Cryptographic Hash Functions

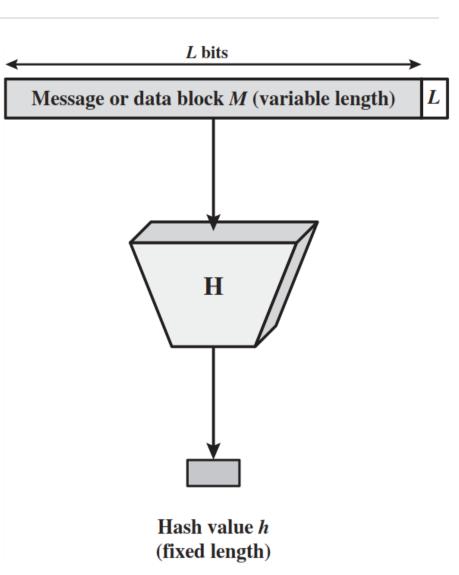


Outline

- Cryptographic Hash Functions
- Applications
- Simple hash functions
- Requirements and security
- Hash functions based on Cipher Block Chaining
- Secure Hash Algorithm (SHA)

Hash Function

- A hash function H accepts a variable-length block of data M as input and produces a fixed-size hash value h = H(M).
- A "good" hash function has the property that the results of applying a change to any bit or bits in M results with high probability, in a change to the hash code.



Applications of Cryptographic Hash Functions

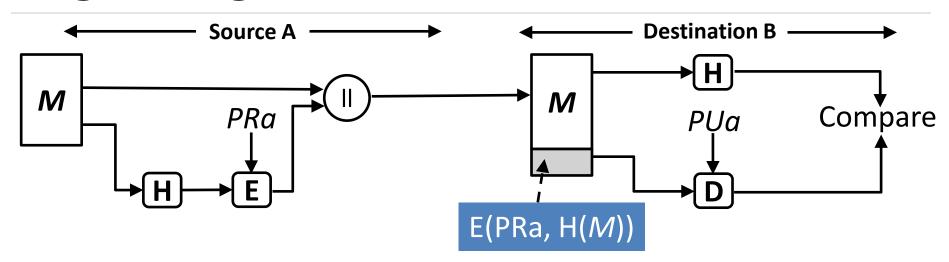
- 1. Message authentication
- 2. Digital Signature
- 3. One-way password file

Some topics of Unit 3/Unit 4 will be covered by the Expert

Digital Signature

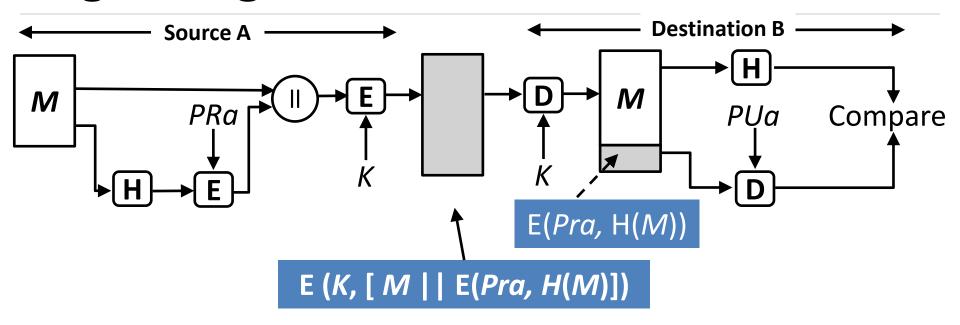
- A digital signature is a mathematical technique used to validate the authenticity and integrity of a message, software or digital document.
- The operation of the digital signature is similar to that of the Message Authentication Code (MAC).
- In the case of the digital signature, the hash value of a message is encrypted with a user's <u>private key</u>.
- Anyone who knows the user's public key can verify the integrity of the message that is associated with the digital signature.

Digital Signature method - 1



- The hash code is encrypted, using public-key encryption with the sender's private key. This provides authentication.
- It <u>also provides a digital signature</u>, because only the sender could have produced the encrypted hash code.

Digital Signature method - 2



If confidentiality as well as a digital signature is desired, then the message plus the private-key-encrypted hash code can be encrypted using a symmetric secret key.

