

The Knot Book

Notes

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1 Introduction

1.1 Introduction

- **Knot:** “A knotted loop of string, except that we think of the string as having no thickness, its cross-section being a single point” (2).
- Do not distinguish between a ‘nice, even’ knot and one that has been deformed through space.
- **Unknot:** “The simplest knot of all... the unknotted circle” (2). *Also known as trivial knot.* See Figure 1.1a.
- **Trefoil knot:** “The next simplest knot” (2). See Figure 1.1b.

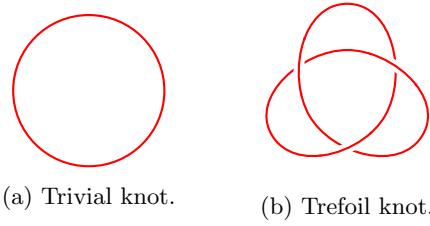


Figure 1.1: Projections of the two simplest knots.

- **Projection:** A picture of a knot, such as those in Figure 1.1.
 - The same knots can have multiple projections (as they are deformed in space).
- **Crossings:** The places in a projection where a knot crosses itself.
 - The trefoil knot in Figure 1.1b is a three-crossing knot because it crosses itself 3 times.
 - Any one-crossing knot is trivial.
 - *Exercise 1.2:* Any two-crossing knot must be trivial because the simplest nontrivial knot is the trefoil knot, which has three crossings.
- Atoms were originally thought to be tangles (knots) in the ether of the universe, but when chemists moved on, mathematicians took up knot theory. In the 1980s, biochemists began to see applications of knot theory in their research (see Section 7).
- **Topology:** “The study of the properties of geometric objects that are preserved under deformations” (6).
 - Knot theory is a subfield of topology (see Section ??).

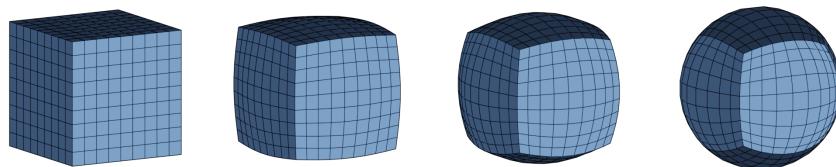


Figure 1.2: Deformation of a cube into a sphere.

- Any knot can have a projection with as many crossings as desired.

- **Alternating knot:** “A knot with a projection that has crossings that alternate between over and under as one travels around the knot in a fixed direction” (7).
 - The trefoil is such a knot.
- *Exercise 1.7**: By changing some of the crossings from over to under or vice versa, any projection of a knot can be made into a projection of the unknot^[1]. See Figure 1.3.

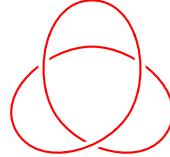


Figure 1.3: A projection of the unknot evoking the trefoil knot.

1.2 Composition of Knots

- **Composition** (of two knots): “A new knot obtained by removing a small arc from two knot projections and then connecting the four endpoints by two new arcs” (7).
 - If two knots are designated J and K , then their composition is denoted $J \# K$.
 - Do not overlap the projections and choose two arcs that are on the outside to avoid new crossings.
 - Make sure that the new arcs do not cross any of the the original knot projections or each other.
- **Composite knot:** A knot that “can be expressed as the composition of two knots, neither of which is the trivial knot” (8).
 - This definition is analogous to composite integers, where an integer is composite if it is the product of positive integers, neither of which is 1.
 - Similarly, if we compose any knot with the unknot, we get the same knot back.
- **Factor knots:** “The knots that make up the composite knot” (8).
- **Prime knot:** “A knot [that] is not the composition of any two nontrivial knots” (9).
- The unknot, trefoil knot, and figure-eight knots are all prime (see Section 4.3).
 - The unknot is not composite for the same reason that 1 is not the product of two integers greater than 1.

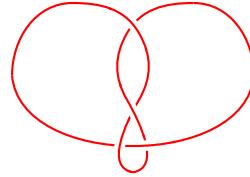


Figure 1.4: The figure-eight knot.

- Similar to integers, “a composite knot factors into a unique set of prime knots” (10).
- *Exercise 1.8*: Using the appendix table, identify the factor knots that make up the composite knot in Figure 1.5.
 - The knot in Figure 1.5 is the composition of two projections of 5_2 .

¹How can I *show* something? How can I do these proofs? What kind of logic solves one of these? See Section 3.1 for a direct proof of/solution to *Exercise 1.7*.

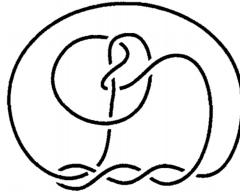
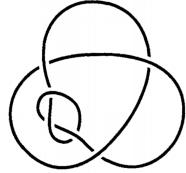


Figure 1.5: The composite knot.

- *Exercise 1.9:* Show that the knot in Figure 1.6a is composite.



(a) The composite knot.



(b) Factors.

Figure 1.6: Factorization of a ‘double trefoil.’

- There is more than one way to take the composition of two knots (by removing different arcs).
 - This is not analogous to multiplication — a break in the pattern.
- **Orientation:** A direction to travel around the knot. Denoted by placing “coherently directed arrows along the projection of the knot in the direction of our choice” (10). A knot with such arrows is **oriented**.
 - All compositions $J \# K$ where the orientations of J and K do match up will yield the same composite knot.
 - J can be ‘slid around’ $J \# K$ until it reaches the second position where the composition was taken.
 - All compositions $J \# K$ where the orientations of J and K do not match up will yield the same composite knot.
 - These two compositions can be distinct.



Figure 1.7: Orientation notation.

- **Invertible:** A knot that can be deformed back to itself so that an orientation on it is sent to the opposite orientation.
 - “In the case that one of the two knots is invertible, say J , we can always deform the composite knot so that the orientation on K is reversed, and hence so that the orientations of J and K always match. Therefore, there is only one composite knot that we can construct from the two knots” (11).
- To determine the possible compositions of knots, it is necessary to know which knots are invertible, but no general technique has yet been discovered.

1.3 Reidemeister Moves

- **Ambient isotropy:** “The movement of the string through three-dimensional space without letting it pass through itself” (12).
- **Planar isotropy:** A deformation of “the projection plane as if it were made of rubber with the projection drawn upon it” (12).
 - Stretching, squeezing, rotating, bending single arcs, etc.
- **Reidemeister move:** “One of three ways to change a projection of the knot that *will* change the relation between the crossings” (13).
- **First Reidemeister move:** “Put in or take out a twist in the knot” (13). See Figure 1.8a. *Also known as type I Reidemeister move.*
- **Second Reidemeister move:** “Either add two crossings or remove two crossings” (13). See Figure 1.8b. *Also known as type II Reidemeister move.*
- **Third Reidemeister move:** “Slide a strand of the knot from one side of a crossing to the other side of the crossing” (13). See Figure 1.8c. *Also known as type III Reidemeister move.*
 - Note that the crossings in Figure 1.8 can be reversed and the move will still be classified under the same category.

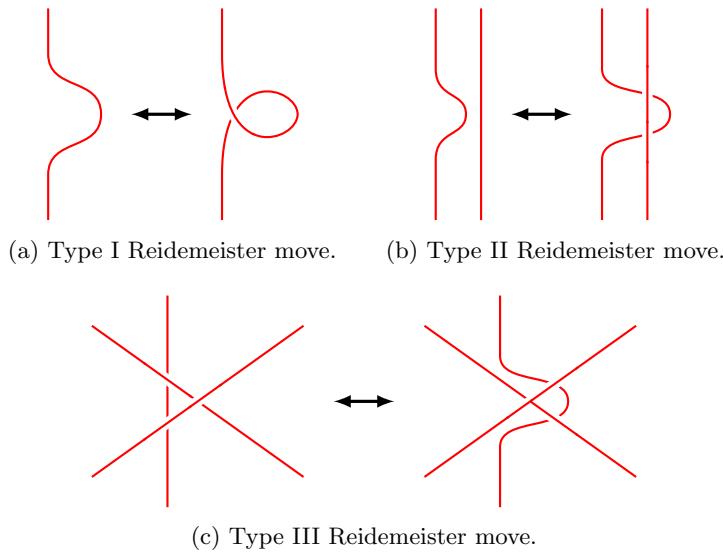
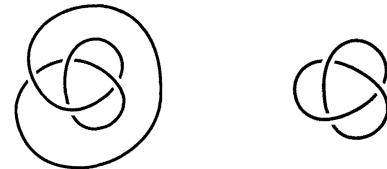


Figure 1.8: Reidemeister moves.

- All Reidemeister moves are ambient isotopies.
- “If we have two distinct projections of the same knot, we can get from the one projection to the other by a series of Reidemeister moves and planar isotopies” (14).
- **Amphicheiral:** A knot that “is equivalent to its mirror image, that is, the knot obtained by changing every crossing... to the opposite crossing” (14-15). *Also known as achiral by chemists.*
 - A knot and its mirror image are distinct unless the knot is amphicheiral.
 - See Sections 6.4 and 7 for more on amphicheirality.

- *Exercise 1.10:* Show that the two projections in Figure 1.9 represent the same knot by finding a series of Reidemeister moves from one to the other.



(a) Initial projection. (b) Final projection.

Figure 1.9: Finding Reidemeister moves.



Figure 1.10: Solution to *Exercise 1.10*.

- *Exercise 1.11*:* Find a sequence of Reidemeister moves to untangle the unknot shown in Figure 1.11.

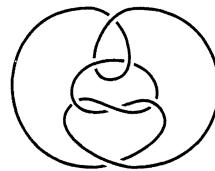


Figure 1.11: Unknot to be untangled.

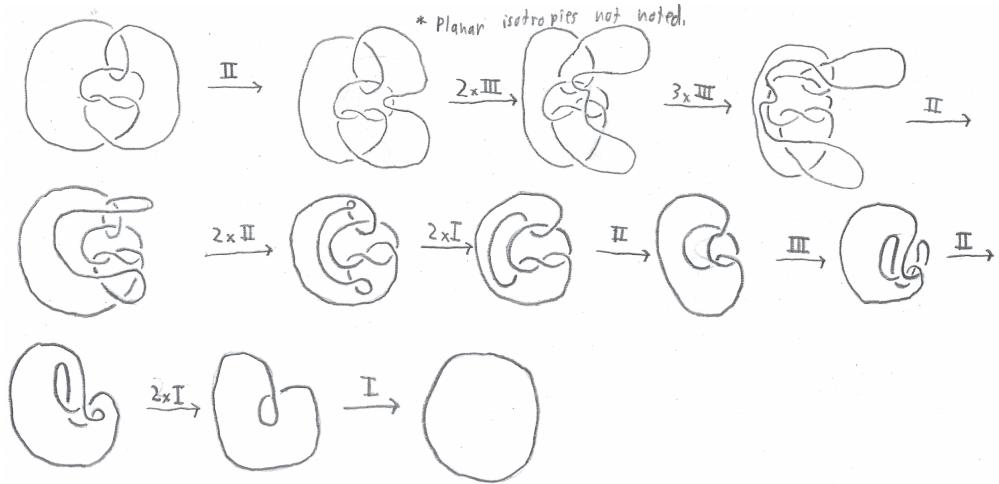
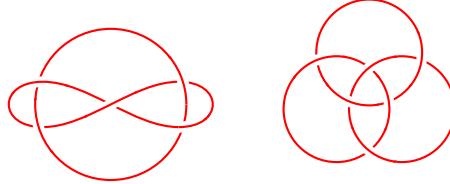


Figure 1.12: Solution to *Exercise 1.11**.

- The bounds on the increase in crossings generated by Reidemeister moves from one projection to another are unknown.

1.4 Links

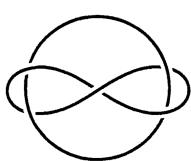
- **Link:** “A set of knotted loops all tangled up together” (17).
- “Two links are considered to be the same if we can deform the one link to the other link without ever having any one of the loops intersect itself or any of the other loops in the process” (17).



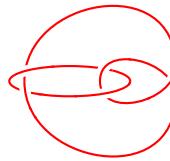
(a) Whitehead link. (b) Borromean rings.

Figure 1.13: Projections of two common links.

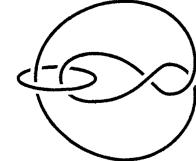
- *Exercise 1.13:* Show that the two projections in Figure 1.14a and 1.14c represent the same link.
 - Untwist the right side of the figure-eight loop, twisting the right side of the circle at the same time. Then add a further twist to the circle loop and flip (rotate along the horizontal axis) the figure-eight loop (which now looks like an ellipse).



(a) Initial projection.



(b) Intermediate projection.



(c) Final projection.

Figure 1.14: Ambient isotopies of the Whitehead link.

- **Link of n components:** A link made up of n loops knotted with each other.
 - The Whitehead link (Figure 1.13a) is a link of 2 components.
 - The Borromean rings (Figure 1.13b) are a link of 3 components.
 - A knot is a link of 1 component.
 - If the number of components in two links differ, then the links are clearly distinct.
- **Splittable:** A link whose components “can be deformed so that they lie on different sides of a plane in three-space” (17).
- *Exercise 1.14:* Show that the link in Figure 1.15 is splittable.

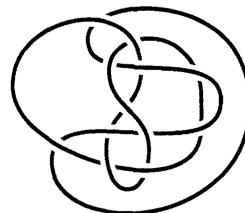


Figure 1.15: Link to be split.

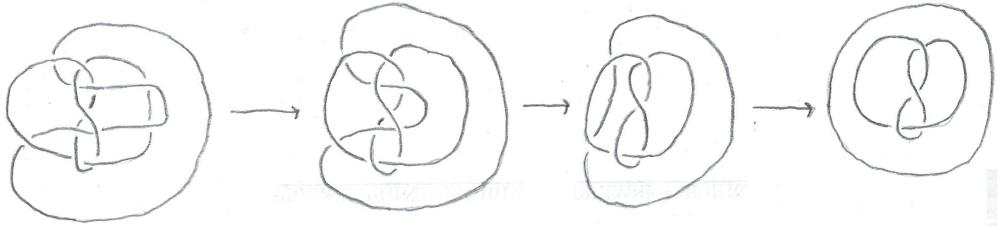


Figure 1.16: Solution to *Exercise 1.14*.

- **Unlink:** One of the two simplest links of two components and the simplest splittable link of two components. *Also known as trivial link.* See Figure 1.17a.
- **Hopf link** The other of the two simplest links of two components and the simplest nonsplittable link of two components. See Figure 1.17b.



Figure 1.17: Projections of the two simplest links of two components.

- **Linking number:** A quantity that numerically measures how linked up two components are.

- If M and N are two components in a link, begin by orienting both of them.
- At each crossing, count a $+1$ if Figure 1.18a holds or count a -1 if Figure 1.18b holds.
 - Note that if you are unsure, do the following: If rotating the bottom strand clockwise lines up the arrows (correlates the two orientations), count $+1$.
 - In the same vein, if rotating the bottom strand counterclockwise lines up the arrows, count -1 .
 - If a component crosses itself, do not count it.

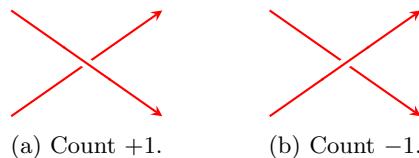


Figure 1.18: Computing linking numbers per crossing.

- Sum all of the $+1$ s and -1 s and divide this sum by 2 to yield the linking number.
 - For example, the linking number for the Hopf link (Figure 1.17b) is ± 1 depending on the chosen orientation.
 - Note that reversing the orientation for one of the two links is equivalent to multiplying the linking number by -1 .
 - As such, the absolute value of the linking number remains constant whatever orientation is chosen.
- *Exercise 1.15:* Compute the linking number of the link pictured in Figure 1.19. Now reverse the direction on one of the components and recompute it.

- Add up all of the numbers in the second row of Table 1.1 to yield 2. Divide by 2 to yield the linking number, 1.

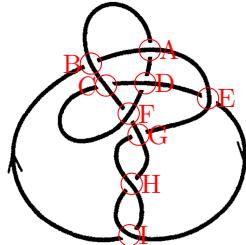


Figure 1.19: Link of linking number n .

A	B	C	D	E	F	G	H	I
-1	-1	0	0	+1	0	+1	+1	+1

Table 1.1: Counting crossings.

- Although a linking number is computed using a single projection of a link, the linking number will be the same for any projection of the link.
 - This can be proven by demonstrating that the Reidemeister moves do not change the linking number. “Since we can get from any one projection of a link to any other via a sequence of Reidemeister moves, none of which will change the linking number, it must be that two different projections of the same link yield the same linking number” (20).
 - Let’s take this case by case.
 - A Type I Reidemeister move generates a self-intersection. Because self-crossings do not count toward the linking number by definition, the first Reidemeister move does not affect the linking number.
 - A Type II Reidemeister move generates two crossings. There are now two cases: Either the new crossings are self-intersections or they are not. If the new crossings are self-intersections, then there is no change to the linking number. If the new crossings are not self-intersections, they will always have opposite linking numbers, and, thus, the two new crossings cancel each other out. Therefore, the second Reidemeister move does not affect the linking number.
 - A Type III Reidemeister move removes two crossings and generates two crossings. There are now two cases: Either the new crossings are self-intersections or they are not. If the new crossings are self-intersections, then there is no change to the linking number. If the new crossings are not self-intersections, both pairs of crossings will always have opposite linking numbers and, thus, both pairs of crossings cancel each other out. Therefore, the third Reidemeister move does not affect the linking number.
 - This logic can be visually confirmed by assigning orientations and counting crossings in Figure 1.8.
- **Invariant:** A quality of a knot or link that, once orientations are chosen, is unchanged by ambient isotopy.
 - Both the linking number and the number of components are link invariants.
- *Exercise 1.16:* Explain why the linking number of a splittable two-component link will always be 0, no matter what projection is used to compute it.

- **SOLUTION 1:** By definition, the linking number numerically measures how linked up two components are. Since splittable links are not joined (or “linked up”) in any way, the linking number must be 0.
- **SOLUTION 2:** If a two-component link is splittable, then there exists a projection of it with no crossings that are not self-intersections (a projection of the split link, as in Figure 1.17b). If all crossings are self-intersections, then the linking number computed for this projection must be 0. Now that it is known that said link has a linking number of 0, any combination of Reidemeister moves can be used to manipulate the link into any other projection. But because Reidemeister moves do not affect the linking number, the linking number will remain 0.
- The absolute value of the linking number, as a link invariant, can be used to distinguish certain distinct links, regardless of orientation.
 - “Any two links with two components that have distinct absolute values of their linking numbers have to be different links” (21).
 - For example, the difference in linking number between the unlink (0; Figure 1.17a) and the Hopf link (1; Figure 1.17b) distinguishes them.
- However, the linking number cannot distinguish between all links.
 - For instance, both the Whitehead link (Figure 1.13a) and the unlink (Figure 1.17a) have linking number 0.
 - One such distinction is discussed in Section 1.5.
- **Brunnian link:** A nontrivial link where the removal of any one component leaves behind a set of trivial unlinked circles.
 - The Borromean rings (Figure 1.13b) are Brunnian^[2].

1.5 Tricolorability

- How do we *prove* that every knot is not just a projection of the unknot?
- **Strand:** “A piece of the link that goes from one undercrossing to another with only overcrossings in between” (23).
- **Tricolorable:** A projection of a knot or link in which each strand “can be colored one of three different colors, so that at each crossing, either three different colors come together or all the same color comes together” (23).



Figure 1.20: Tricolored trefoil.

- *Exercise 1.21:* Determine which of the projections of the three six-crossing knots 6_1 , 6_2 , and 6_3 are tricolorable.
 - 6_1 is tricolorable and the other two knots are not. I determined this through trial and error.
- *Exercise 1.22:* Show that the projection of the knot 7_4 in Figure 1.21a is tricolorable.

²How can I draw a Brunnian link with four components? Or more?



(a) Initial projection.

(b) Tricolored projection.

Figure 1.21: Tricolored 7_4 .

- Reidemeister moves do not affect tricolorability.
 - A Type I Reidemeister move generates a self-intersection. At such a crossing (made out of one original strand), every color will be the same; only one color meets at the crossing.
 - A Type II Reidemeister move generates two crossings. If the strands are the same color, then everything stays the same color. If the strands are different colors, than newly created loop takes on the third color.
 - Type III has many cases (*Exercise 1.23*).
- Since the unknot is not tricolorable and the trefoil knot is, we have just proven that there is at least one other knot besides the unknot.
 - Tricolorability is a knot invariant.
- For links, tricolorability is a bit different — the unlink, for example, is tricolorable.

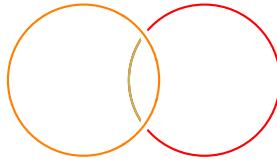


Figure 1.22: Tricolored unlink.

1.6 Knots and Sticks

- We can also consider knots made out of straight sticks glued to each other at each end.
- The first nontrivial knot made out of sticks is the trefoil, with six sticks.
 - We know that the stick trefoil in Figure 1.23 could be made in the real world because, if two vertices lie in the xy -plane, then two lie above and two lie below.
 - This P(lanar), L(ow), H(igh) notation is commonly used.



Figure 1.23: A trefoil knot made of sticks.

- **Stick number:** “The least number of straight sticks necessary to make a knot K ” (29). *Also known as $s(K)$.*

- The stick number for the composition of n trefoil knots is $2n + 4$.
- Many problems listed^[3].

³These could make good topics for a larger paper. Alternatively, some I am not ready to solve and need to learn more math first (those that I will come back to at the end of the book). I do not believe I'm missing anything essential by skipping these for now, but we'll see.

2 Tabulating Knots

2.1 History of Knot Tabulation

- Although many mathematicians (including Gauss) had dabbled with knots, Lord Kelvin's theory that atoms were knotted vortices in the ether (referenced in Section 1.1) kicked off research in earnest.
- Reverend Thomas P. Kirkman began tabulating knots but wrote poorly.
- However, Kirkman's ideas were applied by Scottish physicist Peter Guthrie Tait to list all alternating knots up to 10 crossings.
- C. N. Little of the University of Nebraska was the first to tackle the nonalternating knots.
 - He tabulated up to 10 crossings (43 knots total).
 - However, in 1974, it was discovered that two knots were, in fact, the same and there were really 42 knots.
- Kenneth A. Perko, a part time mathematician and New York lawyer, discovered that the projections were not distinct. As such, the projections are known as the Perko pair (Figure 2.1).

- *Exercise 2.1:* Show that the Perko pair (projected in Figure 2.1) are the same knot.

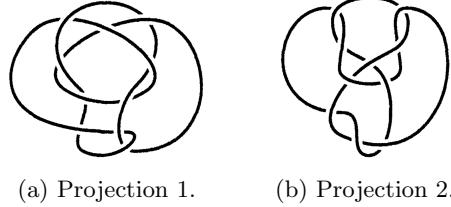


Figure 2.1: The Perko pair.

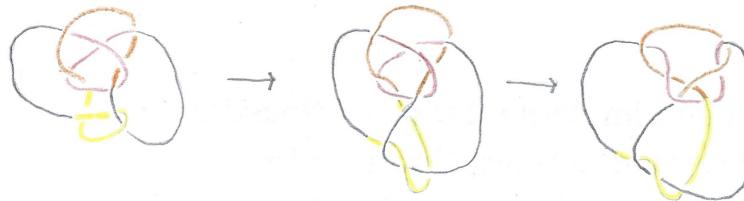


Figure 2.2: Solution to *Exercise 2.1*.

- Little later published a census of 11-crossing alternating knots with 11 omissions and 1 duplication.
- Mary G. Haseman listed all amphicheiral knots of 12 crossings.
- Early on, there were few attempts to *prove* that the tabulated knots were actually distinct.
- In 1927, Alexander and Briggs proved that the knots up to 9 crossings were distinct.
 - They developed the first knot polynomial, the Alexander polynomial.
 - This would remain the only knot polynomial until 1984.
- Once Kurt Reidemeister finished rigorously classifying the knots up to 9 crossings in 1932, tabulation was inactive for a while.

- In 1969, John H. Conway developed a new notation and used it to determine all of the prime knots up to 11 crossings and all prime nonsplittable links up to 10 crossings.
- In 1978, Alain Caudron of the University of Paris corrected Conway's tabulation a bit.
- Meanwhile, a new notation was developed by Hugh Dowker (see Section 2.2) that was implemented as a computer algorithm by Morwen Thistlethwaite.
- In the 1980s-1990s (with computer help), Jim Hoste, Thistlethwaite, and Jeff Weeks (an expert in hyperbolic knots; see Section 5.3) tabulated up to 16-crossing knots.
 - Note that the number of knots of successive numbers of crossings *appears* to grow exponentially (this is unproven).
 - Also note that Hoste et al's listed knots includes those that are not amphicheiral only once, so those knots actually represent 2 distinct knots.
 - Determining amphicheirality is discussed in Section 6
- On drawing a 14-crossing knot.
- Claus Ernst and Dewitt Sumners proved a lower exponential bound on the number of distinct prime knots (see Section 3.2).
- Dominic Welsh proved an upper exponential bound on the number of distinct prime knots.

