

# Distance-based exponential probability models for constrained combinatorial optimization problems.

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# The Question

How do we approach problems with [constraints](#) by means of [EDAs](#)?  
(taking into account their [nature](#)...)

# Constrained Optimization Problems

## Definition

minimising  $f(\mathbf{x}), \quad \mathbf{x} = (x_1, \dots, x_n)$

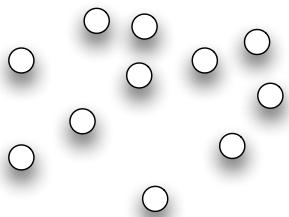
subject to,  $g_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, r$   
 $h_j(\mathbf{x}) = 0, \quad j = r + 1, \dots, m$

### Some examples

- Knapsack Problem
- Graph Colouring Problem
- Maximum Satisfiability Problem
- Capacitated Arc Routing Problem
- ...

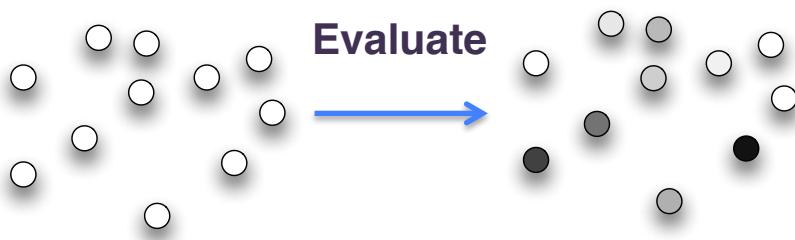
# Estimation of distribution algorithms

**Generate** a set  
of solutions



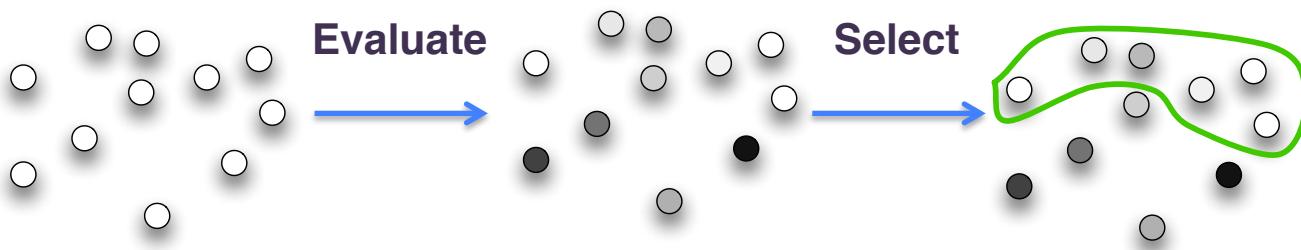
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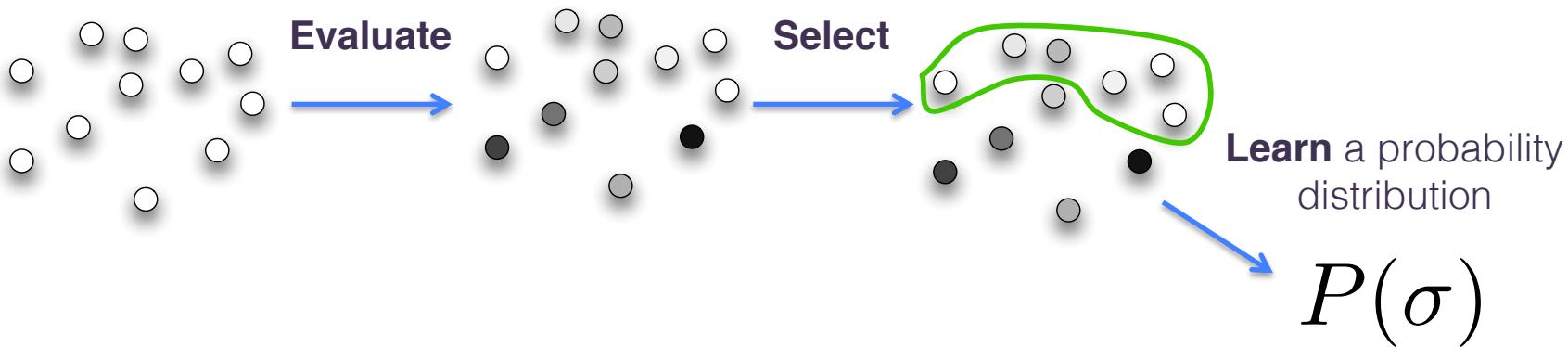
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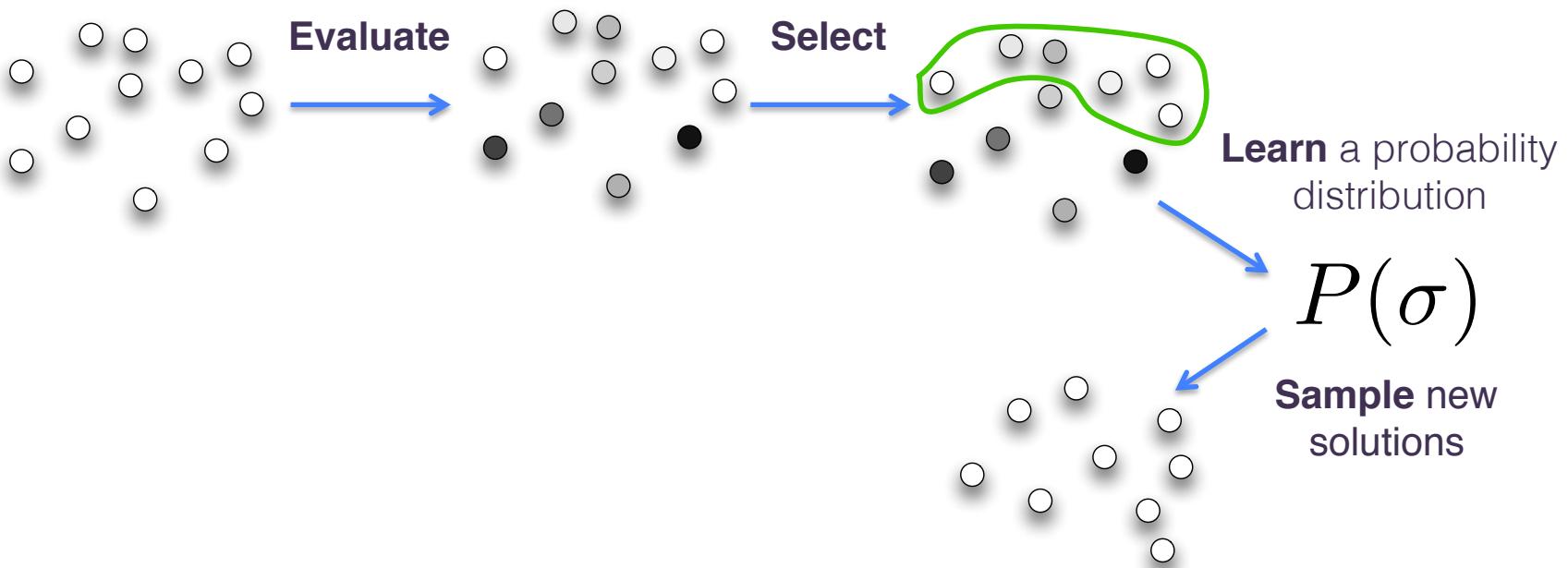
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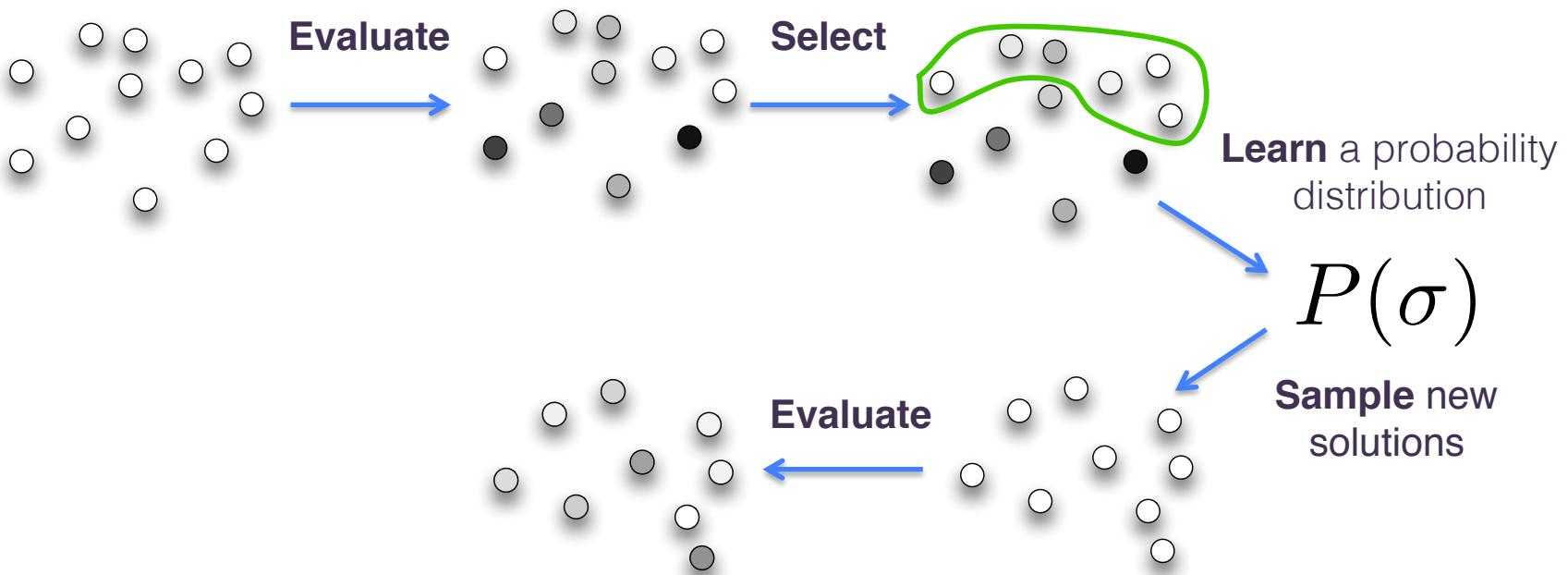
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**Generate** a set  
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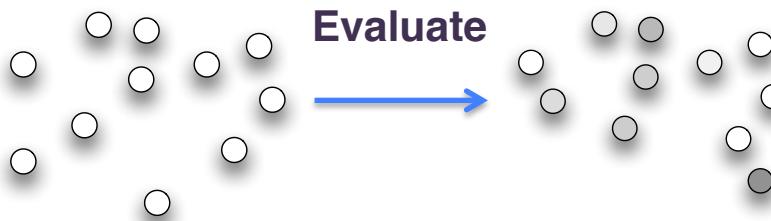
# Estimation of distribution algorithms

