

Unclonable states are necessary as proofs and advice

Rohit Chatterjee, **Srijita Kundu** & Supartha Podder

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3. There is a relational problem which can only be **exactly** solved by a quantum poly-time algorithm with an uncloneable quantum state as quantum **advice**.

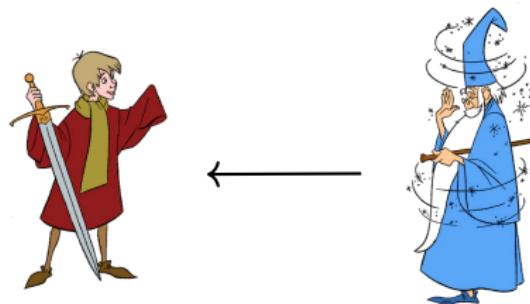
Quantum proofs

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Quantum generalization of NP

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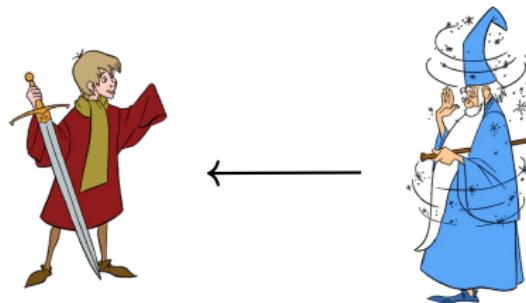
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- ▶ If YES instance, exists witness that can convince Arthur w.h.p.
- ▶ If NO instance, every witness is rejected by Arthur w.h.p.



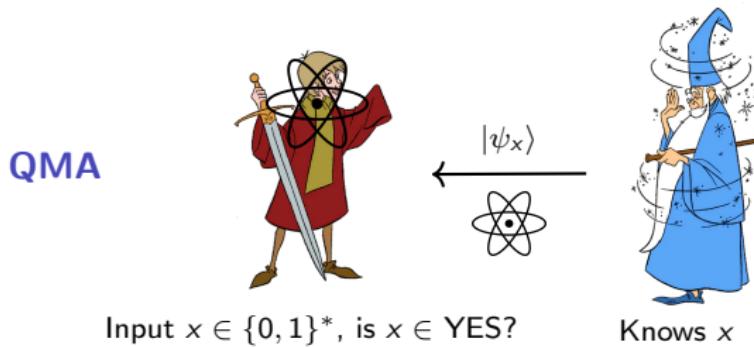
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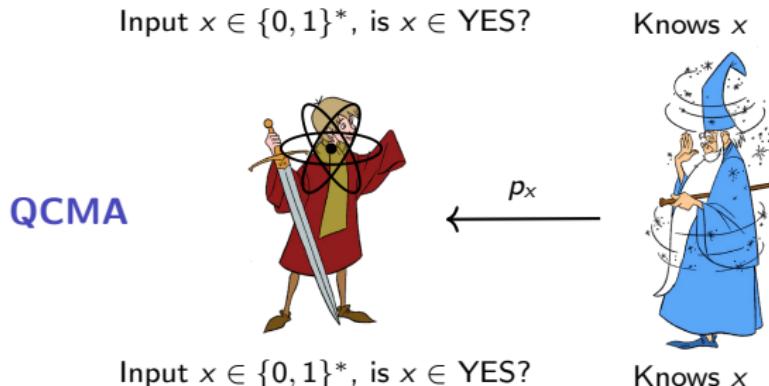
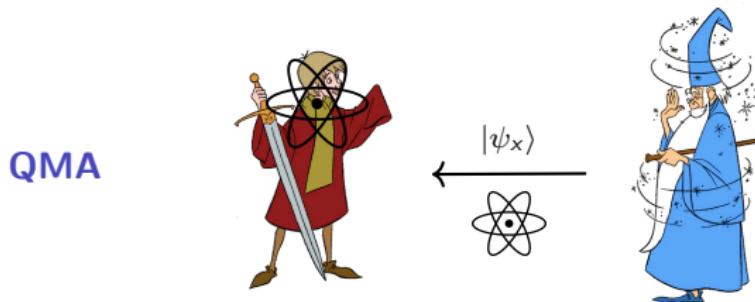
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- ▶ Classical proofs are always cloneable
- ▶ We use the [AK07](#) quantum oracle

