Name:	Shagun Upadhyay
Roll No:	60
Class/Sem:	SE/IV
Experiment No.:	6
Title:	Prim's Algorithm.
Date of Performance:	
Date of Submission:	
Marks:	
Sign of Faculty:	



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Department of Artificial Intelligence & Data Science

Experiment No. 6

Title: Prim's Algorithm.

Aim: To study and implement Prim's Minimum Cost Spanning Tree Algorithm.

Objective: To introduce Greedy based algorithms

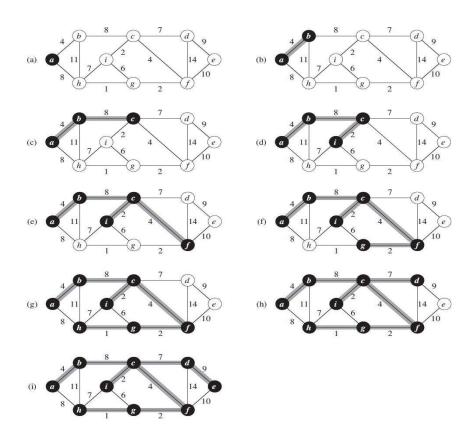
Theory:

Prim's algorithm is a greedy algorithm that finds a minimum spanning tree for a weighted undirected graph. This means it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all the edges in the tree is minimized. The algorithm operates by building this tree one vertex at a time, from an arbitrary starting vertex, at each step adding the cheapest possible connection from the tree to another vertex.

Example:



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Algorithm and Complexity:



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```
Algorithm Prim(E, cost, n, t)
2
    //E is the set of edges in G. cost[1:n,1:n] is the cost
3
    // adjacency matrix of an n vertex graph such that cost[i, j] is
    // either a positive real number or \infty if no edge (i, j) exists.
5
    // A minimum spanning tree is computed and stored as a set of
6
       edges in the array t[1:n-1,1:2]. (t[i,1],t[i,2]) is an edge in
7
       the minimum-cost spanning tree. The final cost is returned.
8
9
         Let (k, l) be an edge of minimum cost in E;
10
         mincost := cost[k, l];
        t[1,1] := k; t[1,2] := l;
11
12
        for i := 1 to n do // Initialize near.
13
             if (cost[i, l] < cost[i, k]) then near[i] := l;
             else near[i] := k;
14
15
        near[k] := near[l] := 0;
         for i := 2 to n-1 do
16
         \{ // \text{ Find } n-2 \text{ additional edges for } t. \}
17
             Let j be an index such that near[j] \neq 0 and
18
19
             cost[j, near[j]] is minimum;
20
             t[i,1] := j; t[i,2] := near[j];
21
             mincost := mincost + cost[j, near[j]];
22
             near[j] := 0;
23
             for k := 1 to n do // Update near[].
24
                  if ((near[k] \neq 0) and (cost[k, near[k]] > cost[k, j]))
25
                      then near[k] := j;
26
27
        return mincost;
28
```

Time Complexity is O(n2), Where, n = number of vertices**Theory:**

Implemenation:

```
#include #include <stdbool.h>
#include <stdio.h>
#include <stdio.h>
#define V 5
int minKey(int key[], bool mstSet[])
{
   int min = INT_MAX, min_index;
for (int v = 0; v < V; v++)
   if (mstSet[v] == false && key[v] < min)
      min = key[v], min_index = v;
   return min_index;
}
int printMST(int parent[], int graph[V][V])</pre>
```

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```
{ printf("Edge \tWeight\n"); for (int i = 1;
 i < V; i++) printf("%d - %d \t%d \n",
 parent[i], i,
    graph[i][parent[i]]);
void primMST(int graph[V][V])
 int parent[V];
 int key[V];
 mstSet[V];
 for (int i = 0; i < V; i++)
  key[i] = INT MAX, mstSet[i] = false;
 key[0] = 0;
 parent[0] = -1;
 for (int count = 0; count < V - 1; count++) {
 int u = minKey(key, mstSet);
 Set mstSet[u] = true;
  for (int v = 0; v < V; v++)
    if (graph[u][v] \&\& mstSet[v] == false
    && graph[u][v] \leq key[v]) parent[v] =
     u, key[v] = graph[u][v];
 }
 printMST(parent, graph);
    code
           int
main()
\{ \text{ int graph}[V][V] = \{ \{ 0, 2, 0, 6, 0 \}, \}
        \{2, 0, 3, 8, 5\},\
        \{0, 3, 0, 0, 7\},\
        \{6, 8, 0, 0, 9\}, \{
        0, 5, 7, 9, 0 \} \};
 primMST(graph);
 return 0;
```

Output:



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Edge	Weight
0 - 1	2
1 - 2	3
0 - 3	6
1 - 4	5

Conclusion: Implementing Prim's algorithm has proven to be effective in generating minimum spanning trees, efficiently connecting all nodes in a graph while minimizing total edge weight. This experiment underscores the algorithm's practical applicability in optimizing network connectivity, demonstrating its importance in various real-world scenarios.