

Problem Definition

Given an integer array, circularly shift the array left side k times.

Example:

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Output:

{3, 4, 5, 1, 2}

Solution 1 (naïve)

shiftLeft (array, k):

if $k < 0$:

return -1

$n \leftarrow$ length of the array

for i from 0 to k :

$\text{temp} \leftarrow \text{array}[0]$

for j from 0 to $n-1$:

$\text{array}[j] \leftarrow \text{array}[j+1]$

end for

$\text{array}[n] \leftarrow \text{temp}$

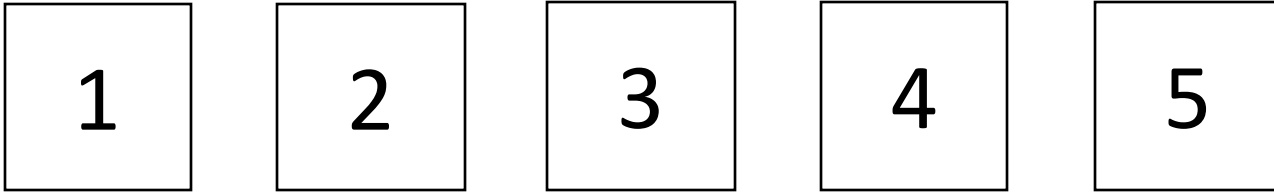
end for

return array

Solution 1 (naïve)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

1

2

3

4

5

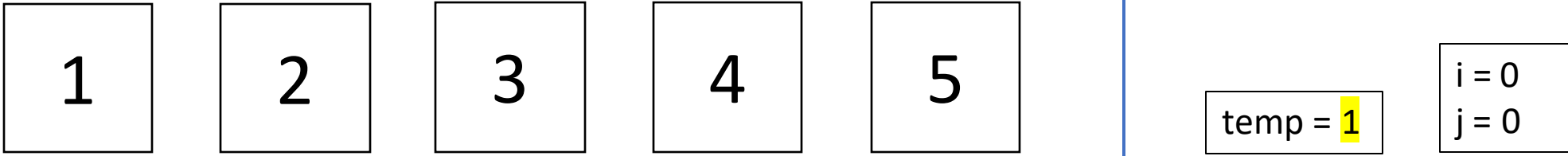


i = 0
j = 0

Solution 1 (naïve)

Inputs:

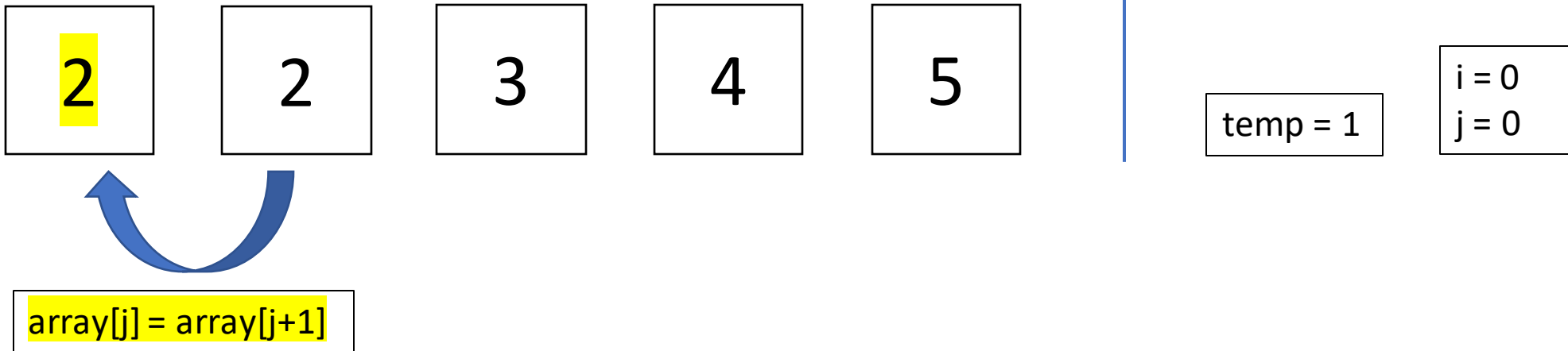
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

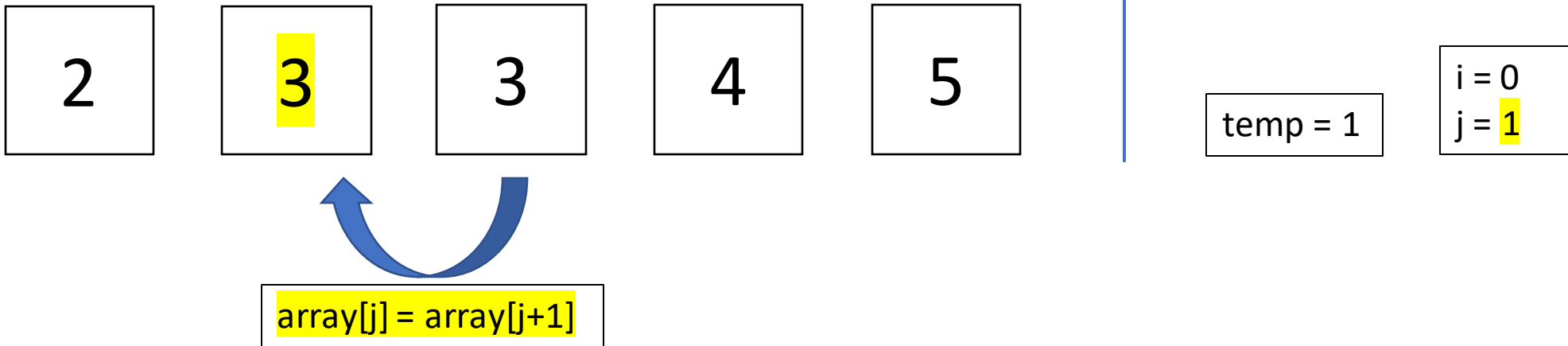
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

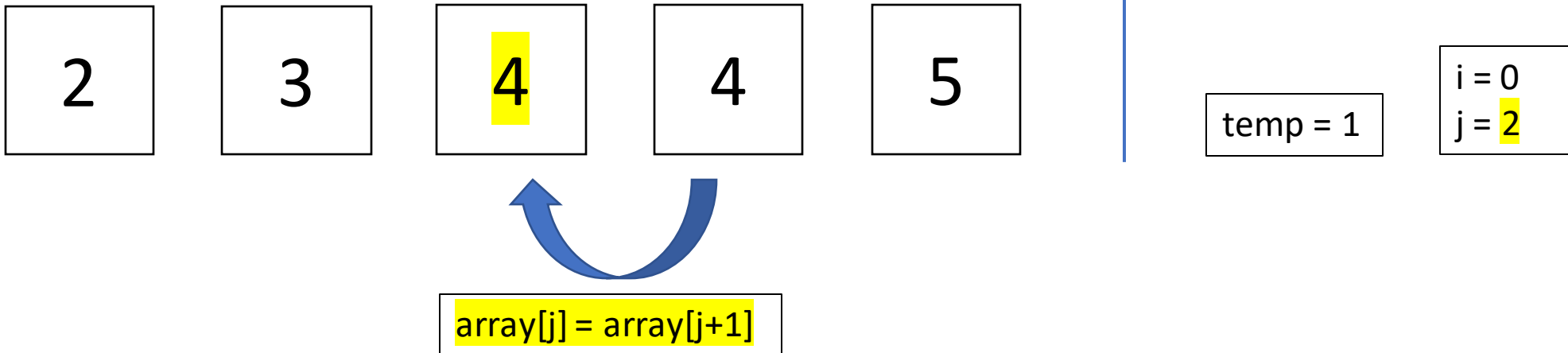
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

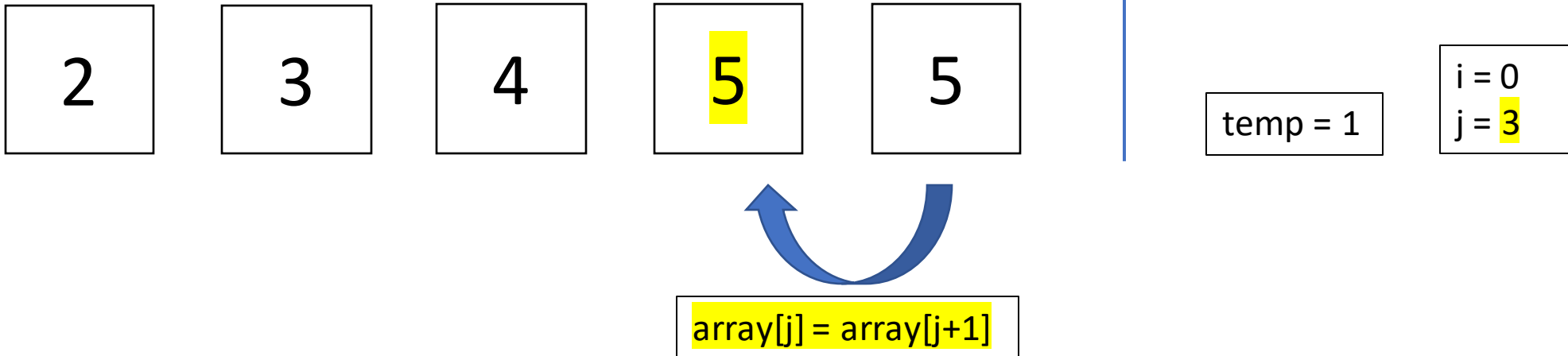
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

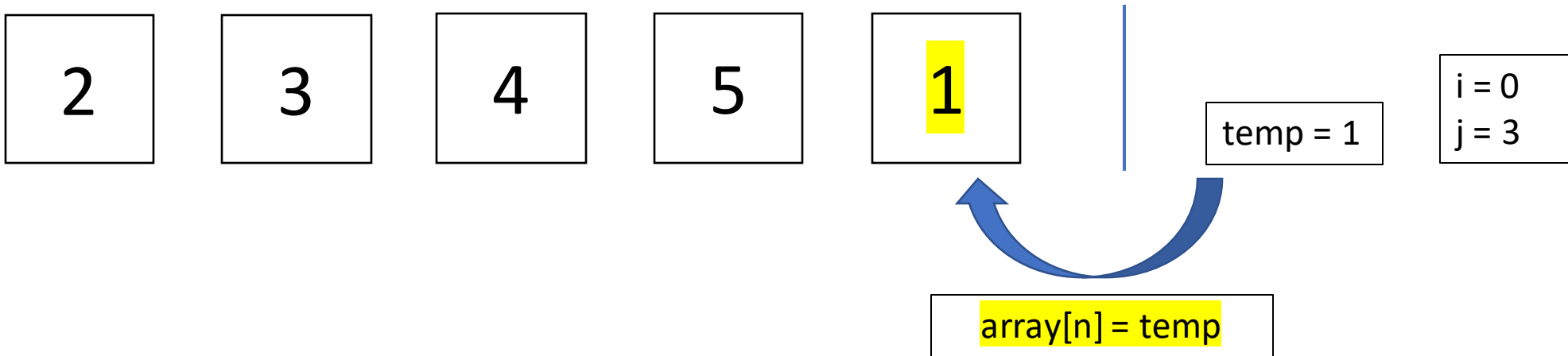
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

