

# The 2nd-shortest Path

朱理真

**Date: 2020-12-4**

## Chapter 1: Introduction

The 2-nd shortest path in a graph is the path from the source to its destination with traceback permitted, which is longer than the shortest path, but is no longer than any other path.

Actually, there has been a lot of algorithms that solve the kth-shortest path problem using Dijkstra, a\*, d\* or other methods to find the shortest paths, and the time complexity are also liner-logarithmic.

Our duty is to find the second shortest path with a given graph and get how long it is.

Input Specification:

Give M(1, 1000)(vertices), N(1, 5000)(edges), and for next each line, provide an edge with the form "vertex1 vertex2 weight"

Output Specification:

print in one line the length of the second-shortest path between node 1 and node M, then followed by the nodes' indices in order. There must be exactly 1 space between the numbers, and no extra space at the beginning or the end of the line.

## Chapter 2: Algorithm Specification

### Data Structures

```
typedef int element; //used for elements in the stack, which are vertex numbers
```

#### NODE

```
typedef struct node // used in ADJ Linked List as a node
{
    int vertex;
    int weight;
    struct node* next;
}*list;
```

#### GRAPH

```
typedef struct g{
    int v; //The number of the vertices in graph
    int e; //The number of edges in graph
    list* AdjList; // ADJ Linked List, with edges storing in it.
}*graph;
```

#### HEAP

```
typedef struct h
```

```
{
    int capacity; //the max number of elements a heap can hold.
    int* vertices; //store the vertices number for every node. vertices
[0] = the number of elements in the heap.
    int* distance; //store the distance of each vertex.
}*heap;
```

STACK

```
typedef struct s
{
    int top; //stack top
    int capacity; //the max number of elements a heap can hold.
    element* container; //store the elements in the stack, that is, a v
ertex number.
}*stack;
```

## Main Algorithm

```
main(){
    Create graph g and read the data from input.
    Create heap h.

    Calculate the shortest path from end vertex to start vertex and
return the shortest length and path information.

    Calculate the second shortest path form start vertex to end vertex
and return the second shortest length and path information.

    Print the second shortest length.

    Using the path information, print the path.
}
```

## Key Algorithms

```
ShortestPath(heap, graph, distances array, source(start point),
destination, path array){
    Create an array named include, recording if each vertex has been
popped from the heap. All elements are not included at the beginning.

    Initialize the distance and path array.
    Make Heap Empty.

    insert the source into the heap.
```

```

while(the heap is not empty){
    vertex = the element popped from the heap top.
    if vertex has already been included
        continue.
    Include the vertex.

    for each neighbor vertex of current vertex{
        if (neighbor is not included yet){
            if (the distance of neighbor is larger
                than the distance of current vertex plus
                the weight between the them){

                distances[neighbor]=distances[vertex]+weight.
                path[neighbor]=vertex.
                Insert the neighbor with its new distance into heap.
            }
        }
    }
}
return the distances[destination]
}

```

**SecondShortestPath**(heap, graph, source, destination, distances array, second\_distance array, path array, second\_path array){

Initialize the distances array, path array.  
Initialize the second\_distances array, second\_path array.

insert the source into the heap.

```

while(the heap is not empty){
    vertex = the element popped from the heap top.
    distance = the distance of the vertex just popped.
    if distance > second_distances[vertex]
        continue.

    for each neighbor vertex of current vertex{
        NewDistance = distance + weight between vertex and neighbor.
        NewVertex = vertex.
        if (NewDistance < distances[neighbour]){

            swap(NewDistance, distances[neighbor])
            swap(NewVertex, path[neighbor])
            Insert the neighbor with distances[neighbor] into the

```

heap.

```

        }
        if (NewDistance > distances[neighbor] and
            NewDistance < second_distances[neighbor]){

            second_distances[neighbor] = NewDistance.
            second_path[neighbor] = NewVertex.
            insert the neighbor along with second_distances[neighbor]
into the heap.
        }
    }
}
return second_distance[destination]
}

```

### **PrintPath(**

```

    stack,
    end vertex,
    shortest length from start vertex to end vertex,
    distances array (stores the shortest distances from start vertex),
    reverse_distances array (shortest distances from end vertex),
    path array,
    reverse_path array,
    second path array
){
    vertex = end vertex
    backed=FALSE
    while(the vertex is not the break point "BEGIN"){
        if(
            the vertex is not in the shortest path OR
            the vertices in the shortest path between the current vertex
and last vertex just pushed are skipped OR
            the last vertex of last vertex is the same as the current
vertex
        )//if you don't understand, see the pictures below.
            backed = TRUE
            push the vertex into stack

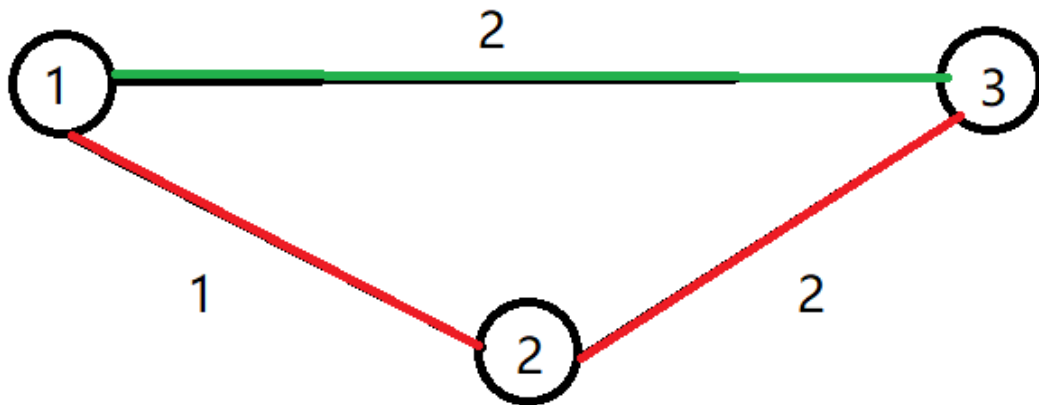
            if backed
                vertex = path[vertex]
            else
                vertex = sec_path[vertex]
        }
    }
}