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- ▶ **FBQP/qpoly**: like BQP/qpoly for polynomially bounded relations $R \subseteq \{0,1\}^* \times \{0,1\}^*, \dots$
- ▶ In fact the **ABK23** construction separates **FEQP/qpoly** and **FBQP/poly**
- ▶ We use the **ABK23** construction

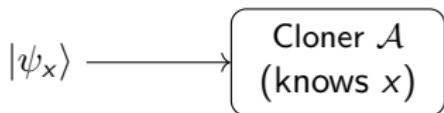
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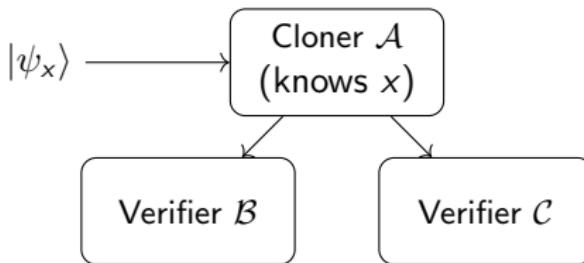
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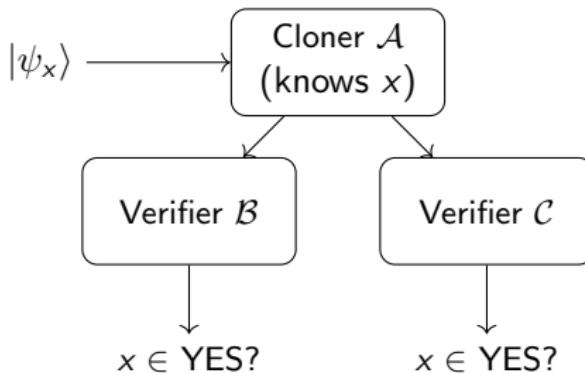
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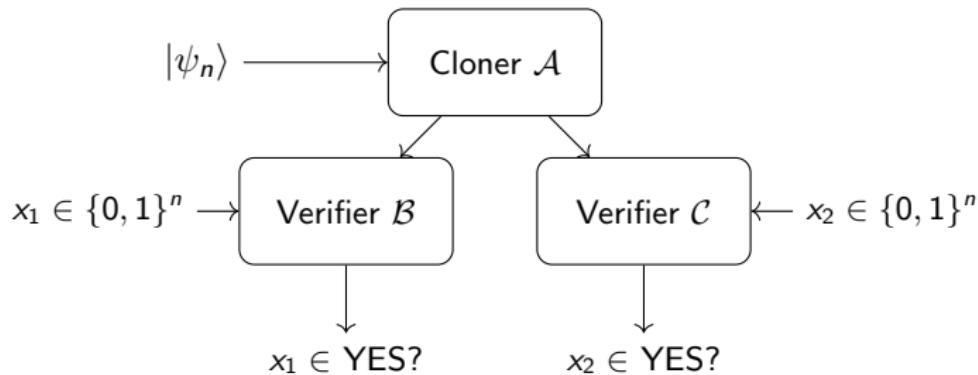
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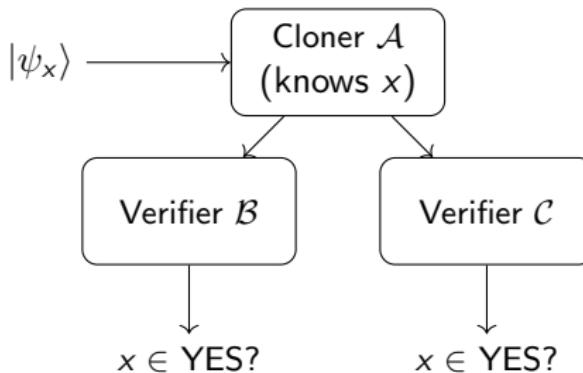
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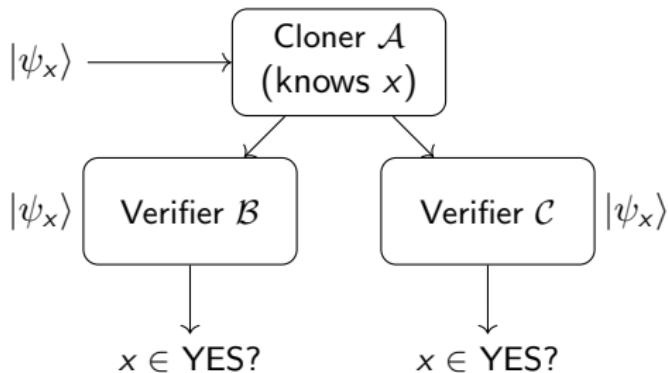
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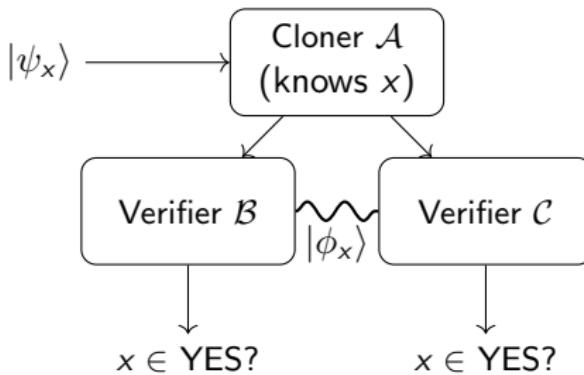
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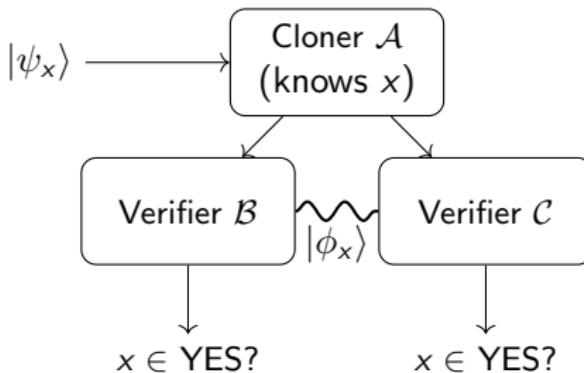
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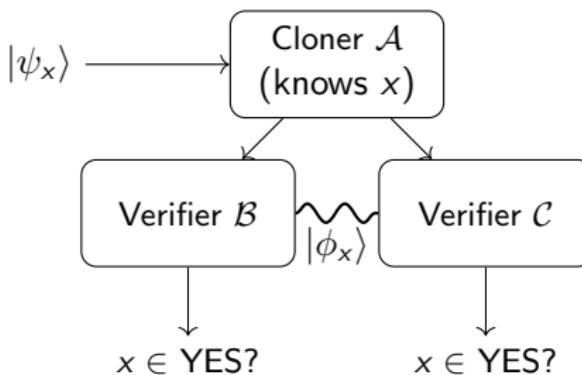
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- Often average-case uncloneability is considered

