

*****Our Template*****

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
using namespace std;

#define ll long long
#define fast ios::sync_with_stdio(false);cin.tie(nullptr)
```

```
template<class T> using special_set = tree<T,
null_type, less<T>, rb_tree_tag,
tree_order_statistics_node_update>;
```

```
// String transformations
```

```
string toupper(string a) {
    for (char &c : a)
        if (islower(c))
            c = toupper(c);
    return a;
}
```

```
string tolower(string a) {
    for (char &c : a)
        if (isupper(c))
            c = tolower(c);
    return a;
}
```

```
// Factorial calculation
```

```
int factorial(int k) {
    int res = 1;
    for (int i = 2; i <= k; ++i)
        res *= i;
    return res;
}
```

```
ll calc_combinatoric(ll a, ll b, ll mod) {
    ll fact_a = factorial(a, mod);
    ll fact_b = factorial(b, mod);
    ll fact_a_b = factorial(a - b, mod);
    ll inv_fact_b = power_MOD(fact_b, mod - 2, mod);
    ll inv_fact_a_b = power_MOD(fact_a_b, mod - 2,
mod);
```

```
    return (fact_a * inv_fact_b % mod * inv_fact_a_b
% mod);
}
```

```
// Comparator for pairs
```

```
bool cmp(pair<int, int> a, pair<int, int> b) {
    if (a.first == b.first) {
        return a.second > b.second;
    } else {
        return a.first < b.first;
    }
}
```

```
// Bitwise operations
```

```
bool is_odd(ll x) {
    return x & 1;
}

bool check_kth_bit(ll x, ll k) {
    return x & (1LL << k);
}

ll set_kth_bit_1(ll x, ll k) {
    return x | (1LL << k);
}

ll set_kth_bit_0(ll x, ll k) {
    return x & ~(1LL << k);
}

ll multiply_by_power_of_2(ll x, ll k) {
    return x << k;
}

ll divide_by_power_of_2(ll x, ll k) {
    return x >> k;
}

int count_set_bits(int x) {
    return __builtin_popcount(x);
}

ll count_set_bits_ll(ll x) {
    return __builtin_popcountll(x);
}

ll toggle_kth_bit(ll x, ll k) {
    return x ^ (1LL << k);
}

bool is_power_of_2(ll x) {
    return x && !(x & (x - 1));
}
```

```
bool find_str(string str1, string str2) {
    if (str1.find(str2) != string::npos) {
        return true;
    } else {
        return false;
    }
}
```

```
ll floor_val(ll n, ll a) {
    return n / a;
}

ll ceil_val(ll n, ll a) {
    return (n + a - 1) / a;
}

ll round_val(ll n, ll a) {
    return (n + a / 2) / a;
}
```

```
void solve() {
    special_set<int> st;
    ...
}
```

```
int main() {
    fast;
    int q;
    cin >> q;
```

```
    while (q--) {
        solve();
    }
```

```
    cerr << "Execution time: " << 1000.f * clock() /
CLOCKS_PER_SEC << " ms." << endl;
```

```
    return 0;
}
```

Math Formula

Parallel :

(1) Average / subsequence / normal number of 1 to n,

$$\Rightarrow \text{Value} = (\text{1st} + \text{last}) / 2$$

$$(2) \text{Nth value} = a + (n - 1) * d$$

a = 1st number

n = Nth

d = normal difference

Nth value = last number

$$(3) \text{Sum} \rightarrow S_n = (n/2)(2a + (n-1)d)$$

$$1+2+3+...+n = (n+1) / 2$$

$$1^2+2^2+...+n^2 = n(n+1)(2n+1) / 6$$

$$1^3+2^3+...+n^3 = n(n+1) / 2$$

$$2+4+6+...+2n = n(n+1)$$

$$1+3+5+...+(2n-1) = n^2$$

$$\text{Harmonic Series : } 1 + (1/2) + (1/3) + (1/4) + ...$$

Multi-dimensional :

$$(1) \text{Nth number} = Ar^{(n-1)}$$

(2) Sum of 1 to n number sum,

$$\text{If } (r > 1) S_n = a(r^n - 1) / (r - 1)$$

$$\text{If } (r < 1) S_n = a(1 - r^n) / (1 - r)$$

SQUARE :

$$1. (a+b)^2 = a^2 + 2ab + b^2$$

$$2. (a+b)^2 = (a-b)^2 - 4ab$$

$$3. (a-b)^2 = a^2 - 2ab + b^2$$

$$4. (a-b)^2 = (a+b)^2 - 4ab$$

$$5. a^2 + b^2 = (a+b)^2 - 2ab$$

$$6. a^2 + b^2 = (a+b)^2 + (a-b)^2 / 2$$

$$7. a^2 - b^2 = (a-b)^2 + 2ab$$

$$8. a^2 - b^2 = (a+b)(a-b)$$

$$9. 2(a^2+b^2) = (a+b)^2 + (a-b)^2$$

$$10. 4ab = (a+b)^2 - (a-b)^2$$

$$11. Ab = (a+b/2)^2 - (a-b/2)^2$$

$$12. (a+b+c)^2 = (a^2+b^2+c^2) + 2ab+2bc+2ca$$

$$13. 2(ab+bc+ca) = (a+b+c)^2 - (a^2+b^2+c^2)$$

CUBE:

$$1. (a+b)^3 = a^3 + 3a^2b + 3ab^2 + b^3$$

$$2. (a-b)^3 = a^3 - 3a^2b + 3ab^2 - b^3$$

$$3. a^3 + b^3 = (a+b)^3 - 3ab(a+b)$$

$$4. a^3 + b^3 = (a+b)(a^2 - ab + b^2)$$

$$5. a^3 - b^3 = (a-b)^3 + 3ab(a-b)$$

$$6. a^3 - b^3 = (a-b)(a^2 + ab + b^2)$$

Ex : 1+2+3+...+n Here, 1 to n all sum,

$$\text{even_cnt} = n / 2$$

$$\text{even sum} = \text{even_cnt} * (\text{even_cnt} + 1)$$

$$\text{odd_cnt} = (n + 1) / 2$$

$$\text{odd sum} = \text{odd_cnt} * \text{odd_cnt}$$

$$\begin{array}{ccccccc} & & 1 & & & & \\ & 1 & & 1 & & & \\ 1 & & 2 & & 1 & & \\ & 1 & & 3 & & 3 & \\ 1 & & 4 & & 6 & & 4 \\ & 1 & & 4 & & 6 & \\ \Rightarrow (x-2)^3 & = & x^3(-2)^0 & + & 3x^2(-2)^1 & + & 3x^1(-2)^2 & + & 1x^0(-2)^3 \end{array}$$

POWER :

$$1. a^m * a^n = a^{m+n}$$

$$2. a^m / a^n = a^{m-n}$$

$$3. (ab)^m = a^m * b^m$$

$$4. (a/b)^m = a^m / b^m$$

$$5. (a^m)^n = a^{mn}$$

$$6. m * \sqrt[n]{a} = a * (1/n)$$

$$7. a \neq 0 \text{ then, } a^0 = 1$$

LOG :

$$1. \text{Log } a^a = 1$$

$$2. \text{Log } a^1 = 0$$

$$3. \text{Log } a^a^n = m \text{ or, } n * \log a^a = m$$

$$4. \text{Log } (m/n) = \log a^m - \log a^n$$

$$5. \text{Log } a^m n = \log a^m + \log a^n$$

$$6. \text{Log } a^m a^n = n * \log a^m$$

$$7. \text{Log } a^m = \log a^m * \log a^b$$

RATIO :

$$(1) a:b = c:d \text{ So, } a/b = c/d \text{ So, } b/a = d/c$$

$$a:b = b+c$$

$$a/b = b/c$$

$$b^2 = ac$$

$$(2) a:b = c:d, a/b = c/d \text{ then, } a/c = b/d$$

Divisibility :

2 = last digit div by 2 or not

3 = sum of digit div by 3 or not

4 = last two div by 4 or not

5 = last dig 0 or 5

6 = If div by both 2 and 3

7 = last dig -> remove double it -> Subtract from remaining -> If by 7 or not

8 = last three dig div by 8 or not

9 = sum of dig div by 9 or not

$$11 = 4123 \% 11$$

$$= ((4*10^3) + (1*10^2) + (2*10^1) + (3*10^0))$$

$$= (4*(-1) + 1 + (2*(-1) + 3) \% 11$$

If this sum is div by 11 or not.

(odd index will mul by -1 and even will 1)

12 = div by 3 and 4 or not

Ex : n! Time d is it div by 3, 7, 9? (d>=1 && d<=9)

$$(1) (n! * d) \% 3 == 0 \text{ (d\%3 == 0) or } (n>=3)$$

$$(2) (d == 7 \&\& n == 3) \text{ OR } (n>2)$$

$$(3) (n < 6) \{$$

ll val = factorial(n);

ll lol = val * d;

If (digitSum(lol) \% 9 == 0) "divisible"

} else {

if(n >= 6) "Divisible"

Formula Area :

$$(1) \text{Square Corner} = n * \sqrt{2}$$

$$(2) \text{Area} = n^2$$

$$(3) \text{Porishima} = 4n$$

$$(4) \text{Rectangle Corner} = \sqrt{a^2+b^2}$$

$$(5) \text{Area} = a*b$$

$$(6) \text{Porishima} = 2(a+b)$$

$$(7) \text{Samakoni / triangle area} = (1/2)*a*b$$

$$(8) \text{Shomobahu} = \sqrt{3}/4 * a^2$$

$$(9) \text{Shomodibahu} = (b/4)\sqrt{4a^2-b^2}$$

$$(10) \text{triangle's area, } A = \sqrt{s(s-a)(s-b)(s-c)}$$

$$\text{semi perimeter of triangle, } S = (a+b+c)/2$$

(11) If we know three length a, b, c of triangle

inside circle area is,

$$R = (abc) / \sqrt{(a+b+c)(b+c-a)(c+a-b)(a+b-c)}$$

$$(12) \text{circle inside triangle} = A / S$$

$$\text{Bhaskara, } x = (-b \pm \sqrt{b^2 - 4ac}) / 2a$$

Modular Arithmetic :

(1) $(a + b) \% m = ((a \% m) + (b \% m)) \% m$
(2) $(a * b) \% m = (1l * (a \% m) * (b \% m)) \% m$
(3) $(a - b) \% m = ((a \% m) - (b \% m) + m) \% m$
(4) $(a^b) \% m = 1, b = 0, (a^{(b/2)})^2, b \text{ even}$
 $a^{(a^{((b-1)/2})^2)}, b \text{ odd}$

Number Theory

(1) Prime Check :

