

The Knot Book

Notes

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1 Introduction

1.1 Introduction

- **Knot:** “A knotted loop of string, except that we think of the string as having no thickness, its cross-section being a single point” (2).
- Do not distinguish between a ‘nice, even’ knot and one that has been deformed through space.
- **Unknot:** “The simplest knot of all... the unknotted circle” (2). *Also known as trivial knot.* See Figure 1.1a.
- **Trefoil knot:** “The next simplest knot” (2). See Figure 1.1b.

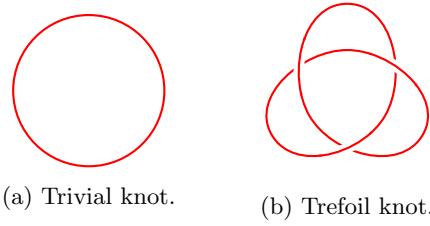


Figure 1.1: Projections of the two simplest knots.

- **Projection:** A picture of a knot, such as those in Figure 1.1.
 - The same knots can have multiple projections (as they are deformed in space).
- **Crossings:** The places in a projection where a knot crosses itself.
 - The trefoil knot in Figure 1.1b is a three-crossing knot because it crosses itself 3 times.
 - Any one-crossing knot is trivial.
 - *Exercise 1.2:* Any two-crossing knot must be trivial because the simplest nontrivial knot is the trefoil knot, which has three crossings.
- Atoms were originally thought to be tangles (knots) in the ether of the universe, but when chemists moved on, mathematicians took up knot theory. In the 1980s, biochemists began to see applications of knot theory in their research (see Section ??).
- **Topology:** “The study of the properties of geometric objects that are preserved under deformations” (6).
 - Knot theory is a subfield of topology (see Section ??).

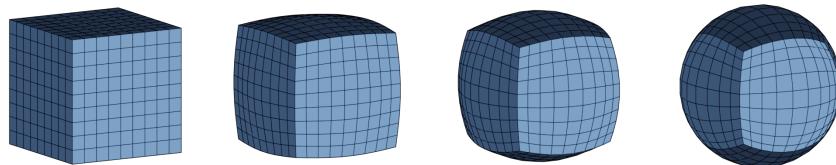


Figure 1.2: Deformation of a cube into a sphere.

- Any knot can have a projection with as many crossings as desired.

- **Alternating knot:** “A knot with a projection that has crossings that alternate between over and under as one travels around the knot in a fixed direction” (7).
 - The trefoil is such a knot.
- *Exercise 1.7**: By changing some of the crossings from over to under or vice versa, any projection of a knot can be made into a projection of the unknot^[1]. See Figure 1.3.

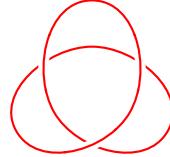


Figure 1.3: A projection of the unknot evoking the trefoil knot.

1.2 Composition of Knots

- **Composition** (of two knots): “A new knot obtained by removing a small arc from two knot projections and then connecting the four endpoints by two new arcs” (7).
 - If two knots are designated J and K , then their composition is denoted $J \# K$.
 - Do not overlap the projections and choose two arcs that are on the outside to avoid new crossings.
 - Make sure that the new arcs do not cross any of the the original knot projections or each other.
- **Composite knot:** A knot that “can be expressed as the composition of two knots, neither of which is the trivial knot” (8).
 - This definition is analogous to composite integers, where an integer is composite if it is the product of positive integers, neither of which is 1.
 - Similarly, if we compose any knot with the unknot, we get the same knot back.
- **Factor knots:** “The knots that make up the composite knot” (8).
- **Prime knot:** “A knot [that] is not the composition of any two nontrivial knots” (9).
- The unknot, trefoil knot, and figure-eight knots are all prime (see Section 4.3).
 - The unknot is not composite for the same reason that 1 is not the product of two integers greater than 1.

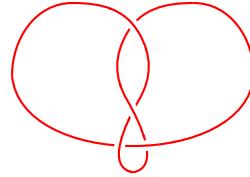


Figure 1.4: The figure-eight knot.

- Similar to integers, “a composite knot factors into a unique set of prime knots” (10).
- *Exercise 1.8:* Using the appendix table, identify the factor knots that make up the composite knot in Figure 1.5.
 - The knot in Figure 1.5 is the composition of two projections of 5_2 .

¹How can I *show* something? How can I do these proofs? What kind of logic solves one of these? See Section 3.1 for a direct proof of/solution to Exercise 1.7.

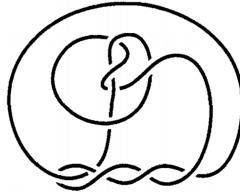
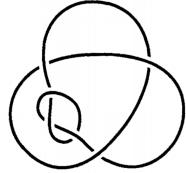


Figure 1.5: The composite knot.

- *Exercise 1.9:* Show that the knot in Figure 1.6a is composite.



(a) The composite knot.



(b) Factors.

Figure 1.6: Factorization of a ‘double trefoil.’

- There is more than one way to take the composition of two knots (by removing different arcs).
 - This is not analogous to multiplication — a break in the pattern.
- **Orientation:** A direction to travel around the knot. Denoted by placing “coherently directed arrows along the projection of the knot in the direction of our choice” (10). A knot with such arrows is **oriented**.
 - All compositions $J \# K$ where the orientations of J and K do match up will yield the same composite knot.
 - J can be ‘slid around’ $J \# K$ until it reaches the second position where the composition was taken.
 - All compositions $J \# K$ where the orientations of J and K do not match up will yield the same composite knot.
 - These two compositions can be distinct.

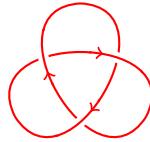


Figure 1.7: Orientation notation.

- **Invertible:** A knot that can be deformed back to itself so that an orientation on it is sent to the opposite orientation.
 - “In the case that one of the two knots is invertible, say J , we can always deform the composite knot so that the orientation on K is reversed, and hence so that the orientations of J and K always match. Therefore, there is only one composite knot that we can construct from the two knots” (11).
- To determine the possible compositions of knots, it is necessary to know which knots are invertible, but no general technique has yet been discovered.

1.3 Reidemeister Moves

- **Ambient isotropy:** “The movement of the string through three-dimensional space without letting it pass through itself” (12).
- **Planar isotropy:** A deformation of “the projection plane as if it were made of rubber with the projection drawn upon it” (12).
 - Stretching, squeezing, rotating, bending single arcs, etc.
- **Reidemeister move:** “One of three ways to change a projection of the knot that *will* change the relation between the crossings” (13).
- **First Reidemeister move:** “Put in or take out a twist in the knot” (13). See Figure 1.8a. *Also known as type I Reidemeister move.*
- **Second Reidemeister move:** “Either add two crossings or remove two crossings” (13). See Figure 1.8b. *Also known as type II Reidemeister move.*
- **Third Reidemeister move:** “Slide a strand of the knot from one side of a crossing to the other side of the crossing” (13). See Figure 1.8c. *Also known as type III Reidemeister move.*
 - Note that the crossings in Figure 1.8 can be reversed and the move will still be classified under the same category.

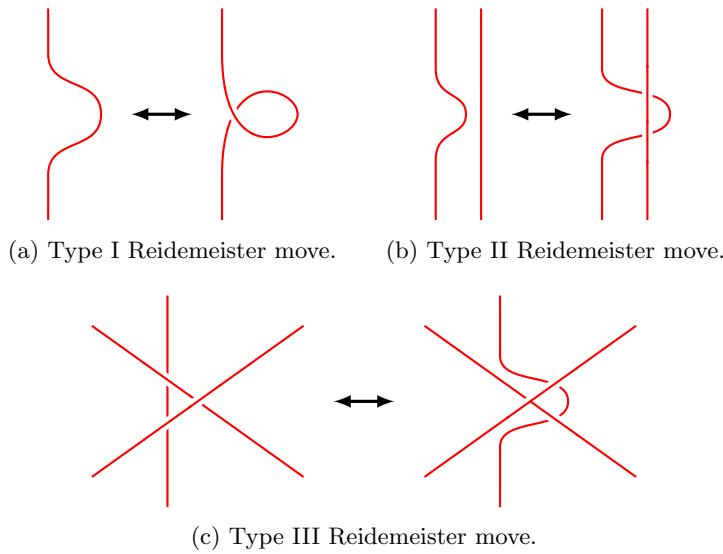
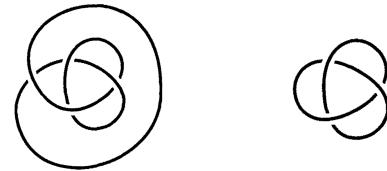


Figure 1.8: Reidemeister moves.

- All Reidemeister moves are ambient isotopies.
- “If we have two distinct projections of the same knot, we can get from the one projection to the other by a series of Reidemeister moves and planar isotopies” (14).
- **Amphicheiral:** A knot that “is equivalent to its mirror image, that is, the knot obtained by changing every crossing... to the opposite crossing” (14-15). *Also known as achiral by chemists.*
 - A knot and its mirror image are distinct unless the knot is amphicheiral.
 - See Section ?? for more on amphicheirality.

- *Exercise 1.10:* Show that the two projections in Figure 1.9 represent the same knot by finding a series of Reidemeister moves from one to the other.



(a) Initial projection. (b) Final projection.

Figure 1.9: Finding Reidemeister moves.



Figure 1.10: Solution to *Exercise 1.10*.

- *Exercise 1.11*:* Find a sequence of Reidemeister moves to untangle the unknot shown in Figure 1.11.

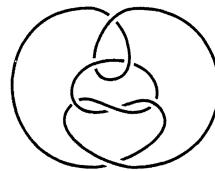


Figure 1.11: Unknot to be untangled.

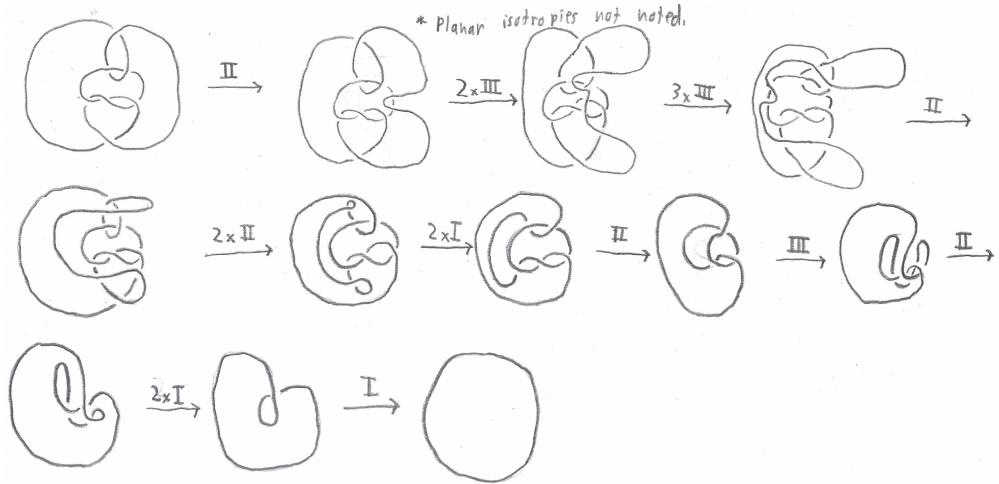
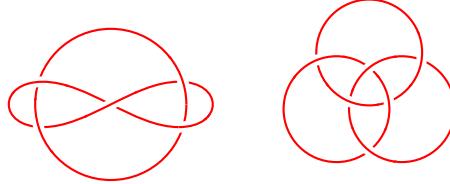


Figure 1.12: Solution to *Exercise 1.11**.

- The bounds on the increase in crossings generated by Reidemeister moves from one projection to another are unknown.

1.4 Links

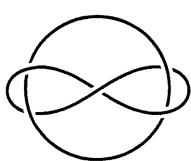
- **Link:** “A set of knotted loops all tangled up together” (17).
- “Two links are considered to be the same if we can deform the one link to the other link without ever having any one of the loops intersect itself or any of the other loops in the process” (17).



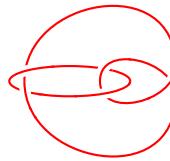
(a) Whitehead link. (b) Borromean rings.

Figure 1.13: Projections of two common links.

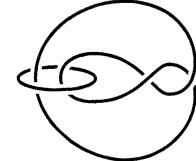
- *Exercise 1.13:* Show that the two projections in Figure 1.14a and 1.14c represent the same link.
 - Untwist the right side of the figure-eight loop, twisting the right side of the circle at the same time. Then add a further twist to the circle loop and flip (rotate along the horizontal axis) the figure-eight loop (which now looks like an ellipse).



(a) Initial projection.



(b) Intermediate projection.



(c) Final projection.

Figure 1.14: Ambient isotopies of the Whitehead link.

- **Link of n components:** A link made up of n loops knotted with each other.
 - The Whitehead link (Figure 1.13a) is a link of 2 components.
 - The Borromean rings (Figure 1.13b) are a link of 3 components.
 - A knot is a link of 1 component.
 - If the number of components in two links differ, then the links are clearly distinct.
- **Splittable:** A link whose components “can be deformed so that they lie on different sides of a plane in three-space” (17).
- *Exercise 1.14:* Show that the link in Figure 1.15 is splittable.