2.2 The Dowker Notation for Knots

- For alternating knots...
 - Begin by choosing an orientation.
 - Label an arbitrary crossing 1.
 - Leaving that crossing on the understrand, label the next crossing 2.
 - Continue through the crossing on the same strand and label the next crossing 3.
 - Continue in this fashion until you have gone all the way around the knot once.
 - Note that this gives each crossing two numbers.

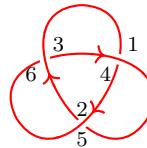


Figure 2.3: Trefoil knot labeled in Dowker notation.

- *Exercise 2.2:* Why does every crossing get one even numbered label and one odd numbered label?
 - Because the knot is alternating — every understrand will add an odd label and every overstrand will add an even number (because you set off on the understrand after labeling your starter crossing 1).

$$\begin{array}{ccc} 1 & 3 & 5 \\ 4 & 6 & 2 \end{array}$$

Table 2.1: Dowker numerical pairing of a trefoil knot.

- From the counting in Figure 2.3, we obtain a pairing between the numbers (Table 2.1).
- Note that the numbers in the top row of Table 2.1 are constantly increasing odd numbers.
- Since the top row numbers are in a predictable pattern, we can shorthand the notation for the trefoil knot in Figure 2.3 to 4 6 2.
- “Thus, from a projection of a knot, we obtain a sequence of even integers, where the number of even integers is exactly the number of crossings in the knot” (36).
- *Exercise 2.3:* Find a sequence of even integers that represents the projection of the knots 6_2 and 6_3 (Figure 2.4). How about a second sequence of even integers that also represents the same projection of 6_3 ?
 - For Figure 2.4a, write 6 8 10 12 2 4.
 - For Figure 2.4b, write 8 6 10 12 4 2.
 - Another sequence of even integers for Figure 2.4b is 4 10 8 12 2 6.

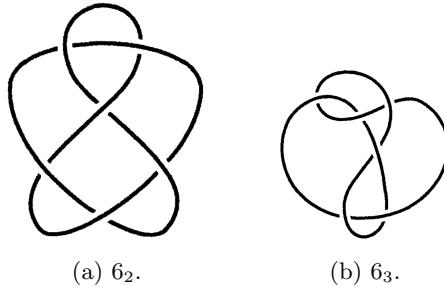


Figure 2.4: Knots 6_2 and 6_3 .

- Knots can also be reconstructed from Dowker notation. See Figure 2.5 for an example of reconstructing an alternating knot from the notation, 8 10 12 2 14 6 4.

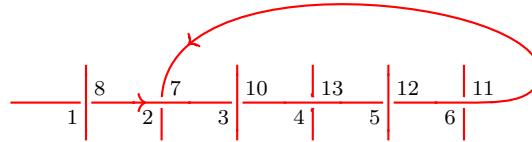


Figure 2.5: Constructing a knot projection from its Dowker notation.

- Our choice in drawing can change the resulting knot.
 - 4 6 2 10 12 8 gives two distinct knots.
 - “Note that the two knots are composite knots, and that this is reflected in the fact that the sequence 4 6 2 10 12 8 is actually a shuffling of the three numbers 2, 4, 6 and then a shuffling of the three numbers 8, 10, and 12” (38).
 - When the Dowker notation can be broken into two subpermutations (as above) the knot is composite (unless one subpermutation is trivial).
 - Any sequence that cannot be split in this way represents either a particular knot or its mirror image.
 - If the knot is amphicheiral, then the sequence only represent one knot.

- A knot and its mirror image are equivalent when projected onto a sphere as opposed to a plane.
- *Exercise 2.6:* Draw two projections given by 10 12 2 14 6 4 8, which are inequivalent as projections in the plane but which are equivalent as projections on the sphere.

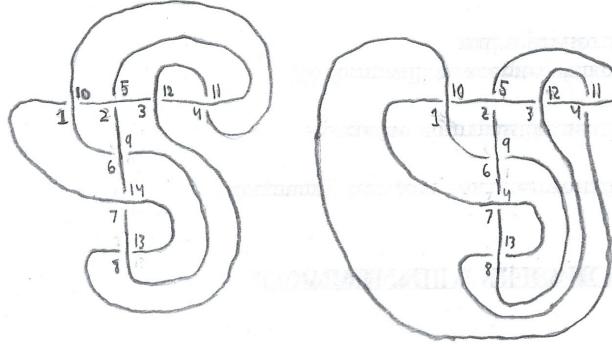
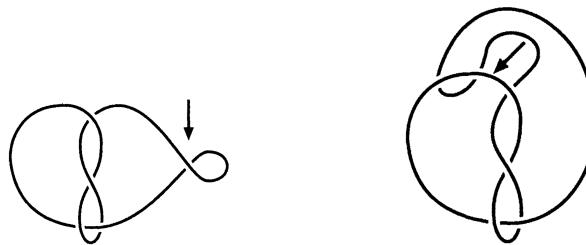


Figure 2.6: Solution to *Exercise 2.6*.

- *Exercise 2.7:* How many different sequences of the integers 2 4 6 8 10 12 14 are there? (This exercise gives us an upper bound on the number of possible alternating knot projections with seven crossings; however, it's far from accurate.)
 - There are $7! = 5040$ sequences.
- There is a slightly modified Dowker notation for nonalternating knots.
 - “If the even integer is assigned to the crossing while we are on the overstrand at that crossing, we leave the even integer positive. But if the even integer is assigned to the crossing while we are on the understrand of that crossing, we make the corresponding even numbers negative” (39).
- *Exercise 2.9:* How can you recognize from the sequence of numbers that a projection has a trivial crossing in it like in Figure 2.7a? How about recognizing a Type II Reidemeister move that will reduce the number of crossings by two like in Figure 2.7b?



(a) Type I Reidemeister move. (b) Type II Reidemeister move.

Figure 2.7: Recognizing extra crossings from the Dowker notation.

- A Type I Reidemeister move can be recognized from the Dowker notation sequence by the pairing of two numbers that have an absolute difference of 1. In Figure 2.7a, a sample Dowker notation is 6 2 8 10 4. Since 2 pairs with 3, and $|2 - 3| = 1$, we know that there is a Type I Reidemeister move. Logically, this makes sense because when moving along the orientation, a trivial crossing will be crossed twice in a row.
- A Type II Reidemeister move can be recognized from the Dowker notation sequence by a specific characteristic of two pairs: If (x_1, y_1) and (x_2, y_2) is such a pair, where the order is listed such

that $|x_2| > |x_1|$ and $|y_1| > |y_2|$, then both of the following must be true.

$$||x_1| - |x_2|| = 1 = ||y_1| - |y_2|| \quad |x_1 - y_2| \neq ||x_1| - |y_2|| \quad (2.1)$$

In Figure 2.7b, a sample Dowker notation is 8 6 10 -2 12 4. Since $(-2,7)$ and $(3,6)$ are pairs that satisfy the above, we know that there is a Type II Reidemeister move. Logically, this makes sense because, when moving along the orientation, such crossings' numbers will be sequential both times crossed (confirmed by the left statement in Equation 2.1). Additionally, the even numbers must have opposite signs or the strand would be linked, as opposed to entirely above or below (confirmed by the right statement in Equation 2.1).

- The most consequential fact of Dowker notation is that it allows us to “feed projections of knots into the computer simply as a sequence of numbers” (40).

2.3 Conway’s Notation

- Particularly suited to calculations involving **tangles**.
- **Tangle:** “A region [in a knot or link] in the projection plane surrounded by a circle such that the knot or link crosses the circle exactly four times” (41).
 - These four crossings will be thought of as occurring in the four ordinal compass directions.
- Tangles can be thought of as building blocks in knot or link projections.
- **Equivalent:** Two tangle projections that can be transformed into each other via a series of Reidemeister moves “while the four endpoints of the strings in the tangle remain fixed and while the strings of the tangle never journey outside the circle defining the tangle” (41).
- Notice that, if a knot is formed by gluing together the tangle’s ends in pairs, the knot is equivalent to other projections of itself as long as the tangles are equivalent (by a series of Reidemeister moves).
- Tangles, such as the one in Figure 2.8c are denoted by the number of left-hand twists (crossings).
 - If we had twisted the tangle the other way, we would have called it -3 .
 - More simply, “for a positive-integer twist, the overstrand always has a positive slope, if we think of it as a small segment of a line [perhaps a cubic, in the case of Figure 2.8c]” (42).

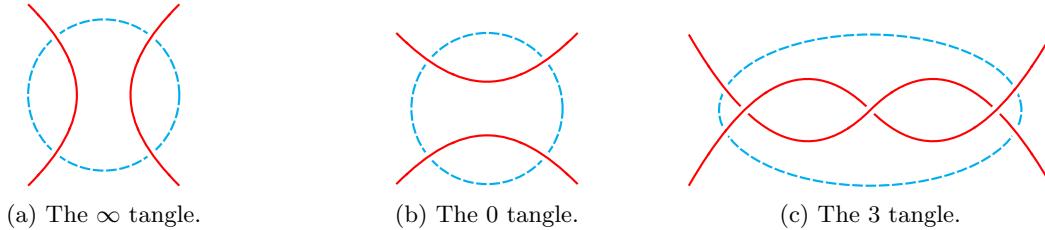


Figure 2.8: Tangles.

- **0 tangle:** The simplest of all tangles (by definition). *Also known as trivial tangle.* See Figure 2.8b.
- If we reflect the 3 tangle across a line connecting the NW and SE corners, and then proceed to extrude the NE and SE strands, twisting positively twice, we yield the tangle, 3 2.
- If we reflect 3 2 across the same line, and then proceed to extrude the NE and SE stands, twisting negatively four times, we yield the tangle, 3 2 -4 .
- **Rational tangle:** Any tangle that can be constructed in the manner described in the above 2 bullet points.

- Note that if the tangle is represented by an even number of integers, it can be thought of as being constructed from the ∞ tangle, and vice versa for the 0 tangle.
- *Exercise 2.10:* Draw the rational tangles corresponding to $2 -3 4 5$ and $3 -1 3 -3 2$.

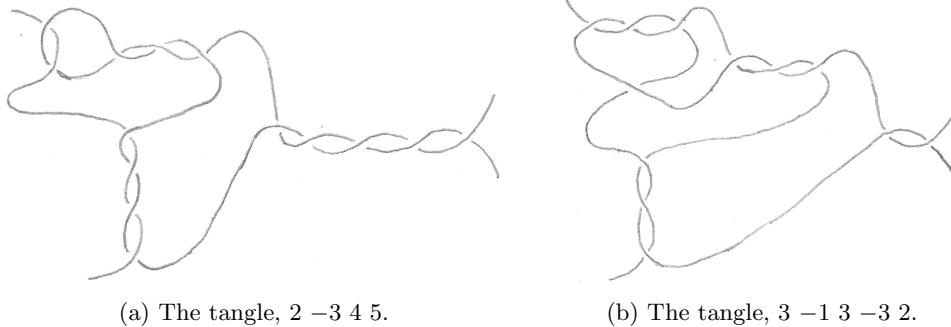


Figure 2.9: Solution to *Exercise 2.10*

- There is an extremely simple way to tell if two rational tangles are equivalent from their notation: **continued fractions**.
- **Continued fraction:** A fraction obtained through the iterative process of summing a number and its reciprocal.

– For example, given the three numbers a , b , and c , the continued fraction would be the following.

$$c + \frac{1}{b + \frac{1}{a}} \quad (2.2)$$

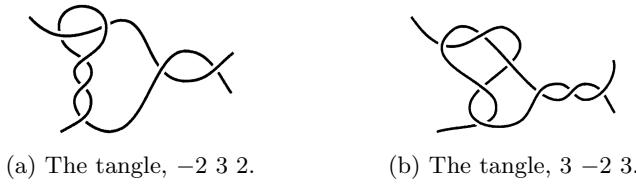


Figure 2.10: Two equivalent tangles.

- Applied to knot theory, we know that the tangles, $-2 3 2$ and $3 -2 3$, (seen in Figure 2.10) are equivalent because of the following.

$$\begin{aligned} 2 + \frac{1}{3 + \frac{1}{-2}} &= 3 + \frac{1}{-2 + \frac{1}{3}} \\ 2 + \frac{1}{\frac{5}{2}} &= 3 + \frac{1}{\frac{-5}{3}} \\ 2 + \frac{2}{5} &= 3 + \frac{-3}{5} \\ \frac{12}{5} &= \frac{12}{5} \end{aligned} \quad (2.3)$$

- *Exercise 2.13:* Determine which of the four rational tangles in Figure 2.11 are equivalent^[4].

⁴Idea for first paper: Prove that a series of NE↔SE and SE↔SW moves is sufficient to generate any rational tangle (like Reidemeister moves). Apply these findings to defining a tangle from a projection. Potentially create a new notation (show how this notation and the current tangle notation convert).

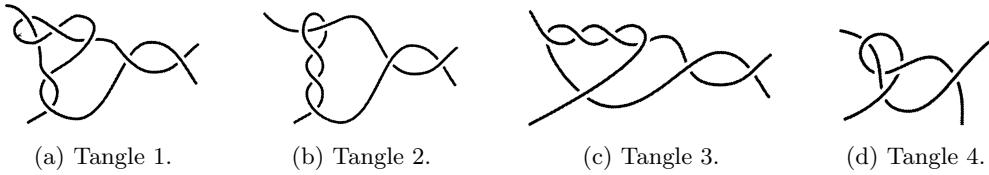


Figure 2.11: Four rational tangles.

– Tangle 1 (Figure 2.11a) can be denoted $-1 -1 -2 -2 -2$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 -2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-1 + \frac{1}{-1}}}}} &= -2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-2}}}} \\
 &= -2 + \frac{1}{-2 + \frac{-2}{5}} \\
 &= -2 + \frac{-5}{12} \\
 &= \frac{-29}{12}
 \end{aligned} \tag{2.4}$$

– Tangle 2 (Figure 2.11b) can be denoted $2 -3 2$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 2 + \frac{1}{-3 + \frac{1}{2}} &= 2 + \frac{-2}{5} \\
 &= \frac{8}{5}
 \end{aligned} \tag{2.5}$$

– Tangle 3 (Figure 2.11c) can be denoted $-4 1 2$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 2 + \frac{1}{1 + \frac{1}{-4}} &= 2 + \frac{4}{3} \\
 &= \frac{10}{3}
 \end{aligned} \tag{2.6}$$

– Tangle 4 (Figure 2.11d) can be denoted $1 1 1 1 1$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1}}}}} &= 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{2}}}} \\
 &= 1 + \frac{1}{1 + \frac{2}{3}} \\
 &= 1 + \frac{3}{5} \\
 &= \frac{8}{5}
 \end{aligned} \tag{2.7}$$

– Therefore, tangles 2 and 4 (Figures 2.11b and 2.11d, respectively) are equivalent by equations 2.5 and 2.7.

- If we close off the ends of a rational tangle, i.e. NE to NW and SE to SW, we form a rational link.

- If the link is a knot, we can denote it by its tangle. For example, the figure-eight knot is a rational knot with **Conway notation** 22 (because of the two twists twice in its center and the gluing of its ends).
- **Conway's notation:** A method of denoting knots by their tangles (all aforementioned theory in this section).
- Multiplying tangles:
 - Reflect the left tangle across its NW to SE diagonal and then glue its new NE and SE ends to the NW and SW ends, respectively, of the adjacent (to the right) tangle.
 - Multiplying a rational tangle by an integer tangle will always generate a rational tangle, e.g. 32 comes from multiplying 3 by 2.
 - To reflect a tangle over its NW-SE line, multiply it by the tangle, 0.

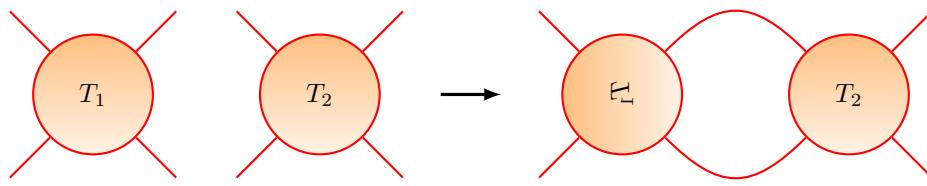


Figure 2.12: Multiplying tangles.

- Adding tangles:
 - Glue the left tangle's NE and SE ends to the right tangle's NW and SW ends, respectively.
 - Note that no reflecting occurs here.
 - “If we multiply each tangle in a sequence of tangles by 0, and then add them all together, we denote the resultant tangle by the sequence of numbers that stand for the original tangles, only now separated by commas” (48).
- **Pretzel knot:** “A knot obtained from a tangle represented by a finite number of integers separated by commas” (48).
- **Algebraic tangle:** “Any tangle obtained by the operations of addition and multiplication on rational tangles” (48).
- **Algebraic link:** “A link formed when we connect the NW string to the NE string and the SW string to the SE string on an algebraic tangle” (48). *Also known as aborescent link.*
 - Denoted the same way as a tangle.
 - Fun fact: The trefoil knot (Figure 1.1b) is an algebraic link of the 3 tangle (Figure 2.8c).
- **Additive identity:** A quantity that can be added to a second quantity without changing that second quantity.
 - 0 is an additive identity for the real numbers.
- *Exercise 2.21:* Is there an additive identity for tangles?
 - The 0 tangle (Figure 2.8b) is an additive identity for tangles. Adding the 0 tangle does the same thing as extruding (performing a planar isotopy on) the NE and SE strands. Such an addition does not change the relative position of any strand's ordinal end(s).
 - Continuing the analogy to the real numbers, it makes sense that the 0 tangle would be analogous to 0, the real number.

- **Multiplicative identity:** A quantity that can be multiplied by a second quantity without changing that second quantity.
 - 1 is a multiplicative identity for the real numbers.
- *Exercise 2.22:* Is there a multiplicative identity for tangles? Is it the same if you multiply a tangle by it on the right side or the left side?
 - The 0 tangle (Figure 2.8b) is a multiplicative identity for tangles. Multiplying the 0 tangle does the same thing as extruding (performing a planar isotopy on) the SE and SW strands. Such a multiplication does not change the relative position of any strand's ordinal end(s).
 - To part 2, yes.
- The analogy to real numbers does not hold in some areas.
 - Multiplication on tangles is not commutative — $ab \neq ba$ for tangles.
 - Multiplication on tangles is not associative — $(ab)c \neq a(bc)$ for tangles.
 - There are no inverse tangles — there is no tangle that when added to T gives back the trivial tangle, 0.
- There are tangles that are not algebraic.
- New knots can be obtained via **mutation**.
- **Mutation:** Severing the connections between a tangle and any adjacent tangles, transforming it (flipping horizontally, vertically, or both), and gluing the strands to their new adjacent counterparts.
 - Knots formed in this manner are called **mutants** of one another.
- *Exercise 2.25:* Show that if we have three tangles as in Figure 2.13a, we can mutate several times in order to permute the tangles. Note that we can then permute n tangles in a row.

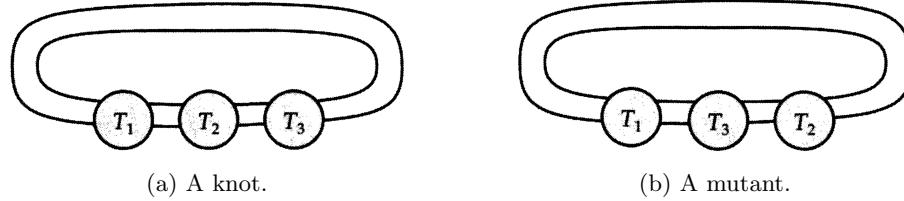


Figure 2.13: Series of mutations.

- Beginning in Figure 2.13a, horizontally flip T_2 .
- Horizontally flip T_3 .
- Horizontally flip the unit of the flipped T_2 and T_3 .
- Mutation cannot turn a nontrivial knot into the trivial knot (Figure 1.1a).
- More on mutant knots can be found in Sections 6 and 7.

2.4 Knots and Planar Graphs

- This section's notation bridges knot theory and graph theory.
 - It is capable of contributing to both branches of mathematics.
- **Graph:** A set of points called vertices that are connected by edges.

- **Planar graph:** A graph that lies in a plane.
- Creating a planar graph from a projection of a knot or link:
 - “Shade every other region of the link projection so that the infinite outermost region is not shaded” (51).
 - “Put a vertex at the center of each shaded region and then connect with an edge any two vertices that are in regions that share a crossing” (52).
 - Assign an orientation.
 - “Label each edge in the planar graph with a + or a –, depending on whether the edge passes through a + crossing or a – crossing” (52). See Figure 2.14. Recall computing link numbers in Section 1.4 (see Figure 1.18, specifically).
 - The result is a **signed planar graph**. See Figure 2.15 for an original example.

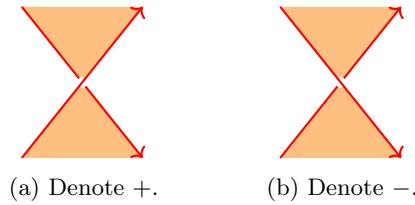


Figure 2.14: Computing signs per crossing.

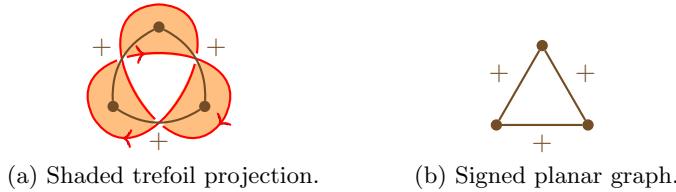


Figure 2.15: A signed planar graph from the standard trefoil knot projection.

- It is also possible return the signed planar graph to a knot projection:
 - “Put an x across each edge” (53).
 - “Connect the edges inside each region of the graph” (53).
 - “Shade those areas that contain a vertex” (53).
 - “At each of the x ’s, put in a crossing corresponding to whether the edge is a + edge or a – edge” (53).
- The significance here, as previously referenced, is that we can turn questions about knots into questions about graphs.
- *Exercise 2.31:* What do the Reidemeister moves become when translated into operations on signed planar graphs? (Make sure you consider what happens when different regions are shaded.)
 - Type I Reidemeister move: Extend an edge from one vertex to a newly created one. Either sign is acceptable (the sign option is analogous to the over-under crossing option).
 - Type II Reidemeister move: Add two edges connecting two, specific vertices and intersecting no other edges. The signs on the new edges must be opposite, but it either can be positive (or negative; this option is analogous to the over-under crossing option).
 - Type III Reidemeister move: I do not know^[5].
- See Section ?? for more on signed planar graphs.

⁵Return to this later.

3 Invariants of Knots

3.1 Unknotting Number

- **Unknotting number:** A value $n \in \mathbb{N}$ specific to a knot K that gives the fewest number of crossing changes needed in any projection to turn it into the unknot. *Also known as $\mathbf{u}(K)$.*
 - It may, however, be a nasty projection of the unknot (see Figure 1.12 and the related Exercise).
- *Exercise 3.1:* Find the unknotting number of the figure-eight knot (Figure 1.4).
 - $u(K) = 1$, where K is the figure-eight knot.
 - K is distinct from the unknot, so $u(K) \neq 0$. The next possibility ($u(K) = 1$) succeeds with the projection in Figure 1.4 (in fact, flipping *any* crossing in the referenced projection trivializes K).
- Note that performing all of the crossing changes on one projection is equivalent to performing one such change, then an ambient isotopy, then another...
 - A proof of this is listed on page 58.
- “The fact that every knot has a finite unknotting number follows from the fact that every projection of a knot can be changed into a projection of the unknot by changing some subset of the crossings in the projection” (58).
 - See *Exercise 1.7*.
- A proof of the above:
 - Pick a starting point (not a crossing) on an arbitrary knot and an orientation in whose direction to traverse the knot.
 - When you arrive at the first crossing, change it so that the strand that you are on is the over strand, if necessary.
 - Continue changing every new crossing to which you come to make the current strand into the overstrand.
 - Do not change crossings that you have already passed through once.
 - At this point, you now have a projection of a knot that was obtained from the original knot by changing crossings and that is the trivial knot. A proof of this triviality follows.
 - Consider the knot in three-space, with the z -axis coming out of the page. Assign to the starting point the z -coordinate, $z = 1$.
 - Decrease the z -coordinate of each consecutive point as you traverse the knot, culminating in the original point being labeled, $z = 0$.
 - Connect the starting and ending points with a vertical bar to preserve the knot as a knot.
 - What you have now obtained is a knot in three-space that, when projected from the top is the knot that we constructed before the line break, but has no crossings when projected from the side (and is therefore the trivial knot).
 - Q.E.D.
- It is very hard, in general, to find the unknotting number of a knot. How do you know that there isn’t a better projection?
- “A knot with unknotting number 1 is prime” (61).
- Therefore, a composite knot cannot be unknotted with a one crossing change.

