{B}$  in (9.1), which is  $\bar{c}_j \geq 0$  for all  $j = 1$  to  $n$ . Thus,  $\bar{B}$  is a dual feasible but primal infeasible basis for (9.5). Using  $\bar{B}$  as an initial basis, we can solve (9.5) by using the dual simplex algorithm. We now discuss how to obtain the inverse tableau corresponding to the basis  $\bar{B}$  for (9.5).

$$\bar{B} = \begin{pmatrix} a_{11} & \dots & a_{1m} \\ \vdots & \ddots & \vdots \\ a_{m1} & \dots & a_{mm} \\ a_{m+1,1} & \dots & a_{m+1,m} \end{pmatrix} \quad \bar{B}^{-1} = \begin{pmatrix} a_{11} & \dots & a_{1m} & 0 \\ \vdots & \ddots & \vdots & \vdots \\ a_{m1} & \dots & a_{mm} & 0 \\ a_{m+1,1} & \dots & a_{m+1,m} & 1 \end{pmatrix}$$

Therefore

$$\bar{B}^{-1} = \begin{pmatrix} \bar{B}^{-1} & 0 \\ \dots & \dots \\ -(a_{m+1,1}, \dots, a_{m+1,m})\bar{B}^{-1} & 1 \end{pmatrix}$$

Basic Variables	Inverse Tableau					Basic Values	Pivot Column ( $x_4$ )
$x_1$	-1	-2	2	0	0	3	-1
$x_2$	1	1	-1	0	0	4	1
$x_3$	-1	0	1	0	0	2	2
$x_7$	5	3	-6	1	0	-12	(-4)
$-z$	-2	4	-1	0	1	-11	1

## Exercises

9.4 Complete the solution of this augmented problem.

9.5 Show that the type of sensitivity analysis discussed here (introducing an additional inequality constraint) is the dual of the one discussed in Section 9.2 (introducing a new sign restricted variable).

9.6 Prove that the set of extreme points of  $K_1$  consists of:

- All extreme points of  $K$  that satisfy (9.3);
- Points of intersection of edges of  $K$  that do not completely lie on the hyperplane  $H_1$  with  $H_1$ , where  $H = \{x: A_{m+1}x = b_{m+1}\}$ .

—K. G. Murty [3.35]

9.7 If every optimum feasible solution of (9.1) violates (9.3), and if  $K_1 \neq \emptyset$ , prove that every optimum feasible solution of the augmented problem satisfies (9.3) as an equation, that is, lies on the hyperplane  $H$  defined in Exercise 9.6.

## 9.4 INTRODUCING AN ADDITIONAL EQUALITY CONSTRAINT

Returning to the LP (9.1), suppose an additional equality constraint has to be introduced:  $A_{m+1}x = b_{m+1}$ . The problem (9.1) together with this additional equality constraint is called the *augmented problem*.

Thus  $\bar{B}^{-1}$  can easily be obtained from  $\bar{B}^{-1}$ . Also obviously  $\bar{\pi} =$  dual solution of (9.5) corresponding to  $\bar{B}$  is equal to  $(\bar{\pi}, 0)$ .

Thus the inverse tableau for (9.5) corresponding to the basis  $\bar{B}$  is easily obtained from the inverse tableau for (9.1) corresponding to the basis  $\bar{B}$ . On the other hand, if (9.1) was solved by computing the canonical tableaux after each pivot step, to obtain the canonical tableau of (9.5) with respect to the basis  $\bar{B}$ , introduce the  $(m+1)$ th constraint row at the bottom of the canonical tableau of (9.1) with respect to the basis  $\bar{B}$ , include  $x_{m+1}$  as the basic variable in that row, and then price out all the basic column vectors of  $\bar{B}$  in it.

## Example 9.3

Returning to the LP in Tableau 9.1, suppose the new constraint,  $x_1 - x_2 + 3x_3 \leq -7$ , has to be imposed. The optimum solution  $\bar{x} = (3, 4, 2, 0, 0)^T$  violates this new constraint. Let  $x_7$  be the slack variable.  $x_7 = -x_1 + x_2 - 3x_3 - 7$ . The inverse tableau for the augmented problem corresponding to the basic vector  $(x_1, x_2, x_3, x_7)$  is given at the top.  $x_7$  is the only negative-valued basic variable here. The updated row vector in which  $x_7$  is the basic variable is given at the bottom. The pivot column is the updated column vector of  $x_4$ , which is entered as the last column in the inverse tableau at the top. The pivot is performed and the dual simplex algorithm applied in a similar manner until termination.

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$-z$	$b$
Updated 4th Row	0	0	0	-4	0	-3	1	0	-12
Updated Cost Row	0	0	0	1	3	8	0	1	-11
$\frac{-c_j}{\bar{a}_{4j}}$ for $\bar{a}_{4j} < 0$				$\frac{1}{4}$		$\frac{8}{3}$			

$$\pi = C\bar{B}^{-1}$$

$$\pi = (3 \ 2 \ -3 \ 0) \begin{pmatrix} -1 & -2 & 2 & 0 \\ 1 & 1 & -1 & 0 \\ -1 & 0 & 1 & 0 \\ 5 & 3 & -1 & 0 \end{pmatrix} = (-4 \ 1 \ 0)$$



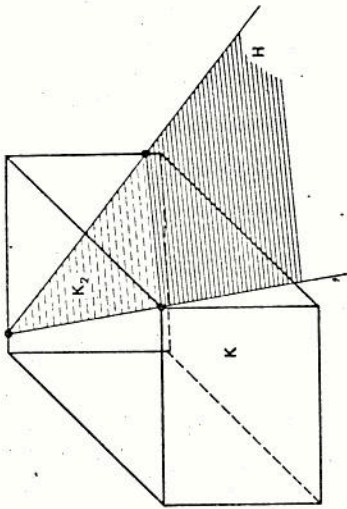


Figure 9.2 Introducing a new equality constraint.