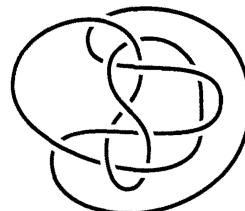


Figure 1.15: Link to be split.

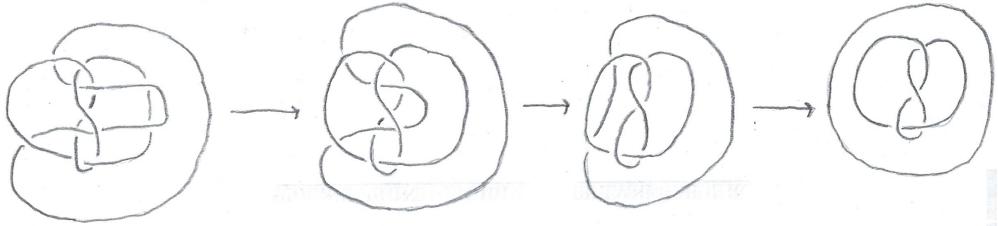


Figure 1.16: Solution to *Exercise 1.14*.

- **Unlink:** One of the two simplest links of two components and the simplest splittable link of two components. *Also known as trivial link.* See Figure 1.17a.
- **Hopf link** The other of the two simplest links of two components and the simplest nonsplittable link of two components. See Figure 1.17b.

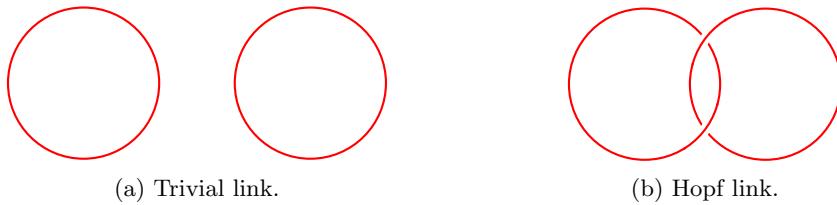


Figure 1.17: Projections of the two simplest links of two components.

- **Linking number:** A quantity that numerically measures how linked up two components are.

- If M and N are two components in a link, begin by orienting both of them.
- At each crossing, count a $+1$ if Figure 1.18a holds or count a -1 if Figure 1.18b holds.
 - Note that if you are unsure, do the following: If rotating the bottom strand clockwise lines up the arrows (correlates the two orientations), count $+1$.
 - In the same vein, if rotating the bottom strand counterclockwise lines up the arrows, count -1 .
 - If a component crosses itself, do not count it.

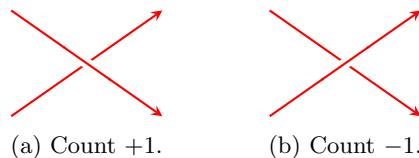


Figure 1.18: Computing linking numbers per crossing.

- Sum all of the $+1$ s and -1 s and divide this sum by 2 to yield the linking number.
 - For example, the linking number for the Hopf link (Figure 1.17b) is ± 1 depending on the chosen orientation.
 - Note that reversing the orientation for one of the two links is equivalent to multiplying the linking number by -1 .
 - As such, the absolute value of the linking number remains constant whatever orientation is chosen.
- *Exercise 1.15:* Compute the linking number of the link pictured in Figure 1.19. Now reverse the direction on one of the components and recompute it.

- Add up all of the numbers in the second row of Table 1.1 to yield 2. Divide by 2 to yield the linking number, 1.

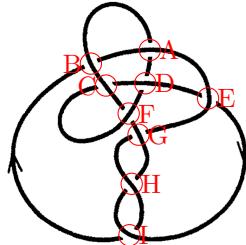


Figure 1.19: Link of linking number n .

A	B	C	D	E	F	G	H	I
-1	-1	0	0	+1	0	+1	+1	+1

Table 1.1: Counting crossings.

- Although a linking number is computed using a single projection of a link, the linking number will be the same for any projection of the link.
 - This can be proven by demonstrating that the Reidemeister moves do not change the linking number. “Since we can get from any one projection of a link to any other via a sequence of Reidemeister moves, none of which will change the linking number, it must be that two different projections of the same link yield the same linking number” (20).
 - Let’s take this case by case.
 - A Type I Reidemeister move generates a self-intersection. Because self-crossings do not count toward the linking number by definition, the first Reidemeister move does not affect the linking number.
 - A Type II Reidemeister move generates two crossings. There are now two cases: Either the new crossings are self-intersections or they are not. If the new crossings are self-intersections, then there is no change to the linking number. If the new crossings are not self-intersections, they will always have opposite linking numbers, and, thus, the two new crossings cancel each other out. Therefore, the second Reidemeister move does not affect the linking number.
 - A Type III Reidemeister move removes two crossings and generates two crossings. There are now two cases: Either the new crossings are self-intersections or they are not. If the new crossings are self-intersections, then there is no change to the linking number. If the new crossings are not self-intersections, both pairs of crossings will always have opposite linking numbers and, thus, both pairs of crossings cancel each other out. Therefore, the third Reidemeister move does not affect the linking number.
 - This logic can be visually confirmed by assigning orientations and counting crossings in Figure 1.8.
- **Invariant:** A quality of a knot or link that, once orientations are chosen, is unchanged by ambient isotopy.
 - Both the linking number and the number of components are link invariants.
- *Exercise 1.16:* Explain why the linking number of a splittable two-component link will always be 0, no matter what projection is used to compute it.

- **SOLUTION 1:** By definition, the linking number numerically measures how linked up two components are. Since splittable links are not joined (or “linked up”) in any way, the linking number must be 0.
- **SOLUTION 2:** If a two-component link is splittable, then there exists a projection of it with no crossings that are not self-intersections (a projection of the split link, as in Figure 1.17b). If all crossings are self-intersections, then the linking number computed for this projection must be 0. Now that it is known that said link has a linking number of 0, any combination of Reidemeister moves can be used to manipulate the link into any other projection. But because Reidemeister moves do not affect the linking number, the linking number will remain 0.
- The absolute value of the linking number, as a link invariant, can be used to distinguish certain distinct links, regardless of orientation.
 - “Any two links with two components that have distinct absolute values of their linking numbers have to be different links” (21).
 - For example, the difference in linking number between the unlink (0; Figure 1.17a) and the Hopf link (1; Figure 1.17b) distinguishes them.
- However, the linking number cannot distinguish between all links.
 - For instance, both the Whitehead link (Figure 1.13a) and the unlink (Figure 1.17a) have linking number 0.
 - One such distinction is discussed in Section 1.5.
- **Brunnian link:** A nontrivial link where the removal of any one component leaves behind a set of trivial unlinked circles.
 - The Borromean rings (Figure 1.13b) are Brunnian^[2].

1.5 Tricolorability

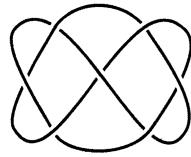
- How do we *prove* that every knot is not just a projection of the unknot?
- **Strand:** “A piece of the link that goes from one undercrossing to another with only overcrossings in between” (23).
- **Tricolorable:** A projection of a knot or link in which each strand “can be colored one of three different colors, so that at each crossing, either three different colors come together or all the same color comes together” (23).



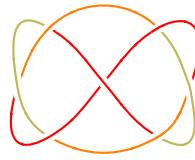
Figure 1.20: Tricolored trefoil.

- *Exercise 1.21:* Determine which of the projections of the three six-crossing knots 6_1 , 6_2 , and 6_3 are tricolorable.
 - 6_1 is tricolorable and the other two knots are not. I determined this through trial and error.
- *Exercise 1.22:* Show that the projection of the knot 7_4 in Figure 1.21a is tricolorable.

²How can I draw a Brunnian link with four components? Or more?



(a) Initial projection.



(b) Tricolored projection.

Figure 1.21: Tricolored 7_4 .

- Reidemeister moves do not affect tricolorability.
 - A Type I Reidemeister move generates a self-intersection. At such a crossing (made out of one original strand), every color will be the same; only one color meets at the crossing.
 - A Type II Reidemeister move generates two crossings. If the strands are the same color, then everything stays the same color. If the strands are different colors, than newly created loop takes on the third color.
 - Type III has many cases (*Exercise 1.23*).
- Since the unknot is not tricolorable and the trefoil knot is, we have just proven that there is at least one other knot besides the unknot.
 - Tricolorability is a knot invariant.
- For links, tricolorability is a bit different — the unlink, for example, is tricolorable.

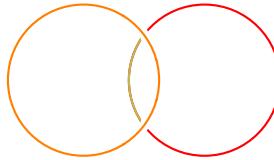


Figure 1.22: Tricolored unlink.

1.6 Knots and Sticks

- We can also consider knots made out of straight sticks glued to each other at each end.
- The first nontrivial knot made out of sticks is the trefoil, with six sticks.
 - We know that the stick trefoil in Figure 1.23 could be made in the real world because, if two vertices lie in the xy -plane, then two lie above and two lie below.
 - This P(lanar), L(ow), H(igh) notation is commonly used.

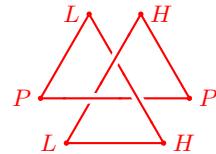


Figure 1.23: A trefoil knot made of sticks.

- **Stick number:** “The least number of straight sticks necessary to make a knot K ” (29). *Also known as $s(K)$.*

- The stick number for the composition of n trefoil knots is $2n + 4$.
- Many problems listed^[3].

³These could make good topics for a larger paper. Alternatively, some I am not ready to solve and need to learn more math first (those that I will come back to at the end of the book). I do not believe I'm missing anything essential by skipping these for now, but we'll see.

2 Tabulating Knots

2.1 History of Knot Tabulation

- Although many mathematicians (including Gauss) had dabbled with knots, Lord Kelvin's theory that atoms were knotted vortices in the ether (referenced in Section 1.1) kicked off research in earnest.
- Reverend Thomas P. Kirkman began tabulating knots but wrote poorly.
- However, Kirkman's ideas were applied by Scottish physicist Peter Guthrie Tait to list all alternating knots up to 10 crossings.
- C. N. Little of the University of Nebraska was the first to tackle the nonalternating knots.
 - He tabulated up to 10 crossings (43 knots total).
 - However, in 1974, it was discovered that two knots were, in fact, the same and there were really 42 knots.
- Kenneth A. Perko, a part time mathematician and New York lawyer, discovered that the projections were not distinct. As such, the projections are known as the Perko pair (Figure 2.1).

- *Exercise 2.1:* Show that the Perko pair (projected in Figure 2.1) are the same knot.

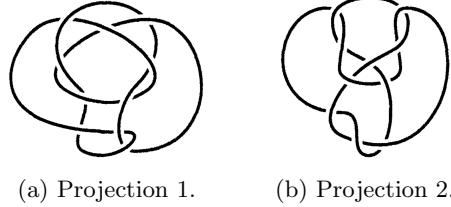


Figure 2.1: The Perko pair.

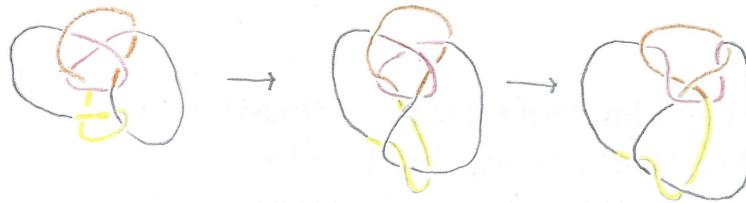


Figure 2.2: Solution to *Exercise 2.1*.

- Little later published a census of 11-crossing alternating knots with 11 omissions and 1 duplication.
- Mary G. Haseman listed all amphicheiral knots of 12 crossings.
- Early on, there were few attempts to *prove* that the tabulated knots were actually distinct.
- In 1927, Alexander and Briggs proved that the knots up to 9 crossings were distinct.
 - They developed the first knot polynomial, the Alexander polynomial.
 - This would remain the only knot polynomial until 1984.
- Once Kurt Reidemeister finished rigorously classifying the knots up to 9 crossings in 1932, tabulation was inactive for a while.

- In 1969, John H. Conway developed a new notation and used it to determine all of the prime knots up to 11 crossings and all prime nonsplittable links up to 10 crossings.
- In 1978, Alain Caudron of the University of Paris corrected Conway's tabulation a bit.
- Meanwhile, a new notation was developed by Hugh Dowker (see Section 2.2) that was implemented as a computer algorithm by Morwen Thistlethwaite.
- In the 1980s-1990s (with computer help), Jim Hoste, Thistlethwaite, and Jeff Weeks (an expert in hyperbolic knots; see Section ??) tabulated up to 16-crossing knots.
 - Note that the number of knots of successive numbers of crossings *appears* to grow exponentially (this is unproven).
 - Also note that Hoste et al's listed knots includes those that are not amphicheiral only once, so those knots actually represent 2 distinct knots.
 - Determining amphicheirality is discussed in Section ??
- On drawing a 14-crossing knot.
- Claus Ernst and Dewitt Sumners proved a lower exponential bound on the number of distinct prime knots (see Section 3.2).
- Dominic Welsh proved an upper exponential bound on the number of distinct prime knots.

