

## Week 9

# Symmetric Group Representation Characteristics

### 9.1 Frobenius Reciprocity; The Branching Theorem

11/27:

- Announcements.
  - OH on Wednesday at 5:30 PM this week; not Tuesday.
  - There will be extra OH next week pre-exam.
    - Roughly like Monday/Wednesday next week.
  - Midterm will be returned on Wednesday; we can pick them up in-person in his office starting then.
  - There are some grade boundaries: Pass/Fail we can do until Friday, withdrawal we can do until 5:00 PM today.
- Let's finish the conversation about induction/restriction and prove the **branching theorem**.
- Reminder to start.
  - We have two mathematical categories,  $G$ -reps and  $H$ -reps where  $H \leq G$ .
  - These categories are related by functors.
    - $\text{Res}_H^G : G\text{-reps} \rightarrow H\text{-reps}$  and vice versa for  $\text{Ind}_H^G$ .
    - See Figure 8.7.
  - Restrictions are stupidly simple.
  - Inductions, most hands-on, we take copies of  $W$  times cosets. Formulaically,

$$\text{Ind}_H^G W = g_1 W \oplus \cdots \oplus g_k W$$

where  $k = (G : H)$  and  $G = \bigsqcup_{i=1}^k g_i H$ .

- In more detail, the action of  $g$  on  $g_i w$  is that of  $g_{\sigma(i)} h_i w$ .
- This is a genuinely hard construction.
- A matrix of this thing will be a block-permutation matrix like

$$\begin{array}{l} g_1 W \\ g_k W \end{array} \left[ \begin{array}{c|c|c} g_1 W & & g_k W \\ \hline \text{////} & 0 & 0 \\ \hline 0 & 0 & \text{////} \\ \hline 0 & \text{////} & 0 \end{array} \right]$$

- As an alternate construction, we have that

$$g_1 W \oplus \cdots \oplus g_k W \cong \text{Hom}_H(\mathbb{C}[G], W)$$

- Recall that elements of the set on the right above are functions  $f : G \rightarrow W$  such that  $f(h(g)) = hf(g)$ .
  - We map between the two via  $f(g) \mapsto f(gx')$ .
- What is nice about induced representations is that  $\dim[\text{Ind}_H^G W] = (\dim W)[G : H]$ .
- There is a very easy statement of the character of an induced representation, the **Frobenius formula**.

- Recall that

$$\tilde{\chi}_W(g) = \begin{cases} 0 & g \notin H \\ \chi_W(g) & g \in H \end{cases}$$

- With this, we average:

$$\chi_{\text{Ind}_H^G W}(g) = \frac{1}{|H|} \sum_{x \in G} \tilde{\chi}_W(xgx^{-1})$$

- Essentially, we're taking a whole bunch of conjugates, summing them up, and dividing to get rid of overcounting.
- We now move onto **Frobenius reciprocity**, which is a relation between the functors/relations  $\text{Ind}_H^G$  and  $\text{Res}_H^G$ .

- The first point where category theory gets interesting is the notion of **adjoint functors**, which we are about to touch on. It is a very subtle notion.
- Here's version 1 of the statement of Frobenius reciprocity.

- Recall that we have a scalar product on the space of class function, given by

$$(\chi_1, \chi_2) = \frac{1}{|G|} \sum_{g \in G} \chi_1(g) \chi_2(g^{-1})$$

where  $\chi_1, \chi_2$  are class functions on  $G$ .

- Recall that if  $\chi_1 = \chi_V$  and  $\chi_2 = \chi_W$ , then

$$(\chi_1, \chi_2) = \dim \text{Hom}_G(V, W) = \dim \text{Hom}_G \left( \bigoplus_{i=1}^k V_i^{n_i}, \bigoplus_{i=1}^k V_i^{m_i} \right) = \sum_{i=1}^k n_i m_i$$

- Then the statement is as follows. If  $V$  is a  $G$ -rep and  $W$  is an  $H$ -rep, then

$$(V, \text{Ind}_H^G W)_G = (\text{Res}_H^G V, W)_H$$

➤ This notation denotes a scalar product in  $G$  and scalar product in  $H$  of the characters of each representation.

- This is similar to the relation between adjoint maps  $V \rightarrow W$  and  $W^* \rightarrow V^*$ .

- Version 2.

- We have that

$$\text{Hom}_G(V, \text{Ind}_H^G W) \cong \text{Hom}_H(\text{Res}_H^G V, W)$$

where the isomorphism is canonical.

- We will not check this last definition; we can tediously do it with definitions, and there's nothing complicated. Rudenko leaves this as an exercise to us.
  - The canonical isomorphism sends a map  $v \mapsto [g \mapsto \varphi(gv)]$  to the map  $\phi : V \rightarrow W$ .
- We now prove Version 1.