- Another way to think about this idea is that changing a crossing to unknot one factor knot would yield the composition of the other factor knot with the unknot, which, according to Section 1.2, is simply the other factor knot (obviously still fully knotted).

$$u(K_1 \# K_2) \leq u(K_1) + u(K_2) \quad (3.1)$$

- The unknotting number can be realized by a projection that is not **minimal**.
 - **Minimal projection:** The projection of a knot with the lowest possible number of crossings.
 - **k -move:** A local change in the projection that replaces two untwisted strings with two strings that twist around each other with k crossings in a right-handed manner. See Figure 3.1.
- If k is positive, then the overstrand has a positive slope.

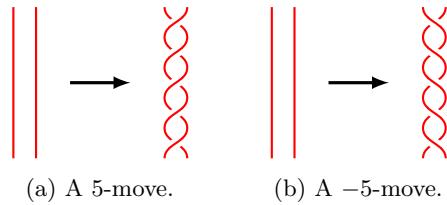


Figure 3.1: Examples of k -moves.

- **k -equivalent:** Two knot or link projections that can be transformed into each other through a series of k -moves and $-k$ -moves.
 - Ambient isotopies are allowed between k -moves.

- **Exercise 3.8:** Show that the knots in Figure 3.2 are each three-equivalent to a trivial link.

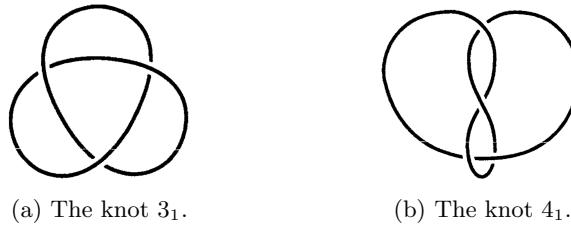


Figure 3.2: Knots that are 3-equivalent to trivial links.

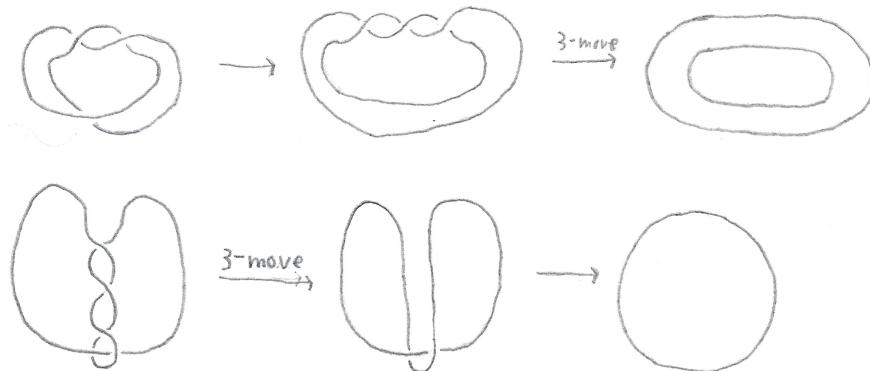


Figure 3.3: Solution to *Exercise 3.8*.

3.2 Bridge Number

- Think of knots as “cutting through the projection plane, rather than lying in it” (64).
 - Every overstrand is entirely above and every understrand is entirely beneath. The points where the strand shifts from over to under are known as **vertices**.
- **Bridge number:** The least number of unknotted arcs lying above the plane in any projection of the knot. *Also known as $b(K)$*
 - The number of maximal overpasses in in the projection.
- **Overpass:** “A subarc of the knot that goes over at least one crossing but never goes under a crossing” (64).
- **Maximal overpass:** “An overpass that could not be made any longer” (65).



Figure 3.4: Overpasses.

- *Exercise 3.10:* Show that if a knot has bridge number 1, it must be the unknot.
 - If a knot has bridge number 1, \exists a projection of the knot such that there is only one arc above the projection plane and 1 arc below the projection plane.
 - This projection, viewed from above, will look like any isotopy of the knot. However, viewed from the side, there will be 0 crossings — therefore, it is the unknot.
- *Exercise 3.11:* Show that the knot 5_2 has bridge number $2^{[6]}$.
 - Since 5_2 is distinct from the unknot, $b(K) \neq 1$.
 - Therefore, the next possibility to check is 2. If \exists a projection of 5_2 with bridge number 2, $b(K) = 2$.

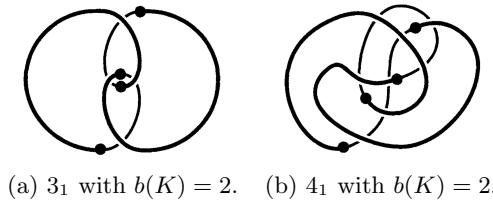


Figure 3.5: Reducing bridge number between projections.

- Before tackling this problem, though, let’s look at some simpler examples.
 - From these, we notice that there are “bridging moves” (distinct combinations of the Reidemeister moves) that can shift the bridge number by ± 1 .
 - Furthermore, long overstrands call to mind a Dowker reconstruction (Figure 2.5).
- Find the Dowker notation of 5_2 .
 - The Dowker notation of the projection listed in the appendix is 6 10 8 2 4.

⁶Idea for a paper: On reducing bridge number. Discuss manipulating a projection to its Dowker reconstruction to create the “killzone” and then using “bridging moves” to make the projection nonalternating. 2 types so far: I \rightarrow III, and 2xII \rightarrow III (in both cases it is the Type III move which climactically changes $b(K)$). It can be necessary to increase $b(K)$ before decreasing it.)

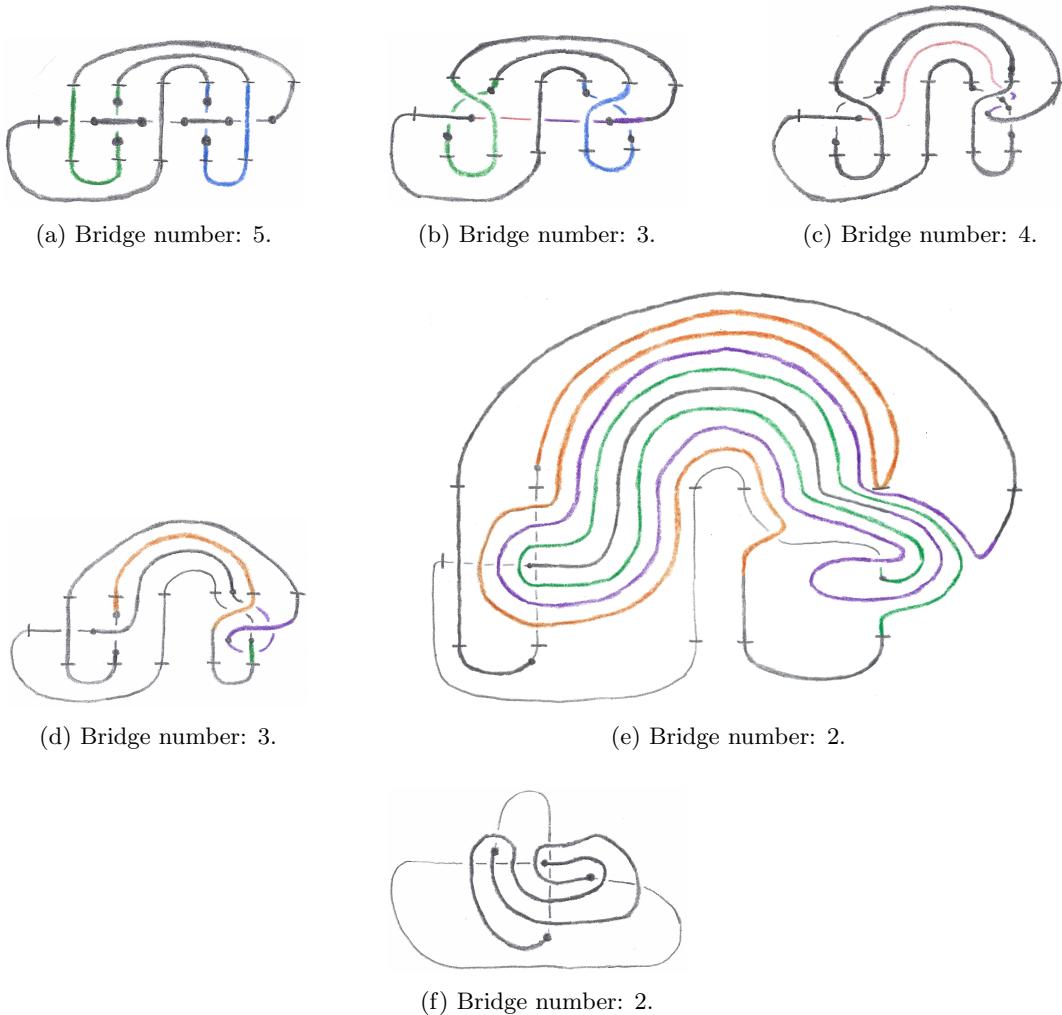


Figure 3.6: Solution to *Exercise 3.11*.

- Note that the following hold for all projections in Figure 3.6.
 - Dashes, when present, represent the boundaries of the “killzone.” They also provide consistent reference points for Figures 3.6a through 3.6e.
 - Dots, when present, represent bounds on the overpasses (and, consequently, the underpasses).
 - Thicker lines represent overpasses. Conversely, thinner lines represent underpasses.
 - Like-colored strands between sequential projections represent the same strand after Reidemeister moves or planar isotopies (colored for clarity, not for any specific purpose such as tricolorability).
- From the previously determined Dowker notation, reconstruct a knot projection. This projection (Figure 3.6a) contains the “killzone,” a concentrated region where every crossing is displayed in an orderly fashion^[7].
- Employ two Type I Reidemeister moves followed respectively by Type III Reidemeister moves to eliminate two overpasses. These moves occur on the green and blue strands, independently. They

⁷Note that from here-down, I do not guarantee that this is the fastest or simplest way to generate a projection that displays $b(K) = 2$, nor can I assure the reader that my solution is the “simplest” projection (perhaps the one with the fewest crossings). However, I do guarantee that it is correct. I know that it is so because all projections are separated only by Reidemeister moves and planar isotopies, and Figure 3.6f visibly has bridge number 2.

yield a projection with bridge number 3 (Figure 3.6b). This projection has several noteworthy properties.

- First, of the three overpasses, two contain 2 crossings while one contains 3 crossings.
- Second, of the three underpasses, two contain 2 crossings while one contains 3 crossings.
- Third, every overpass crosses at least one of its adjacent underpass.
- Because each overpass and underpass has more than one crossing, the bridge number of the projection must increase before it can decrease (one of the overpasses must be split into two [or three]; single overpasses can subsequently be eliminated via “bridging moves”).
- Because each overpass crosses at least one of its adjacent underpasses, no “bridging moves” can eliminate a bridge from this projection directly^[8].
- Employ two Type II Reidemeister moves followed by one Type III Reidemeister move to split one of the 2-crossing overpasses into two 1-crossing overpasses. These moves occur on the red-purple strand, which is “slid along” beneath the central overpass. This change will increase the bridge number to 4 and yield Figure 3.6c^[9].
- Employ a Type I Reidemeister move followed by a Type III Reidemeister move to eliminate one of the new 1-crossing overpasses. These moves occur on the purple strand. This change will decrease the bridge number back to 3. Additionally, perform the reverse of the actions on the green strand between Figures 3.6a and 3.6b to “straighten it out.” This shift will switch the 1-crossing overpass to an adjacent strand. These changes will yield Figure 3.6d. This projection has several noteworthy properties.
 - The 1-crossing overpass has underpasses of greater than 1 crossing on both sides.
 - The 1-crossing overpass passes through a loop, creating a setup similar to the two overpasses eliminated between Figure 3.6a and 3.6b. However, because of the above, a Type I Reidemeister move followed by a Type III Reidemeister move will not eliminate this overpass — it would only shift the position of the overpass.
 - The 1-crossing overpass crosses neither of its adjacent underpasses.
 - The above is key — because of it, a succession of 2xType II Reidemeister moves followed by 1xType III Reidemeister moves can combine the 1-crossing overpass with the partially purple overpass.
 - This set of moves will occur thrice — once for each crossing of the adjacent underpass.
- Three times sequentially, employ two Type II Reidemeister moves followed by one Type III Reidemeister move to (for the first two sets) shift the position of the 1-crossing overpass and then (for the last set) eliminate the 1-crossing overpass. These moves occur on the orange, purple, and green strands, sequentially. They yield a projection with bridge number 2 (Figure 3.6e).
- Q.E.D.
- Figure 3.6e is rather visually unpleasant, however. Figure 3.6f is the result of one Type I Reidemeister move (on the upper part of the orange strand) and planar isotopies intended purely to “clean up” the projection. It remains indistinct from 5_2 and displays $b(K) = 2$.

- **Two-bridge knots:** Knots that have $b(K) = 2$.
- If a two-bridge knot is sliced through the projection plane, there will be two **unknotted untangled** arcs both above and below said plane.
- **Unknotted untangled:** Distinct arcs that have no crossings with each other.
- From 2 such unknotted untangled arcs on each side of the projection plane, it is possible to construct all two-bridge knots.

⁸I am holding this to be true because I cannot demonstrate otherwise, but I do not know that it is true. I'd actually conjecture that it's false — it would seem to make sense that some overpass could be eliminated directly from Figure 3.6b.

⁹Multiple such moves may suffice here to split a 2-crossing bridge. I chose and followed through this one randomly.

- “The tricky part is that although the strings to each side of the plane are individually unknotted, they can twist around each other and themselves” (65).
- All two-bridge knots are rational knots and vice versa.
- Since the two-bridge knots are a very well understood class of knots, properties that are suspected to hold true for all knots are often first proved to hold true for two-bridge knots.
- All two-bridge knots are prime.
- The first tabulated three-bridge knot is 8_{10} .
- *Exercise 3.13:* Find a picture of 8_{10} that shows that it is at most a three-bridge knot.



Figure 3.7: The knot 8_{10} .

- Find the Dowker notation of 8_{10} ^[10].
 - The “maximum” Dowker notation (which, fortunately, is also “ideal”) for 8_{10} is 12 16 10 14
4 6 2 8
- From the previously determined Dowker notation, reconstruct a knot projection.
- There are four possible “type I bridging moves”^[11] allowed by this projection. Perform 3 of these to eliminate 3 bridges.
 - For the “linked loops,” it does matter on which loop the move is performed.
 - This yields a projection with 2x3-crossing bridges, 2x2-crossing bridges, and 1x1-crossing bridge.
- Perform two “type II bridging moves”^[12] to eliminate the 1-crossing bridge.
- Perform four “type II bridging moves” to eliminate a 2-crossing bridge.
- Q.E.D. — see Figure 3.8 for the final projection^[13].
- “In 1954, a German mathematician named H. Shubert proved that $b(K_1 \# K_2) = b(K_1) + b(K_2) - 1$ ” (67).
- *Exercise 3.14:* Explain how Shubert’s result implies that rational knots are all prime.
 - To be distinct from the unknot, a knot must have $b(K) \geq 2$, as proven in *Exercise 3.10*.
 - By Shubert’s theorem, above, and the above fact, any composition of two nontrivial knots must have $b(K) \geq 3$ (at least $2 + 2 - 1 = 3$).
 - Since any knot with $b(K) \geq 3$ is not rational and all composite knots have $b(K) \geq 3$, all rational knots (those with $b(K) = 2$) must be noncomposite, or prime.

¹⁰Note that different Dowker notations will lead to different Dowker reconstructions — some may be easier to manipulate than others. To concentrate the “killzone” in a single (or at least close) to rectangular region, choose Dowker notations with maximum “base strands.” Avoid, in the beginning, passing through at a crossing that will soon be passed through again.

¹¹A Type I Reidemeister move performed on an isolated loop followed by a Type III Reidemeister move to translate the new crossing across the “base strand.”

¹²Two Type II Reidemeister moves performed on the same strand to move it over either side of an adjacent crossing followed by a Type III Reidemeister move to translate the strand over the crossing followed by two Type II Reidemeister moves for simplification purposes.

¹³Note that this figure is drawn with the same rules used in Exercise 3.11, i.e. the original “killzone” is marked as well as the start and end of each bridge. Each “bridging move” has been carried out without simplification, so some thought should be able to reveal the connection between the original Dowker reconstruction (not pictured) and this projection.

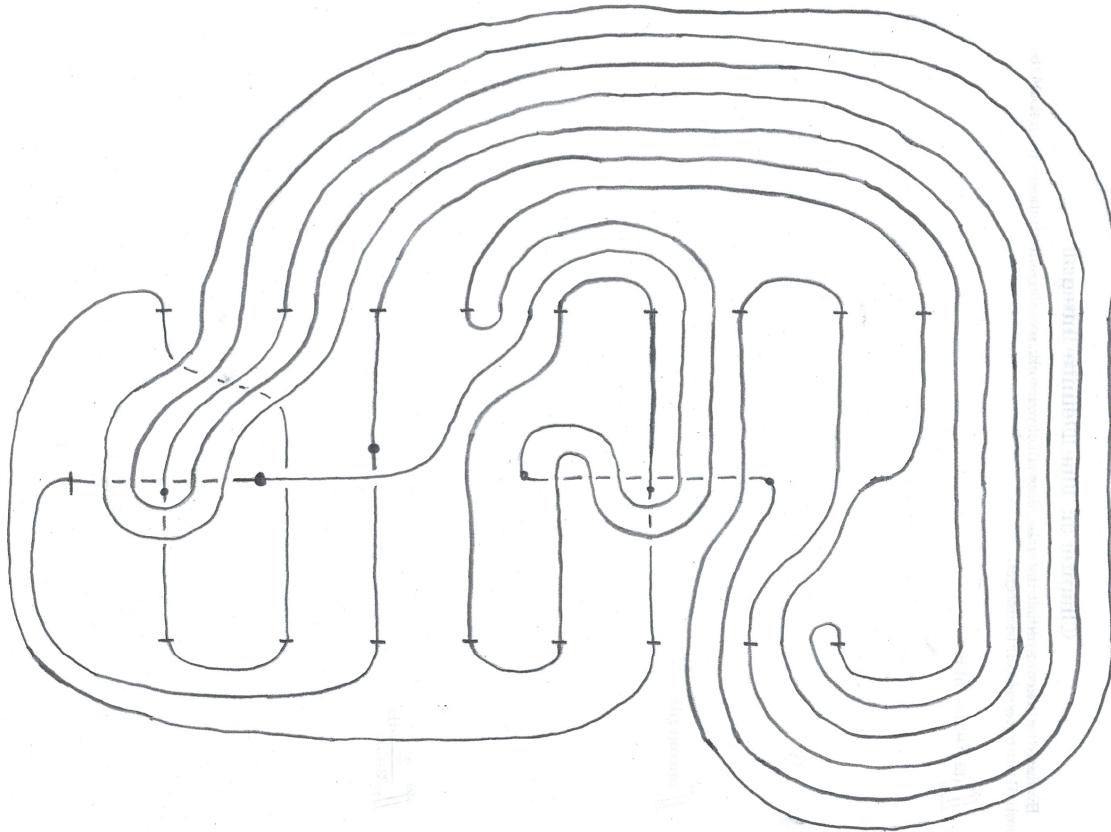


Figure 3.8: Solution to *Exercise 3.13*.

3.3 Crossing Number

- **Crossing number:** “The least number of crossings that occur in any projection of the knot” (67). *Also known as $c(K)$.*
- To find the crossing number of a knot, first find a projection of the knot. Count the crossings, n . At this point, it is clear that $c(K) \leq n$.
 - If all knots with fewer crossing than n are known and K is distinct from all of those knots, then $c(K) = n$.
 - Otherwise, reduce the knot to its minimal projection.
- Proving that a knot is distinct from every knot of fewer crossings is very difficult and requires knot polynomials (see Section 6).
- Determining $c(K)$ is generally very difficult (knots are currently only tabulated up to 13 crossings).
- **Reduced:** A projection of a knot with no easily removed crossings.
 - All “easily removed” crossings in alternating knots are shown in Figure 3.9.
- However, in 1986, Lou Kauffman, Kunio Murasugi, and Morwen Thistlethwaite all independently proved that “an alternating knot in a reduced alternating projection of n crossings has crossing number n ” (68).
 - They utilized the Jones polynomial (see Section 6).

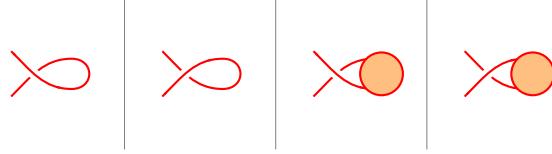


Figure 3.9: Easily removed crossings in an alternating knot projection.

- Because we can visually tell whether or not a projection is reduced (and reduce it if it's not), we can determine $c(K)$ for any alternating knot.
- It is still unknown whether or not $c(K_1 \# K_2) = c(K_1) + c(K_2)$.
 - However, Kaufman, Murasugi, and Thistlethwaite's result proves that the above holds when $K_1 \# K_2$ is alternating.
 - *Exercise 3.15:* Show that if K_1 and K_2 are alternating, then so is $K_1 \# K_2$. Hence, $c(K_1 \# K_2) = c(K_1) + c(K_2)$ holds when K_1 and K_2 are alternating even if $K_1 \# K_2$ does not appear alternating.

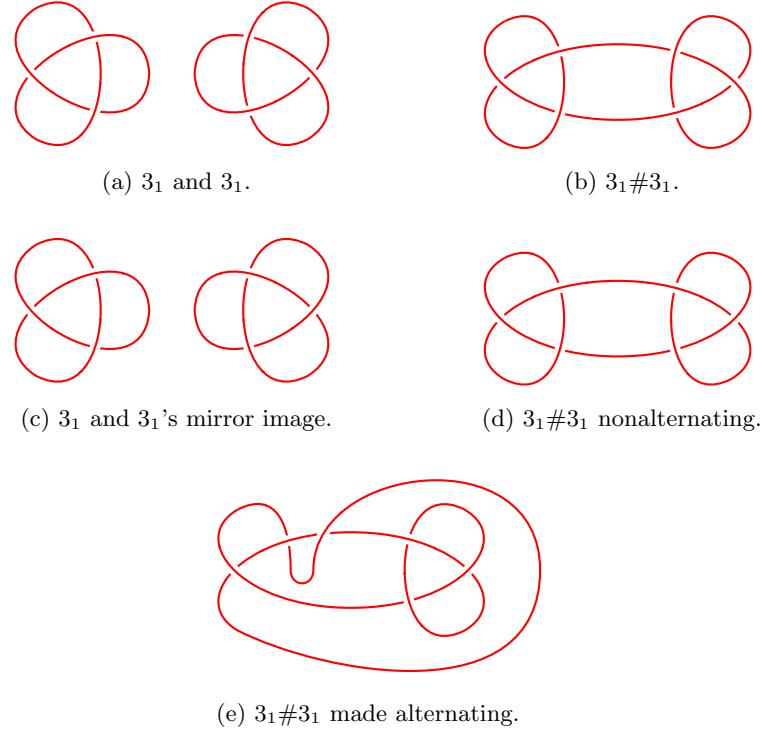


Figure 3.10: Drawing an alternating projection of $K_1 \# K_2$ where K_1 and K_2 are alternating.

- When composing alternating knots, there are two cases, as exemplified by Figure 3.10a and 3.10c, respectively.
 - Case 1: the composite knot will be in an alternating projection.
 - Case 2: the composite knot will not be in an alternating projection.
- For Case 1, the theorem evidently holds. See Figure 3.10b
- For Case 2, first consider the definition of an alternating knot (see Section 1.1), which states that an alternating knot need only have a projection that is alternating, i.e., every projection need not be alternating.

- Therefore, if we can prove that every nonalternating projection obtained through a composition has an alternating equivalent, the theorem will be true. To begin, consider several subcases to determine a pattern in the nonalternation.
 - Because K_1 and K_2 are both alternating, any arc being removed (to compose the two) will necessarily be on an understrand becoming an overstrand or vice versa.
 - Therefore, if we leave K_1 to enter K_2 from an understrand, we will return into K_1 to meet an overstrand and vice versa. The same holds for K_2 .
 - Case A: Leaving from an understrand in K_1 , enter K_2 on an understrand. The next crossing (continuing through K_2) will be an overstrand, and then an understrand, ... Eventually, by the above, we will return from K_2 on an overstrand into an overstrand on K_1 . The next crossing (continuing through K_1) will be an understrand, and then an overstrand, ... Eventually, we will return to where we started.
 - Case B: Leaving from an overstrand in K_1 , enter K_2 on an overstrand. The next crossing (continuing through K_2) will be an understrand, and then an overstrand, ... Eventually, by the above, we will return from K_2 on an understrand into an understrand on K_1 . The next crossing (continuing through K_1) will be an overstrand, and then an understrand, ... Eventually, we will return to where we started.
- In both Cases A and B, the only moments of nonalternation are two consecutive underpasses (or overpasses) followed by a factor knot followed by two consecutive overpasses (or underpasses). Notice this pattern in Figure 3.10d.
- Thus, “slide” one strand of a factor knot over the other knot to eliminate the nonalternation, as in Figure 3.10e.
- Q.E.D.