2.2 The Dowker Notation for Knots

- For alternating knots...
 - Begin by choosing an orientation.
 - Label an arbitrary crossing 1.
 - Leaving that crossing on the understrand, label the next crossing 2.
 - Continue through the crossing on the same strand and label the next crossing 3.
 - Continue in this fashion until you have gone all the way around the knot once.
 - Note that this gives each crossing two numbers.

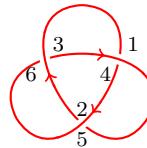


Figure 2.3: Trefoil knot labeled in Dowker notation.

- *Exercise 2.2:* Why does every crossing get one even numbered label and one odd numbered label?
 - Because the knot is alternating — every understrand will add an odd label and every overstrand will add an even number (because you set off on the understrand after labeling your starter crossing 1).

$$\begin{array}{ccc} 1 & 3 & 5 \\ 4 & 6 & 2 \end{array}$$

Table 2.1: Dowker numerical pairing of a trefoil knot.

- From the counting in Figure 2.3, we obtain a pairing between the numbers (Table 2.1).
 - Note that the numbers in the top row of Table 2.1 are constantly increasing odd numbers.
 - Since the top row numbers are in a predictable pattern, we can shorthand the notation for the trefoil knot in Figure 2.3 to 4 6 2.
 - “Thus, from a projection of a knot, we obtain a sequence of even integers, where the number of even integers is exactly the number of crossings in the knot” (36).
 - *Exercise 2.3:* Find a sequence of even integers that represents the projection of the knots 6_2 and 6_3 (Figure 2.4). How about a second sequence of even integers that also represents the same projection of 6_3 ?
 - For Figure 2.4a, write 6 8 10 12 2 4.
 - For Figure 2.4b, write 8 6 10 12 4 2.
 - Another sequence of even integers for Figure 2.4b is 4 10 8 12 2 6.

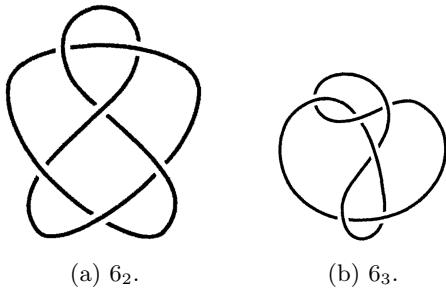


Figure 2.4: Knots 6_2 and 6_3 .

- Knots can also be reconstructed from Dowker notation. See Figure 2.5 for an example of reconstructing an alternating knot from the notation, 8 10 12 2 14 6 4.

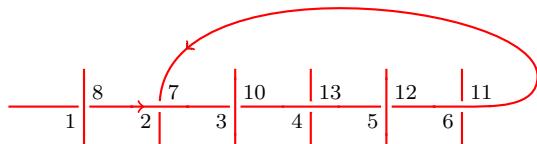


Figure 2.5: Constructing a knot projection from its Dowker notation.

- Our choice in drawing can change the resulting knot.
 - 4 6 2 10 12 8 gives two distinct knots.
 - “Note that the two knots are composite knots, and that this is reflected in the fact that the sequence 4 6 2 10 12 8 is actually a shuffling of the three numbers 2, 4, 6 and then a shuffling of the three numbers 8, 10, and 12” (38).
 - When the Dowker notation can be broken into two subpermutations (as above) the knot is composite (unless one subpermutation is trivial).
 - Any sequence that cannot be split in this way represents either a particular knot or its mirror image.
 - If the knot is amphicheiral, then the sequence only represent one knot.

- A knot and its mirror image are equivalent when projected onto a sphere as opposed to a plane.
- *Exercise 2.6:* Draw two projections given by 10 12 2 14 6 4 8, which are inequivalent as projections in the plane but which are equivalent as projections on the sphere.

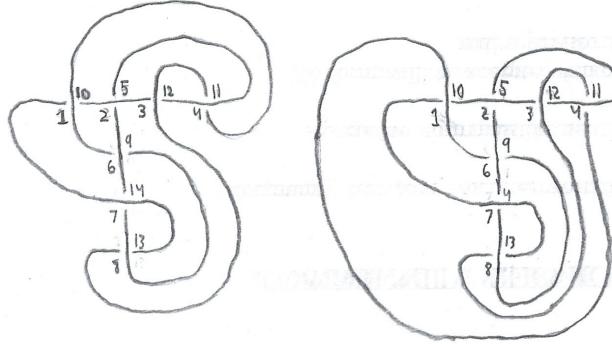
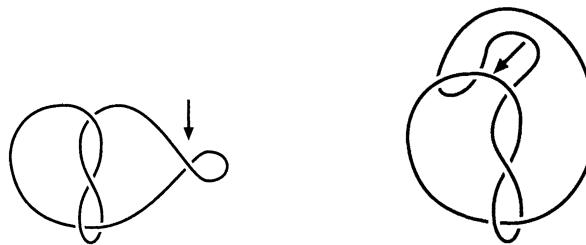


Figure 2.6: Solution to *Exercise 2.6*.

- *Exercise 2.7:* How many different sequences of the integers 2 4 6 8 10 12 14 are there? (This exercise gives us an upper bound on the number of possible alternating knot projections with seven crossings; however, it's far from accurate.)
 - There are $7! = 5040$ sequences.
- There is a slightly modified Dowker notation for nonalternating knots.
 - “If the even integer is assigned to the crossing while we are on the overstrand at that crossing, we leave the even integer positive. But if the even integer is assigned to the crossing while we are on the understrand of that crossing, we make the corresponding even numbers negative” (39).
- *Exercise 2.9:* How can you recognize from the sequence of numbers that a projection has a trivial crossing in it like in Figure 2.7a? How about recognizing a Type II Reidemeister move that will reduce the number of crossings by two like in Figure 2.7b?



(a) Type I Reidemeister move. (b) Type II Reidemeister move.

Figure 2.7: Recognizing extra crossings from the Dowker notation.

- A Type I Reidemeister move can be recognized from the Dowker notation sequence by the pairing of two numbers that have an absolute difference of 1. In Figure 2.7a, a sample Dowker notation is 6 2 8 10 4. Since 2 pairs with 3, and $|2 - 3| = 1$, we know that there is a Type I Reidemeister move. Logically, this makes sense because when moving along the orientation, a trivial crossing will be crossed twice in a row.
- A Type II Reidemeister move can be recognized from the Dowker notation sequence by a specific characteristic of two pairs: If (x_1, y_1) and (x_2, y_2) is such a pair, where the order is listed such

that $|x_2| > |x_1|$ and $|y_1| > |y_2|$, then both of the following must be true.

$$||x_1| - |x_2|| = 1 = ||y_1| - |y_2|| \quad |x_1 - y_2| \neq ||x_1| - |y_2|| \quad (2.1)$$

In Figure 2.7b, a sample Dowker notation is 8 6 10 -2 12 4. Since $(-2,7)$ and $(3,6)$ are pairs that satisfy the above, we know that there is a Type II Reidemeister move. Logically, this makes sense because, when moving along the orientation, such crossings' numbers will be sequential both times crossed (confirmed by the left statement in Equation 2.1). Additionally, the even numbers must have opposite signs or the strand would be linked, as opposed to entirely above or below (confirmed by the right statement in Equation 2.1).

- The most consequential fact of Dowker notation is that it allows us to “feed projections of knots into the computer simply as a sequence of numbers” (40).

2.3 Conway’s Notation

- Particularly suited to calculations involving **tangles**.
- **Tangle:** “A region [in a knot or link] in the projection plane surrounded by a circle such that the knot or link crosses the circle exactly four times” (41).
 - These four crossings will be thought of as occurring in the four ordinal compass directions.
- Tangles can be thought of as building blocks in knot or link projections.
- **Equivalent:** Two tangle projections that can be transformed into each other via a series of Reidemeister moves “while the four endpoints of the strings in the tangle remain fixed and while the strings of the tangle never journey outside the circle defining the tangle” (41).
- Notice that, if a knot is formed by gluing together the tangle’s ends in pairs, the knot is equivalent to other projections of itself as long as the tangles are equivalent (by a series of Reidemeister moves).
- Tangles, such as the one in Figure 2.8c are denoted by the number of left-hand twists (crossings).
 - If we had twisted the tangle the other way, we would have called it -3 .
 - More simply, “for a positive-integer twist, the overstrand always has a positive slope, if we think of it as a small segment of a line [perhaps a cubic, in the case of Figure 2.8c]” (42).

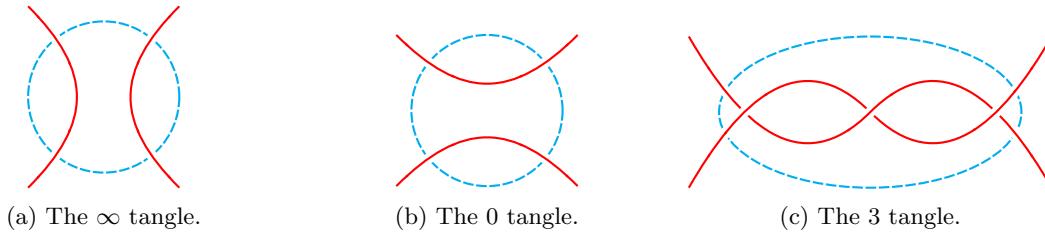


Figure 2.8: Tangles.

- **0 tangle:** The simplest of all tangles (by definition). *Also known as trivial tangle.* See Figure 2.8b.
- If we reflect the 3 tangle across a line connecting the NW and SE corners, and then proceed to extrude the NE and SE strands, twisting positively twice, we yield the tangle, 3 2.
- If we reflect 3 2 across the same line, and then proceed to extrude the NE and SE stands, twisting negatively four times, we yield the tangle, 3 2 -4 .
- **Rational tangle:** Any tangle that can be constructed in the manner described in the above 2 bullet points.

- Note that if the tangle is represented by an even number of integers, it can be thought of as being constructed from the ∞ tangle, and vice versa for the 0 tangle.
- *Exercise 2.10:* Draw the rational tangles corresponding to $2 -3 4 5$ and $3 -1 3 -3 2$.

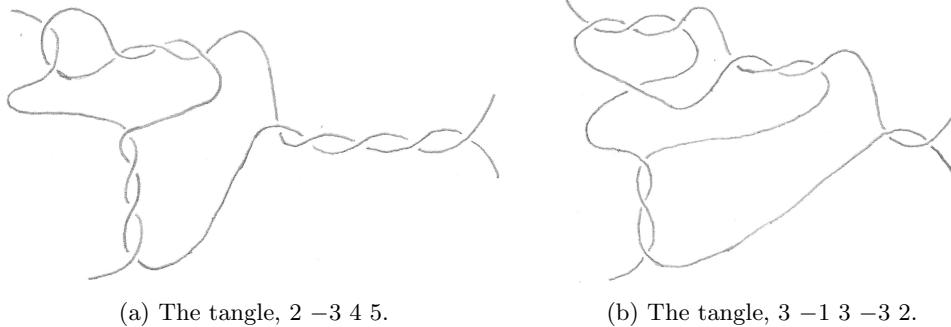


Figure 2.9: Solution to *Exercise 2.10*

- There is an extremely simple way to tell if two rational tangles are equivalent from their notation: **continued fractions**.
- **Continued fraction:** A fraction obtained through the iterative process of summing a number and its reciprocal.

– For example, given the three numbers a , b , and c , the continued fraction would be the following.

$$c + \frac{1}{b + \frac{1}{a}} \quad (2.2)$$

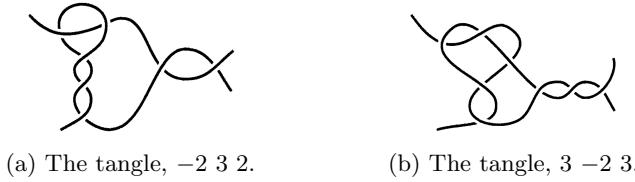


Figure 2.10: Two equivalent tangles.

- Applied to knot theory, we know that the tangles, $-2 3 2$ and $3 -2 3$, (seen in Figure 2.10) are equivalent because of the following.

$$\begin{aligned} 2 + \frac{1}{3 + \frac{1}{-2}} &= 3 + \frac{1}{-2 + \frac{1}{3}} \\ 2 + \frac{1}{\frac{5}{2}} &= 3 + \frac{1}{\frac{-5}{3}} \\ 2 + \frac{2}{5} &= 3 + \frac{-3}{5} \\ \frac{12}{5} &= \frac{12}{5} \end{aligned} \quad (2.3)$$

- *Exercise 2.13:* Determine which of the four rational tangles in Figure 2.11 are equivalent^[4].

⁴Idea for first paper: Prove that a series of NE↔SE and SE↔SW moves is sufficient to generate any rational tangle (like Reidemeister moves). Apply these findings to defining a tangle from a projection. Potentially create a new notation (show how this notation and the current tangle notation convert).

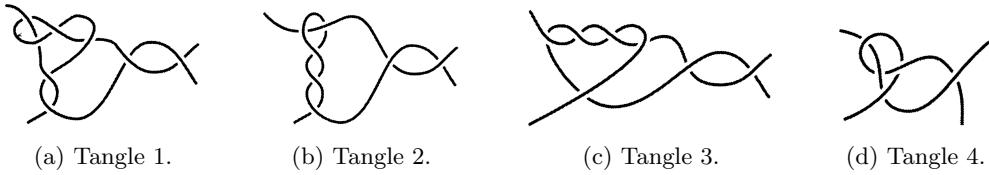


Figure 2.11: Four rational tangles.