*Proof.* We have

$$\begin{aligned}
 (\chi_V, \chi_{\text{Ind}_H^G W})_G &= \frac{1}{|G|} \sum_{g_1 \in G} \chi_V(g_1) \chi_{\text{Ind}_H^G W}(g_1^{-1}) \\
 &= \frac{1}{|G|} \sum_{g_1 \in G} \chi_V(g_1) \left( \frac{1}{|H|} \sum_{g_2 \in G} \tilde{\chi}_W(g_2 g_1^{-1} g_2^{-1}) \right) \\
 &= \frac{1}{|H| \cdot |G|} \sum_{g_1, g_2 \in G} \chi_V(g_1) \tilde{\chi}_W(g_2 g_1^{-1} g_2^{-1}) \\
 &= \frac{1}{|H| \cdot |G|} \sum_{g_1, g_2 \in G} \chi_V(\underbrace{g_2 g_1 g_2^{-1}}_h) \tilde{\chi}_W(\underbrace{g_2 g_1^{-1} g_2^{-1}}_{h^{-1}}) \\
 &= \frac{1}{|H|} \frac{1}{|G|} \sum_{h \in G} |G| \chi_V(h) \tilde{\chi}_W(h^{-1}) \\
 &= (\chi_V|_H, \chi_W)_H \\
 &= (\text{Res}_H^G V, \chi_W)_H
 \end{aligned}$$

From line 4 to line 5: Fix  $h$ ; then  $g_2 g_1 g_2^{-1} = h$  iff  $g_1 = g_2^{-1} h g_2$ , so we have overcounted by  $|G|$  times. From line 5 to line 6:  $\tilde{\chi}_W$  is zero whenever  $h^{-1} \notin H$ , so this ostensible sum over all  $h \in G$  is *de facto* only a sum over all  $h \in H$ ; this is what allows us to consider  $\chi_V$  as “restricted to  $H$ ” in line 6.  $\square$

- We now come to the branching theorem at long last.
- Example first.
  - Consider  $S_n > S_{n-1}$ , where  $S_{n-1}$  is the subgroup that fixes  $n$ . I.e.,  $S_3 > S_2 = \{e, (12)\}$ , and we explicitly omit  $(13), (23), (123), (132)$  because they all move 3.
  - Let  $\lambda \vdash n$ .
  - Let  $\mu \leq \lambda$  be a Young diagram of a partition of  $n-1$ .
  - Then

1. We have

$$\text{Res}_{S_{n-1}}^{S_n} V_\lambda = \bigoplus_{\substack{\mu \leq \lambda \\ |\mu| = n-1}} V_\mu$$

■ Example:

$$\text{Res}_{S_4}^{S_5} \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array} = \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \end{array} \oplus \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \end{array}$$

2. We have

$$\text{Ind}_{S_{n-1}}^{S_n} V_\mu = \bigoplus_{\substack{\mu \leq \lambda \\ |\lambda| = n}} V_\lambda$$

■ Example:

$$\text{Ind}_{S_5}^{S_6} \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array} = \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & & \square \\ \hline \end{array} \oplus \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array} \oplus \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \square & \square & \square & \square \\ \hline \end{array}$$

- The reason that this theorem is called the branching theorem originates from the diagram in Figure 9.1, which (when continued) encapsulates the main idea of the theorem.
  - This graph helps you understand induction and restriction.
  - Dimensions are the number of paths from the bottom to a final Young diagram.

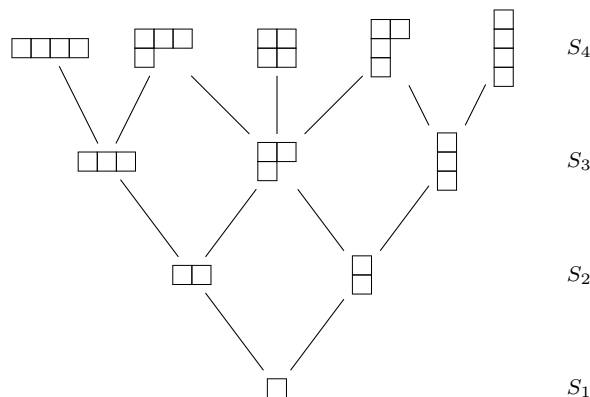


Figure 9.1: The branching theorem.

- For example, the dimension of  $(3, 1)$  is 3 because there are 3 paths to it, listed as follows.
  1.  $(1) \rightarrow (2) \rightarrow (3) \rightarrow (3, 1)$ .
  2.  $(1) \rightarrow (2) \rightarrow (2, 1) \rightarrow (3, 1)$ .
  3.  $(1) \rightarrow (1, 1) \rightarrow (2, 1) \rightarrow (3, 1)$ .
- Number of paths is equivalent to number of standard Young tableaux!
- Theorem (Branching): The following two statements are true.

$$\text{Res}_{S_{n-1}}^{S_n} V_\lambda = \bigoplus_{\substack{\mu \leq \lambda \\ |\mu|=n-1}} V_\mu \quad (9.1)$$

$$\text{Ind}_{S_{n-1}}^{S_n} V_\mu = \bigoplus_{\substack{\mu \leq \lambda \\ |\lambda|=n}} V_\lambda \quad (9.2)$$

*Proof.* We'll talk about the general idea of the proof now, and maybe do the details next time.

(9.1)  $\iff$  (9.2): Suppose first that the left statement above holds true. Then we have that

$$(\text{Res}_{S_{n-1}}^{S_n} V_\lambda, V_\mu) = \begin{cases} 0 & \lambda < \mu \\ 1 & \lambda \geq \mu \end{cases}$$

Thus, by Frobenius reciprocity,

$$(V_\lambda, \text{Ind}_{S_{n-1}}^{S_n} V_\mu) = (\text{Res}_{S_{n-1}}^{S_n} V_\lambda, V_\mu) = \begin{cases} 0 & \lambda < \mu \\ 1 & \lambda \geq \mu \end{cases}$$

Therefore, the second statement holds true. The proof is symmetric in the opposite direction.

(9.1): Let's look at an example. Consider the Young diagram of  $S_8$  shown in Figure 9.2.



Figure 9.2: Proving the branching theorem.

