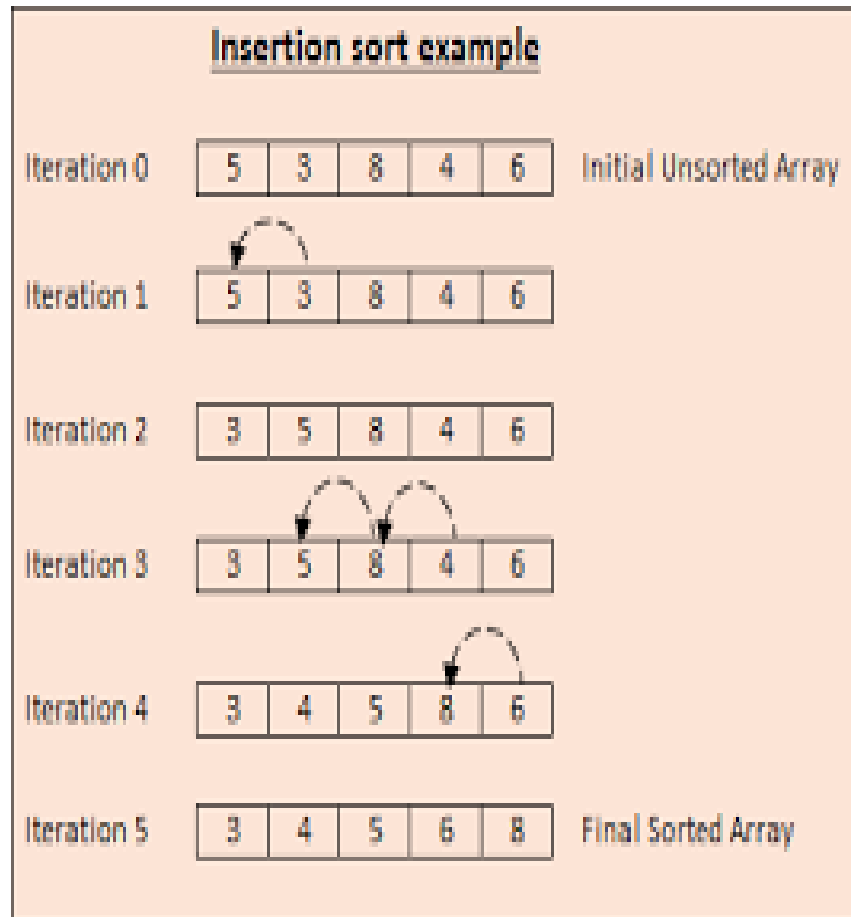


PROJECT REPORT:

INSERTION SORT:

- It is a straightforward Sorting calculation which sorts the cluster by moving components one by one.
- It dependably keeps up a sorted sublist in the lower places of the rundown
- Each new thing is then "embedded" once again into the past sublist to such an extent that the sorted sublist is one thing substantial.



REF: <http://www.programiz.com/sites/tutorial2program/files/Insertion-Sort-AlgorithmProgramming.jpg>

Insertion sorting is a set up sorting calculation. After n cycles in insertion sort, we would have figured out how to sort $(n+1)$ passages. In every emphasis the main outstanding information is evacuated and embedded into its right position.

WORST CASE ANALYSIS:

The worst case and average case time complexity for insertion sort is $O(n^2)$.

SELECTION SORT: The possibility of determination sort is basic: we over and again locate the following biggest (or littlest) component in the exhibit and move it to its last position in the sorted cluster. Choice sort calculation begins by contrasting the initial two components of the exhibit and swapping if essential.

Give us a chance to consider that we are sorting in rising request, then choice sort analyzes the first and second component, if the primary component is more noteworthy than the second component it swaps the first and second component, else its left undisturbed. This procedure goes ahead until the end of the exhibit (i.e. Looking at the first and the last component). In the event that there are an aggregate of n components, then the above procedure is rehashed $n-1$ times so as to get a totally sorted cluster.

42	16	84	12	77	26	53
----	----	----	----	----	----	----

The array, before the selection sort operation begins.

12	16	64	42	77	26	53
----	----	----	----	----	----	----

The smallest number (12) is swapped into the first element in the structure.

12	16	84	42	77	26	53
----	----	----	----	----	----	----

In the data that remains, 16 is the smallest; and it does not need to be moved.

12	16	26	42	77	84	53
----	----	----	----	----	----	----

26 is the next smallest number, and it is swapped into the third position.

12	16	26	42	77	84	53
----	----	----	----	----	----	----

42 is the next smallest number; it is already in the correct position.

12	16	26	42	53	84	77
----	----	----	----	----	----	----

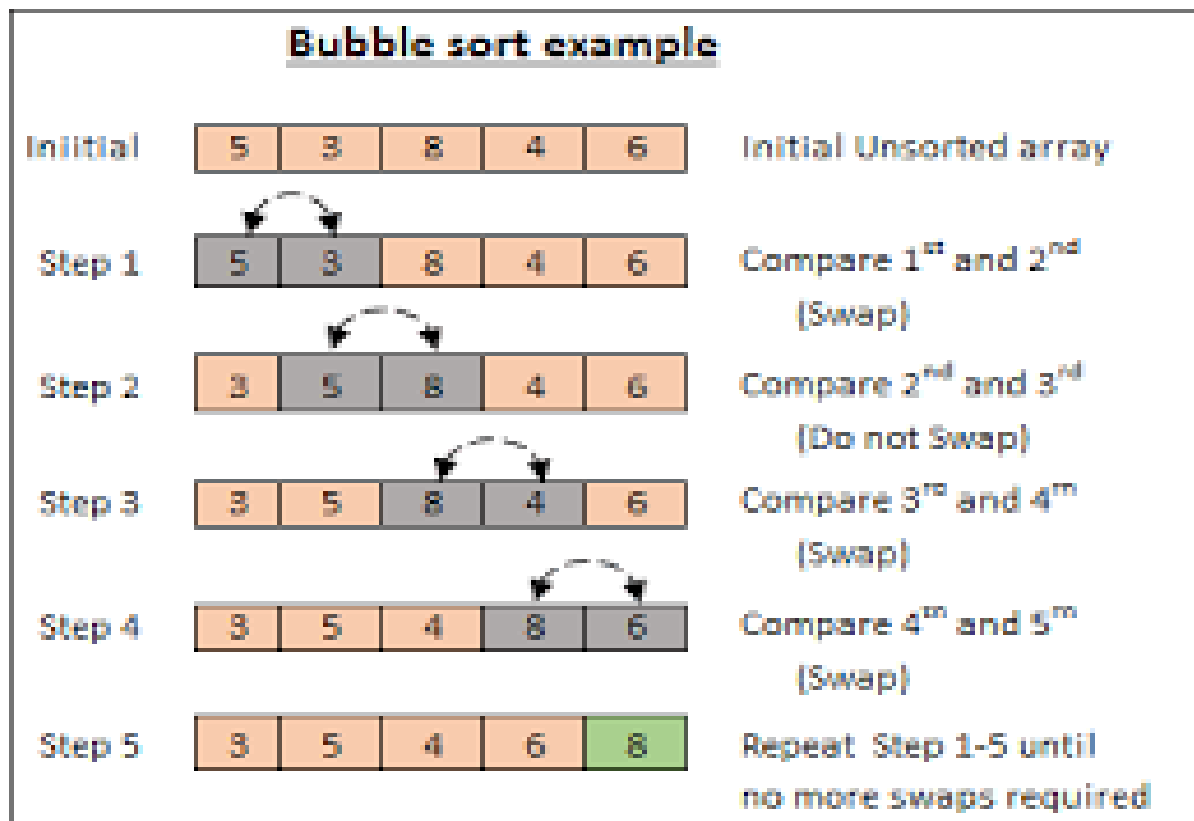
53 is the smallest number in the data that remains; and it is swapped to the appropriate position.

12	16	26	42	53	77	84
----	----	----	----	----	----	----

Of the two remaining data items, 77 is the smaller; the items are swapped. The selection sort is now complete.