Let  $H = \{x: A_{m+1,1}x = b_{m+1,1}\}$ . The set of feasible solutions of the augmented problem is  $K_2 = K \cap H$ . See Figure 9.2 for an example. If the optimum solution of (9.1),  $\bar{x} \in H$ , obviously  $\bar{x}$  is optimal to the augmented problem. Suppose  $\bar{x} \notin H$ , then  $A_{m+1,1}\bar{x} \neq b_{m+1,1}$ . Suppose  $A_{m+1,1}\bar{x} > b_{m+1,1}$ , then change the original problem by adding the constraints  $A_{m+1,1}x = x_{n+1} = b_{m+1,1}$ ,  $x_{n+1} \geq 0$ . (If, on the other hand, it turned out that  $A_{m+1,1}\bar{x} < b_{m+1,1}$ , then the coefficient of  $x_{n+1}$  should be  $+1$  instead.) Here is the new problem:

$$\text{Minimize } Z(x) = cx + Mx_{n+1}$$

$$\text{Subject to } \begin{pmatrix} A_{m+1,1} & 0 \\ \dots & \dots \end{pmatrix} \begin{pmatrix} x \\ \dots \end{pmatrix} = \begin{pmatrix} b_{m+1,1} \\ \dots \end{pmatrix}$$

$$x = \begin{pmatrix} x \\ \dots \end{pmatrix} \geq 0 \quad (9.6)$$

where  $M$  is an arbitrarily large positive number. We will refer to (9.6) as the new problem. The variable  $x_{n+1}$  is an artificial variable in the new problem, and hence it is included in the objective function in (9.6) with a coefficient of  $M$ , as in the Big- $M$  method. Clearly,  $\bar{x}_B = (x_1, \dots, x_m, x_{n+1})^T$  is a basic vector for the new problem, and the basic solution of the new problem corresponding to this basic vector can be verified to be  $\bar{x} = (\bar{x}_1, \dots, \bar{x}_m, \bar{x}_{n+1})^T$ , where  $\bar{x} = (\bar{x}_1, \dots, \bar{x}_n)$  is the BFS of the original problem (9.1) corresponding to the optimal basic vector  $x_B$  for it and  $\bar{x}_{n+1} = A_{m+1,1}\bar{x} - b_{m+1,1}$ . Since  $\bar{x}_{n+1} > 0$  by our

assumptions,  $\bar{x} \geq 0$ , and hence  $\bar{x}_B$  is a feasible basic vector for (9.6). The associated basis for (9.6),  $\bar{B}$ , is

$$\bar{B} = \begin{bmatrix} a_{11} & \dots & a_{1m} & 0 \\ \vdots & \vdots & \vdots & \vdots \\ a_{m1} & \dots & a_{mm} & 0 \\ a_{m+1,1} & \dots & a_{m+1,m} & -1 \end{bmatrix}$$

$$\bar{B}^{-1} = \begin{pmatrix} \dots & \dots & \dots & 0 \\ \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots \\ (a_{m+1,1}, \dots, a_{m+1,m})\bar{B}^{-1} & -1 & \dots & \dots \end{pmatrix}$$

Define  $c_{n+1} = 0$ ,  $c_j^* = 0$ , for  $j = 1$  to  $n$ , and  $c_{n+1}^* = 1$ . Then the objective function in (9.6) is  $Z(x) = \sum_{j=1}^{n+1} (c_j + Mc_j^*)x_j$ . In the original tableau for (9.6) the cost coefficients can be entered in two separate rows, as under the Big- $M$  method discussed in Section 2.7. One row contains the  $c_j$ 's, the cost coefficients in  $z(x)$  from the original objective function. The second row contains  $c_j^*$ 's, the coefficients of  $M$  in  $Z(x)$ . Compute  $\Pi_B = (\Pi_1, \dots, \Pi_{m+1}) = c_B\bar{B}^{-1}$ ,  $\Pi_B^* = (\Pi_1^*, \dots, \Pi_{m+1}^*) = c_B^*\bar{B}^{-1}$ . The inverse tableau for (9.6), corresponding to the basic vector  $\bar{x}_B$ , is

Basic Variables	Inverse Tableau	Basic Values
$x_1$	$\bar{B}^{-1}$	$b_1$
$\vdots$	$\vdots$	$\vdots$
$x_{n+1}$	$\bar{B}^{-1}$	$b_{m+1}$
$-Z$	$-\Pi_1, \dots, -\Pi_{m+1}$ $-\Pi_1^*, \dots, -\Pi_{m+1}^*$	$-Z$ $-Z^*$

Letting  $\bar{c}_j = c_j - \Pi_B A_j$ ,  $\bar{c}_j^* = c_j^* - \Pi_B^* A_j$ , the actual relative cost coefficient of  $x_j$  in (9.6) with respect to the current basic vector is  $\bar{c}_j + M\bar{c}_j^*$ . If  $\bar{c}_j + M\bar{c}_j^* \geq 0$  for each  $j$ , the algorithm terminates. If these termination conditions are not satisfied, let  $J = \{j: \bar{c}_j + M\bar{c}_j^* < 0, \text{ or } \bar{c}_j^* = 0 \text{ and } \bar{c}_j < 0\}$ . The eligible variables at this stage are those  $x_j$  for  $j \in J$ . Select one of the eligible variables as the entering variable, bring it into the basic vector, and continue in this manner until a basic vector satisfying the above termination conditions is obtained. Suppose  $\bar{x} = (\bar{x}_1, \dots, \bar{x}_m, \bar{x}_{n+1})^T$  is the final optimum solution obtained for (9.6). If  $\bar{x}_{n+1} > 0$ ,  $K_2 = \emptyset$ , in this case the additional equality constraint has made the original problem infeasible. On the other hand, if  $\bar{x}_{n+1} = 0$ ,  $\bar{x} = (\bar{x}_1, \dots, \bar{x}_n)^T$  is an optimum feasible solution of the augmented problem.