We want to restrict it down to  $S_7$ . Recall that  $V_\lambda = \text{span}(S_8 : \Delta(x_1, x_2, x_3)(x_4 - x_5)(x_6 - x_7))$ . Now in  $S_7$ , we fix  $x_8$ . Consider subrepresentations of  $V_\lambda$  filtered by degree as follows.

$$\underbrace{\begin{bmatrix} \square \\ \square \\ \square \\ \square \end{bmatrix}}_{\deg_{x_3} \leq 0} \leq \underbrace{\begin{bmatrix} \square \\ \square \\ \square \end{bmatrix}}_{\deg_{x_5} \leq 1} \leq \underbrace{\begin{bmatrix} \square \\ \square \\ \square \end{bmatrix}}_{\deg_{x_8} \leq 2} \leq V_\lambda$$

The proof comes from the fact that if we now take quotients of these subrepresentations, e.g., via

$$\deg = 0, \deg \leq 1 / \deg \leq 0, \deg \leq 2 / \deg \leq 1, \dots$$

then since  $x_8$  can only appear in three boxes, ... □

- Practice with the above example and think it through.

## 9.2 The Character of a Symmetric Group Representation

11/29:

- Announcements.
  - OH today at 5:30.
  - Our midterms are graded; we can look at them in his office whenever (I can do this during OH!).
- Today, we'll formulate the main result he wants to prove next time.
- Goal is still to understand representations of  $S_n$ .
  - We've constructed all of them using Specht modules, but what else do we want?
  - We have dimension, we want characters, etc.
- The main idea is to look at symmetric polynomials once again.
  - Consider  $\mathbb{Q}[x_1, \dots, x_n]^{S_n}$ .
  - We have proven the fundamental theorem that  $\mathbb{Q}[x_1, \dots, x_n]^{S_n} = \mathbb{Q}[\sigma_1, \dots, \sigma_n]$  where  $\sigma_k = \sum_{1 \leq i_1 \leq \dots \leq i_k \leq n} x_{i_1} \cdots x_{i_k}$ .
  - We also proved in PSet 6, Q6 that these rings are equal to  $\mathbb{Q}[p_1, \dots, p_k]$  and  $\mathbb{Q}[h_1, \dots, h_k]$  where

$$p_k = \sum_{i=1}^n x_i^k \qquad h_k = \sum_{1 \leq i_1 \leq \dots \leq i_k \leq n} x_{i_1} \cdots x_{i_k}$$

- Example: If  $n = 3$  and  $k = 2$ , then

$$h_2 = x_1^2 + x_2^2 + x_3^2 + x_1x_2 + x_1x_3 + x_2x_3$$

- Table of bases for  $n, k$ .

$k \setminus n$	1	2	3	4
0	1	1	1	1
1	$x_1$	$x_1 + x_2$	$x_1 + x_2 + x_3$	$\dots$
2	$x_1^2$	$x_1^2 + x_2^2, x_1x_2$	$\sigma_1^2, \sigma_2^2$	$\sigma_1^2, \sigma_2^2$
3	$x_1^3$	...		

Table 9.1: Polynomial bases.

- Now take

$$\Lambda_k = [\mathbb{C}[x_1, \dots, x_k]]_{\deg=k-1} \cong [\mathbb{C}[x_1, \dots, x_{k+1}]]_{\deg=k-1} \cong \dots$$

- Alternatively, we can think of this thing as

$$\Lambda_k = (\mathbb{C}[x_1, \dots])_k$$

with  $\sigma_1^k, \sigma_2\sigma_1^{k-1}, \dots$

- We call  $\Lambda$  the ring of symmetric functions and define it to be equal to

$$\Lambda = \mathbb{Q}[\sigma_1, \sigma_2, \sigma_3, \dots]$$

- In every complete component, only finitely many of the  $\sigma$  will participate, so we get finite things.
- This is a graded ring! We have

$$\Lambda = \bigoplus_{k \geq 0} \Lambda_k$$

and  $\Lambda_k \otimes \Lambda_\ell = \Lambda_{k+\ell}$

- This construction is called the **projective limit**, and we may have encountered it in commutative algebra under the definition

$$\Lambda = \varprojlim \mathbb{C}[x_1, \dots, x_n]^{S_n}$$

- We have identities such as  $p_2 = \sigma_1^2 - 2\sigma_2$ . This means that

$$(x_1 + \dots + x_n)^2 - 2(x_1x_2 + x_1x_3 + \dots) = x_1^2 + x_2^2 + \dots$$

- Observation:  $\dim_{\mathbb{Q}} \Lambda_n$ .

- Now, we need to take a vector space on ring representations; we've done this already with the representation ring.
- Let  $R_n$  be the  $\mathbb{Q}$ -vector space of functions  $\chi : S_n \rightarrow \mathbb{Q}$  such that  $\chi(x\sigma x^{-1}) = \chi(\sigma)$ . This is our favorite space of class functions.
- Theorem (Frobenius characteristic map): There is an isomorphism of vector spaces and of rings called the Frobenius characteristic:  $\text{ch} : \bigoplus_{n \geq 0} R_n \rightarrow \Lambda$ .

*Proof.* Take  $\chi_V \in R_k$ , and  $\chi_W \in R_\ell$ . Let  $V$  an  $S_k$ -rep, and  $W$  an  $S_\ell$ -rep. We know that

$$S_k \times S_\ell = S_{k+\ell}$$

So what we can do is induction  $\text{Ind}_{S_k \times S_\ell}^{S_{k+\ell}} (V \otimes W)$ . Call this operation  $\chi_V \boxtimes \chi_W$ .

