

Graph Theory

By

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Graph Theory

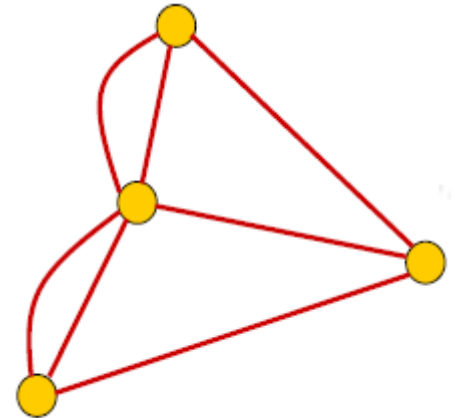
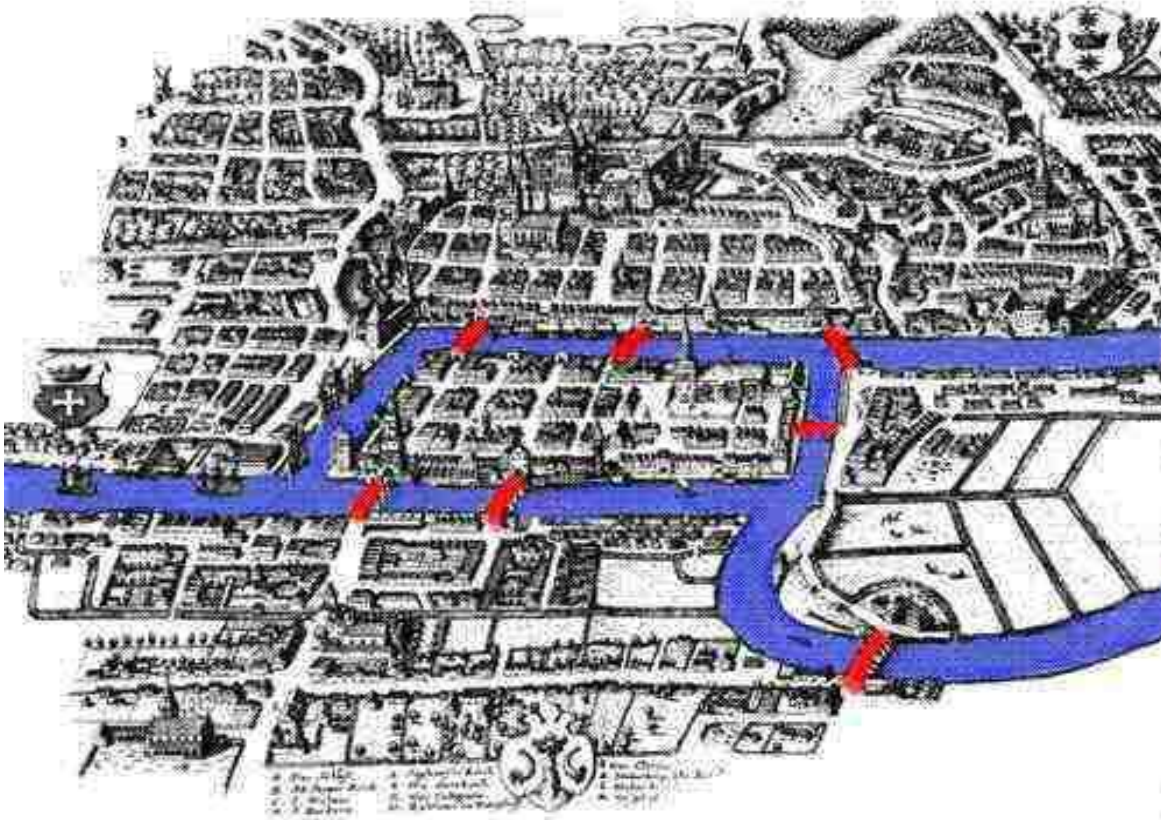
- What about this subject?
- How to study (or discover) it?
- What are the pre-requisites?
- What should be learning outcomes?
- Why you should study this?

Because it is beautiful?

A bit of History...