BUBBLE SORT:

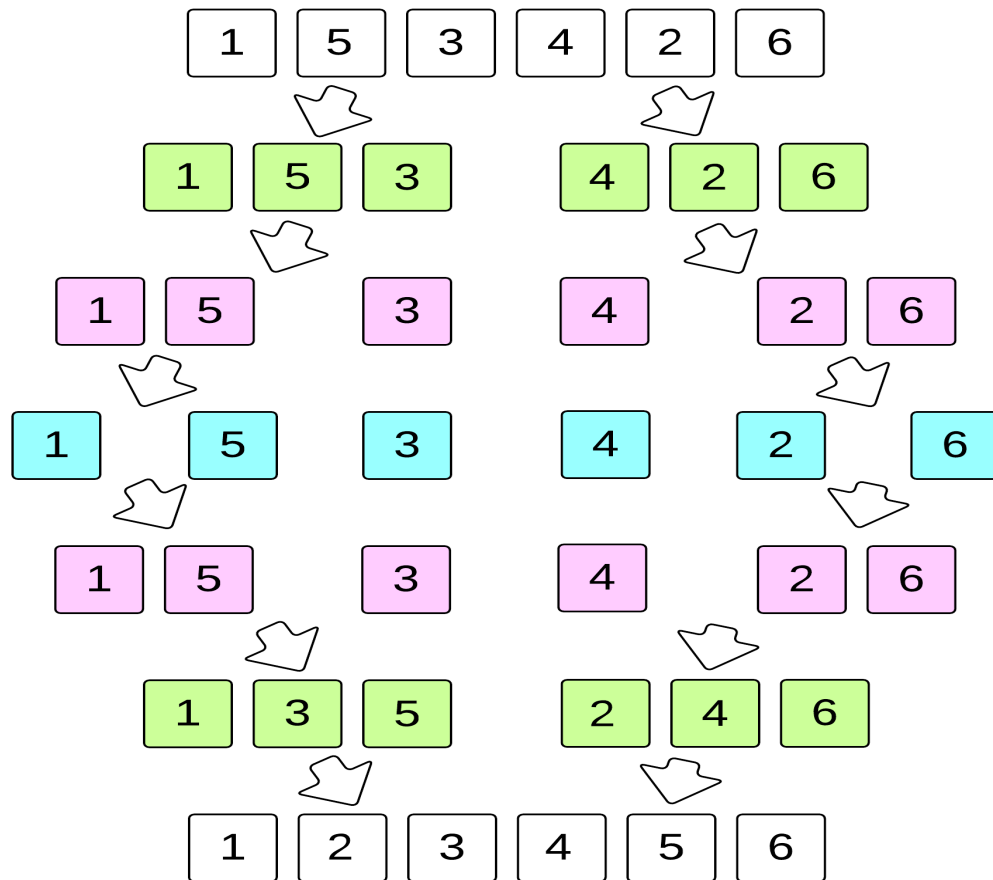
- Bubble deal with calculations cycle a rundown, investigating sets of components from left to right, or start to finish.
- If the furthest left component in the match is not exactly the furthest right component, the combine will stay in a specific order.
- If the furthest right component is not exactly the furthest left component, then the two components will be exchanged. This cycle rehashes from start to finish until a go in which no switch happens.



WORST CASE ANALYSIS:

The worst case and the average case complexity for bubble sort is $O(n^2)$. The space complexity for the algorithm is $O(1)$.

MERGE SORT: Merge Sort is a Divide and Conquer calculation. It separates input cluster in two parts, calls itself for the two parts and afterward combines the two sorted parts. The Merge() capacity is utilized for combining two parts.



WORST CASE ANALYSIS:

The recurrence relation for merge sort is:

$$T(n) = 2T(n/2) + O(n) \text{ if } n > 1$$

$$T(n) = O(1) \text{ if } n = 1$$

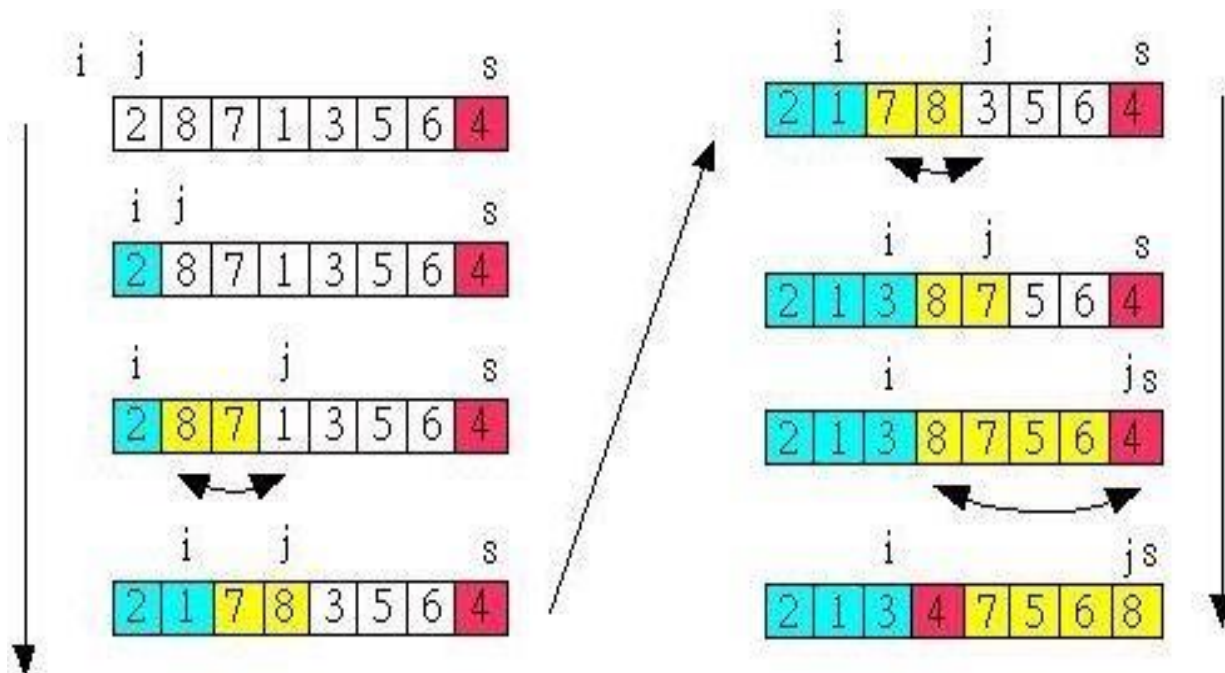
Applying the master's theorem on the above recursive relation, it can be shown that the worst case complexity of merge sort is $O(n \log n)$.

QUICK SORT: The contrast between Quick sort and Merge sort is that consolidation sort does almost no calculation before the recursive calls are made and performs work simply after the recursive calls are registered, whereas Merge sort plays out the greater part of its additional work preceding making the recursive calls.

Quicksort utilizes these means:

1. Pick any component of the exhibit to be the turn.
2. Isolate every single other component (aside from the turn) into two allotments.
 - All components not exactly the rotate must be in the main parcel.
 - All components more noteworthy than the turn must be in the second segment.
3. Utilize recursion to sort both allotments.
4. Join the initially sorted segment, the rotate, and the second sorted parcel.

The best rotate makes segments of equivalent length. The most noticeably awful rotate makes an unfilled segment. The run-time of Quicksort extents from $O(n \log n)$ with the best turns, to $O(n^2)$ with the most noticeably awful rotates, where n is the quantity of components in the cluster.



DATA SETS:

I have used 4 different datasets in order to analyze the time and space complexity of the above mentioned techniques, varying my input size from 1 to 1000 for analyzing time and space complexity.

Of the 4 used datasets, two datasets are Synthetic datasets and the other two datasets are real-time datasets.

REAL DATA SETS:

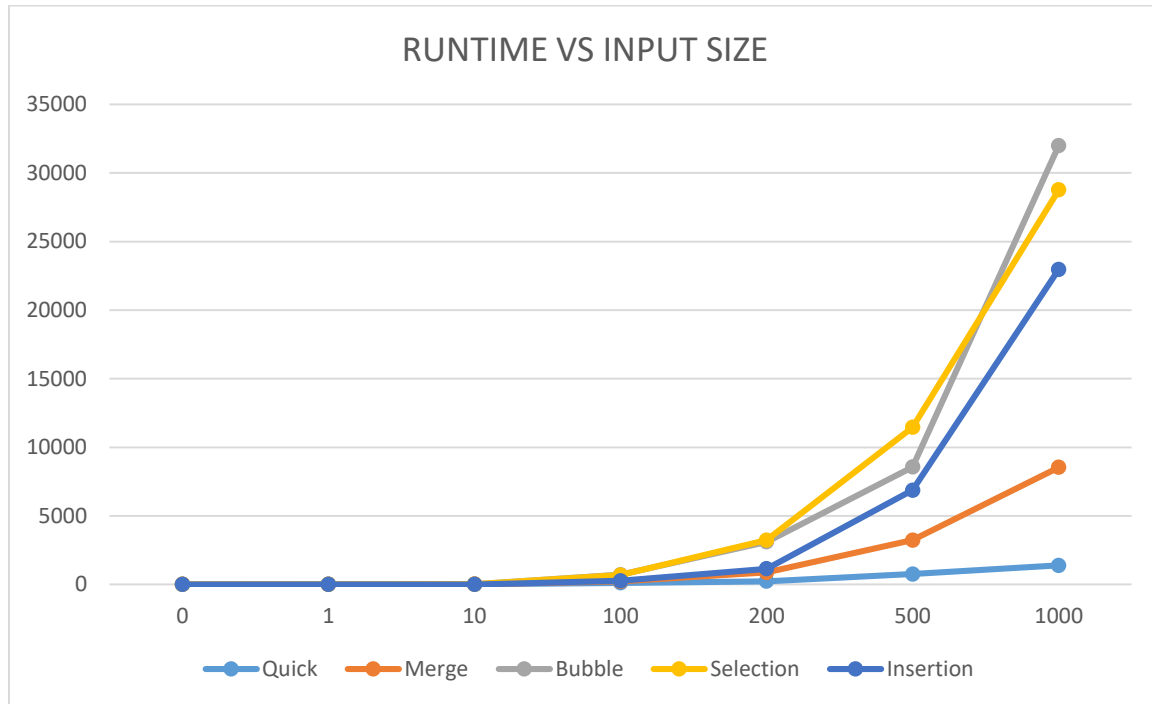
1. Random numbers generated with Java random functions.
2. Generated a **Normal distribution** dataset with mean of 12000 and **Standard deviation** of 1000.