**Generate** a set  
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# Estimation of distribution algorithms

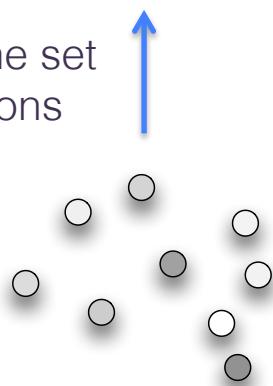
**Generate** a set of solutions



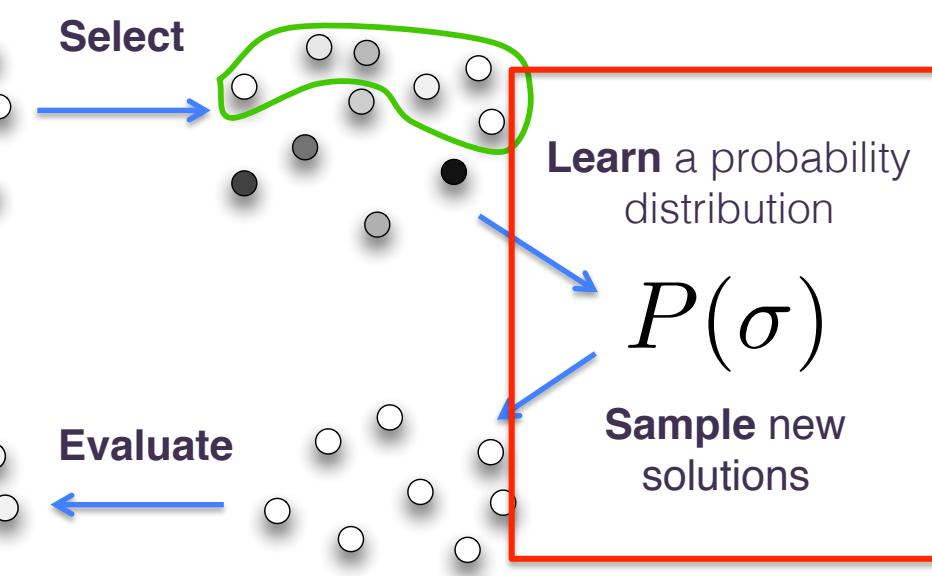
**Evaluate**

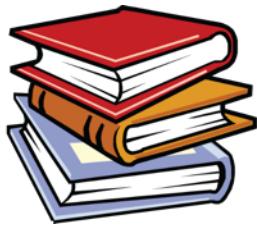
**Select**

**Update** the set of solutions



**Evaluate**





EDAs reported  
in the literature

## Combinatorial Problems

- UMDA [Mühlenbein, 1998]
- MIMIC [DeBonet, 1997]
- FDA [Mühlenbein, 1999]
- EBNA [Etxeberria, 1999]
- BOA [Pelikan, 2000]
- EHBSA [Tsutsui, 2003]
- NHBSA [Tsutsui, 2006]
- TREE [Pelikan, 2007]
- REDA [Romero, 2009]

$\Omega$

$S_n$

## Permutation Problems

- IDEA-ICE [Bosman, 2001]
- MEDA [Ceberio, 2011]
- PLEDA [Ceberio, 2013]
- GMEDA [Ceberio, 2014]
- RKEDA [Ayodele, 2016]

## Continuous Problems

- UMDA<sub>C</sub> [Larrañaga, 2000]
- MIMIC<sub>C</sub> [Larrañaga, 2000]
- EGNA [Larrañaga, 2000]
- EMNA [Larrañaga, 2001]
- IDEA [Bosman, 2000]

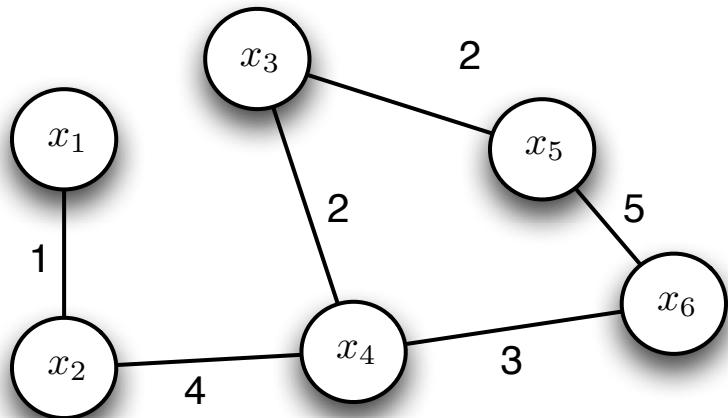
$\mathbb{R}^n$

## Constrained Problems

?

# Graph Partitioning Problem

## Definition



Find a  $k$ -partition of vertices  
minimising the weight of edges  
between sets: **the cut size**

We considered the balanced 2-partition GPP.

Solutions are codified as...