```

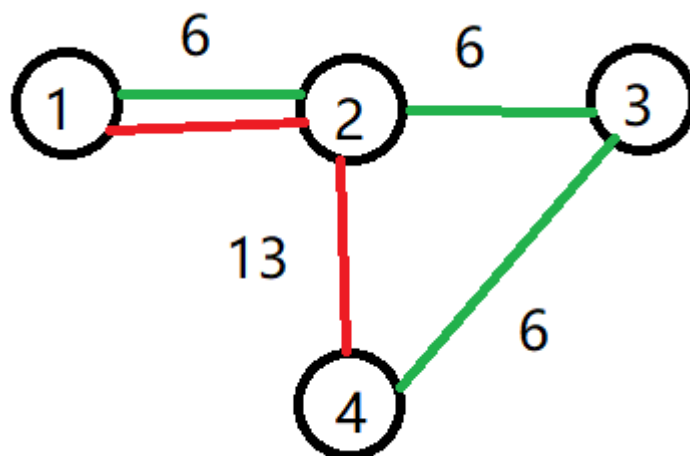
```

while(the stack is not empty){
    pop the vertex from the stack.
    print vertex
}
}

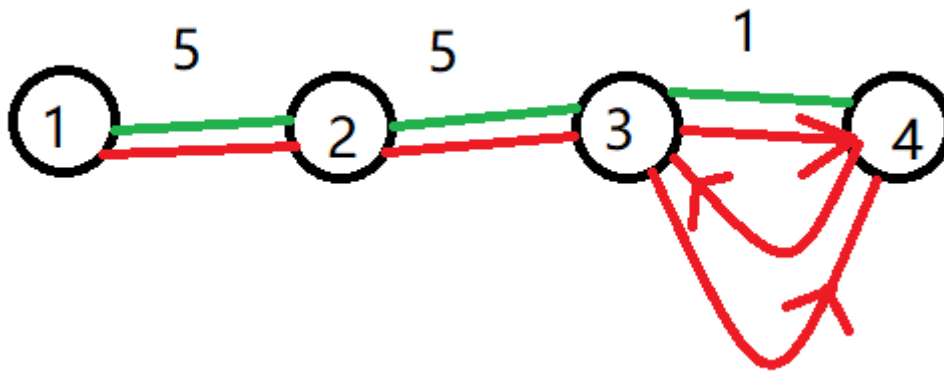
```



the case that one vertex is not in the shortest path



the case that vertices in the second shortest path are all in the shortest path, but some vertices in the shortest path are not included.



The case the last vertex of last vertex of a certain vertex is itself.

## Other Algorithms

The algorithm for heap:

```

decrease(heap, index, NewDistance){
  vertex = vertex of index
  while(vertex is not the root of the min-heap){
    p= parent of index
    if(NewDistance < distance of index ){
      assign the information of p to the index.
      index = p.
    }
    else break.
  }
  Let vertex be the vertex of index, and NewDistance be the distance
  of vertex.
}

```

```

increase(heap, index, NewDistance){
  vertex = vertex of the index

  while(index has children){
    child = the child of index whose distance is the smallest
    if(NewDistance > distance of child){
      assign the information of child to the index
      index = child
    }
    else break.
  }
}

```

```

    Let vertex be the vertex of index, and NewDistance be the distance
of vertex.
}

insert(heap h, vertex, distance){
    increase the size of heap by 1.
    let the last element have vertex as its vertex, distance as its
distance.

    calling function "decrease" to let it go up
}

Minimum(heap){
    if heap has no vertex, return FALSE.
    return the vertex and distance of top element.
    LastLeaf = the last element in the heap.
    decrease the size of heap by 1.
    assign the LastLeaf to the top element.
    calling function "increase" to let it down.
}

```

### Chapter 3: Testing Results

All inputs are also stored in txt file in the folder "document". You should copy the data there to test the program. Inputs here just for read.