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- ▶ Uncloneability is **worst-case** instead of average case

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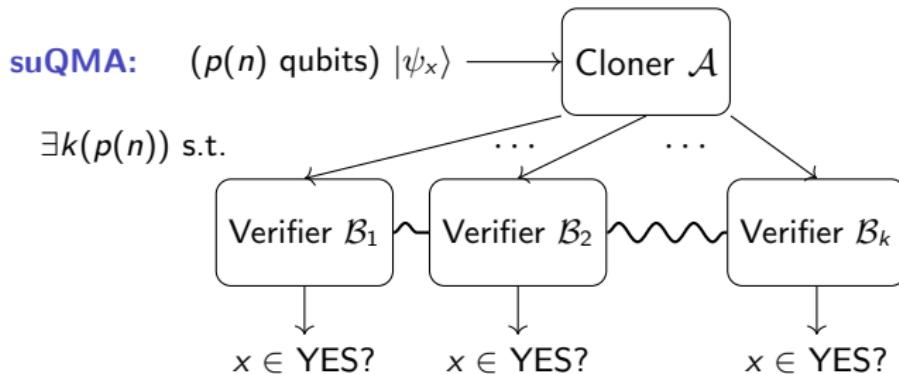
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$\text{QMA} \supseteq \text{suQMA}$:

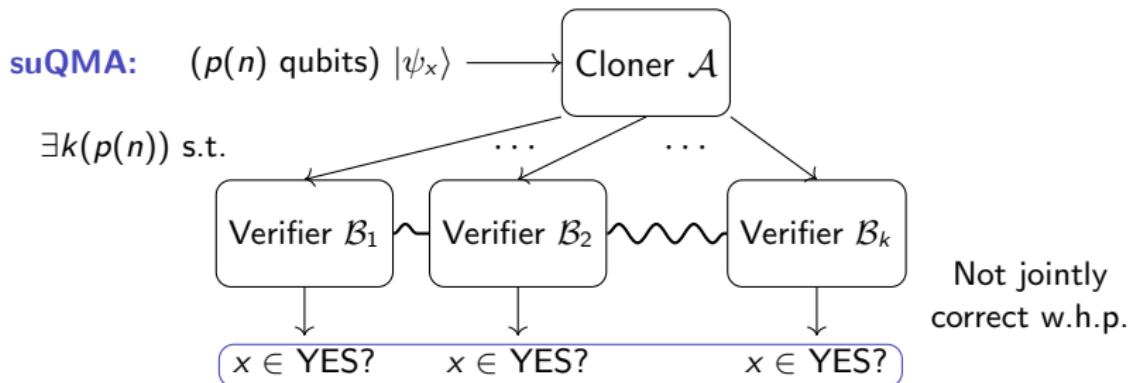
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suQMA: $(p(n) \text{ qubits}) |\psi_x\rangle$