```
bool prime(int n) {
    for (int i = 2; i <= sqrt(n); i++) {
        if (n % i == 0) return false;
    }
    return true;
}
```

(2) sieve of Erothenotenes :

```
vector<bool> is_prime;
vector<ll> primes;

void sieve(ll n) {
    is_prime.resize(n + 1, true);
    is_prime[0] = is_prime[1] = false;

    for (ll i = 2; i * i <= n; i++) {
        if (is_prime[i]) {
            for (ll j = i * i; j <= n; j += i) {
                is_prime[j] = false;
            }
        }
    }

    for (ll i = 2; i <= n; i++) {
        if (is_prime[i]) {
            primes.push_back(i);
        }
    }
}
```

(3) Prime Fact :

```
void primeFact(int n) {
    for (int i = 2; i <= sqrt(n); i++) {
        if (n % i == 0) {
            int cnt = 0;
            while (n % i == 0) {
                cnt++;
                n = n / i;
            }
            cout << i << "^" << cnt;
            if (n > 1) cout << " * ";
        }
    }
    if (n > 1) {
        cout << n << "^1";
    }
}
```

(4) Binary Exponentiation :

```
ll Binary_expo(ll a, ll b, ll mod) {
    if (b == 0) return 1;
    if (b == 1) return a % mod;

    ll temp = power(a, b/2, mod);

    if (b % 2 == 0) {
        return (temp * temp) % mod;
    }
    else {
        return (((temp * temp) % mod) * a) % mod;
    }
}
```

(5) SOD, NOD :

```
void countAndSumDivisors(ll n) {
    set<ll> st;
    ll NOD = 0, SOD = 0;
    for (ll i = 1; i <= sqrt(n); i++) {
        if (n % i == 0) {
```

```
            ll x = i;
            ll y = n / i;
            st.insert(x), st.insert(y);

            SOD += x;
            if (x != y) SOD += y;
        }
    }
    NOD = st.size();
    cout << SOD << " " << NOD << endl;
}
```

(6) Segmented Sieve :

```
vector<int> primes;
// Regular sieve to find all primes up to sqrt(r)
void sieve(int limit) {
    vector<bool> isPrime(limit + 1, true);
    isPrime[0] = isPrime[1] = false;

    for (int i = 2; i * i <= limit; i++) {
        if (isPrime[i]) {
            for (int j = i * i; j <= limit; j += i) {
                isPrime[j] = false;
            }
        }
    }

    for (int i = 2; i <= limit; i++) {
        if (isPrime[i]) {
            primes.push_back(i);
        }
    }
}
// Seg. sieve to find primes in the range [l, r]
void segmentedSieve(ll l, ll r) {
    if (l == 1) l++;

    ll limit = sqrt(r) + 1;
    sieve(limit);

    int size = r - l + 1;
    vector<bool> isPrime(size, true);

    for (int p : primes) {
```

```

    ll start = max((ll)p * p, l + ((p - l % p) % p));
    for (ll j = start; j <= r; j += p) {
        isPrime[j - l] = false;
    }
}
for (int i = 0; i < size; i++) {
    if (isPrime[i]) {
        cout << l + i << endl;
    }
}
}

```

(7) GCD & LCM :

```

ll gcd(ll a, ll b) {
    while (b != 0) {
        ll temp = b;
        b = a % b;
        a = temp;
    }
    return a;
}

ll lcm(ll a, ll b) {
    return (a / gcd(a, b)) * b;
}

```

(8) Big Divisible :

```

void Big_div (string s, int m) {
    int ans = 0;

    for (int i = 0; i < s.size(); i++) {
        ans = ans * 10 + (s[i] - '0') % m;
    }
    if (ans % m == 0) {
        cout << "Yes" << endl;
    } else {
        cout << "No" << endl;
    }
}

```

(9) Big Mod :

```

int bigmod(string base, string expo, int mod) {
    if (mod == 1)
        return 0;

    int res = 1, base = 0, expNum = 0;
    for (int i = 0; i < base.size(); i++) {
        base = (base * 10 + (base[i] - '0')) % mod;
    }
    for (int i = 0; i < expo.size(); i++) {
        expNum = (expNum * 10 + (expo[i] - '0')) %
        (mod - 1);
    }
    while (expNum > 0) {
        if (expNum % 2 == 1) {
            res = (res * base) % mod;
        }
        expNum /= 2;
        base = (base * base) % mod;
    }
    return res;
}

int main() {
    string a, b = "1";
    ll mod;
    cin >> a >> mod; // mod = 1000000007
    cout << bigmod(a, b, mod) << endl;

    return 0;
}

```

(10) Big Subtractor :