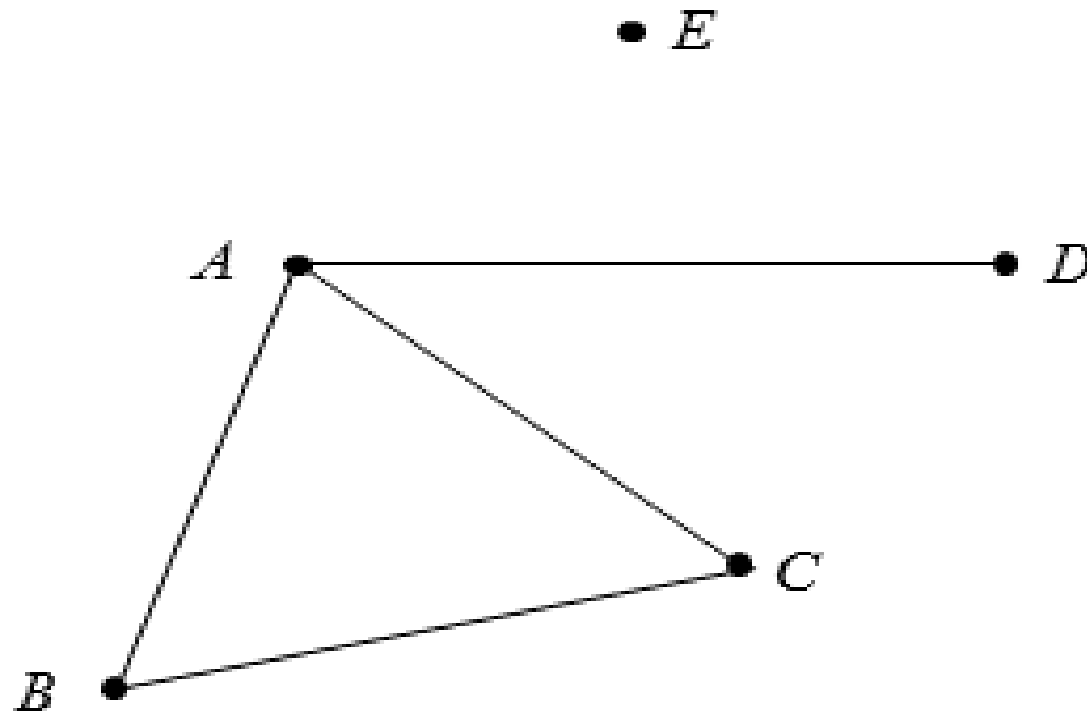
Father of graph theory, Euler

➤ Königsberg bridges problem (1736)



Definitions - Graph

- A graph is a pictorial and mathematical representation of a set of objects where some pairs of objects are connected by links.
- The interconnected objects are represented by points termed as vertices or nodes and the links that connect the vertices are called edges or arcs or lines.
- In other words, a graph is an ordered pair **$G = (V, E)$**



In the above graph,
 $V = \{A, B, C, D, E\}$
 $E = \{AB, BC, CA, AD\}$

Graph Theory

- Graph theory is the sub-field of mathematics and computer science which deals with graphs, diagrams that contain points and lines and which often pictorially represents mathematical truths.
- In short, graph theory is the study of the relationship between edges and vertices.

Fundamental Concepts of **Graph Theory**

- **Point.**

A **point** is a particular position that is located in a space.

A dot is used to represent a point in graph and it is labeled by alphabet, numbers or alphanumeric values.

Example

. P

Here, dot is a point labeled by 'p'.

Fundamental Concepts of **Graph Theory**

Line

Two points are connected to each other through a **line**.

A **line** is a connection between two points.

It is represented by a solid line.

Example



Here, 'A' and 'B' are the points and links between two points is called a line.

Fundamental Concepts of **Graph Theory**

Vertex

A **vertex** is a synonym of point in graph i.e. one of the points on which the graph is defined and which may be connected by lines/edges is called a vertex.

Vertex is also called "node", "point" or "junction". A vertex is denoted by alphabets, numbers or alphanumeric value.

Example

. V

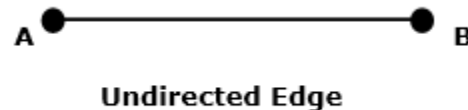
Here, point is the vertex labeled with an alphabet 'v'.