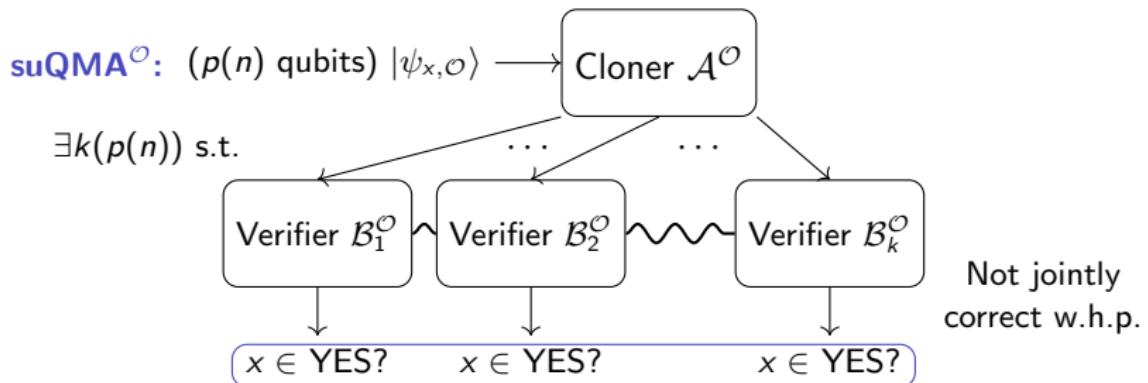
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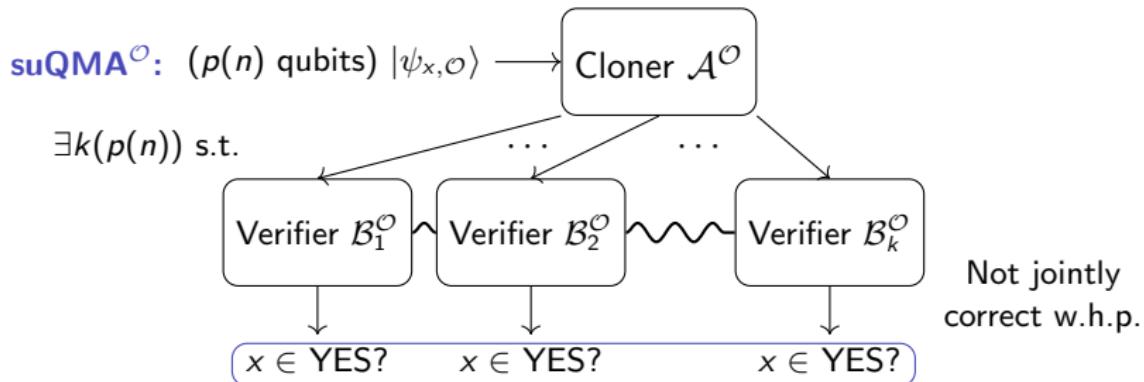
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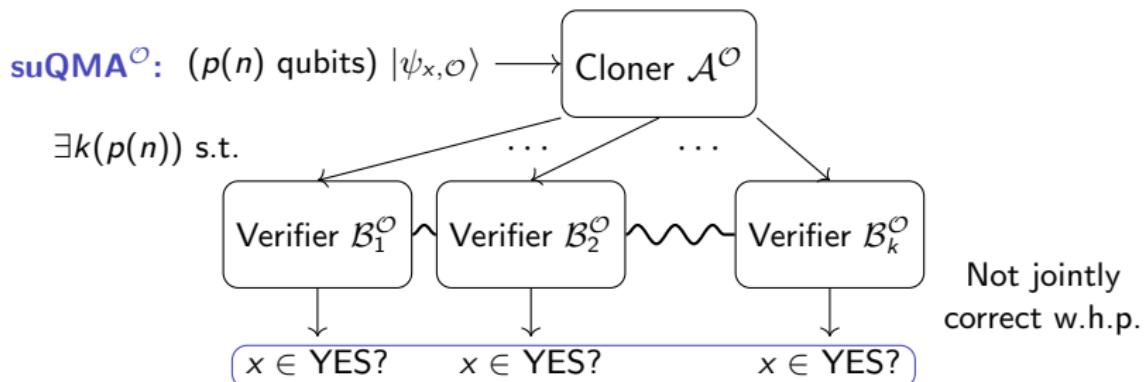