Security Requirements

- 1. Disclosure
- 2. Traffic analysis
- 3. Masquerade
- 4. Content modification
- 5. Sequence modification
- 6. Timing modification
- 7. Source repudiation
- 8. Destination repudiation

Requirements for hash functions

- 1. It can be applied to any *sized message M*.
- 2. It should produce *fixed-length output h*.
- It should be easy to compute h=H(M) for any message M.
- 4. Given the hash value h, it is *infeasible* to find y such that (H(y) = h)
 - One-way property
- 5. For given block x, it is computationally infeasible to find y, $y \neq x$ with H(y) = H(x)
 - Weak collision resistance (i.e., it's hard to find another input 'y' that produces the same hash output)
- 6. It is computationally <u>infeasible to find messages m1 and m2 where their hash values are equal i.e. H(m1) = H(m2)</u>
 - Strong collision resistance

Simple Hash Function

- The input (message, file, etc.) is viewed as a sequence of *n*-bit blocks.
- The input is processed one block at a time in an iterative fashion to produce an n-bit hash.
- One of the simplest hash functions is the bit-by-bit exclusive-OR (XOR) of every block. This can be expressed as

$$C_i = b_{i1} \oplus b_{i2} \oplus ... \oplus b_{im}$$

where

 $C_i = i$ th bit of the hash code, 1 ... i ... n

m = number of n-bit blocks in the input

 $b_{ii} = i$ th bit in jth block

 \bigoplus = XOR operation

SHA

- In recent years, the <u>most widely used hash function</u> has been the Secure Hash Algorithm (SHA) as virtually every other widely used hash function had been found to have substantial cryptanalytic weaknesses
- SHA was developed by the National Institute of Standards and Technology (NIST) and published as a federal information processing standard (FIPS 180) in 1993.
- When weaknesses were discovered in SHA, different versions had been developed subsequently.

SHA - Secure Hash Algorithm

	SHA - 1	SHA - 224	SHA - 256	SHA - 384	SHA - 512
Message Digest Size (bits)	160	224	256	384	512
Message Size (bits)	< 2 ⁶⁴	< 2 ⁶⁴	< 2 ⁶⁴	< 2 ¹²⁸	< 2 ¹²⁸
Block Size (bits)	512	512	512	1024	1024
Word Size (bits)	32	32	32	64	64
No. of Steps	80	64	64	80	80

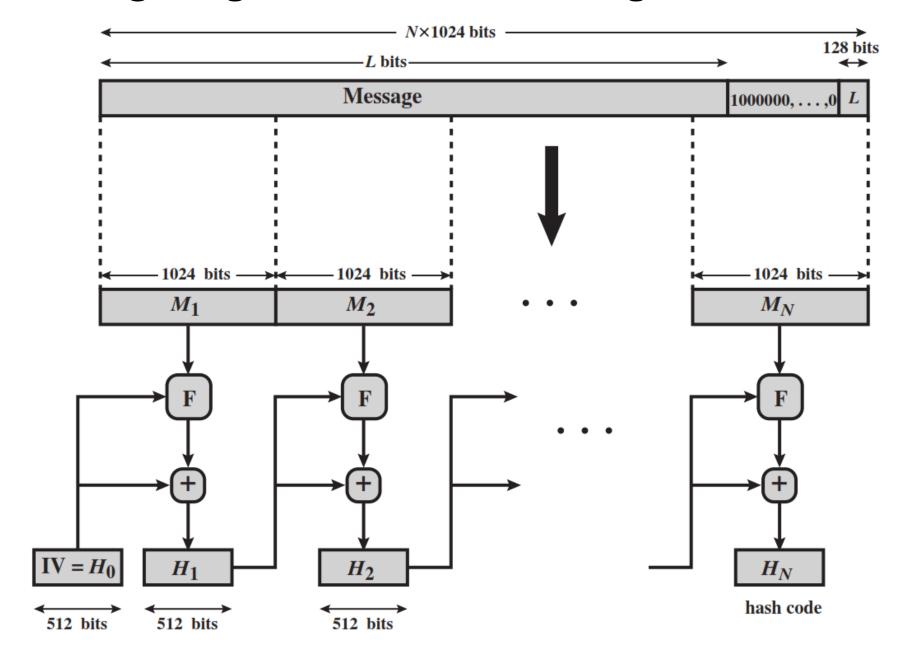
[•]Message digest=> Hash Value

[•]There is technically a limit for Message Size as per the padding scheme requirement

SHA - 512

- The algorithm takes as input a message with a maximum length of less than 2¹²⁸ bits
- It produces output of a 512-bit message digest.
- The input is processed in 1024-bit blocks.

Message Digest Generation using SHA -512

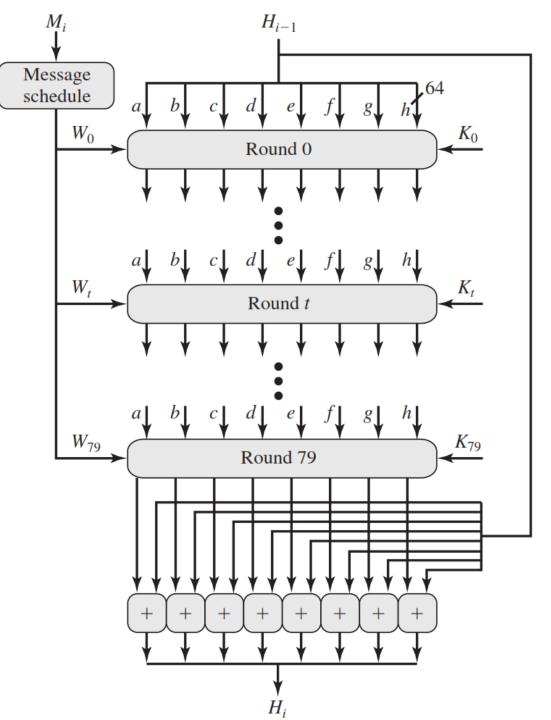


Step -1 Append Padding Bits

- The message is padded so that its length is congruent to 896 modulo 1024 [length ≡ 896(mod 1024)]. (To keep space for Appending Message Length-Check Step 2)
- Padding is always added, even if the message is already of the desired length.
- Thus, the number of padding bits is in the range of 1 to 1024.
- The padding consists of a single 1 bit followed by the necessary number of 0 bits.

