

Chapter 1

Appendix: The Basics of Mathematical Logic

- 6/19:
- **Mathematical logic:** The language one uses to conduct rigorous mathematical proofs.
 - “A logical argument may sometimes look unwieldy, excessively complicated, or otherwise appear unconvincing. The big advantage of writing logically, however, is that one can be absolutely sure that your conclusion will be correct” (305).
 - “Being logical is not the only desirable trait in writing, and in fact sometimes it gets in the way; mathematicians for instance often resort to short informal arguments which are not logically rigorous when they want to convince other mathematicians of a statement without going through all of the long details” (305).
 - “Because logic is innate, the laws of logic that you learn should make sense — if you find yourself having to memorize one of the principles or laws of logic here, without feeling a mental ‘click’ or comprehending why that law should work, then you will probably not be able to use that law of logic correctly and effectively in practice. So, please don’t study this appendix the way you might cram before a final — that is going to be useless. Instead, put away your highlighter pen, and read and understand this appendix rather than merely studying it” (306).

1.1 Mathematical Statements

- **Mathematical statement:** A precise statement “concerning various mathematical objects (numbers, vectors, functions, etc.) and relations between them (addition, equality, differentiation, etc.)” (306).
 - Mathematical statements are either true or false.
 - For example, $2 + 2 = 4$ is true while $2 + 2 = 5$ is false.
- **Ill-formed (statement):** A statement that is neither true nor false (or perhaps not even a statement). *Also known as ill-defined.*
 - For example, $0/0 = 1$.
 - The antonym is **well-formed** or **well-defined** statements.
- A logical argument should include only well-formed statements.
- “It is important, especially when just learning a subject, to take care in keeping statements well-formed and precise. Once you have more skill and confidence, of course you can afford once again to speak loosely^[1], because you will know what you are doing and won’t be in as much danger of veering off into nonsense” (307).

¹This is very similar to the English class/media arts idea of “you have to know the rules to be able to break them.”

- The principle of **proof by contradiction** is “to prove that a statement is true, it suffices to show that it is not false, while to show that a statement is false, it suffices to show that it is not true” (307).
- The axioms of logic become more dubious in very non-mathematical situations. One can attempt to apply logic via a mathematical model, but that’s science or philosophy. There are other models of logic that attempt to account for these scenarios, but they’re beyond the scope of this text.
- Note that statements may be true but not useful ($2 = 2$) or efficient ($4 \leq 4 - 4 = 4$ would be better), and they may be false but still be useful ($\pi = 22/7$). However, we concern ourselves at present with truth, alone, not usefulness or efficiency as those are to some extent matters of opinion.
- **Expression:** “A sequence of mathematical symbols which produces some mathematical object (a number, matrix, function, set, etc.) as its value” (308-09).
 - For example, $2 + 3 * 5$ is an expression while $2 + 3 * 5 = 17$ is a statement.
 - An expression is neither true nor false, but it can be well- or ill-defined.
- **Relation:** A thing that makes statements out of expressions, such as $=$, $<$, \geq , \in , \subset , etc.
- **Property:** A thing that makes statements out of expressions, such as “is prime,” “is continuous,” “is invertible,” etc^[2].
- **Logical connective:** A thing that relates multiple mathematical statements, such as and, or, not, if-then, if-and-only-if, etc.
- **Compound statement:** Two or more mathematical statements joined by logical connectives.
- **Conjunction:** “If X is a statement and Y is a statement, the statement ‘ X and Y ’ is true if X and Y are both true, and is false otherwise” (309).
 - Logician’s notation: “ $X \wedge Y$ ” or “ $X \& Y$.”
 - Note that “ X and Y ” can be reworded “ X and also Y ,” or “Both X and Y are true,” or even “ X , but Y ,” or a multitude of other ways and convey the same statement logically, if not expressively.
- **Disjunction:** “If X is a statement and Y is a statement, the statement ‘ X or Y ’ is true if either X or Y is true, or both” (309).
 - This is called **inclusive or**. There is a such thing as **exclusive or**, but it shows up far less regularly.
 - An exclusive or may be, “Either X or Y is true, but not both,” or “Exactly one of X or Y is true.”
 - To verify a disjunction, it suffices to verify just one case — this comes in handy when it is significantly easier to verify one case over the other.
- **Negation:** “The statement ‘ X is false’... is true if and only if X is false, and is false if and only if X is true” (310).
 - Logician’s notation: “ $\sim X$,” “ $!X$,” or “ $\neg X$.”
 - Negations convert “and” into “or”: The negation of “Jane has black hair and Jane has blue eyes” is “Jane doesn’t have black hair or Jane doesn’t have blue eyes.” Similarly, the negation of “ $x \geq 2$ and $x \leq 6$ ” is “ $x < 2$ or $x > 6$.”
 - Negations convert “or” into “and”: The negation of “John has black hair or brown hair” is “John does not have black hair and does not have brown hair” or, equivalently, “John has neither black nor brown hair.” Similarly, the negation of “ $x < -1$ or $x > 1$ ” is “ $x \geq -1$ and $x \leq 1$.”
 - Negations can produce clearly false statements: the negation of “ x is even or odd” is “ x is neither even nor odd.”