4 Surfaces and Knots

4.1 Surfaces without Boundary

- **Surface:** The two-dimensional exterior of a three-dimensional object with volume.
- **Torus:** “The surface of a donut” (71). *Plural is tori.*
- “At any point on a surface, there is a small region on the surface surrounding and containing the point that looks like a disk” (72).
 - The disk may be deformed (curved), but it must still be a disk.
 - Objects that are not surfaces have points with nondisk “neighborhoods” around them.
- Surfaces can also be called **two-manifolds**.
- **Two-manifold:** “Any object such that every point in that object has a neighborhood in the object that is a (possibly nonflat) disk” (73).
- *Exercise 4.1:* Based on the definition for two-manifolds, decide what the definition should be for a one-manifold. Find two different one manifolds.

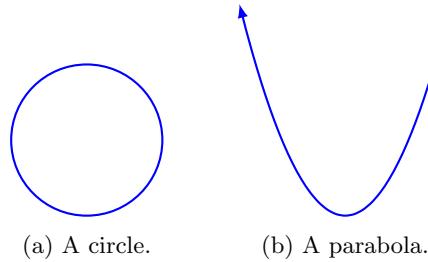


Figure 4.1: Two one-manifolds.

- Any object such that every point in that object has a neighborhood in the object that is a (possibly not straight) line.
- See Section ?? for info on three-manifolds.
- Think of surfaces as being made of rubber. In Figure 1.2, the deformation is an **isotropy** and the left and right surfaces are **isotropic surfaces**.
- **Isotropy:** A continuous deformation of a surface in space; a rubber deformation.
- **Isotropic surfaces:** Two surfaces in space that are equivalent under an isotropy.
- Surfaces that are not isotropic cannot be deformed into one another without some cutting and pasting!^[14]

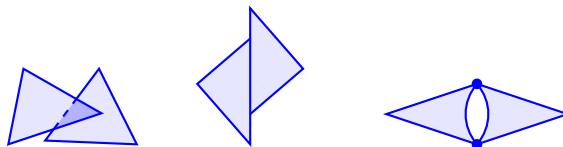


Figure 4.2: Triangles cannot intersect like this.

^[14]Identify surfaces by the number of holes in them?

- **Triangulation:** A division of a surface into deformable triangles.
 - Think of a triangle as a disk “with a boundary made up of three edges connecting three vertices” (75).
 - The triangles must “fit together nicely along their edges so that they cover the entire surface” (74). See Figure 4.2 for a visualization of ways that triangles cannot fit together.
 - Given a triangulation of a surface, the triangles can be separated and we can keep track of the original surface by labeling edges and putting arrows on them.

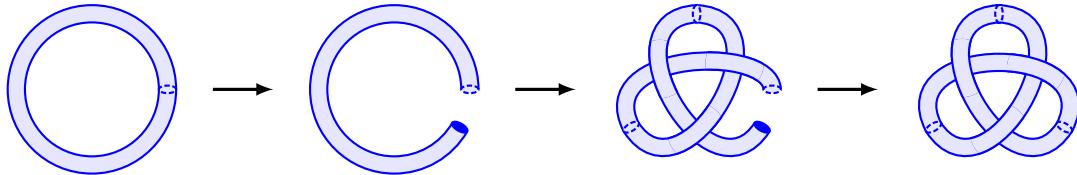


Figure 4.3: Homeomorphic surfaces.

- **Homeomorphic:** Two surfaces such that “one of them can be triangulated, then cut along a subset of the edges into pieces, and then glued back together along the edges according to the instructions given by the orientations and labels on the edges, in order to obtain the second surface” (76).
 - For example, the leftmost and rightmost surfaces in Figure 4.3 are both homeomorphic with a torus.
 - However, proving that the two are not isotropic would take some work.
 - It is not necessary to find a complete triangulation to prove that two surfaces are homeomorphic — it will suffice to circularly cut and put like circles back together, as in Figure 4.3.
 - One could also cut just edges and put like edges back together.
 - A sphere and a torus are clearly not homeomorphic because every closed loop on a sphere cuts it into two pieces. However, there exist closed loops on a torus that do not cut it into two pieces^[15].
 - The sphere, torus, two-hole torus, and three-hole torus are all not homeomorphic to one another.
 - Since we could keep increasing the number of holes, there are infinitely many distinct (nonhomeomorphic) surfaces.
- **Genus** (of a surface): “The number of holes in the doughnut” (78). *Plural is genera.*
 - The sphere has genus 0 while the torus has genus 1.
- **Embedding** (of a surface): “A choice of how to place a surface in space” (78).
 - The leftmost and rightmost surfaces of Figure 4.3 are two distinct embeddings of the torus in three-space.
- **Cube-with-holes:** The surface of a solid object obtained by drilling wormholes out of a cube.
- **Homeomorphism type:** The elementary surface to which a surface can be related.
 - Could be found via the triangulation method/cut-and-paste technique.
 - Could be found using its **Euler characteristic**.

¹⁵This is not an actual proof.

- **Euler characteristic:** Given a triangulation (or subdivision) of a surface, the sum $V - E + F$, where V is the number of vertices, E is the number of edges, and F is the number of triangles (or n -gon faces). Also known as χ .

$$\chi = V - E + F \quad (4.1)$$

- For example, χ for the leftmost sphere in Figure 4.4 is $6 - 12 + 8 = 2$.
- *Exercise 4.5:* Compute the Euler characteristic of the second triangulation of the sphere in Figure 4.4.

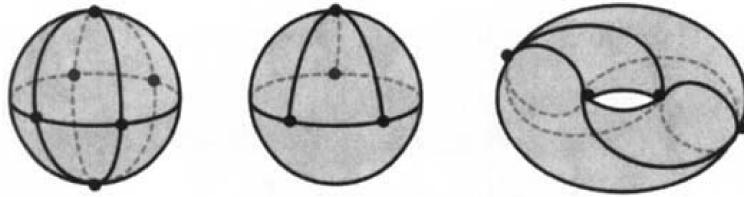


Figure 4.4: Triangulations of the sphere and the torus.

- $\chi = 4 - 6 + 4 = 2$
- *Exercise 4.6:* Compute the Euler characteristic of the triangulation of the torus in Figure 4.4.
 - $\chi = 4 - 12 + 8 = 0$
 - The sameness of the Euler characteristics of the two triangulations of the sphere exemplifies the fact that the Euler characteristic depends only on the surface, not the triangulation.
 - A summary of the proof that any two triangulations have the same Euler characteristic follows.
 - Take T_1 and T_2 , two different triangulations of a surface. We will overlay them and add appropriate edges (to make triangles) to make T_3 . This process will prove that T_1 and T_3 have the same Euler characteristic. Since T_2 could have T_1 overlaid, T_2 and T_3 also have the same Euler characteristic. Therefore, T_1 and T_2 have the same Euler characteristic (transitive property).
 - Each edge of T_2 intersects T_1 a finite number of times, and no vertices are overlaid — this can be assured by shifting where T_2 is overlaid by a slight bit.
 - Add vertices at the intersections of edges. The addition of one vertex and the subsequent splitting of one edge into two cancel out in the Euler characteristic.
 - Add the vertices of T_2 and a corresponding new edge to each. These edges should be a subset of the edges in T_2 . Vertices increase by 1 for each vertex and edges increase by 1 for each edge, so the change likewise cancels out.
 - Add the rest of the edges from T_2 . Note that for every edge added, a face is split in two, canceling out the change.
 - Add edges to split all n -gon faces into triangles. Each added edge splits a face, so the change, again, cancels out.
 - Q.E.D.
 - $\chi(\text{genus} = 0) = 2$.
 - $\chi(\text{genus} = 1) = 0$.
 - For a genus 2 surface, we could find a triangulation and calculate χ , or we could be clever.
 - **Connected sum:** Remove a disk from each of two objects and then glue them together.

- With triangulations of two tori, we can imagine that a genus 2 surface can be created via the connected sum of the two tori minus two matching faces.
- So, $F \rightarrow F - 2$, and because 3 vertices and 3 edges are merged, $V \rightarrow V - 3$ and $E \rightarrow E - 3$. The vertices and edges cancel out, giving the following property.

$$\chi(S) = \chi(S_1) + \chi(S_2) - 2 \quad (4.2)$$

- Therefore, $\chi(\text{genus} = 2) = -2$.
- Exercise 4.8:* Use connected sums to show that the Euler characteristic of a genus 3 surface is -4 .
 - $S_{\text{genus}=2} + S_{\text{genus}=1} = S_{\text{genus}=3}$.
 - $\chi(S_{\text{genus}=3}) = \chi(S_{\text{genus}=2}) + \chi(S_{\text{genus}=1}) - 2 = -2 + 0 - 2 = -4$.
- Exercise 4.9:* Use induction to show that the Euler characteristic of a surface of genus g is $2 - 2g$.
 - Basis step: $2 \stackrel{?}{=} 2 - 2(0)$.
 - Induction hypothesis: $\chi(S_{\text{genus}=g}) = 2 - 2g$
 - Induction step:

$$\begin{aligned} \chi(S_{\text{genus}=g+1}) &\stackrel{?}{=} 2 - 2(g + 1) \\ \chi(S_{\text{genus}=g}) + \chi(S_{\text{genus}=1}) - 2 &\stackrel{?}{=} 2 - 2g - 2 \\ (2 - 2g) + 0 - 2 &\stackrel{?}{=} 0 - 2g \\ 0 - 2g &\stackrel{?}{=} 0 - 2g \end{aligned}$$

- Therefore, the following equation holds true.

$$\chi(g) = 2 - 2g \quad (4.3)$$

- Every surface has a triangulation.
- Compact** (surface): A surface with a triangulation with a finite number of triangles.
- A plane and torus minus a disk are both noncompact surfaces.
- Surfaces appear in knot theory in the space around a knot.
- Complement** (of the knot): The space around the knot $M = \mathbb{R}^3 - K$, where \mathbb{R}^3 is the three-dimensional space around the knot K .
- A splittable link can be defined as a link such that there exists a sphere in the link complement between the two link components.
- Every knot is contained within a torus. Some knots can be surrounded by a torus in unusual ways or by a genus $2+$ surface.
- We are particularly interested in the surfaces that cannot be simplified.
- Compressible** (surface): For a link L , a surface F where there exists “a disk D in $\mathbb{R}^3 - L$ such that D intersects F exactly in its boundary and its boundary does not bound another disk on F . Note that D is not allowed to intersect L ” (86).
- Compression:** The simplifying operation of cutting F along D , gluing D to each open curve, and using an isotopy to “smooth out” the surface.

- *Exercise 4.11:* Show that a compression always increases the Euler characteristic. Use this to show that the genus of the resulting surface or surfaces is always less than the genus of the original surface.

- Think of D as a triangle. During the compression, three vertices and three edges separate into two pairs of three vertices and three edges. This does not change the Euler characteristic.
- However, two faces are also added. This increases the Euler characteristic by 2.
- This change of +2 changes the genus by -1 .

- **Incompressible** (surface): A surface that is not compressible.

- Any time we have a composite knot, there exists an incompressible **swallow-follow torus**.

- **Swallow-follow torus:** A torus that encompasses (swallows) one factor knot and wraps around the strand of (follows) the other factor knot.

4.2 Surfaces with Boundary

- **Boundary component:** Discs (that may have been deformed) that have been removed from surfaces.
- The Euler characteristic of a surface with boundary is that of that surface without boundary minus the number of boundary components.
 - Think of adding each boundary component as removing the face of a triangle.
- **Capping off** (a surface with boundary): “Filling in boundary components by attaching disks” (88).
- *Exercise 4.13:* Find the Euler characteristics of each of the surfaces in Figure 4.5 without triangulating them.

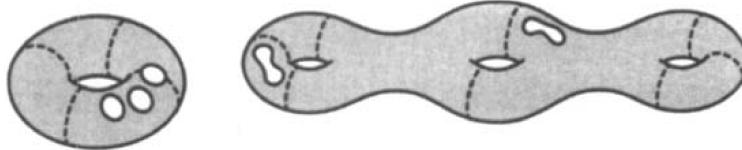


Figure 4.5: Surfaces with boundary.

- The left surface is a genus 1 surface with three boundary components. Therefore, it has $\chi = 0 - 3 = -3$.
- The right surface is a genus 3 surface with two boundary components. Therefore, it has $\chi = -4 - 2 = -6$.
- The Euler characteristic does not definitively distinguish surfaces with boundary.
- There are other ways, such as **orientability**.
- **Orientable** (surface): A surface sitting in \mathbb{R}^3 with “two sides that can be painted different colors, say black and white, so that the black paint never meets the white paint except along the boundary of the surface” (90).
 - For example, a disk and a torus are both orientable.
 - Furthermore, any torus with any number of boundary components is orientable.
- **Möbius band:** One of the simplest surfaces that is not orientable (see Figure 4.6).
- “A surface is nonorientable if and only if it contains a Möbius band within it” (91).

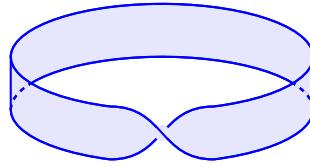


Figure 4.6: A Möbius band.

- A Möbius band must have an odd number of half-twists in it.
- *Exercise 4.15:* Decide which of the two surfaces in Figure 4.7 is orientable and which is nonorientable.

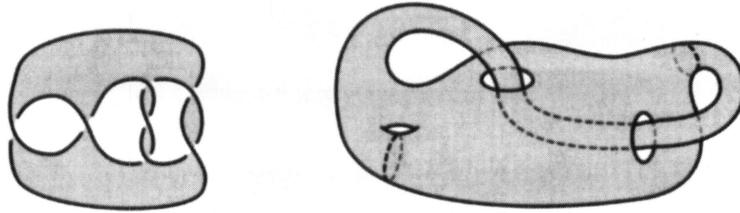
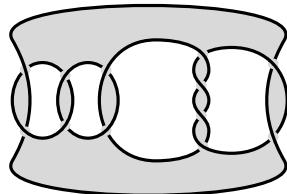
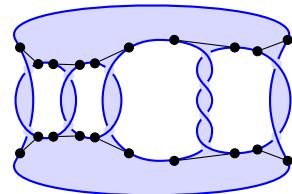


Figure 4.7: Possibly orientable surfaces.

- The surface on the left is nonorientable while the one on the right is orientable.
- Three facts completely determine the homeomorphic type of a surface with boundary.
 1. “Is it orientable or nonorientable?
 2. How many boundary components does it have?
 3. What is its Euler characteristic?” (92).
- My techniques:
 - Choose a subdivision that has one chord for every twist between boundary components.
 - Alternatively, for simpler counting at the expense of bigger numbers, choose two chords at every twist.
 - Count edges along each boundary component and then count the chords.
- *Exercise 4.16:* Use the three criteria to identify the surface with boundary in Figure 4.8a.



(a) A surface.



(b) A subdivision.

Figure 4.8: Identifying a surface with boundary.

- Begin by choosing the subdivision in Figure 4.8b according to the first technique above.
- Count $V = 20$, $E = 30$, and $F = 7$. Count $BC = 5$ where BC is the number of boundary components.

- Thus, $\chi = -3$. Adding +1 to χ for every boundary component, find χ for the original surface without boundary to be 2.
- Therefore, the surface without boundary corresponding to Figure 4.8a is the sphere.
- Define the genus of a surface with boundary to be that of the surface without boundary obtained by capping off each of its boundary components.
- Coming back to knot theory, one way to define knots is as the boundary of surfaces, i.e., the unknot could be defined as the boundary of a disk.
- **Annulus:** A sphere with two boundary components that lies outside one of the factor knots in a composite knot.
 - Note that this definition necessitates that the knot be thickened from a one-manifold.
 - Otherwise, an annulus would be a sphere with two punctures.
- Therefore, another definition of composite knot is “a knot such that there is a sphere in space punctured twice by the knot such that the knot is nontrivial both inside and outside the sphere” (93).
- **Conway sphere:** A sphere with four punctures enclosing a tangle.
 - This is the \mathbb{R}^3 analogy to the cyan circles in the projection plane in Figure 2.8.
 - If the knot is thickened, the Conway sphere becomes a sphere with four boundary components.

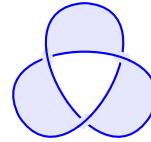


Figure 4.9: The trefoil as a Möbius band.

- Continuing with the idea that a disk defines the unknot, the trefoil knot (for instance) can be defined as the boundary of a Möbius band of three half-twists, as in Figure 4.9.
- Clearly, of particular interest are orientable surfaces with one boundary component such that said boundary component is a knot.

4.3 Genus and Seifert Surfaces

- There is one type of surface that appears in the complement of any knot.
 - In 1934, Herbert Seifert developed “an algorithm so that, given any knot, one can create an orientable surface with one boundary component such that the boundary circle is that knot” (95).
- Seifert’s algorithm is summarized as follows.
 - Take a knot and assign an orientation.
 - Eliminate each crossing by connecting the strand flowing in with the adjacent strand flowing out.
 - This produces a collection of **Seifert circles**.
 - Each circle bounds a disk in a plane. To avoid intersections, raise and lower the disks to different z -heights.
 - Finally, connect the disks to one another at the crossings of the knot by twisted bands.
- **Seifert circle:** A loop that can be deformed into a circle via a planar isotopy.

- *Exercise 4.18:* Show that the surface that we get doesn't depend on the direction we choose for the orientation on the knot.

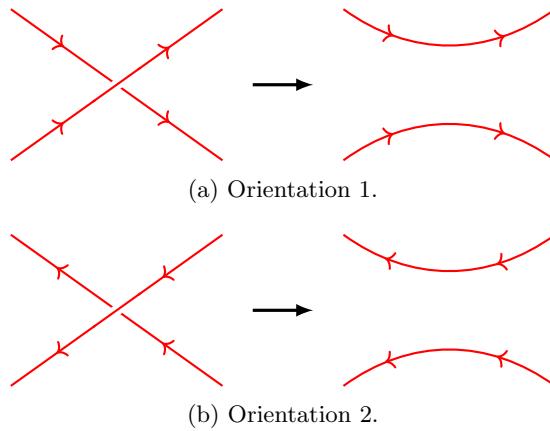


Figure 4.10: The relationship between orientation and crossing changes in Seifert's algorithm.

- At each crossing, two strands flow in and two flow out, regardless of the orientation, as in the left part of both subfigures in Figure 4.10. In effect, any crossing can be reworked to resemble the aforementioned parts of subfigures.
- One of the strands adjacent to a strand flowing in always flows out and vice versa.
- Because of this, reversing the orientation does not change the above fact (compare the left part of Figure 4.10a with that of Figure 4.10b).
- Thus, whatever adjacent strands were paired up before will be paired up again (see the right side of Figures 4.10a and 4.10b).
- Since the Seifert circles determine the final surface, the above proof that said circles will be unchanged verifies that the surfaces will remain unchanged.
- The surface is orientable because of the following thought experiment.
 - For any circle with a clockwise orientation, paint the upper face white and the lower face black.
 - For any circle with a counterclockwise orientation, paint the upper face black and the lower face white.
 - Any two circles above/below one another will have the same orientation, and, thus, the half-twist connects like colored faces.
 - Any two circles adjacent to each other will have opposite orientations, and, thus, the half-twist connects like colored faces.
 - Therefore, the surface is orientable.
- *Exercise 4.19:* Note that if we add an edge and two vertices to the surface across each crossing, cutting each band in half, we cut the surface up into valid faces. Use this fact to show that if c is the number of crossings and s is the number of Seifert circles, then $\chi = s - c$ and the genus of the surface is $g = \frac{c-s+1}{2}$.
 - Every 2-vertices-1-edge addition described above has a predictable effect on V and E .
 - Each vertex adds one new vertex (obviously), so two vertices are added for each 2-vertices-1-edge addition.
 - Each vertex splits an existing edge into two edges. Accounting for the addition of the connecting edge, three edges are added for each 2-vertices-1-edge addition.
 - Since two vertices and three edges are added each time, the quantity $V - E$ changes by -1 with each 2-vertices-1-edge addition.

- Because, by the problem statement, a 2-vertices-1-edge addition occurs at each crossing, there are c such additions.
- Thus, $V - E = -c$ for the entire process.
- The final subdivision, on the other hand, reveals something about F .
 - Each face is separated from the others at the crossings.
 - Thus, since each face is composed of a Seifert circle plus a bit of twisting connection, the number of faces corresponds to the the original number of Seifert circles s , or $F = s$.
- Therefore, the sum of the derived values can be substituted into the sum $V - E + F$ to yield $-c + s$ or $s - c$.
- Since Equation 4.1 gives us $\chi = V - E + F$, the following equation holds by the transitive property.

$$\chi = s - c \quad (4.4)$$

- According to Equation 4.3, $\chi = 2 - 2g$, so $g = \frac{2-\chi}{2}$.
 - However, the genus of a surface with boundary is that of the surface without boundary acquired via capping off said surface with boundary.
 - Since the surface at hand is known to have specifically one boundary component, the χ that should presently be considered is the χ of the previous part plus one, or $\chi_1 = s - c + 1$.
- Substitute: $g = \frac{2-(s-c+1)}{2}$.
- Simplify.

$$\begin{aligned} g &= \frac{2 - (s - c + 1)}{2} \\ &= \frac{2 - s + c - 1}{2} \\ g &= \frac{c - s + 1}{2} \end{aligned} \quad (4.5)$$

- *Exercise 4.21:* Show that Seifert's algorithm always generates at least two Seifert circles.

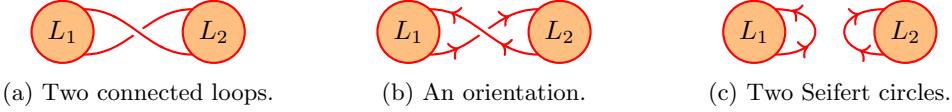


Figure 4.11: Generating Seifert circles from a crossing.

- In any knot with $c(K) > 0$, there exists a region with a crossing between two (possibly knotted) loops L_1 and L_2 , as in Figure 4.11a.
- When assigning an orientation to this crossing, the directions on the subarcs *must* be reversed.
 - If they both flowed from left to right (for example), that would imply that there is some point in L_1 from which the orientation “flows outward” in both directions and some point in L_2 to which the orientation “flows inward” from both directions.
 - This is contrary to the definition of an orientation, and, therefore, cannot be the case.
- Thus, the orientation in Figure 4.11b is valid.
- Given this orientation, the subarcs must separate into two Seifert circles, as in Figure 4.11c, according to Seifert's algorithm.
- This separation would hold for the opposite orientation, as demonstrated by Figure 4.10.
- It is possible to create multiple different-looking surfaces for the same knot by having a different starting projection.

- **Seifert surface** (for K): Given a knot K , “an orientable surface with one boundary component such that the boundary component of the surface is the knot K ” (99).
- **Genus** (of a knot): “The least genus of any Seifert surface for that knot” (99).
- The unknot has genus 0 (and is the only knot with genus 0).
- Seifert’s algorithm yields a Seifert surface of genus 1 for the figure-eight knot, and since the figure-eight knot is nontrivial, no Seifert surface of lower genus can exist for it. Therefore, the genus of the figure-eight knot is 1.
- *Exercise 4.22:* The twist knots are the knots obtained by linking together the two ends of a twisted loop. They include the trefoil and figure-eight knots. Show that all of the twist knots have genus 1.

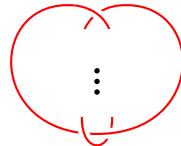


Figure 4.12: The general twist knot.