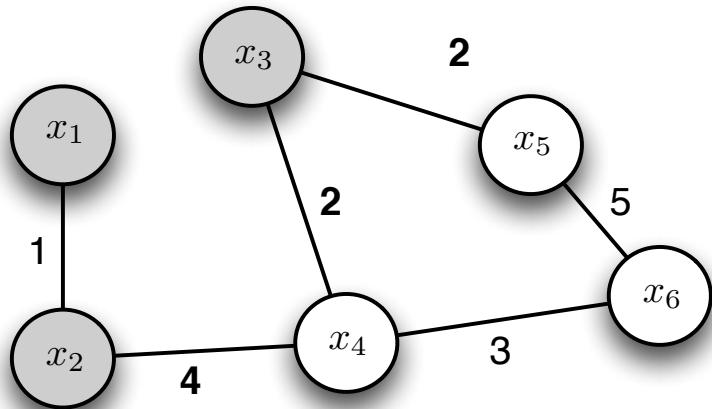
$$\mathbf{x} \in \{0, 1\}^n$$

$$f(\mathbf{x}) = \sum_{i=1}^n \sum_{j=1}^n x_i(1 - x_j)w_{ij}$$

**Objective Function**

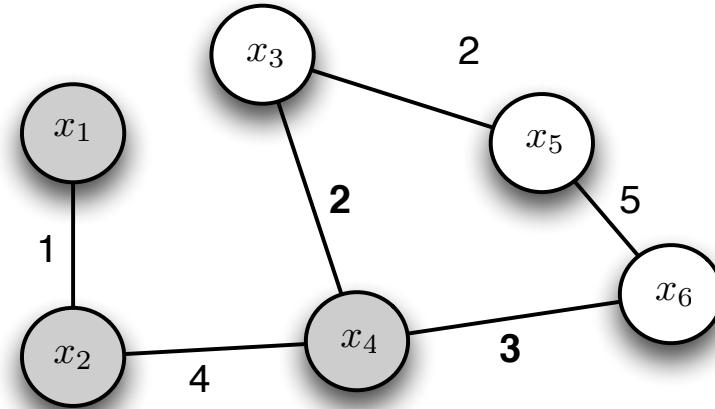
# Graph Partitioning Problem

## Example



$$\mathbf{x}^1 \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 \end{bmatrix} \quad f(\mathbf{x}^1) = 8$$

$x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_5 \quad x_6$



$$\mathbf{x}^2 \begin{bmatrix} 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix} \quad f(\mathbf{x}^2) = 5$$

$x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_5 \quad x_6$

The constraint: equal number of zeros as ones

$$\longrightarrow \sum_{i=1}^n x_i = n/2$$

# Graph Partitioning Problem

Why are they challenging?

The search space of solutions induced by the codification is...

000000	001000	010000	011000	100000	101000	110000	111000
000001	001001	010001	011001	100001	101001	110001	111001
000010	001010	010010	011010	100010	101010	110010	111010
000011	001011	010011	011011	100011	101011	110011	111011
000100	001100	010100	011100	100100	101100	110100	111100
000101	001101	010101	011101	100101	101101	110101	111101
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000100	001100	010100	011100	100100	101100	110100	111100
000101	001101	010101	011101	100101	101101	110101	111101
000110	001110	010110	011110	100110	101110	110110	111110
000111	001111	010111	011111	100111	101111	110111	111111

The majority of the solutions are not **feasible** !

# What happens if we run a UMDA?

**Univariate Marginal Distribution Algorithm (UMDA)**

$$P(x) = \prod_{i=1}^n P(x_i)$$



First order marginals  
No dependencies are considered

# What happens if we run a UMDA?

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$
0	0.6	0.5	0.25	0.66	0.9	0.5
1	0.4	0.5	0.75	0.33	0.1	0.5

↓

0 0 1 1 0 0

1 1 0 0 0 0

0 0 1 1 0 1

0 1 0 1 0 1

0 1 0 1 1 1

# What happens if we run a UMDA?

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$
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↓

0 0 1 1 0 0

1 1 0 0 0 0

0 0 1 1 0 1

0 1 0 1 0 1

0 1 0 1 1 1

Unfeasible solutions are generated...

# Different approaches

## Literature review

### 1. Repair solutions

- **Modify** solutions to hold the constraints

### 2. Penalty functions

- **Punish** solutions to be discarded when selection

### 3. Guarantee feasibility when sampling

- In EDAs, **adapt** the sampling to create feasible solutions

The role of the probability model is somehow denaturalized



# The Idea

How do we approach problems with **constraints** by means of **EDAs**?  
(taking into account their **nature**...)

Conduct the optimisation **entirely**  
**on the set of feasible solutions**...

**4. Use probability distributions  
defined on this set**

# Motivation

## Permutation-based Problems

Combinatorial Optimization Problems

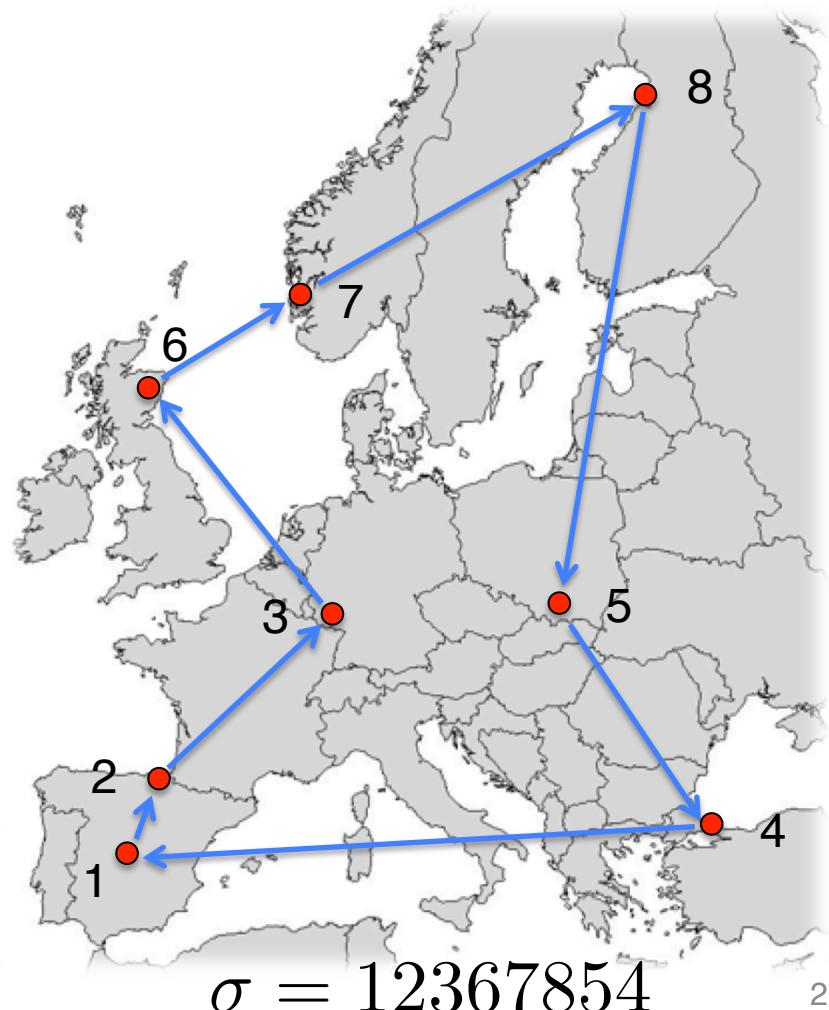
Whose solutions are represented as  
permutations