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Basic Variables	Inverse Tableau						Basic Values	Pivot Column $x_6$	Ratios
$x_1$	-1	-2	2	0	0	0	3	-2	
$x_2$	1	1	-1	0	0	0	4	-2	
$x_3$	-1	0	1	0	0	0	2	1	
$x_7$	-3	-4	4	-1	0	0	4	4	2
$-Z$	-2	4	-1	0	1	0	-11	8	4/2 minimum
	3	4	-4	1	0	1	-4	-2	

#### Example 9.4

Consider the LP in Tableau 9.1. Suppose the additional constraint  $x_1 - 2x_2 + x_3 + 3x_5 = -9$  is to be included.  $\bar{x} = (3, 4, 2, 0, 0)^T$ , the optimum solution of the original problem, violates this new constraint. Let  $\bar{x}_7 = \bar{x}_1 - 2\bar{x}_2 + \bar{x}_3 + 3\bar{x}_5 - (-9) = 4$ . The new problem as discussed above, is:

#### New Problem

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$-Z$	$b$
	1	2	0	1	0	-6	0	0	11
	0	1	1	3	-2	-1	0	0	6
	1	2	1	3	-1	-5	0	0	.13
	1	-2	0	1	3	0	-1	0	-9
$c$	3	2	-3	-6	10	-5	0	1	0
$c^*$	0	0	0	0	0	0	0	0	0

$x_j \geq 0$  for all  $j$ , minimize  $Z$ ,  $x_7$  is the artificial variable.

Let  $\bar{B}$  be the basis for the new problem corresponding to the basic vector  $(x_1, x_2, x_3, x_4, x_7)$ . The inverse tableau for the new problem with respect to the basis  $\bar{B}$  is given at the top.

The updated cost rows are  $\bar{c} = (0, 0, 0, 1, 3, 8, 0)$  and  $\bar{c}^* = (0, 0, 0, 4, -1, -2, 0)$ . So  $x_5$  and  $x_6$  are eligible to enter the basic vector. Suppose we choose  $x_6$  as the entering variable. Continuing the application of the Big-M method, it can be verified that the optimum solution of the new problem is  $X^1 = (0, 8, 7/2, 0, 7/3, 5/6, 0)^T$ . Since  $x_7 = 0$  in this solution,  $x = (0, 8, 7/2, 0, 7/3, 5/6)^T$  is an optimum solution of the augmented problem.

#### Exercises

- 9.8 Consider the LP in Tableau 9.1. Suppose the additional constraint  $-x_1 - x_2 - x_3 + x_4 = 300$  has to be introduced. Apply the algorithm discussed here to solve the augmented problem.

lem. Show that the augmented problem is infeasible.

- 9.9 Returning to the LP in Tableau 9.1, suppose the additional constraint  $x_1 + x_2 + x_3 + x_4 + x_5 + x_6 = 12$  has to be imposed. Obtain an optimum feasible solution of the augmented problem.

- 9.10 Discuss an approach for solving the augmented problem with the additional equality constraint, using the dual simplex algorithm instead of the Big-M approach described here.

- 9.11 Prove that the set of extreme points of  $K_2$  consists of (a) all extreme points of  $K$  that are in  $H$ , and (b) points of intersection of edges of  $K$  that do not completely lie in  $H$  with  $H$ .

- 9.12 Let  $H, K_2$  be defined as in the beginning of this section. Let  $K_1 = K \cap \{x: A_{m+1}x \leq b_{m+1}\}$ .  $K_3 = K \cap \{x: A_{m+1}x \geq b_{m+1}\}$ . If  $K_2 \neq \emptyset$ , prove that  $K_2$  contains an optimum solution of at least one of the following two problems.

(a) Minimize  $z(x) = cx$

$x \in K_1$

(b) Minimize  $z(x) = cx$

$x \in K_3$

#### 9.5 COST RANGING OF A NONBASIC COST COEFFICIENT

Consider again the LP (9.1). Suppose  $x_i$  is a variable that is not in the optimal basic vector  $x_B$ . Assuming that all the other cost coefficients except  $c_i$  remain fixed at their specified values, determine the range of values of  $c_i$  within which  $\bar{B}$  remains an optimal basis. For this, treat  $c_i$  as a parameter.  $\bar{B}$  remains an optimal basis as long as  $\bar{c}_i = c_i - \bar{\pi}A_i \geq 0$ , that is



$c_i \geq \pi A_{ij}$ . The relative cost coefficients of all the other variables are independent of the value of  $c_i$ , and, hence, they remain nonnegative. If the new value of  $c_i$  is not in the closed interval  $[\pi A_{ij}, \infty]$ ,  $x_i$  is the only nonbasic variable that has a negative relative cost coefficient with respect to the basis  $\bar{B}$  in the modified problem. Bring  $x_i$  into the basic vector and continue the application of the simplex algorithm until a terminal basis for the modified problem is obtained.

#### Example 9.5.

Consider the LP in Tableau 9.1. The range of values of  $c_5$  (whose present value is 10) for which the basic vector  $x_B = (x_1, x_2, x_3)$  remains optimal is determined by  $c_5 + (-2, 4, -1)(0, -2, -1)^T \geq 0$ , that is,  $c_5 \geq 7$ . Suppose the value of  $c_5$  has to be modified to 6. When  $c_5$  is changed from 10 to 6, the updated column vector of  $x_5$  becomes  $(2, -1, -1)^T$ . With this as the pivot column, bring  $x_5$  into the basic vector. This leads to the following tableau.