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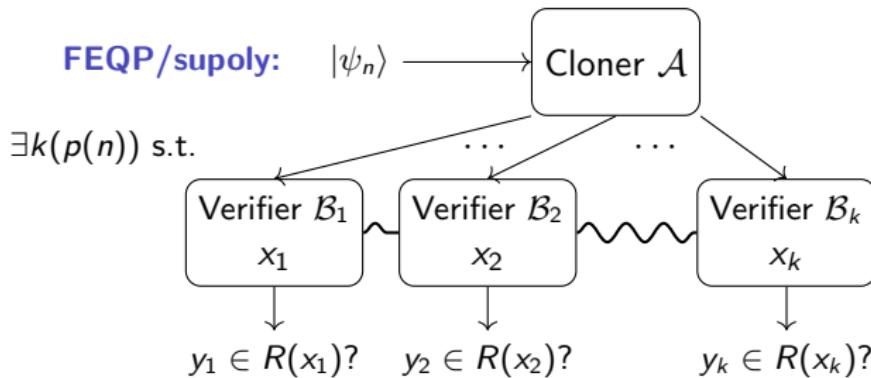
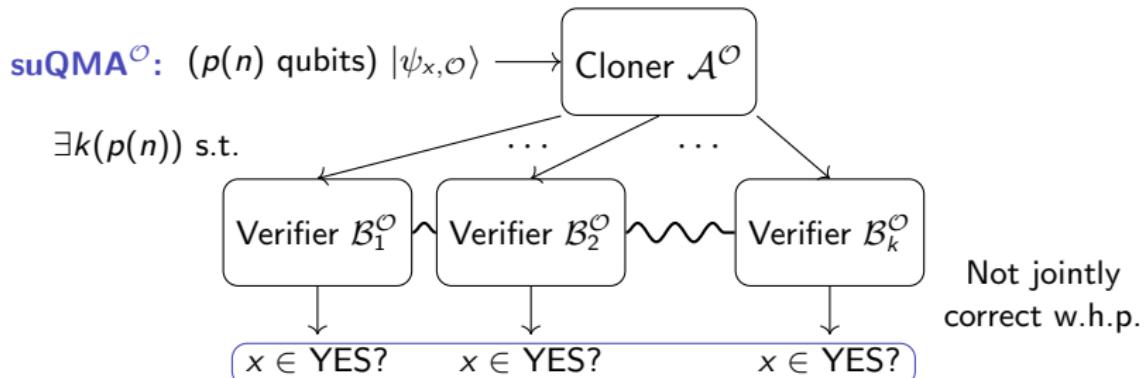
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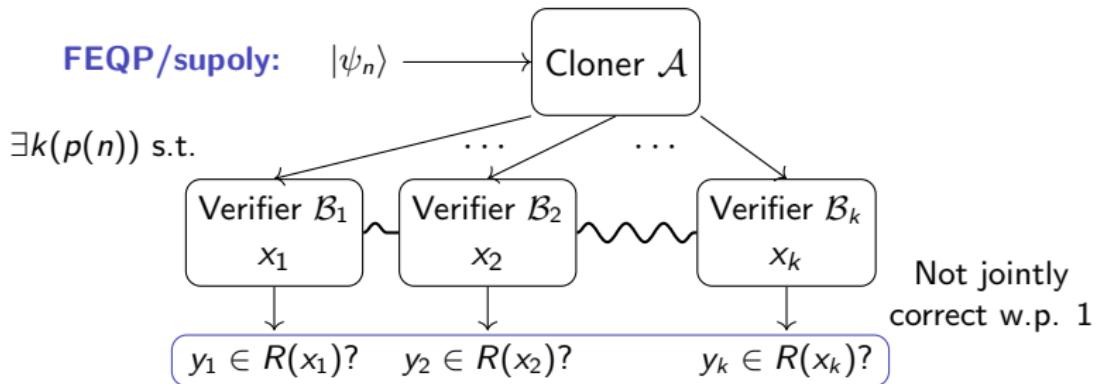
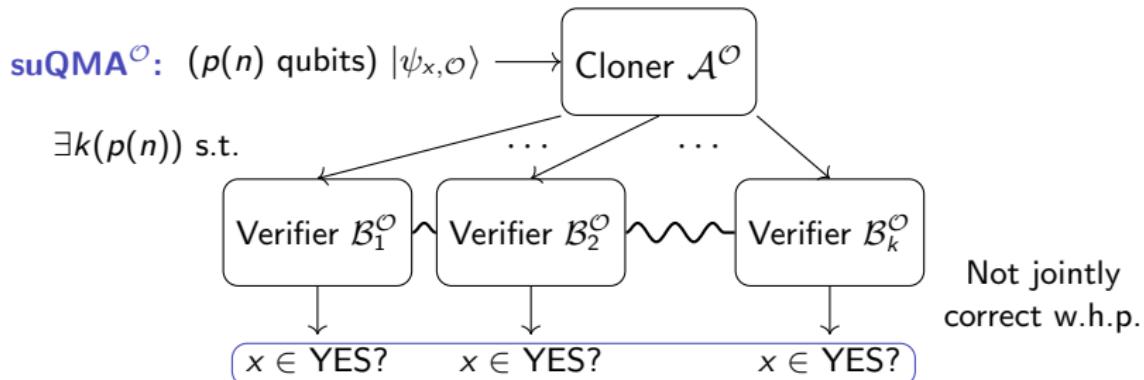


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2. Use uncloneability of canonical proof or advice