Fundamental Concepts of **Graph Theory**

Edge

- **Edge** is the connection between two vertices.
- Edge can either be **directed** or **undirected**

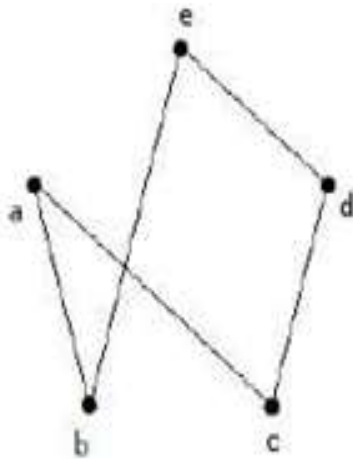
Example



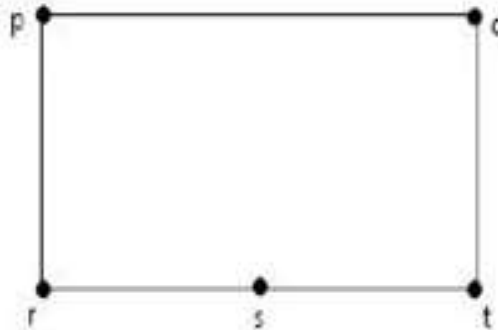
Here, '**A**' and '**B**' are the **vertices** and the link '**AB**' between them is called an **edge**.

Graph

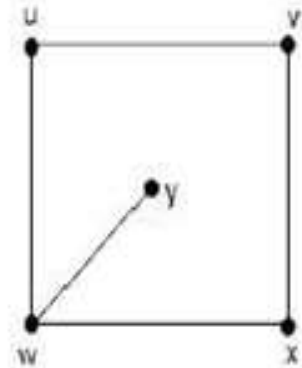
- A graph **G** is defined as $\mathbf{G} = \{\mathbf{V}, \mathbf{E}\}$ where **V** is a set of all vertices or points and **E** is the set of all edges in the graph.



G1



G2



G3

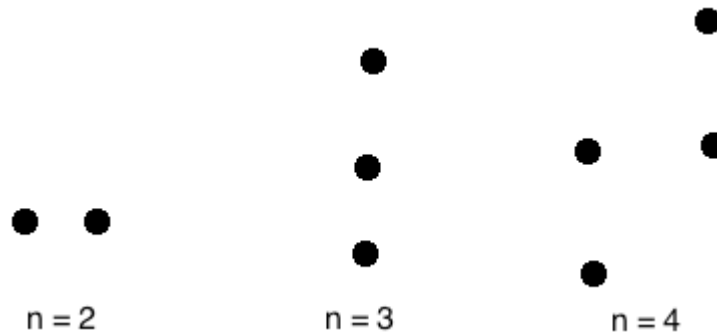
Types of Graph

- Different types of graphs
- depends upon the
 - number of vertices,
 - number of edges,
 - interconnectivity,
 - and their overall structure

Null Graph

A **null graph** is a graph in which there are no edges between its vertices. A null graph is also called empty graph.

Example



A null graph with n vertices is denoted by N_n .

Trivial Graph

A **trivial graph** is the graph which has only one vertex.

Example



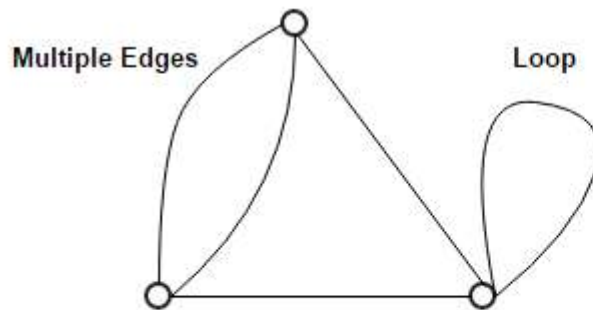
In the above graph, there is only one vertex 'v' without any edge. Therefore, it is a trivial graph.

Simple Graph

A **simple graph** is the undirected graph with **no parallel edges** and **no loops**.

A simple graph which has n vertices, the degree of every vertex is at most $n - 1$.

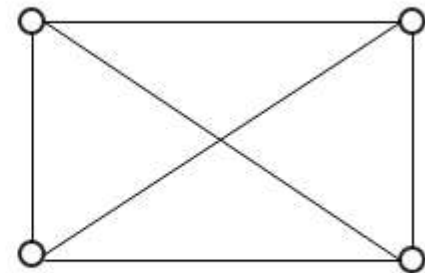
Example



Multiple Edges

Loop

Not a Simple Graph

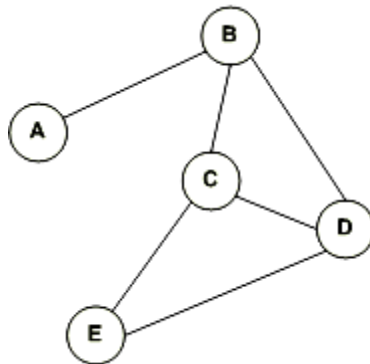


Simple Graph

Undirected Graph

An **undirected graph** is a graph whose edges are **not directed**.