- This demonstration will follow the format of an induction proof. By Equation 4.5, if $g = 1$ for a knot, then $s = c - 1$. This argument will prove by induction that $s = c - 1$.
- Every twist knot can be projected in the mold of Figure 4.12, where the dots are replaced by some number of twists.
- For the trefoil knot, this number of twists is 0 and $s \checkmark c - 1$.
- Each new twist obviously adds one crossing.
- This new crossing, when the knot is assigned an orientation, will perfectly match one of the two left images in Figure 4.10.
 - This is because the orientation weaves its way down all of one of the two “vertical-ish” strands and then changes direction at the bottom loop to weave its way back up.
- Like in Figure 4.10, Seifert’s algorithm will yield one of the two products — separate loops.
- Thus, because every new crossing adds a new place to separate loops, both s and c increase by +1 in each subsequent twist knot.
- Therefore, since $s + 1 = c + 1 - 1$, s always equals $c - 1$.
- *Exercise 4.23:* Show that a minimal genus Seifert surface for a knot K must be incompressible (consider Euler characteristic).
 - Proof by contradiction: Take a compressible Seifert surface F for a knot K of minimal genus g_0 .
 - Compress it, adding +2 to the Euler characteristic χ_0 to get $\chi_1 = \chi_0 + 2$.
 - By Equation 4.3, the following holds.

$$\begin{aligned}\chi_0 &= 2 - 2g_0 \\ \chi_1 &= 2 - 2g_1 \\ \chi_0 + 2 &= 2 - 2g_1 \\ 2 - 2g_0 &= -2g_1 \\ g_1 &= g_0 - 1 \\ g_1 &< g_0\end{aligned}$$

- But since no genus g for K can be lower than g_0 by g_0 's definition as minimal, we have arrived at a contradiction.
- Therefore, the minimal genus Seifert surface must be incompressible.
- “Applying Seifert’s algorithm to an alternating projection of an alternating knot or link does yield a Seifert surface of minimal genus” (100).
- The following equation holds.
$$g(J \# K) = g(J) + g(K) \quad (4.6)$$
 - Clearly, $g(J \# K) \leq g(J) + g(K)$ because taking the composition of the two Seifert surfaces does not add any vertices, edges, or faces.
 - *Isotope surfaces, clean up intersections, perform surgeries to lower the number of intersection curves*^[16].
 - Arrive at $g(J) + g(K) \leq g(J \# K)$.
 - Therefore, $g(J \# K) = g(J) + g(K)$.
- Referencing Section 1.2, the trivial knot cannot be the composition of two nontrivial knots because $g(K) \geq 1$ where K is nontrivial, so by Equation 4.6, $g(J \# K) \geq 2 > 0 \Rightarrow g(J \# K)$ is nontrivial.
- *Exercise 4.27:* Show that the genus 1 knots are prime.
 - For a knot to be nontrivial, its genus must be greater than or equal to 1.
 - For a knot to be composite, its genus must be greater than or equal to 2.
 - Therefore, all knots with genus 1 have a high enough genus to be nontrivial (prime or composite) but not a low enough genus to not be composite (they are only prime).
- *Exercise 4.28:* Prove that if we take the composition of n copies of the same nontrivial knot, calling the result J_n , then as $n \rightarrow \infty$, $c(J_n) \rightarrow \infty$. Hint: Use the fact that there are only a finite number of knots with a given crossing number, but first use the Dowker notation to prove this fact.
 - For $c(K) = n$, the Dowker notation will consist of n even numbers ranging from 2 to $2n$. Therefore, there are n possibilities for the first number, $n - 1$ for the second number, and so on and so forth. This compresses down to $n!$ possible knots of $c(K) = n$.
 - Note that a large number of these knots will not be of this crossing number. This is an upper bound in the most literal sense.
 - One way to prove this would be to use Equation 4.6 to prove that the genus must increase with successive compositions. Since the genus is directly proportional to the quantity $c - s$ (by Equation 4.5), $c - s$ must increase with successive compositions. Since s cannot decrease with successive compositions (and certainly cannot decrease infinitely), c must increase with successive compositions.
 - I believe that the desired way to prove this is as follows. Composite knots are distinct from their factor knots. Generating a new composite knot creates a new knot. Since there are a finite number of knots of $c(K) = n$, eventually, all low-crossing possibilities (which will not all be reached, but regardless) must be used up. This forces $c(K)$ higher and higher as new knots are created.
- *Exercise 4.29:* Prove that if J_n is the composition of n copies of the same nontrivial knot, then $c(J_n) \geq n$. Hint: Use the Euler characteristic of a Seifert surface along with Equations 4.4 and 4.5 and *Exercise 4.21*. Subsequently, prove that $c(J_n) \geq 2n + 1$.
 - See the first solution to Exercise 4.28.

¹⁶Return to at a later date. Perhaps surgeries always increase Euler characteristic, lowering genus, until $g(J)$ and $g(K)$ are both at their minimum. Then their sum must be less than or equal to their composition?

- Seifert's algorithm's applicability to finding genus has been extended from alternating links to **alternative links**, which also include all **torus links** (see Section 5.1).
- There are some knots from which the minimal genus Seifert surface cannot be obtained using Seifert's algorithm on any projection of said knot.
- **Canonical genus** (of a knot): “The minimal genus of any Seifert surface obtained by applying Seifert’s algorithm to a projection of the knot” (106). *Also known as $g_c(K)$.*
- Genus was used to tell apart the Kinoshita-Terasaka mutants (see Section 2.3).

5 Types of Knots

5.1 Torus Knots

- **Torus knot:** A knot that lies on an unknotted torus such that the subarcs do not cross each other on the torus's surface.

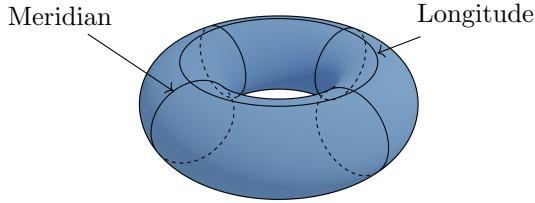


Figure 5.1: Longitude and meridian curves.

- **Meridian curve:** “A curve that runs once the short way around the torus” (108).
- **Longitude curve:** “A curve that runs once around the torus the long way” (108).
- Every torus knot is a (p, q) -torus knot for some pair of positive integers p and q .
 - p and q will always be **relatively prime**.
- **Relatively prime** (numbers): Two numbers whose greatest common divisor is 1.
- On the trefoil’s torus notation.
 - Since the trefoil knot crosses a longitude 3 times, it must wrap around the torus thrice meridionally.
 - Since the trefoil knot crosses a meridian 2 times, it must wrap around the torus twice longitudinally.
 - Therefore, the trefoil knot is a $(3, 2)$ -torus knot.
- *Exercise 5.1:* What torus knot is shown in Figure 5.2?

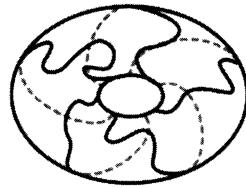


Figure 5.2: A mystery torus knot.

- Since it crosses a longitude curve five times and a meridian curve twice, it is the $(5, 2)$ -torus knot.
- Drawing a (p, q) -torus knot.
 - Begin by drawing two concentric ovals that represent the inner and outer bounds of the torus.
 - Draw p evenly spaced dots around the perimeter of both circles in corresponding locations.
 - Draw a dashed, curved^[17] line between each corresponding point. This makes it so that the upper lines can account for all q intersections with meridians.

^[17]Draw it curved to highlight the roundness of the torus.

- Draw a curved line between a point on the outer circle and a point that is a $\frac{q}{p}$ clockwise turn away on the inner circle.
 - Do this for all points.
- *Exercise 5.2:* Draw a $(4, 3)$ -torus knot.

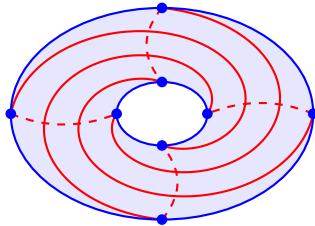


Figure 5.3: Solution to *Exercise 5.2*.

- *Exercise 5.3:* What would happen if we tried to draw a torus knot where p and q were not relatively prime? Say a $(3, 6)$ -torus knot?
 - If p and q are not relatively prime, then there is some natural number n that, multiplied by either p or q will yield the other.
 - Thus, such a torus knot can be denoted as a (p, np) -torus knot or as an (nq, q) -torus knot where $n \in \mathbb{N}$.
 - Let's take this case by case for each of the two above cases.
 - Case 1: What knot is a (p, np) -torus knot?
 - By the previously given method of drawing a knot, each point on the outer ellipse is connected with a point on the inner ellipse via a $\frac{q}{p}$ clockwise turn.
 - If $np = q$, then such points are connected via $\frac{np}{p} = n$ clockwise turns.
 - Thus, the point connected to via the line on the underside of the torus will be the original point.
 - Because the strand has only completed one loop the short way around the torus, it is a meridian, and, thus, the unknot.
 - Since this process is repeated p times, a (p, np) -torus knot is p unknots.
 - Case 2: What knot is an (nq, q) -torus knot?
 - By the previously given method of drawing a knot, each point on the outer ellipse is connected with a point on the inner ellipse via a $\frac{q}{p}$ clockwise turn.
 - If $nq = p$, then such points are connected via a $\frac{q}{nq} = \frac{1}{n}$ clockwise turn.
 - To get back to either the original point will require 1 full turn, or n turns of value $\frac{1}{n}$ (because $n \times \frac{1}{n} = 1$).
 - This process can be repeated q times to yield q such loops.
 - Since these loops are the same distance apart everywhere on the torus (parallel) and have no intersections, they are a set of deformed, radially stacked unknots.
 - Therefore, an (nq, q) -torus is q unknots.
 - Therefore, a torus knot where p and q are not relatively prime yields a number of unknots equal to whichever value (p or q) is smaller.
 - A $(3, 6)$ -torus knot would be 3 unknots.
- *Exercise 5.4:* Show that a (p, q) -torus knot always has a projection with $p(q - 1)$ crossings.
 - There is a potential underpass every $\frac{1}{p}$ turn.

- In going through a $\frac{q}{p}$ turn, a strand begins at one of these potential underpasses and touches $\frac{q/p}{1/p} = q$ more underpasses.
- Therefore, there are $q + 1$ possible crossings.
- However, since the beginning and end “crossings” are just continuations of a line, they are not crossings.
- Thus, each strand generates $q + 1 - 2 = q - 1$ crossings.
- Therefore, for p strands, the number of crossings is the product $p(q - 1)$.
- “Every (p, q) -torus knot is also a (q, p) -torus knot” (110).
 - Remove a disk from the torus that does not intersect the knot.
 - Transform it into two attached bands.
 - Turn each band inside out.
 - Transform it back into a torus.
 - Replace the disk.
- In conjunction with *Exercise 5.4*, this means that a (p, q) -torus knot has a projection with $q(p - 1)$ crossings.
 - Therefore, $c(K) \leq \min(p(q - 1), q(p - 1))$.
 - In fact, Murasugi (1991) recently proved that $c(K) = \min(p(q - 1), q(p - 1))$.
- **Solid torus:** “A doughnut where we include both the interior of the doughnut as well as the surface” (112).
- **Core curve** (of a solid torus): “The trivial knot that runs once around the center of the doughnut” (112).
- **Meridional disk** (of a solid torus): “A disk that has a meridian curve as its boundary” (112).
- *Exercise 5.5:* Let K be a (p, q) -torus knot sitting on a torus that is the boundary of a solid torus. Let J be the core curve at the center of the solid torus. Determine the linking number of J and K . What if K sits on the torus as a (q, p) -torus knot?

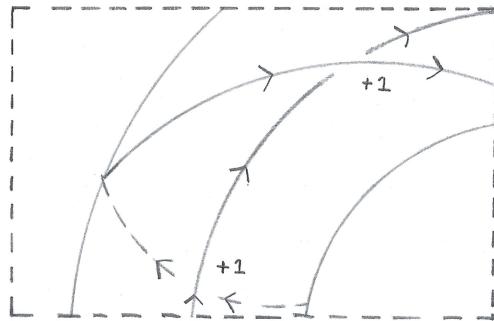


Figure 5.4: p such regions exist in a solid torus.

- For a (p, q) -torus knot, there are p spirals inward and p loops outward.
- The spirals always pass over the core curve while the loops always pass under the core curve.
- For each p , there exists a region similar to the one in Figure 5.4. Each of these regions contributes +1 to the linking number as drawn (see Figure 1.18).
- Therefore, the linking number is p .

- If K sits on the torus as a (q, p) -torus knot, then the linking number is q by the same logic.
- $u((p, q)\text{-torus knot}) = \frac{(p-1)(q-1)}{2}$
- *Exercise 5.7:* Show that in the specific case that $(p, q) = (3, 4)$, the (p, q) -torus knot can be unknotted with $\frac{(p-1)(q-1)}{2}$ crossing changes.

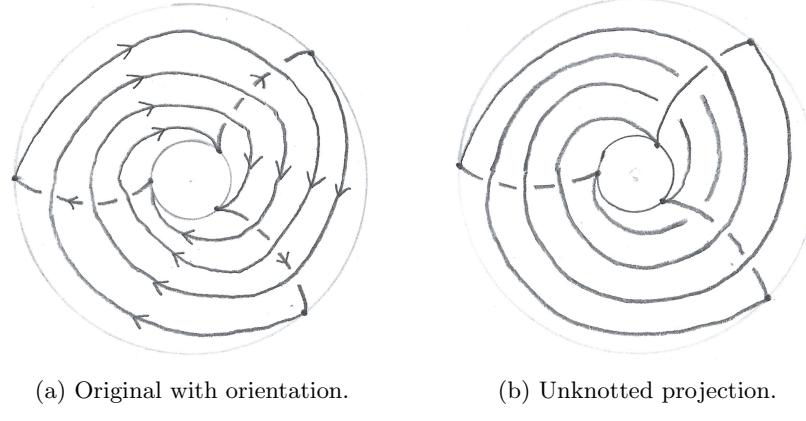


Figure 5.5: Unknotting the $(3, 4)$ -torus knot.

- Assign the orientation shown in Figure 5.5a and start at the leftmost point on the outer edge of the torus.
- Use the method conjectured in *Exercise 1.7* and described in detail in Section 3.1.
- This method will end up changing only $\frac{(3-1)(4-1)}{2} = 3$ crossings and producing a projection of the unknot (Figure 5.5b). Note the difference between Figures 5.5a and 5.5b to see which crossings were changed.
- “Seifert’s algorithm applied to a projection [like Figure 5.3] will yield a minimal genus Seifert surface” (113).
 - Remember that both alternating knots and torus knots qualify as alternative knots, as mentioned in Section 4.3.
- *Exercise 5.8 (a):* Applying the above result, use a standard projection to determine the genus of a (p, q) -torus knot.
 - Apply Seifert’s algorithm to generate s Seifert circles.
 - Assign a clockwise orientation.
 - At every crossing, the entering strand spiraling inward will merge with the leaving strand looping outward.
 - Likewise, the entering strand looping outward will merge with the leaving strand spiraling inward.
 - As this process continues, eventually, there will be a distinct Seifert circle for each of q crossings of a meridian.
 - Therefore, $s = q$.
 - Count the number of crossings c .
 - According to *Exercise 5.4*, the standard projection of a (p, q) -torus knot has $p(q-1)$ crossings.
 - Therefore, $c = p(q-1)$.
 - According to Equation 4.5, $g = \frac{c-s+1}{2}$.

- Therefore, the following holds.

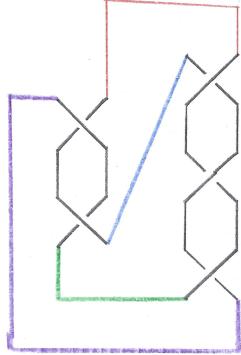
$$\begin{aligned}
g &= \frac{p(q-1) - q + 1}{2} \\
&= \frac{pq - p - q + 1}{2} \\
g &= \frac{(p-1)(q-1)}{2}
\end{aligned} \tag{5.1}$$

- *Exercise 5.8 (b):* Does it matter if the projection comes from the knot represented as a (q,p) -torus knot rather than a (p,q) -torus knot?

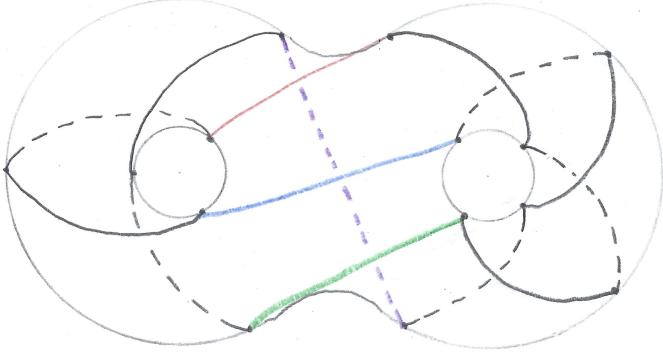
- Substitute $p = q$ and $q = p$ into Equation 5.1 and employ the commutative property of multiplication to prove that the first equation below is equal to Equation 5.1.

$$\begin{aligned}
g &= \frac{(q-1)(p-1)}{2} \\
&= \frac{(p-1)(q-1)}{2}
\end{aligned}$$

- $s((p,p-1)\text{-torus knot}) = 2p$.
 - $s(K)$ is the stick number of a knot K (see Section 1.6).
 - Used the idea of curvature to obtain a lower bound for $p \geq 3$.
- $s((p,q)\text{-torus knot}) = 2p$ if $p > q > \frac{p}{2}$.
 - Used a variation on bridge number called superbridge number to obtain a lower bound.
- **n -braid:** A (p,n) -torus knot.
- **Torus knot** [formally]: “A nontrivial knot that can be placed on the surface of a **standardly embedded** torus without crossing over or under itself on the surface” (114).
- **Standardly embedded** (surface): A surface that is unknotted in space.
- **Two-embeddable knot:** A knot “that cannot be placed on a standardly embedded torus but that can be placed on a standardly embedded genus two surface” (114).
 - The figure-eight knot is a two-embeddable knot.
- **n -embeddable knot:** A knot that “can be placed on a genus n standardly embedded surface without any crossings, but [that] cannot be placed on any standardly embedded surface of lower genus without crossings” (114).
- *Exercise 5.9:* Determine the minimal genus standardly embedded surface that the 5_2 knot can be embedded on, given that it is not a torus knot.
 - This is far from an ideal method, but it does successfully embed the 5_2 knot on a genus 2 surface.
 - Begin with the Conway notation of the 5_2 knot, namely 32.
 - This tells us that the knot can be drawn as the product of two elementary rational tangles (see Figure 5.6a).
 - Draw a standardly embedded genus 2 surface.
 - Draw the 2 tangle on its left handle and the 3 tangle on its right handle. Note that the choice of left versus right is arbitrary.
 - Based on the red, blue, and green connections in Figure 5.6a, draw the corresponding connections in Figure 5.6b.



(a) K_2 as tangles.



(b) K_2 as a two-embeddable knot.

Figure 5.6: Solution to *Exercise 5.9*.

- Since the purple connection cannot leave the torus like in Figure 5.6a, have it pass along the underside of the torus, directly, as in Figure 5.6b. This will add 3 new crossings, but it will embed K_2 on a genus 2 surface.
- A knot with bridge number b is an n -embeddable knot where $n \leq b$.
- Any knot is an n -embeddable knot for some n .
- If K is an n -embeddable knot, then the genus of K is at most $n - 1$ ^[18].

5.2 Satellite Knots

- **Satellite knot:** The knot K_3 produced by taking a knot K_1 , which resides within a solid torus, and then knotting the torus into a distinct knot K_2 .
 - K_1 must lie within the torus such that it hits every meridional disk and cannot be isotroped to miss any of the disks.
- **Companion knot:** The knot K_2 in the above description.
 - K_2 must be a nontrivial knot.
 - The companion knot is the core curve about which K_3 orbits. This is why K_3 is a *satellite* knot.
- For every satellite knot, there exists a knotted torus in the complement of the knot.
 - This knotted torus is always incompressible.
- **Whitehead double** (of the companion knot): A satellite knot where K_1 is the unknot twisted up as in Figure 5.7.
 - So named because K_1 resembles the Whitehead link (see Figure 1.13a).
 - For example, if we take K_1 as the unknot in Figure 5.7 and then knot the solid torus into a trefoil, we will have obtained a Whitehead double of the trefoil knot.
 - If we perform a k -move where k is even^[19] on the unknot in Figure 5.7 and then knot the solid torus into a trefoil, we will have obtained another Whitehead double of the trefoil knot.
 - Both of these Whitehead doubles have the property that if \mathbb{R}^3 is cut along the boundary of the knotted torus, the two pieces will be the solid torus at hand and three-space minus a solid torus knotted as K_2 .

¹⁸These last 3 statements all exist as provable exercises that are feasible with more time.

¹⁹This is equivalent to cutting open the torus along a meridional disk, twisting it some number of full revolutions, and gluing it back together.

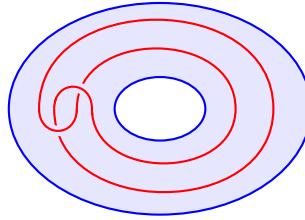


Figure 5.7: K_1 of a Whitehead double.

- **Two-strand cable** (of the companion knot): A satellite knot where K_1 is the unknot twisted up as in Figure 5.8.

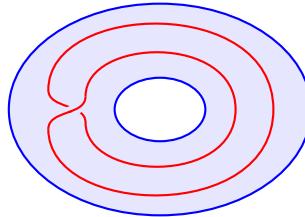


Figure 5.8: K_1 of a two-strand cable.

- Again, the two-strand cable will not be unique (even k -moves).
- If K_1 has only one strand reaching longitudinally around the solid torus, then $K_3 = K_1 \# K_2$.
 - Thus, satellite knots can be thought of as a broader category (a generalization) of composition in that all composite knots are satellite knots but not all satellite knots are composite knots.
 - Note that in this case, the knotted torus is a swallow-follow torus.
- In the same way that every composite knot can be factored into a unique set of prime knots, every satellite knot can be factored into a unique sequence of taking satellites.
- *Exercise 5.14:* Determine how the satellite knot in Figure 5.9 was made. Identify its companion and draw it inside the solid torus before the solid torus is knotted.

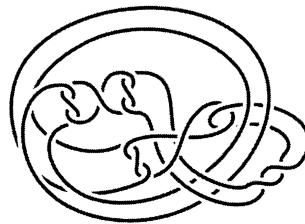


Figure 5.9: A satellite knot.

- I will not answer all parts of this exercise.
- The companion knot is the figure-eight knot, though.
- Furthermore, it appears to be constructed from the composition of 3 trefoil knots, several k -moves, and was originally an unknot similar to the one in Figure 5.7.
- **Cable knot** (of the companion knot): A satellite knot where K_1 is a torus knot.
 - Cable knots can be iterative, i.e., cables on cables on cables on torus knots.

5.3 Hyperbolic Knots

- It is suspected that the vast majority of knots are hyperbolic knots.
 - By “the vast majority of an infinite set,” “we mean that for all of the prime knots of n or fewer crossings, a certain large percentage of them are hyperbolic. As n approaches ∞ , that percentage is expected to approach 100%” (119).
- “William Thurston proved in 1978 that the only knots that are not hyperbolic knots are torus knots and satellite knots” (119).
 - Note that this statement includes all composite knots as satellite knots, a disputed claim.

- *Interesting tidbit on hyperbolic knot history.*
- **Hyperbolic knot:** “A knot that has a complement that can be given a **metric** of constant curvature -1 ” (120).
- **Metric:** A method of measuring distance.
- **Euclidean metric:** The typical way to measure distance in three-space, i.e., through the use of Equation 5.2.

$$d(P_0, P_1) = \sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2 + (z_1 - z_0)^2} \quad (5.2)$$

- On curvature:
 - A sphere has positive curvature — at a point on the sphere, cross sections of the sphere through that point yield circles that all curve away in the same direction.
 - A plane has zero curvature — at a point on the plane, lines on the plane through that point are flat.
 - A saddle has negative curvature — at the central point, one cross section opens up while another opens down.
- The Euclidian metric is a **flat metric**.
- **Flat metric:** A metric with curvature 0.
- **Hyperbolic metric:** A metric with curvature -1 .
- **Hyperbolic geometry:** The geometry that results from applying a hyperbolic metric to three-space.
- **Hyperbolic three-space:** “The simplest example of a three-dimensional space that has a hyperbolic metric” (121). *Also known as \mathbb{H}^3 .*
- The model of \mathbb{H}^3 considered herein is called the **Pointcaré model**.
- The points considered by this model are those of three-space inside the unit ball, or the following in set notation.

$$\mathbb{H}^3 = \{(x, y, z) : x^2 + y^2 + z^2 < 1\} \quad (5.3)$$

- On measuring distance in \mathbb{H}^3 :
 - Refer to Figure 5.10 throughout the following description.
 - Let P_1 and P_2 be two points in \mathbb{H}^3 .
 - If P_1 and P_2 both lie on a diameter, then define the particular path between them to lie along that diameter.
 - Otherwise, P_1 and P_2 lie on a circle that intersects both of them, and then the unit ball perpendicularly.
 - Such a circle is called a **geodesic**.

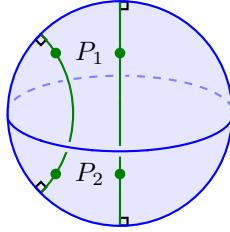


Figure 5.10: Paths between points in \mathbb{H}^3 .

