

Left Recursive:

A grammar is said to be left Recursive if it is of the form.

$A \rightarrow A\alpha$

$A \rightarrow \beta$

Discussion

If we expand A, It will result in recursion as follows

$A \rightarrow A\alpha$

$A \rightarrow AA\alpha$

$A \rightarrow AA A\alpha$

$A \rightarrow AAAA\alpha$

$A \rightarrow AAAAA\alpha$

.....

To remove this problem, the given grammar is re-written as

$A \rightarrow \alpha A'$

$A' \rightarrow \beta/\epsilon$

The two grammars are defining same language i.e. both are same. The 1st one has left recursion and the 2nd one has left recursion.

Examples

Let we have

$E \rightarrow E+T$

$E \rightarrow T$

Now check if the left most "E" repeats itself then there is left recursion.

Here

$A=E$

$A'=E'$

$\alpha=+T$

$\beta=T$

so

$A \rightarrow \beta A'$

$A' \rightarrow \alpha A' / \epsilon$

$E \rightarrow TE'$

$E' \rightarrow +TE'$

Thus the correct grammar is

$E \rightarrow TE'$

$E' \rightarrow +TE'$

Let we have another example:

$X \rightarrow Xby$

$X \rightarrow a$

Eliminate Left Recursion:

Here

$A = X$

$A' = X'$

$\alpha = by$

$\beta = a$

$A \rightarrow \beta A'$

$A' \rightarrow \alpha A' / \epsilon$

$X \rightarrow aX'$

$X' \rightarrow byX' / \epsilon$