Basic Variables	Inverse Tableau				Basic Values
$x_5$	$\frac{1}{-2}$	-1	1	0	3
	$-\frac{1}{2}$				2
$x_2$	$\frac{1}{2}$	0	0	0	$\frac{11}{2}$
					2
$x_3$	$\frac{3}{-2}$	-1	2	0	7
	$-\frac{1}{2}$				2
$-z$	$\frac{5}{-2}$	3	0	1	19
	$-\frac{1}{2}$				-2

It can be verified that this tableau displays an optimum tableau for the modified problem.

### 9.6 COST RANGING OF A BASIC COST COEFFICIENT

Consider the LP (9.1). Suppose all the cost coefficients except  $c_i$ , which is a basic cost coefficient, remain fixed at their present values. Treating  $c_i$  as

a parameter, determine the range of values of  $c_i$  for which  $\bar{B}$  remains an optimal basis. Since  $c_i$  is a basic cost coefficient, any change in  $c_i$  changes the dual solution corresponding to the basis  $\bar{B}$ , and all the relative cost coefficients. For simplicity let  $\gamma_1$  denote the parameter and  $c_i$  its present value. Let  $\pi(\gamma_1)$ ,  $\bar{c}_j(\gamma_1)$ , etc., denote the dual solution and the relative cost coefficient of  $x_j$  corresponding to the basis  $\bar{B}$  as functions of the parameter  $\gamma_1$ . Then  $\pi(\gamma_1) = (\gamma_1, c_2, \dots, c_m)\bar{B}^{-1}$ . Hence, the dual vector is an affine function of  $\gamma_1$ . If we use this dual solution, the relative cost coefficient of  $x_j$  is  $\bar{c}_j(\gamma_1) = c_j - \pi(\gamma_1)A_{ij}$ , for  $j = 1$  to  $n$ . Hence, each  $\bar{c}_j(\gamma_1)$  is itself an affine function of  $\gamma_1$ . The range of values of  $\gamma_1$  for which  $\bar{B}$  remains an optimum basis is the range of  $\gamma_1$  within which every  $\bar{c}_j(\gamma_1)$  remains nonnegative. This range will turn out to be a nonempty closed interval. If it is necessary to modify the cost coefficient of  $x_i$  to some value  $c'_i$  outside this closed interval, fix  $\gamma_1 = c'_i$  and compute the relative cost coefficients  $\bar{c}_j(c'_i)$  using the formulas already developed. Bring one of the variables for which  $\bar{c}_j(c'_i) < 0$  into the basic vector and continue the applications of the simplex algorithm until a terminal basis for the modified problem is obtained.

#### Example 9.6

In the LP in Tableau 9.1, let  $\gamma_1$  represent the cost coefficient of  $x_1$  whose present value is 3. The dual solution corresponding to the basic vector  $(x_1, x_2, x_3)$  as a function of  $\gamma_1$  is  $\pi(\gamma_1) = (\gamma_1, 2, -3)\bar{B}^{-1} = (-\gamma_1 + 5, -2\gamma_1 + 2, 2\gamma_1 - 5)$ . Using  $\bar{c}_j(\gamma_1) = c_j - \pi(\gamma_1)A_{ij}$ , the row of relative cost coefficients is  $\bar{c}(\gamma_1) = (0, 0, 0, \gamma_1 - 2, -2\gamma_1 + 9, 2\gamma_1 + 2) \geq 0$  iff  $2 \leq \gamma_1 \leq 9/2$ . So  $\bar{B}$  is an optimum basis whenever the cost coefficient of  $x_1$  is in the interval  $[2, 9/2]$ , assuming that all the other cost coefficients remain at their present values.

Suppose  $\gamma_1$  has to be changed to 5. Changing the cost coefficient of  $x_1$  from the present value of 3 to 5, the modified dual solution is  $\pi(\gamma_1 = 5) = (0, -8, 5)$ . With this change, the inverse tableau becomes

Basic Variables	Modified Inverse Tableau for Basis $\bar{B}$				Basic Values	Pivot Column (x <sub>1</sub> )
$x_1$	-1	-2	2	0	3	② Pivot Row
$x_2$	1	1	-1	0	4	
$x_3$	-1	0	1	0	2	
$-z$	0	8	-5	1	-17	

The modified relative cost coefficient of  $x_5$  is  $-1$ . Hence, bring  $x_5$  into the basic vector. The updated column vector of  $x_5$  is entered as the pivot column in the inverse tableau. Remember that the cost coefficient of  $x_1$  is now 5, and continue the algorithm.

### Exercises

9.13 Construct a counter example to the following: "If  $B$  is an optimum basis for (9.1), then it remains an optimum basis when some of the entries in  $c_B$  are decreased."

9.14 Is the following true? " $x_1$  is a basic variable in an optimum basic vector for (9.1). If  $c_1$  is decreased and if the modified problem has an optimum feasible solution, then it has an optimum basic vector containing  $x_1$ ."

### 9.7 RIGHT-HAND-SIDE RANGING

In the LP (9.1) suppose we wish to determine the range of values of one of the right-hand-side constants, say  $b_1$ , for which the basis  $\bar{B}$  remains optimal. Treat this right-hand-side constant as a parameter and denote it by  $\beta_1$ ;  $b_1$  is the present value of  $\beta_1$ . We assume that all the other right-hand-side constants stay fixed at their present values.  $\bar{B}$  is an optimum basis when  $\beta_1 = b_1$ . Hence,  $\bar{B}$  is dual feasible. So  $\bar{B}$  is optimal for all values of  $\beta_1$  for which it is primal feasible. The values of the basic variables  $x_B$  are all functions of the parameter  $\beta_1$ , and  $\bar{B}$  is primal feasible as long as all these values are nonnegative.  $x_B(\beta_1) = \bar{B}^{-1}(\beta_1, b_2, \dots, b_m)^T \geq 0$ . Since all these inequalities are linear in  $\beta_1$ , this will determine a closed interval of the form  $\underline{\lambda} \leq \beta_1 \leq \bar{\lambda}$ . For all values of  $\beta_1$  in this closed interval,  $\bar{B}$  remains an optimal basis. If it is necessary to modify the value of  $\beta_1$  from its present value  $b_1$  to a value  $b_1'$  outside the closed interval  $[\underline{\lambda}, \bar{\lambda}]$ , fix  $\beta_1$  at  $b_1'$  and obtain the modified values of the basic variables.  $\bar{B}$  is still dual feasible, but primal infeasible. Starting with  $\bar{B}$ , apply the dual simplex routine until a new terminal basis is obtained.