²With the introduction of properties, it is worth officially noting that mathematical statements can contain English words.

- Negations can get unwieldy — be careful!
- Note that “ X is true” can be abbreviated to simply “ X .”
- **If and only if:** “If X is a statement, and Y is a statement, we say that “ X is true if and only if Y is true,” [if] whenever X is true, Y has to be also, and whenever Y is true, X has to be also (i.e., X and Y are ‘equally true’)” (311). *Also known as iff.*
 - Logician’s notation: “ $X \leftrightarrow Y$.”
 - Denotes “logically equivalent statements.”
 - For example, “ $x = 3$ if and only if $2x = 6$ ” is true.
 - False statements can be logically equivalent: “ $2 + 2 = 5$ if and only if $4 + 4 = 10$.”
 - Sometimes, it is of interest to show that more than two statements are logically equivalent (See Exercise 1.1.5 and 1.1.6).

Exercises

6/22: 1. What is the negation of the statement, “either X is true, or Y is true, but not both?”

Proof. X and Y or neither X nor Y . □

2. What is the negation of the statement, “ X is true if and only if Y is true?” There may be multiple ways to phrase this negation.

Proof. X is true if and only if Y is false. □

3. Suppose that you have shown that whenever X is true, then Y is true, and whenever X is false, then Y is false. Have you now demonstrated that X and Y are logically equivalent? Explain.

Proof. Yes — although we have only described the dependence of Y on X , the dependence of X on Y is implied (suppose X is false when Y is true — but since X is false, Y is false, a contradiction; the same holds the other way around). □

4. Suppose that you have shown that whenever X is true, then Y is true, and whenever Y is false, then X is false. Have you now demonstrated that X is true if and only if Y is true? Explain.

Proof. No — X could be false and Y could still be true without leading to any contradictions. □

5. Suppose you know that X is true if and only if Y is true, and you know that Y is true if and only if Z is true. Is this enough to show that X, Y, Z are all logically equivalent? Explain.

Proof. Yes — by the supposition, $X \leftrightarrow Y$ and $Y \leftrightarrow Z$. Since X true implies Y true implies Z true, X false implies Y false implies Z false, Z true implies Y true implies X true, and Z false implies Y false implies X false, $X \leftrightarrow Z$. Thus, X, Y, Z are logically equivalent. □

6. Suppose you know that whenever X is true, then Y is true; that whenever Y is true, then Z is true; and whenever Z is true, then X is true. Is this enough to show that X, Y, Z are all logically equivalent? Explain.

Proof. Yes — clearly X, Y, Z are equally true. If they were not equally false (say if only one or two were false), then we would have a contradiction, as one true implies all the others are true. Thus, they are equally false, too. Therefore, they are logically equivalent. □

1.2 Implication

6/24:

- **Implication:** “If X is a statement, and Y is a statement, then ‘if X , then Y ’ is the implication from X to Y ” (312).
 - Logician’s notation: “ $X \implies Y$.”
 - Can be reworded “when X is true, Y is true,” “ X implies Y ,” “ Y is true when X is,” or “ X is true only if Y is true.”
 - “What this statement ‘if X , then Y ’ means depends on whether X is true or false. If X is true, then ‘if X , then Y ’ is true when Y is true, and false when Y is false. If however X is false, then ‘if X , then Y ’ is *always* true, regardless of whether Y is true or false!” (312).
 - “When X is true, the statement ‘if X , then Y ’ implies that Y is true. But when X is false, the statement ‘if X , then Y ’ offers no information about whether Y is true or not; the statement is true, but **vacuous**” (312).
- **Vacuous** (statement): A statement that “does not convey any new information” (312).
 - Vacuous statements can still be helpful in proofs, though.
- The only way to disprove an implication (to show that it’s false) is to show that the hypothesis (first statement) is true *and* the conclusion (second statement) is false.
- Whereas “if and only if” asserts that X and Y are equally true, an implication only asserts that Y is at least as true as X .
- Vacuously true implications are often used in a situation where the conclusion and hypothesis are both false, but the implication is true regardless, e.g., “if John had left work at 5pm, then he would be here by now.”
- They can also be used to facilitate a proof by contradiction: We know that “if John had left work at 5pm, then he would be here by now” and that “John is not here now.” Suppose that John left work at 5pm. Then he would be here by now, a contradiction. Thus, John did not leave work at 5pm.
- Implications can be true even when there is no causal link between the hypothesis and conclusion^[3], e.g., if $1 + 1 = 2$, then Washington, D.C. is the capital of the United States is true.
 - Using such acausal implications in a logical argument is typically frowned upon, since it will likely cause unneeded confusion.
- To prove, “if X , then Y ,” assume X is true and then deduce Y from X and other known hypotheses.
 - It does not matter whether or not X is true; an implication can be proved true irrespective of whether or not X is true.
 - It would be incorrect, though, to assume Y and deduce X . One cannot assume the conclusion and deduce the hypothesis.
- With regard to vacuously true statements, one need not be concerned that some hypotheses in their argument might not be correct, as long as their argument is still structured to give the correct conclusion regardless of whether those hypotheses were true or false (proving $n(n + 1)$ is even when $n \in \mathbb{N}$).
- **Converse** (of “if X , then Y ”): The statement “if Y , then X .”
 - “ X if and only if Y ” implies that “if X , then Y ” *and* its converse are true.
- **Inverse** (of “if X , then Y ”): The statement “if X is false, then Y is false.”
 - The inverse of a true implication is not necessarily a true implication.

³This is really weird — I’ll need to return to this later.

- **Contrapositive** (of “if X , then Y ”): The statement “if Y is false, then X is false.”
 - The contrapositive and the original statement are equally true.
- **Proof by contradiction**: “To show that something must be false, assume first that it is true, and show that this implies something which you know to be false (e.g., that a statement is simultaneously true and not true)” (316). *Also known as* **reductio ad absurdum**.
 - Particularly (but not exclusively) useful for proving “negative” statements, e.g., X is false, $a \neq b$, etc.
- On logician’s notation: general-purpose mathematicians do not often use them because English words are often more readable and don’t take up that much more space. That being said, \implies is a possible exception.

1.3 The Structure of Proofs

6/25:

- **Direct approach**: Assume the hypothesis and work one’s way toward a conclusion.
- For example, to prove “ A implies B ,” show A implies C , implies D , implies B . Done.
- We can also work backwards: “To show B , it would suffice to show D . Since C implies D , we just need to show C . But C follows from A ” (318).
 - Note that this *is* different from starting with the conclusion and proving the hypothesis.
- One can also move forward and backward at will.
 - “Thus there are many ways to write the same proof down; how you do so is up to you, but certain ways of writing proofs are more readable and natural than others, and different arrangements tend to emphasize different parts of the argument” (319).
- Multiple hypotheses and conclusions can split the proof into cases, making it more complicated.
- “There are several things to try when attempting a proof. With experience, it will become clearer which approaches are likely to work easily, which ones will probably work but require much effort, and which ones are probably going to fail” (319).
- “There may definitely be multiple ways to approach a problem, so if you see more than one way to begin a problem, you can just try whichever one looks the easiest, but be prepared to switch to another approach if it begins to look hopeless” (320).