```

string BigSub(string x, string y) {
    reverse(x.begin(), x.end());
    reverse(y.begin(), y.end());

    int len1 = x.size(), len2 = y.size();

    if (len1 > len2) {
        int gap = len1 - len2;

```

```

        while (gap--) {
            y.push_back('0');
        }
    } else if (len2 > len1) {
        int gap = len2 - len1;
        while (gap--) {
            x.push_back('0');
        }
    }
    int n = x.size(), carry = 0;
    string sub = "";

    for (int i = 0; i < n; i++) {
        int p = (x[i] - '0') - carry;
        int q = (y[i] - '0');

        if (p < q) {
            p += 10;
            carry = 1;
        } else {
            carry = 0;
        }
        int r = p - q;
        char lastDigit = r + '0';
        sub.push_back(lastDigit);
    }
    while (sub.size() > 1 && sub.back() == '0') {
        sub.pop_back();
    }
    reverse(sub.begin(), sub.end());
    return sub;
}

```

(11) Big Summation :

```

string BigSum(string x, string y) {
    reverse(x.begin(), x.end());
    reverse(y.begin(), y.end());
    int len1 = x.size(), len2 = y.size();

    if (len1 > len2) {
        int gap = len1 - len2;
        while (gap--) {
            y.push_back('0');

```

```

    }
} else {
    int gap = len2 - len1;
    while (gap--) {
        x.push_back('0');
    }
}
int n = x.size(). carry = 0;
string sum = "";

for (int i = 0; i < n; i++) {
    int p = (x[i] - '0') + carry;
    int q = (y[i] - '0');
    int r = p + q;
    char lastDigit = (r % 10) + '0';
    sum.push_back(lastDigit);
    carry = r / 10;
}
if (carry > 0) {
    sum.push_back(carry + '0');
}
reverse(sum.begin(), sum.end());

return sum;
}

```

(12) Big Multiplication :

```

string multiply(string num1, string num2) {
    int len1 = num1.size(), len2 = num2.size();
    string result(len1 + len2, '0');

    for (int i = len1 - 1; i >= 0; --i) {
        int carry = 0;
        for (int j = len2 - 1; j >= 0; --j) {
            int product = (num1[i] - '0') * (num2[j] - '0')
+ carry + (result[i + j + 1] - '0');
            carry = product / 10;
            result[i + j + 1] = (product % 10) + '0';
        }
        result[i] += carry;
    }
    result.erase(0, result.find_first_not_of('0'));
}

```

```

if (result.empty()) {
    return "0";
}
return result;
}

```

(13) Find Factorial :

```

ll multiply(ll x, ll res[], ll res_size) {
    ll carry = 0;
    for (int i = 0; i < res_size; i++) {
        ll prod = res[i] * x + carry;
        res[i] = prod % 10;
        carry = prod / 10;
    }
    while (carry) {
        res[res_size] = carry % 10;
        carry = carry / 10;
        res_size++;
    }
    return res_size;
}

void factorial(ll n) { // if N <= 15000
    ll res[5000];
    res[0] = 1;
    ll res_size = 1;

    for (int x = 2; x <= n; x++) {
        res_size = multiply(x, res, res_size);
    }

    cout << "Factorial of " << n << " is: ";
    for (int i = res_size - 1; i >= 0; i--) {
        cout << res[i];
    }
    cout << endl;
}

```

```

int main() {
    int N;
    cout << "Enter a number: ";
}

```

```

cin >> N;

factorial(N);
return 0;
}

```

Problems

(1) Least Prime Factor, Greatest Prime Factor, Distinct Prime Factor, Total Prime Factor, SOD, NOD.

```

const int N = 1e6+9;
int spf[N];
// Precompute the smallest prime factor (spf) for
every number up to N
void sieve() {
    for (int i = 2; i <= N; i++) {
        spf[i] = i;
    }
    for (int i = 2; i <= N; i++) {
        if (spf[i] == i) { // Only prime numbers
            for (int j = i * 2; j <= N; j += i) {
                spf[j] = min(spf[j], i);
            }
        }
    }
}
// Function to get the least prime factor (lpf) of a
number

int LPF(int x) {
    return spf[x];
}
// Function to get the greatest prime factor (gpf)
of a number

int GPF(int x) {
    int ans = 0;
    while (x > 1) {
        int p = spf[x];
    }
}

```

```

        ans = max(ans, p);
        while (x % p == 0) {
            x /= p;
        }
    }
    return ans;
}
// Function to get the number of distinct prime
factors

int distinctPrimeFactors(int x) {
    int distinct_count = 0;
    while (x > 1) {
        int p = spf[x];
        distinct_count++;
        while (x % p == 0) {
            x /= p;
        }
    }
    return distinct_count;
}
// Function to get the total number of prime
factors (including their powers)
int totalPrimeFactors(int x) {
    int total_count = 0;
    while (x > 1) {
        int p = spf[x];
        while (x % p == 0) {
            total_count++;
            x /= p;
        }
    }
    return total_count;
}
// Function to calculate the number of divisor of x

int numberOfDivisors(int x) {
    int num_divisors = 1;
    while (x > 1) {
        int p = spf[x];
        int power_of_p = 0;
        while (x % p == 0) {
            power_of_p++;
            x /= p;
        }
    }
}

```

```

    }
    num_divisors *= (power_of_p + 1);
}
return num_divisors;
}
// Function to calculate the sum of divisors of x

// sumOfDivisors(int x) {
//     int sum_divisors = 1;
//     while (x > 1) {
//         int p = spf[x];
//         int prime_power = 1;
//         while (x % p == 0) {
//             prime_power *= p;
//             x /= p;
//         }
//         sum_divisors *= (1LL * (prime_power * p - 1)
// / (p - 1));
//     }
//     return sum_divisors;
// }
int main() {
    sieve();

    int n;
    cin >> n;

    while (n--) {
        int q;
        cin >> q;
        int lpf = LPF(q);
        int gpf = GPF(q);
        int dist_prime_fact =
distinctPrimeFactors(q);
int total_prime_fact = totalPrimeFactors(q);
int NOD = numberOfDivisors(q);
// SOD = sumOfDivisors(q);

cout << lpf << " " << gpf << " " << dist_prime_fact
<< " " << total_prime_fact << " " << NOD << " "
<< SOD << endl;
    }
    return 0;
}

```

(2) Given an array of N length. Q query and in each query have given [L, R]. we need to print gcd excluding part from range [L, R] $1 \leq N$, $Q \leq 10^5$.

```

int main () {
    int tc;
    cin >> tc;

    while (tc--) {
        int n, q;
        cin >> n >> q;
        int arr[n+3];

        for (int i = 1; i <= n; i++) {
            cin >> arr[i];
        }
        int pre[n+2], post[n+2];
        pre[0] = 0, post[0] = 0;

        for (int i = 1; i <= n; i++) {
            pre[i] = gcd(pre[i-1], arr[i]);
        }

        for (int i = 1; i <= n; i++) {
            post[i] = gcd(post[i+1], arr[i]);
        }
        while (q--) {
            int l, r;
            cin >> l >> r;
            int ans = gcd(pre[l-1], post[r+1]);
            cout << ans << endl;
        }
    }
    return 0;
}

```

(3) Given A:B1, B2:C, Find A:C?

```

// a, b1, b2, c;
cin >> a >> b1 >> b2 >> c;

```

```

// g1 = gcd(a, b1);
a/= g1, b1 /= g1;

```

```

ll g2 = gcd(b2, c);
b2 /= g2, c /= g2;

if (b1 == b2) {
    cout << a << " " << c << endl;
} else {
    ll lcm_b = lcm(b1, b2);
    a *= (lcm_b / b1);
    c *= (lcm_b / b2);
    ll g = gcd(a, c);
    a /= g, c /= g;

    cout << a << " " << c << endl;
}

```

(4) You are given a positive integer N. Check whether the number N is representable as the sum of the cubes of two positive integers. such that, $a^3 + b^3 = N$.

```

void solve() {
    ll n, k = 0; cin >> n;
    for (ll i = 1; i <= cbrt(n); i++) {
        ll p = n - i * i * i;
        k = cbrt(p);
        if (k != 0 && k * k * k == p) {
            cout << "YES" << endl;
            break;
        }
        k = 0;
    }
    if (k == 0) {
        cout << "NO" << endl;
    }
}

```

(5) Given A, B find maximum gcd of x, y. $A \leq x, y \leq B$.

