```
Minclude <string.h>
Fdefine MAXPAROLA 30
#define MAXRIGA 80
   int freq[MAXPAROLA]; /* vettore di condatori
delle frequenze delle lunghezze delle parole
   char riga[MAXRIGA] ;
lint i, inizio, lunghezza
```

# **System and Device Programming**

### **Asynchronous I/O**

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# Synchronous I/O

- All standard I/O operations are synchronous
  - ➤ I/O is **blocking** and the task waits until the I/O operation completes
  - > Select delivers a synchronous form of notification
- Unfortunately, synchronous operations are inherently slow compared to other processing
  - Delays may be caused by
    - Hardware device, e.g., track and sector seek time on random access, etc.
    - Relatively slow data transfer rate between a physical device and the system memory
    - Network transfer using file servers, storage area networks, etc.

# **Asynchronous I/O**

- Threads can perform asynchronous I/O
  - A task can continue without waiting for an I/O operation to complete
- There are two techniques
  - Multithread I/O
  - POSIX asynchronous I/O

### **Asynchronous I/O**

### Multithread I/O

- > We use multiple threads
- ➤ Each thread within a process (or in different processes) may use a synchronous model
  - Each thread is responsible for a sequence of one or more synchronous, blocking I/O operations
  - Each thread should have its own file or pipe handle
- > Other threads can continue execution
  - The threads run asynchronous to each other
- > This is the most general technique

# **Asynchronous I/O**

### Asynchronous I/O

- We incur additional complexity when we use the POSIX asynchronous I/O interfaces
  - We have to worry about sources of errors for every asynchronous operation
  - The interfaces involve a lot of extra setup and processing rules compared to their conventional counterparts
  - Recovering from errors can be difficult
    - For example, if we submit multiple asynchronous writes and one fails, how should we proceed? If the writes are related, we might have to undo the ones that succeeded

- ❖ POSIX gives us a consistent way to perform asynchronous I/O, regardless of the type of file
- The interfaces use AIO control blocks to describe I/O operations
  - > The aiocb structure defines an AIO control block
  - ➤ It contains at least the fields shown in the following structure (implementations might include additional fields)
  - Note
    - Insert the <aio.h> library in the C file
    - Compile the C file with the realtime library (librt.a),
       i.e., -lrt

```
#include <aio.h>

struct aiocb {
   int aio_fildes;
   off_t aio_offset;
   volatile void *aio_buf;
   size_t aio_nbytes;
   int aio_reqprio;
   struct sigevent aio_sigevent;
   int aio_lio_opcode;
};
```

#### Parameters

The aio\_fildes field is the file descriptor open for the file to be read or written

- Read or writes start at the offset specified by aio\_offset
  - For a read Data is copied to the buffer that begins at the address specified by aio\_buf
  - For a write, data is copied from this buffer
- The aio\_nbytes field contains the number of bytes to read or write

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- The aio\_reqprio field is a hint that gives applications a way to suggest an ordering for the asynchronous I/O requests
  - The system has only limited control over the exact ordering, however, so there is no guarantee that the hint will be honored

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  int aio_lio_opcode;
};
```

- ➤ The aio\_sigevent field controls how the application is notified about the completion of the I/O event
  - It is described by a sigevent structure

```
struct sigevent {
  int sigev_notify;
  int sigev_signo;
  union sigval sigev_value;
  void (*sigev_notify_function) (union sigval);
  pthread_attr_t *sigev_notify_attributes;
};
```

- The sigev\_notify field controls the type of notification
  - SIGEV\_NONE the process is not notified when the asynchronous I/O request completes
  - SIGEV\_SIGNAL the signal specified by the sigev\_signo field is generated when the asynchronous I/O request completes
  - SIGEV\_THREAD the function specified by the sigev\_notify\_function field is called when the asynchronous I/O request completes address of a pthread attribute

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  struct sigevent aio_sigevent;
  int aio_lio_opcode;
};
```

#### **Read and Write**

```
#include <aio.h>
int aio_read(struct aiocb *aiocb);
int aio_write(struct aiocb *aiocb);
```

- To perform asynchronous I/O, we need to
  - > Initialize an AIO control block
  - Call either aio\_read or the aio\_write
- When these functions return success, the asynchronous I/O request has been queued for processing by the operating system
  - ➤ The return value has no relation to the result of the actual I/O operation

#### Guidelines

- While the I/O operation is pending, we have to be careful to ensure that the AIO control block and data buffer remain stable
  - Their underlying memory must remain valid and we cannot reuse them until the I/O operation completes

```
int aio_read(struct aiocb *aiocb);
int aio_write(struct aiocb *aiocb);
```

# **Synchronization**

```
#include <aio.h>
int aio_fsync (int op, struct aiocb *aiocb);
```

- We can use this function to force all pending asynchronous writes to persistent storage without waiting
- The aio\_fsync operation returns when the synch is scheduled
  - ➤ The data will not be persistent until the asynchronous synch completes
  - > The AIO control block controls how we are notified

# **Synchronization**

- The aio\_fildes field in the AIO control block indicates the file whose asynchronous writes are synched
- If the op argument is set to
  - O\_DSYNC, then the operation behaves like a call to fdatasync
  - > O\_SYNC, the operation behaves like a call to fsync

### Suspension

```
#include <aio.h>
int aio_suspend(const struct aiocb *const list[],
  int nent, const struct timespec *timeout);
```

- We use asynchronous I/O when we have other processing to do and we don't want to block while performing the I/O operation
- However, when we have completed the processing and find that we still have asynchronous operations outstanding, we can call the aio\_suspend function to block until an operation completes

#### **Cancel**

```
#include <aio.h>
int aio_cancel (int fd, struct aiocb *aiocb);
```

When we have pending asynchronous I/O operations that we no longer want to complete, we can attempt to cancel them with the aio\_cancel function