1.4 Variables and Quantifiers

- **Propositional logic**: “Starting with primitive statements (such as ‘ $2 + 2 = 4$ ’ or ‘John has black hair’), then forming compound statements using logical connectives, and then using various laws of logic to pass from one’s hypotheses to one’s conclusions” (320). *Also known as* **Boolean logic**.
 - There are a dozen or so laws of propositional logic, but [?] excludes them because memorizing such a list is not how one should learn how to do logic. I will, however, include the list (from [?]) for reference.
 - Let P , Q , and R be statements. Let t be a formula that is a tautology and let f be a formula that is a contradiction. Then:
 1. **Law of Double Negation**: $\neg\neg P \equiv P$.
 2. **Associative Law for Conjunction**: $P \vee (Q \vee R) \equiv (P \vee Q) \vee R$.
 3. **Associative Law for Disjunction**: $P \wedge (Q \wedge R) \equiv (P \wedge Q) \wedge R$.

4. **Commutative Law for Conjunction:** $P \vee Q \equiv Q \vee P$.
 5. **Commutative Law for Disjunction:** $P \wedge Q \equiv Q \wedge P$.
 6. **First Distributive Law:** $P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)$.
 7. **Second Distributive Law:** $P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)$.
 8. **Idempotent Law for Conjunction:** $P \vee P \equiv P$.
 9. **Idempotent Law for Disjunction:** $P \wedge P \equiv P$.
 10. **Identity Law for Tautologies:** $P \wedge t \equiv P$.
 11. **Identity Law for Contradictions:** $P \vee f \equiv P$.
 12. **Inverse Law for Tautologies:** $P \vee \neg P \equiv t$.
 13. **Inverse Law for Contradictions:** $P \wedge \neg P \equiv f$.
 14. **Domination Law for Tautologies:** $P \vee t \equiv t$.
 15. **Domination Law for Contradictions:** $P \wedge f \equiv f$.
 16. **De Morgan's Law 1:** $\neg(P \vee Q) \equiv \neg P \wedge \neg Q$.
 17. **De Morgan's Law 2:** $\neg(P \wedge Q) \equiv \neg P \vee \neg Q$.
- Propositional logic is insufficient to do all math, though, because it does not incorporate **variables**.
 - **Mathematical logic:** Propositional logic plus **variables**.
 - **Variable:** “A symbol, such as n or x , which denotes a certain type of mathematical object — an integer, a vector, a matrix, that kind of thing” (320). They denote “quantities which are unknown, or set to some value, or assumed to obey some property” (320).
 - In almost all circumstances, the type of object that the variable represents should be declared; variables should have an explicit type.
 - $x = x$ is true regardless of the type, but $x + y = y + x$ is only true if x and y are variables for which addition makes sense and are commutative.
 - “The truth of a statement involving a variable may depend on the **context** of the statement” (321).
 - **Free variable:** A variable with a specified type but an unspecified value.
 - “Statements with free variables might not have a definite truth value” (321).
 - $x + 3 = 5$ does not have a definite truth value if x is a free variable.
 - $(x + 1)^2 = x^2 + 2x + 1$ does have a definite truth value even with x as a free variable.
 - **Set** (a free variable): Define a value to a free variable via a statement such as “Let $x = 2$ ” or “Set x equal to 2.”
 - **Bound variable:** A variable with a specified type that has been set to some value.
 - “Statements involving only bound variables and no free variables do have a definite truth value” (321).
 - A free variable can be bounded without being set to a *specific* value, per se, with the use of “for all” or “for some.”
 - “ $(x + 1)^2 = x^2 + 2x + 1$ ” is indefinite; “ $(x + 1)^2 = x^2 + 2x + 1$ for all $x \in \mathbb{R}$ ” is definite.
 - “ $x + 3 = 5$ ” is indefinite; “ $x + 3 = 5$ for some $x \in \mathbb{R}$ ” is definite.
 - **Universal quantifiers:** A quantifier of the form “for all x of type T .”
 - Logician's notation: “ \forall .”

- “ $P(x)$ is true for all x of type T ” is equivalent to the implication, “if x is of type T , then $P(x)$ is true.”
- Prove such a statement by letting x be any element of type T and showing that $P(x)$ holds.
- Disprove it with just a single counterexample (an element $x \in T$ for which $P(x)$ is false).
- If there are no variables x of type T , then “ $P(x)$ is true for all x of type T ” is **vacuously true**.
- **Existential quantifiers:** A quantifier of the form “for some x of type T .”
 - Logician’s notation: “ \exists .”
 - There exists at least one (but possibly numerous) x of type T for which $P(x)$ is true.
 - To indicate uniqueness, use “for exactly one $x \dots$ ” instead of “for some $x \dots$ ”
 - Prove such a statement by providing a single example of such an x .
 - A vacuously true “for all” statement will be false as a “for some” statement.
 - “ $6 < 2x < 4$ for all $3 < x < 2$ is true, but $6 < 2x < 4$ for some $3 < x < 2$ is false” (324).

