COMP 352
Data Structures and Algorithms

BUCKET & RADIX SORT

Chapter 12

Bucket-Sort

- Let be S be a sequence of n (key, element) items with keys in the range $[0, N-1]$

- Bucket-sort uses the keys as indices into an auxiliary array B of sequences (buckets)
  - Phase 1: Empty sequence S by moving each entry (k, o) into its bucket $B[k]$
  - Phase 2: For $i = 0, \ldots, N-1$, move the entries of bucket $B[i]$ to the end of sequence S
- **Animation:** *https://www.cs.usfca.edu/~galles/visualization/BucketSort.html*

Bucket-Sort

Algorithm bucketSort(S):
Input: Sequence S of entries with integer keys in the range [0, N – 1]
Output: Sequence S sorted in non-decreasing order of the keys let B be an array of N sequences, each of which is initially empty

```
for each entry e in S do  //phase 1
   k = the key of e
   remove e from S
   insert e at the end of bucket B[k]
for i = 0 to N–1 do      //phase 2
   for each entry e in B[i] do
      remove e from B[i]
      insert e at the end of S
```

Bucket-Sort

Analysis:

- Phase 1 takes $O(n)$ time
- Phase 2 takes $O(n + N)$ time

Bucket-sort takes $O(n + N)$ time

Properties

Key-type Property

- The keys are used as indices into an array and cannot be arbitrary objects
- No external comparator

Stable Sort Property

- The relative order of any two items with the same key is preserved after the execution of the algorithm

Extensions

Extensions

- **Integer** keys in the range [a, b]
  - Put entry (k, o) into bucket $B[k - a]$
- **String** keys from a set D of possible strings, where D has constant size (e.g., names of the 50 U.S. states)
  - Sort D and compute the rank $r(k)$ of each string k of D in the sorted sequence
  - Put entry (k, o) into bucket $B[r(k)]$

Lexicographic Order

A d-tuple is a sequence of d keys $(k_1, k_2, \ldots, k_d)$, where key k_i is said to be the i-th dimension of the tuple

Example:

- The Cartesian coordinates of a point in space are a 3-tuple

The lexicographic order of two d-tuples is recursively defined as follows

$$(x_1, x_2, \ldots, x_d) < (y_1, y_2, \ldots, y_d)$$
$$\Leftrightarrow$$
$$x_1 < y_1 \ \vee \ x_1 = y_1 \wedge (x_2, \ldots, x_d) < (y_2, \ldots, y_d)$$

I.e., the tuples are compared by the first dimension, then by the second dimension, etc.

Lexicographic-Sort

- Let C_i be the comparator that compares two tuples by their i-th dimension
- Let $stableSort(S, C)$ be a stable sorting algorithm that uses comparator C
- Lexicographic-sort sorts a sequence of d-tuples in lexicographic order by executing d times algorithm $stableSort$, one per dimension
- Lexicographic-sort runs in $O(dT(n))$ time, where $T(n)$ is the running time of $stableSort$

Which value do we sort by?

1st, 2nd then 3rd?

Example:

(7, 4, 6) (5, 1, 5) (2, 4, 6) (2, 1, 4) (3, 2, 4)

(2, 4, 6) (2, 1, 4) (3, 2, 4) (5, 1, 5) (7, 4, 6)

(2, 1, 4) (5, 1, 5) (3, 2, 4) (2, 4, 6) (7, 4, 6)

(2, 1, 4) (3, 2, 4) (5, 1, 5) (2, 4, 6) (7, 4, 6)

Which value do we sort by?

3rd, 2nd, then 1st?

Example:

(7,4,6) (5,1,5) (2,4,6) (2, 1, 4) (3, 2, 4)

(2, 1, 4) (3, 2, 4) (5,1,5) (7,4,6) (2,4,6)

(2, 1, 4) (5,1,5) (3, 2, 4) (7,4,6) (2,4,6)

(2, 1, 4) (2,4,6) (3, 2, 4) (5,1,5) (7,4,6)

Lexicographic-Sort

Algorithm *lexicographicSort*(S)
  Input sequence S of d-tuples
  Output sequence S sorted in lexicographic order

  for $i \leftarrow d$ **downto** 1
    stableSort(S, C_i)

Radix-Sort

- Radix-sort is a specialization of lexicographic-sort that uses bucket-sort as the stable sorting algorithm in each dimension
- Radix-sort is applicable to tuples where the keys in each dimension i are integers in the range $[0, N-1]$
- Radix-sort runs in time $O(d(n+N))$

Radix-Sort

Algorithm *radixSort*(S, N)

Input sequence S of d-tuples such that $(0, \ldots, 0) \leq (x_1, \ldots, x_d)$ and $(x_1, \ldots, x_d) \leq (N-1, \ldots, N-1)$ for each tuple $(x_1, \ldots, x_d)$ in S

Output sequence S sorted in lexicographic order

for $i \leftarrow d$ **downto** 1

bucketSort(S, N)