– Tangle 1 (Figure 2.11a) can be denoted $-1 -1 -2 -2 -2$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 -2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-1 + \frac{1}{-1}}}}} &= -2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-2 + \frac{1}{-2}}}} \\
 &= -2 + \frac{1}{-2 + \frac{-2}{5}} \\
 &= -2 + \frac{-5}{12} \\
 &= \frac{-29}{12}
 \end{aligned} \tag{2.4}$$

– Tangle 2 (Figure 2.11b) can be denoted $2 -3 2$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 2 + \frac{1}{-3 + \frac{1}{2}} &= 2 + \frac{-2}{5} \\
 &= \frac{8}{5}
 \end{aligned} \tag{2.5}$$

– Tangle 3 (Figure 2.11c) can be denoted $-4 1 2$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 2 + \frac{1}{1 + \frac{1}{-4}} &= 2 + \frac{4}{3} \\
 &= \frac{10}{3}
 \end{aligned} \tag{2.6}$$

– Tangle 4 (Figure 2.11d) can be denoted $1 1 1 1 1$.

■ Its continued fraction simplifies as follows.

$$\begin{aligned}
 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1}}}}} &= 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{2}}}} \\
 &= 1 + \frac{1}{1 + \frac{2}{3}} \\
 &= 1 + \frac{3}{5} \\
 &= \frac{8}{5}
 \end{aligned} \tag{2.7}$$

– Therefore, tangles 2 and 4 (Figures 2.11b and 2.11d, respectively) are equivalent by equations 2.5 and 2.7.

- If we close off the ends of a rational tangle, i.e. NE to NW and SE to SW, we form a rational link.

- If the link is a knot, we can denote it by its tangle. For example, the figure-eight knot is a rational knot with **Conway notation** 22 (because of the two twists twice in its center and the gluing of its ends).
- **Conway's notation:** A method of denoting knots by their tangles (all aforementioned theory in this section).
- Multiplying tangles:
 - Reflect the left tangle across its NW to SE diagonal and then glue its new NE and SE ends to the NW and SW ends, respectively, of the adjacent (to the right) tangle.
 - Multiplying a rational tangle by an integer tangle will always generate a rational tangle, e.g. 32 comes from multiplying 3 by 2.
 - To reflect a tangle over its NW-SE line, multiply it by the tangle, 0.

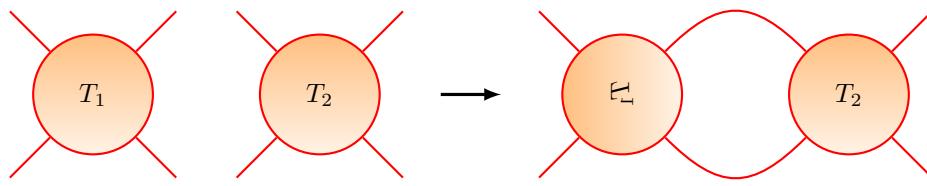


Figure 2.12: Multiplying tangles.

- Adding tangles:
 - Glue the left tangle's NE and SE ends to the right tangle's NW and SW ends, respectively.
 - Note that no reflecting occurs here.
 - “If we multiply each tangle in a sequence of tangles by 0, and then add them all together, we denote the resultant tangle by the sequence of numbers that stand for the original tangles, only now separated by commas” (48).
- **Pretzel knot:** “A knot obtained from a tangle represented by a finite number of integers separated by commas” (48).
- **Algebraic tangle:** “Any tangle obtained by the operations of addition and multiplication on rational tangles” (48).
- **Algebraic link:** “A link formed when we connect the NW string to the NE string and the SW string to the SE string on an algebraic tangle” (48). *Also known as aborescent link.*
 - Denoted the same way as a tangle.
 - Fun fact: The trefoil knot (Figure 1.1b) is an algebraic link of the 3 tangle (Figure 2.8c).
- **Additive identity:** A quantity that can be added to a second quantity without changing that second quantity.
 - 0 is an additive identity for the real numbers.
- *Exercise 2.21:* Is there an additive identity for tangles?
 - The 0 tangle (Figure 2.8b) is an additive identity for tangles. Adding the 0 tangle does the same thing as extruding (performing a planar isotopy on) the NE and SE strands. Such an addition does not change the relative position of any strand's ordinal end(s).
 - Continuing the analogy to the real numbers, it makes sense that the 0 tangle would be analogous to 0, the real number.

- **Multiplicative identity:** A quantity that can be multiplied by a second quantity without changing that second quantity.
 - 1 is a multiplicative identity for the real numbers.
- *Exercise 2.22:* Is there a multiplicative identity for tangles? Is it the same if you multiply a tangle by it on the right side or the left side?
 - The 0 tangle (Figure 2.8b) is a multiplicative identity for tangles. Multiplying the 0 tangle does the same thing as extruding (performing a planar isotopy on) the SE and SW strands. Such a multiplication does not change the relative position of any strand's ordinal end(s).
 - To part 2, yes.
- The analogy to real numbers does not hold in some areas.
 - Multiplication on tangles is not commutative — $ab \neq ba$ for tangles.
 - Multiplication on tangles is not associative — $(ab)c \neq a(bc)$ for tangles.
 - There are no inverse tangles — there is no tangle that when added to T gives back the trivial tangle, 0.
- There are tangles that are not algebraic.
- New knots can be obtained via **mutation**.
- **Mutation:** Severing the connections between a tangle and any adjacent tangles, transforming it (flipping horizontally, vertically, or both), and gluing the strands to their new adjacent counterparts.
 - Knots formed in this manner are called **mutants** of one another.
- *Exercise 2.25:* Show that if we have three tangles as in Figure 2.13a, we can mutate several times in order to permute the tangles. Note that we can then permute n tangles in a row.

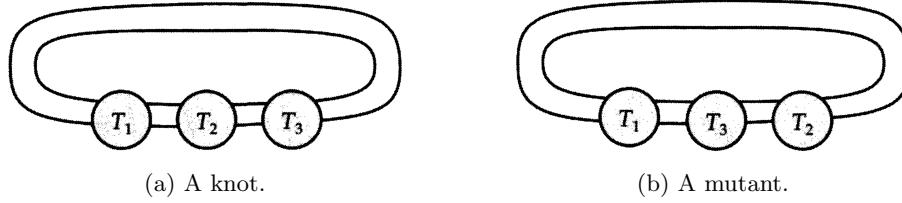


Figure 2.13: Series of mutations.

- Beginning in Figure 2.13a, horizontally flip T_2 .
- Horizontally flip T_3 .
- Horizontally flip the unit of the flipped T_2 and T_3 .
- Mutation cannot turn a nontrivial knot into the trivial knot (Figure 1.1a).
- More on mutant knots can be found in Sections ?? and ??.

2.4 Knots and Planar Graphs

- This section's notation bridges knot theory and graph theory.
 - It is capable of contributing to both branches of mathematics.
- **Graph:** A set of points called vertices that are connected by edges.

- **Planar graph:** A graph that lies in a plane.
 - Creating a planar graph from a projection of a knot or link:
 - “Shade every other region of the link projection so that the infinite outermost region is not shaded” (51).
 - “Put a vertex at the center of each shaded region and then connect with an edge any two vertices that are in regions that share a crossing” (52).
 - Assign an orientation.
 - “Label each edge in the planar graph with a + or a –, depending on whether the edge passes through a + crossing or a – crossing” (52). See Figure 2.14. Recall computing link numbers in Section 1.4 (see Figure 1.18, specifically).
 - The result is a **signed planar graph**. See Figure 2.15 for an original example.

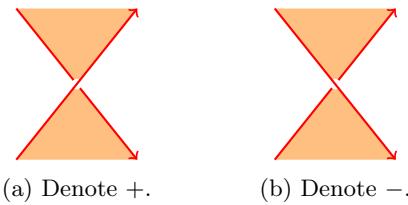


Figure 2.14: Computing signs per crossing.

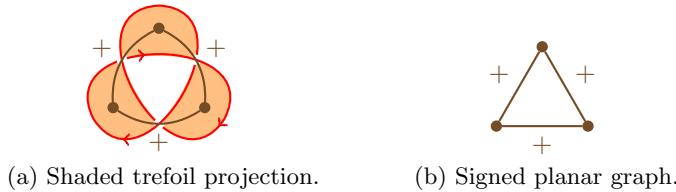


Figure 2.15: A signed planar graph from the standard trefoil knot projection.

- It is also possible return the signed planar graph to a knot projection:
 - “Put an x across each edge” (53).
 - “Connect the edges inside each region of the graph” (53).
 - “Shade those areas that contain a vertex” (53).
 - “At each of the x ’s, put in a crossing corresponding to whether the edge is a $+$ edge or a $-$ edge” (53).
 - The significance here, as previously referenced, is that we can turn questions about knots into questions about graphs.
 - *Exercise 2.31:* What do the Reidemeister moves become when translated into operations on signed planar graphs? (Make sure you consider what happens when different regions are shaded.)
 - Type I Reidemeister move: Extend an edge from one vertex to a newly created one. Either sign is acceptable (the sign option is analogous to the over-under crossing option).
 - Type II Reidemeister move: Add two edges connecting two, specific vertices and intersecting no other edges. The signs on the new edges must be opposite, but it either can be positive (or negative; this option is analogous to the over-under crossing option).
 - Type III Reidemeister move: I do not know^[5].
 - See Section ?? for more on signed planar graphs.

⁵Return to this later.

3 Invariants of Knots

3.1 Unknotting Number

- **Unknotting number:** A value $n \in \mathbb{N}$ specific to a knot K that gives the fewest number of crossing changes needed in any projection to turn it into the unknot. *Also known as $\mathbf{u}(K)$.*
 - It may, however, be a nasty projection of the unknot (see Figure 1.12 and the related Exercise).
- *Exercise 3.1:* Find the unknotting number of the figure-eight knot (Figure 1.4).
 - $u(K) = 1$, where K is the figure-eight knot.
 - K is distinct from the unknot, so $u(K) \neq 0$. The next possibility ($u(K) = 1$) succeeds with the projection in Figure 1.4 (in fact, flipping *any* crossing in the referenced projection trivializes K).
- Note that performing all of the crossing changes on one projection is equivalent to performing one such change, then an ambient isotopy, then another...
 - A proof of this is listed on page 58.
- “The fact that every knot has a finite unknotting number follows from the fact that every projection of a knot can be changed into a projection of the unknot by changing some subset of the crossings in the projection” (58).
 - See Exercise 1.7.
- A proof of the above:
 - Pick a starting point (not a crossing) on an arbitrary knot and an orientation in whose direction to traverse the knot.
 - When you arrive at the first crossing, change it so that the strand that you are on is the over strand, if necessary.
 - Continue changing every new crossing to which you come to make the current strand into the overstrand.
 - Do not change crossings that you have already passed through once.
 - At this point, you now have a projection of a knot that was obtained from the original knot by changing crossings and that is the trivial knot. A proof of this triviality follows.
 - Consider the knot in three-space, with the z -axis coming out of the page. Assign to the starting point the z -coordinate, $z = 1$.
 - Decrease the z -coordinate of each consecutive point as you traverse the knot, culminating in the original point being labeled, $z = 0$.
 - Connect the starting and ending points with a vertical bar to preserve the knot as a knot.
 - What you have now obtained is a knot in three-space that, when projected from the top is the knot that we constructed before the line break, but has no crossings when projected from the side (and is therefore the trivial knot).
 - Q.E.D.
- It is very hard, in general, to find the unknotting number of a knot. How do you know that there isn’t a better projection?
- “A knot with unknotting number 1 is prime” (61).
- Therefore, a composite knot cannot be unknotted with a one crossing change.

- Another way to think about this idea is that changing a crossing to unknot one factor knot would yield the composition of the other factor knot with the unknot, which, according to Section 1.2, is simply the other factor knot (obviously still fully knotted).

$$u(K_1 \# K_2) \leq u(K_1) + u(K_2) \quad (3.1)$$

- The unknotting number can be realized by a projection that is not **minimal**.
 - **Minimal projection:** The projection of a knot with the lowest possible number of crossings.
 - **k -move:** A local change in the projection that replaces two untwisted strings with two strings that twist around each other with k crossings in a right-handed manner. See Figure 3.1.
- If k is positive, then the overstrand has a positive slope.

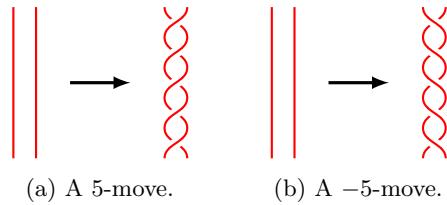


Figure 3.1: Examples of k -moves.

- **k -equivalent:** Two knot or link projections that can be transformed into each other through a series of k -moves and $-k$ -moves.
 - Ambient isotopies are allowed between k -moves.

- **Exercise 3.8:** Show that the knots in Figure 3.2 are each three-equivalent to a trivial link.