1.5 Nested Quantifiers

- 6/28: • Two or more quantifiers can be combined to produce a type of compound statement.
- Negating a universal statement produces an existential statement, and vice versa, as alluded to above.
- 6/29: • **Aristotlean logic:** “Deals with objects, their properties, and quantifiers such as ‘for all’ and ‘for some’” (326).
- A subset of mathematical logic since it lacks the concepts of logical connectives and binary relations (e.g., = or <).
 - **Syllogism:** A typical line of reasoning in Aristotlean logic consisting of a major and minor premise, and a conclusion.
 - Swapping the order of quantifiers may or may not make a difference.
 - “For all real numbers a , and for all real numbers b , we have $(a + b)^2 = a^2 + 2ab + b^2$ ” is logically equivalent to “for all real numbers b , and for all real numbers a , we have $(a + b)^2 = a^2 + 2ab + b^2$.”
 - “For every integer n , there exists an integer m which is larger than n ” is not logically equivalent to “There exists an integer m which is larger than n for every integer n .”
 - “The reason why the order of quantifiers is important is that the inner variables may possibly depend on the outer variables, but not vice versa” (327).

Exercises

- 7/2: 1. What does each of the following statements mean, and which of them are true? Can you find gaming metaphors for each of these statements?
- (a) For every positive number x , and every positive number y , we have $y^2 = x$.
- Proof.* It means that for all positive numbers x , the statement “ $y^2 = x$ for all positive number y ” is true. However, this is a false statement ($2^3 \neq 3$, for example). As a gaming metaphor, suppose you play a game where your opponent picks two positive numbers x and y . If $y^2 = x$, then you win. If you can always win, then the statement is true. \square
- (b) There exists a positive number x such that for every positive number y , we have $y^2 = x$.

Proof. It means that for some positive number x , the statement “ $y^2 = x$ for all positive numbers y ” is true. However, this is a false statement (no positive number exists such that every positive number squared equals it). As a gaming metaphor, suppose you play a game where you first pick a positive number x , and then your opponent picks a positive number y . If you can pick an x such that any y picked by your opponent satisfies $y^2 = x$, then you win. If you can win once, then the statement is true. \square

- (c) There exists a positive number x , and there exists a positive number y , such that $y^2 = x$.

Proof. It means that the statement “ $y^2 = x$ ” is true for some pair of positive numbers x, y . This statement is true. As a gaming metaphor, suppose you play a game where you pick two positive numbers x and y . If $y^2 = x$, then you win. If you can win once, then the statement is true. \square

- (d) For every positive number y , there exists a positive number x such that $y^2 = x$.

Proof. It means that for all positive numbers y , the statement “ $y^2 = x$ for some positive number x ” is true. This statement is true. As a gaming metaphor, suppose you play a game where your opponent first picks a positive number y , and then you pick a positive number x . If $y^2 = x$, then you win. If you can always win, then the statement is true. \square

- (e) There exists a positive number y such that for every positive number x , we have $y^2 = x$.

Proof. It means that for some positive number y , the statement “ $y^2 = x$ for all positive numbers x ” is true. However, this is a false statement (no positive number exists such that the square root of every positive number equals it). As a gaming metaphor, suppose you play a game where you first pick a positive number y , and then your opponent picks a positive number x . If you can pick a y such that any x picked by your opponent satisfies $y^2 = x$, then you win. If you can win once, then the statement is true. \square

1.6 Some Examples of Proofs and Quantifiers

- This section proves some simple results involving the “for all” and the “for some” quantifiers to demonstrate how the quantifiers are arranged and how the proofs are structured. Let’s begin.

- **Proposition 1.1.** *For every $\varepsilon > 0$ there exists a $\delta > 0$ such that $2\delta < \varepsilon$.*

Proof. Let $\varepsilon > 0$ be arbitrary. We have to show that there exists a $\delta > 0$ such that $2\delta < \varepsilon$. We only need to pick one such δ ; choosing $\delta := \frac{\varepsilon}{3}$ will work, since one then has $2\delta = \frac{2}{3}\varepsilon < \varepsilon$. \square

- Since we are proving something for *every* ε , we must make ε arbitrary.
- On the other hand, δ can be chosen as we wish, because we only need to show that there exists *some* δ which does what we want.
- The δ can depend on ε because “the δ -quantifier is nested inside the ε -quantifier” (328).