- Example

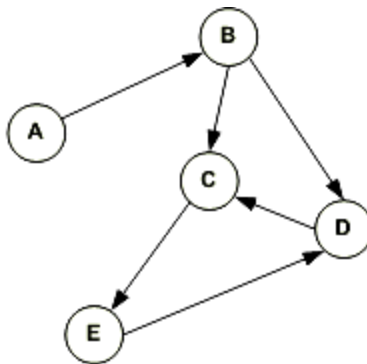


Directed Graph

A **directed graph** is a graph in which the **edges are directed** by arrows.

Directed graph is also known as **digraphs**.

- Example

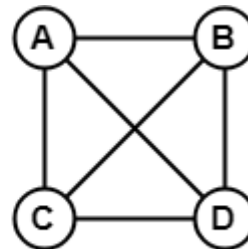
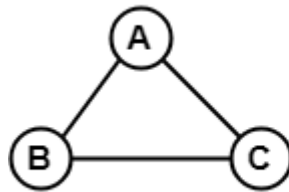


- A directed edge has an arrow from A to B, means A is related to B, but B is not related to A.

Complete Graph

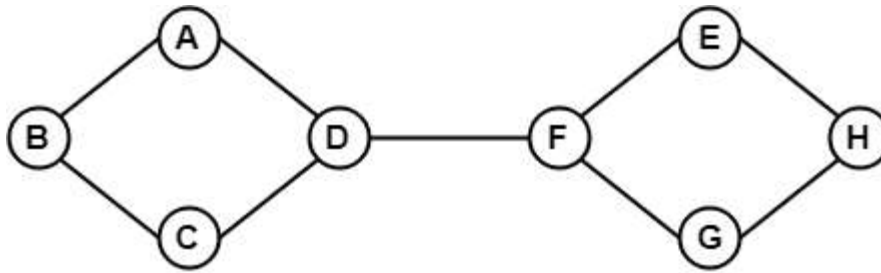
- A graph in which every pair of vertices is joined by exactly one edge is called **complete graph**. It contains all possible edges.

Example



Connected Graph

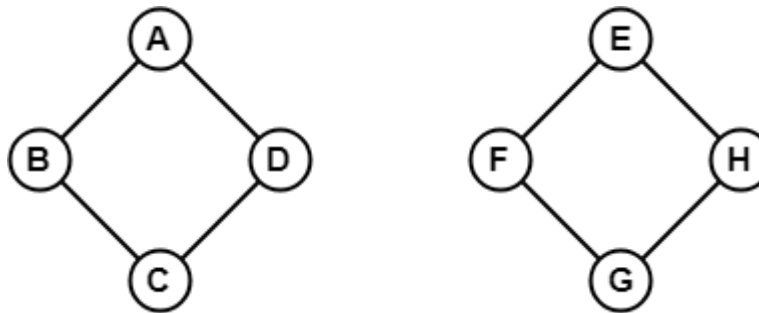
- A **connected graph** is a graph in which we can visit from any one vertex to any other vertex. In a connected graph, at least one edge or path exists between every pair of vertices.
- Example



Disconnected Graph

- A **disconnected graph** is a graph in which any path does not exist between every pair of vertices.

- Example

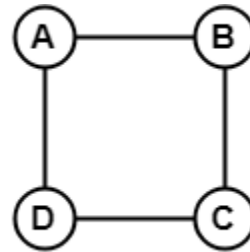
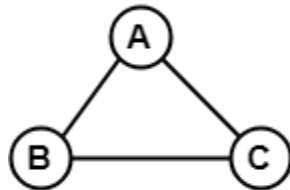


- The above graph consists of two independent components which are disconnected.

Regular Graph

- A **Regular graph** is a graph in which degree of all the vertices is same.
- If the degree of all the vertices is k , then it is called k -regular graph.