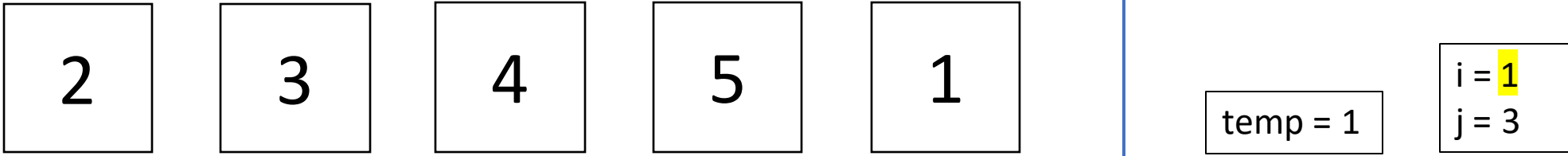
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

2

3

4

5

1

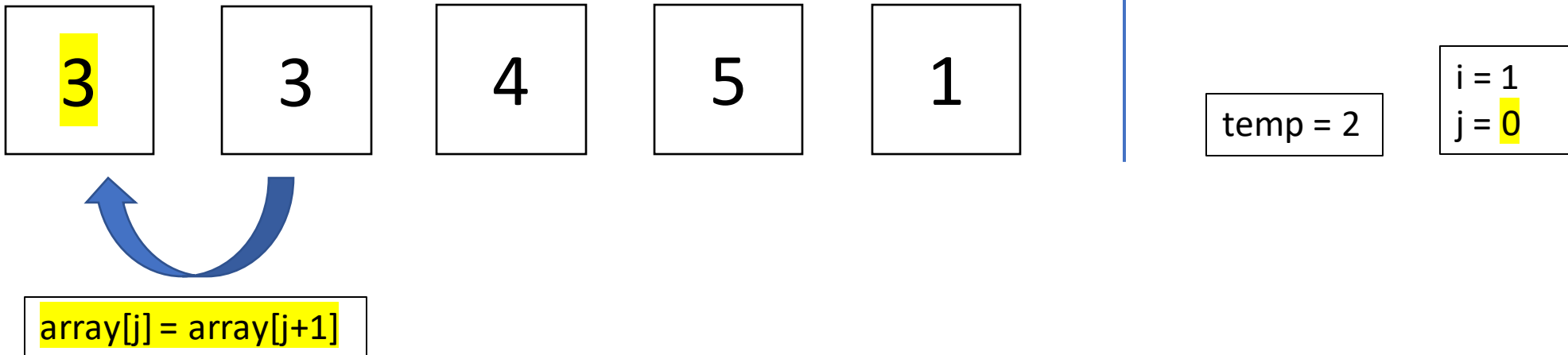
temp = 2

i = 1
j = 3

Solution 1 (naïve)

Inputs:

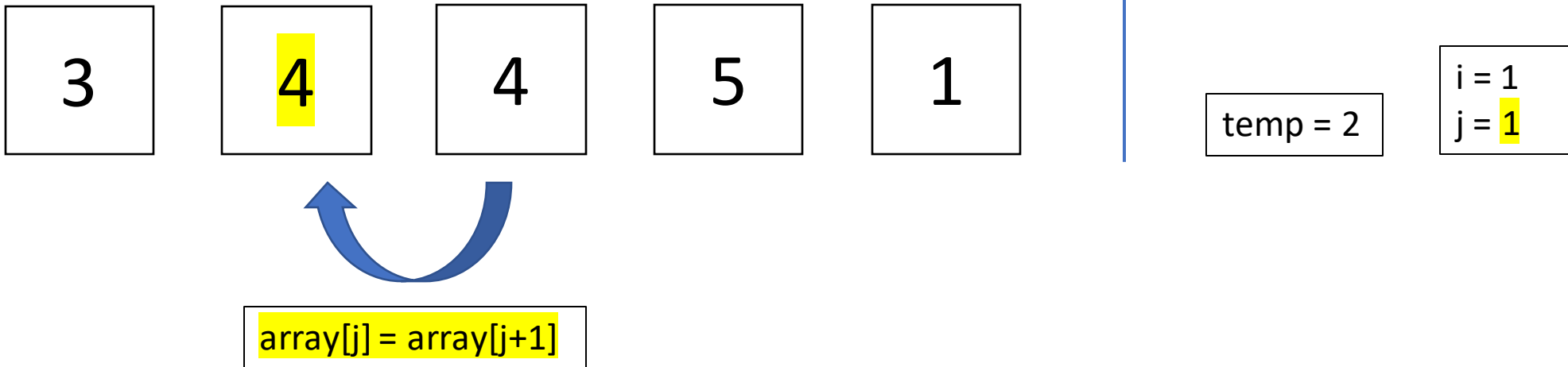
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

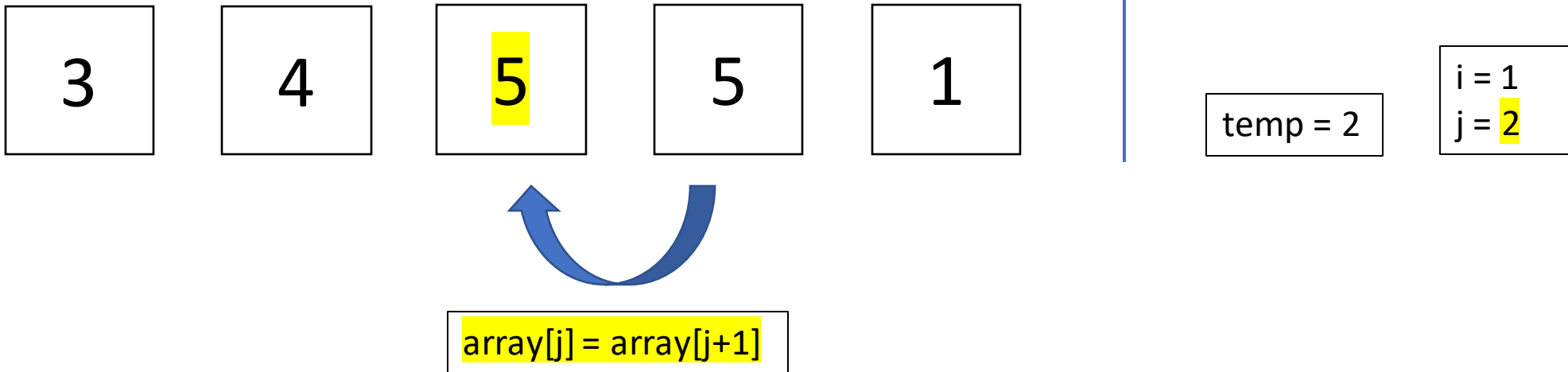
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

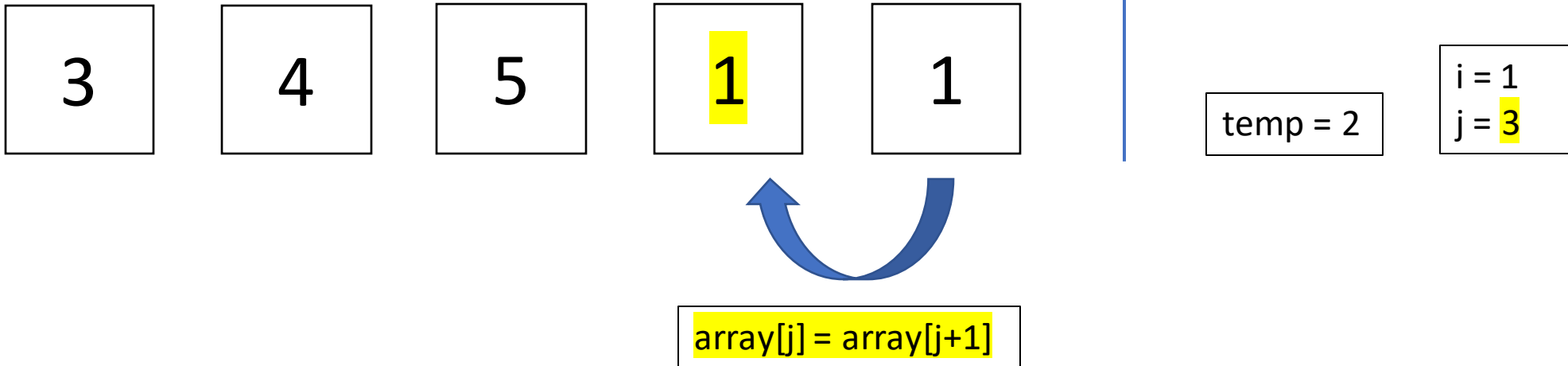
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

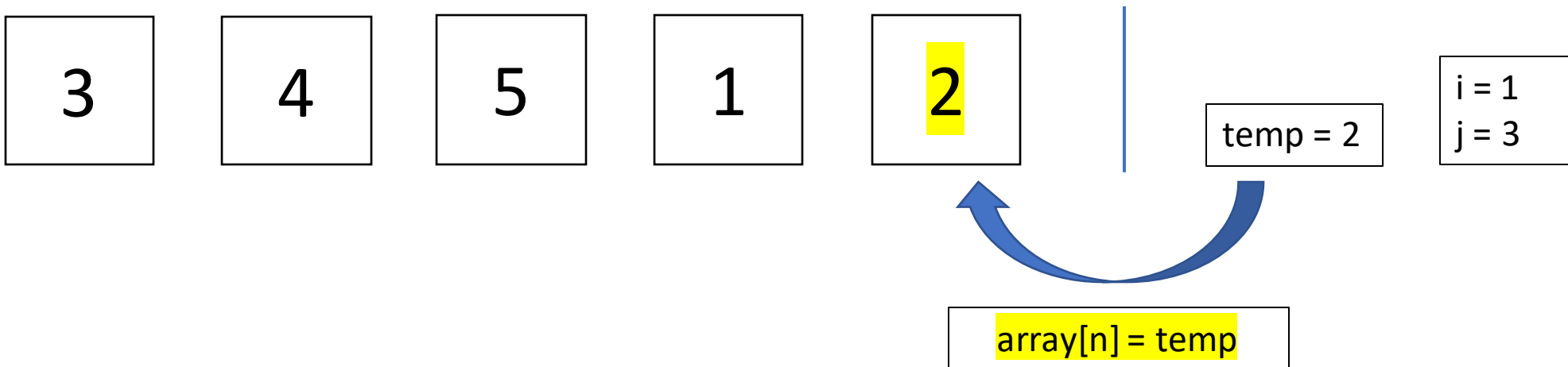
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

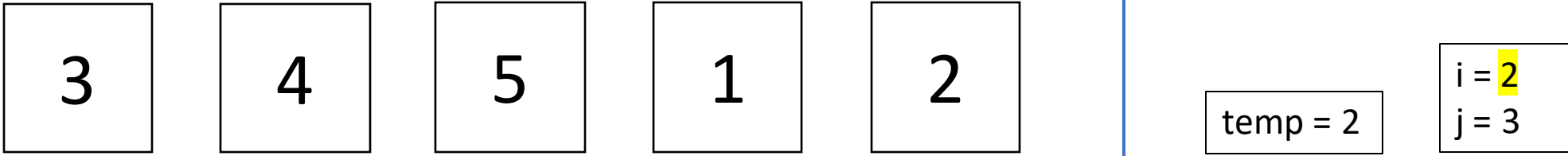
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

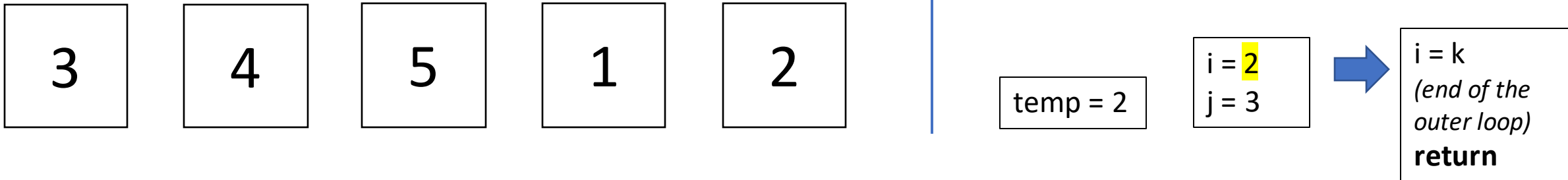
array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2



Solution 1 (naïve)

Time Complexity?