SYNTHETIC DATA SETS REFERENCE:

<http://perso.telecom-paristech.fr/.eagan/class/igr204/datasets>

GRAPHS:

RUNTIME VS ARRAY SIZE: (RANDOM NUMBERS DATASETS)



QUICK SORT:

Input	T1	T2	T3	T4	T5	T(avg)
0	0	0	0	0	0	0
1	14	13	17	18	14	15.2
10	16	15	13	20	13	15.4
100	104	176	106	113	103	120.4
200	227	245	232	239	240	236.6
500	729	823	731	721	789	758.6
1000	1322	1466	1317	1448	1424	1395.4

MERGE SORT:

Input	T1	T2	T3	T4	T5	T(avg)
0	0	0	0	0	0	0
1	1	3	2	1	3	2
10	26	32	24	31	29	28.4
100	206	253	196	244	267	233.2
200	861	842	881	893	920	879.4
500	3119	3362	3261	2924	3496	3232.4
1000	8372	8217	8416	9007	8721	8546.6

BUBBLE SORT:

Input	T1	T2	T3	T4	T5	T(avg)
0	0	0	0	0	0	0
1	3	4	5	4	6	4.4
10	18	12	17	16	17	16
100	708	672	829	652	701	712.4
200	2381	3142	3619	2816	3599	3111.4
500	8659	8266	8913	8765	8216	8563.8
1000	34100	28945	33633	31153	32113	31988.8

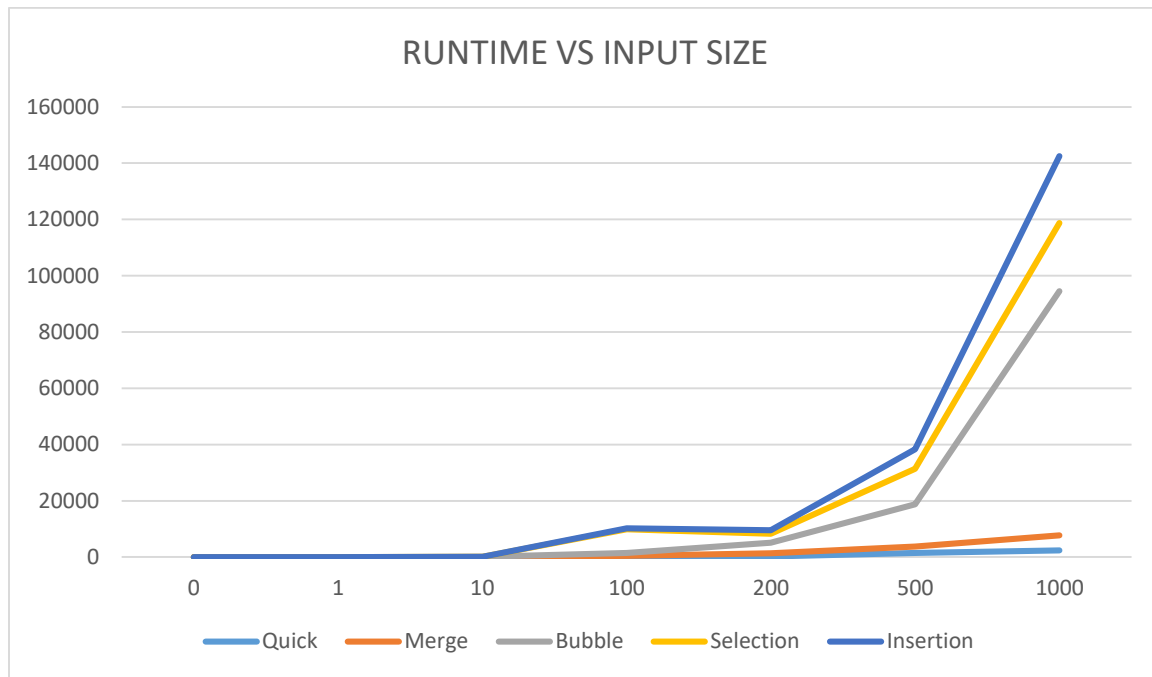
SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(avg)
0	0	0	0	0	0	0
1	2	2	3	2	3	2.4
10	14	17	17	15	14	15.4
100	678	691	742	572	769	690.4
200	3129	3254	3310	2879	3589	3232.3
500	11129	10986	11679	12789	10673	11451.2
1000	28674	26842	28427	28780	30863	28761.4

INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(avg)
0	0	0	0	0	0	0
1	2	1	2	2	1	1.6
10	8	7	11	7	9	8.4
100	262	247	295	287	303	278.8
200	1043	1361	985	1193	1172	1150.8
500	6723	7709	7021	6727	6192	6874.4
1000	23699	22943	23274	21245	23679	22968

RUNTIME VS ARRAY SIZE: (NORMAL DISTRIBUTION)



INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	2	3	3	2	2	2.4
10	10	7	8	11	8	8.8
100	346	305	366	354	351	344.4
200	1262	1055	1501	1395	1044	1251.4
500	6119	6982	7064	7695	6986	6969.2
1000	22815	22474	20696	23675	29098	23751.6

SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	2	3	1	2	1	1.8
10	13	11	10	12	14	12
100	816	832	805	876	849	835.6
200	3125	3096	3167	3255	3297	3188
500	12745	12626	12071	12918	12613	12594.6
1000	25552	23318	22962	25910	23361	24220.6

BUBBLE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	3	2	3	4	3	3
10	12	15	14	12	16	13.8
100	976	1005	1097	1018	1047	1028.6
200	3731	3946	3826	3731	3598	3766.4
500	15305	14625	14749	15316	14894	14977.8
1000	87146	87693	81497	88611	88789	86750.8

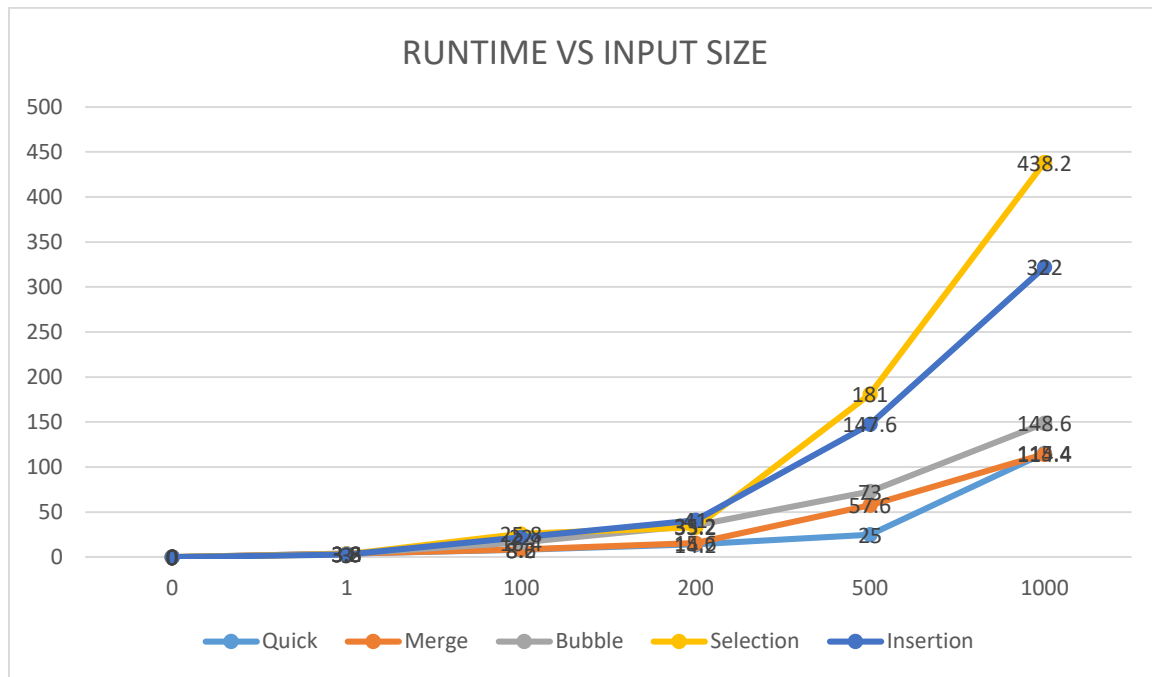
MERGE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	4	3	3	4	3	3.4
10	23	31	22	23	32	26.2
100	321	345	319	327	344	331.2
200	1179	1045	1292	1186	1177	1175.8
500	2365	2492	2292	2254	2292	2339
1000	5286	5164	5244	5286	5760	5348

