GADTs Meet Their Match:

Pattern-Matching Warnings That Account for GADTs, Guards, and Laziness

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Pattern Syntax $K \in \mathsf{Con}$ $x, y, a, b \in Var$ $\tau, \sigma \in \mathsf{Type}$ $e \in Expr$ $:= x : \tau$ $K \overline{a} \overline{\gamma} \overline{e : \tau}$ TyCt ::= $\tau_1 \sim \tau_2 \mid ...$ $\gamma \in$:= let $x : \tau = e$ Grd $a \in$ $K \overline{a} \overline{y} \overline{y : \tau} \leftarrow x$ **Oracle Syntax** $:= \emptyset \mid \Gamma, x : \tau \mid \Gamma, a$ Context $:= \times | \checkmark | \Delta, \delta | \Delta_1 \vee \Delta_2$ Delta $:= \times | \checkmark | \Delta, \delta | \Delta_1 \vee \Delta_2$ Delta $:= \gamma \mid K \ \overline{x : \tau} \leftarrow y \mid x \not\approx K \mid x \approx \bot \mid x \not\approx \bot \mid x \approx e$ Constraints

TODO LIST

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GADTs Meet Their Match: 1:3

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Clause tree
                                                                                                   \mathcal{T}[r] ::= Rhs
                                                                                                                                       Many \overline{r}
                                                                                        t_G \in \mathsf{Gdt} ::= \mathcal{T}[t_G]
                                                                                                                         | Guard q t_G
                                                                                        t_A \in \mathsf{Ant} ::= \mathcal{T}[t_A]
                                                                                                                         \mid MayDiverge t_A
                                                                                                                                       Inaccessible t_A
                                                                                                       Checking guard trees
                                                                                                                   \mathcal{U}(\Delta, Gdt) = \Delta
\mathcal{U}(\Delta, \mathsf{Rhs})
\mathcal{U}(\Delta, \mathsf{Many}\ \bar{t})
                                                                                                    = \mathcal{U}(...\mathcal{U}(\mathcal{U}(\Delta, t_1), t_2)..., t_n)
\mathcal{U}(\Delta,\operatorname{Guard}(!x)\ t) = \mathcal{U}(\Delta \oplus (x \not\approx \bot),t)

\mathcal{U}(\Delta,\operatorname{Guard}(\operatorname{let} x = e)\ t) = \mathcal{U}(\Delta \oplus (x \approx e),t)
\mathcal{U}(\Delta,\mathsf{Guard}\ (K\ \overline{a}\ \overline{\gamma}\ \overline{y:\tau}\leftarrow x)\ t) = (\Delta\oplus(x\not\approx K)\oplus(x\not\approx\bot))\cup\mathcal{U}(\Delta\oplus(K\ \overline{y:\tau}\leftarrow x)\oplus\overline{\gamma},gs)
                                                                                                              \mathcal{A}(\Delta, Gdt) = Ant
                                                                                                   = \begin{cases} \text{Inaccessible Rhs}, & \mathcal{V}(\Gamma, \Delta) = \emptyset \\ \text{Rhs}, & \text{otherwise} \end{cases}
= \text{Many} \left(\mathcal{A}(\Delta_0', t_1), \dots, \mathcal{A}(\Delta_{n-1}', t_n)\right) \text{ where } \begin{cases} \Delta_0' &= \Delta \\ \Delta_{n+1}' &= \mathcal{U}(\Delta_n', t_{n+1}) \end{cases}
\mathcal{A}(\Delta, Rhs)
\mathcal{A}(\Delta, \mathsf{Many}\ \overline{t})
 \begin{array}{lll} \mathcal{A}(\Delta,\operatorname{Guard}\ (!x)\ t) & = & \mathcal{D}(\Delta,\mathcal{A}(\Delta\oplus(x\not\approx\bot),t)) \\ \mathcal{A}(\Delta,\operatorname{Guard}\ (\text{let}\ x=e)\ t) & = & \mathcal{A}(\Delta\oplus(x\approx e),t) \end{array} 
\mathcal{A}(\Delta,\operatorname{Guard}(K\ \overline{a}\ \overline{\gamma}\ \overline{y:\tau}\leftarrow x)\ t) \quad = \quad \mathcal{D}(\Delta,\mathcal{A}(\Delta\oplus(K\ \overline{y:\tau}\leftarrow x)\oplus\overline{\gamma},t))
                                                        \mathcal{D}(\Delta,\mathsf{Ant}) = \mathsf{Ant} \mathcal{D}(\Delta,t) = \begin{cases} t, & \mathcal{V}(\Gamma,\Delta\oplus(x\approx\bot)) = \varnothing \\ \mathsf{MayDiverge}\ t & \mathsf{otherwise} \end{cases}
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Test if Oracle state Delta is unsatisfiable

$$\frac{\not\vdash_{SAT} \Gamma \vdash \Delta}{\not\vdash_{SAT} \Gamma \vdash fvs\Gamma \triangleright \Delta}$$

$$\not\vdash_{SAT} \Gamma \vdash \Delta$$

Test a list of SAT roots for inhabitants

$$\begin{array}{c|c}
 & \swarrow_{\text{SAT}} \Gamma \vdash \overline{x} \triangleright \Delta \\
 & \swarrow_{\text{SAT}} \Gamma \vdash x_i \triangleright \Delta \\
 & \swarrow_{\text{SAT}} \Gamma \vdash \overline{x} \triangleright \Delta
\end{array}$$

Test a single SAT root for inhabitants

Add a single equality to Δ

$$\nvdash_{\mathsf{SAT}} \Gamma \vdash \oplus \Delta \delta$$

Term stuff: Bottom, negative info, positive info + generativity, positive info + univalence

$$x \not\approx sth \in \Delta \qquad x \approx K \overline{y} \in \Delta$$

$$\nearrow_{SAT} \Gamma \vdash \oplus \Delta x \approx \bot \qquad \nearrow_{SAT} \Gamma \vdash \oplus \Delta x \approx \bot$$

$$x \not\approx K \in \Delta \qquad x \approx K_i \overline{y} \in \Delta \quad i \neq j \quad K_i \text{ and } K_j \text{ generative}$$

$$\nearrow_{SAT} \Gamma \vdash \oplus \Delta x \approx K \overline{y} \qquad \nearrow_{SAT} \Gamma \vdash \oplus \Delta x \approx K_j \overline{z}$$

$$x \approx K \overline{\tau} \overline{y} \in \Delta \qquad \nearrow_{SAT} \Gamma \vdash \oplus \Delta \tau_i \sim \sigma_i \qquad x \approx K \overline{\tau} \overline{y} \in \Delta \qquad \nearrow_{SAT} \Gamma \vdash \oplus \Delta y_i \approx z_i$$

$$\nearrow_{SAT} \Gamma \vdash \oplus \Delta x \approx K \overline{\sigma} \overline{z} \qquad \nearrow_{SAT} \Gamma \vdash \oplus \Delta x \approx K \overline{\sigma} \overline{z}$$

Type stuff: Hand over to unspecified type oracle

 τ_1 and τ_2 incompatible to Givens in Δ according to type oracle

$$\not\vdash_{SAT} \Gamma \vdash \oplus \Delta \tau_1 \sim \tau_2$$

Mixed: Instantiate K and see if that leads to a contradiction TODO: Proper instantiation

$$\frac{\cancel{\nvdash}_{SAT} \ \Gamma \vdash y \triangleright \Delta \cup y \not\approx \bot}{\cancel{\nvdash}_{SAT} \ \Gamma \vdash \oplus \Delta x \approx K \ \overline{y}}$$

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