The search space consist of  $n!$  solutions

$$8! = 40320$$

$$20! = 2.43 \times 10^{18}$$

Travelling Salesman Problem (TSP)



# Motivation

## Permutation-based Problems

The space of permutations can be seen as a constrained space  
of the integers space

$n = 3$

111	211	311
112	212	312
113	213	313
121	221	321
122	222	322
123	223	323
131	231	331
132	232	332
133	233	333

# Motivation

## Permutation-based Problems

The space of permutations can be seen as a constrained space  
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# Motivation

## Probability Models on Rankings

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- M. A. Fligner and J. S. Verducci (1998), Multistage Ranking Models, *Journal of the American Statistical Association*, vol. 83, no. 403, pp. 892-901.
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- D. R. Luce (1959), Individual Choice Behaviour, *Wiley*.
- R. A. Bradley AND M. E. Terry (1952), Rank Analysis of Incomplete Block Designs: I. The Method of Paired Comparisons, *Biometrika*, vol. 39, no. 3, pp. 324-345.
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# Motivation

## Probability Models on Rankings

$$P(\sigma) = \frac{1}{\psi(\theta)} e^{-\theta D(\sigma, \sigma_0)}$$

Mallows

Distance-based

$$P(\sigma) = \frac{1}{\psi(\theta)} e^{-\sum_{j=1}^{n-1} \theta_j S_j(\sigma, \sigma_0)}$$

Generalized Mallows

$$P(\sigma) = \prod_{i=1}^{n-1} \frac{w_{\sigma(i)}}{\sum_{j=i}^n w_{\sigma(j)}}$$

Plackett-Luce

Order statistics

$$P(\sigma) = \prod_{i=1}^{n-1} \prod_{j=i+1}^n \frac{w_{\sigma(i)}}{w_{\sigma(i)} + w_{\sigma(j)}}$$

Bradley-Terry

# The Idea

Do probability models for constrained spaces exist?

No...

at least, we do not know them...



