

CS 300 Project One: Pseudocode & Runtime Analysis

This document contains finalized pseudocode for three data structures—Vector, Hash Table, and Binary Search Tree (BST)—to power ABCU Advising’s course tool. It also includes a Big-O runtime and memory analysis and a data-structure recommendation. The input file uses comma-separated values: courseNumber, name, prerequisite1..N.

Course Object (common across all designs):

STRUCT Course:

```
courseNumber : STRING  
  
name      : STRING  
  
prerequisites: LIST<STRING>
```

END STRUCT

File Open, Read, Parse (common logic used by each structure):

FUNCTION LoadFromFile(path: STRING) -> LIST<Course>:

```
courses = EMPTY LIST<Course>  
  
file = OPEN(path, "r")  
  
FOR EACH line IN file:  
  
    IF line IS EMPTY: CONTINUE  
  
    fields = SPLIT(line, ",")  
  
    IF LENGTH(fields) < 2:  
  
        PRINT "Format error: missing course number or name"; CONTINUE  
  
    course = NEW Course  
  
    course.courseNumber = TRIM(fields[0])
```

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    course.name      = TRIM(fields[1])

    course.prerequisites = EMPTY LIST

    FOR i FROM 2 TO LENGTH(fields)-1:

        prereq = TRIM(fields[i])

        IF prereq != "": APPEND(course.prerequisites, prereq)

    APPEND(courses, course)

CLOSE(file)

RETURN courses

END FUNCTION

Vector Design:

GLOBAL vecCourses : VECTOR<Course>

PROCEDURE BuildVector(path):

    vecCourses = LoadFromFile(path)

FUNCTION FindCourseVector(courseNum: STRING) -> Course|NULL:

    FOR EACH c IN vecCourses:

        IF c.courseNumber == courseNum: RETURN c

    RETURN NULL

PROCEDURE PrintCourseInfoVector(courseNum: STRING):

    c = FindCourseVector(courseNum)

    IF c == NULL: PRINT "Course not found"; RETURN

```

```

PRINT c.courseNumber + ": " + c.name

IF LENGTH(c.prerequisites) == 0:

    PRINT "Prerequisites: None"

ELSE:

    PRINT "Prerequisites: " + JOIN(c.prerequisites, ", ")

```

PROCEDURE PrintSortedListVector():

```

temp = COPY(vecCourses)

SORT temp BY courseNumber ASC // O(n log n)

FOR EACH c IN temp: PRINT c.courseNumber + ", " + c.name

```

Hash Table Design (keyed by courseNumber):

```

GLOBAL htCourses : HASH_TABLE<STRING, Course>

```

PROCEDURE BuildHash(path):

```

htCourses = NEW HASH_TABLE

list = LoadFromFile(path)

FOR EACH c IN list:

    INSERT htCourses[c.courseNumber] = c

```

FUNCTION FindCourseHash(courseNum: STRING) -> Course|NULL:

```

IF EXISTS htCourses[courseNum]: RETURN htCourses[courseNum]

RETURN NULL

```

PROCEDURE PrintCourseInfoHash(courseNum: STRING):

 c = FindCourseHash(courseNum)

 IF c == NULL: PRINT "Course not found"; RETURN

 PRINT c.courseNumber + ": " + c.name

 IF LENGTH(c.prerequisites) == 0:

 PRINT "Prerequisites: None"

 ELSE:

 PRINT "Prerequisites: " + JOIN(c.prerequisites, ", ")

PROCEDURE PrintSortedListHash():

 keys = KEYS(htCourses)

 SORT keys ASC // O(n log n)

 FOR EACH k IN keys:

 c = htCourses[k]

 PRINT c.courseNumber + ", " + c.name

Binary Search Tree Design (ordered by courseNumber):

STRUCT Node:

 key : STRING // courseNumber

 value : Course

 left : Node

 right : Node

END STRUCT

GLOBAL root : Node = NULL

FUNCTION InsertBST(root, key, value) -> Node:

IF root == NULL:

node = NEW Node; node.key = key; node.value = value; RETURN node

IF key < root.key: root.left = InsertBST(root.left, key, value)

ELSE IF key > root.key: root.right = InsertBST(root.right, key, value)

ELSE: root.value = value

RETURN root

PROCEDURE BuildBST(path):

root = NULL

list = LoadFromFile(path)

FOR EACH c IN list: root = InsertBST(root, c.courseNumber, c)

FUNCTION FindBST(root, key) -> Course|NULL:

WHILE root != NULL:

IF key == root.key: RETURN root.value

IF key < root.key: root = root.left

ELSE: root = root.right

RETURN NULL

PROCEDURE PrintCourseInfoBST(courseNum: STRING):

