

5-lecture

Process States

Process creation,
Fork system call.

Ready Queue:-

⇒ All the processes that are ready to run will be in "Ready Queue"

⇒ Har device ki apni apni queue hoti hai.

Process states.

Mode switch

User Running

Sys call
or
interrupt

Kernel Running

⇒ Ye joh kernel mode mn
gaya ~~hoga~~ bhi. And ye Ausi
ek same process ki baat
horati, we can't say that
"Kernel apna kaam shuru
kardeg".

Created state.

⇒ Jab process pahli taraf
ata hai toh seetha to
rossly mn nhi jaega nq.

⇒ Aus waqt woh created
state mn ata. (Job memory
mapping ho yahi hoti hain)

⇒ Jab OS ausa initial resources
(Memory wagar) all of kardeg
hai, aus waqt hum

keh sakte hain. Ye ready state mn hoi

⇒ "House keeping"

⇒ Process jab khatam hota hai
toh exit kp ek call ah
hai. exp: int 91, ~~exit~~ exit

⇒ Ye saare ke saare ("syscall")
hain. Inki wajah se hum
jate hain ~~boot~~ kernel
mode m

⇒ Phir kernel usq clean
kardeta hai.

⇒ To the point def of (House keeping).

Jitna bhi apne pass resources
allocate kiya gaya thi
bhi cheez ko (In this case)
process. Aun saare ko
free up karne ko House keeping
kaha jata hai.

OS → Housekeeping (jhadu pherdega
hai)

Our House keeping (We do it with
no data loss)

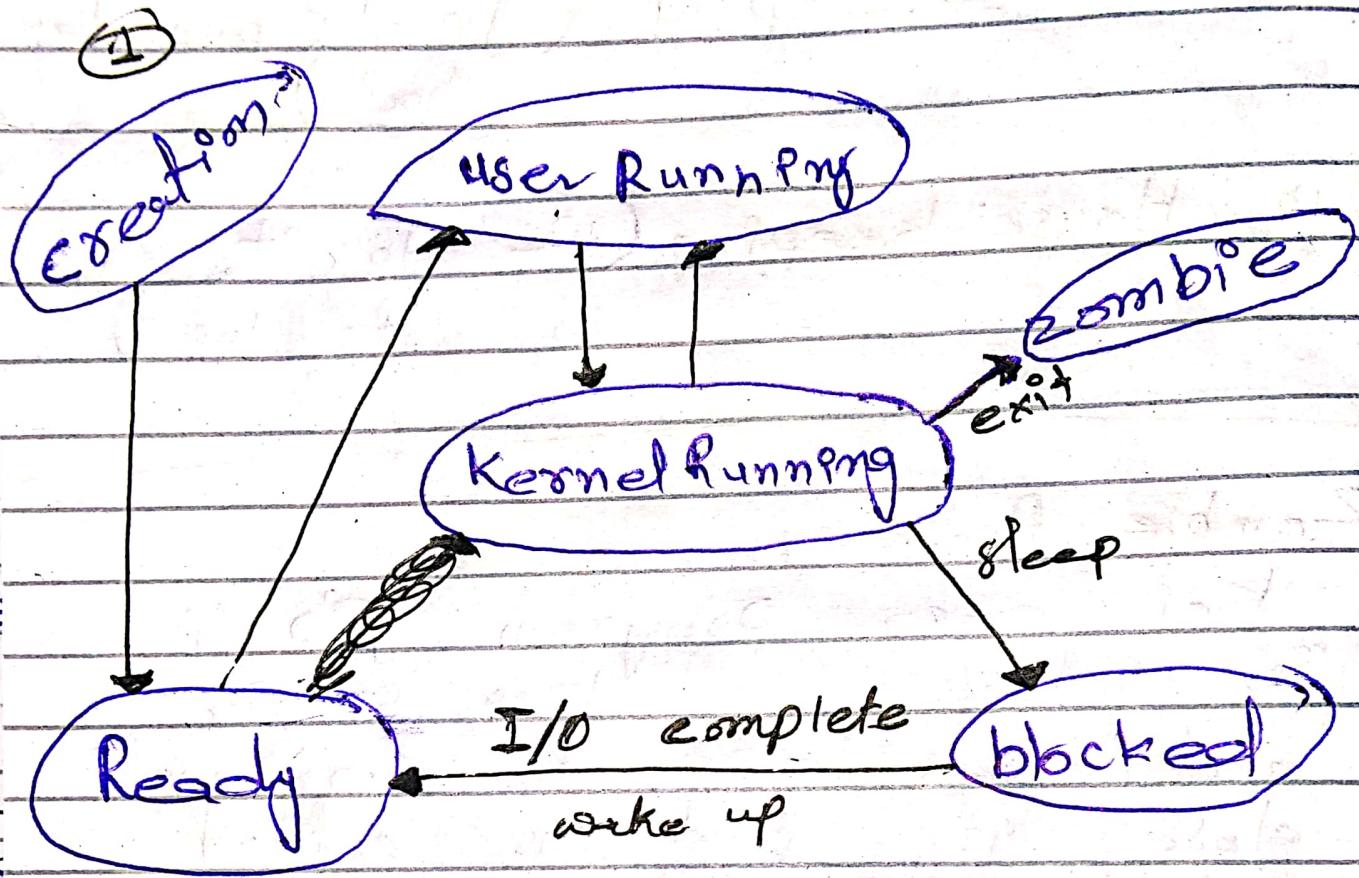
Zombie Process:-

Koi aesa program ~~kr~~ jise
OS close na kar sake
and oe have to go and
close qt, had is called
zombie process.