```

cin >> a >> b;
for (i = b; ; i--) {
    if ((a + i - 1) / i < (b / i)) {
        cout << i << endl;
    }
}

```

```

        break;
    }
}

(6) Count number of pair indices(i, j) such that
i < j and  $a_j - a_i = j - i$ ;

int n;
cin >> n;
ll arr[n+1];
map<ll, ll> mp;
for (int i = 1; i <= n; i++) {
    cin >> arr[i];
    cnt += mp[arr[i] - i]++;
}
cout << cnt << endl;

```

(7) Given array n & value m. we need to find how many consecutive subset are divisible by m. (Prefix Sum)

```

input -> n, m;
unordered_map<ll, ll> freq;
freq[0] = 0;
ll pref = 0, cnt = 0;
for (int i = 0; i < n; i++) {
    pref += v[i];
    ll rem = pref % m;
    if (rem < 0) rem += m;
    if (freq.count(rem) > 0) {
        cnt += freq[rem];
    }
    freq[rem]++;
}
cout << cnt << endl;

```

(8) Solve $x^2 + \sqrt{x} = c$ find value of X.

```

cin >> c;
double l = 0, r = 1e18;
for (int i = 0; i <= 200; i++) {
    double mid = (l + r) / 2.0;

```

```

    if ((mid * mid) + sqrt(mid) <= c) {
        l = mid;
        ans = mid;
    }
    else {
        r = mid;
    }
}
cout << setprecision(12) << ans << endl;

(9) A, B given now find the maximum gcd of x,
y. ( $A \leq x \leq y \leq B$ )

cin >> a >> b;

for (int i = b; ; i--) {
    if ((a + i - 1) / i < (b / i)) {
        cout << i << endl;
        break;
    }
}

```

Some Tricks :

Divisor :

1. How to find which number has odd number of divisor.

sol-> if ($\sqrt{n} * \sqrt{n} == n$) yes;

2. Print 1 to $1e12$ those have odd divisor.

sol-> $1^2, 2^2, 3^2, \dots$

3. given N and need to print number of divisor. sol-> N's prime factors each value power + 1.

->Precompute smallest prime factors (SPF) up to N (using sieve).

->Factorize N using SPF (take $O(n \log n)$)

->Calculate the number of divisors using the

formula.(total *= (count + 1))

4. N have x divisors, task is sum of these divisor,

sol->

$-2^3 * 5^2$
 $-1(1+5+5^2)+$
 $2(1+5+5^2)+$
 $2^2(1+5+5^2)+$
 $2^3(1+5+5^2)$
 $-(1+2+2^2+2^3) + (1+5+5^2)$

if : $p1^{e1} * p2^{e2} * p3^{e3}..$

formula : $((p1^{e1} - 1) / (p1 - 1)) * ((p2^{e2} - 1) / (p2 - 1)) * ...$

Prime :

1. Print 1 to N all number those have 1, 2, 3 divisor.

sol-> 1 have 1 divisor, prime number have 2 divisor, prime² have 3 divisor so on.

2. Need to print 1 to N all divisors and sum divisors of all value.

Code :

```
int N = 100;
int d[104], sum[104];
// O(N*log(logN))
```

```
for (int i=1; i<=N; i++) {
    for (int j=i; j<=n; j+=i) {
        d[j]++;
        sum[j] += i;
    }
}
```

3. For q query take N and print N's prime fact q, n in worst 1e6. (Sieve Factorization)

code : O(logN)
const int N = 1e6+9;

```
int spf[N];
```

```
int main() {
    for (int i=2; i<N; i++) {
        spf[i] = i;
    }
    for (int i=2; i<N; i++) {
        for (int j=i; j<N; j+=i) {
            spf[j] = min(spf[j], i);
        }
    }
    int q; cin >> q;

    while (q--) {
        int n; cin >> n;
        vector<ll> ans;
        while (n > 1) {
            ans.pb(spf[n]);
            n /= spf[n];
        }
        for (auto x : ans) {
            cout << x << " ";
        }
        cout << endl;
    }
}
```

Tricks :

1. How many Trailing zero have in N!

sol-> In N! have $10^a * 10^b$ for we find trailing zero. So, 10 mean $2^a * 5^b$. ans is min(a, b).

code :

```
int countTrailingZeros(int N) {
    int count = 0;
    for (int i = 5; N / i >= 1; i *= 5) {
        count += N / i;
    }
    return count;
}
```

2. How many divisor have in N!

sol-> find prime fact of N then apply,
 $N! = x^n * y^m * z^o = (n+1)*(m+1)*(o+1)$

3. Make a array where for all subarray product's are divisible by it's length.

sol-> if subarray length is divisible by any value inside subarray then it can.

ex : 1, 2, 3, 4 is divisible by 4.

4. L to R how many number divisible by m.

sol-> 1 to R divisible by m is cnt1, 1 to L-1 divisible by m is cnt2. now ans is (cnt1 - cnt2)

5. Given an array size N, $M=10^5$, $a_i \leq 10^9$.

print how many pair is $(a_i + a_j) \% M = 0$.

sol-> here, $(a_i + a_j) \% M = 0$ means, $a_i \% M = a_j \% M$.

we will go ahead and check how many present $(M - a[i])$ before $a[i]$.

0, 1, 2, 3, 1, 2, 5 (m = 5)

$a_0 = 0$,

$a_1 = 1$, $5-1 = 4$ not present

$a_2 = 2$, $5-2 = 3$ not present

$a_3 = 3$, $5-3 = 2$ present cnt = 1

$a_4 = 1$, $5-1 = 4$ not present

$a_5 = 2$, $5-2 = 3$ present cnt = 2

6. Given array need to print how many subarray are divisible by m.

sol-> first using prefix sum we will make sum array also will mod by m. then i will check how many have,

$p_j - p_i = 0$, ($p_j = p_i - 1$)

7. In N how many odd divisor exist.

sol-> Firstlt we will find N's prime fact then will just remove 2's part and others are ans.

8. Given N and need to say can we show n equal sum of 2 prime.

sol-> firstly we will find out all prime of N.
then, we will check $a+b = n$, $b = n-a$.

```
code :
sieve() {
...
}
bool ok (int n) {
    for (int i=2; i<n; i++) {
        if (spf[i]==i && spf[n-i]==n-i) {
            return true;
        }
    }
    return false;
}
```

9. Given N and need to present N in minimum prime number sum.

sol-> if N is prime then ans is 1, else we will take 2 and remain part $N-2 = X$.
if X not prime then we will try to make X into 2 prime.

10. How many digit in a number N!

sol-> if $n = 10^5$ then,
-log(n!) is huge. so,
-log₁₀($1*2*3*...*n$)
-log₁₀¹ + log₁₀² + ... + log₁₀ⁿ. check brute force then add 1. and show ans in floor value.

11. n, ai = 10^5 need to print maximum subset of number that's gcd is 1. else print -1.

sol-> we have to find just 2 number so that their gcd become 1.
hint : $2^0 + 2^1 + 2^2 + 2^3 + 2^4...+2^n < 2^{n+1}$.

12. Given N we need to show minimum power of 2. like, 10 : $2^1 + 2^3$, $2^1 + 2^2 + 2^2$. (ans is $2^1 + 2^3$)

sol-> we will find max power then remaining was 1 like, for 10 - $2^3 = 8$ remain 2^1 .

13. Given N. need to print a, b. so that Lcm(a, b) is smallest between $1 \leq a, b \leq n$ all possible pair.
($a+b == n$) $n \leq 10^9$. (the lcm of a, b is as small as possible).
sol-> if prime (1, n-1), if even ($n/2, n/2$), else find smallest (s),
then largest divisor ($d = n/s$) ans pair will be (d, n-d).

14. Given 2 array a, b. our task is output num of minimum inversion or swap so that a become b.
sol-> count at each index value that how many value are strictly greater then that previously value.

Graph

Graph representation,