- **Proposition 1.2.** *There exists an $\varepsilon > 0$ such that $\sin x > \frac{x}{2}$ for all $0 < x < \varepsilon$.*

Proof. We pick $\varepsilon > 0$ to be chosen later, and let $0 < x < \varepsilon$. Since the derivative of $\sin x$ is $\cos x$, we see from the mean-value theorem we have

$$\frac{\sin x}{x} = \frac{\sin x - \sin 0}{x - 0} = \cos y$$

for some $0 \leq y \leq x$. Thus, in order to ensure that $\sin x > \frac{x}{2}$, it would suffice to ensure that $\cos y > \frac{1}{2}$. To do this, it would suffice to ensure that $0 \leq y < \frac{\pi}{3}$ (since the cosine function takes the value of 1 at 0, takes the value of $\frac{1}{2}$ at $\frac{\pi}{3}$, and is decreasing in between). Since $0 \leq y \leq x$ and $0 < x < \varepsilon$, we see that $0 \leq y < \varepsilon$. Thus if we pick $\varepsilon := \frac{\pi}{3}$, then we have $0 \leq y < \frac{\pi}{3}$ as desired, and so we can ensure that $\sin x > \frac{x}{2}$ for all $0 < x < \varepsilon$. \square

- Normally, when proving a “there exists (ε) such that X is true” statement, one proceeds by selecting ε carefully, and then showing that X is true for that ε . However, when the selection would require a lot of foresight, we defer it until later in the argument, as above.
- Note that since the value of ε didn’t depend on the nested variables x and y , the proof is legitimate.

1.7 Equality

- **Equality:** The most important relation. It links two objects x, y of the same type T . “Given two such objects x and y , the statement $x = y$ may or may not be true; it depends on the value of x and y and also on how equality is defined for the class of objects under consideration” (329).
 - For example, $12 \neq 2$ in ordinary arithmetic, but $12 = 2$ in modular arithmetic with modulus 10.
- For the purposes of logic, equality obeys the following four **axioms of equality**.
- **Reflexive axiom:** Given any object x , we have $x = x$.
- **Symmetry axiom:** Given any two objects x and y of the same type, if $x = y$, then $y = x$.
- **Transitive axiom:** Given any three objects x, y, z of the same type, if $x = y$ and $y = z$, then $x = z$.
- **Substitution axiom:** Given any two objects x and y of the same type, if $x = y$, then $f(x) = f(y)$ for all functions or operations f . Similarly, for any property $P(x)$ depending on x , if $x = y$, then $P(x)$ and $P(y)$ are equivalent statements.
 - The first three axioms assert that equality is an **equivalence relation**.
- “Thus, from the point of view of logic, we can define equality on a class of objects however we please, so long as it obeys the reflexive, symmetry, and transitive axioms, and is consistent with all other operations on the class of objects under discussion in the sense that the substitution axiom was true for all of those operations” (330).

Exercises

1. Suppose you have four real numbers a, b, c, d and you know that $a = b$ and $c = d$. Use the above four axioms to deduce that $a + d = b + c$.

Proof. Let $f(x) = x + c$. By the substitution axiom, since $a = b$, $f(a) = f(b) \Rightarrow a + c = b + c$. Let $g(x) = a + x$. By the substitution axiom, since $c = d$, $f(c) = f(d) \Rightarrow a + c = a + d$. By the symmetry axiom, $a + d = a + c$. By the transitive axiom, $a + d = a + c = b + c$. \square

1.8 Misc. Notes

- “From a logical point of view, there is no difference between a lemma, proposition, theorem, or corollary — they are all claims waiting to be proved. However, we use these terms to suggest different levels of importance and difficulty. A lemma is an easily proved claim which is helpful for proving other propositions and theorems, but is usually not particularly interesting in its own right. A proposition is a statement which is interesting in its own right, while a theorem is a more important statement than a proposition which says something definitive on the subject, and often takes more effort to prove than a proposition or lemma. A corollary is a quick consequence of a proposition or theorem that was proven recently.” (25).