TestCase 1	case purpose
	sample
	expected result
	240 1 2 4 5
	actual behavior
	240 1 2 4 5
	possible cause
	current status
	pass
	Input
	5 6



	1 2 50 2 3 100 2 4 150 3 4 130 3 5 70 4 5 40
TestCase 2	<b>case purpose</b>
	the case that vertices in the second shortest path are all in the shortest path, but some vertices in the shortest path are not included.
	<b>expected result</b>
	8 1 4
	<b>actual behavior</b>
	8 1 4
	<b>possible cause</b>
	<b>current status</b>
	pass
	<b>Input</b>
	4 4 1 2 6 1 4 8 2 3 5 2 4 1
TestCase 3	<b>case purpose</b>
	The case the last vertex of last vertex of a certain vertex in the second shortest path is itself.
	<b>expected result</b>
	14 1 2 3 2 4 5
	<b>actual behavior</b>
	14 1 2 3 2 4 5

	possible cause
	current status
	pass
	input
	5 5 1 2 6 2 3 1 2 4 2 2 5 10 4 5 4
TestCase 4	case purpose
	Minimum size
	expected result
	9 1 2 1 2
	actual behavior
	9 1 2 1 2
	possible cause
	current status
	pass
	input
	2 1 1 2 3
TestCase 5	case purpose
	the case that one vertex is not in the shortest path
	expected result
	7 1 4 5
	actual behavior
	7 1 4 5
	possible cause

	current status
	pass
	input
	5 6
	1 2 1
	1 4 3
	2 3 2
TestCase 6	2 4 4
	3 5 3
	4 5 4
	case purpose
	Larger Size
	expected result
	271 1 7 1 20
	actual behavior
	271 1 7 1 20
	possible cause
	current status
	pass
	Input
	20 87
	1 2 268
	1 7 64
	1 8 2982
	1 9 2325
	1 10 3595
	1 12 3249
	1 15 1573

	1 17 2210
	1 20 143
	2 3 1190
	2 4 3265
	2 5 3459
	2 6 1767
	2 7 2413
	2 11 620
	2 14 4003
	2 15 2802
	2 18 2069
	2 19 4060
	2 20 2899
	3 4 1408
	3 6 2622
	3 7 1311
	3 13 174
	3 14 4029
	3 15 4460
	3 16 3225
	3 20 4922
	4 5 3549
	4 9 3920
	4 10 3541
	4 12 3637
	4 18 560
	5 7 3056
	5 9 207
	5 12 1216

	5 13 411
	5 14 2393
	5 15 131
	5 16 3143
	5 17 4734
	5 20 3079
	6 7 4109
	6 10 4544
	6 13 525
	6 17 2376
	6 18 1839
	6 19 2795
	6 20 1130
	7 8 1488
	7 9 882
	7 13 887
	7 14 3136
	7 17 1850
	7 18 2703
	8 10 1913
	8 13 4768
	8 14 2745
	8 15 3214
	8 18 1751
	8 20 891
	9 11 581
	9 12 4115
	9 19 2800
	10 14 2282

	10 18 1856 10 19 1341 10 20 1943 11 19 3613 12 13 3368 12 14 3994 12 15 3390 12 16 4304 12 17 2919 12 19 2441 13 16 731 13 18 4167 13 20 629 14 15 4034 14 17 4475 14 19 3299 14 20 1077 15 16 2699 15 18 1122 15 20 3792 16 18 174 16 19 140
TestCase 7	case purpose
	Larger Size
	expected result
	410 1 8 23 8 28 40
	actual behavior
	410 1 8 23 8 28 40
	possible cause

	current status
	pass
	input
	40 381 1 2 3494 1 3 3362 1 8 19 1 12 1263 1 14 3086 1 15 2690 1 20 1717 1 23 4151 1 24 2576 1 26 4244 1 31 369 1 36 3459 1 37 2113 1 38 2761 1 39 1468 2 3 2431 2 4 1822 2 6 362 2 7 3191 2 8 215 2 9 3078 2 10 4999 2 11 790 2 12 3271 2 13 1142

2 18 355
2 25 3831
2 27 2300
2 28 3643
2 29 818
2 31 914
2 32 1731
2 33 4758
2 37 786
2 39 2674
3 4 3590
3 6 954
3 7 1211
3 9 1856
3 10 3925
3 13 3725
3 14 834
3 16 1470
3 18 3926
3 19 3589
3 20 159
3 21 4786
3 22 827
3 25 4688
3 26 2241
3 27 1030
3 29 1123
3 31 4854
3 34 4261



	3 35 3214
	3 37 3017
	3 39 4269
	4 5 3496
	4 6 2802
	4 9 1036
	4 11 1533
	4 12 4156
	4 15 3737
	4 21 2294
	4 27 4506
	4 28 4676
	4 29 4787
	4 31 2705
	4 32 866
	4 33 3565
	4 34 2956
	4 37 922
	4 39 4005
	5 6 982
	5 9 2147
	5 12 3799
	5 15 4332
	5 17 3142
	5 21 3954
	5 27 2943
	5 29 1517
	5 31 4248
	5 32 1231

5 35 4232
5 38 512
5 39 4863
5 40 1187
6 8 1071
6 9 207
6 13 3007
6 15 4315
6 16 1977
6 20 4340
6 21 2908
6 27 4808
6 28 293
6 30 793
6 32 4333
6 33 2126
6 34 1685
6 36 4080
6 37 2890
6 39 2029
6 40 107
7 8 2191
7 9 2800
7 10 4590
7 11 3365
7 12 1691
7 13 940
7 16 3433
7 17 3583