```
list,
vector<int> graph[10000];
while (e--) {
    int u, v;
    cin >> u >> v;
    graph[u].push_back(v);
    graph[v].push_back(u);
}

for (int i = 1; i <= n; i++) { // highlight graph
    cout << i;
    for (int j = 0; j < graph[i].size(); j++) {
        cout << "->" << graph[i][j];
    }
    cout << endl;
}
```

Matrix,
int graph[n+1][n+1];

```
for (int i=1; i<=n; i++) {
    for (int j=1; j<=n; j++) {
        if (i == j) {
            graph[i][j] = 0;
        }
        else {
            graph[i][j] = -1;
        }
    }
}
while (e--) {
    int u, v; cin >> u >> v;
    graph[u][v] = 1;
    graph[v][u] = 1;
}
for (int i=1; i<=n; i++) {
    for (int j=1; j<=n; j++) {
        cout << graph[i][j] << " ";
    }
    cout << endl;
}
```

DFS

```
vector<int> graph[1002];
int vis [1003];

void DFS (int node){
    vis[node] = 1;
    cout << node << " -> ";

    for(int i = 0; i < graph[node].size(); i++){
        int child = graph[node][i];

        if(vis[child] == 0){
            DFS (child);
        }
    }
}
```

BFS

```
vector<int> graph[1002];
int vis[1002] = {0};
```

```
void BFS(int root) {
    queue<int> q;
    q.push(root);
```

```
    while(!q.empty()) {
        int node = q.front();
        q.pop();
```

```
        if(vis[node] != 0) {
            continue;
        }
        vis[node] = 1;
        cout << node << " -> ";
```

```
        for (int i=0; i<graph[node].size(); i++) {
            int child = graph[node][i];
```

```
            if (vis[child] == 0) {
                q.push(child);
            }
        }
    }
}
```

Connected Component

```
vector <int> v[100005];
int vis[100005];
```

```
void DFS (int node) {
    vis[node] = 1;
    for (int child : v[node]) {
        if (vis[child] == 0) {
            DFS(child);
        }
    }
}
```

```
int main () {
    ...
    for (int i = 1; i <= n; i++) {
        if (vis[i] == 0) {
            cnt++;
            DFS(i);
        }
    }
    cout << cnt << endl;
}
```

Bipartite Check

```
vector <int> graph[10001];
int vis[10001], col[10001];
```

```
bool dfs(int node, int c) {
    vis[node] = 1;
    col[node] = c;

    for (auto child : graph[node]) {
        if (vis[child] == 0) {
            if (dfs(child, c^1) == false) {
                return false;
            }
        }
        else {
            if (col[node] == col[child]) {
                return false;
            }
        }
    }
    return true;
}
```

```
int main () {
    ...
    if (dfs(1, 0) == true) {
        cout << "bipartite" << endl;
    }
    else {
        cout << "not bipartite" << endl;
    }
}
```

```
}
```

Cycle Detection

```
vector <int> graph[10001];
int vis[10001];
```

```
bool dfs (int node, int par) {
    vis[node] = 1;
```

```
    for (auto child : graph[node]){
        if (vis[child] == 0) {
            if (dfs(child, node) == true){
                return true;
            }
        }
        else {
            if (child != par) {
                return true;
            }
        }
    }
    return false;
}
```

```
int main () {
    ...
    if (dfs(1, 0) == false) {
        cout << "No cycle exist"<<endl;
    } else {
        cout << "Cycle exist"<<endl;
    }
}
```

Single source shortest path

```
vector<int> graph[10001];
int vis[10001], dist[10001];
```

```
void bfs(int node) {
    queue<int> q; q.push(node);
    vis[node] = 1, dist[node] = 0;
```

```

while (!q.empty()) {
    int current = q.front();
    q.pop();

    for (auto child : graph[current]) {
        if (!vis[child]) {
            vis[child] = 1;
            dist[child] = dist[current] + 1;
            q.push(child);
        }
    }
}

int main() {
    ...
    bfs(1);
}

```

Finding Diameter of Tree / Longest Path

```

vector<int> graph[10001];
int vis[10001];
int maxD, maxnode;

void dfs(int node, int d) {
    vis[node] = 1;

    if (d > maxD) {
        maxD = d;
        maxnode = node;
    }

    for (auto child : graph[node]) {
        if (vis[child] == 0) {
            dfs(child, d+1);
        }
    }
}

int main () {
    ...
    maxD = -1;
}

```

```

dfs(1, 0);

for (int i=1; i<=n; i++) {
    vis[i] = 0;
}

maxD = -1;
dfs(maxnode, 0);

cout << maxD << endl;
return 0;
}

```

Disjoint Set Union

```

const int N = 1e5 + 7;
int parent[N];
int component_size[N];

void make(int v) {
    parent[v] = v;
    component_size[v] = 1;
}

int find(int v) {
    if (v == parent[v]) {
        return v;
    }
    return parent[v] = find(parent[v]);
} // Path Compression

void Union(int a, int b) {
    a = find(a);
    b = find(b);

    if (a != b) { // Union by component_size
        if (component_size[a] <
            component_size[b]) {
            swap(a, b);
        }
        parent[b] = a;
        component_size[a] +=
            component_size[b];
    }
}

```

```

}
}

int main() {
    int n, q, cc = 0;
    cin >> n >> q;

    for (int i = 1; i <= n; i++) {
        make(i);
    }

    while (q--) {
        int t, u, v;
        cin >> t >> u >> v;

        if(t == 0) {
            Union(u, v);
        }
        else {
            if (find(u) == find(v)) {
                cout << 1 << endl;
            }
            else {
                cout << 0 << endl;
            }
        }
    }

    for (int i = 1; i <= n; i++) {
        if (find(i) == i) {
            cc++;
        }
    }

    cout << cc << endl;

    return 0;
}

```

Prim's Algorithm MST

```

vector<pair<int, int>> graph[1002];

```

```

int vis[1002] = {0};

int MST(int root) {
    priority_queue<pair<int, int>, vector<pair<int,
int>>, greater<pair<int, int>>> pq;
    int sum = 0;
    pq.push({0, root});

    while (!pq.empty()) {
        auto it = pq.top();
        pq.pop();
        int node = it.second;
        int wt = it.first;

        if (vis[node]) {
            continue;
        }
        vis[node] = 1;
        sum += wt;

        for (auto adj : graph[node]) {
            int adjNode = adj.first;
            int edgeWeight = adj.second;

            if (!vis[adjNode]) {
                pq.push({edgeWeight, adjNode});
            }
        }
    }
    return sum;
}

int main() {
    ...
    int result = MST(1);
    cout << "Total Cost : " << result << endl;
}

```

Array / Grid Direction

```
int dx[] = { +1, 0, -1, 0, +1, +1, -1, -1};
```

```

int dy[] = {0, +1, 0, -1, +1, -1, +1, -1};

int dx[] = { +0, +0, +1, -1, -1, +1, -1, +1};
//King's Move
int dy[] = { -1, +1, +0, +0, +1, +1, -1, -1};
//king's Move

int dx[] = { -2, -2, -1, -1, +1, +1, +2, +2};
//knight'sMove
int dy[] = { -1, +1, -2, +2, -2, +2, -1, +1};
///knight's Move

```

Example Code :

```

const int N = 507;
char graph[N][N];
bool vis [N][N];
int n, m, k;

int dx[] = {0, 0, -1, 1};
int dy[] = {1, -1, 0, 0};

bool valid(int x, int y) {
    return (x>=0 && x<n && y>=0 && y<m
&& !vis[x][y] && graph[x][y]!='.');
```

