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①

$\therefore T(n) = n$ for $n \leq 5$

① $T(n) = T(n-3) + 2$

Solution

$$T(n) = T(n-3) + 2$$

So, $T(n-3) = [T(n-3-3) + 2] + 2$

$$T(n-3) = [T(n-6) + 2] + 2$$

$$T(n-3) = T(n-6) + 4$$

Now Put $T(n-3)$ in eq.

$$T(n) = [T(n-3-6) + 4] + 2$$

$$T(n) = [T(n-9) + 4] + 2$$

$$T(n) = T(n-9) + 6$$

$$T(n) = T(n-3K) + 2K$$

As we can observe that it is expanding by a rate of $(n-K) + K$ Method # 1, 2

$$T(n) = T(n-K) + K$$

Now let's say $n-K=0$, $n=K$

$$T(n) = T(n-n) + n$$

$$= T(0) + n$$

As we know $T(n \leq 5) = n$

$$= n + n$$

$$= 2n \text{ (Ignore constants)}$$

$$T(n) = n \Rightarrow T(n) = \Theta(n) \Rightarrow T(n) = O(n)$$

$$T(n) = T\left(n - \frac{n}{3}\right) + 2\left(\frac{n}{3}\right)$$

$$T(n) = T\left(n - \frac{n}{3}\right) + 2\left(\frac{n}{3}\right)$$

$$T(n) = T(0) + 2\left(\frac{n}{3}\right)$$

$$T(n) = 0 + 2\left(\frac{n}{3}\right)$$

$$T(n) = O(n)$$

Tightest Bound.

$$\textcircled{2} \quad T(n) = T(n/10) + 1 \quad \textcircled{2}$$

Sol:

$$a=1, \quad b=10/9, \quad f(n)=1$$

$$n^{\log_b a} \quad n^{\log_{10/9} 1} = n^0 = n^1$$

$$n^{\log_b a} \neq f(n) \quad \text{but } n^0 = 1 \quad \text{so } 0 \neq 1.$$

Case #02 implies

$$O(n^{\log_b a} \log n) = O(1 \log n)$$

Choosing highest bound

$$\textcircled{O} O(\log n).$$

$$\textcircled{3} \quad T(n) = T\left(\frac{n^{1/2}}{2}\right) + 1$$

$$\text{let } n = 2^m$$

$$T(n^m) = \frac{2^{m/2}}{2} + 1$$

$$T(2^m) = S(m)$$

$$S(m) = T(m/2 \cdot 2) + 1$$

$$= T(m/4) + 1$$

$$m^{\log_b a} \Rightarrow m^{\log_2 2} = m^0 = 1$$

$$\left(n^{\log_b a} (\log n) \right) = m^{\log_2 2} (\log m) = 1 \log(m)$$

$$\text{but, } n = 2^m \quad \text{take log}$$

$$\log(n) = m \log_2 2 = m(1) \text{ Ans} = O(\log(\log(n))).$$

③

$$\textcircled{4} T(n) = 3T(n-1)$$

Sol:

$$T(n-1) = 3T(n-2)$$

$$\therefore T(n-2) = 3T(n-3)$$

$$T(n) = 3T(3T(n-2))$$

$$T(n-2) = 3T(n-3)$$

$$T(n) = 3T(3T(n-3))$$

expanding until K .

$$T(n) = 3^K T(n-K)$$

$$n-K = 0$$

$$n = K$$

$$= 3^K T(n-n)$$

$$= 3^K T(0) \quad \therefore T(n \leq 5) = n$$

$$= 3^K \cdot n$$

$$T(n) = n 3^K$$

$$T(n) = O(n 3^K)$$

(4)

$$(5) \quad T(n) = 2T(n/2) + 2$$

Sol:

$$n^{\log_b a} = n^{\log_2 2} = n^1 \Rightarrow n^{1-1} = n^0 = 1$$

So it is base $\neq 1$

$$n^{\log_b a - \epsilon} = f(n)$$

$$n^{1-1} = 1$$

$$1 = 1$$

$$T(n) = O(n) \therefore n \Rightarrow n^{\log_b a} \underline{n^1}$$

$$(7) \quad T(n) = 5T(n/3) + n$$

$$n^{\log_b a} = n^{\log_3 5} = n^{1.5}$$

$$n^{\log_b a - \epsilon} = f(n)$$

$$n^{1.5 - \epsilon} = f(n)$$

$$n = n$$

$$\text{So } = O(n^{\log_b a})$$

$$= O(n^{1.5})$$

$$= O(n^{3/2}) \text{ Ans.}$$

$$(8) \quad T(n) = 8T(n/4) + n^2$$

$$n^{\log_b a} = n^{\log_4 8} = n^{1.5}$$

$$n^{\log_b a - \epsilon} = f(n)$$

$$n^{1.5 - \epsilon} = f(n)$$

$$af(n/b) \leq cf(n)$$

$$8(n/4)^2 \leq cn^2$$

$$1/2 n^2 \leq \epsilon n^2$$

$$n^2/2 \leq n^2/2 \checkmark$$

So

$$T(n) = O(n^2)$$

⑤

$$\textcircled{9} \quad T(n) = 2T(n/9) + n \log n$$

$$\textcircled{a} \quad n^{\log_b a} = n^{\log_9 27} = n^{1.5}$$

$$n^{\log_b a - \epsilon} = f(n)$$

$$n \approx n \log(n)$$

$$T(n) = O(n^{3/2}) \quad \Rightarrow n^1 \log_b a$$

So, (10) $T(n) = 4T\left(\frac{n}{4}\right) + \frac{n}{\log}$ - (1)

f.l. $T(n) = O(n \log (\log n))$

$T(n) \leq c \cdot n \log (\log n)$ - (2)

$n = \frac{n}{4}$

$T\left(\frac{n}{4}\right) \leq c \frac{n}{4} \log \left(\log \left(\frac{n}{4}\right)\right)$

\Rightarrow (1)

$= c \frac{n}{4} \log \left(\log \left(\frac{n}{4}\right)\right)$

$T(n) = c n \log \left(\log \frac{n}{4}\right) + \frac{n}{4}$

$= c n \log (\log n - \log 4) + \frac{n}{\log(n)}$

$\log (\log n - \log 4) \approx \log (\log n)$

$= c n \log (\log n) + \frac{n}{\log(n)}$

$= n \log (\log n)$ proved

$$\textcircled{6} \quad T(n) = (T(\sqrt{n}))^2$$

Sol: Master theorem Not applied,

$$T(\sqrt{n}) = O(\log \sqrt{n}) = O\left(\frac{1}{2} \log n\right)$$

$$T(n) = \left(O\left(\frac{1}{2} \log n\right)\right)^2 = O\left(\left(\frac{1}{2} \log n\right)^2\right) = O(\log n)^2$$