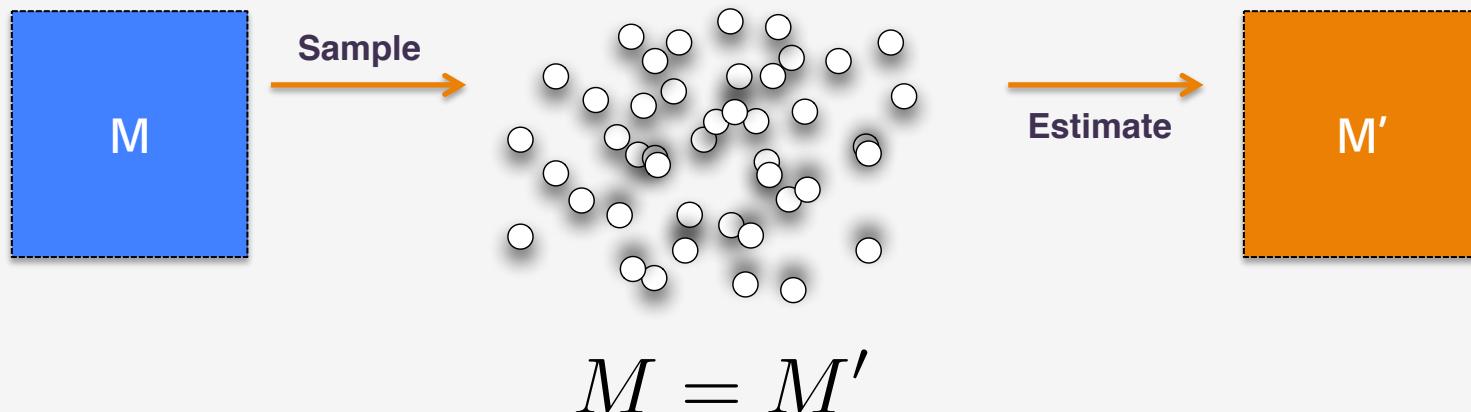
# Designing probability models

## Requirements

$$1 \quad \forall x \in \Omega, \quad 0 \leq P(x) \leq 1$$

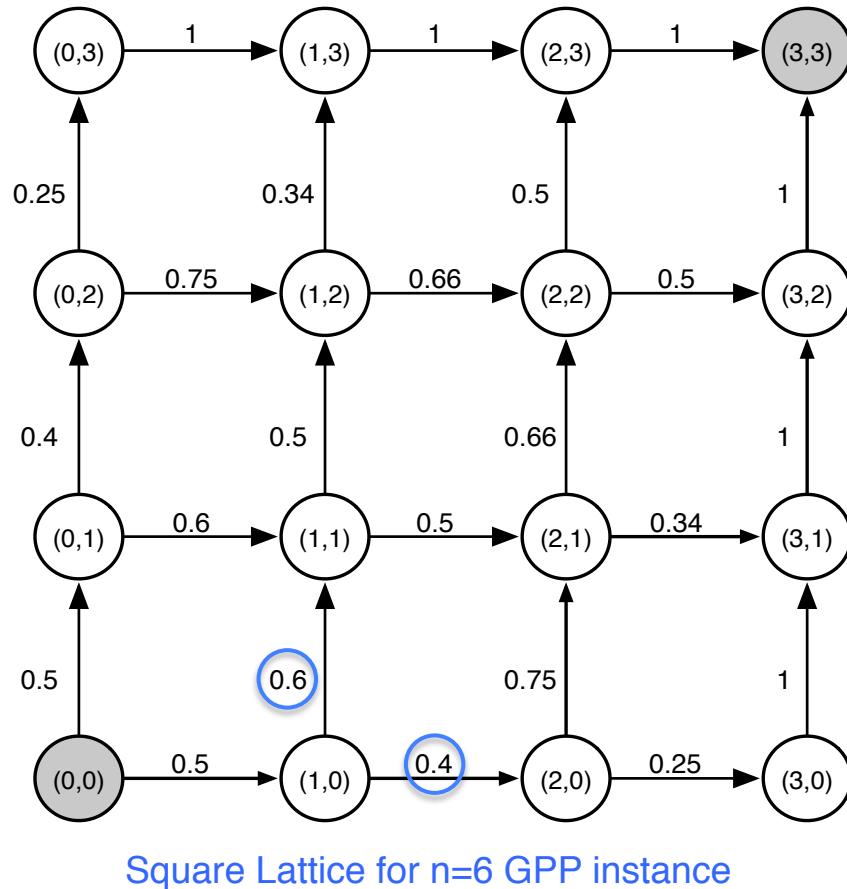
$$2 \quad \sum_{i=1}^{|\Omega|} P(x_i) = 1$$

3 Develop efficient **learning** and **sampling** methods



# Previously: A Square Lattice

## The Probability Model



$$\mathbf{x} = (1, 0, 0, 1, 0, 1)$$

$$\pi = ((0, 1), (1, 1), (2, 1), (2, 2), (3, 2), (3, 3))$$

Solutions are modelled as paths on a square lattice of  $(n/2+1)^2$  vertices

Vertices: the number of ones and zeros at that stage

Edges: the probability of moving from one vertex to another.

$$P(\pi) = \prod_{i=1}^n P_{(h^i, k^i), (h^{i+1}, k^i)}^{(h^{i+1} - h^i)} P_{(h^i, k^i), (h^i, k^{i+1})}^{(k^{i+1} - k^i)}$$

Horizontal step      Vertical step

# Previously: A Square Lattice

## The Probability Model



**Suffer from...**

**1**

Redundant model

**2**

The order?

**3**

Too many parameters

**4**

Difficult to analyze

Modelled as  
lattice of  
sites

Number of  
sites that

Probability of  
a vertex to

$k^{i+1} - k^i)$   
 $(h^i, k^i), (h^i, k^i + 1)$

Vertical step

# The Idea

## A different model

**Distance-based exponential probability models**

$$P(x) = \frac{e^{-\theta d(x, \bar{x})}}{\sum_{y \in \Omega} e^{-\theta d(y, \bar{x})}}$$

Diagram illustrating the components of the formula:

- Spread parameter** ( $\theta$ ): Indicated by a red arrow pointing to the term  $e^{-\theta d(x, \bar{x})}$ .
- Central solution** ( $\bar{x}$ ): Indicated by a red arrow pointing to the term  $d(x, \bar{x})$ .
- A distance-metric**: Indicated by a blue arrow pointing to the term  $d(x, \bar{x})$ .

# The Challenges

A distance-metric is needed

$$x, y \in \{0, 1\}^n$$

$$d(x, y) ?$$

000000	001000	010000	011000	100000	101000	110000	111000
000001	001001	010001	011001	100001	101001	110001	111001
000010	001010	010010	011010	100010	101010	110010	111010
000011	001011	010011	011011	100011	101011	110011	111011
000100	001100	010100	011100	100100	101100	110100	111100
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000110	001110	010110	011110	100110	101110	110110	111110
000111	001111	010111	011111	100111	101111	110111	111111

Hamming Distance

$$d_H(x, \bar{x}) = \sum_{i=1}^n \mathbf{1}_{[x(i) \neq \bar{x}(i)]}$$

$x$	1	1	0	0	1	0
-----	---	---	---	---	---	---

$$d_H(x, \bar{x}) = ?$$

$\bar{x}$	1	0	0	0	1	1
-----------	---	---	---	---	---	---

# The Challenges

A distance-metric is needed

$$x, y \in \{0, 1\}^n$$

$$d(x, y) ?$$

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000111	001111	010111	011111	100111	101111	110111	111111

Hamming Distance

$$d_H(x, \bar{x}) = \sum_{i=1}^n \mathbf{1}_{[x(i) \neq \bar{x}(i)]}$$

$x$	1	1	0	0	1	0
-----	---	---	---	---	---	---

$$d_H(x, \bar{x}) = 2$$

$\bar{x}$	1	0	0	0	1	1
-----------	---	---	---	---	---	---

# The Challenges

A distance-metric is needed

$$x, y \in \{0, 1\}^n$$

$$d(x, y) ?$$

000000	001000	010000	011000	100000	101000	110000	111000
000001	001001	010001	011001	100001	101001	110001	111001
000010	001010	010010	011010	100010	101010	110010	111010
000011	001011	010011	011011	100011	101011	110011	111011
000100	001100	010100	011100	100100	101100	110100	111100
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Hamming Distance

$$d_H(x, \bar{x}) = \sum_{i=1}^n \mathbf{1}_{[x(i) \neq \bar{x}(i)]}$$

$x$	1	1	0	0	1	0
-----	---	---	---	---	---	---

$$d_H(x, \bar{x}) = ?$$

$\bar{x}$	0	0	1	1	0	1
-----------	---	---	---	---	---	---

# The Challenges

A distance-metric is needed

$$x, y \in \{0, 1\}^n$$

$$d(x, y) ?$$

000000	001000	010000	011000	100000	101000	110000	111000
000001	001001	010001	011001	100001	101001	110001	111001
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Hamming Distance