QUICK SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	19	12	16	17	19	16.6
10	20	19	17	20	19	15.4
100	123	156	167	179	63	157.6
200	224	227	221	223	225	224
500	1367	1467	1492	1517	1478	1464.2
1000	2459	2366	2519	2319	2502	2432.4

RUNTIME VS ARRAY SIZE: (CAMERA-DATA SETS)



INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	3	2	3	2	5	3
100	21	23	22	21	23	22
200	43	38	41	40	43	41
500	149	142	147	144	156	147.6
1000	324	339	321	298	332	322.8

SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	3	3	4	3	3	3.2
100	24	19	25	33	28	25.8
200	29	42	31	35	29	33.2
500	177	169	167	176	167	181
1000	462	456	424	438	411	438.2

BUBBLE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	3	4	3	5	3	3.6
100	16	15	18	16	17	16.4
200	32	34	33	39	38	35.2
500	70	69	72	82	72	73
1000	141	147	146	159	150	148.6

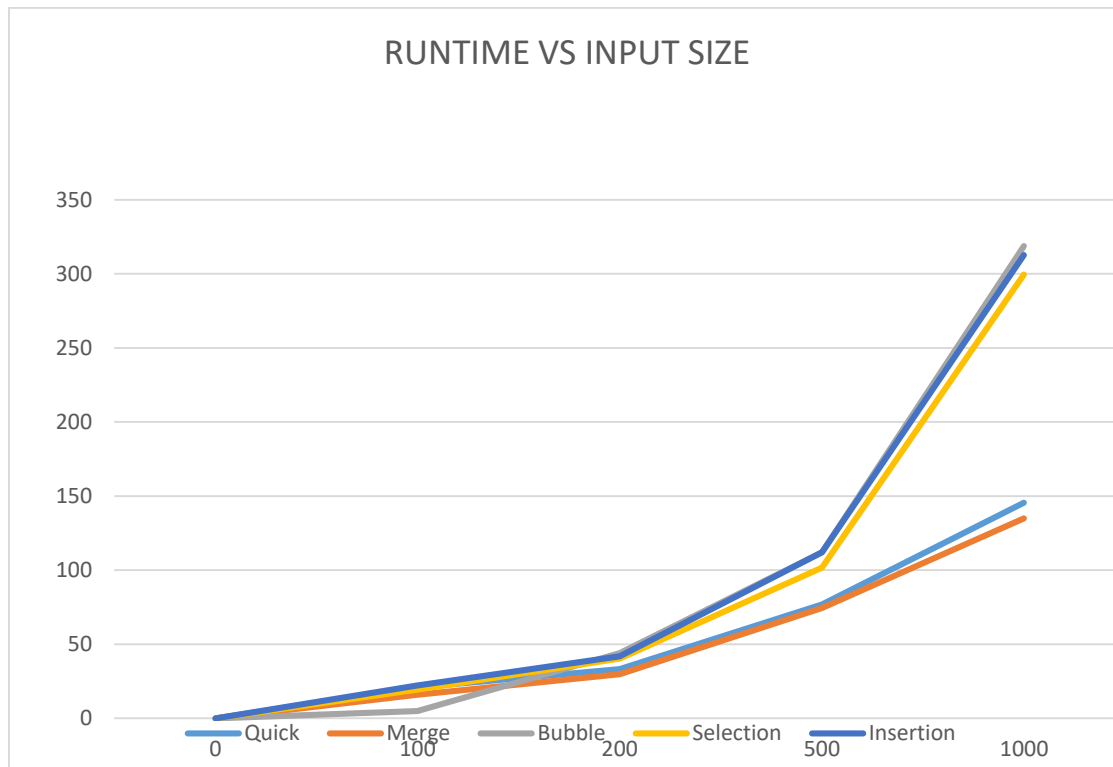
MERGE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	3	4	5	3	4	3.8
100	8	9	11	7	8	8.6
200	16	11	18	17	16	15.6
500	61	59	62	54	52	57.6
1000	112	119	119	102	121	114.4

QUICK SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
1	3	5	4	3	4	3.8
100	8	11	9	6	7	8.2
200	15	16	11	18	11	14.2
500	24	28	25	26	22	25
1000	112	119	123	102	121	115.4

RUNTIME VS ARRAY SIZE: (PRICE DATA-SET)



INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	21	23	22	24	21	22.2
200	40	45	44	42	38	41.2
500	98	117	114	115	116	112
1000	310	322	312	296	324	312.8

SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	22	16	25	18	16	19.4
200	32	46	39	48	36	40.2
500	98	110	102	103	95	101.6
1000	303	318	293	285	299	299.6

BUBBLE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	4	6	3	4	7	4.8
200	42	46	49	40	43	44
500	123	109	108	114	106	112
1000	326	319	325	312	302	318.8

MERGE SORT:

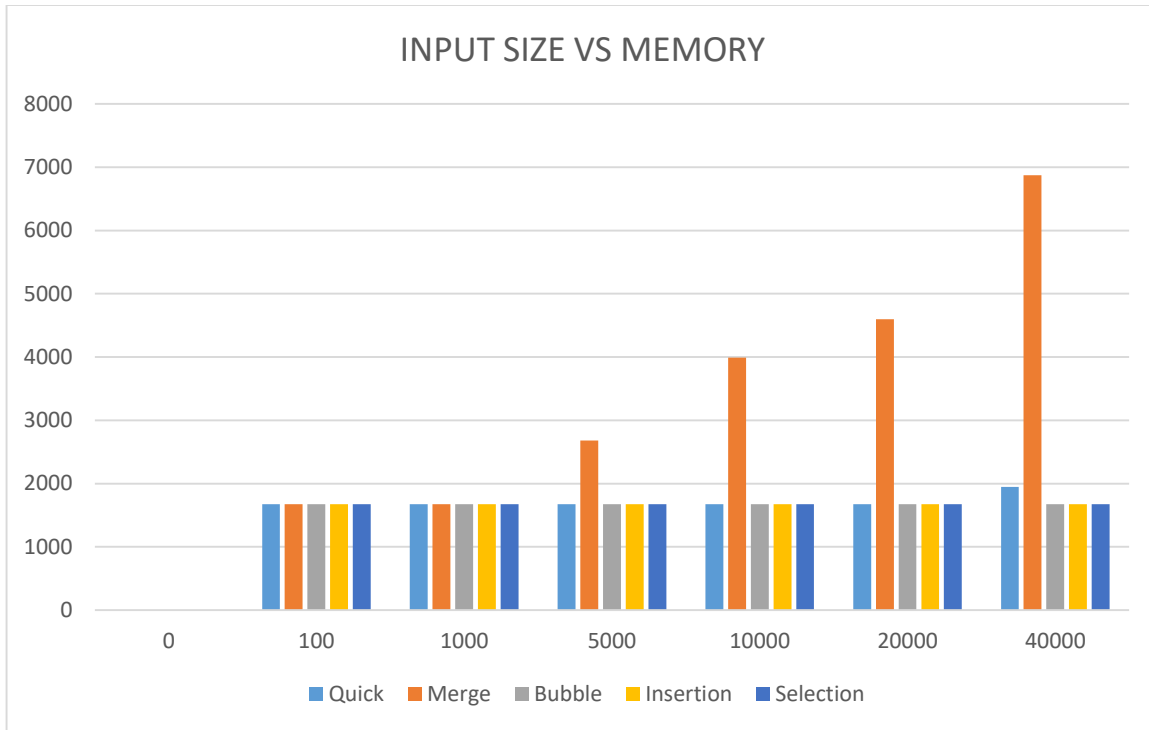
Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	16	15	17	17	14	15.8
200	29	30	28	31	30	29.6
500	68	72	86	79	67	74.4
1000	147	135	132	131	130	135

QUICK SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	19	25	17	21	22	20.8
200	32	31	36	35	32	33.2
500	82	78	81	72	71	76.8
1000	151	146	147	142	141	145.4

From graphs we conclude that:

It is clear from the bends portrayed over that the three Sorting calculation to be specific Insertion sort, bubble sort, Selection sort have a bend that is expanding at a quadratic rate. Thus the most pessimistic scenario many-sided quality of these three calculations are $O(n^2)$. The Merge sort is sub-direct with the pattern of $\log n$. Quick sort is again sub-direct with the pattern of $\log n$. The most pessimistic scenario of quicksort is $O(n^2)$ which is the same as that of Insertion, bubble, selection sorting calculations. Thus from the above graphs we can conclude that quick sort and merge sort perform better than Insertion, bubble, selection sorting algorithms for large input sizes.