The best and the worst cases of the algorithm depend on **k**. Shifting left once requires n swapping operations. We have to repeat this k times, therefore the time complexity is $O(n*k)$.

Specifically:

- If $k = 0$, it means no shifting and therefore the best case time complexity becomes $\Omega(1)$.
- The array might be shifted to left at most $n-1$ times. Because *shifting it to left side 3 times or $n+3$ times forms the same output*. So, if we edit the outer loop's upper limit at $k\%n$, the loop will be iterated at most $n-1$ times. Which makes the worst case time complexity $O(n*(n-1)) = O(n^2)$. But if we don't edit the outer loop, complexity might grow drastically according to k .

Space Complexity?

The algorithm uses two iterators and a temporary variable. Space complexity is constant.

Solution 2 (auxiliary array)

```
shiftLeft(array, k):  
  if  $k \leq 0$ :  
    return -1  
  end if  
   $k \leftarrow k \% n$   
   $n \leftarrow \text{length of the array}$   
  aux_array  $\leftarrow []$   
  for i from 0 to k:  
    append array[i] to aux_array  
  end for  
  for i from k to n:  
    array[i-k]  $\leftarrow$  array[i]  
  end for  
  for i from n-k to n:  
    array[i]  $\leftarrow$  aux_array [i - (n - k)]  
  end for  
  return array
```

Solution 2 (auxiliary array)

shiftLeft(array, k):

if $k \leq 0$:

return -1

end if

$k \leftarrow k \% n$

$n \leftarrow \text{length of the array}$

 aux_array $\leftarrow []$

1st loop { **for** i from 0 to k:
 append array[i] to aux_array
 end for

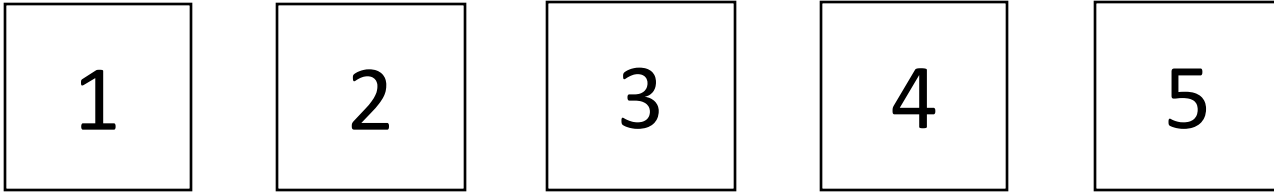
2nd loop { **for** i from k to n:
 array[i-k] \leftarrow array[i]
 end for

3rd loop { **for** i from n-k to n:
 array[i] \leftarrow aux_array [i - (n - k)]
 end for
 return array

Solution 2 (auxiliary array)

Inputs:

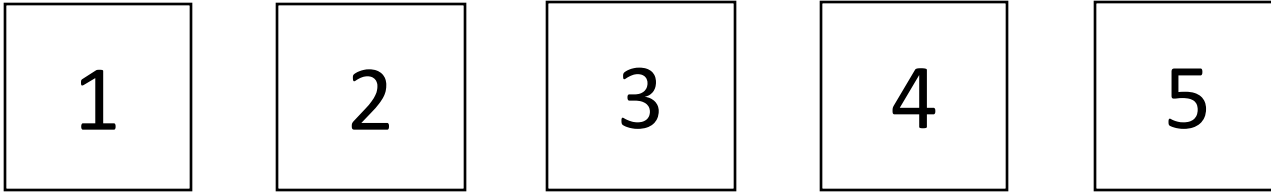
array = {1, 2, 3, 4, 5} and k = 2



Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2



aux_array = {}

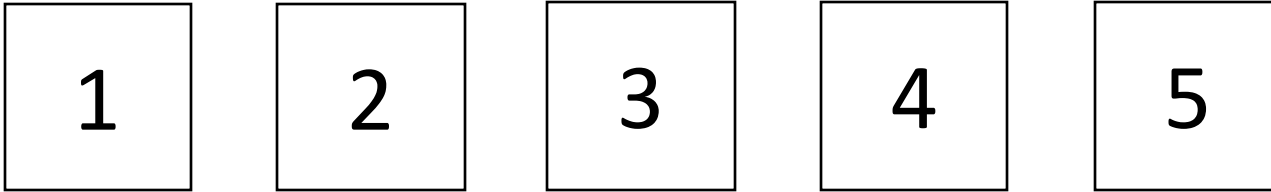
n = 5

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st loop:



aux_array = {}

n = 5

1st loop:

for i from 0 to k

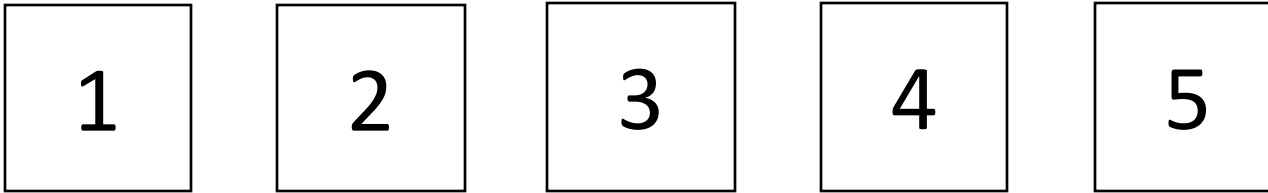
i = 0

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st loop:



aux_array = {1}

n = 5

1st loop:

for i from 0 to k

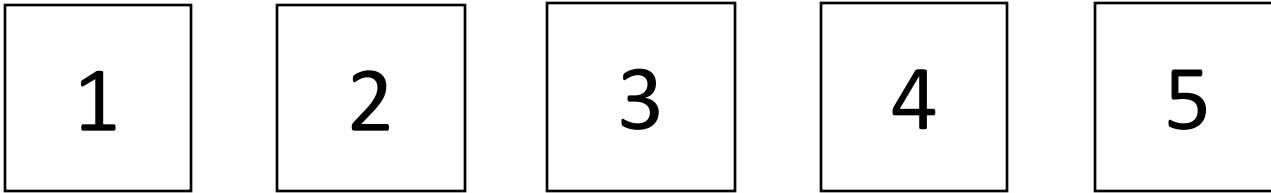
i = 0 → append array[0] to aux_array

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st loop:



aux_array = {1}

n = 5

1st loop:

for i from 0 to k

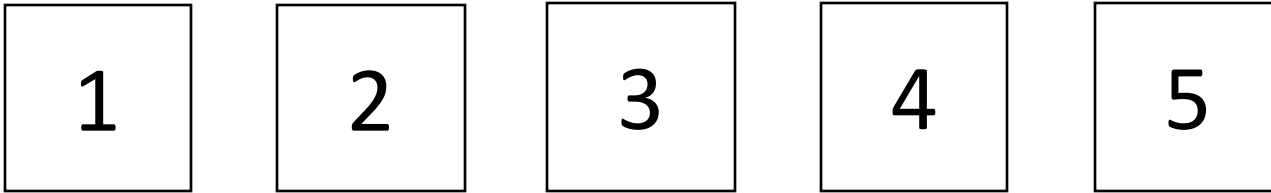
i= 1

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st loop:



aux_array = {1, 2}

n = 5

1st loop:

for i from 0 to k

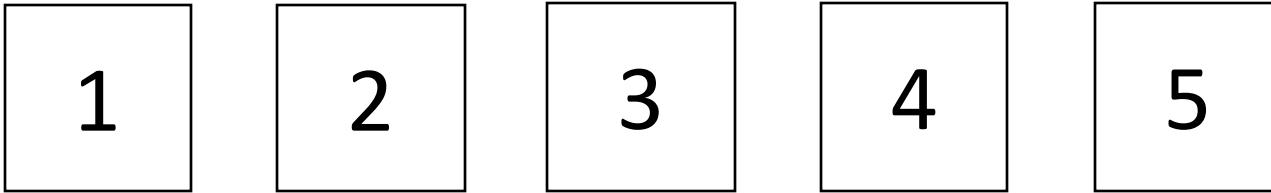
i = 1 → append array[1] to aux_array

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st loop:



aux_array = {1, 2}

n = 5

1st loop:

for i from 0 to k

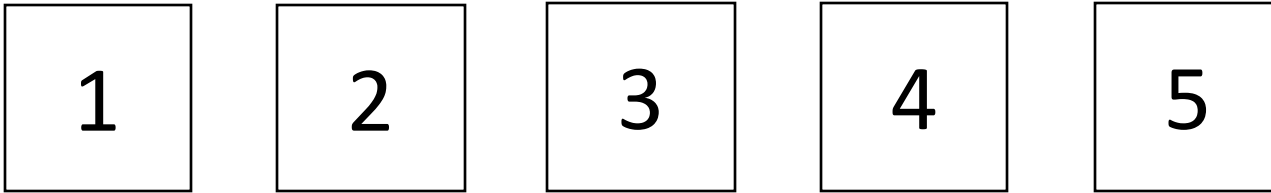
i= 2

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st loop:



aux_array = {1, 2}

n = 5

1st loop:

for i from 0 to k

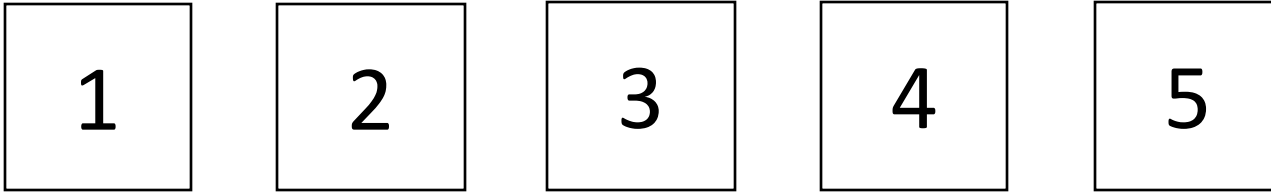
i = 2 → i = k → end of the loop

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

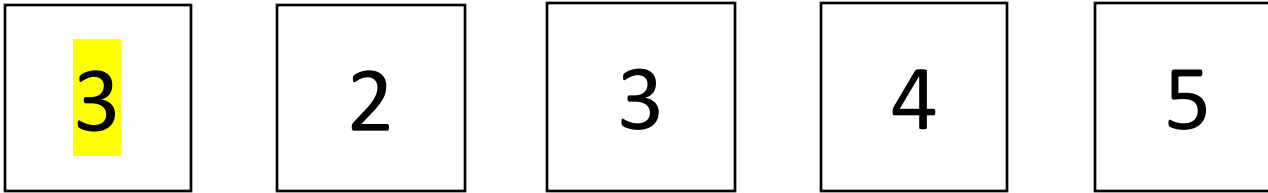
i= 2

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

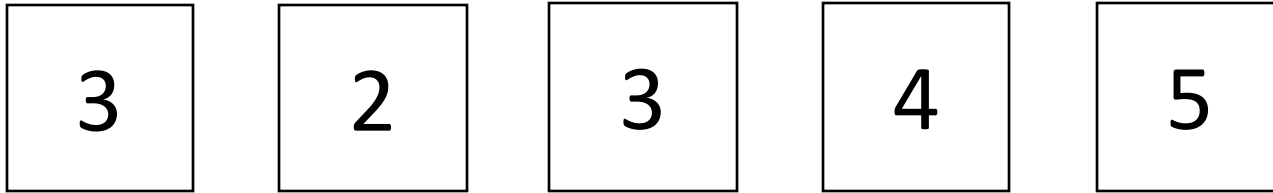
i = 2 → update array[i-k] as array[i]
(array[0] = array[2])

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

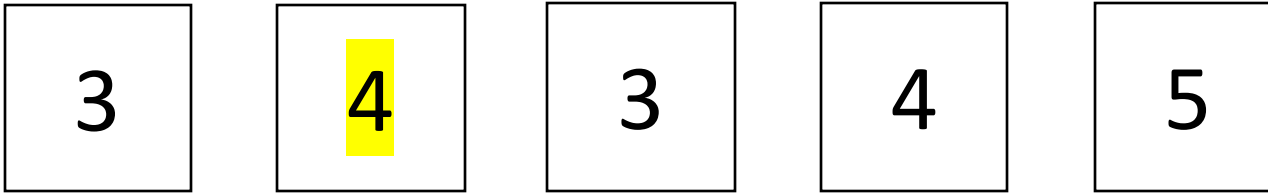
i= 3

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

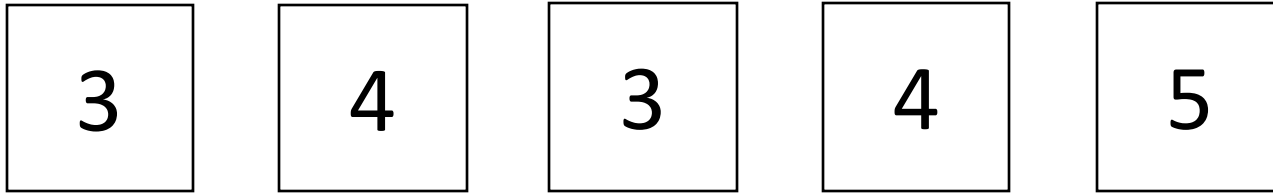
i = 3 → update array[i-k] as array[i]
(array[1] = array[3])

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

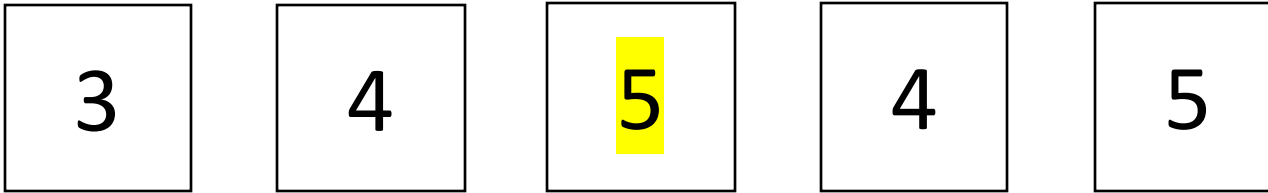
i= 4

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

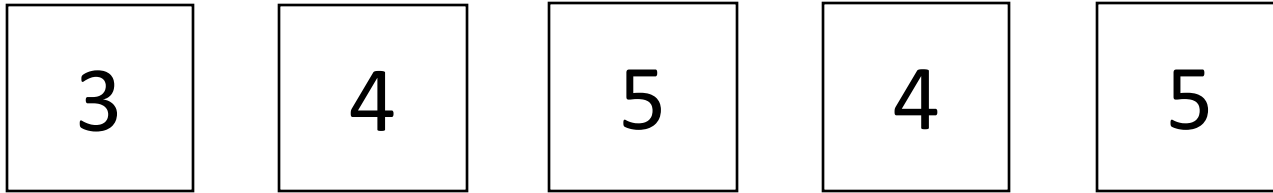
i = 4 → update array[i-k] as array[i]
(array[2] = array[4])

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

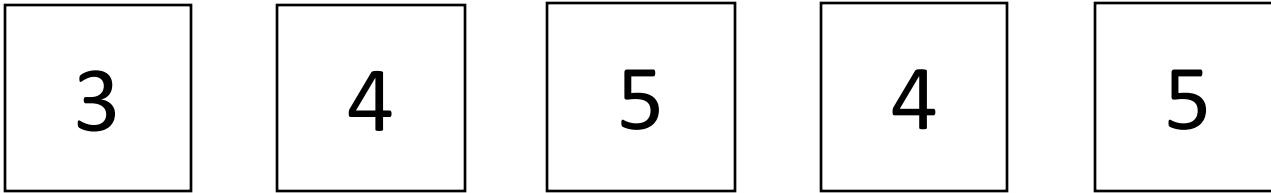
i= 5

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd loop:



aux_array = {1, 2}

n = 5

2nd loop:

for i from k to n

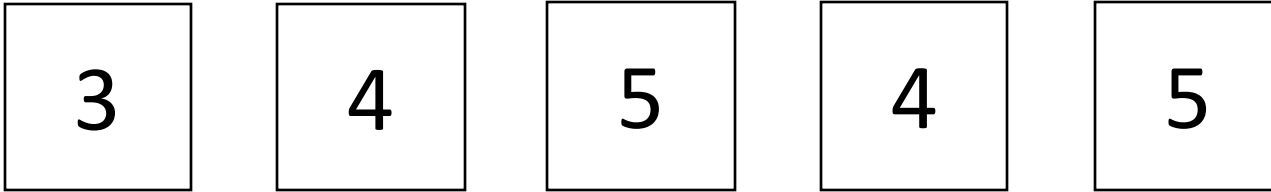
i = 5 → i = n → end of the 2nd loop

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd loop:



aux_array = {1, 2}

n = 5

3rd loop:

for i from n-k to n

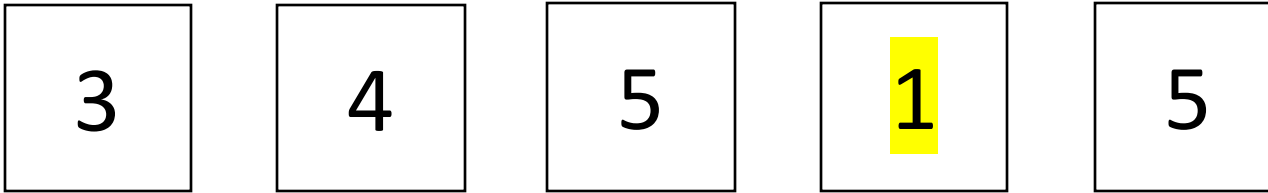
i= 3

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd loop:



aux_array = {1, 2}

n = 5

3rd loop:

for i from n-k to n

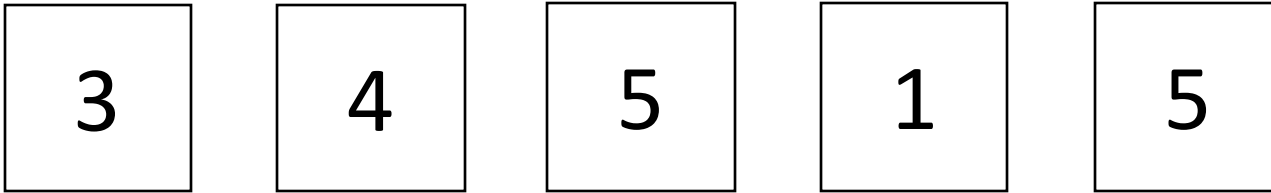
i = 3 → update array[i] as aux_array[i - (n - k)]
(array[3] = aux_array[0])

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd loop:



aux_array = {1, 2}

n = 5

3rd loop:

for i from n-k to n

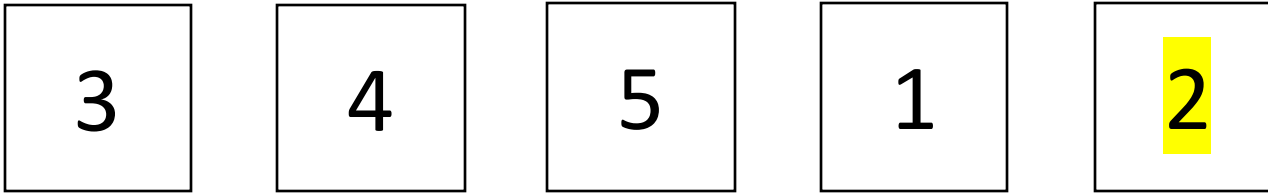
i= 4

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd loop:



aux_array = {1, 2}

n = 5

3rd loop:

for i from n-k to n

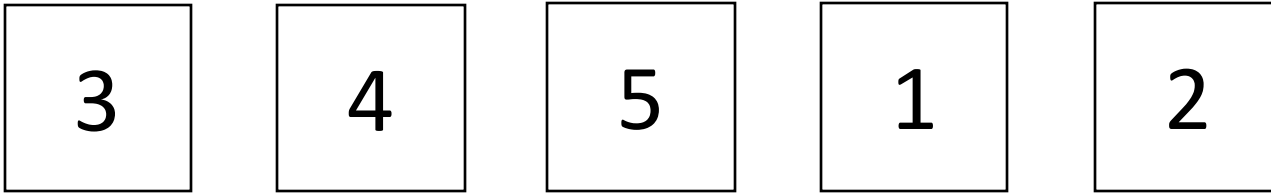
i = 4 → update array[i] as aux_array[i - (n - k)]
(array[4] = aux_array[1])