$$d_H(x, \bar{x}) = \sum_{i=1}^n \mathbf{1}_{[x(i) \neq \bar{x}(i)]}$$

$x$	1	1	0	0	1	0
-----	---	---	---	---	---	---

$\bar{x}$	0	0	1	1	0	1
-----------	---	---	---	---	---	---

~~$d_H(x, \bar{x}) = 6$~~

$d(x, \bar{x}) = 0$

# The Challenges

A distance-metric is needed

$$x, y \in \{0, 1\}^n$$

$$d(x, y) ?$$

000000	001000	010000	011000	100000	101000	110000	111000
000001	001001	010001	011001	100001	101001	110001	111001
000010	001010	010010	011010	100010	101010	110010	111010
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000110	001110	010110	011110	100110	101110	110110	111110
000111	001111	010111	011111	100111	101111	110111	111111

New Distance (based on Hamming)

$$d(x, \bar{x}) = \min\{d_H(x, \bar{x}), d_H(\neg x, \bar{x})\}$$

Max distance  
 $\lfloor n/2 \rfloor$

$$\binom{n/2}{k/2}^2 \leftarrow$$

The number of  
solutions at distance  $k$

# The Challenges

What about the normalization constant?

## Distance-based exponential probability models

$$P(x) = \frac{e^{-\theta d(x, \bar{x})}}{\sum_{y \in \Omega} e^{-\theta d(y, \bar{x})}}$$

Diagram annotations:

- Spread parameter:  $\theta$  (circled in red)
- Central solution:  $\bar{x}$  (circled in blue)
- A distance-metric:  $d(x, \bar{x})$  (circled in red)
- Normalization function:  $\sum_{y \in \Omega} e^{-\theta d(y, \bar{x})}$  (boxed in blue)

The size of the search space. GPP:

$$K = \binom{n}{n/2}$$

$$\sum_{l=0}^{K/2} \binom{n/2}{l}^2 e^{-\theta 2l}$$

# The Challenges

How to estimate MLE parameters? In two steps...

$$L(\theta, \bar{x} | \mathbf{x}) = \prod_{i=1}^N \frac{e^{-\theta d(x_i, \bar{x})}}{\psi(\theta)} \quad \leftarrow \text{Find the parameters that maximize it}$$

$$\log L(\theta, \bar{x} | \mathbf{x}) = -\theta N \bar{d} - N \log \left[ \sum_{l=0}^{K/2} \binom{n/2}{l}^2 e^{-\theta 2l} \right]$$

1

Calculate  $\bar{x}$   
as average  
solution

1	1	0	0	1	0
0	1	0	1	1	0
1	0	0	0	1	1
1	1	0	0	1	0

# The Challenges

How to estimate MLE parameters? In two steps...

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1

Calculate  $\bar{x}$   
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1	1	0	0	1	0
0	1	0	1	1	0
1	0	0	0	1	1
1	1	0	0	1	0

0. 75	0. 75	0	0. 25	1	0. 25
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# The Challenges

How to estimate MLE parameters? In two steps...

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1

Calculate  $\bar{x}$  as average solution

1	1	0	0	1	0
0	1	0	1	1	0
1	0	0	0	1	1
1	1	0	0	1	0

1	1	0	0	1	0
---	---	---	---	---	---

2

Approximate  $\theta$  numerically

$$\sum_{l=0}^{K/2} \binom{n/2}{l}^2 (2l - \bar{d}) e^{-\theta 2l} = 0$$

**Newton-Raphson...**

# The Challenges

How to sample new solutions? Also in two steps...

- 1 A distance  $k$  at which generate a solution is sampled from

$$P(k) = \sum_{x|d(x, \bar{x})=k} P(x) = \binom{n/2}{k/2}^2 \frac{e^{-\theta k}}{\psi(\theta)}$$

- 2 Generate a solution at distance  $k$

$\bar{x}$	1	1	0	0	1	0
-----------	---	---	---	---	---	---

$$k = 4$$

$x$	1	1	0	0	1	0
-----	---	---	---	---	---	---

# The Challenges

How to sample new solutions? Also in two steps...

- 1 A distance  $k$  at which generate a solution is sampled from

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$\bar{x}$	1	1	0	0	1	0
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# The Challenges

How to sample new solutions? Also in two steps...

- 1 A distance  $k$  at which generate a solution is sampled from

$$P(k) = \sum_{x|d(x, \bar{x})=k} P(x) = \binom{n/2}{k/2}^2 \frac{e^{-\theta k}}{\psi(\theta)}$$

- 2 Generate a solution at distance  $k$

$\bar{x}$	1	1	0	0	1	0
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$$k = 4$$

$x$	1	0	1	0	0	1
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# Experimental Study

## Experimental Setting

UMDA \*, TREE\*<sup>1</sup> ,  
Lattice, Exp

(\*) adapted sampling

Algorithms

Pop-size:  $10n$

Sel-size:  $5n$

Off-size:  $10n$

Max evals.:  $100n^2$

10 repetitions

Parameters

22 Instances  
(Johnson et al.)

G-type and  
U-type

$n=124, 250, 500,$   
 $1000$

Benchmarks

<sup>1</sup>M. Pelikan, S. Tsutsui, and R. Kalapala, *Dependency Trees, Permutations and Quadratic Assignment Problem*, Medal Report No. 2007003 Tech. Rep. 2007.