- For any two points on a geodesic, the shortest distance between them also lies on the curve.
- Play the same role as straight lines in Euclidean space.
- “The shortest path in hyperbolic three-space from P_1 to P_2 is the path within the arc of a circle or vertical line from P_1 to P_2 ” (121).
 - Call this path w .
- To measure the distance between P_1 and P_2 , use the following path integral^[20] where r is the distance to the origin.

$$d(P_1, P_2) = \int_w \frac{2}{1 - r^2} ds \quad (5.4)$$
- **Geodesic** (in \mathbb{H}^3): “Any arc of a circle or diameter in \mathbb{H}^3 that is perpendicular to the unit sphere” (121).
- “The angles of any triangle in hyperbolic three-space add up to less than 180° ^[21]” (122).
- Use tetrahedral pieces of \mathbb{H}^3 to obtain a hyperbolic manifold.
 - The edges of such tetrahedra are geodesics, and the faces are geodesic planes (instead of the circular intersection in Figure 5.10, use a sphere that satisfies the same conditions).
- Similar to gluing together triangles (Figure 4.2) to form surfaces, glue together tetrahedra in \mathbb{H}^3 to form a knot complement.
 - “If the tetrahedra that we glue together are actually hyperbolic tetrahedra, in that they sit inside hyperbolic space, and if we glue them together along their faces so that their faces match without distortion, in order that the hyperbolic method of measuring distances within the individual tetrahedra match, the result is a hyperbolic knot complement” (123).
- We can then extrapolate the distance metric of \mathbb{H}^3 (applicable to the edges of a tetrahedron) to measure the entire knot complement hyperbolically.
- If this can be done, then the knot is a hyperbolic knot.
- The measurements facilitated by this process allow for the calculation of the **hyperbolic volume**.
- **Hyperbolic volume:** A positive real number that is “the sum of the volumes of the individual hyperbolic tetrahedra that make up the knot complement” (123).
 - Although a knot complement might seem like it’s infinite, measured hyperbolically, it’s finite.
- Hyperbolic volume is a knot invariant.
 - In fact, it is incredibly useful in distinguishing knots.
 - However, it obviously cannot distinguish knots that are not hyperbolic.

^[20]Perhaps a fun topic to explore at a later date.

^[21]Similar to the discussion of Non-Euclidian geometry and the parallel postulate in *Journey Through Genius*.

- It struggles with k -moves sometimes.
- Mutants have the same hyperbolic volume.
- The figure-eight knot does, indeed, have the smallest hyperbolic volume (Meyerhoff and Cao, 1997).
- To compute the hyperbolic volume (very surface-level):
 - Cut the complement into tetrahedra.
 - Place the tetrahedra in hyperbolic three-space.
 - In order to ensure that they properly mesh in \mathbb{H}^3 , a set of polynomial equations must be satisfied.
 - “If there are n tetrahedra, we obtain n polynomial equations in n variables” (125).
 - These variables are complex variables.
 - Use numerical methods to solve the system.
 - The solution determines the hyperbolic metric on each tetrahedron.
 - The volume can then be computed for each tetrahedron individually.
 - The individual volumes can be summed to get the hyperbolic volume.
- This process can be carried out by hand, but there is also a computer program to do it.

5.4 Braids

- **Braid:** “A set of n strings, all of which are attached to a horizontal bar at the top and at the bottom” (127).
 - See Figure 5.11a.
- Each strand can intersect a horizontal plane between the horizontal bars only once.
 - Each strand must constantly head down — it cannot loop back upward.
- **Equivalent** (braids): Two braids such that one can be rearranged to look like the other “without passing any strings through one another or themselves while keeping the bars fixed and keeping the strings attached to the bars” (128).
 - The strings cannot be pulled over the upper bar or beneath the bottom bar in rearranging a braid.
 - Think of the bars as infinite sheets to discourage this urge.



(a) A braid.



(b) The closure of the braid.

Figure 5.11: A braid and its closure.

- **Closure** (of a braid): The link formed by pulling the bottom bar around and gluing it to the top bar.
 - See the difference between the braid in Figure 5.11a and the closure of that braid in Figure 5.11b.

- Imagine an axis perpendicular to the page through the center of Figure 5.11b.
 - The closure is wrapped around this axis.
- **Closed braid representation** (of a knot): A projection where there exists “a choice of orientation on the knot so that, as we traverse the knot in that direction, we always travel clockwise around the axis without any backtracking” (128).
- *Exercise 5.16:* Draw closed braid representations of the knots in Figure 5.12. Hint: You may want to make them out of string or an extention cord and then see if you can rearrange them to travel clockwise around a pencil. Another option is to make a choice of an axis and then start to have the knot travel around the axis clockwise. Whenever it starts to travel counterclockwise, pass that portion of the knot through the axis to fix the problem. It is recommended that you do not try to use the technique that follows this problem in order to do it.

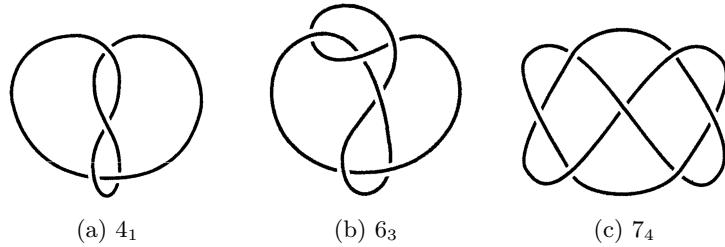


Figure 5.12: Draw these knots as closed braids.

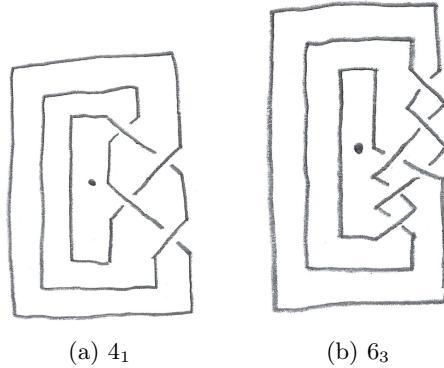


Figure 5.13: Solution to *Exercise 5.16* (incomplete).

- “Every knot or link is a closed braid” (129).
- This was proven in 1923 by J. W. Alexander of the Alexander polynomial (see Section 6).
- Let’s take a look at this proof.
 - This proof utilizes the idea of bridges from Section 3.2.
 - Let L be a knot or link.
 - Assign an orientation to L .
 - For each subarc (strand between crossings), choose a point along said subarc.
 - Choose a point and label it P_1 .
 - Proceed along the orientation of the knot, labeling each successive point P_i through P_n . These points will be vertices (the places where overpasses and underpasses pass through the projection plane).

- Notice that the strand between P_1 and P_2 is an overpass while the strand between P_2 and P_3 is an underpass. In fact, the strand between P_{2i-1} and P_{2i} is always an overpass while the strand between P_{2i} and P_{2i+1} is always an underpass.
 - Isotope the projection so that the understrands are lined up vertically next to each other. Make sure not to add any new crossings.
 - Let A be a straight line in the projection plane that is perpendicular to each underpass.
 - Each of the underpasses obviously cross A only once, but make sure that the overpasses only cross A once, too.
 - Every overpass crosses A an odd number of times (because one must cross at least once to move from P_{2i-1} to P_{2i} and any other crossings could be thought of as Type II Reidemeister moves, which add 2 crossings each time).
 - To resolve any issues, go to the problematic overpass's second crossing of A along the orientation.
 - Push it beneath A , reating a new underpass, a bridge that crosses A once, and a bridge that crosses A two fewer times.
 - Repeat this process for all problem areas.
 - For all underpasses, start at P_{2i} , decrease z until the crossing with A , and then increase z to intersect with P_{2i+1} .
 - For all overpasses, first make sure that all strands lie over the region contained by the underpasses. If the overpasses have to cross one another, that is fine. Now, start at P_{2i-1} , increase z until the crossing with A , and then decrease z to intersect P_{2i} .
 - The result is a closed braid with A as an axis.
- **Braid index:** “The least number of strings in a braid corresponding to a closed braid representation of the link” (132). *Also known as $b(L)$.*

Knot	unknot	3_1	4_1
Braid index	1	2	3

Table 5.1: Some braid indices.

- *Exercise 5.17:* Describe all of the knots and links of braid index 2.
 - All of the knots and links of braid index have crossings that can be undone by a single k -move, i.e., all such knots and links are k -equivalent to the trivial link of 2 components.
 - All such knots and links are $(k, 2)$ -torus knots or links.
 - If k is even, the manifold is a link of 2 components of linking number $\frac{k}{2}$.
 - If k is odd, the manifold is the k_1 knot (Alexander-Briggs notation).
- The braid index is a link invariant.
 - Upper bounds are easy to find.
 - Lower bounds, however, are difficult to find. One way uses polynomials (see Section 6).
- “The braid index of a knot is equal to the least number of Seifert circles in any projection of the knot” (132).
- Given a nonsplittable link L with crossing number $c(L)$ and braid index $b(L)$, Equation 5.5 holds.

$$2(b(L) - 1) \leq c(L) \tag{5.5}$$

- To describe a braid, use a **word**.
- **Word** (for a braid): Given a braid with n strings, first ensure that all crossings are at different vertical heights. Working one's way down the braid from the top, denote the i^{th} string crossing over the $(i+1)^{\text{th}}$ string by σ_i and the i^{th} string crossing under the $(i+1)^{\text{th}}$ string by σ_i^{-1} . Concatenate all such notations.

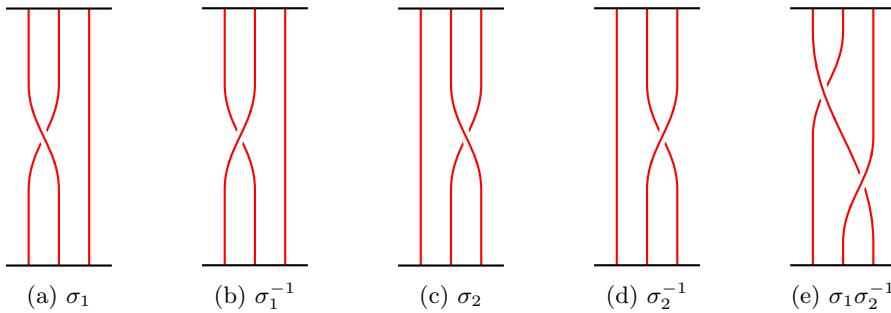


Figure 5.14: Denoting crossings in a braid.

- *Exercise 5.18:* Draw the five-string braid given by the word $\sigma_1\sigma_2\sigma_3\sigma_2^{-1}\sigma_4\sigma_1^{-1}$.

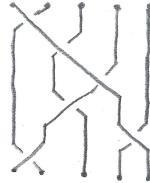


Figure 5.15: Solution to *Exercise 5.18*.

- *Exercise 5.19:* Identify the knot that has four-string braid word $(\sigma_1^{-1}\sigma_2\sigma_3^{-1})^2\sigma_3^3$. It is one of the six-crossing knots.
 - After drawing the closure of the braid, reduce crossings down to six.
 - At this point, a 4 tangle and a 2 tangle can be picked out.
 - Therefore, this is the 6_1 knot.
- *Exercise 5.20:* Find a word such that the closure of the corresponding braid gives the 6_3 knot.
 - Use Figure 5.13b, which gives a closed braid representation of 6_3 . From this projection, it is possible to pick find the word.
 - Therefore, the word is $\sigma_2\sigma_2\sigma_1^{-1}\sigma_2\sigma_1^{-1}\sigma_1^{-1}$.
- *Exercise 5.21:* Show that the closure of the n -string braid $(\sigma_1\sigma_2\cdots\sigma_{n-1})^m$ is a knot if and only if n and m are relatively prime.
 - Given this particular word, every strand “shifts” m places from top to bottom.
 - Consider a scenario where $n = 3$ and $m = 1$.
 - The left strand at the top becomes the rightmost strand at the bottom by the definition of the word.
 - This means that the right strand at the top becomes the middle strand at the bottom and the middle strand at the top becomes the left strand at the bottom.
 - Now consider when $n = 3$ and $m = 2$.

- The above shifting of strands takes place, but then it takes place again, i.e., every strand shifts twice.
- Thus, from top to bottom, left \rightarrow right \rightarrow middle, middle \rightarrow left \rightarrow right, and right \rightarrow middle \rightarrow left.
- Take a strand that is i^{th} from the left at the top.
- Assuming that $i > m$, the strand will be $(i - m)^{\text{th}}$ from the left at the bottom.
- Therefore, it will reenter the top of the braid in the $(i - m)^{\text{th}}$ from the left position.
- It may then leave at the $(i - 2m)^{\text{th}}$ from the left position (or it may shift back to the right side).
- Regardless, it will pass through some number of strands of the braid in this fashion.
- Only if n and m are relatively prime will one loop pass through every strand in a nontrivial fashion.
- *Exercise 5.22:* Check that the closure of the n -string braid $(\sigma_1 \sigma_2 \cdots \sigma_{n-1})^m$ is simply the (m, n) -torus knot (assuming m and n are relatively prime).
 - Indeed, this is the case.
- The word notation has some useful advantages.

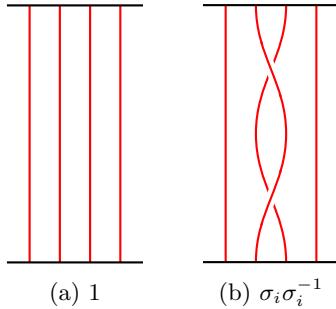


Figure 5.16: Braids equivalent via a Type II Reidemeister move.

- If a $\sigma_i \sigma_i^{-1}$ term appears somewhere within a word, it corresponds to the crossings shown in Figure 5.16.
- These crossings can easily be undone via a Type II Reidemeister move.
- In the same way, $\sigma_i \sigma_i^{-1} = 1$ by the definition of an inverse, so the term reduces out of the word.
- *Exercise 5.23:* If a given word represents a knot with a particular orientation, what word gives the same knot but with the opposite orientation?
 - Take the terms of the word and reverse their order.
- To multiply two n -string braids, just stack them on top of each other.
 - This is reflected by the notation — stacking tangles means that the word of the top one will be immediately followed by the word of the bottom one.
- *Exercise 5.24 (a):* Describe an n -string braid I_n that acts as an identity when it multiplies any other n -string braid. That is to say, multiplying a braid B by this braid yields the braid B back again.
 - I_n is the trivial braid of n -strings.
 - Multiplication extends the bottom strands, but they can be isotoped back to being shorter.
- *Exercise 5.24 (b):* Describe an inverse braid for a given braid B . That is to say, find a braid B^{-1} that has the property that $BB^{-1} = I_n = B^{-1}B$.

- B^{-1} is just B but with the inverse of the terms of the word listed in reverse order.
- For example, if $B = \sigma_1\sigma_2^{-1}\sigma_2$, then $B^{-1} = \sigma_2^{-1}\sigma_2\sigma_1^{-1}$.
- *Exercise 5.24 (c):* Show that the multiplication of n -string braids is associative, namely $B_1(B_2B_3) = (B_1B_2)B_3$ for any three n -string braids B_1 , B_2 , and B_3 .
 - Since the terms of a word follow the normal rules of multiplication, so do words, themselves^[22].
- B_n will denote the **group** of n -string braids (developed throughout *Exercise 5.24*).
- **Group:** A set of elements and a way to multiply the elements such that:
 1. “There exists an identity element that, when multiplying any element, doesn’t change it;
 2. Every element has an inverse;
 3. The multiplication is associative, that is, $a(bc) = (ab)c$ for any three elements a , b , and c of the group” (135).
- A particular element of B_n is an equivalent class of braids.
- *Exercise 5.26:* Using a picture, show that the two braids $\sigma_i\sigma_{i+1}\sigma_i$ and $\sigma_{i+1}\sigma_i\sigma_{i+1}$ are equivalent.

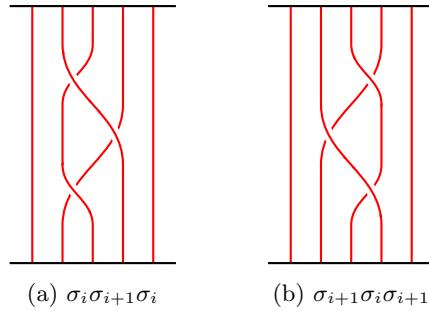


Figure 5.17: Braids equivalent via a Type III Reidemeister move.

- If $\sigma_i\sigma_j$ is changed to $\sigma_j\sigma_i$ when $|i - j| > 1$, equivalence is maintained.

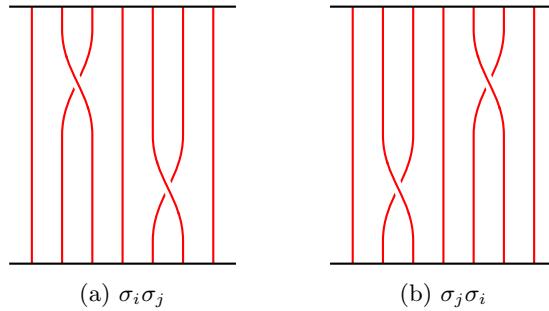


Figure 5.18: Braids equivalent via a shifting of unconnected strings.

- Not based on a Type I Reidemeister move.
- Just changes the relative height of two distinct twists.
- “Two words w_1 and w_2 represent the same braid if and only if we can get from the one word to the other by a sequence of these operations” (136).

²²This is like proving that a vector space is closed under addition and scalar multiplication using the fact the regular numbers (components of vectors) are obey addition and scalar multiplication.

- *Exercise 5.27:* Show that the two words $w_1 = \sigma_2\sigma_1\sigma_2\sigma_1\sigma_2^{-1}\sigma_4$ and $w_2 = \sigma_2^2\sigma_3\sigma_1\sigma_3^{-1}\sigma_4$ represent the same five-string braid, up to equivalence.

$$\begin{aligned}
 w_1 &= \sigma_2\sigma_1\sigma_2\sigma_1\sigma_2^{-1}\sigma_4 \\
 \xrightarrow{\text{Rule 2}} &\sigma_2\sigma_2\sigma_1\sigma_2\sigma_2^{-1}\sigma_4 \\
 \xrightarrow{\text{Rule 1}} &\sigma_2\sigma_2\sigma_1\sigma_4 \\
 \xrightarrow{\text{Condense}} &\sigma_2^2\sigma_1\sigma_4 \\
 \xrightarrow{\text{Rule 1}} &\sigma_2^2\sigma_1\sigma_3\sigma_3^{-1}\sigma_4 \\
 \xrightarrow{\text{Rule 3}} &\sigma_2^2\sigma_3\sigma_1\sigma_3^{-1}\sigma_4 = w_2
 \end{aligned}$$

- **Markov equivalent** (braids): Two braids whose “closures yield the same oriented link” (137).
- **Markov’s theorem:** Two braids are Markov equivalent if and only if they can be transformed into one another via the previous three operations, **conjugation**, and **stabilization**.
- **Conjugation:** Multiply the word by σ_j at the beginning and σ_j^{-1} at the end (or vice versa).

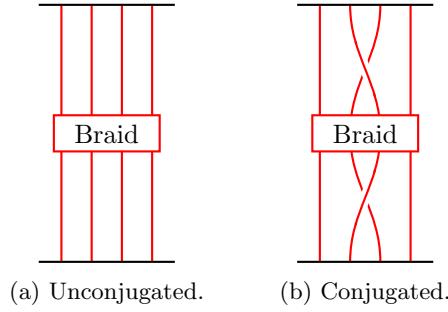


Figure 5.19: Braids equivalent via conjugation by σ_j .

- **Stabilization:** Multiply the word w corresponding to an n -string braid by σ_n to yield $w\sigma_n$. Alternatively, given an $(n+1)$ -string braid $w\sigma_n$ where σ_n is only present as the last term, remove σ_n to yield w . Multiplying by or removing σ_n^{-1} is also valid.

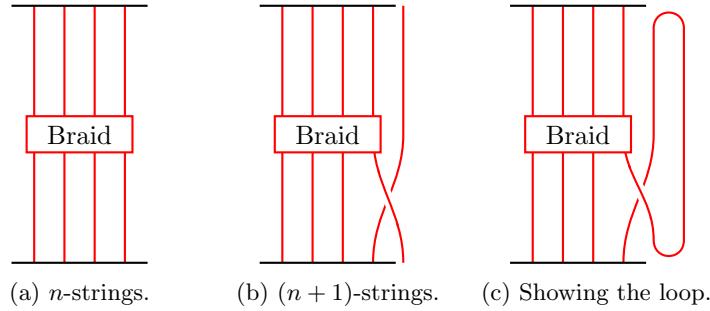


Figure 5.20: Braids equivalent via stabilization by σ_n .

- This move changes the number of strings/loops in a closed braid.
- It is kind of like a Type I Reidemeister move.
- **Letter:** An individual $\sigma_i^{\pm 1}$ term in a word.
- **Markov moves:** Formally, the five listed operations (Figures 5.16, 5.17, 5.18, 5.19, and 5.20).

5.5 Almost Alternating Knots

- **Almost alternating knot:** A nonalternating knot that has a projection where one crossing change would make the knot alternating.
- “Of the 393 nonalternating knots and links of eleven or fewer crossings, all but at most three are almost alternating” (140).
- “An algebraic knot that contains no negative signs in its Conway notation is in fact an alternating knot” (141).
- “An algebraic link that has exactly one negative sign in its Conway notation has an almost alternating projection^[23]” (141).
 - This statement proves that all but 17 of Conway’s nonalternating knots are almost alternating.
- In this section, results known for alternating knots (many are listed in Section 6) will be generalized to almost alternating knots.
- “Every alternating knot or link has an almost alternating projection” (141).
 - A Type II Reidemeister move on the alternating projection produces an almost alternating projection.
- “A prime almost alternating knot is either a torus knot or a hyperbolic knot” (142).
 - Thus, no satellite knot is an almost alternating knot.
 - This is an extention of the previously known fact that a prime alternating link is either a torus link or a hyperbolic link.
 - Because the only alternating torus links are the closures of braids where $b(L) = 2$ (see *Exercise 5.17*), almost all prime alternating links are hyperbolic.
- **Tongue move:** A Type II Reidemeister move followed by a Type III Reidemeister move about the nonalternating crossing^[24].

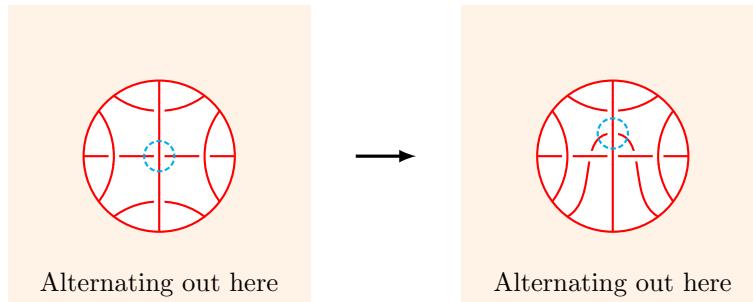


Figure 5.21: A tongue move.

- **m -almost alternating knot:** “A knot that has a projection where m crossing changes would make the projection alternating, and the knot has no projection that could be made alternating in fewer crossing changes” (144).
 - Alternating knots are 0-almost alternating knots and almost alternating knots are 1-almost alternating knots.
- The simplest Whitehead double of the trefoil knot is 2-almost alternating.

²³Available as an exercise.

²⁴This is a Type II bridging move!

- It is not alternating.
 - It is a satellite knot, and, thus, is not almost alternating.
 - There exists a projection of it where two crossing changes would make it alternating.
- *Exercise 5.33:* Show that every knot is m -almost alternating for some m .
- Take a projection of a knot.
 - Assign an orientation.
 - Choose a crossing from which to begin traversing the knot.
 - Leave on the understrand.
 - If the next crossing is not passed through on the overstrand, flip the crossing. Otherwise, continue along.
 - If the crossing after that is not passed through on the understrand, flip the crossing. Otherwise, continue along.
 - Continue in this manner until you arrive at the original crossing on the understrand. Make sure to not change any single crossing more than once.
 - Tally up the crossing changes. This is an upper bound on m .
- *Exercise 5.34:* Show that if a knot has an m -almost alternating projection, then it has an $(m+1)$ -almost alternating projection.
- Take the initial projection and perform any Type II Reidemeister move that adds two crossings (see above).
 - This new projection will be $(m+1)$ -almost alternating.
- *Exercise 5.35:* Show that if a knot K has a projection with n crossings, then it is m -almost alternating for some $m \leq \frac{n}{2}$.
- The worst case scenario (the projection that would need the most crossing changes) is a projection where it is possible to travel through every crossing on the understrand before reaching an overpass (or vice versa).
 - In this case, flip every other crossing ($\frac{n}{2}$ crossings if n is even or $\frac{n-1}{2}$ crossings if n is odd).
 - Therefore, $m \leq \frac{n}{2}$.
- m is similar to the unknotting number.
- The two are extreme opposites: $u(K)$ measures the distance from the “simplest” version of the knot while m measures the distance from the “most complicated” version of the knot.
- **Toroidally alternating knot:** A knot that can be projected onto a torus so that it is alternating when viewed from the outside of the torus. Furthermore, every closed curve on the torus that doesn’t bound a disk^[25] must intersect the projection.
- It is easy to tell if a projection on a torus is alternating. The second qualifier can be determined via this process:
- Treat the crossings as vertices and subarcs as edges so that the knot projection becomes a graph.
 - Cut along the edges to yield disks.
 - “Any closed curve on the torus that doesn’t intersect the projection must lie in one of these disks. But then it will bound a disk on the torus. Hence, this is a toroidally alternating knot” (145).
- Any almost alternating knot is toroidally alternating.
- Any alternating knot is toroidally alternating^[26].
- “Toroidally alternating knots behave similarly to alternating and almost alternating knots” (146).