7 19 863
7 21 4631
7 22 2227
7 23 4905
7 24 3381
7 26 296
7 28 3212
7 31 2717
7 36 3427
7 38 1367
7 40 2799
8 9 2749
8 12 1200
8 14 1070
8 16 2544
8 19 1830
8 20 4803
8 22 1528
8 23 1
8 24 1629
8 25 4712
8 28 117
8 30 4843
8 31 2485
8 32 2443
8 33 1254
8 34 2589
8 38 665
8 40 2416

9 13 3845
9 14 3874
9 15 2949
9 18 4833
9 19 1217
9 20 747
9 22 823
9 25 4295
9 31 1353
9 33 4007
9 36 3120
9 37 186
10 11 415
10 16 3222
10 20 2165
10 21 647
10 22 3957
10 24 2149
10 25 4057
10 26 4067
10 27 3050
10 28 1379
10 31 4319
11 12 4433
11 13 4875
11 14 2558
11 16 3314
11 17 2780
11 18 51

	11 21 2901
	11 23 356
	11 24 4743
	11 25 1881
	11 27 3424
	11 28 805
	11 29 2211
	11 31 4524
	11 34 2621
	11 35 1766
	11 36 332
	11 38 4494
	11 40 4345
	12 13 4969
	12 14 1051
	12 15 1531
	12 22 4705
	12 23 3119
	12 26 3530
	12 27 2556
	12 28 4044
	12 29 4966
	12 30 524
	12 32 4017
	12 33 1065
	12 38 2358
	12 40 4125
	13 14 2794
	13 16 2055

	13 18 4228
	13 19 2318
	13 20 4021
	13 23 4139
	13 24 1381
	13 25 739
	13 26 3104
	13 27 4259
	13 28 2327
	13 30 4266
	13 32 1531
	13 34 544
	13 38 2407
	13 39 4449
	13 40 80
	14 16 44
	14 17 253
	14 19 4390
	14 24 4477
	14 25 3213
	14 26 2761
	14 27 3535
	14 28 2273
	14 32 757
	14 33 3980
	14 37 4192
	14 39 3705
	14 40 1278
	15 16 776

	15 18 4852
	15 20 2529
	15 22 2251
	15 23 3093
	15 24 1323
	15 26 3185
	15 27 678
	15 28 1793
	15 34 3574
	15 36 4997
	15 37 59
	15 39 2767
	16 18 1535
	16 19 3066
	16 24 4837
	16 27 4149
	16 29 411
	16 37 2432
	16 38 1055
	16 39 3018
	16 40 3232
	17 19 2474
	17 21 1896
	17 23 2386
	17 25 2931
	17 31 4373
	17 32 2993
	17 34 156
	17 35 1205

	17 36 3022
	17 37 3004
	17 38 1947
	17 40 1669
	18 24 1696
	18 26 3640
	18 28 1733
	18 29 3342
	18 30 1376
	18 32 1830
	18 33 2148
	18 35 1835
	18 36 3808
	18 37 3435
	18 39 4547
	19 21 1866
	19 23 3032
	19 28 2102
	19 29 501
	19 34 3342
	19 39 3598
	19 40 4951
	20 21 1462
	20 23 1183
	20 25 2857
	20 26 86
	20 28 3679
	20 29 1024
	20 31 3256



	20 32 4474
	20 34 2381
	20 35 2093
	20 37 1512
	20 39 4673
	20 40 3488
	21 22 249
	21 25 2445
	21 27 2423
	21 28 4412
	21 29 2258
	21 30 4107
	21 32 755
	21 34 758
	21 36 1320
	21 38 962
	21 40 4177
	22 25 1912
	22 28 4026
	22 29 271
	22 30 451
	22 35 2914
	22 37 2490
	22 40 4681
	23 24 1009
	23 25 2182
	23 26 4956
	23 28 1499
	23 29 1584

	23 30 4925
	23 31 4771
	23 37 1742
	23 40 518
	24 28 570
	24 34 1435
	24 36 84
	24 37 3058
	24 39 3136
	25 28 2155
	25 33 4863
	25 40 3398
	26 27 3418
	26 28 480
	26 30 1333
	26 33 3278
	26 34 2722
	26 36 775
	26 37 3734
	26 38 100
	26 39 516
	27 29 4390
	27 31 634
	27 33 264
	27 34 3126
	27 35 3240
	27 38 4076
	28 29 540
	28 34 2278

	28 37 382
	28 38 840
	28 39 1979
	28 40 272
	29 30 3333
	29 31 1049
	29 32 2315
	29 33 825
	29 35 4879
	29 39 2167
	30 32 4516
	30 35 1209
	30 36 4116
	30 38 862
	30 39 1038
	31 36 2696
	31 37 807
	31 39 1483
	31 40 2651
	32 33 1748
	32 34 4640
	32 35 4828
	32 36 1098
	32 38 523
	32 39 2861
	32 40 4579
	33 34 3464
	33 35 4697
	34 35 3074

	34 36 1122 34 37 3923 34 38 2749 34 40 1241 35 36 2423 35 37 1009 35 38 4844 38 40 3005
TestCase 8	case purpose
	MAXSIZE(v=1000,e=5041)
	expected result
	2564 1 796 892 173 552 792 234 364 234 364 108 392 782 1000
	actual behavior
	2564 1 796 892 173 552 792 234 364 234 364 108 392 782 1000
	possible cause
	current status
	pass
	input
	Too much to show.
	The input data are put in the text file “t8” in folder “document”

## Chapter 4: Analysis and Comments

Using the detailed pseudo code to analyze.