INPUT SIZE VS MEMORY USAGE: (RANDOM DATA SETS)

INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1672	1672	1672	1672	1672	1672

SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1672	1672	1672	1672	1672	1672

BUBBLE SORT:

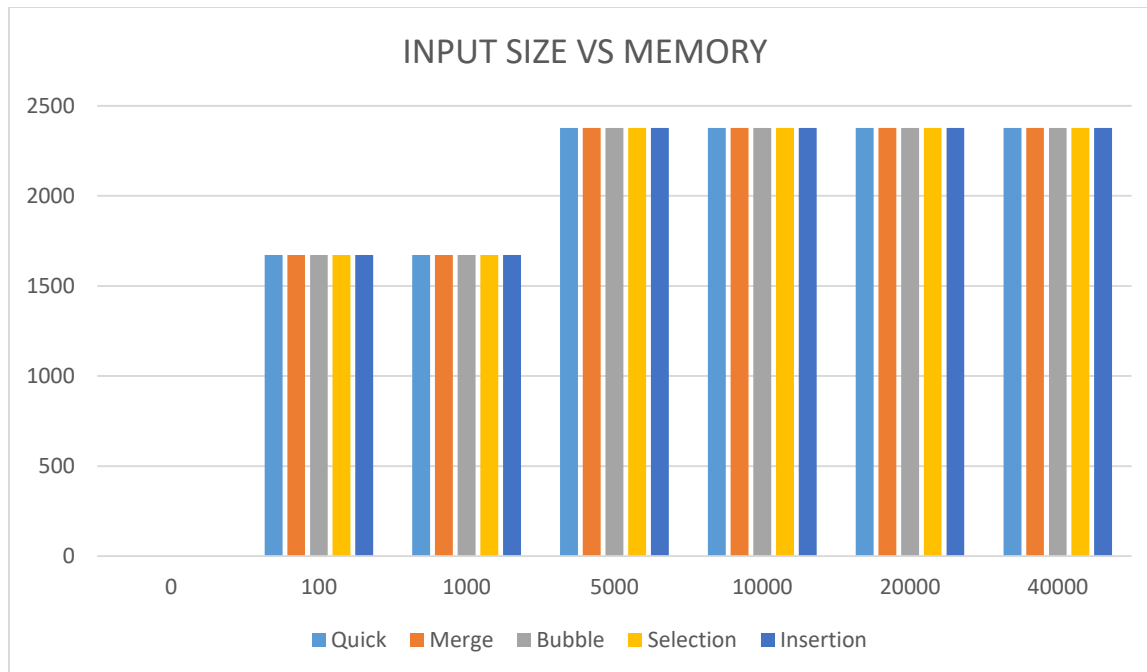
Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1672	1672	1672	1672	1672	1672

MERGE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	2678	2678	2678	2678	2678	2678
10000	3989	3989	3989	3989	3989	3989
20000	4598	4598	4598	4598	4598	4598
40000	6827	6827	6827	6827	6827	6827

QUICK SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1948	1948	1948	1948	1948	1948



INPUT SIZE VS MEMORY USAGE: (NORMAL DISTRIBUTION)

INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1672	1672	1672	1672	1672	1672

SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1672	1672	1672	1672	1672	1672

BUBBLE SORT:

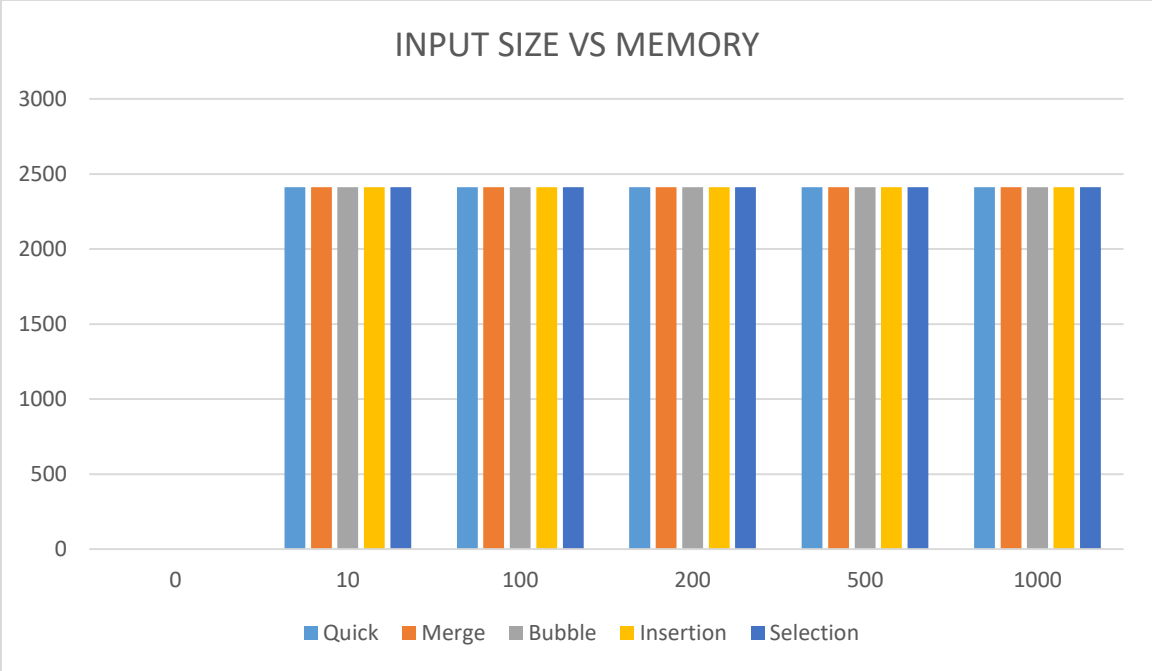
Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1672	1672	1672	1672	1672	1672

QUICK SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	1672	1672	1672	1672	1672	1672
10000	1672	1672	1672	1672	1672	1672
20000	1672	1672	1672	1672	1672	1672
40000	1672	1672	1672	1672	1672	1672

MERGE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	0	0	0	0	0	0
100	1672	1672	1672	1672	1672	1672
1000	1672	1672	1672	1672	1672	1672
5000	2378	2378	2378	2378	2378	2378
10000	2378	2378	2378	2378	2378	2378
20000	2378	2378	2378	2378	2378	2378
40000	2378	2378	2378	2378	2378	2378



INPUT SIZE VS MEMORY USAGE: (CAMERA DATA SETS)

INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412
5000	2412	2412	2412	2412	2412	2412
10000	2412	2412	2412	2412	2412	2412
20000	2412	2412	2412	2412	2412	2412
40000	2412	2412	2412	2412	2412	2412

SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412
5000	2412	2412	2412	2412	2412	2412
10000	2412	2412	2412	2412	2412	2412
20000	2412	2412	2412	2412	2412	2412
40000	2412	2412	2412	2412	2412	2412

BUBBLE SORT:

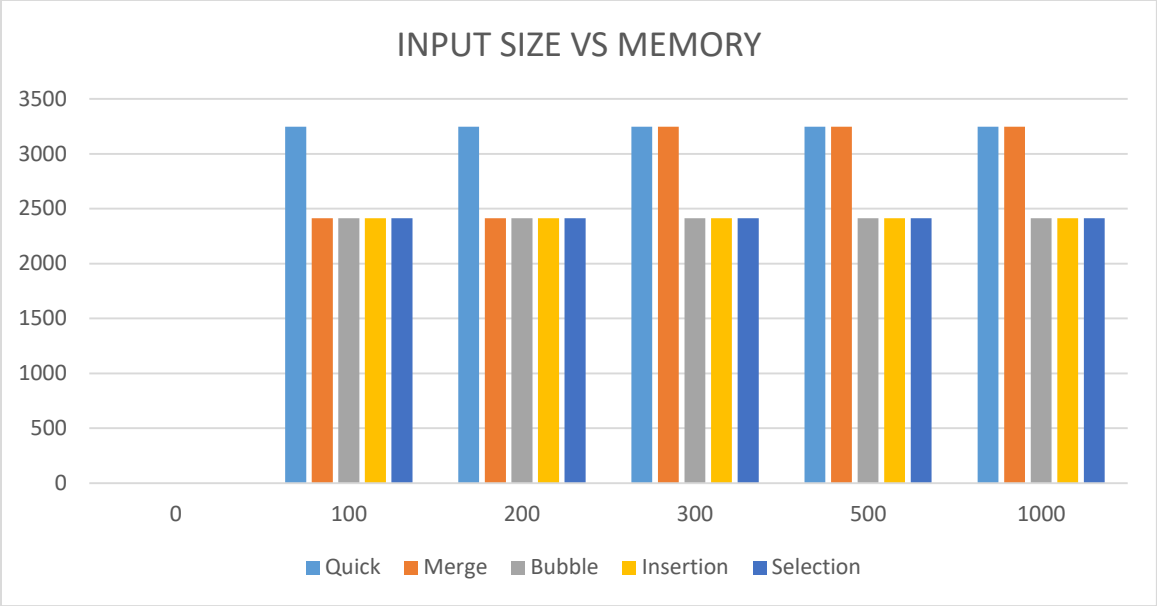
Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412
5000	2412	2412	2412	2412	2412	2412
10000	2412	2412	2412	2412	2412	2412
20000	2412	2412	2412	2412	2412	2412
40000	2412	2412	2412	2412	2412	2412