- Example



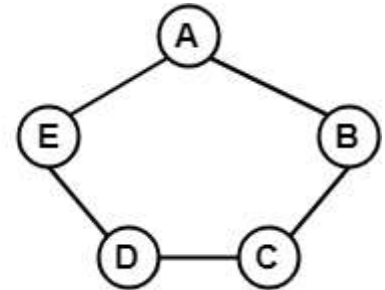
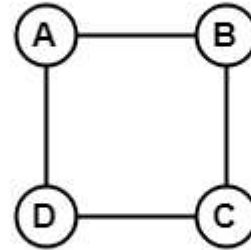
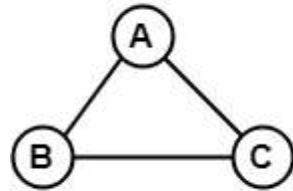
- In the above example, all the vertices have degree 2. Therefore they are called 2- **Regular graph**.

Cyclic Graph

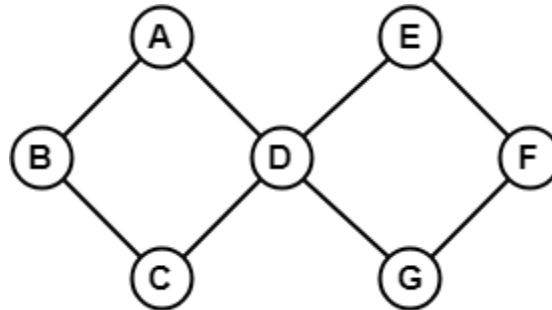
- A graph with ' n ' vertices (where, $n \geq 3$) and ' n ' edges forming a cycle of ' n ' with all its edges is known as **cycle graph**.
- A graph containing at least one cycle in it is known as a **cyclic graph**.
- In the cycle graph, degree of each vertex is 2.
- The cycle graph which has n vertices is denoted by C_n .

Cyclic Graph

- Example 1



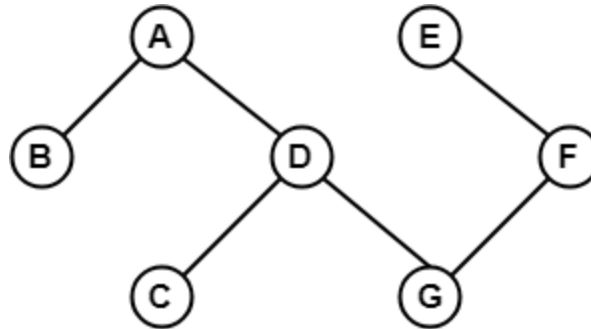
- In the above example, all the vertices have degree 2. Therefore they all are cyclic graphs.
- Example 2



- Since, the above graph contains two cycles in it therefore, it is a cyclic graph.

Acyclic Graph

- A graph which does not contain any cycle in it is called as an **acyclic graph**.
- Example



- Since, the above graph does not contain any cycle in it therefore, it is an acyclic graph.

Bipartite Graph

- A **bipartite graph** is a graph in which the vertex set can be partitioned into two sets such that edges only go between sets, not within them.
- A graph $G(V, E)$ is called bipartite graph if its vertex-set $V(G)$ can be decomposed into two non-empty disjoint subsets $V_1(G)$ and $V_2(G)$ in such a way that each edge $e \in E(G)$ has its one end point in $V_1(G)$ and other end point in $V_2(G)$.

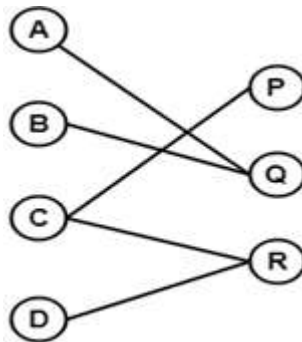
Bipartite Graph (Cont...)

- The partition $V = V1 \cup V2$ is known as bipartition of G .

- Example 1



- Example 2



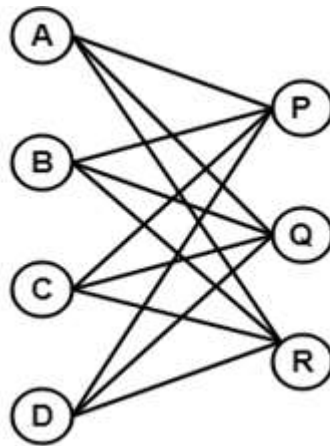
Complete Bipartite

- A **complete bipartite graph** is a bipartite graph in which each vertex in the first set is joined to each vertex in the second set by exactly one edge.
- A complete bipartite graph is a bipartite graph which is complete.