Step -2 Append Length

- A block of 128 bits is appended to the message.
- This block is treated as an unsigned 128-bit integer (most significant byte first) and contains the length of the original message (before the padding).
- The <u>outcome</u> of the first two steps yields a message that is an <u>integer multiple of 1024</u> bits in length.
- In Figure, the expanded message is represented as the sequence of 1024-bit blocks M_1 , M_2 , ..., M_N , so that the <u>total length of the expanded message is N * 1024 bits</u>.



SHA-512 Processing of a Single 1024-Bit Block

Step -3 Initialize hash buffer

- The outcome of the first two steps produces a message that is an integer multiple of 1024 bits in length.
- the expanded message is represented as the sequence of 1024-bit blocks M_1 , M_2 , ..., M_N , so that the total length of expanded message is $N \times 1024$ bits.
- A 512-bit buffer is used to hold intermediate and final results of the hash function. The buffer can be represented as eight 64-bit registers (a, b, c, d, e, f, g, h).

Step -4 Process message in 1024-bit (128-word) blocks

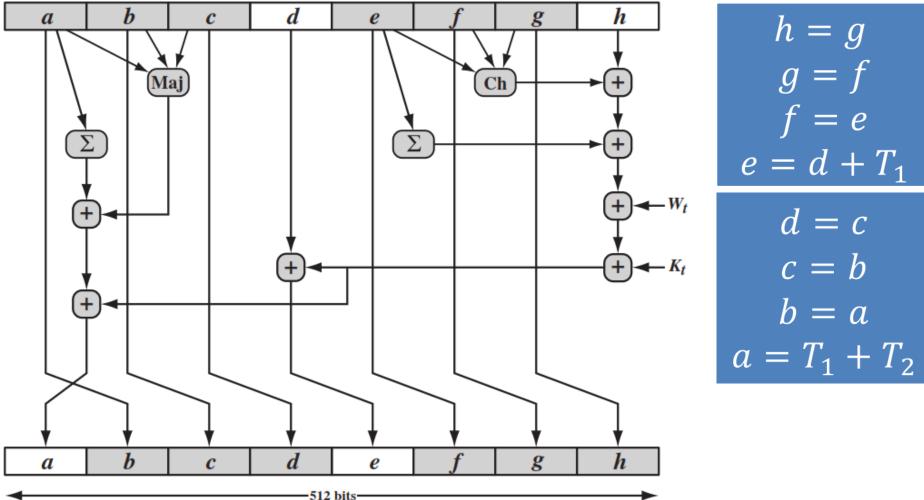
The heart of the algorithm is a module that consists of 80 rounds; this module is labelled F

SHA-512 Processing of a Single 1024-Bit Block

- Each round takes as input the 512-bit buffer value, abcdefgh, and updates the contents of the buffer.
- At input to the first round, the buffer has the intermediate hash value, H_{i-1} .
- Each round t makes use of a 64-bit value Wt, derived from the current 1024-bit block being processed.
- The output of the eightieth round is added to the input to the first round (H_i-1) to produce H_i .

Step – 5 Output

After all N 1024-bit blocks have been processed, the <u>output from</u>
 the Nth stage is the 512-bit message digest



$$h = g$$

$$g = f$$

$$f = e$$

$$e = d + T_1$$

$$d = c$$

$$c = b$$

$$b = a$$

$$T_1 = h + \text{Ch}(e, f, g) + \left(\sum_{1}^{512} e\right) + W_t + K_t$$

$$T_2 = \left(\sum_{0}^{512} a\right) + \text{Maj}(a, b, c)$$

SHA-512 Round **Function**

SHA-512 Round Function Elements

- Maj(a,b,c) = (a AND b) XOR (a AND c) XOR (b AND c) returns a result based on majority value among these inputs
- Σ (a) = ROTR(a,28) **XOR** ROTR(a,34) **XOR** ROTR(a,39) (ROTR(a,28) means rotate right by 28 positions)
- Σ (e) = ROTR(e,14) **XOR** ROTR(e,18) **XOR** ROTR(e,41)
- \blacksquare + = addition modulo 2^{64}
- $K_t = a 64$ -bit additive constant
- W_t = a 64-bit word derived from the current input block.