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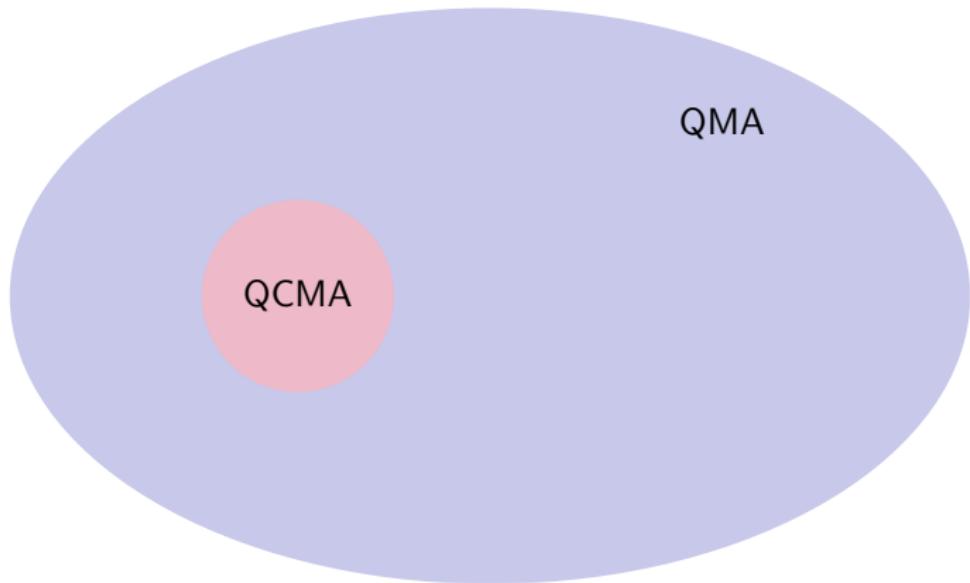
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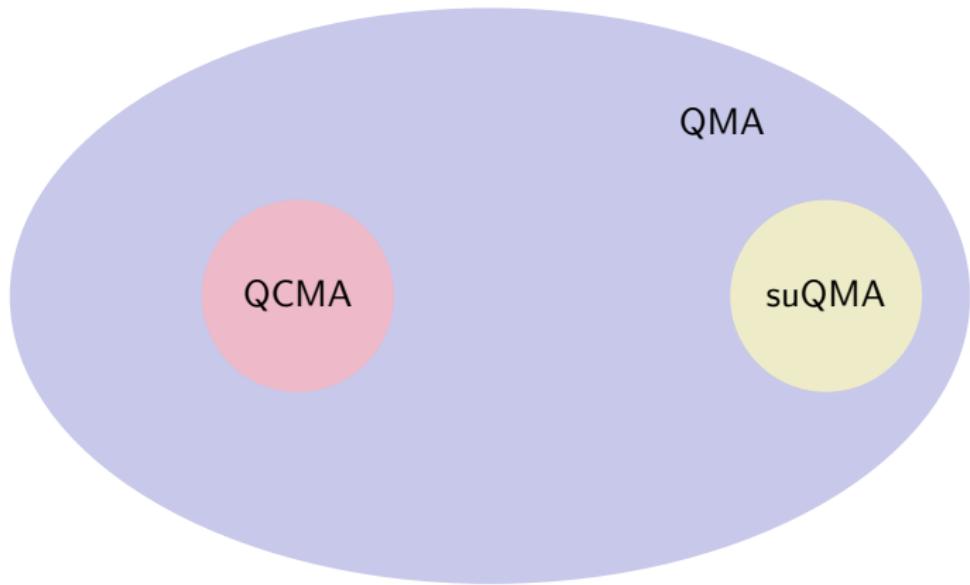
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- ▶ For FEQP/supoly, we show:
 - Proofs $|\psi_f\rangle$ that works jointly for $\mathcal{B}_1, \dots, \mathcal{B}_k$ must have inner product $|\langle\psi_{f'}|\psi_f\rangle| = 2^{-\Omega(k)}$;
 - In order to achieve the inner products, the states need $\Omega(k)$ qubits.