Complete Bipartite graph = Bipartite graph + Complete graph

Complete Bipartite (Cont...)

- Example



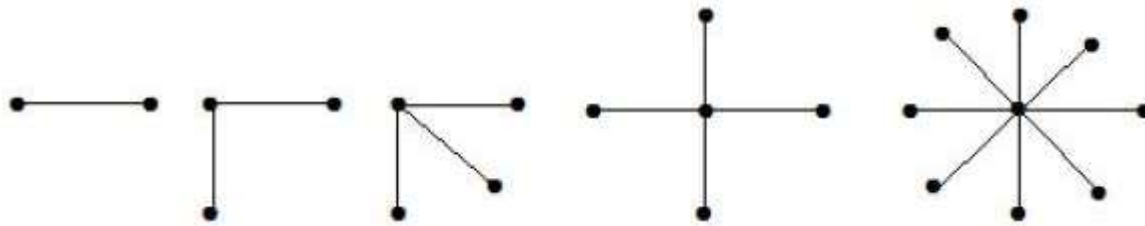
- The above graph is known as $K_{4,3}$.

Star Graph

- A **star graph** is a complete bipartite graph in which $n-1$ vertices have degree 1 and a single vertex have degree $(n - 1)$.
- This exactly looks like a star where $(n - 1)$ vertices are connected to a single central vertex.
- A star graph with n vertices is denoted by S_n .

Star Graph (Cont...)

- Example



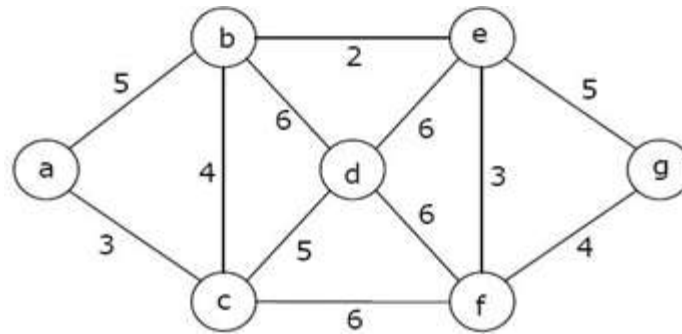
- In the above example, out of n vertices, all the $(n-1)$ vertices are connected to a single vertex. Hence, it is a star graph.

Weighted Graph

- A **weighted graph** is a graph whose edges have been labeled with some weights or numbers.
- The length of a path in a weighted graph is the sum of the weights of all the edges in the path.

Weighted Graph (cont...)

- Example

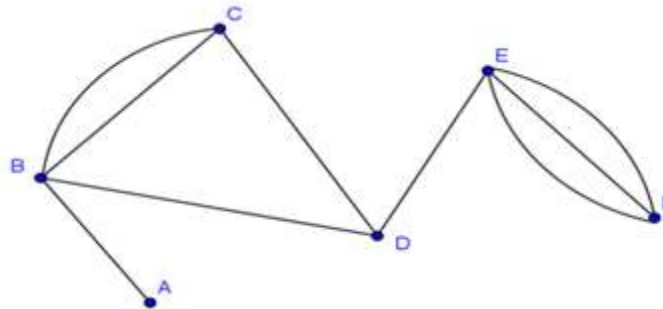


- In the above graph, if path is $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e \rightarrow g$ then the length of the path is $5 + 4 + 5 + 6 + 5 = 25$.

Multi-graph

- A graph in which there are multiple edges between any pair of vertices or there are edges from a vertex to itself (loop) is called a **multi - graph**.

- Example

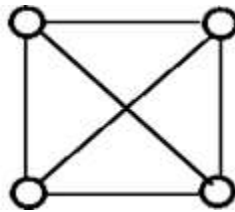


- vertex-set B and C are connected with two edges. Similarly, vertex sets E and F are connected with 3 edges. Therefore, it is a multi-graph.

Planar Graph

- A **planar graph** is a graph that we can draw in a plane in such a way that no two edges of it cross each other except at a vertex to which they are incident.