²⁵So then, a longitude?

²⁶These statements exist as exercises.

6 Polynomials

6.1 The Bracket Polynomial and the Jones Polynomial

- A knot polynomial is one of the most successful ways to tell knots apart.
 - From a projection, we can find a polynomial.
 - If two projections yield the same polynomial, they represent the same knot.
 - If two projections yield different polynomials, they represent different knots.
- Therefore, knot polynomials are knot invariants.
- **Laurent polynomial:** A polynomial with both positive and negative powers of t .
- J. Alexander introduced the first knot polynomial in 1928.
- John Conway found a way to calculate the Alexander polynomial of a link using a **skein relation** in 1969.
- **Skein relation:** An equation that relates the polynomial of a link L_1 to the polynomial of a link L_2 obtained by changing the crossings in a projection L_1 .
- Vaughan Jones found a new polynomial for knots and links in 1984.
- Four months later, the HOMFLY polynomial was announced.
 - So named for its discoverers: Hoste, Ocneanu, Millet, Freyd, Lickorish, and Yetter.
- Since then, many new polynomials have been discovered.
- Many of these discoveries have introduced areas of mathematics previously unrelated to knot theory to knot theory.
- The first polynomial developed will be the **bracket polynomial**.
 - We will take the approach of a mathematician trying to discover a polynomial invariant for knots and links.
- Let $\langle K \rangle$ denote the bracket polynomial of a knot K .
- Begin by listing basic requirements that we want the polynomial to have.
 - The polynomial should be 1 for the trivial knot.

$$\text{Rule 1 : } \langle \textcircled{O} \rangle = 1$$

- The polynomial for a link should be able to be obtained in terms of the polynomials of simpler links.
 - Use a skein relation.
 - Split open a crossing of L_1 vertically and horizontally in order to obtain two new link projections L_2 and L_3 with one fewer crossing.
 - The bracket polynomial of L_1 can be given as a linear combination of those for L_2 and L_3 .
Use A and B as placeholder coefficients.
 - Note that the second equation is just the first equation rotated 90° .

$$\begin{aligned} \text{Rule 2 : } & \langle \text{X} \rangle = A \langle \text{O} \rangle + B \langle \text{X} \rangle \\ & \langle \text{X} \rangle = A \langle \text{X} \rangle + B \langle \text{O} \rangle \end{aligned}$$

- The polynomial should be able to account for adding in a trivial component to a link.

$$\text{Rule 3 : } \langle L \cup \textcolor{red}{O} \rangle = C \langle L \rangle$$

- The most important criterion is that this polynomial be an invariant. Therefore, it must not change under the Reidemeister moves.
- Begin with a Type II Reidemeister move's effect.

- In other words, we want $\langle \textcolor{red}{X} \rangle = \langle \textcolor{red}{O} \rangle$. See the following expansion by the three previous rules.

$$\begin{aligned}
\langle \textcolor{red}{X} \rangle &\stackrel{?}{=} A \langle \textcolor{red}{X} \rangle + B \langle \textcolor{blue}{Y} \rangle \\
&\stackrel{?}{=} A(A \langle \textcolor{red}{X} \rangle + B \langle \textcolor{red}{X} \rangle) + B(A \langle \textcolor{red}{O} \rangle + B \langle \textcolor{red}{X} \rangle) \\
&\stackrel{?}{=} A(A \langle \textcolor{red}{X} \rangle + BC \langle \textcolor{red}{X} \rangle) + B(A \langle \textcolor{red}{O} \rangle + B \langle \textcolor{red}{X} \rangle) \\
&\stackrel{?}{=} (A^2 + ABC + B^2) \langle \textcolor{red}{X} \rangle + BA \langle \textcolor{red}{O} \rangle \\
&\stackrel{?}{=} \langle \textcolor{red}{O} \rangle
\end{aligned} \tag{6.1}$$

- To make the polynomial unchanged, BA must equal 1 and $A^2 + ABC + B^2$ must equal 0.

- Thus, define $B = A^{-1}$.

- This definition makes the following equal to Equation 6.1.

$$(A^2 + C + A^{-2}) \langle \textcolor{red}{X} \rangle + \langle \textcolor{red}{O} \rangle \tag{6.2}$$

- Now define $C = -A^2 - A^{-2}$.

- This definition makes the following equal to Equation 6.2 (and, by extention, Equation 6.1).

$$\langle \textcolor{red}{O} \rangle$$

- Therefore, the equality associated with Equation 6.1 holds.

- Based on the substitutions for B and C , Rules 1-3 become the following.

$$\begin{aligned}
\text{Rule 1 : } \langle \textcolor{red}{O} \rangle &= 1 \\
\text{Rule 2 : } \langle \textcolor{red}{X} \rangle &= A \langle \textcolor{red}{O} \rangle + A^{-1} \langle \textcolor{red}{X} \rangle \\
&\quad \langle \textcolor{red}{X} \rangle = A \langle \textcolor{red}{X} \rangle + A^{-1} \langle \textcolor{red}{O} \rangle \\
\text{Rule 3 : } \langle L \cup \textcolor{red}{O} \rangle &= (-A^2 - A^{-2}) \langle L \rangle
\end{aligned}$$

- Now let's look at the Type III Reidemeister move.

$$\begin{aligned}
\langle \textcolor{red}{X} \textcolor{red}{X} \rangle &= A \langle \textcolor{red}{X} \rangle + A^{-1} \langle \textcolor{red}{X} \textcolor{red}{X} \rangle \\
&= A \langle \textcolor{red}{X} \rangle + A^{-1} \langle \textcolor{red}{X} \textcolor{red}{X} \rangle \\
&= \langle \textcolor{red}{X} \textcolor{red}{X} \rangle
\end{aligned}$$

- Note that the transition from the first to the second line above is allowed because the bracket polynomial has been proven to be unchanged by Type II Reidemeister moves.

- Calculate the polynomial for a trivial link of two components.

$$\begin{aligned}
\langle \textcolor{red}{O} \cup \textcolor{red}{O} \rangle &= -(A^2 + A^{-2}) \langle \textcolor{red}{O} \rangle \\
&= -(A^2 + A^{-2}) 1 \\
&= -A^2 - A^{-2}
\end{aligned} \tag{6.3}$$

- *Exercise 6.1:* What would the bracket polynomial of the usual projection of the trivial link of n components be?

- It will be given by Equation 6.4. This could be proved via mathematical induction.

$$(-A^2 - A^{-2})^{n-1} \quad (6.4)$$

- Calculate the bracket polynomial for the Hopf link (see Figure 1.17b).

$$\begin{aligned}
<\text{H}\text{L}> &= A <\text{O}\text{O}> + A^{-1} <\text{G}\text{G}> \\
&= A (A <\text{O} \cup \text{O}> + A^{-1} <\text{O}>) + A^{-1} (A <\text{O}> + A^{-1} <\text{O} \cup \text{O}>) \\
&= A (A (-A^2 - A^{-2}) + A^{-1}(1)) + A^{-1} (A(1) + A^{-1} (-A^2 - A^{-2})) \\
&= A (-A^3 - A^{-1} + A^{-1}) + A^{-1} (A - A - A^{-3}) \\
&= A (-A^3) + A^{-1} (-A^{-3}) \\
&= -A^4 - A^{-4}
\end{aligned} \quad (6.5)$$

- *Exercise 6.2:* Find the bracket polynomial for the trefoil knot (see Figure 1.1b).



(a) Trefoil closure.

$$<\text{3}\text{C}> = A <\text{G}\text{G}> + A^{-1} <\text{U}> \quad (1)$$

$$= A \left(A <\text{O}\text{O}> + A^{-1} <\text{O}> \right) + A^{-1} <\text{O}\text{O}> \quad (2)$$

$$= A \left(A \left(A <\text{O}\text{O}> + A^{-1} <\text{O}> \right) + A^{-1} \left(A <\text{O}\text{O}> + A^{-1} <\text{O}> \right) \right) + A^{-1} <\text{O}\text{O}> \quad (3)$$

$$= A \left(A \left(A (-A^2 - A^{-2})^{2-1} + A^{-1} (-A^2 - A^{-2}) \right) + A^{-1} \left(A (-A^2 - A^{-2}) + A^{-1} (1) \right) \right) + A^{-1} (-A^4 - A^{-4}) \quad (4)$$

$$= A \left(A \left(A (A^4 + 2 + A^{-4}) - A - A^{-3} \right) + A^{-1} (-A^2 - A^{-1} + A^{-1}) \right) - A^3 - A^{-5} \quad (5)$$

$$= A \left(A (A^5 + 2A + A^{-3} - A - A^{-3}) - A^2 \right) - A^3 - A^{-5} \quad (6)$$

$$= A (A^6 + 2A^2 - A^2 - A^2) - A^3 - A^{-5} \quad (7)$$

$$= A^7 - A^3 - A^{-5} \quad (8)$$

(b) Derivation of the bracket polynomial.

Figure 6.1: Solution to *Exercise 6.2*.

- To ensure that all of the crossings are aligned the same way, view the trefoil as its closed braid representation (Figure 6.1a).
 - For simplicity's sake, only the braid (not its closure) is represented in the equations of Figure 6.1b.
- Start in Equation 1 of Figure 6.1b by employing Rule 2 on the top crossing.
 - At this point, the left term is A times the unknot while the right term is A^{-1} times the Hopf link (see Figure 1.17b).
 - While the right term can be substituted for the Hopf link, the left term must continue to be split using Rules 1-3 because the effect of the Type I Reidemeister move on the Bracket polynomial has yet to be dealt with.
- From Equation 1 to 2 of Figure 6.1b, substitute the right term for the Hopf link and employ Rule 2 to the middle (now upper) crossing of the left term.
- From Equation 2 to 3 of Figure 6.1b, employ Rule 2 on the bottom (now only) crossing of each of the left two terms.
- From Equation 3 to 4 of Figure 6.1b, substitute knot polynomials in for all of the $\langle K \rangle$ terms.
 - The leftmost term is an unlink of 3 components, so substitute in Equation 6.4 with $n = 3$.
 - The second-from-the-left and middle terms are unlinks of 2 components, so substitute in Equation 6.3.
 - The second-from-the-right term is the unknot, so substitute in 1 by Rule 1.
 - The rightmost term is the Hopf link, so substitute in Equation 6.5.
- From Equation 4 to 8 of Figure 6.1b, simplify the algebra.
- Equation 8 of Figure 6.1b is the final answer.

- Now let's look at the Type I Reidemeister move.

$$\begin{aligned}
 \langle \text{unknot} \rangle &= A \langle \text{unknot} \rangle + A^{-1} \langle \text{unknot} \rangle \\
 &= A(-A^2 - A^{-2}) \langle \text{unknot} \rangle + A^{-1} \langle \text{unknot} \rangle \\
 &= -A^3 \langle \text{unknot} \rangle
 \end{aligned}$$

$$\begin{aligned}
 \langle \text{unknot} \rangle &= A \langle \text{unknot} \rangle + A^{-1} \langle \text{unknot} \rangle \\
 &= A \langle \text{unknot} \rangle + A^{-1} (-A^2 - A^{-2}) \langle \text{unknot} \rangle \\
 &= -A^{-3} \langle \text{unknot} \rangle
 \end{aligned} \tag{6.6}$$

- Based on the above, the bracket polynomial *is* changed by a Type I Reidemeister move.
- Therefore, the bracket polynomial is not a knot invariant.
- However, it can still be modified as follows.
- Assign an orientation to the knot or link.
- At every crossing, assign a +1 or -1 as directed by Figure 1.18.
- **Writhe:** The sum of the +1s and -1s at every crossing. *Also known as $w(L)$.*
 - Note that this is different from the linking number in that every crossing contributes, not just those between different link components.
- *Exercise 6.3:* Calculate the writhe of the oriented link projection in Figure 6.2.
 - Table 6.1 displays the crossings (numbered from left to right and with a top/bottom subscript when necessary) and their corresponding value.
 - Summing these values gives $w(L) = -2$.

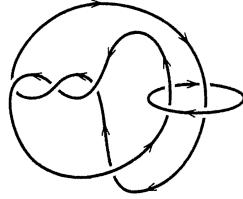


Figure 6.2: Determine the writhe of this link projection.

Crossing number	1	2	3	4	5_t	5_b	6_t	6_b
Writhe value	-1	-1	-1	+1	-1	-1	+1	+1

Table 6.1: Writhe values.

- *Exercise 6.4:* Show that the writhe of a link projection is invariant under Reidemeister moves II and III.
 - Consider both strands of Figure 1.8b to have a parallel orientation.
 - In either case, the top crossing is +1 and the bottom crossing is -1, canceling out.
 - If both strands have an antiparallel orientation, the top crossing is -1 and the bottom crossing is +1, again canceling out.
 - A similar (albeit longer) process can be carried out for every case in Figure 1.8c.
- On the other hand, a Type I Reidemeister move always changes the writhe by ± 1 .
- To counter this, define a new polynomial called the **X polynomial** as the following.

$$X(L) = (-A^3)^{-w(L)} \langle L \rangle \quad (6.7)$$

- Since both $\langle L \rangle$ and $w(L)$ are invariant under Type II and Type III Reidemeister moves, so is $X(L)$.
- Let's see how the X polynomial handles a Type I Reidemeister move.
 - Let the right half of Figure 1.8a be L and the left half be L' . The orientation will be top to bottom^[27].
 - First, note that $w(L') = w(L) - 1$.
 - Based on this and Equation 6.6, we find the following equality.

$$\begin{aligned} X(L) &= (-A^3)^{-w(L)} \langle L \rangle \\ &= (-A^3)^{-w(L')+1} (-A^{-3}) \langle L' \rangle \\ &= (-A^3)^{-w(L')} (-A^3) (-A^{-3}) \langle L' \rangle \\ &= (-A^3)^{-w(L')} \langle L' \rangle \\ &= X(L') \end{aligned}$$

- Thus, the X polynomial does not change under any of the three Reidemeister moves.
- “Therefore, $X(L)$ is an invariant for knots and links” (153).
- *Exercise 6.5:* Compute $X(L)$ for the trefoil knot (Figure 1.1b) and for the Hopf link (Figure 1.17b). What happens if we change the orientation from the original used to calculate $X(L)$?

²⁷Note that if it was the opposite orientation, the math would still work out.

- By Equation 8 of Figure 6.1b, $\langle 3_1 \rangle = A^7 - A^3 - A^{-5}$.
- By Equation 6.7, $X(L) = (-A^3)^{-w(L)} \langle L \rangle$.
- An examination of Figure 6.1a reveals that $w(3_1) = -3$.
- Thus, $X(3_1) = (-A^3)^3 (A^7 - A^3 - A^{-5})$.
- Therefore, the following equation holds.

$$X(3_1) = -A^{16} + A^{12} + A^4 \quad (6.8)$$

- By Equation 6.5, $\langle 2_1^2 \rangle = -A^4 - A^{-4}$.
- An examination of Figure 1.17b reveals that $w(2_1^2) = 2$.
- Thus, $X(2_1^2) = (-A^3)^{-2} (-A^4 - A^{-4})$.
- Therefore, the following equation holds.

$$X(2_1^2) = -A^{-2} - A^{-10}$$

- Nothing would happen for the trefoil.
- Changing the orientation of one component of the Hopf link would cause the writhe to be multiplied by -1 . This would give an alternative X polynomial, namely, one with positive exponents instead of negative and the whole thing multiplied by -1 .

- *Exercise 6.6 should have a fascinating solution.*
- “The **Jones polynomial** is obtained from the X polynomial by replacing each A by $t^{-\frac{1}{4}}$ ” (154).
- *Exercise 6.7:* Use the solution to *Exercise 6.5* to write the Jones polynomial of the trefoil knot.
- If $A = t^{-\frac{1}{4}}$ in Equation 6.8, then the Jones polynomial can be written as follows.

$$V(t) = -t^{-4} + t^{-3} + t^{-1}$$

- Does the Jones polynomial distinguish every other knot from the unknot?
- The Jones polynomial of an (m, n) -torus knot is given by the following.

$$V(t) = t^{\frac{(m-1)(n-1)}{2}} \left(\frac{1 - t^{m+1} - t^{n+1} + t^{m+n}}{1 - t^2} \right)$$

- The Jones polynomial satisfies its own Skein relation.
 - “Let L_+ , L_- , and L_0 be three oriented link projections that are identical except where they appear in Figure [6.3]” (155).
 - This was the original Skein relation recognized by Vaughan Jones (see Section ??).
- *Exercise 6.8:* Use the skein relation of the bracket polynomial in order to show that the Jones polynomial of the three links in Figure 6.3 are related through Equation 6.9.

$$t^{-1}V(L_+) - tV(L_-) + \left(t^{-\frac{1}{2}} - t^{\frac{1}{2}} \right) V(L_0) = 0 \quad (6.9)$$

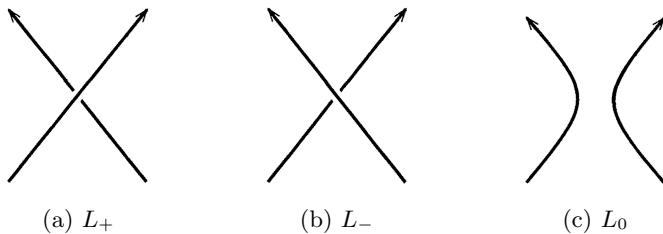


Figure 6.3: Three almost identical link projections.

- This is not a pretty solution. It was reverse engineered from Equation 6.9, so the algebra may not follow any logical pattern, but it is correct.
- The substitutions come from Rule 2, the definition of the X polynomial, and the fact that $A = t^{-\frac{1}{4}}$.

$$\begin{aligned}
 & <\cancel{\text{X}}> + (A^{-2} + A^2) <)> = <\cancel{\text{X}}> + (A^2 + A^{-2}) <)> \\
 & A^{-1} (A <\cancel{\text{X}}> + A^{-1} <)> + A^2 <)> = A (A <)> + A^{-1} <\cancel{\text{X}}> + A^{-2} <)> \\
 & \quad A^{-1} <\cancel{\text{X}}> + A^2 <)> = A <\cancel{\text{X}}> + A^{-2} <)> \\
 & \quad -A <\cancel{\text{X}}> + A^2 <)> = -A^{-1} <\cancel{\text{X}}> + A^{-2} <)> \\
 & A^4 (-A^{-3}) <\cancel{\text{X}}> + A^2 <)> = A^{-4} (-A^3) <\cancel{\text{X}}> + A^{-2} <)> \\
 & A^4 (-A^3)^{-w(L_+)} <\cancel{\text{X}}> + A^2 (-A^3)^{-w(L_0)} <)> = A^{-4} (-A^3)^{-w(L_-)} <\cancel{\text{X}}> + A^{-2} (-A^3)^{-w(L_0)} <)> \\
 & \quad A^4 V(L_+) + A^2 V(L_0) = A^{-4} V(L_-) + A^{-2} V(L_0) \\
 & \quad t^{-1} V(L_+) + t^{-\frac{1}{2}} V(L_0) = t V(L_-) + t^{\frac{1}{2}} V(L_0) \\
 & t^{-1} V(L_+) - t V(L_-) + \left(t^{-\frac{1}{2}} - t^{\frac{1}{2}} \right) V(L_0) = 0
 \end{aligned}$$

Figure 6.4: Solution to *Exercise 6.8*.

6.2 Polynomials of Alternating Knots

- Goal (from here to Equation 6.10): develop a second way to consider the bracket polynomial.
- Four regions of the projection plane (each bounded by two strands) intersect at a crossing (see Figure 6.5).

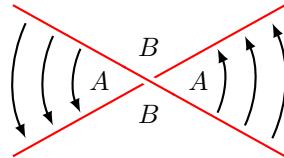


Figure 6.5: Labeling a crossing with A and B .

- Imagine rotating the overstrand counterclockwise over the understrand so that two of the regions disappear.
- Label each of these disappearing regions A .
- Label each region that did not disappear B .
- In Rule 2 for calculating the bracket polynomial, each crossing is split open two ways.
- **A -split:** The Rule 2 split that combines the two regions labeled A .
- **B -split:** The Rule 2 split that combines the two regions labeled B .
- Suppose L is a link of n crossings.
- Rule 2 allows the bracket polynomial of L to be calculated in terms of the bracket polynomials of L_1 and L_2 .
 - L_1 and L_2 each have one fewer crossing than L .
 - Note that L_1 and L_2 are obtained via an A -split and then a B -split.

- Continuing splitting crossings eventually gives the bracket polynomial of L in terms of the bracket polynomials of 2^n links.
 - Each of these links is an unlink of m components where $m \leq n$.
 - Each of these links is a **state**.
 - Denote one of these links L' .
- **State:** “A choice of how to split all of the n crossings in the projection of L ” (157).
- Let $|S|$ be the number of loops in L' .
- By Equation 6.4, the bracket polynomial of L' is $(-A^2 - A^{-2})^{|S|-1}$.
- However, each split requires that the bracket polynomial of the simpler link be multiplied by either A or A^{-1} .
 - Every time an A -split occurs, A is the factor multiplied.
 - Every time a B -split occurs, A^{-1} is the factor multiplied.
- Thus, the factor multiplied by L' is $A^{a(S)} A^{-b(S)}$.
 - Let $a(S)$ denote the number of A -splits needed to create S .
 - Let $b(S)$ denote the number of B -splits needed to create S .
- Thus, the contribution of L' to the bracket polynomial of L is $A^{a(S)} A^{-b(S)} (-A^2 - A^{-2})^{|S|-1}$.
- Therefore, the bracket polynomial of L is given by Equation 6.10.

$$\langle L \rangle = \sum_s \left(A^{a(S)} A^{-b(S)} (-A^2 - A^{-2})^{|S|-1} \right) \quad (6.10)$$

- Recall reduced alternating projections (see Figure 3.9).
- Prove that two reduced alternating projections of the same knot have the same number of crossings.
- First, define **span**.
- **Span** (of a polynomial): “The difference between the highest power that occurs in the polynomial and the lowest power that occurs in the polynomial” (159).
- The span of the bracket polynomial is a knot invariant.
 - This is because the bracket polynomial does not change under Type II and III Reidemeister moves and Type I Reidemeister moves change all exponents by ± 3 , which doesn’t affect the span (the delta of the exponents).
- Now, prove a lemma^[28].
- Lemma: If K has a reduced alternating projection of n crossings, then $\text{span}(\langle K \rangle) = 4n$.
 - Since the span doesn’t depend on the projection, choose the reduced alternating projection.
 - First, find the highest power. Then, find the lowest power.
 - Each state contributes a term in the form $A^{a(S)} A^{-b(S)} (-A^2 - A^{-2})^{|S|-1}$.
 - Upon expanding, the highest power of A occurs with $A^{a(S)} A^{-b(S)} (A^2)^{|S|-1}$.
 - In other words, we need to make $a(S) - b(S) + 2(|S| - 1)$ as large as possible.

²⁸**Lemma:** A subsidiary or intermediate theorem in an argument or proof (Google).

- When $a(S) = n$ and $b(S) = 0$, what is $|S|$?
- In an alternating projection, every region has only A s or only B s (see how white regions are A regions and blue regions are B regions in Figure 4.9).
- When all of the A -channels are opened, white “waters” flood the crossings, leaving only blue islands in Figure 4.9.
- Let D be the number of dark (shaded/blue) regions in the original projection. Thus, $|S| = D$.
- The highest power of A is $n + 2(D - 1)$.
- Every other power is lower than this because of the following.
 - Let S_1 be the state with all A -splits.
 - No other state besides S_1 has all A -splits, so the closest is an S_2 with one B -split.
 - However, this implies that $a(S_2) - b(S_2) + 2(|S_2| - 1) = (a(S_1) - 1) - (b(S_1) + 1) + 2(|S_2| - 1)$.
 - Since changing one A -split can only change $|S|$ by ± 1 (which would either counteract the change to $a(S)$ and $b(S)$ or double it), no power can be higher.
 - Note that another state could only be equal by this logic.
 - By *Exercise 6.10*^[29], “the highest power of A that occurs in an all- A -split state is strictly greater than the highest power of A that occurs in a state with exactly one B -split” (161).
- By a similar argument, the lowest power of A is $-n - 2(W - 1)$, where W is the number of white regions in the original projection.
- Thus, the following holds (see *Exercise 6.11* for the transition from step 2 to 3).

$$\begin{aligned}\text{span}(\langle K \rangle) &= n + 2(D - 1) - (-n - 2(W - 1)) \\ &= 2n + 2(D + W - 2) \\ &= 2n + 2n \\ &= 4n\end{aligned}$$

– Q.E.D.