Now we write down the formula:

$$\text{ch}(\chi) = \frac{1}{n!} \sum_{\sigma \in S_n} \chi(\sigma) p_1^{\lambda_1(\sigma)} \dots p_k^{\lambda_k(\sigma)}$$

where  $\lambda_1(\sigma), \lambda_2(\sigma), \dots$  represent the cycle structure of  $\sigma$ ; each  $\lambda_i$  is a number of cycles of length  $1, 2, \dots$ .  $\square$

- Examples.

1.  $S_1$ .

- Sends the YD (1) to  $p_1 = x_1 + x_2 + x_3 + \dots$ .

2.  $S_2$ .

- Sends (2) to  $\frac{1}{2!}(p_1^2 + p_2) = \frac{1}{2}((x_1 + x_2)^2 + x_1^2 + x_2^2) = x_1^2 + x_2^2 + x_1x_2 = h_2$ .
- It also sends (1, 1) to  $\frac{1}{2!}(p_1^2 - p_2) = \frac{1}{2}((x_1 + x_2)^2 - x_1^2 - x_2^2) = x_1x_2 = \sigma_2$ .
- Let's check our formula. What is  $\text{Ind}_{S_1 \times S_1}^{S_2} (1) \otimes (1)$ ? Since the induction of the trivial representation is the regular representation, which we can decompose, we know that this induction equals  $(1, 1) \oplus (2)$ . It follows that  $p_1^2 = x_1^2 + x_2^2 + x_1x_2 + x_1x_2 = (x_1 + x_2)^2$ .

3.  $S_3$ .

- Sends (3) to

$$\begin{aligned}\frac{1}{3!}(p_1^3 + 3p_1p_2 + 2p_3) &= \frac{1}{6}[(x_1 + x_2 + x_3)^3 + 3(x_1 + x_2 + x_3)(x_1^2 + x_2^2 + x_3^2) + 2(x_1^3 + x_2^3 + x_3^3)] \\ &= \frac{1}{6}[6(x_1^3 + x_2^3 + x_3^3) + 6(x_1^2x_2 + x_1x_2^2 + x_1x_3^2 + x_1^2x_3 + \cdots) + 6x_1x_2x_3] \\ &= h_3\end{aligned}$$

- Sends (1, 1, 1) to

$$\begin{aligned}\frac{1}{3!}(p_1^3 - 3p_1p_2 + 2p_3) &= \frac{1}{6}[(x_1 + x_2 + x_3)^3 - 3(x_1 + x_2 + x_3)(x_1^2 + x_2^2 + x_3^2) + 2(x_1^3 + x_2^3 + x_3^3)] \\ &= x_1x_2x_3 \\ &= \sigma_3\end{aligned}$$

- Sends (2, 1) to

$$\begin{aligned}\frac{1}{3!}(2p_1^3 - p_3) &= \frac{1}{6}[2(x_1^3 + x_2^3 + x_3^3) + 6(x_1^2x_2 + \cdots) + 12x_1x_2x_3] \\ &= (x_1^2 + \cdots) + 2x_1x_2x_3\end{aligned}$$

- Again, we can check that

$$\text{Ind}_{S_2 \times S_1}^{S_3}[(1, 1) \otimes (1)] = \sigma_1\sigma_2$$

- We compute  $\text{Ind}_{S_2}^{S_3}(1, 1) = (1, 1, 1) \oplus (2, 1)$  via the branching formula: There are only two ways to add a box!
- We have  $\sigma_1\sigma_2 - \sigma_3 = (x_1 + x_2 + x_3)(x_1x_2 + x_1x_3 + x_2x_3) - x_1x_2x_3$ .

- Do we need to be fluent in the techniques by which you expanded all of the polynomials above??
- Thus, we have two conjectures:

$$\text{ch}[(n)] = h_n \qquad \text{ch}[\underbrace{(1, \dots, 1)}_{n \text{ times}}] = \sigma_n$$

- The theorem is cool because it sends all of representation theory to some symmetric polynomial game!
- How do we compute  $\text{ch}(V_\lambda)$ ?
  - We say it equals  $S_\lambda$ , where  $S_\lambda$  is a Schur polynomial.
  - Take the YT of  $\lambda$ . Recall standard YTs.
  - **Semistandard** (YT): Things strictly increase in columns, but only monotonically increase in rows. *draw picture!*
  - The six semistandard ones give us the Schur polynomial.
  - Relation to RSK correspondence.
- Proving why this stuff is true is not hard.
- To understand *why* this is true, Google the **Schur-Weyl duality**.