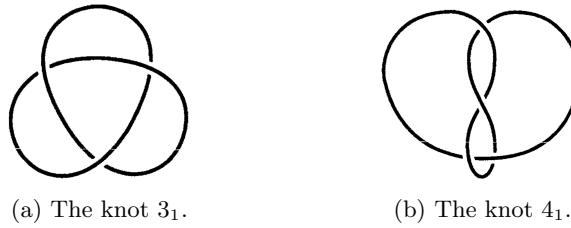


Figure 3.2: Knots that are 3-equivalent to trivial links.

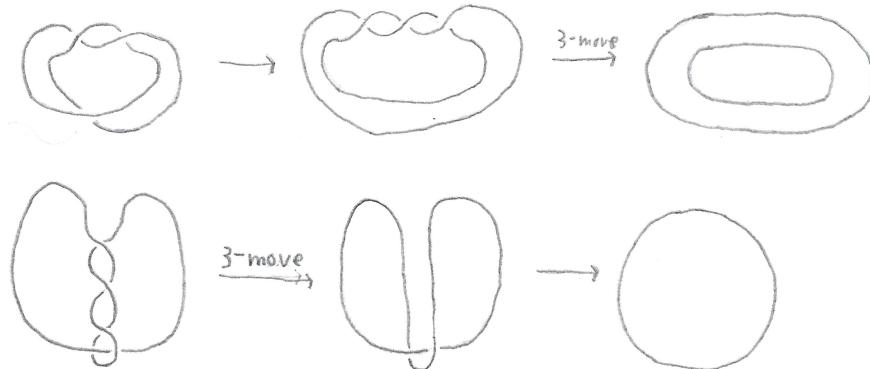


Figure 3.3: Solution to *Exercise 3.8*.

3.2 Bridge Number

- Think of knots as “cutting through the projection plane, rather than lying in it” (64).
 - Every overstrand is entirely above and every understrand is entirely beneath. The points where the strand shifts from over to under are known as **vertices**.
- **Bridge number:** The least number of unknotted arcs lying above the plane in any projection of the knot. *Also known as $b(K)$*
 - The number of maximal overpasses in in the projection.
- **Overpass:** “A subarc of the knot that goes over at least one crossing but never goes under a crossing” (64).
- **Maximal overpass:** “An overpass that could not be made any longer” (65).



Figure 3.4: Overpasses.

- *Exercise 3.10:* Show that if a knot has bridge number 1, it must be the unknot.
 - If a knot has bridge number 1, \exists a projection of the knot such that there is only one arc above the projection plane and 1 arc below the projection plane.
 - This projection, viewed from above, will look like any isotopy of the knot. However, viewed from the side, there will be 0 crossings — therefore, it is the unknot.
- *Exercise 3.11:* Show that the knot 5_2 has bridge number $2^{[6]}$.
 - Since 5_2 is distinct from the unknot, $b(K) \neq 1$.
 - Therefore, the next possibility to check is 2. If \exists a projection of 5_2 with bridge number 2, $b(K) = 2$.

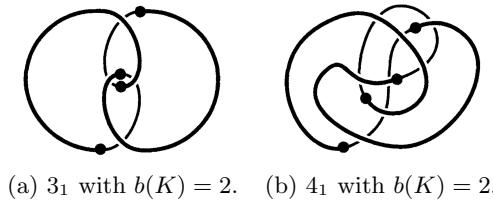


Figure 3.5: Reducing bridge number between projections.

- Before tackling this problem, though, let’s look at some simpler examples.
 - From these, we notice that there are “bridging moves” (distinct combinations of the Reidemeister moves) that can shift the bridge number by ± 1 .
 - Furthermore, long overstrands call to mind a Dowker reconstruction (Figure 2.5).
- Find the Dowker notation of 5_2 .
 - The Dowker notation of the projection listed in the appendix is 6 10 8 2 4.

⁶Idea for a paper: On reducing bridge number. Discuss manipulating a projection to its Dowker reconstruction to create the “killzone” and then using “bridging moves” to make the projection nonalternating. 2 types so far: I \rightarrow III, and 2xII \rightarrow III (in both cases it is the Type III move which climactically changes $b(K)$). It can be necessary to increase $b(K)$ before decreasing it.)

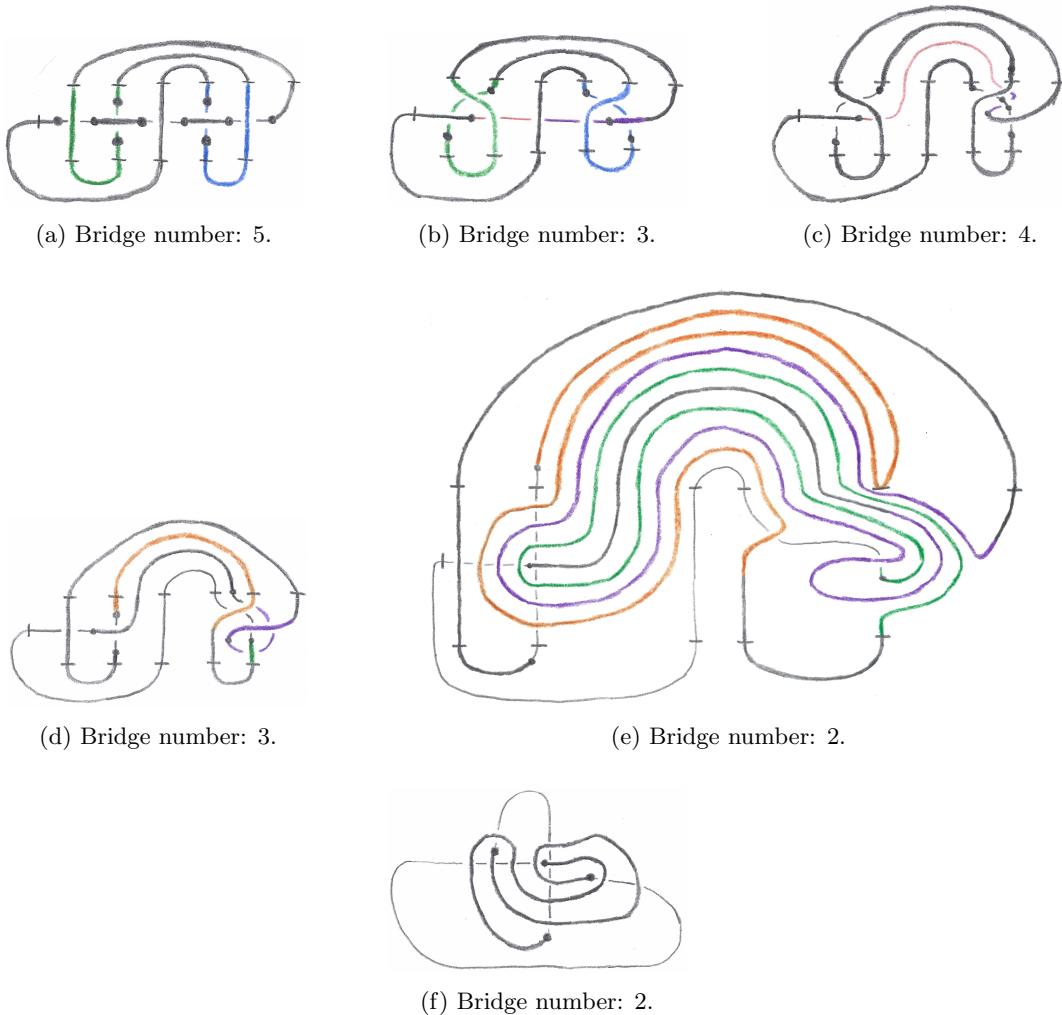


Figure 3.6: Solution to *Exercise 3.11*.

- Note that the following hold for all projections in Figure 3.6.
 - Dashes, when present, represent the boundaries of the “killzone.” They also provide consistent reference points for Figures 3.6a through 3.6e.
 - Dots, when present, represent bounds on the overpasses (and, consequently, the underpasses).
 - Thicker lines represent overpasses. Conversely, thinner lines represent underpasses.
 - Like-colored strands between sequential projections represent the same strand after Reidemeister moves or planar isotopies (colored for clarity, not for any specific purpose such as tricolorability).
- From the previously determined Dowker notation, reconstruct a knot projection. This projection (Figure 3.6a) contains the “killzone,” a concentrated region where every crossing is displayed in an orderly fashion^[7].
- Employ two Type I Reidemeister moves followed respectively by Type III Reidemeister moves to eliminate two overpasses. These moves occur on the green and blue strands, independently. They

⁷Note that from here-down, I do not guarantee that this is the fastest or simplest way to generate a projection that displays $b(K) = 2$, nor can I assure the reader that my solution is the “simplest” projection (perhaps the one with the fewest crossings). However, I do guarantee that it is correct. I know that it is so because all projections are separated only by Reidemeister moves and planar isotopies, and Figure 3.6f visibly has bridge number 2.

yield a projection with bridge number 3 (Figure 3.6b). This projection has several noteworthy properties.

- First, of the three overpasses, two contain 2 crossings while one contains 3 crossings.
- Second, of the three underpasses, two contain 2 crossings while one contains 3 crossings.
- Third, every overpass crosses at least one of its adjacent underpass.
- Because each overpass and underpass has more than one crossing, the bridge number of the projection must increase before it can decrease (one of the overpasses must be split into two [or three]; single overpasses can subsequently be eliminated via “bridging moves”).
- Because each overpass crosses at least one of its adjacent underpasses, no “bridging moves” can eliminate a bridge from this projection directly^[8].
- Employ two Type II Reidemeister moves followed by one Type III Reidemeister move to split one of the 2-crossing overpasses into two 1-crossing overpasses. These moves occur on the red-purple strand, which is “slid along” beneath the central overpass. This change will increase the bridge number to 4 and yield Figure 3.6c^[9].
- Employ a Type I Reidemeister move followed by a Type III Reidemeister move to eliminate one of the new 1-crossing overpasses. These moves occur on the purple strand. This change will decrease the bridge number back to 3. Additionally, perform the reverse of the actions on the green strand between Figures 3.6a and 3.6b to “straighten it out.” This shift will switch the 1-crossing overpass to an adjacent strand. These changes will yield Figure 3.6d. This projection has several noteworthy properties.
 - The 1-crossing overpass has underpasses of greater than 1 crossing on both sides.
 - The 1-crossing overpass passes through a loop, creating a setup similar to the two overpasses eliminated between Figure 3.6a and 3.6b. However, because of the above, a Type I Reidemeister move followed by a Type III Reidemeister move will not eliminate this overpass — it would only shift the position of the overpass.
 - The 1-crossing overpass crosses neither of its adjacent underpasses.
 - The above is key — because of it, a succession of 2xType II Reidemeister moves followed by 1xType III Reidemeister moves can combine the 1-crossing overpass with the partially purple overpass.
 - This set of moves will occur thrice — once for each crossing of the adjacent underpass.
- Three times sequentially, employ two Type II Reidemeister moves followed by one Type III Reidemeister move to (for the first two sets) shift the position of the 1-crossing overpass and then (for the last set) eliminate the 1-crossing overpass. These moves occur on the orange, purple, and green strands, sequentially. They yield a projection with bridge number 2 (Figure 3.6e).
- Q.E.D.
- Figure 3.6e is rather visually unpleasant, however. Figure 3.6f is the result of one Type I Reidemeister move (on the upper part of the orange strand) and planar isotopies intended purely to “clean up” the projection. It remains indistinct from 5_2 and displays $b(K) = 2$.

- **Two-bridge knots:** Knots that have $b(K) = 2$.
- If a two-bridge knot is sliced through the projection plane, there will be two **unknotted untangled** arcs both above and below said plane.
- **Unknotted untangled:** Distinct arcs that have no crossings with each other.
- From 2 such unknotted untangled arcs on each side of the projection plane, it is possible to construct all two-bridge knots.

⁸I am holding this to be true because I cannot demonstrate otherwise, but I do not know that it is true. I'd actually conjecture that it's false — it would seem to make sense that some overpass could be eliminated directly from Figure 3.6b.

⁹Multiple such moves may suffice here to split a 2-crossing bridge. I chose and followed through this one randomly.

- “The tricky part is that although the strings to each side of the plane are individually unknotted, they can twist around each other and themselves” (65).
- All two-bridge knots are rational knots and vice versa.
- Since the two-bridge knots are a very well understood class of knots, properties that are suspected to hold true for all knots are often first proved to hold true for two-bridge knots.
- All two-bridge knots are prime.
- The first tabulated three-bridge knot is 8_{10} .
- *Exercise 3.13:* Find a picture of 8_{10} that shows that it is at most a three-bridge knot.



Figure 3.7: The knot 8_{10} .

- Find the Dowker notation of 8_{10} ^[10].
 - The “maximum” Dowker notation (which, fortunately, is also “ideal”) for 8_{10} is 12 16 10 14
4 6 2 8
- From the previously determined Dowker notation, reconstruct a knot projection.
- There are four possible “type I bridging moves”^[11] allowed by this projection. Perform 3 of these to eliminate 3 bridges.
 - For the “linked loops,” it does matter on which loop the move is performed.
 - This yields a projection with 2x3-crossing bridges, 2x2-crossing bridges, and 1x1-crossing bridge.
- Perform two “type II bridging moves”^[12] to eliminate the 1-crossing bridge.
- Perform four “type II bridging moves” to eliminate a 2-crossing bridge.
- Q.E.D. — see Figure 3.8 for the final projection^[13].
- “In 1954, a German mathematician named H. Shubert proved that $b(K_1 \# K_2) = b(K_1) + b(K_2) - 1$ ” (67).
- *Exercise 3.14:* Explain how Shubert’s result implies that rational knots are all prime.
 - To be distinct from the unknot, a knot must have $b(K) \geq 2$, as proven in *Exercise 3.10*.
 - By Shubert’s theorem, above, and the above fact, any composition of two nontrivial knots must have $b(K) \geq 3$ (at least $2 + 2 - 1 = 3$).
 - Since any knot with $b(K) \geq 3$ is not rational and all composite knots have $b(K) \geq 3$, all rational knots (those with $b(K) = 2$) must be noncomposite, or prime.