MERGE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412
5000	2412	2412	2412	2412	2412	2412
10000	2412	2412	2412	2412	2412	2412
20000	2412	2412	2412	2412	2412	2412
40000	2412	2412	2412	2412	2412	2412

QUICK SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412
5000	2412	2412	2412	2412	2412	2412
10000	2412	2412	2412	2412	2412	2412
20000	2412	2412	2412	2412	2412	2412
40000	2412	2412	2412	2412	2412	2412



INPUT SIZE VS MEMORY USAGE: (PRICE DATA SETS)

INSERTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
200	2412	2412	2412	2412	2412	2412
300	2412	2412	2412	2412	2412	2412
500	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412

SELECTION SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
200	2412	2412	2412	2412	2412	2412
300	2412	2412	2412	2412	2412	2412
500	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412

BUBBLE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
200	2412	2412	2412	2412	2412	2412
300	2412	2412	2412	2412	2412	2412
500	2412	2412	2412	2412	2412	2412
1000	2412	2412	2412	2412	2412	2412

QUICK SORT:

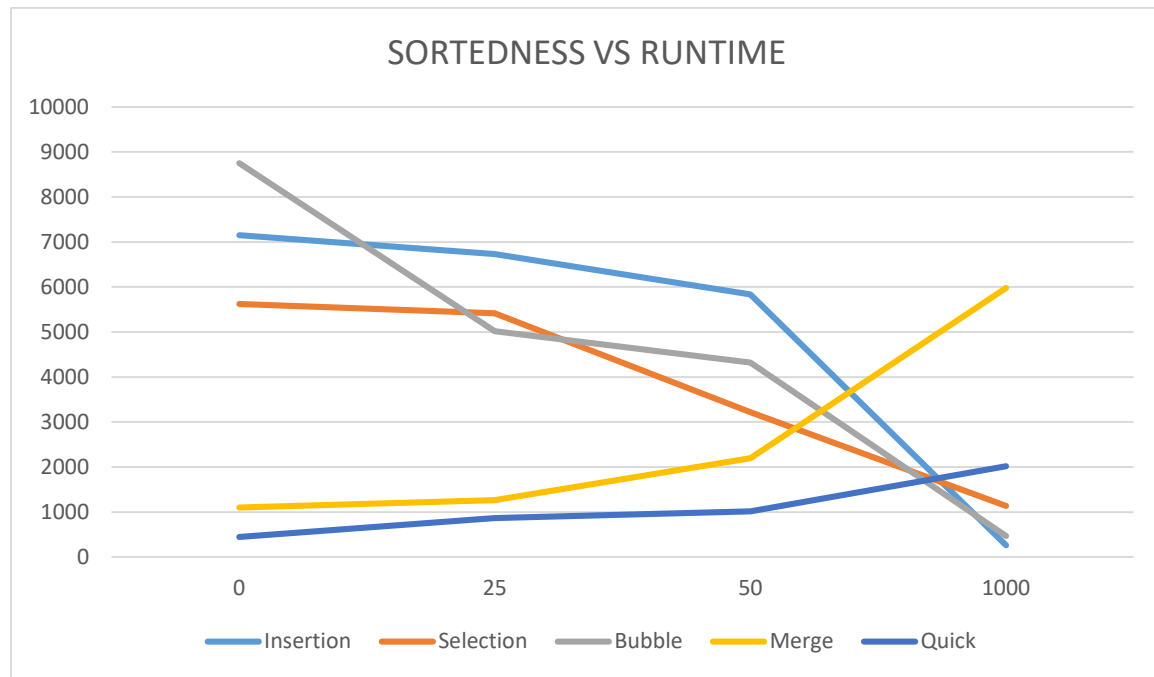
Input	T1	T2	T3	T4	T5	T(Avg)
0	3248	3248	3248	3248	3248	3248
100	3248	3248	3248	3248	3248	3248
200	3248	3248	3248	3248	3248	3248
300	3248	3248	3248	3248	3248	3248
500	3248	3248	3248	3248	3248	3248
1000	3248	3248	3248	3248	3248	3248

MERGE SORT:

Input	T1	T2	T3	T4	T5	T(Avg)
0	2412	2412	2412	2412	2412	2412
100	2412	2412	2412	2412	2412	2412
200	2412	2412	2412	2412	2412	2412
300	3248	3248	3248	3248	3248	3248
500	3248	3248	3248	3248	3248	3248
1000	3248	3248	3248	3248	3248	3248

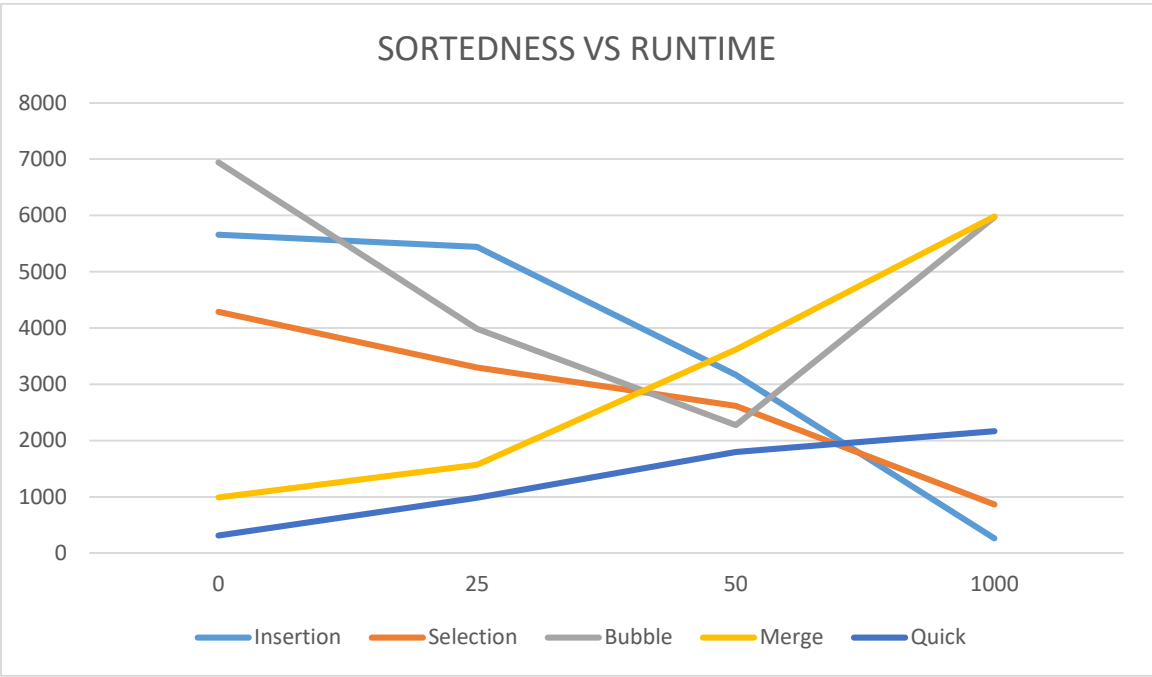
From graphs we conclude that:

The three sorting calculations Insertion sort, selection and bubble sort devour consistent memory, concerning the information estimate. Merge sort and quicksort just possess abundance memory when the information size is expansive. Merge sort performs more terrible than quicksort regarding memory, as it takes in an additional spaceduring the separation step.



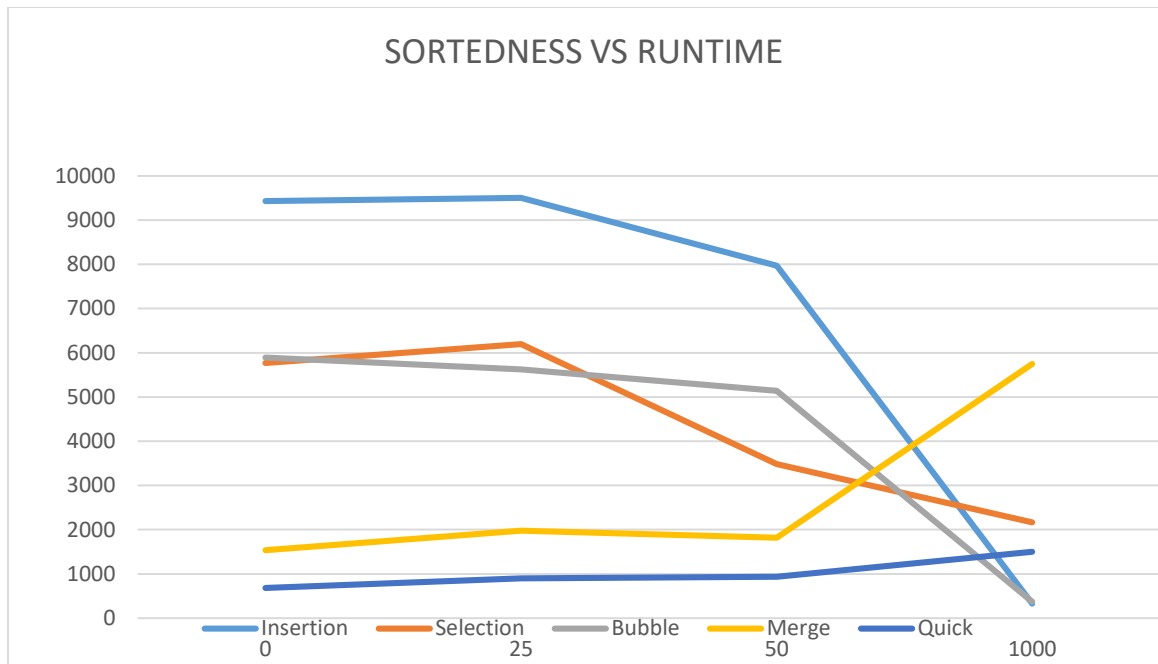
SORTEDNESS RANDOM VS RUNTIME: (NORMAL)