#### Example 9.7

In the LP in Tableau 9.1 the values of the basic variables in the basic vector  $(x_1, x_2, x_3)$  as functions of

$\beta_1$  are

$$\bar{B}^{-1} \begin{pmatrix} \beta_1 \\ 6 \\ 13 \end{pmatrix} = \begin{pmatrix} -1 & -2 & 2 \\ 1 & 1 & -1 \\ -1 & 0 & 1 \end{pmatrix} \begin{pmatrix} \beta_1 + 14 \\ \beta_1 - 7 \\ -\beta_1 + 13 \end{pmatrix}$$

All these are nonnegative iff  $7 \leq \beta_1 \leq 13$ . Hence, this is the interval within which the basic vector  $(x_1, x_2, x_3)$  remains optimal. In this range, an optimum solution of the problem is  $x = (14 - \beta_1, -7 + \beta_1, 13 - \beta_1, 0, 0, 0)^T$ , with an optimum objective value of  $-11 + 2\beta_1$  and an associated optimum dual solution of  $\pi = (2, -4, 1)$ .

Now, suppose it is required to solve the problem for  $\beta_1 = 15$ . The inverse tableau for the modified problem corresponding to the present basic vector  $(x_1, x_2, x_3)$  is:

Basic Variables	Inverse Tableau			Basic Values	Pivot Column
$x_1$	-1	-2	2	0	1
$x_2$	1	1	-1	0	8
$x_3$	-1	0	1	0	-2
$-z$	-2	4	-1	1	-19
					3

Applying the dual simplex algorithm starting with this inverse tableau, it can be verified that the modified problem is infeasible.

### Exercises

9.15 Obtain the new optimum feasible solution of the LP in Tableau 9.1 after  $b_1$  is changed from its present value of 11 to 6.

### 9.8 CHANGES IN THE INPUT-OUTPUT COEFFICIENTS IN A NONBASIC COLUMN VECTOR

Let  $x_j$  be a variable that is not in the optimum basic vector of  $x_B$  of (9.1). And  $a_{ij}$  is one of the input-output coefficients in the column vector of  $x_j$ . If all the other coefficients in the problem except  $a_{ij}$  remain fixed at their present values, what is the range of values of  $a_{ij}$  within which  $\bar{B}$  remains an optimal basis?

Treat  $a_{ij}$  as a parameter. To avoid confusion, call this parameter,  $\alpha_{ij}$ , and let  $a_{ij}$  be its present value. Since  $x_j$  is a nonbasic variable, a change in  $\alpha_{ij}$  does not affect the primal feasibility of  $\bar{B}$ . A change in  $\alpha_{ij}$  can only change the relative cost coefficient of  $x_j$



and this is obtained by  $\bar{c}_j(x_{ij}) = c_j + (\sum_{i=1}^m x_{ij}(-\bar{a}_{ij})a_{ij}) - \bar{a}_{ij}x_{ij}$ . As long as  $\bar{c}_j(x_{ij}) \geq 0$ ,  $\bar{B}$  remains an optimal basis. This determines a closed interval for  $x_{ij}$ , and as long as  $x_{ij}$  is in this interval  $\bar{B}$  remains an optimal basis. If  $x_{ij}$  has to be changed to a value  $a'_{ij}$  outside this interval, make the change, bring  $x_{ij}$  into the basic vector, and continue with the application of the simplex algorithm until a new terminal basis is obtained.

#### Example 9.8

In the LP in Tableau 9.1, the present value of  $x_{25}$  in the column vector of  $x_5$  is  $-2$ . The relative cost coefficient  $\bar{c}_5(x_{25}) = 10 + (-2, 4, -1)(0, x_{25} - 1)^T = 11 + 4x_{25} \geq 0$  iff  $x_{25} \geq -11/4$ . Thus  $\bar{B}$  remains an optimal basis for  $x_{25} \geq -11/4$ . Suppose we change  $x_{25}$  from its present value of  $-2$  to  $-3$ . The updated column vector of  $x_5$  changes to  $(4, -2, -1, -1)^T$ . With this pivot column bring  $x_5$  into the basic vector and continue until termination is reached.

### 9.9 CHANGES IN A BASIC INPUT-OUTPUT COEFFICIENT

Let  $x_1$  be a basic variable in the optimum basic vector  $x_B$  for (9.1). Suppose we have to modify one input-output coefficient, say  $a_{11}$ , in the column vector of  $x_1$  to  $a'_{11}$ . The modified column vector of  $x_1$  will be  $A'_1 = (a'_{11}, a_{21}, \dots, a_{m1})^T$ . Let  $x'_1$  indicate the level of this activity corresponding to the new column vector  $A'_1$ . The previous column vector  $A_1$  is no longer a part of the problem and it should be eliminated. Physically  $x'_1$  replaces  $x_1$  in the original tableau. We will refer to the altered problem by the name *modified problem*.