¹⁰Note that different Dowker notations will lead to different Dowker reconstructions — some may be easier to manipulate than others. To concentrate the “killzone” in a single (or at least close) to rectangular region, choose Dowker notations with maximum “base strands.” Avoid, in the beginning, passing through at a crossing that will soon be passed through again.

¹¹A Type I Reidemeister move performed on an isolated loop followed by a Type III Reidemeister move to translate the new crossing across the “base strand.”

¹²Two Type II Reidemeister moves performed on the same strand to move it over either side of an adjacent crossing followed by a Type III Reidemeister move to translate the strand over the crossing followed by two Type II Reidemeister moves for simplification purposes.

¹³Note that this figure is drawn with the same rules used in Exercise 3.11, i.e. the original “killzone” is marked as well as the start and end of each bridge. Each “bridging move” has been carried out without simplification, so some thought should be able to reveal the connection between the original Dowker reconstruction (unpictured) and this projection.

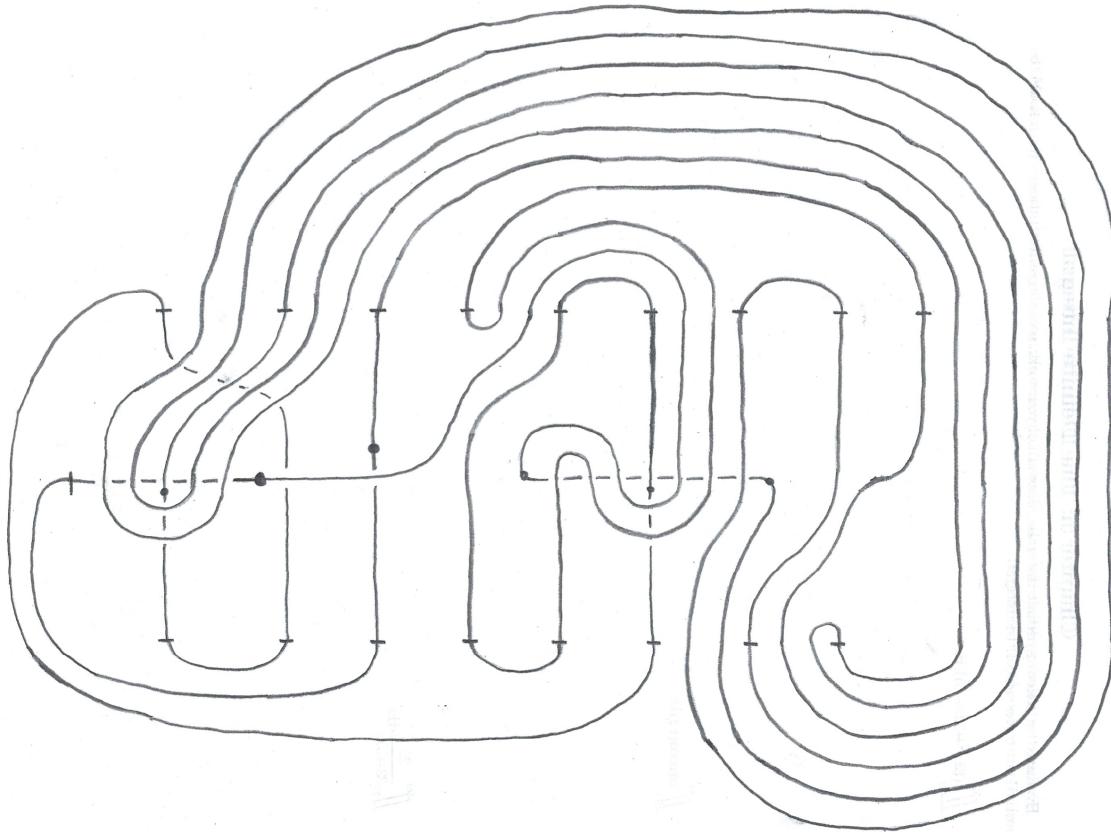


Figure 3.8: Solution to *Exercise 3.13*.

3.3 Crossing Number

- **Crossing number:** “The least number of crossings that occur in any projection of the knot” (67). *Also known as $c(K)$.*
- To find the crossing number of a knot, first find a projection of the knot. Count the crossings, n . At this point, it is clear that $c(K) \leq n$.
 - If all knots with fewer crossing than n are known and K is distinct from all of those knots, then $c(K) = n$.
 - Otherwise, reduce the knot to its minimal projection.
- Proving that a knot is distinct from every knot of fewer crossings is very difficult and requires knot polynomials (see Section ??).
- Determining $c(K)$ is generally very difficult (knots are currently only tabulated up to 13 crossings).
- **Reduced:** A projection of a knot with no easily removed crossings.
 - All “easily removed” crossings in alternating knots are shown in Figure 3.9.
- However, in 1986, Lou Kauffman, Kunio Murasugi, and Morwen Thistlethwaite all independently proved that “an alternating knot in a reduced alternating projection of n crossings has crossing number n ” (68)..
 – They utilized the Jones polynomial (see Section ??).

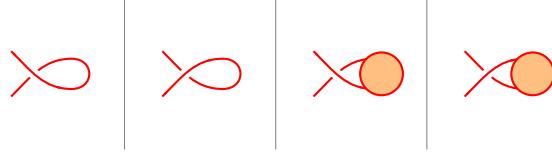


Figure 3.9: Easily removed crossings in an alternating knot projection.

- Because we can visually tell whether or not a projection is reduced (and reduce it if it's not), we can determine $c(K)$ for any alternating knot.
- It is still unknown whether or not $c(K_1 \# K_2) = c(K_1) + c(K_2)$.
 - However, Kaufman, Murasugi, and Thistlethwaite's result proves that the above holds when $K_1 \# K_2$ is alternating.
 - *Exercise 3.15:* Show that if K_1 and K_2 are alternating, then so is $K_1 \# K_2$. Hence, $c(K_1 \# K_2) = c(K_1) + c(K_2)$ holds when K_1 and K_2 are alternating even if $K_1 \# K_2$ does not appear alternating.

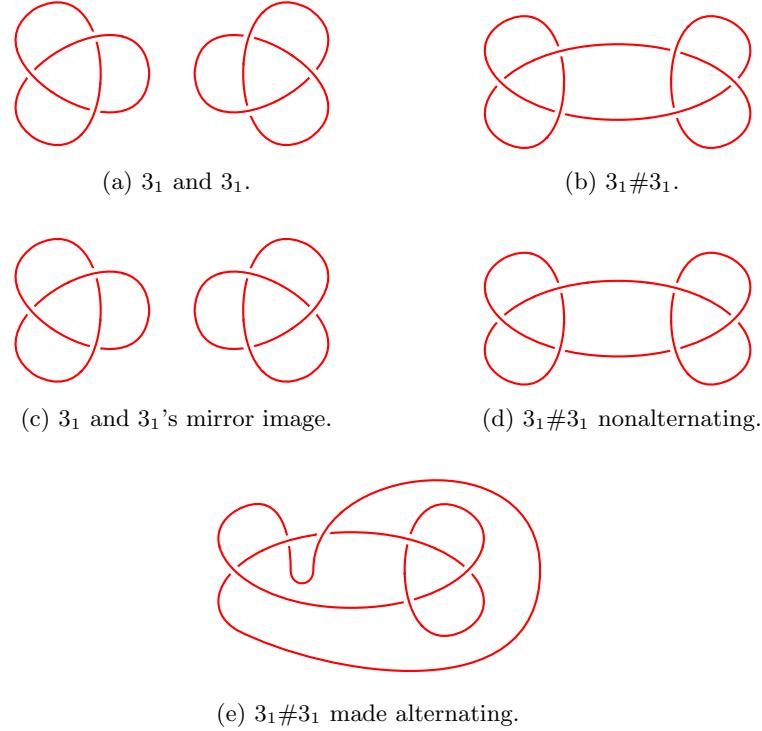


Figure 3.10: Drawing an alternating projection of $K_1 \# K_2$ where K_1 and K_2 are alternating.

- When composing alternating knots, there are two cases, as exemplified by Figure 3.10a and 3.10c, respectively.
 - Case 1: the composite knot will be in an alternating projection.
 - Case 2: the composite knot will not be in an alternating projection.
- For Case 1, the theorem evidently holds. See Figure 3.10b
- For Case 2, first consider the definition of an alternating knot (see Section 1.1), which states that an alternating knot need only have a projection that is alternating, i.e., every projection need not be alternating.

- Therefore, if we can prove that every nonalternating projection obtained through a composition has an alternating equivalent, the theorem will be true. To begin, consider several subcases to determine a pattern in the nonalternation.
 - Because K_1 and K_2 are both alternating, any arc being removed (to compose the two) will necessarily be on an understrand becoming an overstrand or vice versa.
 - Therefore, if we leave K_1 to enter K_2 from an understrand, we will return into K_1 to meet an overstrand and vice versa. The same holds for K_2 .
 - Case A: Leaving from an understrand in K_1 , enter K_2 on an understrand. The next crossing (continuing through K_2) will be an overstrand, and then an understrand, ... Eventually, by the above, we will return from K_2 on an overstrand into an overstrand on K_1 . The next crossing (continuing through K_1) will be an understrand, and then an overstrand, ... Eventually, we will return to where we started.
 - Case B: Leaving from an overstrand in K_1 , enter K_2 on an overstrand. The next crossing (continuing through K_2) will be an understrand, and then an overstrand, ... Eventually, by the above, we will return from K_2 on an understrand into an understrand on K_1 . The next crossing (continuing through K_1) will be an overstrand, and then an understrand, ... Eventually, we will return to where we started.
- In both Cases A and B, the only moments of nonalternation are two consecutive underpasses (or overpasses) followed by a factor knot followed by two consecutive overpasses (or underpasses). Notice this pattern in Figure 3.10d.
- Thus, “slide” one strand of a factor knot over the other knot to eliminate the nonalternation, as in Figure 3.10e.
- Q.E.D.

4 Surfaces and Knots

4.1 Surfaces without Boundary

- **Surface:** The two-dimensional exterior of a three-dimensional object with volume.
- **Torus:** “The surface of a donut” (71). *Plural is tori.*
- “At any point on a surface, there is a small region on the surface surrounding and containing the point that looks like a disk” (72).
 - The disk may be deformed (curved), but it must still be a disk.
 - Objects that are not surfaces have points with nondisk “neighborhoods” around them.
- Surfaces can also be called **two-manifolds**.
- **Two-manifold:** “Any object such that every point in that object has a neighborhood in the object that is a (possibly nonflat) disk” (73).
- *Exercise 4.1:* Based on the definition for two-manifolds, decide what the definition should be for a one-manifold. Find two different one manifolds.

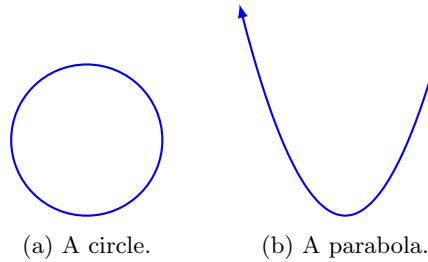


Figure 4.1: Two one-manifolds.

- Any object such that every point in that object has a neighborhood in the object that is a (possibly not straight) line.
- See Section ?? for info on three-manifolds.
- Think of surfaces as being made of rubber. In Figure 1.2, the deformation is an **isotropy** and the left and right surfaces are **isotropic surfaces**.
- **Isotropy:** A continuous deformation of a surface in space; a rubber deformation.
- **Isotropic surfaces:** Two surfaces in space that are equivalent under an isotropy.
- Surfaces that are not isotropic cannot be deformed into one another without some cutting and pasting!^[14]

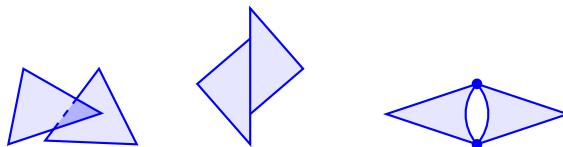


Figure 4.2: Triangles cannot intersect like this.

¹⁴Identify surfaces by the number of holes in them?

- **Triangulation:** A division of a surface into deformable triangles.
 - Think of a triangle as a disk “with a boundary made up of three edges connecting three vertices” (75).
 - The triangles must “fit together nicely along their edges so that they cover the entire surface” (74). See Figure 4.2 for a visualization of ways that triangles cannot fit together.
 - Given a triangulation of a surface, the triangles can be separated and we can keep track of the original surface by labeling edges and putting arrows on them.