# Experimental Study

## Results - Performance

Instance	Best Fitness	ARPD			Tree
		Exp	Lattice	UMDA	
G124.02	13	0,32	0,32	0,61	0,19
G124.16	449	0,04	0,02	0,05	0,01
G250.01	31	0,40	0,33	0,49	0,20
G250.02	118	0,10	0,07	0,14	0,06
G250.04	360	0,04	0,04	0,10	0,03
G250.08	830	0,05	0,01	0,05	0,01
G500.005	61	0,38	0,30	0,40	0,08
G500.01	234	0,14	0,09	0,21	0,07
G500.02	642	0,06	0,03	0,11	0,03
G500.04	1754	0,05	0,02	0,06	0,02
G1000.0025	131	1,16	2,96	3,20	0,74
G1000.005	496	0,52	1,22	1,28	0,88
G1000.01	1420	0,28	0,56	0,66	0,62
G1000.02	3450	0,18	0,35	0,40	0,39
U500.05	23	1,34	1,17	1,89	0,57
U500.10	61	1,50	1,05	1,12	0,57
U500.20	185	0,89	0,56	0,87	0,44
U500.40	412	0,81	0,41	0,38	0,28
U1000.05	77	4,82	1,62	12,83	2,39
U1000.10	170	5,09	1,67	11,67	3,73
U1000.20	352	5,23	1,67	10,58	4,94
U1000.40	862	4,09	1,53	3,24	2,29

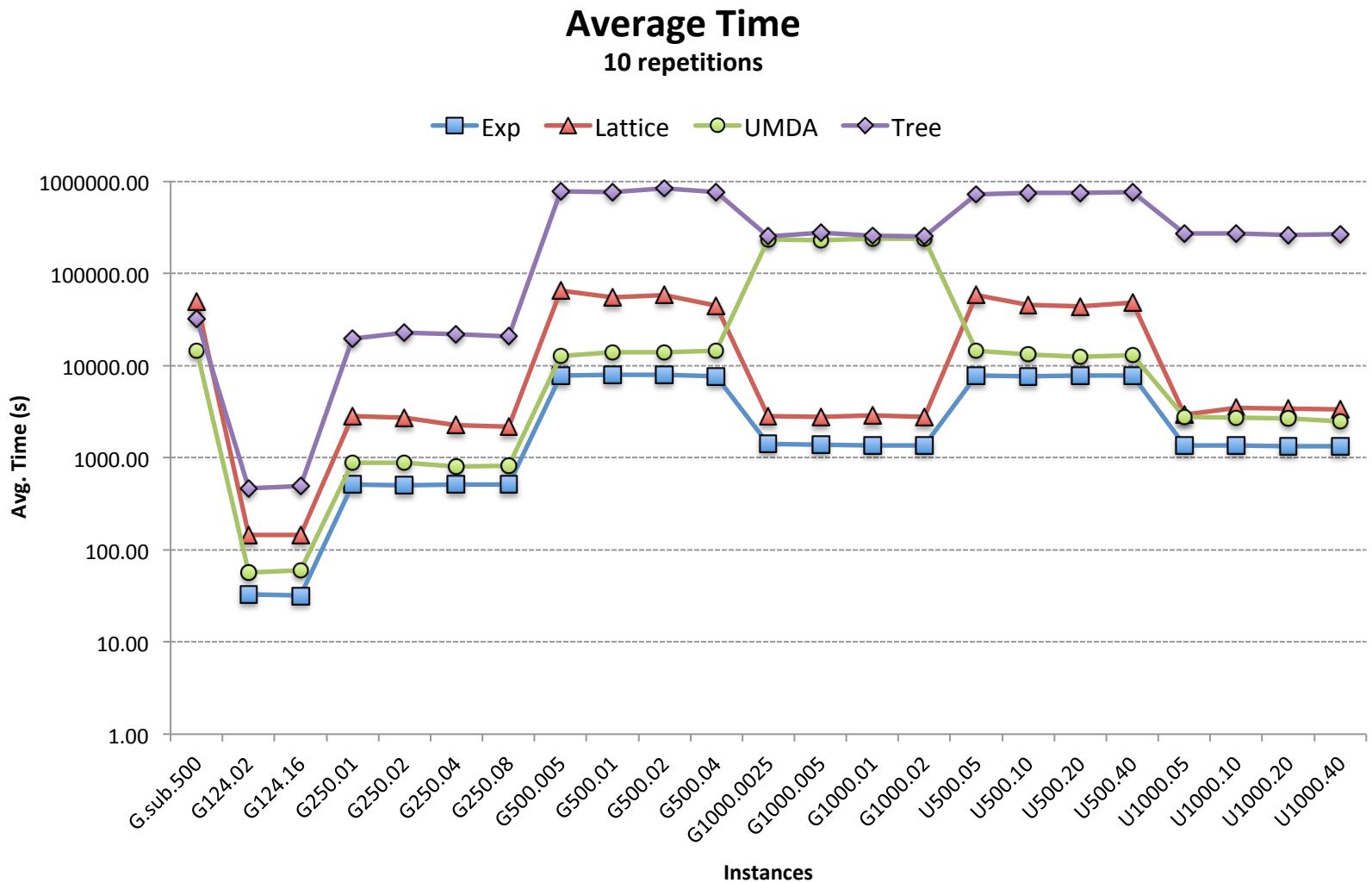
# Experimental Study

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# Experimental Study

## Results - Performance



# Conclusions

The experiments support the validity of our research line:  
Designing probability models exclusively on the set of feasible solutions

**However, many issues still pending...**

# Future Work

Many aspects to be faced!

Consider other  
methods to estimate  
 $\bar{x}$

Analyze regularity of  
the solutions regarding  
the distance

Propose other  
distances

**The probability model**

Larger benchmarks  
Understand the  
dynamics of the EDA  
for different problem  
sizes

**Experimentation**

New probability  
models

**EDA**

# Distance-based exponential probability models for constrained combinatorial optimization problems.

Thank you for your attention!!!



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