## 9.3 Office Hours

- I got a 68/100 on the midterm: 30, 24, 0, 14.
  - I would have needed to show my work (or at least one example of a calculation) to get full credit for 2, even though it just said “find.”
  - Rudenko did not expect that finding conjugacy classes would be so difficult for us; he will adjust for this difficulty on the final.
- Week 3, Lecture 2: You proved that  $\langle \chi_V, \chi_W \rangle = \delta_{VW}$ . To do so, you used a projection function  $p = (1/|G|) \sum_{g \in G} gv$ . You began your proof by proving that  $p$  is a  $G$ -morphism and then never used this result again, as far as I can tell. Did you use it again? See pp. 45-47, 58 (it needs to be a morphism of  $G$ -representations to map between the representations  $V, V^G$ ?).
- Week 3, Lecture 2: Same proof. To prove that  $\text{Im}(P) = V^G$ , do we need more than  $p^2 = p$ ? I think so, but you didn't do it explicitly. See pp. 46-47.
- Week 3, Lecture 2: Same proof. What's up with the trivial special case? See p. 48.
- \*Week 3, Lecture 3: Cube thing (see picture from 10/13)?
  - It's just a depiction of two different 3-coordinate bases of the same space. It was drawn to illustrate a possible relation between the orthonormal basis  $\chi_1, \chi_2, \chi_3$  (cube) and the orthogonal basis  $\chi_{C_1}, \chi_{C_2}, \chi_{C_3}$ .
- Week 3, Lecture 3: Why did we talk about the infinite-dimensional regular representation here? See p. 50.
- \*Week 3, Lecture 3: What is the point of the misc. calculations involved in computing the  $S_4$  character table? See p. 52.
  - Just to check that we were on the right path and shown an example of using the orthogonality relations.
- \*Week 3, Lecture 3: Proof of the second orthogonality relation your way? It's in Serre (1977), but I don't think that's the way you proved it. See p. 52.
  - To begin, note that it is a *highly* nontrivial statement that if  $A, B$  are matrices such that  $AB = I$ , then  $BA = I$ . It seems so simple to us, but think about it! For an arbitrary matrix  $A, B$ ,  $AB$  looks nothing like  $BA$ ! We have two entirely different systems of equations.
  - However, using this fact, basically it is possible to translate the orthogonality relation for the *columns* into the orthogonality relation about the *rows*.
- \*Week 3, Lecture 3: All the talk about the exceptional homomorphisms? See p. 52, 61 (the final representation has something to do with an **involution** of trace 2, and is a representation of a quotient group?).
  - So the representation is  $\rho : S_4 \twoheadrightarrow S_3 \xrightarrow{\tilde{\rho}} GL_n$ , where  $\tilde{\rho} : S_3 \rightarrow GL_n$  is the representation of  $\rho$  corresponding to the character  $(2, 0, 1)$ .
- \*Week 4, Lecture 1: Alternate construction of  $R(G)$ ? See p. 63.
- \*Week 4, Lecture 1: Extension of scalars with the representation ring? See p. 64.
  - We don't need to know anything about this stuff.
  - What it is though is basically analogous to extending the real numbers into a subset of the complex numbers by treating every  $x \in \mathbb{R}$  as  $x + 0i \in \mathbb{C}$ . Very trivial, silly concept.
  - There is also such a thing as a **reduction of scalars**.



- \*Week 4, Lecture 1: Does multiplying a column vector in the basis  $\{\delta_{C_i}\}$  by the character table put it in the basis  $\{\chi_{V_i^*}\}$ , or vice versa? See p. 66.

- Derive it for yourself.
- Example: Consider the character table for  $S_3$  (Table 3.1) represented as the following matrix.

$$A = \begin{bmatrix} 1 & 1 & 1 \\ 1 & -1 & 1 \\ 2 & 0 & 1 \end{bmatrix}$$

- Denote the conjugacy classes of  $S_3$  by  $e, (xx), (xxx)$ .
- Interpretation 1.

- We see that the standard representation is the class function sending

$$e \mapsto 2 \qquad (xx) \mapsto 0 \qquad (xxx) \mapsto -1$$

- Additionally, we have

$$\begin{array}{llll} e \mapsto 1 & (xx) \mapsto 0 & (xxx) \mapsto 0 & (\delta_e) \\ e \mapsto 0 & (xx) \mapsto 1 & (xxx) \mapsto 0 & (\delta_{(xx)}) \\ e \mapsto 0 & (xx) \mapsto 0 & (xxx) \mapsto 1 & (\delta_{(xxx)}) \end{array}$$

- Thus, we can express  $\chi_{(2,1)}$  as a linear combination of the  $\delta$ 's via

$$\begin{aligned} \chi_{(2,1)} &= (2)\delta_e + (0)\delta_{(xx)} + (-1)\delta_{(xxx)} \\ &= \chi_{(2,1)}(e)\delta_e + \chi_{(2,1)}(xx)\delta_{(xx)} + \chi_{(2,1)}(xxx)\delta_{(xxx)} \\ &= \sum_{C_i} \chi_{(2,1)}(C_i)\delta_{C_i} \end{aligned}$$

- It follows in particular that if we represent the  $\delta_{C_i}$ 's as the standard column vector basis of  $\mathbb{C}^3$ , then

$$\chi_{(2,1)} = A^T \delta_{(xxx)}$$

- Interpretation 2.

- If we multiply  $A$  by the column vector equal to each representation weighted by  $|C_i|$ , then we recover the  $\delta$  basis:

$$\begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} = A \begin{bmatrix} 1 \cdot 1 \\ 3 \cdot 1 \\ 1 \cdot 1 \end{bmatrix} \qquad \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} = A \begin{bmatrix} 1 \cdot 1 \\ 3 \cdot -1 \\ 1 \cdot 1 \end{bmatrix} \qquad \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} = A \begin{bmatrix} 1 \cdot 2 \\ 3 \cdot 0 \\ 1 \cdot -1 \end{bmatrix}$$

- This interpretation is also what is expressed by the following formula from Lecture 4.2.

$$\delta_{C_j}(g) = \sum_{V_i} \frac{|C_j| \bar{\chi}_{V_i}(g)}{|G|} \cdot \chi_{V_i}(g)$$

- \*Week 4, Lecture 2: Isotypical components example. See p. 68.
- \*\*Week 4, Lecture 3: Proof the  $\mathbb{C}$  is the only finite-dimensional division algebra? See p. 71.
  - Let  $A$  be an arbitrary finite-dimensional division algebra over  $\mathbb{C}$ .
  - To prove that  $A = \mathbb{C}$ , we will use a bidirectional inclusion proof.
  - Naturally,  $\mathbb{C} \subset A$ .
  - To prove the reverse implication, start by letting  $a \in A$  be arbitrary.
  - Define the left-multiplication operator  $L_a : A \rightarrow A$  by  $x \mapsto ax$  for all  $x \in A$ .