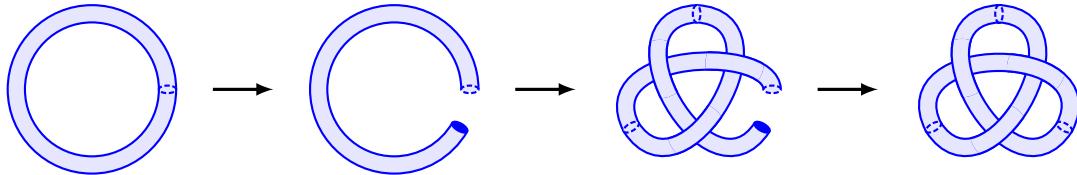


Figure 4.3: Homeomorphic surfaces.

- **Homeomorphic:** Two surfaces such that “one of them can be triangulated, then cut along a subset of the edges into pieces, and then glued back together along the edges according to the instructions given by the orientations and labels on the edges, in order to obtain the second surface” (76).
 - For example, the leftmost and rightmost surfaces in Figure 4.3 are both homeomorphic with a torus.
 - However, proving that the two are not isotropic would take some work.
 - It is not necessary to find a complete triangulation to prove that two surfaces are homeomorphic — it will suffice to circularly cut and put like circles back together, as in Figure 4.3.
 - One could also cut just edges and put like edges back together.
 - A sphere and a torus are clearly not homeomorphic because every closed loop on a sphere cuts it into two pieces. However, there exist closed loops on a torus that do not cut it into two pieces^[15].
 - The sphere, torus, two-hole torus, and three-hole torus are all not homeomorphic to one another.
 - Since we could keep increasing the number of holes, there are infinitely many distinct (nonhomeomorphic) surfaces.
- **Genus** (of a surface): “The number of holes in the doughnut” (78). *Plural is genera.*
 - The sphere has genus 0 while the torus has genus 1.
- **Embedding** (of a surface): “A choice of how to place a surface in space” (78).
 - The leftmost and rightmost surfaces of Figure 4.3 are two distinct embeddings of the torus in three-space.
- **Cube-with-holes:** The surface of a solid object obtained by drilling wormholes out of a cube.
- **Homeomorphism type:** The elementary surface to which a surface can be related.
 - Could be found via the triangulation method/cut-and-paste technique.
 - Could be found using its **Euler characteristic**.

^[15]This is not an actual proof.

- **Euler characteristic:** Given a triangulation (or subdivision) of a surface, the sum $V - E + F$, where V is the number of vertices, E is the number of edges, and F is the number of triangles (or n -gon faces). Also known as χ .

$$\chi = V - E + F \quad (4.1)$$

- For example, χ for the leftmost sphere in Figure 4.4 is $6 - 12 + 8 = 2$.
- *Exercise 4.5:* Compute the Euler characteristic of the second triangulation of the sphere in Figure 4.4.

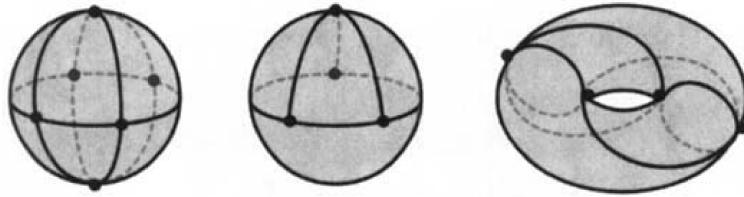


Figure 4.4: Triangulations of the sphere and the torus.

- $\chi = 4 - 6 + 4 = 2$
- *Exercise 4.6:* Compute the Euler characteristic of the triangulation of the torus in Figure 4.4.
 - $\chi = 4 - 12 + 8 = 0$
 - The sameness of the Euler characteristics of the two triangulations of the sphere exemplifies the fact that the Euler characteristic depends only on the surface, not the triangulation.
 - A summary of the proof that any two triangulations have the same Euler characteristic follows.
 - Take T_1 and T_2 , two different triangulations of a surface. We will overlay them and add appropriate edges (to make triangles) to make T_3 . This process will prove that T_1 and T_3 have the same Euler characteristic. Since T_2 could have T_1 overlaid, T_2 and T_3 also have the same Euler characteristic. Therefore, T_1 and T_2 have the same Euler characteristic (transitive property).
 - Each edge of T_2 intersects T_1 a finite number of times, and no vertices are overlaid — this can be assured by shifting where T_2 is overlaid by a slight bit.
 - Add vertices at the intersections of edges. The addition of one vertex and the subsequent splitting of one edge into two cancel out in the Euler characteristic.
 - Add the vertices of T_2 and a corresponding new edge to each. These edges should be a subset of the edges in T_2 . Vertices increase by 1 for each vertex and edges increase by 1 for each edge, so the change likewise cancels out.
 - Add the rest of the edges from T_2 . Note that for every edge added, a face is split in two, canceling out the change.
 - Add edges to split all n -gon faces into triangles. Each added edge splits a face, so the change, again, cancels out.
 - Q.E.D.
 - $\chi(\text{genus} = 0) = 2$.
 - $\chi(\text{genus} = 1) = 0$.
 - For a genus 2 surface, we could find a triangulation and calculate χ , or we could be clever.
 - **Connected sum:** Remove a disk from each of two objects and then glue them together.

- With triangulations of two tori, we can imagine that a genus 2 surface can be created via the connected sum of the two tori minus two matching faces.
- So, $F \rightarrow F - 2$, and because 3 vertices and 3 edges are merged, $V \rightarrow V - 3$ and $E \rightarrow E - 3$. The vertices and edges cancel out, giving the following property.

$$\chi(S) = \chi(S_1) + \chi(S_2) - 2 \quad (4.2)$$

- Therefore, $\chi(\text{genus} = 2) = -2$.
- Exercise 4.8:* Use connected sums to show that the Euler characteristic of a genus 3 surface is -4 .
 - $S_{\text{genus}=2} + S_{\text{genus}=1} = S_{\text{genus}=3}$.
 - $\chi(S_{\text{genus}=3}) = \chi(S_{\text{genus}=2}) + \chi(S_{\text{genus}=1}) - 2 = -2 + 0 - 2 = -4$.
- Exercise 4.9:* Use induction to show that the Euler characteristic of a surface of genus g is $2 - 2g$.
 - Basis step: $2 \stackrel{?}{=} 2 - 2(0)$.
 - Induction hypothesis: $\chi(S_{\text{genus}=g}) = 2 - 2g$
 - Induction step:

$$\begin{aligned} \chi(S_{\text{genus}=g+1}) &\stackrel{?}{=} 2 - 2(g + 1) \\ \chi(S_{\text{genus}=g}) + \chi(S_{\text{genus}=1}) - 2 &\stackrel{?}{=} 2 - 2g - 2 \\ (2 - 2g) + 0 - 2 &\stackrel{?}{=} 0 - 2g \\ 0 - 2g &\stackrel{?}{=} 0 - 2g \end{aligned}$$

- Therefore, the following equation holds true.

$$\chi(g) = 2 - 2g \quad (4.3)$$

- Every surface has a triangulation.
- Compact** (surface): A surface with a triangulation with a finite number of triangles.
- A plane and torus minus a disk are both noncompact surfaces.
- Surfaces appear in knot theory in the space around a knot.
- Complement** (of the knot): The space around the knot $M = \mathbb{R}^3 - K$, where \mathbb{R}^3 is the three-dimensional space around the knot K .
- A splittable link can be defined as a link such that there exists a sphere in the link complement between the two link components.
- Every knot is contained within a torus. Some knots can be surrounded by a torus in unusual ways or by a genus $2+$ surface.
- We are particularly interested in the surfaces that cannot be simplified.
- Compressible** (surface): For a link L , a surface F where there exists “a disk D in $\mathbb{R}^3 - L$ such that D intersects F exactly in its boundary and its boundary does not bound another disk on F . Note that D is not allowed to intersect L ” (86).
- Compression:** The simplifying operation of cutting F along D , gluing D to each open curve, and using an isotopy to “smooth out” the surface.

- *Exercise 4.11:* Show that a compression always increases the Euler characteristic. Use this to show that the genus of the resulting surface or surfaces is always less than the genus of the original surface.

- Think of D as a triangle. During the compression, three vertices and three edges separate into two pairs of three vertices and three edges. This does not change the Euler characteristic.
- However, two faces are also added. This increases the Euler characteristic by 2.
- This change of +2 changes the genus by -1 .

- **Incompressible** (surface): A surface that is not compressible.

- Any time we have a composite knot, there exist an incompressible **swallow-follow torus**.

- **Swallow-follow torus:** A torus that encompasses (swallows) one factor knot and wraps around the strand of (follows) the other factor knot.

4.2 Surfaces with Boundary

- **Boundary component:** Discs (that may have been deformed) that have been removed from surfaces.
- The Euler characteristic of a surface with boundary is that of that surface without boundary minus the number of boundary components.
 - Think of adding each boundary component as removing the face of a triangle.
- **Capping off** (a surface with boundary): “Filling in boundary components by attaching disks” (88).
- *Exercise 4.13:* Find the Euler characteristics of each of the surfaces in Figure 4.5 without triangulating them.

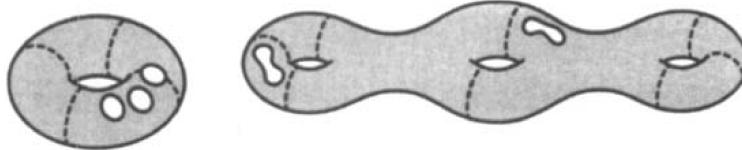


Figure 4.5: Surfaces with boundary.

- The left surface is a genus 1 surface with three boundary components. Therefore, it has $\chi = 0 - 3 = -3$.
- The right surface is a genus 3 surface with two boundary components. Therefore, it has $\chi = -4 - 2 = -6$.
- The Euler characteristic does not definitively distinguish surfaces with boundary.
- There are other ways, such as **orientability**.
- **Orientable** (surface): A surface sitting in \mathbb{R}^3 with “two sides that can be painted different colors, say black and white, so that the black paint never meets the white paint except along the boundary of the surface” (90).
 - For example, a disk and a torus are both orientable.
 - Furthermore, any torus with any number of boundary components is orientable.
- **Möbius band:** One of the simplest surfaces that is not orientable.
- “A surface is nonorientable if and only if it contains a Möbius band within it” (91).

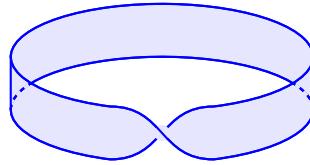


Figure 4.6: A Möbius band.

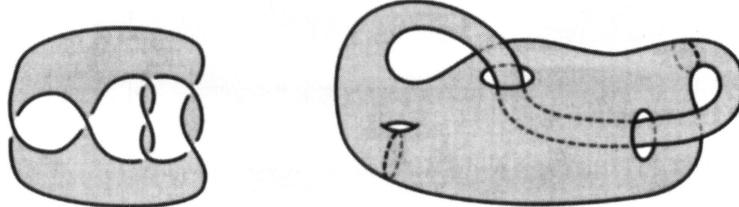


Figure 4.7: Possibly orientable surfaces.

- A Möbius band must have an odd number of half-twists in it.
- *Exercise 4.15:* Decide which of the two surfaces in Figure 4.7 is orientable and which is nonorientable.
 - The surface on the left is nonorientable while the one on the right is orientable.
- Three facts completely determine the homeomorphic type of a surface with boundary.
 1. “Is it orientable or nonorientable?
 2. How many boundary components does it have?
 3. What is its Euler characteristic?” (92).
- My techniques:
 - Choose a subdivision that has one chord for every twist between boundary components.
 - Alternatively, for simpler counting at the expense of bigger numbers, choose two chords at every twist.
 - Count edges along each boundary component and then count the chords.
- *Exercise 4.16:* Use the three criteria to identify the surface with boundary in Figure 4.8a.

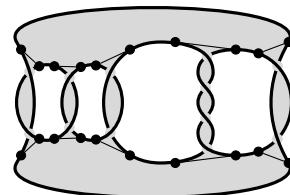
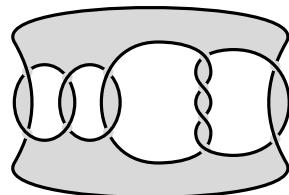


Figure 4.8: Identifying a surface with boundary.

- Begin by choosing the subdivision in Figure 4.8b according to the first technique above.
- Count $V = 20$, $E = 30$, and $F = 7$. Count $BC = 5$ where BC is the number of boundary components.

- Thus, $\chi = -3$. Adding +1 to χ for every boundary component, find χ for the original surface without boundary to be 2.
- Therefore, the surface without boundary corresponding to Figure 4.8a is the sphere.
- Define the genus of a surface with boundary to be that of the surface without boundary obtained by capping off each of its boundary components.
- Coming back to knot theory, one way to define knots is as the boundary of surfaces, i.e., the unknot could be defined as the boundary of a disk.
- **Annulus:** A sphere with two boundary components that lies outside one of the factor knots in a composite knot.
 - Note that this definition necessitates that the knot be thickened from a one-manifold.
 - Otherwise, an annulus would be a sphere with two punctures.
- Therefore, another definition of composite knot is “a knot such that there is a sphere in space punctured twice by the knot such that the knot is nontrivial both inside and outside the sphere” (93).
- **Conway sphere:** A sphere with four punctures enclosing a tangle.
 - This is the \mathbb{R}^3 analogy to the cyan circles in the projection plane in Figure 2.8.
 - If the knot is thickened, the Conway sphere becomes a sphere with four boundary components.