### Time Complexity:

**Analyze:**

**In function main:**

Normal sequential execution statements take  $O(1)$  time.

Function CreateGraph, CreateHeap( $O(1)$ ), ShortestPath, SecondShortestPath and PrintPath are called.

Therefore,  $T(\text{main}) = T(\text{CreateGraph}) + T(\text{ShortestPath}) + T(\text{SecondShortestPath}) + \text{PrintPath}$

The following is the time complexity analysis of the functions mentioned above.

**In function CreateGraph:**

Normal sequential execution statements take  $O(1)$  time.

In while statement:

```
while (e--)  
{  
    scanf("%d%d%d",&va,&vb,&w);  
    va--;  
    vb--;  
    g->AdjList[va]=CreateNode(vb,w,g->AdjList[va]); //Let the new node  
    be the first element of the ADJ list.  
    g->AdjList[vb]=CreateNode(va,w,g->AdjList[vb]);  
}
```

Every statement in while loop takes  $O(1)$  time.

Therefore, the  $T(\text{CreateGraph}) = O(e)$ .

**In function ShortestPath:**

Normal sequential execution statements take  $O(1)$  time.

In outer while loop:

it takes  $O(e)$  time to visit every edge, and to insert and delete them (by a heap),  $O(\log v)$  time is taken (for each edge). Therefore, the while loop takes  $O(e \cdot \log v)$  time

Hence,  $T(\text{ShortestPath}) = O(e \cdot \log v)$

**In function SecondShortestPath:**

The function has a similar structure with function ShortestPath, and they have the same time complexity  $O(e \cdot \log v)$ , but it may have a larger constant for every edge will be inserted and deleted twice, one for shortest path, the other one for the second shortest path.

Hence,  $T(\text{SecondShortestPath}) = O(e \cdot \log v)$

**In function PrintPath:**

Normal sequential execution statements take  $O(1)$  time.

Let  $L$  = the length of the second shortest path.

Clearly in the first while loop, statements will be executed  $O(L+1)$  times, the same as the second while loop.

Therefore,  $T(\text{PrintPath}) = O(L)$

**Conclusion:**

Above all,  $T = T(\text{main}) = O(e) + O(\log v) + O(\log v) + O(L)$

$= O(\log v) + O(L) = O(\log v) + O(e+2) = O(\log v) + O(e) = O(\log v)$

(

$L$  be the length of the second shortest path,

$e$  be the number of edges,

$v$  be the number of vertices.

)

## Space Complexity:

**Analyze:****In function main:**

graph  $g$  created, taking  $O(e)$  space

heap  $h$  created, taking  $O(e)$  space

stack  $s$  created, taking  $O(e)$  space

6 arrays created, taking  $O(v)$  space

And in function ShortestPath, array "include" is created taking another  $O(v)$  space.

Other local variables take  $O(1)$  space

And there's no recursion in the program.

Above all,  $S = S(\text{main}) = O(e) + O(v) = O(e + v)$

( $e$  is the number of edges,  $v$  is the number of vertices)

## Comments:

**Possible Improvement:**

1)

put the vertices and distances together forming a struct.

I regret that I didn't put the vertices field and distance field of structure heap into a smaller struct, and it is a little inconvenient when I pop and insert the element into the heap, for I need two parameters for every function, and I have to use pointers to get the result instead of "return".

2)

make the names of variable more understandable.

Following the tradition, I named the variable "distance" as "distance", which save the shortest distance from a vertex to the source, but when you also use the second shortest path, it may be confusing for still using this word, because you don't know which "distance" it is.

It is the same as variable "path" and "second\_path". And I worried that the "reverse\_dst" is also hard to understand.