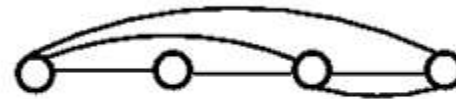
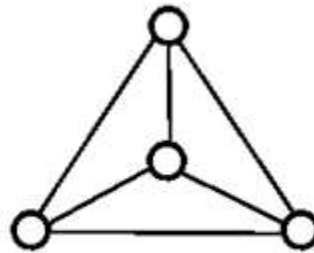
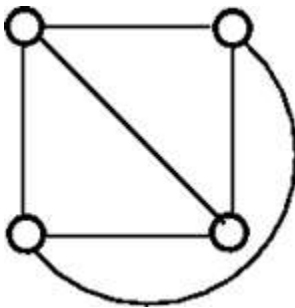
Example



- The above graph may not seem to be planar because it has edges crossing each other. But we can redraw the above graph.

Planar Graph (cont...)

- The three plane drawings of the above graph are:

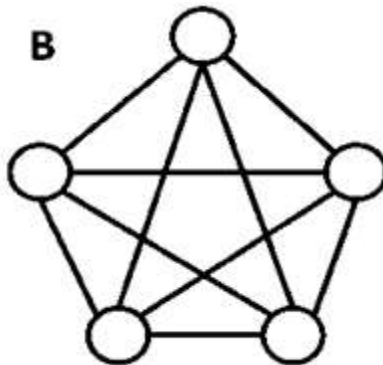


- The above three graphs do not consist of two edges crossing each other and therefore, all the above graphs are planar.

Non - Planar Graph

- A graph that is not a planar graph is called a non-planar graph. In other words, a graph that cannot be drawn without at least one pair of its crossing edges is known as non-planar graph.

Example



Applications of Graph Theory

Computer Science

➤ **study of algorithms** like:

- Dijkstra's Algorithm
- Prim's Algorithm
- Kruskal's Algorithm

➤ **flow of computation.**

➤ **networks of communication.**

➤ **data organization.**

➤ to find **shortest path in road** or a network.

➤ **Google Maps,**

Applications of Graph Theory (cont...)

- **Electrical Engineering**

- designing of circuit connections.

- **Linguistics**

- parsing of a language tree and grammar of a language tree.

- Semantics networks are used within **lexical semantics**

Applications of Graph Theory (cont...)

- Methods in **phonology** (e.g. theory of optimality, which uses lattice graphs) and **morphology** (e.g. morphology of finite - state, using finite-state transducers) are common in the analysis of language as a graph.

Applications of Graph Theory (cont...)

- **Physics and Chemistry**

- study molecules.
- The **3D structure of complicated simulated atomic structures** can be studied quantitatively by gathering statistics on graph-theoretic properties related to the topology of the atoms.
- **Statistical physics** also uses graphs

Applications of Graph Theory (cont...)

- to express **the micro-scale channels** of porous media, in which the vertices represent the pores and the edges represent the smaller channels connecting the pores.
- Graph is also helpful in **constructing the molecular structure** as well as lattice of the molecule.

Applications of Graph Theory (cont...)

- **Computer Network**

- the **relationships among interconnected computers** within the network, follow the principles of graph theory.
- **network security.**
- We can use the vertex coloring algorithm to find a proper **coloring of the map** with four colors.
- Vertex coloring algorithm may be used for assigning at most four different frequencies for any **GSM (Grouped Special Mobile) mobile phone networks.**

Applications of Graph Theory (cont...)

- **Social Sciences**

- Graph theory is also used in sociology. For example, to explore rumor spreading, or to measure actors' prestige notably through the use of social network analysis software.
- Acquaintanceship and friendship graphs describe whether people know each other or not.
- In influence graphs model, certain people can influence the behavior of others.
- In collaboration graphs model to check whether two people work together in a particular way, such as acting in a movie together.

Applications of Graph Theory (cont...)

- **Biology**

- Nodes in biological networks represent biomolecules such as genes, proteins or metabolites, and edges connecting these nodes indicate functional, physical or chemical interactions between the corresponding biomolecules.
- Graph theory is used in transcriptional regulation networks.
- It is also used in Metabolic networks.
- In PPI (Protein - Protein interaction) networks graph theory is also useful.
- Characterizing drug - drug target relationships.

Applications of Graph Theory (cont...)

- **Mathematics**

- In mathematics, operational research is the important field. Graph theory provides many useful applications in operational research. Like:
 - Minimum cost path.
 - A scheduling problem.