Construct a new problem by augmenting the present original tableau with the new variable  $x'_1$  with its column vector  $A'_1$  and cost coefficient  $c_1$ , and changing the cost coefficient of  $x_1$  to  $M$ , where  $M$  is a very large positive number. This leads to the problem, which we call the *new problem*. In this problem  $x_1$  plays the role of an artificial variable, associated with a cost coefficient of  $M$ , a very large positive number. The objective function  $Z$ , in the new problem is  $\sum_{j=2}^n c_j x_j + c_1 x'_1 + M x_1$ . As in the Big- $M$  method, the cost coefficients are entered in two rows.  $\bar{B}$  is still a feasible basis to this new problem. However, in the inverse tableau corresponding to the basis  $\bar{B}$ , the dual vector has to be recomputed, since the cost

coefficient of  $x_1$  has been changed to  $M$ . The new dual solution corresponding to  $\bar{B}$  is  $\pi(M) = (M, c_2, \dots, c_n)\bar{B}^{-1} = \pi + M\pi^*$ , where  $\pi = (0, c_2, \dots, c_n)\bar{B}^{-1}$ ,  $\pi^* = (1, 0, \dots, 0)\bar{B}^{-1}$ . Construct the inverse tableau for the new problem, corresponding to the basic vector  $x_B$  and solve it as discussed in Section 9.4, and the normal termination conclusions of the big- $M$  method apply.

### 9.10 PRACTICAL APPLICATIONS OF SENSITIVITY ANALYSIS

When studying a system using a linear programming model, techniques of sensitivity analysis can be used to evaluate new products (as in Section 9.2), to determine the breakeven selling price of a new product, at which point it becomes competitive with the existing list of products in terms of profitability (as in Sections 9.2 and 9.5); to evaluate new technologies or processes for making products (as in Sections 9.8 and 9.9); to assess how profitable it is to acquire additional resources and to determine which resources to acquire in what quantities (using the ideas discussed in Sections 4.6.3, 9.7, and Chapter 8); to evaluate the effects of changes in the costs (as in Sections 9.5, 9.6, and Chapter 8), and to determine optimal policies for handling new constraints that might arise. When used in this manner, the linear programming model not only determines an optimum solution to implement but becomes a tool for optimal planning. Examples of some of these uses are discussed in the problems that follow.

Optimization models other than linear programming models (e.g., integer programming models and nonlinear programming models) do not lend themselves that readily to a marginal analysis or sensitivity analysis as the linear programming model. That's why if a practical problem can be approximated reasonably closely by a linear programming model, it is so much easier to study it than otherwise.

#### Exercises

9.16 Consider the family's diet problem discussed in Chapter 4. The original tableau for it is given at the top of page 319. Here  $x_1$  to  $x_6$  are the kilograms of the primary foods 1 to 6 in the family's diet, and  $x_7$ ,  $x_8$  are the slack variables representing the excess amounts of the nutrients, vitamins A and C, in the diet over the minimum requirements. The basis  $B_1$

Original Tableau for the Family's Diet Problem

$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$	$-z$	$b$
1	0	2	2	1	2	-1	0	0	9
0	1	3	1	3	2	0	-1	0	19
35	30	60	50	27	22	0	0	1	0 (minimize $z$ )

$x_j \geq 0$  for all  $j$ .

associated with the basic vector  $(x_5, x_6)$  is optimal to this problem. The optimum inverse tableau is:

Basic Variables	Inverse Tableau					Basic Values
$x_5$	$1$	$-\frac{1}{2}$	$\frac{1}{2}$	$0$	$5$	
$x_6$	$3$	$4$	$1$	$0$	$2$	
$-z$	$-3$	$-8$	$1$	$1$	$-179$	

Answer each of the following questions with respect to this original problem:

- (a) Suppose a new primary food, food 7, is available in the market at 88 cents/kg. One kilogram of this food contains two units of vitamin A and four units of vitamin C. Should the family include this new food in its diet? If not, how much should the cost of this food decrease before the family can consider including it in its diet?

At this breakeven cost show that there is an optimum diet that includes food 7 and another that does not.

What is the optimum diet if food 7 is actually available at 32 cents/kg?

- (b) Consider the original family's diet problem again. The family has just read an article in a health magazine that says that another nutrient, vitamin E, is very important for the family's health. The vitamin E content of the six primary foods are 2, 3, 5, 2, 1, and 1 units/kg, respectively. The minimum requirement of vitamin E is 10 units. The family wants to include this additional constraint in the problem. How does this change its optimum diet?

- (c) Referring to the original problem, in

addition to meeting the minimum requirements on vitamins A and C, suppose the family decides to have the diet consist of a total of 2000 calories exactly. The calorie contents of the six primary foods are 160, 20, 500, 280, 300, and 350/kg, respectively. How does this additional requirement change its optimal diet?

- (d) What is the marginal effect of increasing the minimum requirement of vitamin C on the cost of the optimal diet in the original family's diet problem? How much can the minimal requirement of vitamin C increase before the basis  $B_1$  becomes nonoptimal to the problem? What is an optimal diet when this requirement is 39 units?

- (e) In the original family's diet problem, suppose that each additional unit of vitamin A in the diet is expected to bring an average savings of 10 cents in medical expenses. How does this alter the optimal diet?

- (f) For what range of cost per kilogram of primary foods 4 and 5 does the basis  $B_1$  remain optimal to the original family's diet problem?

- (g) What happens to the optimal diet in the original family's diet problem if the vitamin C content of food 4 changes? For what range of values of this quantity does the basis  $B_1$  remain optimal? Suppose a richer version of food 4 containing  $1 + x$  units of vitamin C per kilogram is available at a cost of  $50 + 4x$  cents/kg for any  $x \geq 0$ . What is the minimum value of  $x$  at which it becomes attractive for the family to include it in its diet?