Input	Insertion	Selection	Bubble	Merge	Quick
0	7149.2	5618.75	8751.2	1103.27	445.3
25	6734	5416	5018.9	1267	867.5
50	5832	3218.25	4318	21986	1018.9
1000	264	1137	467	5978	2019.5



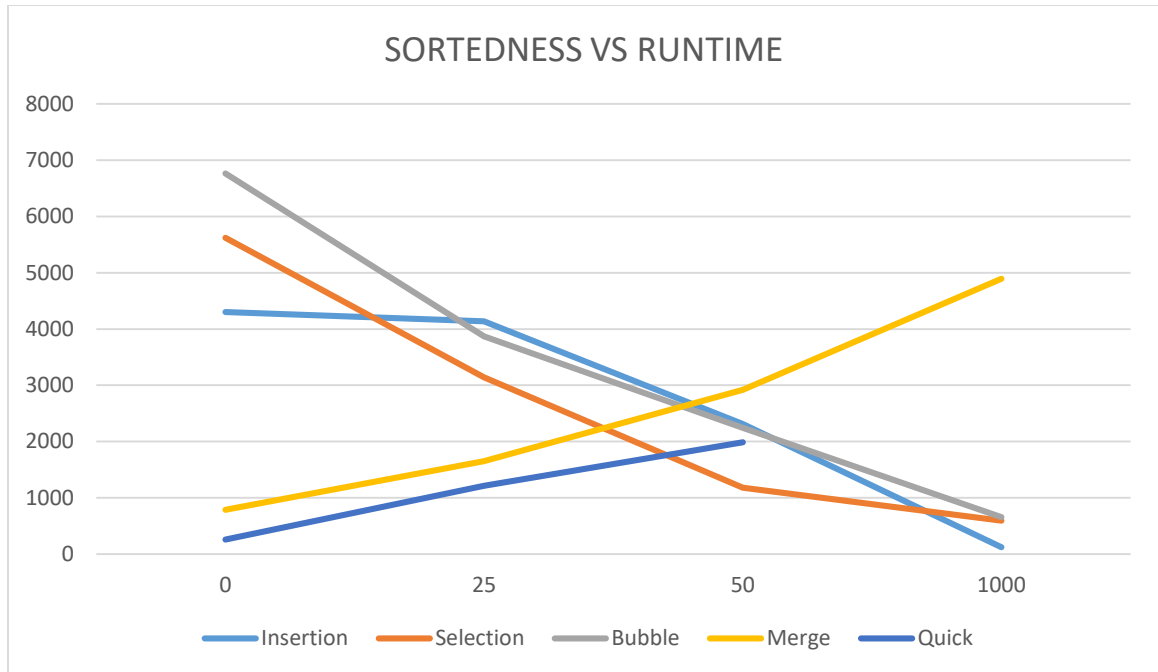
SORTEDNESS RANDOM VS RUNTIME: (CAMERA -DATASETS)

Input	Insertion	Selection	Bubble	Merge	Quick
0	5658	4289.7	6943	987	312
25	5438.9	3296.8	3987.95	1568.25	986.25
50	3168.2	2617	2274.25	3618.75	1793.8
1000	179	866	5968	5982	2168.9



SORTEDNESS RANDOM VS RUNTIME: (RANDOM DATASETS)

Input	Insertion	Selection	Bubble	Merge	Quick
0	9429	5796.25	5893	1534.25	683.25
25	9501.25	6195.75	5627.25	1976	896.75
50	7959	3482.65	5138.75	1817.65	937
1000	327	2167.2	373	5750	1498

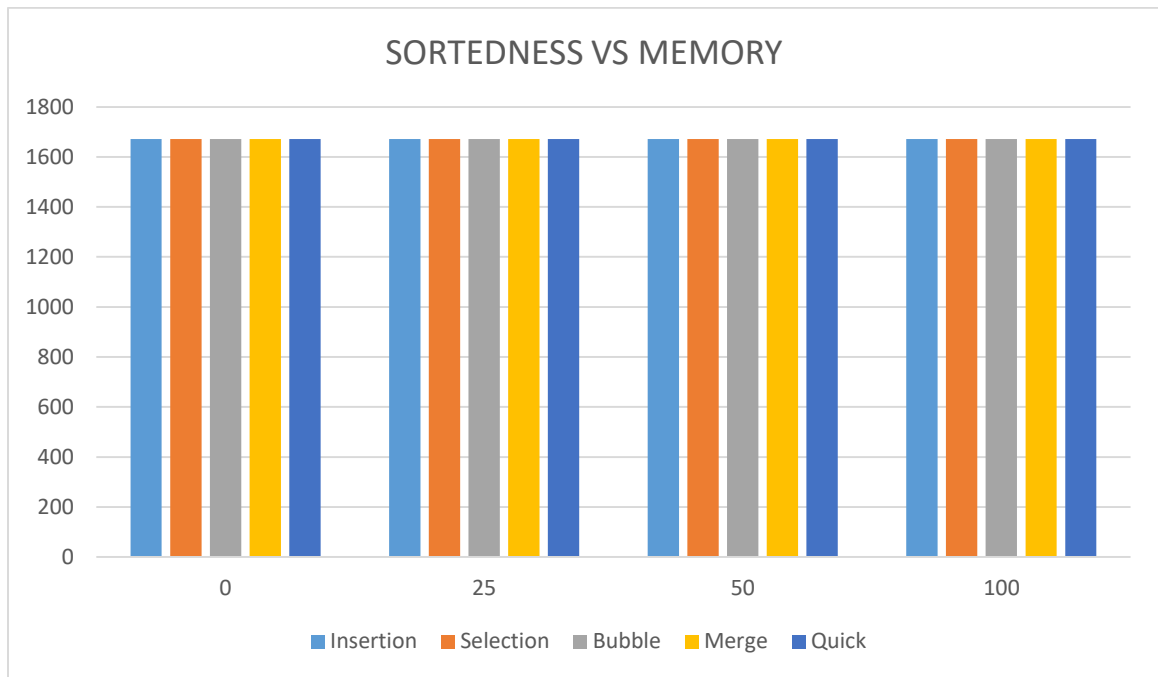


SORTEDNESS RANDOM VS RUNTIME: (PRICE DATASETS)

Input	Insertion	Selection	Bubble	Merge	Quick
0	4301	5618.7	6761	789.2	259
25	4134.2	3139.8	3872.1	1653.8	1213.8
50	2312.5	1179.25	2243.2	2918.5	1988.5
1000	123	595	656	4891.2	2894.6

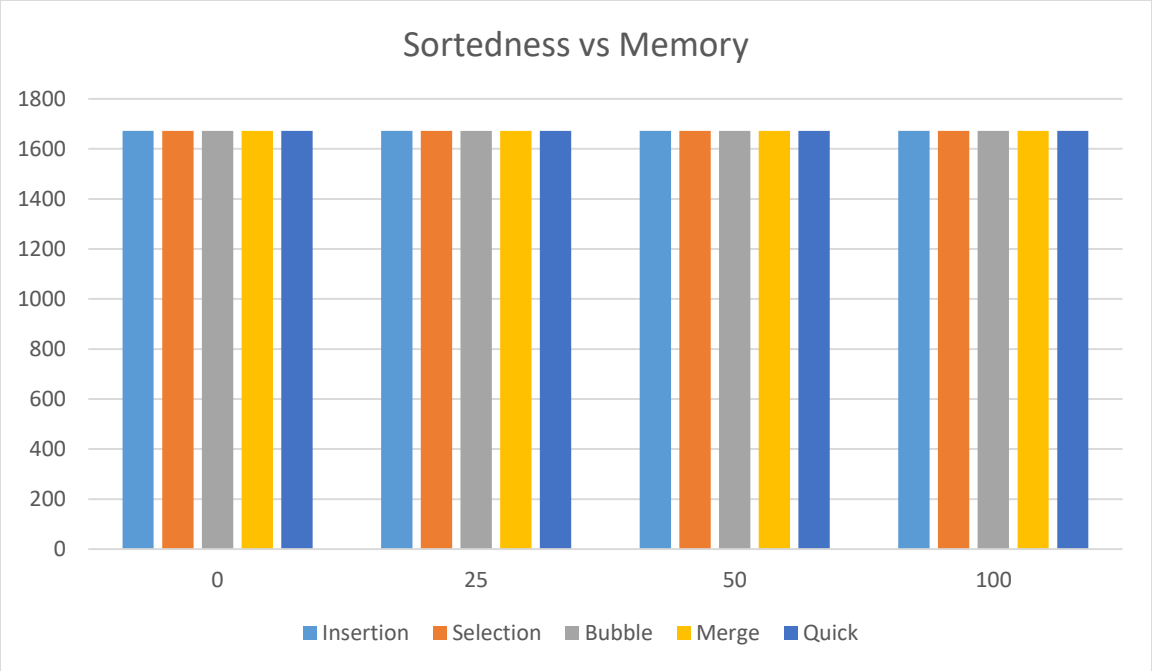
From the graphs we conclude that:

If an exhibit is not sorted, the three sorting calculations insertion sort, selection sort, bubble sort have higher runtime when contrasted with merge sort and quick sort on differing the information estimate. As the level of sortedness increments (25%,50%) the runtime of the three calculations insertion sort, bubble sort, selection sort diminishes, whereas the runtime of union sort and quicksort keep on increasing. When the cluster is totally sorted, merge sort and quick sort have the greatest runtime at the point when contrasted with bubble sort, insertion sort and selection sort. Runtime of merge sort increments radically when the exhibit is totally sorted.



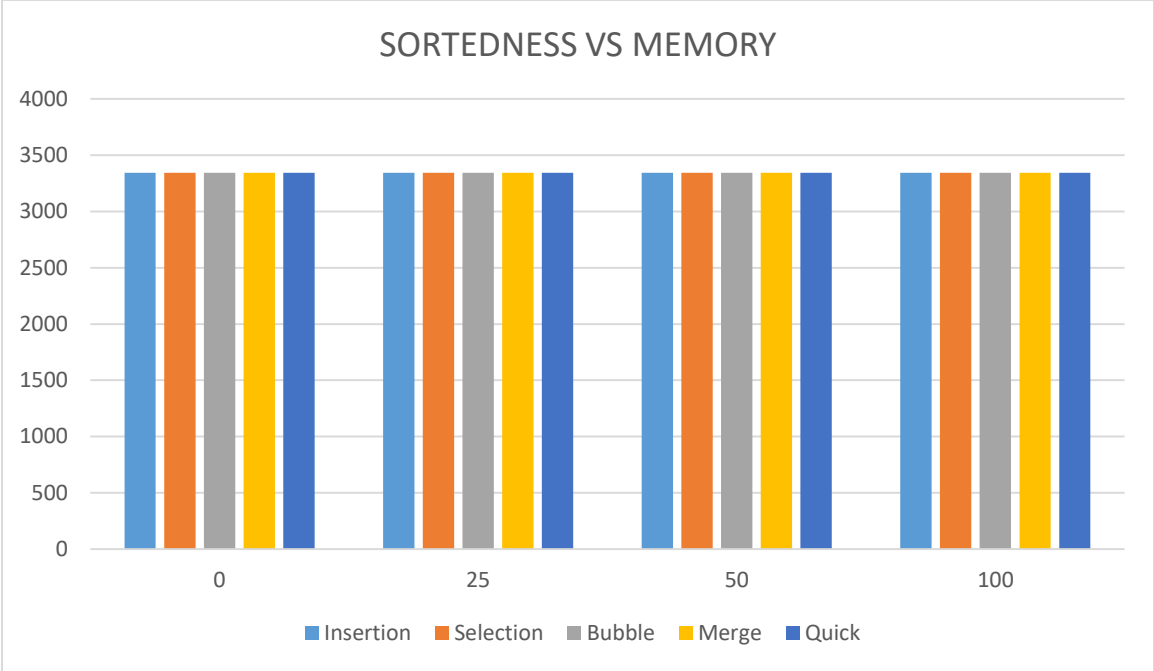
SORTEDNESS VS MEMORY: (RANDOM)

Input	Insertion	Selection	Bubble	Merge	Quick
0	1672	1672	1672	1672	1672
25	1672	1672	1672	1672	1672
50	1672	1672	1672	1672	1672
100	1672	1672	1672	1672	1672



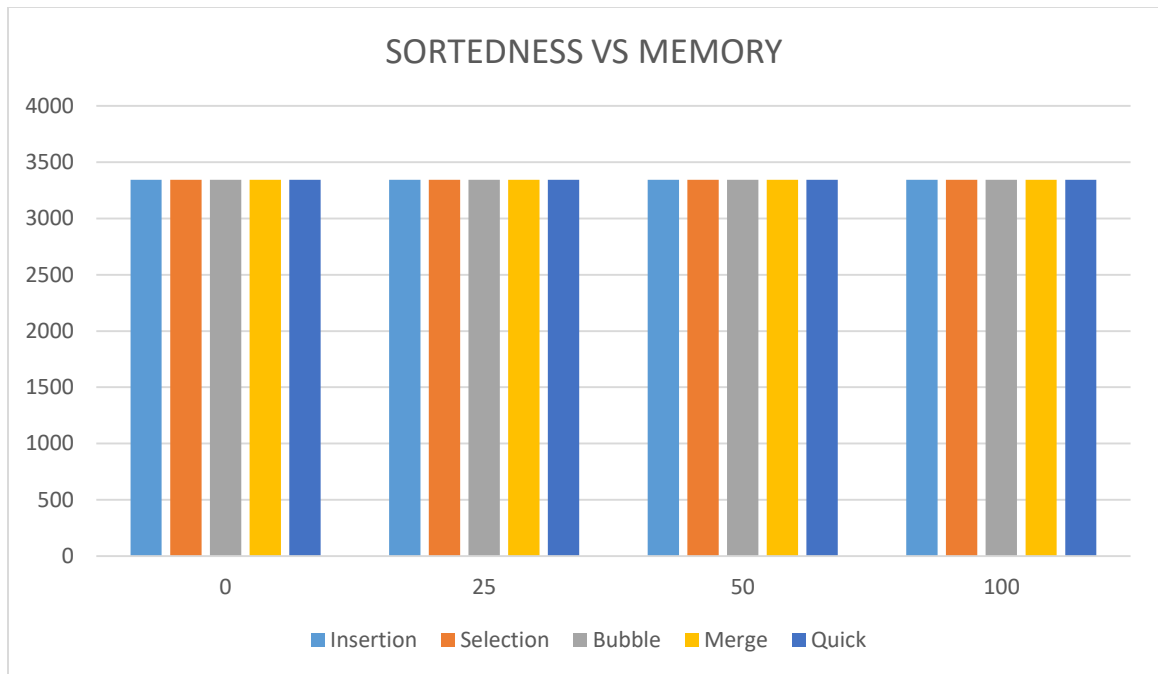
SORTEDNESS VS MEMORY: (NORMAL)

Input	Insertion	Selection	Bubble	Merge	Quick
0	1672	1672	1672	1672	1672
25	1672	1672	1672	1672	1672
50	1672	1672	1672	1672	1672
100	1672	1672	1672	1672	1672



SORTEDNESS VS MEMORY: (CAMERA -DATASETS)

Input	Insertion	Selection	Bubble	Merge	Quick
0	3344	3344	3344	3344	3344
25	3344	3344	3344	3344	3344
50	3344	3344	3344	3344	3344
100	3344	3344	3344	3344	3344



SORTEDNESS VS MEMORY: (PRICE DATASETS)

Input	Insertion	Selection	Bubble	Merge	Quick
0	3344	3344	3344	3344	3344
25	3344	3344	3344	3344	3344
50	3344	3344	3344	3344	3344
100	3344	3344	3344	3344	3344

From the graphs we conclude that:

The memory of the 5 specified calculations does not rely on the level of sortedness. For all estimations of level of sortedness the calculation is as yet expending a similar steady measure of memory. The memory depends just on the info size of the algorithm.

SUMMARY OF SORTEDNESS MEASURE:

- With high degree of sortedness, the merge sort and quick sort perform worse than bubble sort, insertion sort and selection sort.
- The space complexity of Merge sort and quick sort are greater than any other sorting algorithm as they are not in-place sorting algorithms. Other sorting algorithms such as insertion sort, bubble sort and selection sort are in-place sorting algorithms and their space complexity is $O(1)$.
- The runtime of quick sort highly depends on choosing the pivot element. Consistent poor choices of pivot element can result in drastically slow $O(n^2)$ complexity for quick sort. Choosing median as the pivot results in $O(n \log n)$ complexity.
- When we have large amount of data the merge sort and quick sort perform a lot better than insertion sort, bubble sort and selections sort.
- Quick sort and merge sort work on the fundamental principle of divide and conquer. This makes them highly scalable sorting algorithms