## Appendix: Source Code (in C)

```
//TheSecondShortestPath.c

/*header files*/
#include <stdio.h>
#include <assert.h>
#include <stdlib.h>
#include <string.h>
#include <limits.h>

/*Macro Definition*/
#define New(type,n) (type*)malloc(sizeof(type)*n) //Alloc a new space with specific size and use.
#define NewArray(type,n) (type*)calloc(n,sizeof(type)) //Alloc a new space served as an array with specific size and use.
#define NULLPOS -
1 //which implies a variable is not existing, often used in stack, queue, heap and other structures.
#define INFINITY INT_MAX //Used to initialize the distance of vertices from the start point.
#define ROOT 1 //point to the minimum element of the heap.
#define Parent(i) (i/2) //the parent of the node i in the heap
#define Left(i) (2*i) //the left child of the node i in the heap
#define Right(i) (Left(i)+1) //the right child of the node i in the heap
```

```

#define TRUE 1 //Used to tell the status of a function or a variable, the same as FALSE
#define FALSE 0
#define BEGIN -2 //Used to simulate the last vertex of the start vertex

typedef int element; //used for elements in the stack, which are vertex numbers

typedef struct node // used in ADJ linked list as a node
{
    int vertex;
    int weight;
    struct node* next;
}*list;

typedef struct g{ //The structure of the whole graph.
    int v; //The number of the vertices in graph
    int e; //The number of edges in graph
    list* AdjList; // ADJ linked list, with edges storing in it.
}*graph;

typedef struct h //The structure of a heap
{
    int capacity; //the max number of elements a heap can hold.
    int* vertices; //store the vertices number for every node. vertices[0] = the number of elements in the heap.
    int* distance; //store the distance of each vertex.
}*heap;

typedef struct s //The structure of a stack
{
    int top; //stack top
    int capacity; //the max number of elements a heap can hold.
    element* container; //store the elements in the stack, that is, a vertex number.
}*stack;

/* Function: Top_k
 * Used to get the k_th element away
 * from the stack top.
 *
 * k: the k_th number, 0 is on the top.
 * s: the stack need to provide
 */

```



```

int Top_k(int k,stack s);

/* Function: Minimum
 * Used to get top element of the heap.
 *
 * h: heap structure need to provide
 *
 * return_dst: provide the pointer,
 * and the function will put the
 * distance of the element in it.
 *
 * return_vertex:provide the pointer,
 * and the function will put the
 * vertex number of the element in it.
 *
 * If there's no element in the heap,
 * it will return FALSE;
 * else, return TRUE.
 */
int Minimum(heap h,int* return_dst,int* return_vertex);

/* Function: ShortestPath
 * Used to calculate the shortest
 * path in the graph and its path
 * given a source and a destination.
 * with Dijkstra Algorithm.
 *
 * h: heap you need to provide
 *
 * g: graph you need to provide
 *
 * dst: the function will save the shortest
 * distances of each vertex from the source
 * in it, please make sure it points
 * to a address in the memory with enough
 * space.
 *
 * source: start point
 *
 * dest: end point
 *
 * path: the function will save the shortest
 * paths of each vertex from the source
 * in it, please make sure it points

```

```

* to a address in the memory with enough
* space.
*/
int ShortestPath(heap h,graph g,int dst[],int source,int dest,int path[
]);

/* Function: ShortestPath
* Used to calculate the second shortest
* path in the graph and its path
* given a source and a destination
* with Dijkstra Algorithm.
*
* h: heap you need to provide
*
* g: graph you need to provide
*
* dst: the function will save the shortest
* distances of each vertex from the source
* in it, please make sure it points
* to a address in the memory with enough
* space.
*
* source: start point
*
* dest: end point
*
* dst_second: the function will save the
* second shortest distances of each vertex
* from the source in it, please make sure
* it points to a address in the memory with
* enough space.
*
* path: the function will save the shortest
* paths of each vertex from the source
* in it, please make sure it points
* to a address in the memory with enough
* space.
*
* sec_path: the function will save the second
* shortest paths of each vertex from the source
* in it, please make sure it points
* to a address in the memory with enough
* space.
*/

```

```

int SecondShortestPath(heap h,graph g,int source,int dest,int dst[],int
    dst_second[],int path[],int sec_path[]);

/* Function: ShortestPath
 * Used to make the heap empty.
 * h: the heap you need to provide.
 */
void dispose(heap h);

/* Function: exchange
 * Swap the values of two
 * integer.
 * a: operand 1
 * b: operand 2
 */
void exchange(int* a,int* b);

/* Function: push
 * push the element in
 * the stack.
 *
 * x: the element needed to
 * be provide, often being a
 * vertex number.
 */
void push(stack s,element x);

/* Function: insert
 * insert an element in heap
 *
 * h: heap
 * vertex: the vertex field of the element put in
 * dst: the distance field of the element put in
 */
void insert(heap h,int vertex,int dst);

/* Function: increase
 * increase the distance of the element in
 * heap, and let it sink in the heap, so that
 * we maintain a min-heap's property
 *
 * h: heap
 * vertex: the index of the element to increase
 * NewDistance: refresh the distance value

```

```

    */
void increase(heap h,int index,int NewDistance);

/* Function: decrease
 * decrease the distance of the element in
 * heap, and let it get up from the heap, so that
 * we maintain a min-heap's property
 *
 * h: heap
 * vertex: the index of the element to increase
 * NewDistance: refresh the distance value
 */
void decrease(heap h,int index,int NewDistance);

/* Function: SetSource
 * Often used before a Dijkstra Algorithm, to
 * Initialize the distance and path arrays given
 * a start point.
 *
 * v: the number of vertices
 * source: the start point
 * dst: distance array
 * path: path array
 */
void SetSource(int v,int source,int dst[],int path[]);

/* Function: PrintPath
 * print the second shortest path
 *
 * s: stack used in this procedure
 *
 * end: the end point
 *
 * short_path: the length of the shortest path
 *
 * dst: the shortest distances of vertices
 * away from the source (start point)
 *
 * reverse_dst: the shortest distances of vertices
 * away from the end (end point)
 *
 * path: the shortest path for source(basically, source=0)
 * reverse_path: the shortest path end(basically, end=v-1)
 * sec_path: the second shortest path for source