```

}

void dfs(int x, int y) {
    vis[x][y] = true;

    for (int i=0; i<4; i++) {
        int next_x = dx[i] + x;
        int next_y = dy[i] + y;

        if (valid(next_x, next_y)) {
            dfs(next_x, next_y);
        }
    }
    if (k > 0) {
        graph[x][y] = 'X';
        --k;
    }
}

```

```

}

int main() {
    cin >> n >> m >> k;

    for (int i=0; i<n; i++) {
        for (int j=0; j<m; j++) {
            cin >> graph[i][j];
        }
    }

    for (int i=0; i<n; i++) {
        for (int j=0; j<m; j++) {
            if (!vis[i][j] && graph[i][j]!='.') {
                dfs(i, j);
            }
        }
    }
}

```

Dijkstra

```

#define INF 1000000000
vector<pair<int, int>> graph[1001];
vector<int> dist(1009, INF);

int main() {
    int n, m;
    cin >> n >> m;

    while (m--) {
        int u, v, w;
        cin >> u >> v >> w;
        graph[u].push_back({v, w});
        graph[v].push_back({u, w});
    }
    priority_queue<pair<int, int>, vector<pair<int,
int>>, greater<pair<int, int>>> pq;
    pq.push({0, 1});
    dist[1] = 0;
}

```

```

while (!pq.empty()) {
    int node = pq.top().second;
    int curr_d = pq.top().first;
    pq.pop();

    for (auto child : graph[node]) {
        int child_node = child.first;
        int child_wt = child.second;

        if (dist[child_node] > dist[node] +
            child_wt) {
            dist[child_node] = dist[node] +
            child_wt;
            pq.push({dist[child_node],
            child_node});
        }
    }

    for (int i = 1; i <= n; i++) {
        cout << dist[i] << " ";
    }
    cout << endl;

    return 0;
}

```

Floyd warshall

```
const ll INF = 1e18;
```

```

int main() {
    ll n, m, q;
    cin >> n >> m >> q;
    ll graph[n+1][n+1];

    for(int i=1;i<=n;i++) {
        for(int j=1;j<=n;j++) {
            graph[i][j] = (i==j)?0:INF;
        }
    }

    for(int i=0; i<m; i++){

```

```

    ll v1, v2, w;
    cin >> v1 >> v2 >> w;

    graph[v1][v2] = min (graph[v1][v2], w);
    graph[v2][v1] = min (graph[v2][v1], w);
}
for(int k = 1; k<=n; k++) {
    for(int i=1; i<=n; i++) {
        for(int j=1; j<=n; j++) {
            if(graph[i][k] != INF && graph[k][j] !=
            INF) {
                graph[i][j] = min(graph[i][j],
                graph[i][k]+graph[k][j]);
            }
        }
    }
}
while (q--) {
    int u, v;
    cin >> u >> v;
    if(graph[u][v] >= INF) {
        cout << "-1" << endl;
    }
    else {
        cout << graph[u][v] << endl;
    }
}
for(int i=0;i<nodes;i++){
    for(int j=0;j<nodes;j++) {
        cout << graph[i][j] << " ";
    }
    cout << endl;
}
}

```

Bellman Ford

```

const long long INF = 1e18;
const int MAX_N = 1e5 + 5;

// Global variables
vector<pair<int, int>> adj[MAX_N];

```

```

vector<bool> vis(MAX_N, false);
vector<long long> dist(MAX_N, INF);

bool bellmanFord(int n, int src) {
    dist[src] = 0;
    // Relax all edges n-1 times
    for (int i = 1; i < n; i++) {
        for (int u = 1; u <= n; u++) {
            if (dist[u] == INF) continue;
            for (auto edge : adj[u]) {
                int v = edge.first;
                int w = edge.second;
                if (dist[u] + w < dist[v]) {
                    dist[v] = dist[u] + w;
                }
            }
        }
    }

    // Check for negative weight cycles
    for (int u = 1; u <= n; u++) {
        if (dist[u] == INF) continue;
        for (auto edge : adj[u]) {
            int v = edge.first, w = edge.second;
            if (dist[u] + w < dist[v]) {
                return true; // Negative cycle detected
            }
        }
    }

    return false; // No negative cycle
}

int main() {
    cin >> n >> m >> q;

    while(m--) {
        int v1, v2, w;
        cin >> v1 >> v2 >> w;
        adj[v1].push_back({v2, w});
        adj[v2].push_back({v1, w});
    }

    while (q--) {
        int u, v; cin >> u >> v;

```

```

fill(dist.begin(), dist.end(), INF);
fill(vis.begin(), vis.end(), false);
bool neg_cyc = bellmanFord(n, u);

if (neg_cyc) {
    cout << "contains negative cycle\n";
} else if (dist[v] == INF) {
    cout << "-1\n";
} else {
    cout << dist[v] << "\n";
}
}
}

```

Example :

1. Three musketeers. Choose three musketeers knowing each other and what is the minimum possible sum of their recognitions.

Code :

```

const int N = 1e4+7;
vector<int> g[N];
int vis[N], D[N], ans = INT_MAX;

void dfs(int node){
    vis[node] = 1;

    for (auto child1: g[node]) {
        if (child1 > node) {
            for (auto child2: g[child1]) {
                auto it = find(g[node].begin(), g[node].end(),
child2);
                if (child2 > child1 && it != g[node].end()) {
                    int rec = D[node]+D[child1]+D[child2];
                    ans = min(ans, rec);
                }
            }
        }
    }
}

int main() {

```

```

int n, m;
cin >> n >> m;
for (int i=1; i<=m; i++) {
    int a, b; cin >> a >> b;
    g[a].push_back(b);
    g[b].push_back(a);
    D[a]++, D[b]++;
}
for (int i=1; i<=n; i++) {
    if (!vis[i]) dfs(i);
}
if (ans != INT_MAX) {
    cout << ans - 6 << endl;
} else {
    cout << -1 << endl;
}
return 0;
}

```

2. A subtree of the tree is called balanced if the number of white vertices equals the number of black vertices. Count the number of balanced subtrees. output a single integer — the number of balanced subtrees.

code :

```

vector<bool> vis;
vector<vector<ll>> graph;
string s;
ll ans;

pair<ll, ll> dfs(ll head) {
    vis[head] = true;
    ll b = 0, w = 0;
    if (s[head - 1] == 'B') b++;
    else w++;

    for (int i=0; i<graph[head].size(); i++) {
        ll child = graph[head][i];
        if (vis[child]) {
            pair<ll, ll> temp = dfs(child);
            b += temp.first;

```

```

        w += temp.second;
    }
}
if (b == w) ans++;
return {b, w};
}

void solve() {
    ll n; cin >> n;
    vis.assign(n+1, false);
    graph.assign(n+1, vector<ll>());

    for (int i=2; i<=n; i++) {
        ll x; cin >> x;
        graph[x].push_back(i);
        graph[i].push_back(x);
    }
    cin >> s;
    ans = 0;
    dfs(1);
    cout << ans << endl;
}

```

3. Erdos p. Number (BFS)

code :

```

map<string, vector<string>> graph;
map<string, int> dist;
map<string, bool> vis;

void bfs(string start) {
    queue<string> q;
    dist[start] = 0;
    vis[start] = true;
    q.push(start);

    while (!q.empty()) {
        string curr = q.front();
        q.pop();

        for (auto child : graph[curr]) {

```

```

    if (!vis[child]) {
        vis[child] = true;
        dist[child] = dist[curr] + 1;
        q.push(child);
    }
}
}

int tc, papers, names;
string line;
cin >> tc;

for (int cs = 1; cs <= tc; cs++) {
    cin >> papers >> names;
    cin.ignore();

    graph.clear();
    dist.clear();
    vis.clear();

    for (int i = 0; i < papers; i++) {
        getline(cin, line);

        int colon_pos = line.find(':');
        if (colon_pos != string::npos) {
            line = line.substr(0, colon_pos);
        }
        int comma_count = 0;

        for (int j = 0; j < line.size(); j++) {
            if (line[j] == ',') {
                comma_count++;
                if (comma_count % 2 == 0) {
                    line[j] = ' ';
                }
            }
        }
    }
    vector<string> authors;
    stringstream ss(line);
    string last_name, initials;

```