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd loop:



aux_array = {1, 2}

n = 5

3rd loop:

for i from n-k to n

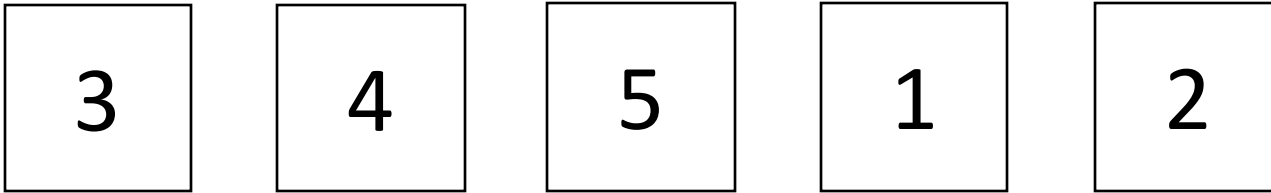
i= 5

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd loop:



aux_array = {1, 2}

n = 5

3rd loop:

for i from n-k to n

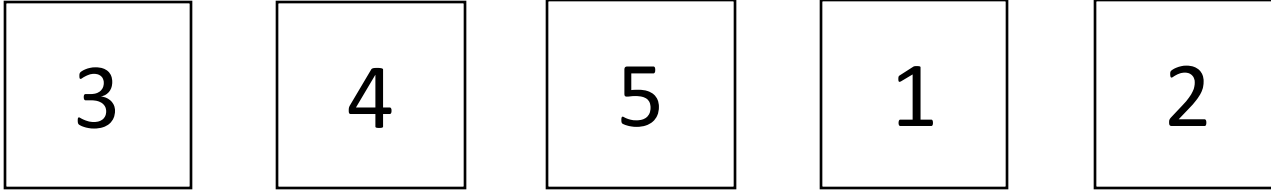
i = 5 → i = n → end of the 3rd loop

Solution 2 (auxiliary array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

The algorithm ends.



Solution 2 (auxiliary array)

Time Complexity?

The algorithm consists of 3 consecutive loops, each can have at most n iterations. If we analyze the algorithm we see that the array is divided into 2 distinct parts and first 2 loops work on these parts separately.

1st loop works on the first part which is `array[0:k]` # `array[0:k]` indicates the elements of the array between indexes 0 and k
2nd loop works on the second part which is `array[k:n]`

If we minimize the number of iterations of 1st loop, the 2nd array will get closer to n iterations. So, in the best case, we can make them both work with $n/2$ iterations. In this case the number of iterations of the 3rd loop becomes $n/2$ too. As a result, we end up with $\Omega(3 \cdot (n/2)) = \Omega(n)$.

Obviously, none of the loops takes more than $O(n)$ time. Therefore, the worst case is the same as the best case: $O(n)$.

Space Complexity?

The algorithm uses an additional array and an iterator. The array's size is k and it dominates the size of the iterator. Therefore, the space complexity is $O(k)$.

Solution 3 (reversing the array)

shiftLeft (array, k):

if $k \leq 0$:

return -1

end if

$k \leftarrow k \% n$

$n \leftarrow \text{length of the array}$

 array \leftarrow reverse(array, 0, k-1)

 array \leftarrow reverse(array, k, n-1)

 array \leftarrow reverse(array, 0, n-1)

return array

reverse(array, left, right)

while left < right:

swap (array[left] and array[right])

 left \leftarrow left + 1

 right \leftarrow right - 1

end while

return array

Solution 3 (reversing the array)

shiftLeft (array, k):

if $k \leq 0$:

return -1

end if

$k \leftarrow k \% n$

$n \leftarrow$ length of the array

 array \leftarrow reverse(array, 0, k-1) } **1st reverse operation**

 array \leftarrow reverse(array, k, n-1) } **2nd reverse operation**

 array \leftarrow reverse(array, 0, n-1) } **3rd reverse operation**

return array

reverse(array, left, right)

while left < right:

swap (array[left] and array[right])

 left \leftarrow left + 1

 right \leftarrow right - 1

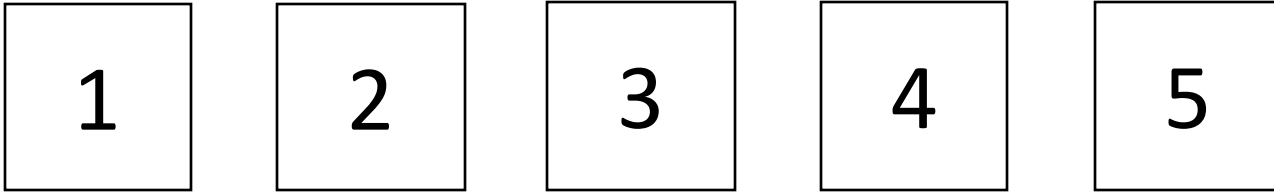
end while

return array

Solution 3 (reversing the array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

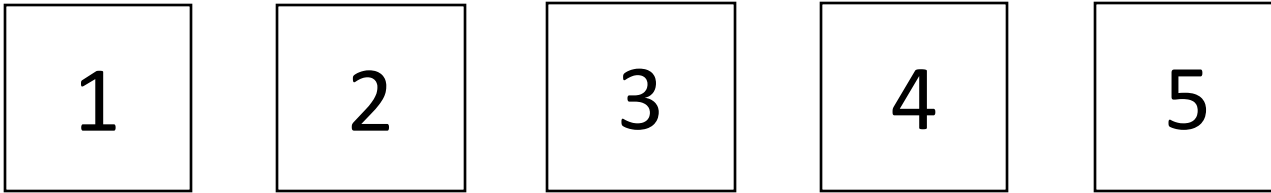


Solution 3 (reversing the array)

Inputs:

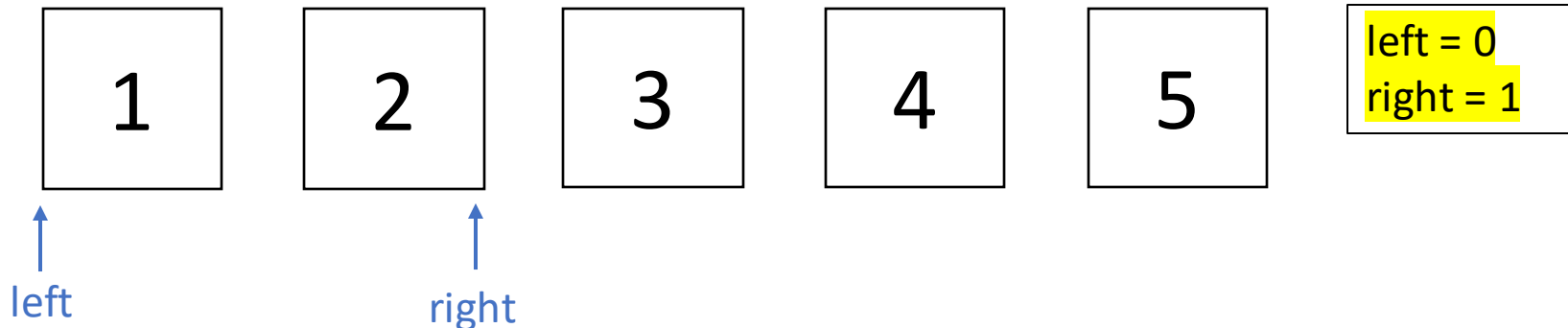
array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st reverse operation:



function call:

↓ reverse(array, 0 , k-1) ↓

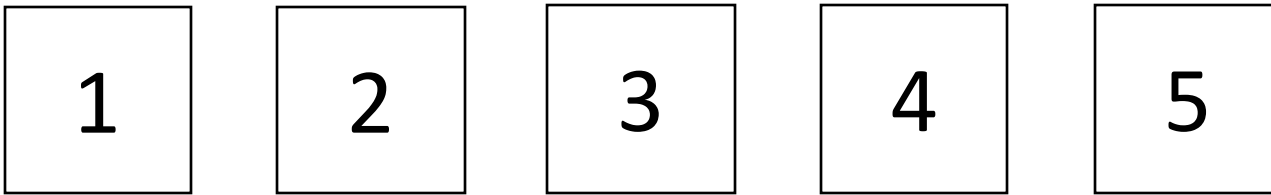


Solution 3 (reversing the array)

Inputs:

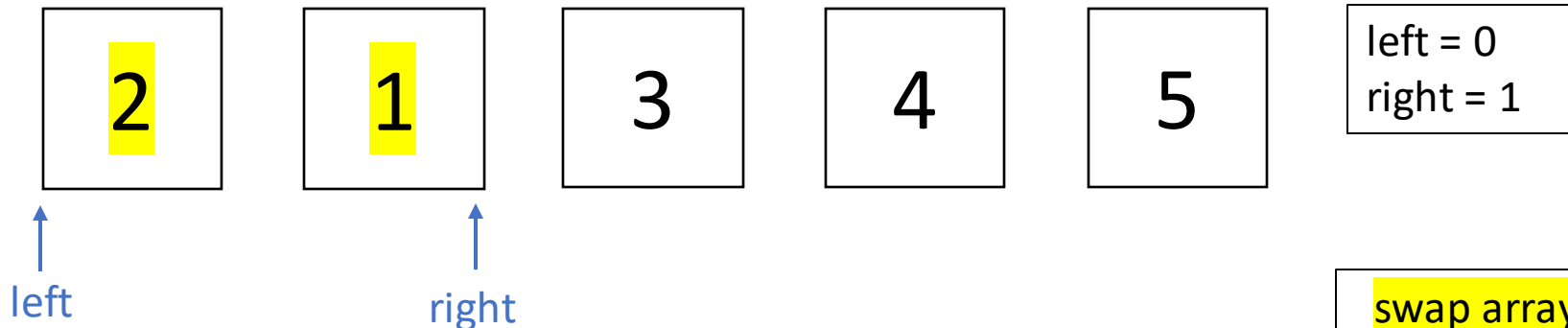
array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st reverse operation:



function call:

↓ reverse(array, 0, k-1) ↓



left = 0
right = 1

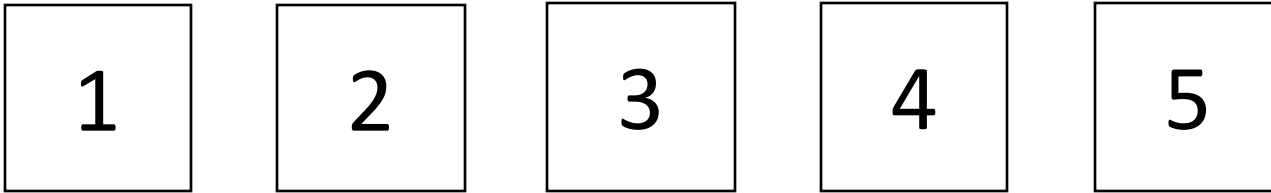
swap array[left] and array[right]

Solution 3 (reversing the array)

Inputs:

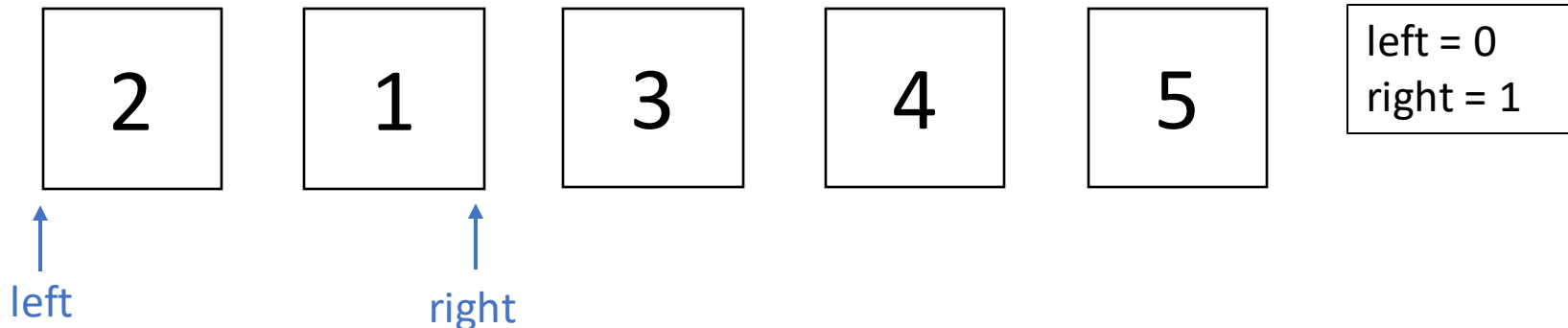
array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st reverse operation:



function call:

↓ reverse(array, 0, k-1) ↓

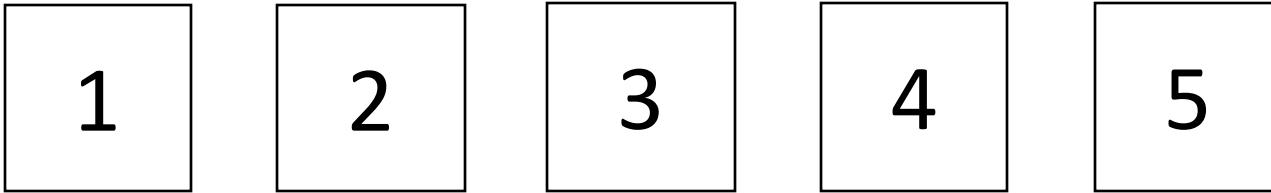


Solution 3 (reversing the array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

Implementing 1st reverse operation:



function call:

↓ reverse(array, 0, k-1) ↓



↑
right

↑
left

left = 1
right = 0

→ left <? right
False → end of the while loop
return array

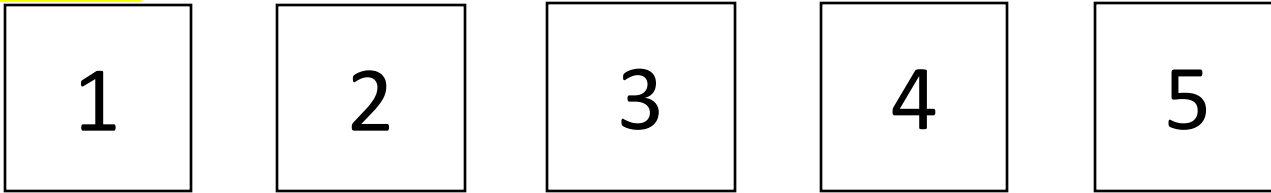
Solution 3 (reversing the array)

Inputs:

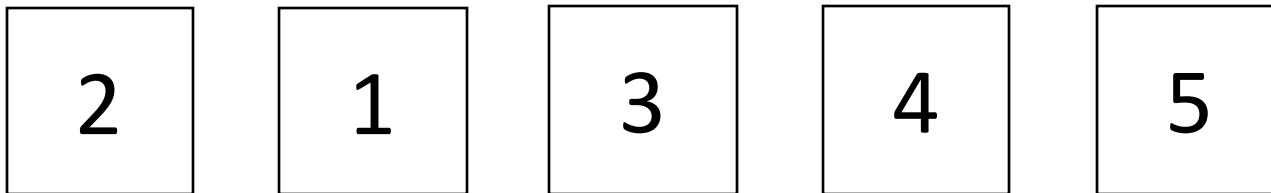
array = {1, 2, 3, 4, 5} and k = 2

Array is updated after the 1st reverse operation:

Before:



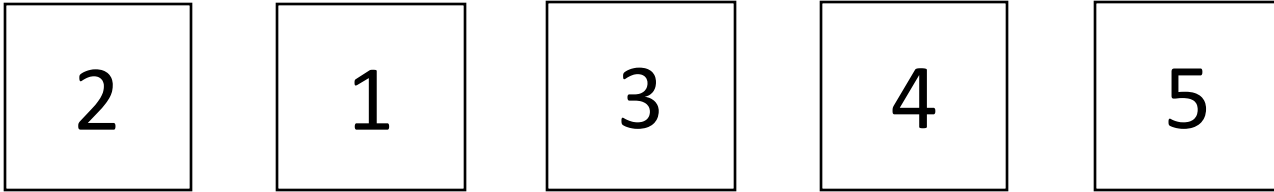
After:



Solution 3 (reversing the array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

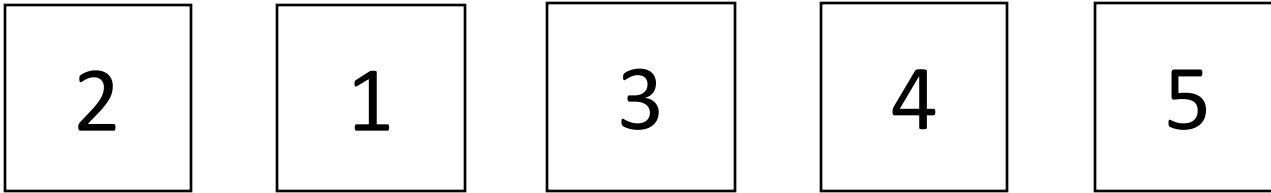


Solution 3 (reversing the array)

Inputs:

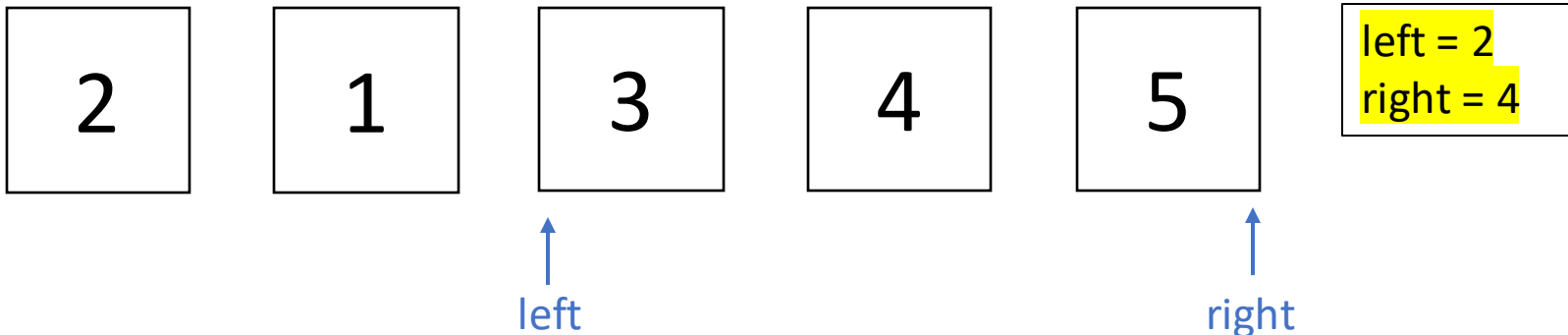
array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd reverse operation:



function call:

↓ reverse(array, k, n-1) ↓ (n = 5)

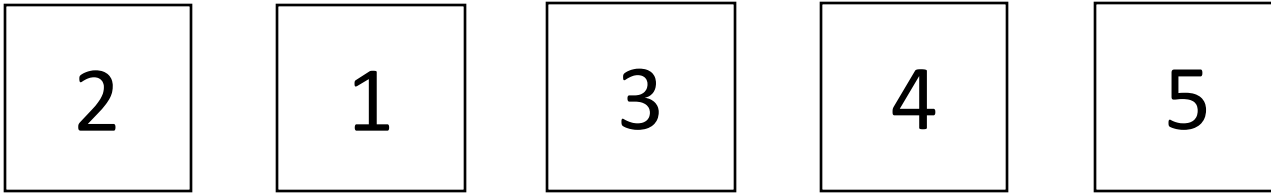


Solution 3 (reversing the array)

Inputs:

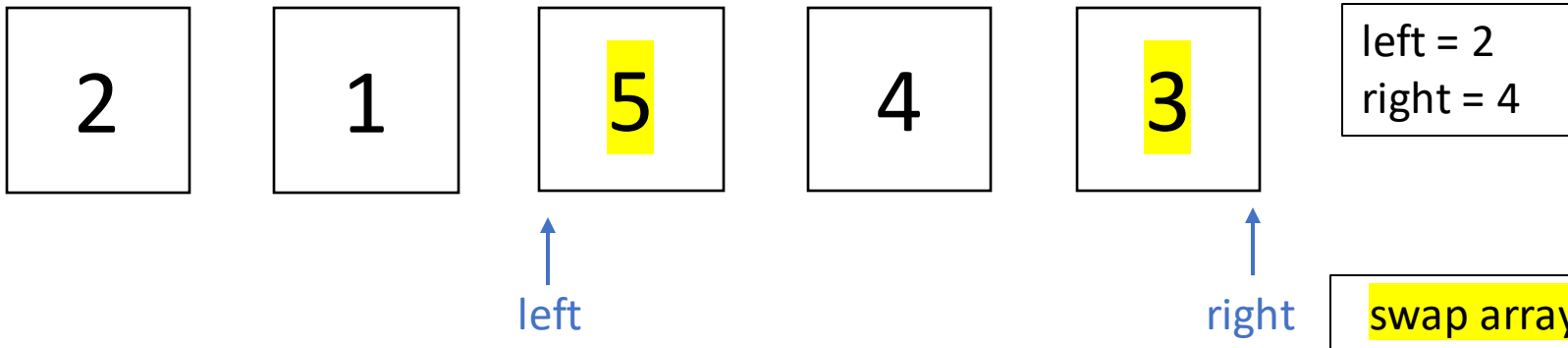
array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd reverse operation:



function call:

↓ reverse(array, k , n-1) ↓ (n = 5)

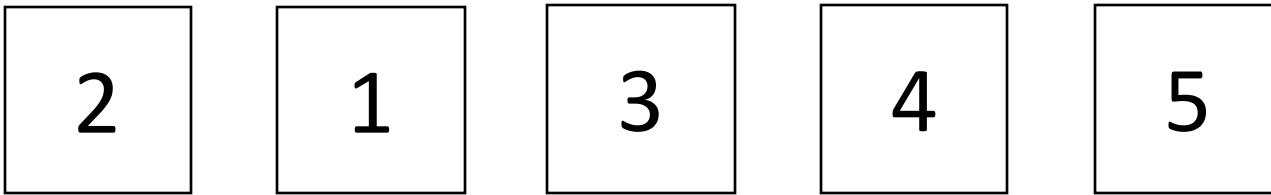


Solution 3 (reversing the array)