Radix-Sort for Binary Numbers

- Consider a sequence of n b-bit integers

  $$x = x_{b-1} \ldots x_1 x_0$$

- We represent each element as a b-tuple of integers in the range $[0, 1]$ and apply radix-sort with $N = 2$
- This application of the radix-sort algorithm runs in $O(bn)$ time
- For example, we can sort a sequence of 32-bit integers in linear time

Radix-Sort for Binary Numbers

Algorithm *binaryRadixSort*(S)

Input sequence S of b-bit integers

Output sequence S sorted replace each element x of S with the item $(0, x)$

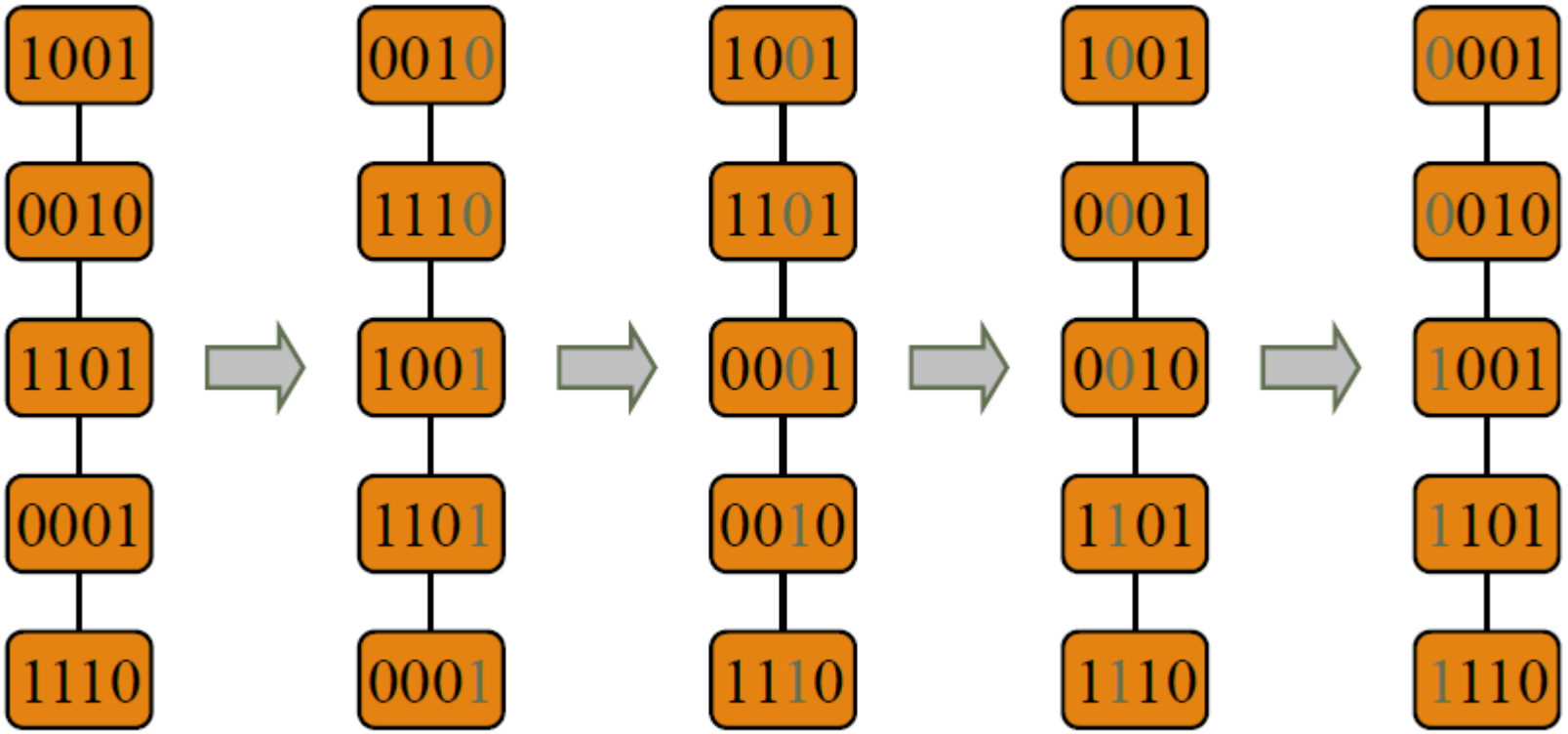
for $i \leftarrow 0$ **to** $b - 1$

replace the key k of each item (k, x) of S with bit x_i of x

bucketSort(S, 2)

Example

Sorting a sequence of 4-bit integers



References

These slides has been extracted, modified and updated from original slides of :

1. Data Structures and Algorithms in Java, 6th edition. John Wiley& Sons,
2. Introduction to Algorithms, 3rd Edition. Thomas H. Cormen and Charles E. Leiserson