- 9.17 Data on different solvents are given in the table at the top of page 320. Let  $x_j$  be the



	Solvent				Chemical Requirement in Blend per kg
	1	2	3	4	
Chemical 1 content (units/kg)	180	120	90	60	$\geq 90$
Chemical 2 content (units/kg)	3	2	6	5	$\leq 4$
Cost (cent/kg)	16	12	10	11	

proportion of solvent type  $j$  in the blend,  $j = 1$  to 4. Variables  $x_5$  and  $x_6$  are the slack variables associated with chemicals 1 and 2 requirements, respectively. The basis  $B_1$  associated with the basic vector  $(x_2, x_3, x_5)$  is an optimum basis for the problem. Compute  $B_1^{-1}$ , and using it, answer the following parts, each of which is independent of the others.

- Write down the dual problem and the complementary slackness conditions for optimality. Obtain an optimum feasible solution for the dual problem from the above information.
- How much does the optimum objective value change if the minimal chemical 1 requirement is changed to 88 units? Why?
- Let  $\beta_1$  be the minimal chemical 1 requirement. Its present value is 90. For what range of values of  $\beta_1$  does  $B_1$  remain optimal to the problem? When  $\beta_1$  is 114, what is an optimum feasible solution to the problem?
- How much can the cost per kilogram of solvent 3 change before  $B_1$  becomes nonoptimal to the problem? When the cost per kilogram of solvent 3 becomes 11 cents, what is an optimum solution?

- 9.18 Consider the LP given at the bottom. Variables  $x_5, x_6, x_7$  are the slack variables corresponding to the various inequalities. The basis  $B_1$  corresponding to the basic vector  $(x_1, x_3, x_2)$  is optimal to the problem. Compute  $B_1^{-1}$ .

$$\begin{array}{ll}
 \text{Minimize} & z(x) = -2x_1 - 4x_2 - x_3 - x_4 \\
 \text{Subject to} & x_1 + 3x_2 + x_4 \leq 8 = \text{available amount of raw material 1} \\
 & 2x_1 + x_2 \leq 6 = \text{available amount of raw material 2} \\
 & x_2 + 4x_3 + x_4 \leq 6 = \text{available amount of raw material 3} \\
 & x_j \geq 0 \quad \text{for all } j
 \end{array}$$

- If the availability of only one of the raw materials can be marginally increased, which one should be picked? Why?

- For what range of values of  $b_1$  (the amount of raw material 1 available) does the basis  $B_1$  remain optimal? What is an optimal solution to the problem if  $b_1 = 20$ ?

- If seven more units of raw material 1 can be made available (over the present 8 units), what is the maximum you can afford to pay for it? Why?

- The company has an option to produce a new product. Let  $x_8$  be the number of units of this product manufactured. The input-output vector of  $x_8$  will be  $(10, 20, 24 - 3\lambda, -25 + \lambda/4)^T$ , where  $\lambda$  is a parameter that can be set anywhere from 0 to 6. What is the minimum value of  $\lambda$  at which it becomes profitable to produce the product? What is an optimum solution when  $\lambda = 4$ ?

- 9.19 Consider the diet problem with data given at the top of page 321. Let  $x_1, x_2$ , and  $x_3$  be the amounts of greens, potatoes, and corn included in the diet, respectively. Let  $x_4, x_5$ , and  $x_6$  be the slack variables representing the excess of vitamins A, C, and D, respectively, in the diet. The basis  $B_1$  associated with the basic vector  $(x_4, x_1, x_6)$  is optimal to this problem.
- Find the optimum primal and dual solutions associated with the basis  $B_1$ .

Nutrient	Nutrient Content in Foods Available (units/kg)			Minimum Daily Requirement (MDR) for Nutrient
	Greens	Potatoes	Corn	
Vitamin A	10	1	9	5
C	10	10	10	50
D	10	11	11	10
Cost (cent/kg)	50	100	51	

- (b) A new food (milk) has become available. One liter of milk contains 0, 10, and 20 units, respectively, of vitamins A, C, and D and costs 40 cents. Would you recommend including it in the diet? Why? What is the highest price of milk at which it is still attractive to include it in the diet?
- (c) Consider the original problem again. A nutrition specialist claims that the actual MDRs for vitamins A, C, and D should really be 5, 50 + 10 $\lambda$ , and 10 + 15 $\lambda$ , where  $\lambda$  is a nonnegative parameter; and an experiment is proposed to estimate  $\lambda$ . Assuming that the claim is correct, find  $\lambda$ , the maximum value of  $\lambda$  at which  $B_1$  is still an optimum basis to the problem. What is an optimum solution if  $\lambda = \bar{\lambda} + 1$ ?
- 9.20 Consider the LP (9.1) and its dual. Discuss what effects the following have on the primal and dual feasible solution sets and the respective optimal objective values.
- Introducing a new nonnegative primal variable.
  - Introducing a new inequality constraint in the primal problem.
  - Introducing a new equality constraint in the primal problem.
- 9.21 Consider the LP given at the bottom. Construct the inverse tableau for this problem corresponding to the basic vector  $(x_1, x_2, x_3)$  and verify that it is an optimum basic vector. Using

$$\begin{aligned}
 \text{Minimize } z(x) &= -4x_1 + 9x_3 + 12x_4 - 3x_5 + 7x_6 \\
 \text{Subject to } & x_1 + 2x_2 - x_3 + x_4 + 2x_5 + 2x_6 = 3 \\
 & x_2 + x_3 + 2x_4 + x_5 + 2x_6 = 6 \\
 & -x_1 + 2x_3 + 2x_4 - 2x_5 + x_6 = 5 \\
 & x_j \geq 0 \quad \text{for all } j = 1 \text{ to } 6.
 \end{aligned}$$

it, answer the following questions. Each of these questions is independent of the others, and each refers to the original problem given above.