```

        while (ss >> last_name >> initials) {
            authors.push_back(last_name + " " +
initials);
        }

        for (int j = 0; j < authors.size(); j++) {
            for (int k = j + 1; k < authors.size(); k++) {
                graph[authors[j]].push_back(authors[k]);
                graph[authors[k]].push_back(authors[j]);
            }
        }
    }

    bfs("Erdos, P.");

    cout << "Scenario " << cs << endl;
    for (int i = 0; i < names; i++) {
        getline(cin, line);
        if (dist.find(line) == dist.end()) {
            cout << line << " infinity" << endl;
        }
        else {
            cout << line << " " << dist[line] << endl;
        }
    }
}

```

Dynamic Programming

(1) Minimum Number of Coins to Make the Target

```

int n, k; cin >> n >> k;
vector<int> coins(n);
for (int i = 0; i < n; i++) {
    cin >> coins[i];
}
vector<int> dp(k + 1, INT_MAX);
dp[0] = 0; // 0 coins needed for amount 0

for (int i = 1; i <= k; i++) {

```

```

    for (int coin : coins) {
        if (coin <= i && dp[i - coin] != INT_MAX) {
            dp[i] = min(dp[i], 1 + dp[i - coin]);
        }
    }
}
if (dp[n] == INT_MAX) {
    cout << -1 << endl;
} else {
    cout << dp[n] << endl;
}

```

(2) Number of Ways to Make the Target

```

const int MOD = 1e9 + 7;
int n, k;
cin >> n >> k;

vector<int> coins(n);
for (int i = 0; i < n; i++) cin >> coins[i];
vector<int> dp(k + 1, 0);
dp[0] = 1; // 1 way to make amount 0

for (int coin : coins) {
    for (int i = coin; i <= k; i++) {
        dp[i] = (dp[i] + dp[i - coin]) % MOD;
    }
}
cout << dp[k] << endl;

```

(3) Minimum sum path in grid

```

int m, n; cin >> m >> n;
vector<vector<int>> grid(m, vector<int>(n));

for (int i = 0; i < m; i++) {
    for (int j = 0; j < n; j++) {
        cin >> grid[i][j];
    }
}
vector<vector<int>> dp(m, vector<int>(n, 0));
dp[0][0] = grid[0][0];
for (int i = 1; i < m; i++) { // first column

```

```

    dp[i][0] = dp[i-1][0] + grid[i][0];
}
for (int j = 1; j < n; j++) { // first row
    dp[0][j] = dp[0][j-1] + grid[0][j];
}
for (int i = 1; i < m; i++) { // remaining cells
    for (int j = 1; j < n; j++) {
        dp[i][j] = min(dp[i-1][j], dp[i][j-1]) +
        grid[i][j];
    }
}
cout << dp[m-1][n-1] << endl;

```

(4) Monkey Banana (Diamond shape)

```

ll arr[400][400];
ll dp[400][400];

ll monkey (ll i, ll j, ll k) {
    if (arr[i][j] == 0) {
        return 0;
    }
    if (i == k-1) {
        return arr[i][j];
    }
    if (dp[i][j] != -1) {
        return dp[i][j];
    }
    ll r1=0, r2=0;

    r1 = arr[i][j] + monkey(i+1, j, k);
    r2 = arr[i][j] + monkey(i+1, j+1, k);

    dp[i][j] = max(r1, r2);

    return dp[i][j];
}

memset(arr, 0, sizeof(arr));
memset(dp, -1, sizeof(dp));

ll n;

```

```

cin >> n;
ll k = 2 * n - 1;

for (int i = 0; i < n; i++) {
    for (int j = 0; j <= i; j++) {
        cin >> arr[i][j];
    }
}
ll l = 1;

for (int i = n; i < k; i++) {
    for (int j = l; j < n; j++) {
        cin >> arr[i][j];
    }
    l++;
}
ll ans = monkey(0, 0, k);
printf("Case %d: %d\n", it, ans);

```

(5) Greedy Algo Knapsack :

```

int n, w;
cin >> n >> w;
vector<int> weight(n), profit(n);
vector<double > ratio;

for (int i=0; i<n; i++) {
    cin >> weight[i];
}
for (int i=0; i<n; i++) {
    cin >> profit[i];
}
for (int i=0; i<n; i++) {
    double x = (double)profit[i] / weight[i];
    ratio.push_back(x);
}

for (int i=0; i<n-1; i++) {
    for (int j=i+1; j<n; j++) {
        if (ratio[i] < ratio[j]) {
            swap(ratio[i], ratio[j]);

```

```

        swap(profit[i], profit[j]);
        swap(weight[i], weight[j]);
    }
}

int sum=0, total_w=0;
for (int i=0; i<n; i++) {
    if (weight[i] < w) {
        sum += profit[i];
        w -= weight[i];
    }
    else {
        sum += w * ratio[i];
        w = 0;
    }
}
cout << sum << endl;

```

(6) 0/1 Knapsack

```

int n, w;
cin >> n >> w;

vector<ll> wt(n), val(n);

for (int i=0; i<n; i++) {
    cin >> wt[i];
}
for (int i=0; i<n; i++) {
    cin >> val[i];
}
vector<vector<ll>> dp(n+1, vector<ll>(w+1, 0));

for (int i=1; i<n+1; i++) {
    for (int j=0; j<w+1; j++) {
        dp[i][j] = dp[i-1][j];

        if (j-wt[i-1] >= 0) {
            dp[i][j] = max(dp[i][j], dp[i-1][j-wt[i-1]] +
            val[i-1]);

```



```

    }
}
cout << dp[n][w] << endl;

```

BIT Masking

1. You are given an array arr of 20 integer and another integer S. How many subset with sym equal to S?
code : (Approach: check all possible subset)

```

int n;
cin >> n;
vector<ll> v(n);

for (int i=0; i<n; i++) {
    cin >> v[i];
}
for (int mask=0; mask < (1<<n); mask++) {
    int sum = 0;

    for(int i=0; i<n; i++) {
        if ((mask >> i) && 1) {
            sum += v[i];
        }
    }
    if (sum == s) ans++;
}
cout << ans << endl;

```

Basic Stuff :

Sum of all pair,

1. XOR

```

ll Pair_Xor_Sum (int arr[], int n) {
    ll total_sum = 0;

```

```

    for (int bit=0; bit<32; bit++) {
        int cnt_set = 0;
        for (int i=0; i<n; i++) {
            if (arr[i] & (1 << bit)) {
                cnt_set++;
            }
        }
        int cnt_Unset = n - cnt_set;
        total_sum += (1LL << bit) * cnt_set *
cnt_Unset;
    }
    return total_sum;
}

```

2. OR

```

ll Pair_OR_Sum (int arr[], int n) {
    ll total_sum = 0;

    for (int bit=0; bit<32; bit++) {
        int cnt_set = 0;
        for (int i=0; i<n; i++) {
            if (arr[i] & (1 << bit)) {
                cnt_set++;
            }
        }
        total_sum += (1LL << bit) * cnt_set * n;
    }
    return total_sum;
}

```

3. AND

```

ll Pair_AND_Sum (int arr[], int n) {
    ll total_sum = 0;

    for (int bit=0; bit<32; bit++) {
        int cnt_set = 0;
        for (int i=0; i<n; i++) {
            if (arr[i] & (1 << bit)) {
                cnt_set++;
            }
        }
    }
}

```

```

    }
    total_sum += (1LL << bit) * (cnt_set *
(cnt_set - 1)) / 2;
}
return total_sum;
}

```

Sum of all subset, 1. AND