Inputs:

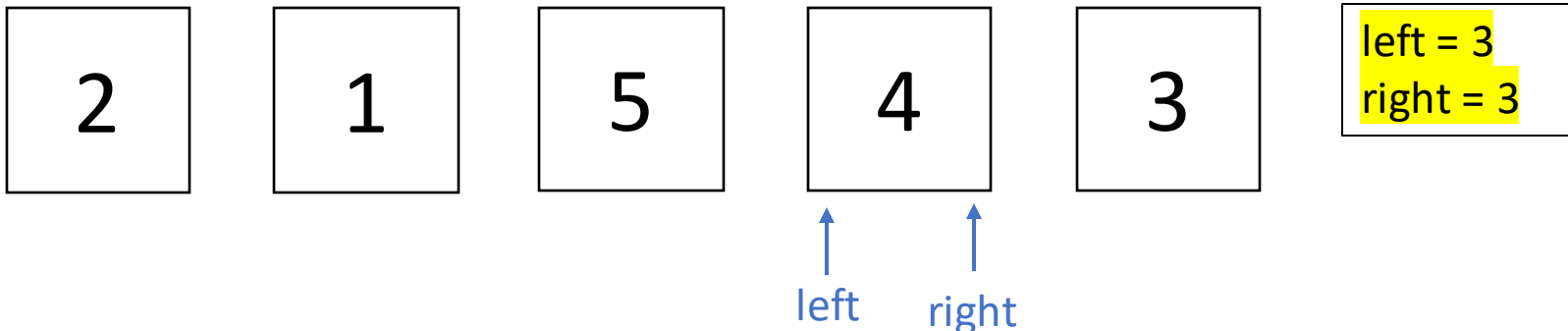
array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd reverse operation:



function call:

↓ reverse(array, k, n-1) ↓ (n = 5)

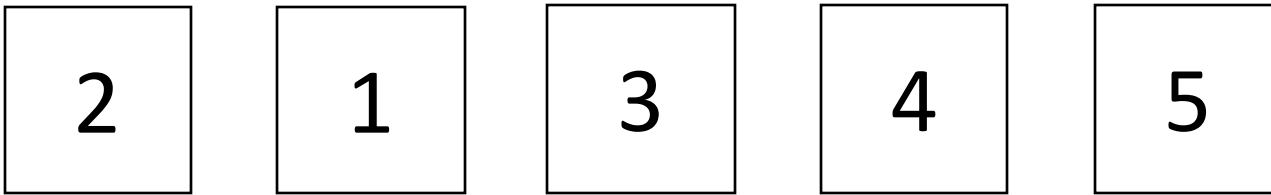


Solution 3 (reversing the array)

Inputs:

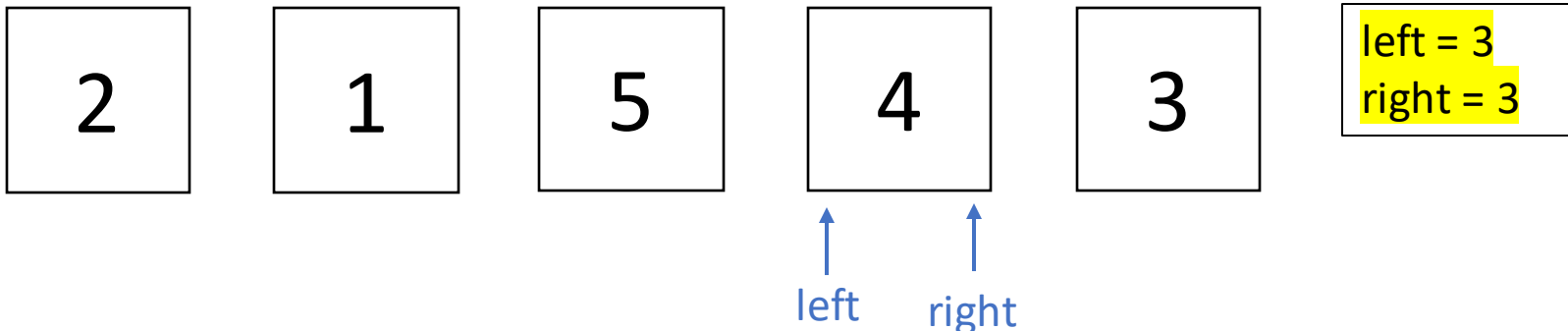
array = {1, 2, 3, 4, 5} and k = 2

Implementing 2nd reverse operation:



function call:

↓ reverse(array, k , n-1) ↓ (n = 5)



→ left <? right
False → end of the while loop
return array

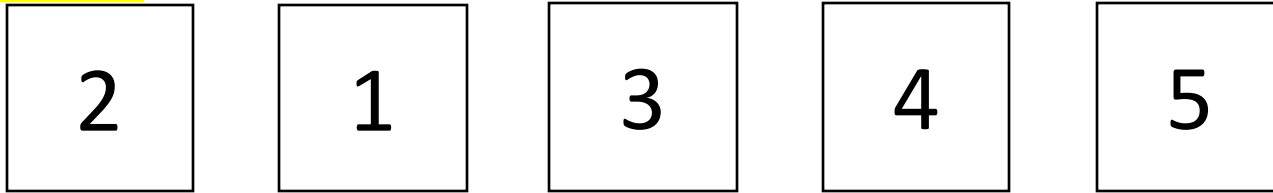
Solution 3 (reversing the array)

Inputs:

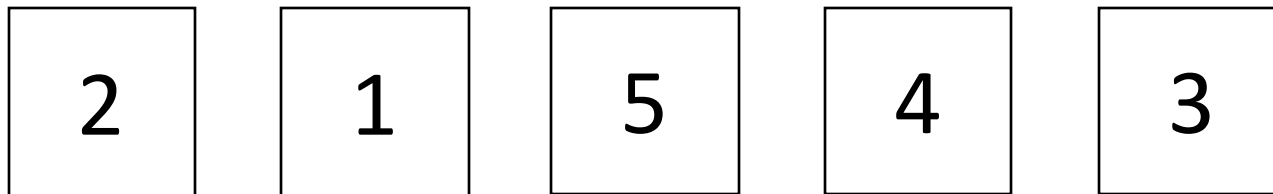
array = {1, 2, 3, 4, 5} and k = 2

Array is updated after the 2nd reverse operation:

Before:



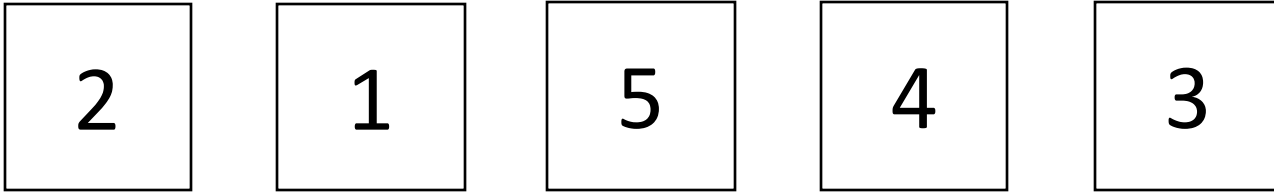
After:



Solution 3 (reversing the array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

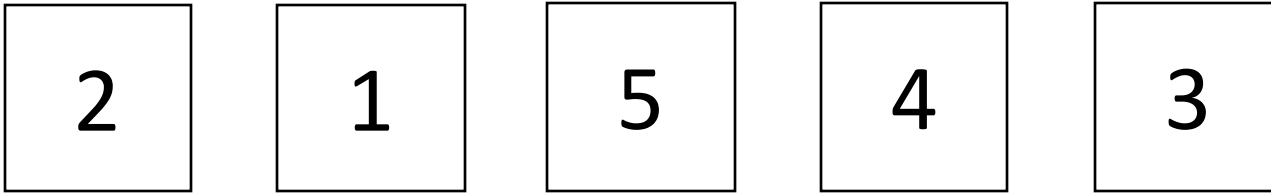


Solution 3 (reversing the array)

Inputs:

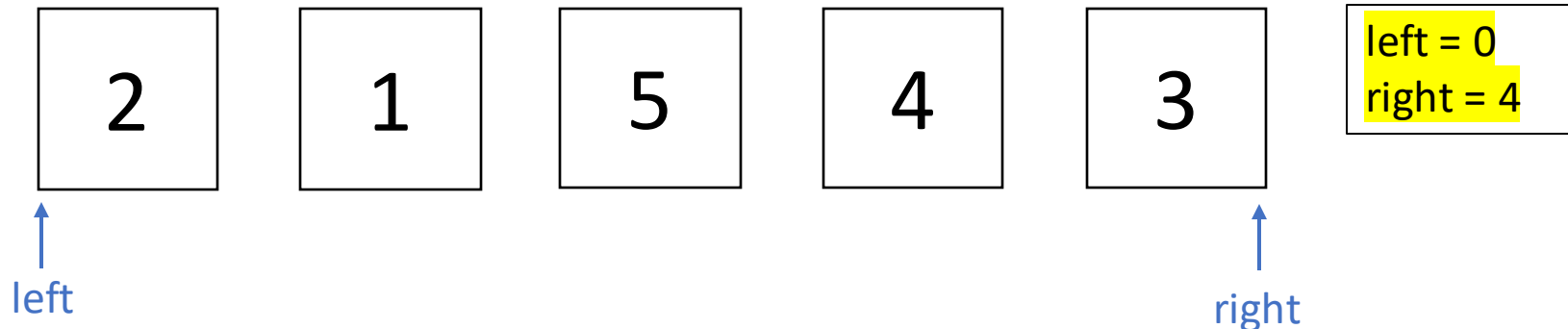
array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd reverse operation:



function call:

↓ reverse(array, 0, n-1) ↓ (n = 5)

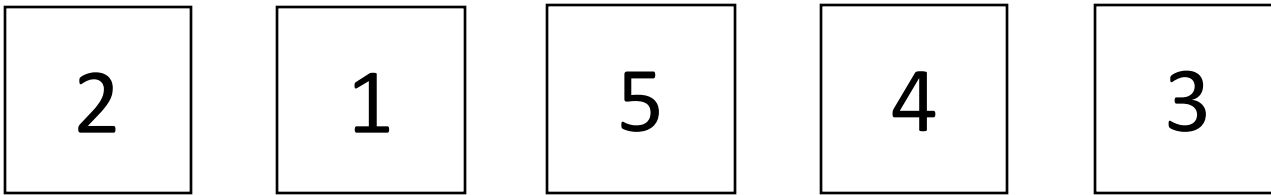


Solution 3 (reversing the array)

Inputs:

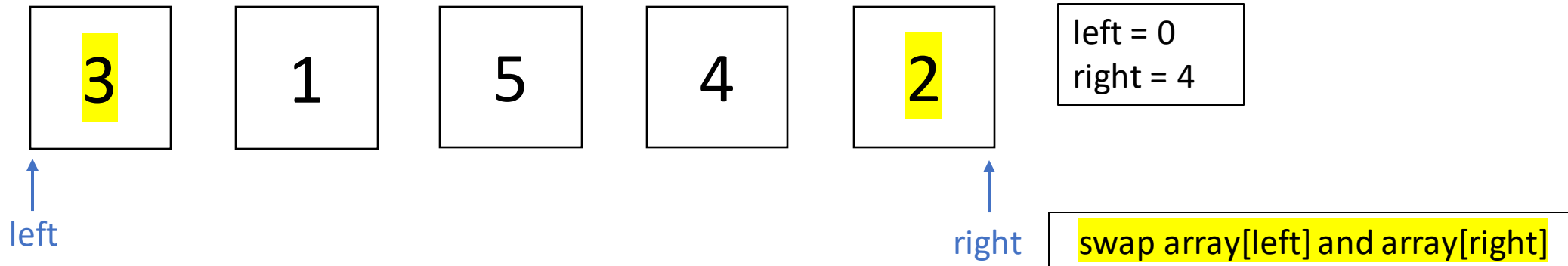
array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd reverse operation:



function call:

↓ reverse(array, 0, n-1) ↓ (n = 5)

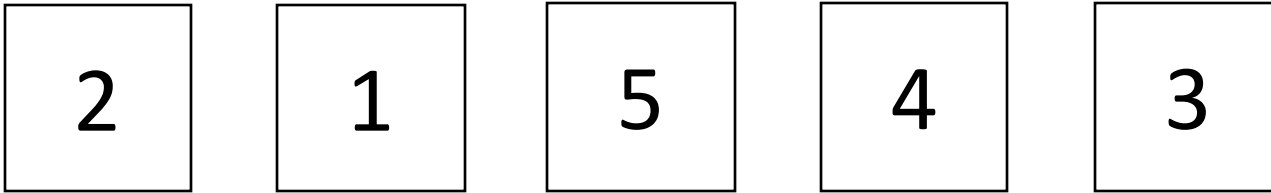


Solution 3 (reversing the array)

Inputs:

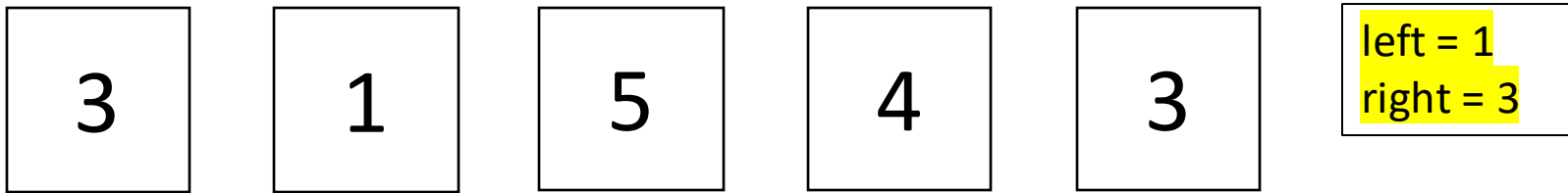
array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd reverse operation:



function call:

↓ reverse(array, 0, n-1) ↓ (n = 5)



↑
left

↑
right

left = 1
right = 3

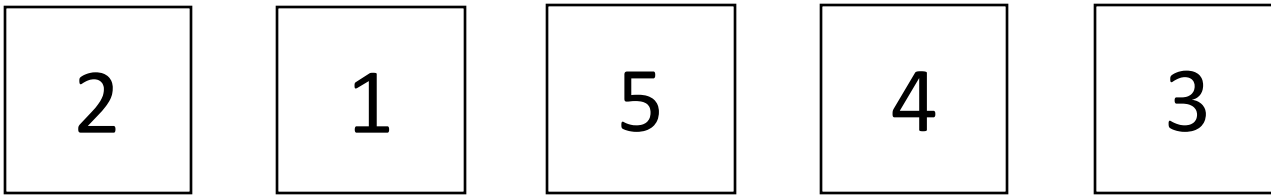
→ left <? right
True → continue

Solution 3 (reversing the array)

Inputs:

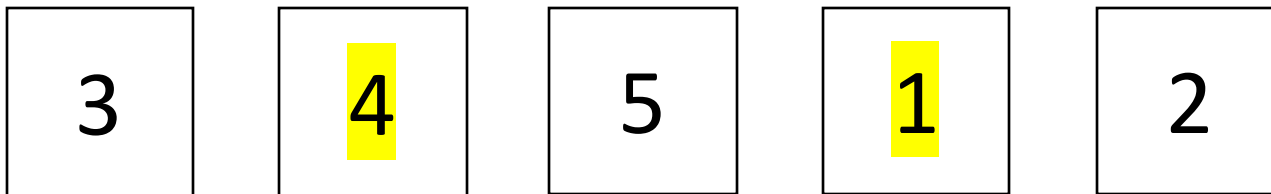
array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd reverse operation:



function call:

↓ reverse(array, 0, n-1) ↓ (n = 5)



↑
left

↑
right

left = 1
right = 3

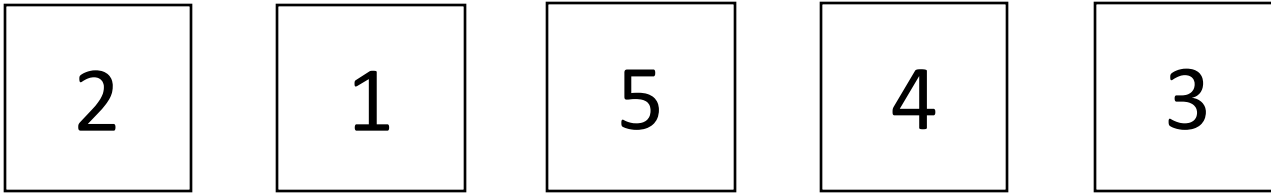
swap array[left] and array[right]

Solution 3 (reversing the array)

Inputs:

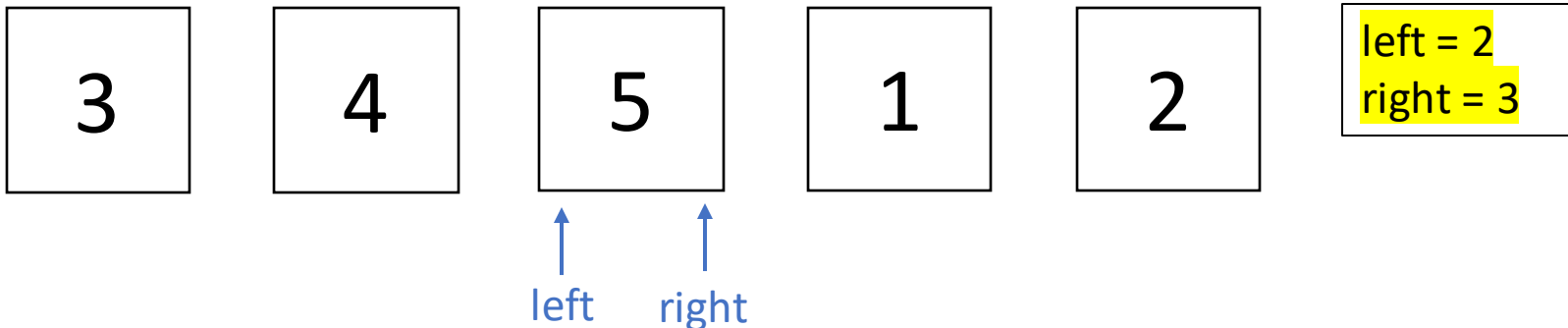
array = {1, 2, 3, 4, 5} and k = 2

Implementing 3rd reverse operation:



function call:

↓ reverse(array, 0, n-1) ↓ (n = 5)



→ left <? right
False → end of the while loop
return array

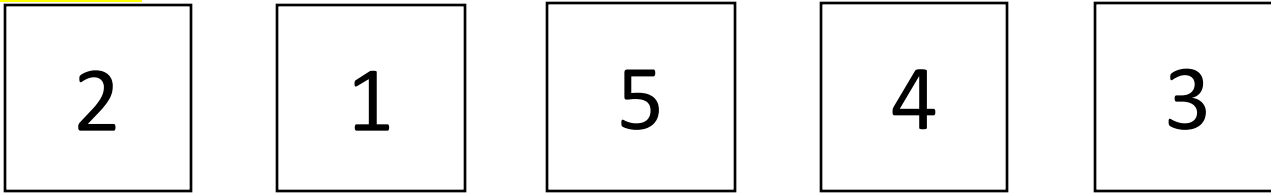
Solution 3 (reversing the array)

Inputs:

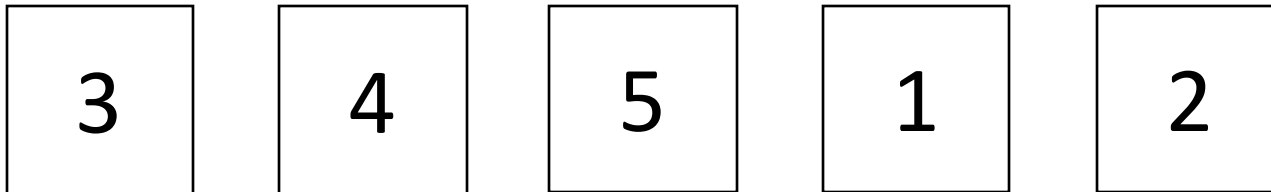
array = {1, 2, 3, 4, 5} and k = 2

Array is updated after the 3rd reverse operation:

Before:



After:

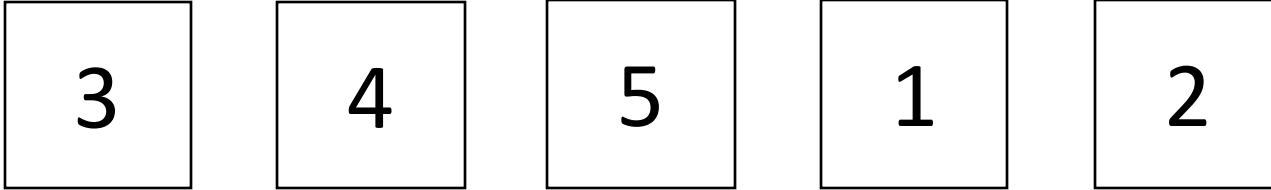


Solution 3 (reversing the array)

Inputs:

array = {1, 2, 3, 4, 5} and k = 2

After the 3rd reverse operation, the algorithm ends.



Solution 3 (reversing the array)

Time Complexity?

The algorithm implements an operation called **reverse** 3 times. This operation uses a while loop. As in Solution 2, this operation takes at most $O(n)$ time. And if we try to minimize one of these 3 operations, the other one will get closer to n . So, best case and worst case the time complexities are the same: $O(n)$.

Space Complexity?

The algorithm doesn't use an additional array as in Solution 2. It only uses a temporary variable for swapping operations and iterators. Which makes the space complexity constant.

Comparison of the solutions

Solution	Best Case Time Complexity	Worst Case Time Complexity	Space Complexity
Solution 1 (naïve)	$\Omega(1)$	$O(n^2)$	$O(1)$
Solution 2 (auxiliary array)	$\Omega(n)$	$O(n)$	$O(k)$
Solution 3 (reverse array)	$\Omega(n)$	$O(n)$	$O(1)$