- Find the range of values of the parameter  $c_2$  for which the solution obtained above remains optimal to the problem. If the new value of  $c_2$  is slightly greater than the upper bound of the optimality range computed above, which variables become eligible to enter the basic vector? Why?
- Find the range of values of  $b_1$  for which the basic vector  $(x_1, x_2, x_3)$  remains optimal to the problem. What is the slope of the optimum objective value as a function of  $b_1$  in this optimality range? Why?
- Consider the original LP back again. Suppose a new activity leads to a new available. This activity leads to a new nonnegative variable  $x_7$ , with the data  $A_7 = (0, 2, 3)^T$ ,  $c_7 = 18$ , in the original tableau for the problem. Is it worthwhile for the decision maker to perform this new activity? Why? If not, determine the value to which  $c_7$  should decrease (assuming that all the other data, including the entries in  $A_7$ , remain unaltered) before this new activity becomes economically competitive. Find an optimum solution to the augmented problem assuming that  $A_7$  is as given above, but that  $c_7 = 12$ , using the methods of sensitivity analysis.



$$\begin{aligned}
 &\text{Minimize Cost} = 28x_1 + 67x_2 + 12x_3 + 35x_4 \\
 &\text{Product 1 output} = x_1 + 2x_2 + x_4 \geq 17 \\
 &\text{Product 2 output} = 2x_1 + 5x_2 + x_3 + 2x_4 \geq 36 \\
 &\text{Product 3 output} = x_1 + x_2 + 3x_4 \geq 8 \\
 &x_j \geq 0 \quad \text{for all } j
 \end{aligned}$$

- 9.22 In the LP given at the top,  $x_5, x_6, x_7$  are the slacks in that order. Vector  $(x_1, x_2, x_7)$  is an optimum basic vector. (1) What is the most critical product for the company? Why? (2) Determine the optimality range of  $c_1$  for the present basis. (3) Determine the optimality range of  $b_3$ . (4) Find the new optimum when  $b_3$  changes to 16 from 8.

- 9.23 In (9.1), let  $g(b_1)$  denote the optimum objective value as a function of  $b_1$  when all other data remain fixed. Prove that  $g(b_1)$  is piecewise linear convex. Discuss how to minimize  $g(b_1)$  over  $b_1$ . From this, discuss how to find the new optimum when a constraint from (9.1) is eliminated.

- 9.24 Let  $K = \{x: Ax = b, x \geq 0\}$ . Let  $f_1(x) = cx$ ,  $f_2(x) = dx$ . If both  $f_1(x)$  and  $f_2(x)$  are  $> 0$  on  $K$ , develop an algorithm for minimizing  $f_1(x)/f_2(x)$  over  $x \in K$ , using a parametric right-hand-side LP. Generalize to the case where  $f_1(x), f_2(x)$  may not be positive on  $K$ .

—[Y. P. Aneja, V. Aggarwal, and K. P. K. Nair, "On a class of quadratic programs," *University of New Brunswick, Fredericton, N. B., Canada*]

- 9.25 Let  $(P)$  be the LP: minimize  $cx$ , subject to  $A_i x \geq b_i, i = 1$  to  $m$ . Let  $x^1$  be optimal to the LP: minimize  $cx$ , subject to  $A_i x \geq b_i, i = 2$  to  $m$ . Let  $x^2$  be optimal to the LP: minimize  $cx$ , subject to  $A_i x = b_i, i = 2$  to  $m$ . Prove that one of the points  $x^1$  or  $x^2$  must be optimal to  $(P)$ .

- 9.26 Let  $K$  be the set of feasible solutions of (9.1), and  $K_1 = \{x: x \in K, x_n = 0\}$ . Assume that (9.1) is nondegenerate and that  $K, K_1$  are both nonempty and bounded. All data are fixed except  $c_n$ . Prove that there exists a  $\gamma$  such

that for all  $c_n > \gamma$ ,  $x_n$  is a nonbasic variable in all optimum bases for (9.1), and for all  $c_n < \gamma$ ,  $x_n$  is a basic variable in all optimum bases for (9.1). How can  $\gamma$  be computed?

- 9.27 Consider the following variant of the simplex algorithm for solving the LP (9.1). It begins with a feasible basic vector  $x_B$  for (9.1). Let  $\bar{x}$  be the associated BFS, and let  $\bar{c} = (c_j)$  be the vector of relative cost coefficients with respect to  $x_B$ . If  $\bar{c} \geq 0$ ,  $\bar{x}$  is optimal to (9.1). Let  $\bar{c} \not\geq 0$ , let  $J = \{j: \bar{c}_j < 0\}$ . Let  $A_{n+1} = \sum_{j \in J} A_j w_j, c_{n+1} = \sum_{j \in J} c_j w_j$ , where  $w_j > 0$  for each  $j \in J$  are positive weights. Introduce a new variable  $x_{n+1}$  into (9.1) with its column vector  $A_{n+1}$  and cost coefficient  $c_{n+1}$ . Let  $X = (x_1, \dots, x_n; x_{n+1})$ . Bring  $x_{n+1}$  into the basic vector  $x_B$  for this augmented problem and suppose this leads to the new solution  $\tilde{x} = (\tilde{x}_1, \dots, \tilde{x}_n; \tilde{x}_{n+1})$ . Remembering that  $x_{n+1} = \sum_{j \in J} w_j x_j$ , obtain the feasible solution  $\hat{x}$  of the original problem (9.1) corresponding to  $\tilde{x}$  for the augmented problem. In general,  $\hat{x}$  may not be a BFS for (9.1). Using the method discussed in the proof of Theorem 3.3 of Section 3.5.5, obtain a BFS  $x^*$  for (9.1) whose objective value is  $\leq$  the objective value at  $\hat{x}$ . Repeat the procedure, starting with the new BFS  $x^*$ .

Possible choices for the weights in this procedure are either  $w_j = 1$  for all  $j \in J$ , or  $w_j = -\bar{c}_j$  for each  $j \in J$ , or some other positive values.

Compare this procedure for solving (9.1), starting with a feasible basic vector for it, with the usual simplex algorithm, in terms of computational efficiency. How does this procedure differ from the simplex algorithm geometrically?