```

```

    */
void PrintPath(stack s,int end,int short_path,int dst[],int reverse_dst
[],int path[],int reverse_path[],int sec_path[]);

/* Function: push
 * pop the element from
 * the stack.
 *
 * return the element popped.
 */
element pop(stack s);

/* Function: push
 * Create a heap.
 * capacity: the maximum size of the heap
 * return the heap variable.
 */
heap CreateHeap(int capacity);

/* Function: push
 * Create a graph
 * v: the number of vertices
 * e: the number of edges
 * return the graph created.
 */
graph CreateGraph(int v,int e);

/* Function: CreateStack
 * Create a stack.
 * capacity: the maximum size of the stack
 * return the stack
 */
stack CreateStack(int capacity);

/* Function: CreateStack
 * Create a node.
 */
list CreateNode(int vertex,int weight,list next);

int main(){
    int v,e;
    scanf("%d%d",&v,&e);
    graph g=CreateGraph(v,e);
    heap h=CreateHeap(4*e+1);

```

```

    int* path=New(int,v);
    int* reverse_path=New(int,v);//the shortest path from end point, namely v-1.
    int* sec_path=New(int,v);
    int* dst=New(int,v);
    int* reverse_dst=New(int,v);//the shortest distances away from end point, namely v-1.
    int* sec_dst=New(int,v);
    assert(path&&reverse_dst&&reverse_path&&sec_path&&dst&&reverse_dst&&sec_dst);//assert all space allocated are valid.

    int shortest_path=ShortestPath(h,g,reverse_dst,v-1,0,reverse_path);
    int sec_shortest_path=SecondShortestPath(h,g,0,v-1,dst,sec_dst,path,sec_path);

    printf("%d",sec_shortest_path);

    stack s=CreateStack(e+3);
    PrintPath(s,v-1,shortest_path,dst,reverse_dst,path,reverse_path,sec_path);

    return 0;
}

int Top_k(int k,stack s){
    if(s->top-1-k<0) return NULLPOS;//the number of stack is too few.
    return s->container[s->top-1-k];
}

void PrintPath(stack s,int end,int short_path,int dst[],int reverse_dst[],int path[],int reverse_path[],int sec_path[]){
    /* Since the path[] only
    * store the last vertex,
    * to print the path, we
    * need a stack to invert
    * the order we print.
    */
    int vertex=end;//start from end point.
    int backed=FALSE;//if backed=TRUE, we can straightly head to the source with shortest path.
    int tmp;

    while (vertex!=BEGIN){

```

```

        if(
            //the case that the vertex is not in the shortest path.
            dst[vertex]+reverse_dst[vertex]!=short_path ||
            //the case that vertices in the shortest path between the v
            ertex and last vertex are skipped.
            (tmp=Top_k(0,s))!=NULLPOS && tmp !=reverse_path[vertex] ||
            //the case that the last vertex of last vertex is the same
            as the current vertex.
            Top_k(1,s) == vertex
        )
            backed=TRUE;//the detour has finished.

    push(s,vertex);

    if(backed){
        vertex=path[vertex];//go straight in the shoretest path.
    }
    else{
        vertex=sec_path[vertex];
    }

}

//print the vertices in the path.
while((vertex=pop(s))!=NULLPOS){
    printf(" %d",vertex+1);
}
}

void dispose(heap h){
    h->vertices[0]=0;
}

void push(stack s,element x){
    assert(s->top<=s->capacity);
    s->container[s->top++]=x;
}

element pop(stack s){
    if(!(s->top)) return NULLPOS;
    return s->container[--(s->top)];
}

```

```

void exchange(int* a,int* b){
    int tmp;
    tmp=*a;
    *a=*b;
    *b=tmp;
}

void SetSource(int v,int source,int dst[],int path[]){
    int i;
    for(i=0;i<v;i++){
        dst[i]=INFINITY;
        path[i]=NULLPOS;
    }

    dst[source]=0;
    path[source]=BEGIN;
}

list CreateNode(int vertex,int weight,list next){
    list p = New(struct node,1);
    assert(p);
    p->vertex=vertex;
    p->weight=weight;
    p->next=next;
    return p;
}

graph CreateGraph(int v,int e){
    graph g = New(struct g,1);
    assert(g);
    g->e=e;
    g->v=v;
    g->AdjList=NewArray(list,v);
    assert(g->AdjList);

    int va,vb,w;

    while (e--){
        {
            scanf("%d%d%d",&va,&vb,&w);
            va-
-; //Because the C array begins from 0, we substract every vertex by 1 to
save the space.
            vb--;