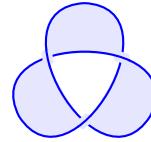


Figure 4.9: The trefoil as a Möbius band.

- Continuing with the idea that a disk defines the unknot, the trefoil knot (for instance) can be defined as the boundary of a Möbius band of three half-twists, as in Figure 4.9.
- Clearly, of particular interest are orientable surfaces with one boundary component such that said boundary component is a knot.

4.3 Genus and Seifert Surfaces

- There is one type of surface that appears in the complement of any knot.
 - In 1934, Herbert Seifert developed “an algorithm so that, given any knot, one can create an orientable surface with one boundary component such that the boundary circle is that knot” (95).
- Seifert’s algorithm is summarized as follows.
 - Take a knot and assign an orientation.
 - Eliminate each crossing by connecting the strand flowing in with the adjacent strand flowing out.
 - This produces a collection of **Seifert circles**.
 - Each circle bounds a disk in a plane. To avoid intersections, raise and lower the disks to different z -heights.
 - Finally, connect the disks to one another at the crossings of the knot by twisted bands.
- **Seifert circle:** A loop that can be deformed into a circle via a planar isotopy.

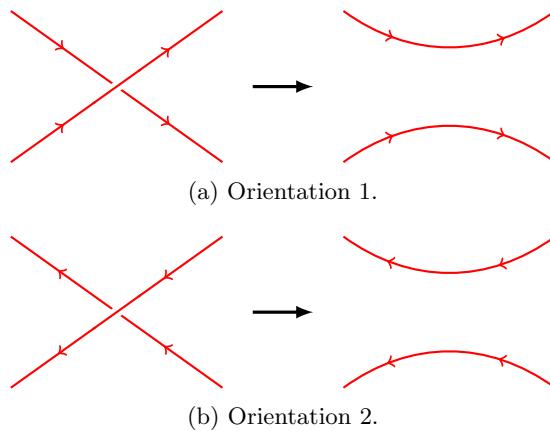


Figure 4.10: The relationship between orientation and crossing changes in Seifert's algorithm.

- *Exercise 4.18:* Show that the surface that we get doesn't depend on the direction we choose for the orientation on the knot.
 - At each crossing, two strands flow in and two flow out, regardless of the orientation, as in the left part of both subfigures in Figure 4.10. In effect, any crossing can be reworked to resemble the aforementioned parts of subfigures.
 - One of the strands adjacent to a strand flowing in always flows out and vice versa.
 - Because of this, reversing the orientation does not change the above fact (compare the left part of Figure 4.10a with that of Figure 4.10b).
 - Thus, whatever adjacent strands were paired up before will be paired up again (see the right side of Figures 4.10a and 4.10b).
 - Since the Seifert circles determine the final surface, the above proof that said circles will be unchanged verifies that the surfaces will remain unchanged.
- The surface is orientable because of the following thought experiment.
 - For any circle with a clockwise orientation, paint the upper face white and the lower face black.
 - For any circle with a counterclockwise orientation, paint the upper face black and the lower face white.
 - Any two circles above/below one another will have the same orientation, and, thus, the half-twist connects like colored faces.
 - Any two circles adjacent to each other will have opposite orientations, and, thus, the half-twist connects like colored faces.
 - Therefore, the surface is orientable.
- *Exercise 4.19:* Note that if we add an edge and two vertices to the surface across each crossing, cutting each band in half, we cut the surface up into valid faces. Use this fact to show that if c is the number of crossings and s is the number of Seifert circles, then $\chi = s - c$ and the genus of the surface is $g = \frac{c-s+1}{2}$.
 - Every 2-vertices-1-edge addition described above has a predictable effect on V and E .
 - Each vertex adds one new vertex (obviously), so two vertices are added for each 2-vertices-1-edge addition.
 - Each vertex splits an existing edge into two edges. Accounting for the addition of the connecting edge, three edges are added for each 2-vertices-1-edge addition.

- Since two vertices and three edges are added each time, the quantity $V - E$ changes by -1 with each 2-vertices-1-edge addition.
- Because, by the problem statement, a 2-vertices-1-edge addition occurs at each crossing, there are c such additions.
- Thus, $V - E = -c$ for the entire process.
- The final subdivision, on the other hand, reveals something about F .
 - Each face is separated from the others at the crossings.
 - Thus, since each face is composed of a Seifert circle plus a bit of twisting connection, the number of faces corresponds to the the original number of Seifert circles s , or $F = s$.
- Therefore, the sum of the derived values can be substituted into the sum $V - E + F$ to yield $-c + s$ or $s - c$.
- Since Equation 4.1 gives us $\chi = V - E + F$, the following equation holds by the transitive property.

$$\chi = s - c \quad (4.4)$$

- According to Equation 4.3, $\chi = 2 - 2g$, so $g = \frac{2-\chi}{2}$.
 - However, the genus of a surface with boundary is that of the surface without boundary acquired via capping off said surface with boundary.
 - Since the surface at hand is known to have specifically one boundary component, the χ that should presently be considered is the χ of the previous part plus one, or $\chi_1 = s - c + 1$.
 - Substitute: $g = \frac{2-(s-c+1)}{2}$.
 - Simplify.

$$\begin{aligned} g &= \frac{2 - (s - c + 1)}{2} \\ &= \frac{2 - s + c - 1}{2} \\ g &= \frac{c - s + 1}{2} \end{aligned} \quad (4.5)$$

- *Exercise 4.21:* Show that Seifert's algorithm always generates at least two Seifert circles.

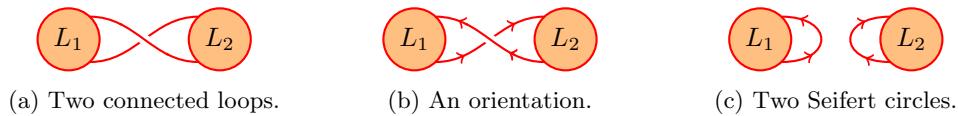


Figure 4.11: Generating Seifert circles from a crossing.

- In any knot with $c(K) > 0$, there exists a region with a crossing between two (possibly knotted) loops L_1 and L_2 , as in Figure 4.11a.
- When assigning an orientation to this crossing, the directions on the subarcs *must* be reversed.
 - If they both flowed from left to right (for example), that would imply that there is some point in L_1 from which the orientation “flows outward” in both directions and some point in L_2 to which the orientation “flows inward” from both directions.
 - This is contrary to the definition of an orientation, and, therefore, cannot be the case.
- Thus, the orientation in Figure 4.11b is valid.
- Given this orientation, the subarcs must separate into two Seifert circles, as in Figure 4.11c, according to Seifert's algorithm.
- This separation would hold for the opposite orientation, as demonstrated by Figure 4.10.

- It is possible to create multiple different-looking surfaces for the same knot by having a different starting projection.
- **Seifert surface** (for K): Given a knot K , “an orientable surface with one boundary component such that the boundary component of the surface is the knot K ” (99).
- **Genus** (of a knot): “The least genus of any Seifert surface for that knot” (99).
- The unknot has genus 0 (and is the only knot with genus 0).
- Seifert’s algorithm yields a Seifert surface of genus 1 for the figure-eight knot, and since the figure-eight knot is nontrivial, no Seifert surface of lower genus can exist for it. Therefore, the genus of the figure-eight knot is 1.
- *Exercise 4.22:* The twist knots are the knots obtained by linking together the two ends of a twisted loop. They include the trefoil and figure-eight knots. Show that all of the twist knots have genus 1.

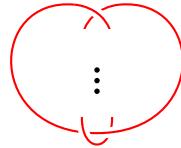


Figure 4.12: The general twist knot.

- This demonstration will follow the format of an induction proof. By Equation 4.5, if $g = 1$ for a knot, then $s = c - 1$. This argument will prove by induction that $s = c - 1$.
- Every twist knot can be projected in the mold of Figure 4.12, where the dots are replaced by some number of twists.
- For the trefoil knot, this number of twists is 0 and $s \leq c - 1$.
- Each new twist obviously adds one crossing.
- This new crossing, when the knot is assigned an orientation, will perfectly match one of the two left images in Figure 4.10.
 - This is because the orientation weaves its way down all of one of the two “vertical-ish” strands and then changes direction at the bottom loop to weave its way back up.
- Like in Figure 4.10, Seifert’s algorithm will yield one of the two products — separate loops.
- Thus, because every new crossing adds a new place to separate loops, both s and c increase by +1 in each subsequent twist knot.
- Therefore, since $s + 1 = c + 1 - 1$, s always equals $c - 1$.
- *Exercise 4.23:* Show that a minimal genus Seifert surface for a knot K must be incompressible (consider Euler characteristic).
 - Proof by contradiction: Take a compressible Seifert surface F for a knot K of minimal genus g_0 .
 - Compress it, adding +2 to the Euler characteristic χ_0 to get $\chi_1 = \chi_0 + 2$.
 - By Equation 4.3, the following holds.

$$\begin{aligned}\chi_0 &= 2 - 2g_0 \\ \chi_1 &= 2 - 2g_1 \\ \chi_0 + 2 &= 2 - 2g_1 \\ 2 - 2g_0 &= -2g_1 \\ g_1 &= g_0 - 1 \\ g_1 &< g_0\end{aligned}$$

- But since no genus g for K can be lower than g_0 by g_0 's definition as minimal, we have arrived at a contradiction.
- Therefore, the minimal genus Seifert surface must be incompressible.
- “Applying Seifert’s algorithm to an alternating projection of an alternating knot or link does yield a Seifert surface of minimal genus” (100).
- The following equation holds.
$$g(J \# K) = g(J) + g(K) \quad (4.6)$$
 - Clearly, $g(J \# K) \leq g(J) + g(K)$ because taking the composition of the two Seifert surfaces does not add any vertices, edges, or faces.
 - *Isotope surfaces, clean up intersections, perform surgeries to lower the number of intersection curves*^[16].
 - Arrive at $g(J) + g(K) \leq g(J \# K)$.
 - Therefore, $g(J \# K) = g(J) + g(K)$.
- Referencing Section 1.2, the trivial knot cannot be the composition of two nontrivial knots because $g(K) \geq 1$ where K is nontrivial, so by Equation 4.6, $g(J \# K) \geq 2 > 0 \Rightarrow g(J \# K)$ is nontrivial.
- *Exercise 4.27:* Show that the genus 1 knots are prime.
 - For a knot to be nontrivial, its genus must be greater than or equal to 1.
 - For a knot to be composite, its genus must be greater than or equal to 2.
 - Therefore, all knots with genus 1 have a high enough genus to be nontrivial (prime or composite) but not a low enough genus to not be composite (they are only prime).
- *Exercise 4.28:* Prove that if we take the composition of n copies of the same nontrivial knot, calling the result J_n , then as $n \rightarrow \infty$, $c(J_n) \rightarrow \infty$. Hint: Use the fact that there are only a finite number of knots with a given crossing number, but first use the Dowker notation to prove this fact.
 - For $c(K) = n$, the Dowker notation will consist of n even numbers ranging from 2 to $2n$. Therefore, there are n possibilities for the first number, $n - 1$ for the second number, and so on and so forth. This compresses down to $n!$ possible knots of $c(K) = n$.
 - Note that a large number of these knots will not be of this crossing number. This is an upper bound in the most literal sense.
 - One way to prove this would be to use Equation 4.6 to prove that the genus must increase with successive compositions. Since the genus is directly proportional to the quantity $c - s$ (by Equation 4.5), $c - s$ must increase with successive compositions. Since s cannot decrease with successive compositions (and certainly cannot decrease infinitely), c must increase with successive compositions.
 - I believe that the desired way to prove this is as follows. Composite knots are distinct from their factor knots. Generating a new composite knot creates a new knot. Since there are a finite number of knots of $c(K) = n$, eventually, all low-crossing possibilities (which will not all be reached, but regardless) must be used up. This forces $c(K)$ higher and higher as new knots are created.
- *Exercise 4.29:* Prove that if J_n is the composition of n copies of the same nontrivial knot, then $c(J_n) \geq n$. Hint: Use the Euler characteristic of a Seifert surface along with Equations 4.4 and 4.5 and *Exercise 4.21*. Subsequently, prove that $c(J_n) \geq 2n + 1$.
 - See the first solution to Exercise 4.28.

¹⁶Return to at a later date. Perhaps surgeries always increase Euler characteristic, lowering genus, until $g(J)$ and $g(K)$ are both at their minimum. Then their sum must be less than or equal to their composition?

- Seifert's algorithm's applicability to finding genus has been extended from alternating links to **alternative links**, which also include all **torus links** (see Section ??).
- There are some knots from which the minimal genus Seifert surface cannot be obtained using Seifert's algorithm on any projection of said knot.
- **Canonical genus** (of a knot): “The minimal genus of any Seifert surface obtained by applying Seifert’s algorithm to a projection of the knot” (106). *Also known as $g_c(K)$.*
- Genus was used to tell apart the Kinoshita-Terasaka mutants (see Section 2.3).