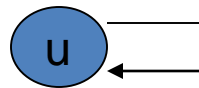
Applications of Graph Theory (cont...)

- **General**

Graphs are used to represent the routes between the cities. With the help of tree that is a type of graph, we can create hierarchical ordered information such as **family tree**.

Definitions – Edge Type

- **Loop:** A loop is an edge whose endpoints are equal i.e., an edge joining a vertex to it self is called a loop. Represented as $\{u, u\} = \{u\}$



- **Multiple Edges:** Two or more edges joining the same pair of vertices.

Adjacent Vertices

- A vertex u is said to be adjacent to vertex v if there is an edge $\{u, v\}$ in graph G .

Degree of a vertex

- Number of edges connected to a vertex x is known as the degree of vertex x in an undirected graph G .

Star graph

- A graph of p vertices in which one vertex has degree equal to $p-1$ while every other vertex has a degree equal to 1.
- It is known as a star graph because it looks like a star with rays of light coming out?

Walk

- You can walk on the edges of graph G edges starting from vertex u and ending at vertex v traversing different edges and vertices.
- In a walk you can traverse an edge more than once.
- A walk is open if vertex u and v are different. It is a closed walk if vertex u and v are the same.

Walk (cont...)

- Please note that in an un-directed graph you can traverse an edge in both directions but in a directed graph you can traverse an edge in only one direction - that is the direction of that edge.

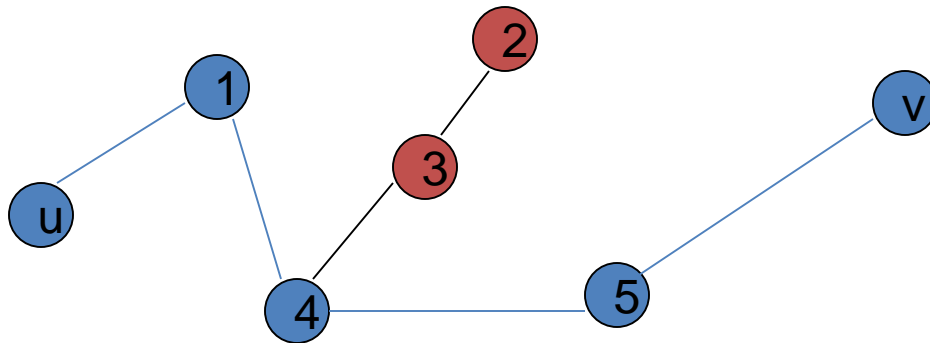
Trail and Circuit

- If no edge is repeated in a walk from a vertex u to a vertex v then the walk is known as a trail.
- A closed trail - that is when vertex u and v are the same then it is known as a circuit.
- Please note that it is possible to traverse a vertex more than once but an edge should not be traversed more than once in a trail or in a circuit.
- Thus a trail is always a walk but it is not the other way round.

Path

- If neither an edge nor a vertex is repeated in a walk starting from a vertex u and ending at vertex v then the walk is known as a path.
- It is known as a $u - v$ path.
- A path is always a trail (or a walk) but it is not the other way round.

- Representation example: $G = (V, E)$,
- Path P represented, from u to v is $\{\{u, 1\}, \{1, 4\}, \{4, 5\}, \{5, v\}\}$

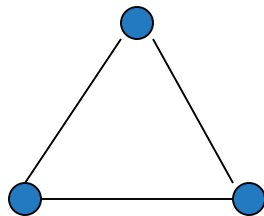


Shortest path

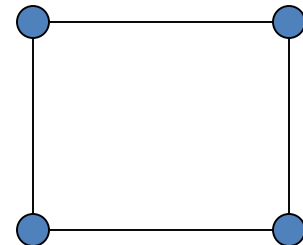
- Among all paths between vertex u and vertex v , the one with minimum length is known as the shortest path between vertex u and vertex v .
- In an unweighted graph G the minimum length is measured in terms of number of edges encountered in the $u - v$ path.
- In a weighted graph G the minimum length is measured in terms of sum of weights of all edges in the $u - v$ path.

Cycle

- If a path is closed that means you come back to the vertex from where you have started then that path is known as a cycle. A cycle is a circuit but a circuit may not be a cycle as no vertex should be repeated in a cycle.



C_3



C_4