Other consequences

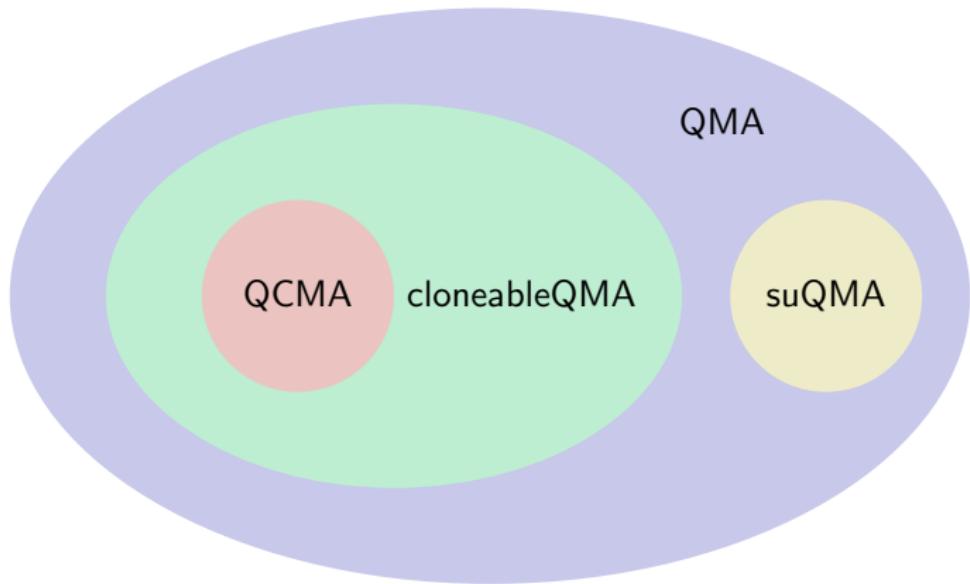
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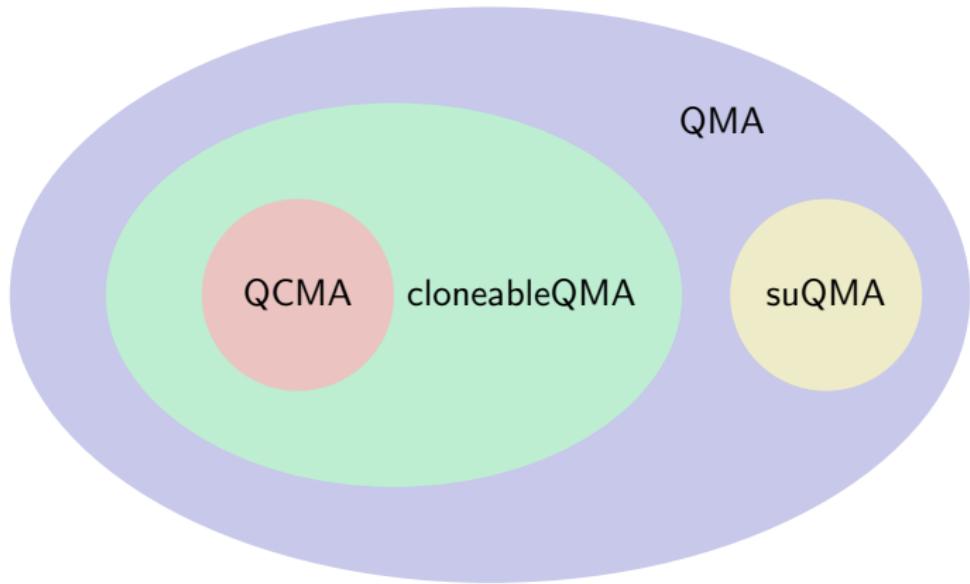


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From definition, $\text{suQMA} \subseteq \text{QMA} \setminus \text{cloneableQMA}$
⇒ w.r.t. a quantum oracle, $\text{QMA} \neq \text{cloneableQMA}$

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Thanks!