- Recall that  $A$  is a complex vector space in addition to being an algebra, the same way a ring is also a group. Thus,  $L_a$  is a linear operator on a complex vector space.
- It follows by the theorem of linear algebra that  $L_a$  has an eigenvalue  $\lambda \in F = \mathbb{C}$  and corresponding eigenvector  $b \in A$ .
- Consequently, by the cancellation lemma,

$$L_a b = \lambda b$$

$$ab = \lambda b$$

$$a = \lambda$$

- Therefore,  $a \in A$  implies  $a \in \mathbb{C}$ , so  $A = \mathbb{C}$ .
- \*Week 6, Lecture 2: Proof that  $\sqrt{2}/2$  is not an algebraic integer using Gauss's lemma? See p. 87.
  - Let  $\alpha := \sqrt{2}/2$  for the sake of notation.
  - Suppose for the sake of contradiction that  $\alpha$  is an algebraic integer.
  - Then there exists a monic polynomial  $p(x) \in \mathbb{Z}[x]$  such that  $p(\alpha) = 0$ .
  - Observe that the minimal polynomial in  $\mathbb{Z}[x]$  that annihilates  $\alpha$  is  $2x^2 - 1$ .
  - Thus, by polynomial division,

$$p(x) = q(x) \cdot (2x^2 - 1) + r(x)$$

for some  $q, r \in \mathbb{Q}[x]$  such that  $\deg r \leq 2 - 1$ .

- We have that

$$r(\alpha) = p(\alpha) - q(\alpha) \cdot (2\alpha^2 - 1) = 0 - q(\alpha) \cdot 0 = 0$$

- Additionally, since  $r \in \mathbb{Q}[x]$  and  $\deg r \leq 1$ , we know that  $r(x) = ux + v$  for some  $u, v \in \mathbb{Q}$ .
- We now prove that  $u = v = 0$ .
  - Suppose for the sake of contradiction that either  $u$  or  $v$  was not equal to zero.
  - Combining the previous two claims reveals that

$$0 = r(\alpha)$$

$$= u\alpha + v$$

$$-\frac{v}{u} = \alpha$$

- If  $u = 0$ , then  $\alpha$  is undefined and we have arrived at a contradiction. Thus,  $u \neq 0$ .
  - Thus,  $\alpha \in \mathbb{Q}$ . But since  $\alpha \notin \mathbb{Q}$  by definition, we have arrived at a contradiction.
  - Therefore,  $u = v = 0$ .
- Having established that  $r = 0$ , we know that  $p = (2x^2 - 1)q$ , i.e.,  $2x^2 - 1$  divides  $p$ .
- Now define  $N$  to be the least common multiple of the denominators of the coefficients of  $q$ .
- Consider

$$Np = (Nq)(2x^2 - 1)$$

- It follows by Gauss's lemma that

$$c(Np) = c[(Nq)(2x^2 - 1)]$$

$$N = c(Nq) \cdot c(2x^2 - 1)$$

$$= 1 \cdot 1$$

$$= 1$$

where  $c$  denotes the **content**.

- But if  $N = 1$ , then  $q \in \mathbb{Z}[x]$ , so leading term of  $p$  — equal to the product of  $2x^2$  and the leading term of  $q$  — has a coefficient that is a multiple of 2, i.e., is *not* equal to 1 as is required of a monic polynomial, a contradiction.
- \*Week 6, Lecture 3: Questions about Lemma 1 of the proof of Burnside's theorem. See p. 92.
  - The roots  $a_1, \dots, a_k$  of the minimal polynomial of the algebraic integer  $a$  are known as **conjugate algebraic integers**.
  - The conjugate algebraic integers of a root of unity are also roots of unity.
    - Suppose  $\varepsilon$  is a root of unity.
    - Then the minimal polynomial of  $\varepsilon$  is  $x^n - 1$  for some  $n \in \mathbb{N}$ .
    - Naturally, the roots of this polynomial (the conjugate algebraic integers to  $\varepsilon$ ) are all of the other roots of unity of order  $n$ .
  - The conjugate algebraic integers of a sum of roots of unity is a sum of roots of unity.
    - It can be shown that the minimal polynomial for  $\varepsilon_1 + \varepsilon_2$  is