```



```

        g->AdjList[va]=CreateNode(vb,w,g->AdjList[va]); //Let the new node
        be the first element of the ADJ list.
        g->AdjList[vb]=CreateNode(va,w,g->AdjList[vb]);
    }
    return g;
}

heap CreateHeap(int capacity){
    heap h = New(struct h,1);
    assert(h);
    h->capacity=capacity;

    h->vertices=New(int,capacity+1);
    assert(h->vertices);

    h->distance=NewArray(int,capacity+1);
    assert(h->distance);

    h->vertices[0]=0; // h->vertices[0] = the number of elements in the
    heap.
    return h;
}

void decrease(heap h,int index,int NewDistance){ //we make the distances
    the key field of elements in the heap.
    int p;
    int vertex = h->vertices[index];
    while(index > ROOT){
        p=Parent(index);
        if(NewDistance<h->distance[p]){
            h->vertices[index]=h->vertices[p];
            h->distance[index]=h->distance[p];
        }
        else break;
        index = p;
    }
    h->vertices[index]=vertex;
    h->distance[index]=NewDistance;
}

void increase(heap h,int index,int NewDistance){
    int child;
    int vertex = h->vertices[index];
    int LastLeaf=h->vertices[0];

```

```

        while(index <= Parent(LastLeaf)){
            child = Left(index);
            if(child!=LastLeaf && h->distance[child]>h->distance[Right(index)])
                child=Right(index);

            if(NewDistance>h->distance[child]){
                h->vertices[index]=h->vertices[child];
                h->distance[index]=h->distance[child];
            }
            else break;
            index=child;
        }
        h->vertices[index]=vertex;
        h->distance[index]=NewDistance;
    }

void insert(heap h,int vertex,int dst){
    int LastLeaf = ++(h->vertices[0]);
    assert(LastLeaf<=h->capacity);

    h->vertices[LastLeaf]=vertex;
    decrease(h,LastLeaf,dst);
}

int Minimum(heap h,int* return_dst,int* return_vertex){
    if(h->vertices[0]==0) return FALSE;//There's no element in the heap
    .
    * return_vertex = h->vertices[ROOT];
    * return_dst = h->distance[ROOT];

    int LastLeaf = h->vertices[0]--;
    h->vertices[ROOT]=h->vertices[LastLeaf];

    increase(h,ROOT,h->distance[LastLeaf]);
    return TRUE;
}

int SecondShortestPath(heap h,graph g,int source,int dest,int dst[],int
dst_second[],int path[],int sec_path[]){
    int v=g->v;
    int vertex,distance,NewDistance,neighbor,weight,NewVertex;

```

```

dispose(h);

SetSource(v,source,dst,path);
SetSource(v,source,dst_second,sec_path);

/* The second shortest distance from source
 * to source need to be initialized as INFINITY,
 * otherwise the algorithm won't work because the
 * if-statement in line 497.
 */
dst_second[source]=INFINITY;

insert(h,source,dst[source]);

list p;

while(Minimum(h,&distance,&vertex)){
    if(distance <= dst_second[vertex]){
        p=g->AdjList[vertex];
        while(p){
            neighbor=p->vertex;
            NewDistance=p->weight+distance;
            NewVertex=vertex;

            if(NewDistance < dst[neighbor]){
                /* if the original distance is larger,
                 * it can be used for the second
                 * shortest distance. Therefore, we use
                 * a "exchange" to assign the NewDistance
                 * with original shortest distance for the
                 * next if-statement.
                 */
                exchange(&dst[neighbor],&NewDistance);
                exchange(&NewVertex,&path[neighbor]);
                insert(h,neighbor,dst[neighbor]);
            }

            if(NewDistance > dst[neighbor] && NewDistance < dst_sec
ond[neighbor]){
                dst_second[neighbor]=NewDistance;
                sec_path[neighbor]=NewVertex;
                //The second shortest distance should also be put i
nto the heap.
                insert(h,neighbor,dst_second[neighbor]);
            }
        }
    }
}

```

```

        }

        p=p->next;
    }
}

return dst_second[dest];
}

stack CreateStack(int capacity){
    stack s=New(struct s,1);
    assert(s);
    s->top=0;
    s->capacity=capacity;
    s->container=New(element,capacity);
    assert(s->container);
    return s;
}

int ShortestPath(heap h,graph g,int dst[],int source,int dest,int path[
]){
    int vertex = source;
    int neighbor;
    int distance;
    int v=g->v;
    list p = NULL;

    int *include=New(int,v);//The array used to tell if an element has
    been popped from the heap.

    assert(include);

    SetSource(v,source,dst,path);
    memset(include,FALSE,sizeof(int)*v);//All elements are not included
    at the beginning.

    dispose(h);

    insert(h,source,dst[source]);

    while(Minimum(h,&distance,&vertex)){
        if(!include[vertex]){
            include[vertex]=TRUE;

```

```

        p=g->AdjList[vertex];
        while(p){
            neighbor = p->vertex;

            if(!include[neighbor]){

                if(dst[neighbor]>distance+p->weight){
                    dst[neighbor]=distance+p->weight;
                    path[neighbor]=vertex;
                    insert(h,neighbor,dst[neighbor]);
                }
            }
            p=p->next;
        }
    }

    free(include);
    return dst[dest];
}

```

## Declaration

*I hereby declare that all the work done in this project titled " The 2nd-shortest Path" is of my independent effort.*