```

ll Subset_AND_Sum (int arr[], int n) {
    ll total_sum = 0;

    for (int bit=0; bit<32; bit++) {
        int cnt_set = 0;
        for (int i=0; i<n; i++) {
            if (arr[i] & (1 << bit)) {
                cnt_set++;
            }
        }
        total_sum += (1LL << bit) * ((1LL << cnt_set) - 1);
    }
    return total_sum;
}

```

2. OR

```

ll Subset_OR_Sum (int arr[], int n) {
    ll total_sum = 0;

    for (int bit=0; bit<32; bit++) {
        int subset = 0;
        for (int i=0; i<n; i++) {
            if (bit & (1 << i)) {
                subset |= arr[i];
            }
        }
        total_sum += subset;
    }
    return total_sum;
}

```

3. XOR

```

ll Subset_XOR_Sum (int arr[], int n) {
    ll total_sum = 0;

    for (int bit=0; bit<32; bit++) {
        int cnt_set = 0;
        for (int i=0; i<n; i++) {
            if (arr[i] & (1 << bit)) {
                cnt_set++;
            }
        }
        total_sum += (1LL << bit) * cnt_set * (1LL <<
(n-1));
    }
    return total_sum;
}

```

Searching & Sorting

1. Two Sum

```

int n = nums.size();
unordered_map<ll, ll> mp;

for (int i=0; i<n; i++) {
    ll diff = target - nums[i];

    if (mp.find(diff) != mp.end()) {
        return {mp[diff], i};
    }
    mp[nums[i]] = i;
}

```

Another,

given array, target
vector<pair<ll, ll>> ans;

```

for (int i=0; i<v.size(); i++) {
    ans.push_back({v[i], i});
}
sort(ans);
ll l = 0, r = ans.size()-1;

```

```

while (l < r) {
    ll curr = ans[l].first + ans[r].first;

    if (curr == target) {
        return {min(ans[l].second, ans[r].second),
max(ans[l].second, ans[r].second)
};
    }
    else if (curr > target) {
        r--;
    }
    else {
        l++;
    }
}

```

2. Three Sum

```

int n = nums.size();
sort(nums);
vector<vector<ll>> ans;

for (int i=0; i<n; i++) {
    if (i > 0 && nums[i] == nums[i-1]) {
        continue;
    }
    ll j = i+1, k = n-1;

    while (j < k) {
        ll sum = nums[i] + nums[j] + nums[k];
        if (sum > target) k--;
        else if (sum < target) j++;
        else {
            ans.push_back({nums[i], nums[j],
nums[k]});
            j++;
            k--;
            while(j<k && nums[j] == nums[j-1]) {
                j++;
            }
        }
    }
}

```

```

return ans;

```

3. Four Sum

```

int n = nums.size();
vector<vector<int>> ans;
sort(nums.begin(), nums.end());

for (int i = 0; i < n; i++) {
    if (i > 0 && nums[i] == nums[i - 1]) {
        continue;
    }
    for (int j = i + 1; j < n; j++) {
        if (j > i + 1 && nums[j] == nums[j - 1]) {
            continue;
        }
        int p = j + 1, q = n - 1;
        while (p < q) {
            long long sum =
(ll)nums[i]+(ll)nums[j]+(ll)nums[p]+(ll)nums[q];
            if (sum > target) {
                q--;
            } else if (sum < target) {
                p++;
            } else {
                ans.push_back({nums[i], nums[j],
nums[p], nums[q]});
                p++;
                q--;
                while (p < q && nums[p] == nums[p - 1]) {
                    p++;
                }
                while (p < q && nums[q] == nums[q + 1]) {
                    q--;
                }
            }
        }
    }
}
return ans;

```

4. We need to produce n copies of a document

using two copiers with different speeds. The first copier takes x seconds per copy, while the second takes y seconds. We can use both copiers simultaneously and can make copies from either the original or any existing copies. The challenge is to determine the minimum time required to make all n copies.

```
bool good(ll t) {
    if (t < min(x, y)) {
        return false;
    }
    ll total = 1;
    t -= min(x, y);
    total += floor(t/x) + floor(t/y);
    return total >= n;
}

void solve () {
    cin >> n >> x >> y;
    ll l=0, r = max(x,y)*n;

    while (l+1 < r) {
        ll mid = (l + (r-l)/2);
        if (good(mid)) {
            r = mid;
        }
        else {
            l = mid;
        }
    }
    cout << r << endl;
}
```

5. There are n rectangles of the same size: w in width and h in length. It is required to find a square of the smallest size into which these rectangles can be packed. Rectangles cannot be rotated.

```
ll w, h, n;

bool good (ll x) {
    return (x / w) * (x / h) >= n;
}
```

```
void solve() {
    cin >> w >> h >> n;
    ll l = 0, r = 1;

    while (!good(r)) {
        r *= 2;
    }
    while (r > l + 1) {
        ll mid = (l + r) / 2;
        if (good(mid)) {
            r = mid;
        } else {
            l = mid;
        }
    }
    cout << r << endl;
}
```

6. There are n ropes, you need to cut k pieces of the same length from them. Find the maximum length of pieces you can get.

```
ll n, k;
vector<ll> v;

bool good (double x) {
    int seg = 0;
    for (int i=0; i<n; i++) {
        seg += floor(v[i] / x);
    }
    return seg >= k;
}
```

```
void solve() {
    cin >> n >> k;
    v.resize(n);

    for (int i=0; i<n; i++) {
        cin >> v[i];
    }
    double l=0, r = 1e8;

    for (int i=0; i<100; i++) {
        double mid = (l + r) / 2;
```

```
        if (good(mid)) {
            l = mid;
        } else {
            r = mid;
        }
    }
    cout << setprecision(20) << l << endl;
}
```

Comparator Sort

```
// second value descending
(1) bool cmp(pair<int, int> a, pair<int, int> b) {
    return a.second > b.second;
}

// second value sort
(2) bool cmp(pair<int, int> a, pair<int, int> b) {
    return a.first < b.first;
}

// second value sort
(3) bool cmp(pair<int, int> a, pair<int, int> b) {
    if (a.first == b.first) {
        return a.second > b.second;
    }
}
```

Algorithms

1. KMP Algorithm

// O(m+n)
 // find how many times pattern exist in text
 // the task is to print the indexes of all the occurrences of pattern string in the text string. for printing, starting index of a string should be taken as 1.

```
#include <bits/stdc++.h>
```

```

using namespace std;

void search (string pat, string txt) {
    vector <int> lps(pat.size(), 0);

    for (int i = 1; i < pat.size(); i++) {
        int j = lps[i-1];

        while (j > 0 && pat[j] != pat[i]) {
            j = lps[j-1];
        }

        if (pat[j] == pat[i])j++;

        lps[i] = j;
    }

    int m = pat.size();
    int n = txt.size();
    int i = 0, j = 0, cnt = 0;
    vector <int> ans;

    while (i < n) {
        if (pat[j] == txt[i]) {
            i++;
            j++;
        }
        if (j == m) {
            cnt++;
            ans.push_back(i-(m-1)); // 3 - (-2) = 1
            j = lps[j-1];
        }
        else if (pat[j] != txt[i]) {
            if (j == 0)i++;
            j = lps[j-1];
        }
    }

    cout << "cnt : " << cnt << endl;

    for (auto it : ans) {
        cout << it << " ";
    }
}

```

```

    }
    cout << endl;
}

int main () {
    string txt, pat;
    cin >> txt >> pat;
    search(pat, txt);

    return 0;
}

```

2. Function to find the maximum subarray sum using Kadane's Algorithm

```

int maxSubarraySum(vector<int> &arr) {
    int max_curr = arr[0];
    int max_global = arr[0];

    for (int i=1; i<arr.size(); i++) {
        max_curr = max(arr[i], max_curr + arr[i]);

        if (max_curr > max_global) {
            max_global = max_curr;
        }
    }
    return max_global;
}

```