$$p(x) = \prod_{i,j=1}^n (x - \varepsilon_1^i - \varepsilon_2^j)$$

- Evidently, the above polynomial is symmetric under permutations of  $\varepsilon_1^i, \varepsilon_2^j$ , and we'd generate the same polynomial with any  $\pm \varepsilon_1^i \pm \varepsilon_2^j$  as starting material.
- Explicit example.
  - $\pm\sqrt{2}$  are conjugate algebraic integers, as solutions to  $x^2 - 2$ . Similarly,  $\pm\sqrt{3}$  are conjugate algebraic integers as solutions to  $x^2 = 3$ .
  - Thus, we expect the minimal polynomial for  $\sqrt{2} + \sqrt{3}$  to be
$$p(x) = (x - \sqrt{2} - \sqrt{3})(x - \sqrt{2} + \sqrt{3})(x + \sqrt{2} - \sqrt{3})(x + \sqrt{2} + \sqrt{3})$$
  - Expanding, we obtain
$$\begin{aligned} p(x) &= (x^2 - (\sqrt{2} + \sqrt{3})^2)(x^2 - (\sqrt{2} - \sqrt{3})^2) \\ &= x^4 - [(\sqrt{2} + \sqrt{3})^2 + (\sqrt{2} - \sqrt{3})^2]x^2 + (\sqrt{2} + \sqrt{3})^2(\sqrt{2} - \sqrt{3})^2 \\ &= x^4 - 10x^2 + 1 \end{aligned}$$
  - Indeed, the above polynomial is a monic polynomial
  - From the definition, this polynomial is evidently also the minimal polynomial for  $\sqrt{2} - \sqrt{3}$ ,  $-\sqrt{2} + \sqrt{3}$ , and  $-\sqrt{2} - \sqrt{3}$ .
  - Thus, the conjugate algebraic integers of  $\sqrt{2} + \sqrt{3}$  are the four sums of all individual algebraic integers.
- How do we extend this argument to the case in the problem?? What about when  $\varepsilon_1 = -1$  and  $\varepsilon_2 = i$  so that simple powers don't access every combination as the  $p(x)$  formula does?
- We know that  $\prod_{i=1}^n a_i \in \mathbb{Z}$  because of Vieta's formula.
  - In particular, this tells us that  $x_1 \cdots x_n$  is equal to the last coefficient in the minimal polynomial which, by definition, is an integer.
- Don't worry too much about all this, though: Burnside's theorem will no longer be on the final because Rudenko changed his mind.
- \*Week 7, Lecture 2: Symmetric polynomials and roots of symmetric polynomials. See p. 101.
- \*Week 7, Lecture 2: Word in blackboard picture? See p. 102.
  - “Remain” to show...

- \*Week 7, Lecture 2: What is  $d$  in the proof of the alternating polynomials theorem? See p. 103.
  - $d = n - 1$ .
- \*Week 8, Lecture 2: What is  $d$  in the definition on p. 111.
  - Consider the Specht polynomial corresponding to  $(2, 2, 1)$ .


 Figure 9.3: Young diagram for  $(2, 2, 1)$ .

- Since  $(2, 2, 1)' = (3, 2)$ , the Specht polynomial is  $(x_1 - x_2)(x_1 - x_3)(x_2 - x_3) \cdot (x_4 - x_5)$ .
- $\Delta_{123}$  is of degree  $3 = \binom{3}{2}$  because looking at the first column of the YD, which corresponds to  $\lambda'_1 = 3$ , out of the 3 boxes, we must choose 2 for each of the three terms  $(x_1 - x_2), (x_1 - x_3), (x_2 - x_3)$ . Then we just add this to the 2 choose 2 for the second column of the YD.
- \*Week 9, Lecture 2: Do we need to be fluent in the techniques you used to expand and reduce the various polynomial powers? How did you do that again?

## 9.4 The Frobenius Characteristic Map

12/1:

- Proving the theorem.
- The statement is that there exists a function

$$\text{ch} : \underbrace{\bigoplus_{n \geq 0} \mathbb{Q}_{\text{cl}}(S_n)}_{\{f: S_n \rightarrow \mathbb{Q} : f(\sigma = \sigma^{-1}) = f(i)\}} \rightarrow \bigoplus_{n \geq 0} \Lambda_n$$

where  $\Lambda_n = \mathbb{Q}[\sigma_1, \sigma_2, \sigma_3, \dots]_{\deg=n} = [\mathbb{Q}[x_1, \dots, x_n]^{S_n}]_{\deg=n}$ . Note that by convention,  $\Lambda_0 = \mathbb{Q}$ . This function is given by

$$\text{ch}(\chi) = \frac{1}{n!} \sum_{\sigma \in S_n} \chi(\sigma) \cdot p_1^{\lambda_1(\sigma)} \cdot p_n^{\lambda_n(\sigma)}$$

where  $p_k = x_1^k + x_2^k + \dots$  and  $\lambda_i(\sigma)$  is the number of cycles in  $\sigma$  of length  $i$ . Moreover,  $\text{ch}$  is an isomorphism of  $\mathbb{Q}$ -algebras. To create a product of  $V$  a  $S_n$ -rep and  $W$  an  $S_m$ -rep, we map

$$V \boxtimes W = \text{Ind}_{S_n \times S_m}^{S_{n+m}} (V \otimes W)$$

*Proof.* To prove that  $\text{ch}$  is a ring isomorphism, we need...

1.  $\boxtimes$  is associative;
2.  $\text{ch}(\chi_1 \boxtimes \chi_2) = \text{ch}(\chi_1) \cdot \text{ch}(\chi_2)$ ;
3.  $\text{ch}((n)) = h_n$ .

1,2,3 imply the theorem because 2,3 imply that  $\text{ch}$  is surjective,  $\Lambda_n$  has a  $\mathbb{Q}$ -basis  $h_{\lambda_1} \cdots h_{\lambda_n}$  for  $\lambda_1 \leq \dots \leq \lambda_n$  and  $\lambda_1 + \lambda_2 + \dots + \lambda_n = n$ . For example, for  $\Lambda_5$ , we have  $h_1^5, h_1^3 h_2, h_1 h_2^2, h_2^2 h_3, h_1^2 h_3, h_1 h_4, h_5$ . This surjectivity combined with the fact that  $\dim \mathbb{Q}_{\text{cl}}[S_n] = \dim \Lambda_n$  implies that  $\text{ch}$  is an isomorphism of rings.

Last thing:  $\text{ch}((n)) = \frac{1}{n!} \sum p_1^{c_1(\sigma)} \cdots p_n^{c_n(\sigma)}$  where  $c_i(\sigma)$  denotes the number of cycles of length  $i$  and hence  $\sum i c_i = n$ . Denote  $p_1^{c_1(\sigma)} \cdots p_n^{c_n(\sigma)}$  by  $p^{c(\sigma)}$ .