- *Exercise 6.11:* Show that the number of regions R in a connected knot or link projection is always two more than the number of crossings (including the region outside the knot). Hint: Either use the fact that the Euler characteristic of a disk is always 1 or simply draw a knot, keeping count of the number of regions created whenever you create a new crossing^[30].
 - The unknot divides space into two regions with 0 crossings.
 - If it is twisted to generate 1 crossing and a edge that resembles the infinity symbol, space is now divided into three regions.
 - Subsequent crossings behave in a similar way.
- *Exercise 6.12:* As a simple corollary^[31] to the lemma, show that if K has a reduced alternating projection of n crossings, then the span of its Jones polynomial is exactly n . Note that this implies that $c(K)$ for an alternating knot is exactly the span of its Jones polynomial.
 - Since $\text{span}(\langle K \rangle) = 4n$ and the transition from the bracket polynomial to the Jones polynomial involves dividing each exponent by 4, $\text{span}(V(K)) = n$.
- Theorem: “Two reduced alternating projections of the same knot have the same number of crossings” (162).
 - If the first projection P_1 has n crossings, the lemma implies that $\text{span}(P_1) = 4n$.

²⁹Possibly within my grasp.

³⁰Could be done rigorously with more time.

³¹**Corollary:** A proposition that follows from (and is often appended to) one already proved.

- Since the span of the bracket polynomial is a knot invariant, the span of the second projection P_2 is also $4n$.
- Therefore, by the transitive property, both projections have n crossings.
- William Menasco proved in 1983 that if $K_1 \# K_2$ is alternating, then it appears composite in any alternating projection.
 - Appears composite: In any alternating projection, a circle that intersects the knot twice can fully enclose one alternating factor knot.
- Menasco's result implies the following corollary.
- Corollary: $c(K_1 \# K_2) = c(K_1) + c(K_2)$ for all $K_1 \# K_2$ alternating knots.
 - Choose a reduced alternating projection of $K_1 \# K_2$.
 - By Menasco's result, K_1 and K_2 are both parts of this projection.
 - Thus, both factor knots are in their minimal-crossing projections.
 - Therefore, $c(K_1 \# K_2)$ is a simple sum of $c(K)$ for the two factor knots.
- The above corollary holds true for **adequate knots**.
- Furthermore, “for an alternating graph with an even number of edges hitting each vertex, the minimal crossing number is realized in any reduced alternating projection such that every edge has a crossing” (163).
 - Think of one of these graphs as having a single vertex and edge for each component in the knot or link. The edge exits the vertex, loops around and knots in space, and returns to the vertex.
 - “More generally, any graph consisting of a set of vertices connected by edges can be embedded in space. The edges can twist, curve and knot in space. The vertices can have any number of edges coming out of them” (163).
- On the **Tait Flyping Conjecture**:

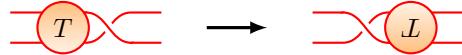


Figure 6.6: A flype.

- Conjectured by Peter Guthrie Tait (1831-1901) in the 1890s.
- If true, it would allow the generation of all possible reduced alternating projections of a given alternating knot.
- This process generates a finite number of projections.
- In 1990, William Menasco and Morwen Thistlethwaite together proved the Tait Flyping Conjecture using a blend of geometric techniques and new polynomials.
- **Tait Flyping Conjecture:** Two reduced alternating projections of the same knot are equivalent when projected on a sphere (see Section 2.2) if and only if they are related through a series of **flypes**.
- **Flype:** A 180° rotation of a tangle (see Figure 6.6).

6.3 The Alexander and HOMFLY Polynomials

- The **Alexander polynomial** Δ was invented in 1928 and was the first polynomial to be invented.
- It is a polynomial for oriented links in t .
- J. W. Alexander initially defined it in rather abstract mathematical concepts. In fact, it wasn't until 1969 that John Conway showed that the Alexander polynomial could be computed using just two rules.
- First is the usual rule, except that this holds true for any projection of the trivial knot instead of just the zero-crossing one.

$$\text{Rule 1 : } \Delta(\textcircled{O}) = 1$$

- Consider Figure 6.3 to understand the second rule.

$$\text{Rule 2 : } \Delta(L_+) - \Delta(L_-) + \left(t^{\frac{1}{2}} - t^{-\frac{1}{2}}\right) \Delta(L_0) = 0$$

- It will not be proven herein, but these two rules suffice to ensure that the Alexander polynomial is an invariant for knots and links.
- *Exercise 6.15:* The Alexander polynomial of a splittable link is always 0. Hint: Picture the splittable link as L_0 .
- Let's calculate the Alexander polynomial for the trefoil knot.
 - Orient the trefoil and consider the upper-left crossing in the leftmost term of Equation 6.11 to be L_+ . Employ Rule 2.

$$\Delta(\textcircled{\text{K}}) - \Delta(\textcircled{\text{K}}) + \left(t^{\frac{1}{2}} - t^{-\frac{1}{2}}\right) \Delta(\textcircled{\text{T}}) = 0 \quad (6.11)$$

- Employ Rule 1 on the second term of Equation 6.11.

$$\Delta(\textcircled{\text{T}}) = \Delta(\textcircled{O}) = 1$$

- Employ Rule 2 on the third term of Equation 6.11, considering the upper crossing in the leftmost term of Equation 6.12 to be L_+ .

$$\Delta(\textcircled{\text{T}}) - \Delta(\textcircled{\text{T}}) + \left(t^{\frac{1}{2}} - t^{-\frac{1}{2}}\right) \Delta(\textcircled{O}) = 0 \quad (6.12)$$

- By *Exercise 6.15*, $\Delta(\textcircled{\text{T}}) = 0$, so the following holds.

$$\Delta(\textcircled{\text{T}}) = -t^{\frac{1}{2}} + t^{-\frac{1}{2}}$$

- Therefore, the following holds, giving Equation 6.13 as the final Alexander polynomial for the trefoil knot.

$$\begin{aligned} \Delta(\textcircled{\text{T}}) &= \left(t^{\frac{1}{2}} - t^{-\frac{1}{2}}\right)^2 + 1 \\ &= t - 1 + t^{-1} \end{aligned} \quad (6.13)$$

- “Unlike the Jones polynomial, there are known examples of nontrivial knots with Alexander polynomial equal to 1” (166).
 - One such example is the $(-3, 5, 7)$ -pretzel knot (see Figure 6.7).
- Like the Jones polynomial, the Alexander polynomial can always be calculated.

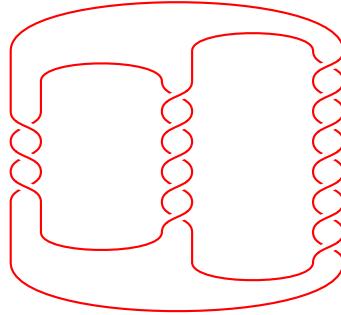


Figure 6.7: The $(-3, 5, 7)$ -pretzel knot.

- Choose the crossing that is modified by the skein relation to be one of the crossings that would be changed to unknot the knot in question (see Section 3.1).
- Rule 2 takes the crossing as either L_+ or L_- , so Rule 2 produces a knot with one less crossing and a knot that is one crossing closer to being the unknot.
- Iteratively, this process will once again produce several series of trivial links.
- Displayed visually, this process forms a **resolving tree**.
- **Resolving tree:** A tree of links formed through the “process of repeatedly choosing a crossing, and then applying the skein relation to obtain two simpler links” (167).

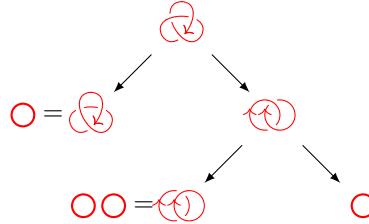


Figure 6.8: The resolving tree for the trefoil knot.

- **Depth** (of a resolving tree): “The number of levels of links in the tree, not including the initial level at the top” (168).
 - For example, the resolving tree in Figure 6.8 has depth 2.
- **Depth** (of a link L): The minimal depth for any resolving tree of L .
 - A knot invariant.
 - Measures the complexity of the calculation of the Alexander polynomial.
 - Links of depth 0 are trivial links.
 - A knot of depth 1 is trivial.
 - A link of depth 1 is the Hopf link, possibly with some extra disentangled trivial components.
 - The links of depth 2 have been classified.
- On the other hand, the **HOMFLY polynomial** P is a two-variable Laurent polynomial in m and l satisfying the following two rules.

$$\text{Rule 1 : } P(\textcircled{O}) = 1$$

$$\text{Rule 2 : } lP(L_+) + l^{-1}P(L_-) + mP(L_0) = 0$$

- Like the Alexander polynomial, Rule 1 holds for any projection of the trivial knot.
- Note that once again, Rule 2 is very similar to that of the Alexander polynomial and of the Jones polynomial's skein relation. Refer to Figure 6.3.
- Find $P(L)$ for the trivial link of two components (Figure 6.9c).

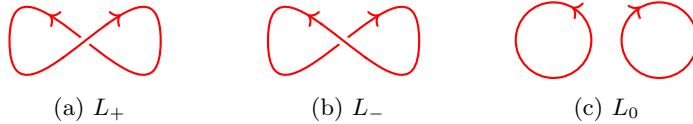


Figure 6.9: Finding the HOMFLY polynomial of the trivial link of two components.

- By Figure 6.9 and Rule 2, $lP(L_+) + l^{-1}P(L_-) + mP(L_0) = 0$.
- Because Figures 6.9a and 6.9b are representations of the unknot, $P(L_+) = P(L_-) = 1$.
- Thus, $mP(L_0) = -(l + l^{-1})$.
- Therefore, the HOMFLY polynomial $P(L_0)$ for the trivial link of two components is given by the following. Note that Equation 6.14 is an acceptable way to leave the HOMFLY polynomial.

$$P(L_0) = -m^{-1}(l + l^{-1}) \quad (6.14)$$

- *Exercise 6.19:* Show that the HOMFLY polynomial of a knot is identical to the HOMFLY polynomial of the same knot but with the opposite orientation.
 - Consider how the HOMFLY polynomial is constructed (Figure 6.3).
 - Take each of the three projections in Figure 6.3 and reverse the orientations.
 - The orientations on both strands must change because the knot is continuous.
 - Immediately, this may look like something has changed. However, if these new projections are rotated 180° about an axis through the page, they will be identical to the original pictures.
 - Thus, since reversing the orientation does not change which crossings correspond to L_+ , L_- , and L_0 , the new HOMFLY polynomial is calculated the same way, and is therefore identical to the original.
 - With links, however, changing some of the orientations can have an effect.
 - The HOMFLY polynomial *can* always be computed, but the computations can take a long time.
 - A 100 crossing knot would resolve into 2^{100} unlinks.
 - However, Hugh Morton and Hamish Short created an algorithm applicable to closed braids of nine or fewer strings that can successfully compute the polynomial of knots of several hundred crossings.
 - **Split union** (of two links): A link $L_1 \cup L_2$ where L_1 and L_2 can be separated by a sphere.
 - The HOMFLY polynomial of a split union is given by Equation 6.15^[32].

$$P(L_1 \cup L_2) = -m^{-1}(l + l^{-1})P(L_1)P(L_2) \quad (6.15)$$

- This does lead to second interesting property (Equation 6.16).

$$P(L_1 \# L_2) = P(L_1)P(L_2) \quad (6.16)$$

³²The proof of this is uses the rules for calculating the HOMFLY polynomial. However, it does get messy.

- Note that all distinct composite links with the same factor knots will have the same HOMFLY polynomial.
- Prove Equation 6.16 using Equation 6.15.

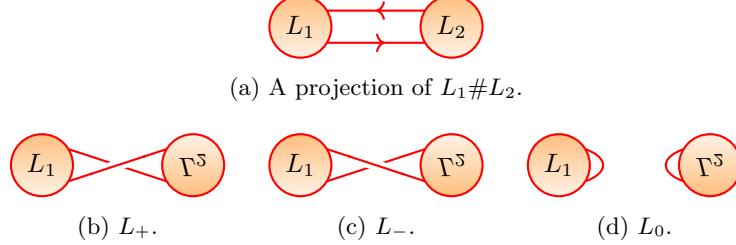


Figure 6.10: Proving the formula for the HOMFLY polynomial of a composite knot.

- Figure 6.10a shows any composite link, and Figures 6.10b, 6.10c, and 6.10d show the connecting strands twisted in such a way as to match Figure 6.3.
- Equation 6.15 governs Figure 6.10d since Figure 6.10d is a split union.
- Rule 2 connects Figures 6.10b, 6.10c, and 6.10d via the following.

$$lP(L_1 \# L_2) + l^{-1}P(L_1 \# L_2) + mP(L_1 \cup L_2) = 0$$

- Substitute in Equation 6.15.

$$lP(L_1 \# L_2) + l^{-1}P(L_1 \# L_2) + m(-m^{-1}(l + l^{-1})P(L_1)P(L_2)) = 0$$

- Simplify and solve for $P(L_1 \# L_2)$.

$$\begin{aligned} (l + l^{-1})P(L_1 \# L_2) - (l + l^{-1})P(L_1)P(L_2) &= 0 \\ (l + l^{-1})P(L_1 \# L_2) &= (l + l^{-1})P(L_1)P(L_2) \\ P(L_1 \# L_2) &= P(L_1)P(L_2) \end{aligned}$$

- The property described by Equation 6.16 continues the analogy of composition to multiplication introduced in Section 1.2.
- The HOMFLY polynomial is not a **complete invariant** for knots.
 - The HOMFLY polynomial cannot distinguish between mutant knots.
- **Complete invariant:** A characteristic that distinguishes all entities in a set.
- The HOMFLY polynomial is better at telling apart knots and links than either the Jones or Alexander polynomials because these other polynomials are simply a special case of the HOMFLY polynomial.
 - To convert a HOMFLY polynomial into a Jones polynomial, make the following substitutions where $i = \sqrt{-1}$.

$$\begin{aligned} l &= it^{-1} \\ m &= i(t^{-\frac{1}{2}} - t^{\frac{1}{2}}) \end{aligned}$$
 - To convert a HOMFLY polynomial into an Alexander polynomial, make the following substitutions where $i = \sqrt{-1}$ (*Exercise 6.21*).

$$\begin{aligned} l &= i \\ m &= i(t^{\frac{1}{2}} - t^{-\frac{1}{2}}) \end{aligned}$$

- “Let E be the largest exponent of l in the HOMFLY polynomial of the oriented link L and e be the smallest exponent of l ” (174). In this case, the braid index n of the link L satisfies the inequality in Equation 6.17.

$$n \geq \frac{E - e}{2} + 1 \quad (6.17)$$

- This inequality was proven by Robert Williams, John Franks, and (independently) Hugh Morton.
- It is **sharp** for “all but five prime knots of 10 or fewer crossings (the exceptions being 9_{42} , 9_{49} , 10_{132} , 10_{150} , and 10_{156})” (174).
- **Sharp** (inequality): An inequality where both sides are equal.
- Several other polynomials exist, but they will not be explored herein^[33].

6.4 Amphicheirality

- See Section 1.3 for a refresher on amphicheirality.
- Let K^* be the mirror image of K .
- *Exercise 6.22:* Show that the bracket polynomial of K^* is just the bracket polynomial of K where the variable A is replaced everywhere by the variable A^{-1} . Show that the same is true for the X polynomial.
 - If every crossing is flipped from K to K^* , then the version of Rule 2 used to split every crossing switches as well.
 - Thus, every split assigned A is now assigned A^{-1} and vice versa.
 - Thus, the coefficient before $(-A^2 - A^{-2})$ that was A^n becomes A^{-n} .
 - Furthermore, every unlink is eventually resolved with $(-A^2 - A^{-2})$.
 - So if we were multiplying A^n through the above to yield $-A^{n+2} - A^{n-2}$, and now we’re multiplying A^{-n} , we get $-A^{-(n-2)} - A^{-(n+2)}$, which is identical to replacing all of the A s in the first expression with A^{-1} .
 - Considering the X polynomial as Equation 6.7 relates it to the bracket polynomial, flipping every crossing has the effect of multiplying the former writhe by -1 .
 - However, this is offset by the substitution of $-A^3$ for $-A^{-3}$ since a power to a power multiplies the exponents and $-1 \times -1 = 1$.
- If K is amphicheiral, then K^* is simply a different projection of K . Thus, $X_K(A) = X_{K^*}(A)$ ^[34].
- By *Exercise 6.22*, $X_K(A) = X_{K^*}(A^{-1})$.
- Therefore, if K is amphicheiral, then Equation 6.18 must hold.

$$X_K(A) = X_{K^*}(A^{-1}) = X_K(A^{-1}) \quad (6.18)$$

- To satisfy Equation 6.18, the X polynomial must be **palindromic**.
- **Palindromic** (polynomial): A polynomial where the absolute values of the ordered coefficients read the same backwards and forwards. Replacing every variable by the inverse of that variable gives the original polynomial.
- Section 1.3 showed that the figure-eight knot was amphicheiral. This can be further shown by the fact that the X polynomial of 4_1 is $A^8 + A^4 - 3 + A^{-4} + A^{-8}$, which is clearly palindromic.
- Similarly, the fact that the X polynomial of the trefoil knot is not palindromic likely means that the trefoil knot is not amphicheiral.

³³Idea for paper: Find a polynomial as a function of a knot’s Dowker notation.

³⁴Let the notation $X_K(A)$ denote the X polynomial of K in the variable A .



(a) The left-hand trefoil.



(b) The right-hand trefoil.

Figure 6.11: The two versions of the trefoil knot.

- The two variations of the trefoil knot have the common sense names highlighted by Figure 6.11.
- There exist some implications of this for alternating knots.
 - It was previously shown that $\max \deg \langle K \rangle = n + 2(D - 1)$ and $\min \deg \langle K \rangle = -n - 2(W - 1)$.
 - This can be adapted for the X polynomial. Namely, through $\max \deg X(K) = n + 2(D - 1) - 3w(K)$ and $\min \deg X(K) = -n - 2(W - 1) - 3w(K)$.
 - Since $X_K(A) = X_K(A^{-1})$ for an amphicheiral knot (an amphicheiral knot is palindromic), then $\max \deg X(K) = -\min \deg X(K)$.
 - Thus, $n + 2(D - 1) - 3w(K) = -(-n - 2(W - 1) - 3w(K))$.
 - Therefore, $3w(K) = D - W$ for any reduced alternating projection of an amphicheiral knot.
- For alternating knots, if $\frac{c(K)}{2} \notin \mathbb{N}$, then the knot is not amphicheiral.
 - This conjecture does not hold true for nonalternating knots or links.
- Equation 6.18 is also applicable to the HOMFLY polynomial (since the HOMFLY carries within it the information of the X polynomial). However, it is not infallible at determining amphicheirality, e.g., in the case of 9_{42} .
 - In other words, if a knot is amphicheiral, Equation 6.18 holds true.
 - However, the converse of the above statement is false.

7 Biology, Chemistry, and Physics

7.1 DNA

- Begins with a general, elementary introduction to DNA.
- **Topo-isomerase:** An enzyme that manipulates DNA topologically.
 - Used to expose various sequences for replication, transcription, and recombination.
- Topo-isomerases perform actions that range from simple to very sophisticated.
- Some such enzymes take two strands of DNA and replace them with a nontrivial tangle.
 - If DNA is linear, a knot, once formed, could easily “slip off” one end or the other.
 - Thus, scientists allow specific topo-isomerases to act on circular DNA, because if the enzyme in question causes knotting, the knot will be captured in the circle.
- Recently, scientists have learned to artificially create cyclic DNA. It is this DNA that is subjected to various topo-isomerases.
- Focus on duplex cyclic DNA.

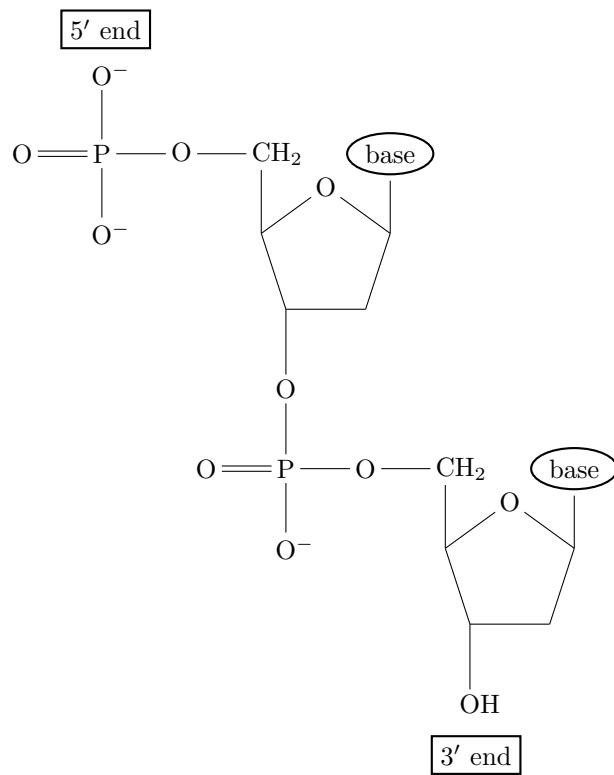


Figure 7.1: A segment of a strand of DNA.

- Each phosphate is bonded to two different sugars. This forms the backbone.
- These sugars are, in turn, bonded to cytosine, thymine, guanine, or adenine bases.
- A single phosphate binds to the 3' site of one sugar and the 5' site of another.
- This means that one end of the molecule will have an open 5' site while the other will have an open 3' site.

- Thus, an orientation is created based on the convention of moving from 5' to 3'.
- However, because linear duplex DNA has an open 3' and 5' end on both sides, it cannot be oriented (only linear single-stranded DNA can be oriented).
- Thus, when the ends of a linear strand of DNA are joined, each end must join its own end, giving two linked strands in space as opposed to one strand that loops around twice.
 - In other words, there must be “an even number of half-twists in the cyclic duplex DNA” (184).
- This cyclic duplex can be modeled as a ribbon in space, whose edges are bounded by the cyclic duplex (because of the above, it is never a Möbius band; is orientable).
- **Axis:** The curve that runs along the center of the ribbon.
 - The axis does not correspond to a specific part of the molecule, but if it and the edges are all similarly oriented, it conveys the extent of the molecule’s contortion in space.
- **Twist:** A measure of how much the ribbon twists around its axis. *Also known as $\text{Tw}(R)$.*
 - The definition listed below is only applicable when the ribbon is static/rigid and lies within the projection plane.
 - Take one of the two link components bounding the ribbon (it doesn’t matter which one) and the axis.
 - Note +1 or -1 for every crossing (of the link component and axis) according to Figure 1.18.
 - Sum these values and divide by 2.
 - If we wish to calculate $\text{Tw}(R)$ for a malleable ribbon in space, we must sum “the incremental twist of the ribbon about the axis, integrated as we traverse the axis once” (185).
 - This value need not be an integer.
- **Writhe:** A measure of how contorted the axis of the ribbon is in space. *Also known as $\text{Wr}(R)$.*
 - Take the average of the **signed crossover number** over every possible projection of the axis, i.e., “the planar pictures that would result as we looked at the fixed axis from all possible vantage points on a [unit] sphere that surrounds it” (185).
 - Determined using integrals ($\int dA = 4\pi$ since 4π is the surface area of the unit sphere):

$$\begin{aligned} \text{Wr}(R) &= \frac{\int \text{signed crossover number } dA}{\int dA} \\ &= \frac{\int \text{signed crossover number } dA}{4\pi} \end{aligned}$$
 - The writhe also need not be an integer.
- **Signed crossover number:** For any particular projection of the axis, the sum of all the ± 1 s occurring at crossings where the axis crosses itself assigned according to Figure 1.18.
- If the ribbon axis lies in a plane, $\text{Wr}(R) = 0$ because the only vantage points with a nonzero number of crossings (technically undefined crossings) is the equator of the sphere, which makes up 0% of the sphere’s area.
- Linking number: See Section 1.4. *Also known as $\text{Lk}(R)$.*
 - This characteristic does not depend on the specific projection, unlike twist and writhe.
- James White, Brock Fuller, and G. Călugăreanu all independently discovered Equation 7.1.

$$\text{Lk}(R) = \text{Tw}(R) + \text{Wr}(R) \quad (7.1)$$