```

c = FindBST(root, courseNum)

IF c == NULL: PRINT "Course not found"; RETURN

PRINT c.courseNumber + ": " + c.name

IF LENGTH(c.prerequisites) == 0:

    PRINT "Prerequisites: None"

ELSE:

    PRINT "Prerequisites: " + JOIN(c.prerequisites, ", ")

```

PROCEDURE InOrder(node):

```

IF node == NULL: RETURN

InOrder(node.left)

PRINT node.value.courseNumber + ", " + node.value.name

InOrder(node.right)

```

PROCEDURE PrintSortedListBST():

```

InOrder(root)      // O(n) once built

```

Menu (shared UX over any backing store):

PROCEDURE Main():

```

dsType = PROMPT("Choose data structure: 1=Vector, 2=Hash, 3=BST")

path = PROMPT("Enter input file path")

IF dsType == 1: BuildVector(path)

ELSE IF dsType == 2: BuildHash(path)

ELSE: BuildBST(path)

```

REPEAT:

PRINT "1) Load file 2) Print course list 3) Print course info 9) Exit"

option = PROMPT("Select:")

IF option == 1:

IF dsType == 1: BuildVector(path)

ELSE IF dsType == 2: BuildHash(path)

ELSE: BuildBST(path)

PRINT "Data loaded."

ELSE IF option == 2:

IF dsType == 1: PrintSortedListVector()

ELSE IF dsType == 2: PrintSortedListHash()

ELSE: PrintSortedListBST()

ELSE IF option == 3:

key = PROMPT("Enter course number (e.g., CSCI200):")

IF dsType == 1: PrintCourseInfoVector(key)

ELSE IF dsType == 2: PrintCourseInfoHash(key)

ELSE: PrintCourseInfoBST(key)

ELSE IF option == 9:

BREAK

ELSE:

PRINT "Invalid option."

UNTIL FALSE

END PROCEDURE

Runtime & Memory Analysis (worst-case, n courses, p avg prereqs):

Load & Parse (common): $O(n \cdot p)$ to split lines and collect prerequisites.

Vector:

- Build: $O(n)$ (copy from parsed list). Memory: $O(n)$.
- Find course: $O(n)$ linear scan.
- Print sorted list: $O(n \log n)$ (sort by courseNumber) each time.
- Strengths: simple, cache-friendly; easy to implement.
- Weaknesses: linear search; must re-sort or maintain order for option 2.

Hash Table:

- Build: $O(n)$ average inserts. Memory: $O(n)$ + overhead for buckets.
- Find course: $O(1)$ average, $O(n)$ worst under heavy collisions.
- Print sorted list: extract keys then $O(n \log n)$ sort each time.
- Strengths: fastest lookups for option 3.
- Weaknesses: requires separate sort for option 2; tuning load factor; non-deterministic order.

Binary Search Tree (unbalanced vs. balanced):

- Build: $O(n \log n)$ average inserts; $O(n^2)$ worst if unbalanced (sorted input).
- Find course: $O(\log n)$ average; $O(n)$ worst if unbalanced.
- Print sorted list: $O(n)$ via in-order traversal (no extra sort).
- Memory: $O(n)$ nodes + pointers.

- Strengths: naturally ordered; very fast repeated sorted listings.
- Weaknesses: worst-case degradation unless self-balancing (AVL/Red-Black).

Recommendation:

Use a self-balancing BST (e.g., AVL or Red-Black). Rationale: advisors routinely need an alphanumerically sorted list (option 2), which is $O(n)$ to output from a BST after $O(\log n)$ inserts/lookups. While a hash table gives $O(1)$ average lookups, it still needs an $O(n \log n)$ sort each time the full list is printed. If sorted listing is infrequent and random single-course queries dominate, a hash table is a strong alternative.