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	Algorithms - HW4	
Ans 1-		
71110	The state of the s	1
a	x*	
- u/	$S \longrightarrow \epsilon \mid \times S$	
	0110 5 301110 5 01110 5 0110 5 010 6 80 62	
1	_+	
6>	$S \rightarrow z \mid zS$	- FINA
	9 020 b56 020 6-2	· Ka
c	1-albertandelle Adizadisendelle ex	
-/-	s → a b	
d	01 /9 X2X <-2	1.4
~/	$s \rightarrow oA$	
isdic X	MA -> I as & XXXXX & & XXXXX & & XXX & & XXX & & XXX & & & XXX & & & XXX & & & & XXX &	
e>	(01)*	JASA
	S -> OA E	
	$A \rightarrow 1S$	-
		1
f>	(01) 1 <tent>> < +201 > = : < +201 > ></tent>	Fre K
J,	S-OA E 1 STANDED : STANDE	
	A -> 1B KAMMARAN KAMMARAN = :: KHINS	4
	$B \rightarrow 0A \mid E$	
	Kelsegant > :: = " " < Apath > "	
Ansz-	Chory of Springly of First Property	
(a)	S-EXS	
	XX = 2XX = 2X = 2	
	The state of the s	
6>	S-aBE	
	CB -> bs	
	$S \Rightarrow aB \Rightarrow abS \Rightarrow abaB \Rightarrow ababS \Rightarrow abab$	

Kreator

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S-> OB B -> 1B OC $C \rightarrow \epsilon$ S → 0B ⇒ 01B ⇒ 011B ⇒ 01110C ⇒ 01110 Ans 3-S -> asa bsb csc f S ⇒ aSa ⇒ ab Sba ⇒ abb Sbba ⇒ abbcSc bba ⇒ abbcb Sb cbba => abb c bb c bba S-XSX E S⇒XSX ⇒ aSX ⇒ axSXX → abSXX ⇒ abXX ⇒ abcX ⇒ abca < hostport> := < host > < exthost> Ans 4-<exthost > :: = ": " < port> | " " <port > :: = <digits> <digits > : = <digit > | <digit > <digit > < digit > :: = "0" "1" "2" "3" "4" "5" "6" "7" "8" "9"

a) < syntax > is the start rule for porsing BNF. b) Non-terminal must begin with an angle bracket (<) small brockets

First letter of the non-terminal must begin with capital or snat

c) After the first character, a scale name can have letters

letters, digits or "-"

d) Rule name can be in the control of the control of the control of the capital or snate letters. d) Rule name can have infinitely many characters. They should have a minimum of one character. e) hiterals are a set of characters (letters, symbols and digita).

These can be written inside single quotation 'OR can
be written inside " " quotation marks. g> "\"" OR """" a) In parsing expression granamers, choice operator is ordered wrlike BNF where any of the choices from the choice operator can be chosen). If the first alternative succeeds, the second alternative is ignored. Hence in this case it is always the first "if" then "else" # expression that will be matched. If this match fails, only then will the second expression be matched. Hence, with PEG'S, we do a not care about have a froblem if alternatives have the same frefix.

S = ("if" C "then" S "else" S) / ("if" C "then" S)

S => "if" (test 1) "then" S "else" S

=> "if" (test 1) "then" "if" C "then" S "else" S

=> "if" (test 1) "then" "if" (test 2) "then" S "else" S

=> "if" (test 1) "then" "if" (test 2) "then" S "else" S

=> "if" (test 1) "then" "if" (test 2) "then" frint (3) "else" frint (6)

=> "if" (test 1) "then" "if" (test 2) "then" frint (3) "else" frint (6)

() The start symbol 'S' in the first line has the else clause When the derivation starts, the first step involves:

S => "if" C "then" S "else" S

All citations:

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