*Proof.* Let

$$\begin{aligned}
 \sum h_n t^n &= \sum \left( \sum_{i_1 \leq \dots \leq i_n} x_{i_1} \cdots x_{i_n} t^n \right) \\
 &= \frac{1}{1-x_1 t} \cdot \frac{1}{1-x_2 t} \cdots \frac{1}{1-x_n t} \\
 &= \exp \left( \log \left( \prod_{i=1}^n \frac{1}{1-x_i t} \right) \right) \\
 &= \exp \left( \sum_{i=1}^n -\log(1-x_i t) \right) \\
 &= \exp \left( x_1 + \frac{x_1^2 t^2}{2} + \frac{x_1^3 t^3}{3} + \cdots + x_2 + \frac{x_2^2 t^2}{2} + \cdots \right) \\
 &= \exp \left( p_1 + \frac{p_2 t^2}{2} + \frac{p_3 t^3}{3} + \cdots \right) \\
 &= \prod_{m \geq 1} \exp \left( \frac{p_m t^m}{m} \right) \\
 &= *
 \end{aligned}$$

We get the second equality because each  $1/(1-x_i t) = 1 + x_i + x_i^2 t^2 + \cdots$ . We need the power series  $-\log(1-t) = t + t^2/2 + \cdots$  and  $\exp(t) = 1 + t + t^2/2! + \cdots$ . Thus,  $\exp(\log(1-t)) = 1-t$ . Now note that

$$\begin{aligned}
 \sum_{n \geq 0} \text{ch}[(n)] \cdot t^n &= \sum_{n \geq 0} \frac{t^n}{n!} \sum_{\sigma \in S_n} p_1^{c_1(\sigma)} \cdots p_n^{c_n(\sigma)} \\
 &= \sum_{n \geq 0} \frac{t^n}{n!} \left( \sum_{a_1 + 2a_2 + \cdots + na_n = n} p_1^{a_1} \cdots p_n^{a_n} \right) \cdot \frac{n!}{1^{a_1} a_1! 2^{a_2} a_2! \cdots n^{a_n} a_n!} \\
 &= \sum_{\substack{n \geq 0 \\ a_1, \dots, a_n : a_1 + 2a_2 + \cdots + na_n = n}} \frac{1}{a_1!} \left( \frac{p_1}{1} \right)^{a_1} t^{a_1} \frac{1}{a_2!} \left( \frac{p_2}{2} \right)^{a_2} t^{a_2} \cdots \frac{1}{a_n!} \left( \frac{p_n}{n} \right)^{a_n} t^{a_n} \\
 &= \prod_{k=1}^n \left( \sum_{a_k=1}^{\infty} \frac{(p_k)^{a_k} t^{a_k k}}{a_k! k^{a_k}} \right) \\
 &= \prod_{k \geq 1} \exp \left( \frac{p_k t^k}{k} \right) \\
 &= *
 \end{aligned}$$

We're overcounting because we can cyclically permute cycles (i.e.,  $(12) = (21)$ ), hence the correction factor in the second line above.

Note: This exponent/logarithm trick is a common computational trick in combinatorics, varieties, etc.  $\square$

Now we prove the part 3, i.e., that  $\boxtimes$  is associative. We do this by direct computation.

$$\underbrace{\text{Ind}_{S_{n+m} \times S_\ell}^{S_{n+m+\ell}} \left[ \text{Ind}_{S_n \times S_m}^{S_{n+m}} (\chi_1 \otimes \chi_2) \right] \otimes \chi_3}_{(\chi_1 \boxtimes \chi_2) \boxtimes \chi_3} = \text{Ind}_{S_n \times S_m \times S_\ell}^{S_{n+m+\ell}} (\chi_1 \otimes \chi_2 \otimes \chi_3)$$

...

Then proving 2 (homomorphism bit) is the hardest. We have

$$\begin{aligned}
 \text{ch}(\text{Ind}_{S_n \times S_m}^{S_{n+m}} (\chi_1 \otimes \chi_2)) &= \frac{1}{n!} \sum_{\sigma \in S_n} (\text{Ind}_{S_n \times S_m}^{S_{n+m}} \chi_1 \otimes \chi_2)(\sigma) \underbrace{p_1^{c_1(\sigma)} \cdots p_{n+m}^{c_{n+m}(\sigma)}}_{\psi} \\
 &= \left\langle \text{Ind}_{S_n \times S_m}^{S_{n+m}} (\chi_1 \otimes \chi_2), \psi \right\rangle_{S_{n+m}} \\
 &= \left\langle \chi_1 \otimes \chi_2, \text{Res}_{S_n \times S_m}^{S_{n+m}} \psi \right\rangle \\
 &= \sum_{\substack{\sigma_1 \in S_n \\ \sigma_2 \in S_m}} \chi_1(\sigma_1) \chi_2(\sigma_2) p_1^{c_1(\sigma_1)} \cdots p_n^{c_n(\sigma_1)} p_1^{c_1(\sigma_2)} \cdots p_m^{c_m(\sigma_2)} \\
 &= \text{ch}(\chi_1) \text{ch}(\chi_2)
 \end{aligned}$$

We use Frobenius reciprocity somewhere in here. We also have  $\psi : S_n \rightarrow \Lambda_n$  and  $\psi(\tau\sigma\tau^{-1}) = \psi(\sigma)$ .  $\square$

- After another 10 years of trying to understand the representations of the symmetric group, we'll be here.
- At this point, we can study compact Lie groups.