Zombie state:-

Whenever the process dies or
exits , it goes into ~~the~~
~~the~~ zombie state. Then some
house keeping is done, Then
the process is finished.
Tata, see u, bye :)

Big Picture



⇒ Has Process ki Information
in Terci cheegon mn store
hodi ha

Memory Map CPU PCB

(Process Creation)

⇒ Ek process pehle OS create karta hai, which is init.

⇒ When we click on Application us tym OS process ko create nhi karta. Woh kese hoga hain. Woh hum idhar dekhna chahiye hain.

Creating a Process under POSIX

⇒ Fork system call.

process execute

= =

= =

= =
Fork();



Kernel mode

is process kg

duplicate bongay
and Add

, dono processes
"Ready state"
mn chalega ye.

What happens??

① Check for available resources.

(Resources hain ke be nhi)



(Resources hain ke be nhi)

Ex:- ek user 500 process
create kar sakte hai

② Allocate a new PCB

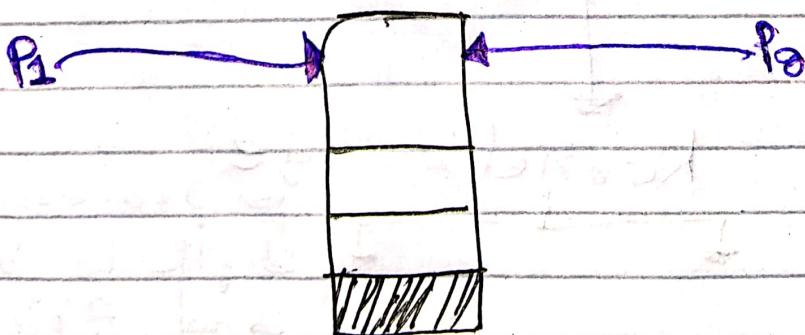
③ Assign a unique PID

④ Check process limit user

⑤ Set child state to "created"

⑥ Copy data from parent PCB slot
to child (Except "PID")

Abb yahan pe dono ki
memory Mapping bhi same
hoga kuch is tarha



⇒ Is it good thing ~~for multiple~~
both processes have same
Memory Mapping.

~~Ques~~

Ans 102

⇒ For Reading its fine coz.
 do data kii kya zarorat hai
 just read it from ones.

⇒ But For writing it's very Bad
 if ek process ~~ch~~ kuch change
 karega tbh dusra mn bhi
 automatically changes hojaegi.

⑦ Copy ON write (COW)
~~we need hardware support for this~~

⇒ Abb hoga ye jab thi ye
 read kar raha hoga, its
 fine, lekin ye jese
 e write korne ki koshish
 karega us tym, isko ek new
 data allocate kardiya jaegi. :)

⑧ Increment counts on current
 directory and open file &

???

⑨ Set child state to
 "ready to run".

~~Ques~~

⑯ Wait for the scheduler to run the process.

→ Scheduler ki manzi hai konse process ko pehla run karne hai

→ Hosakta hai parent pehla run hojai, maybe child. :)

⇒

Ex-1

```
int main()
```

```
{ int pid;
```

```
pid = fork();
```

```
if (pid == 0) you are its child  
    printf("child");
```

```
else if (pid == -1)
```

```
    printf("error");
```

```
else {
```

```
    printf("Parent of %d", pid);
```

You're the parent, and its value
is the pid of its child
only non-zero integer

Mn parent hu, lekin apko

pid child ke bhejo

gaya hai. :)

(Gret smashers)

main()

?

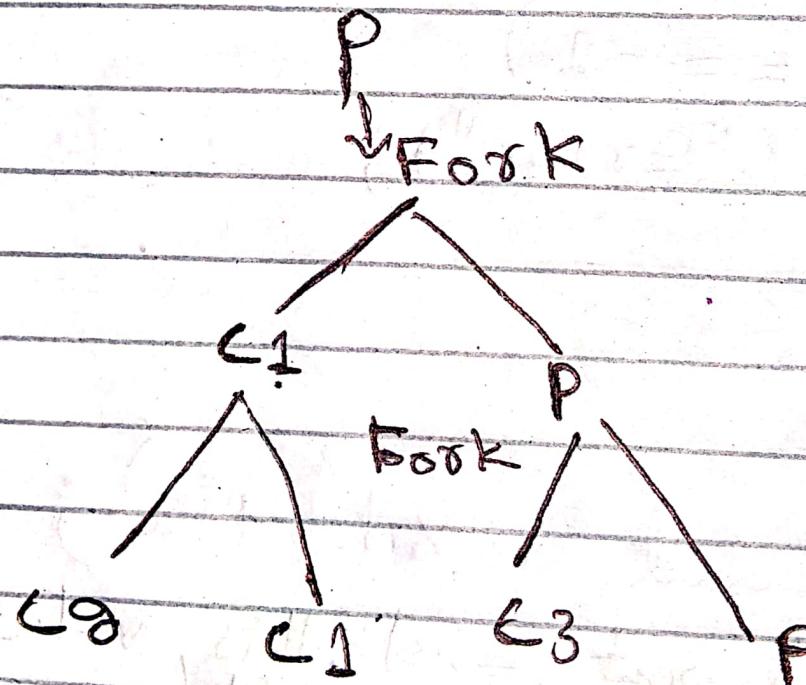
\$

Fork();

Fork();

printf("hello");

⇒ ?

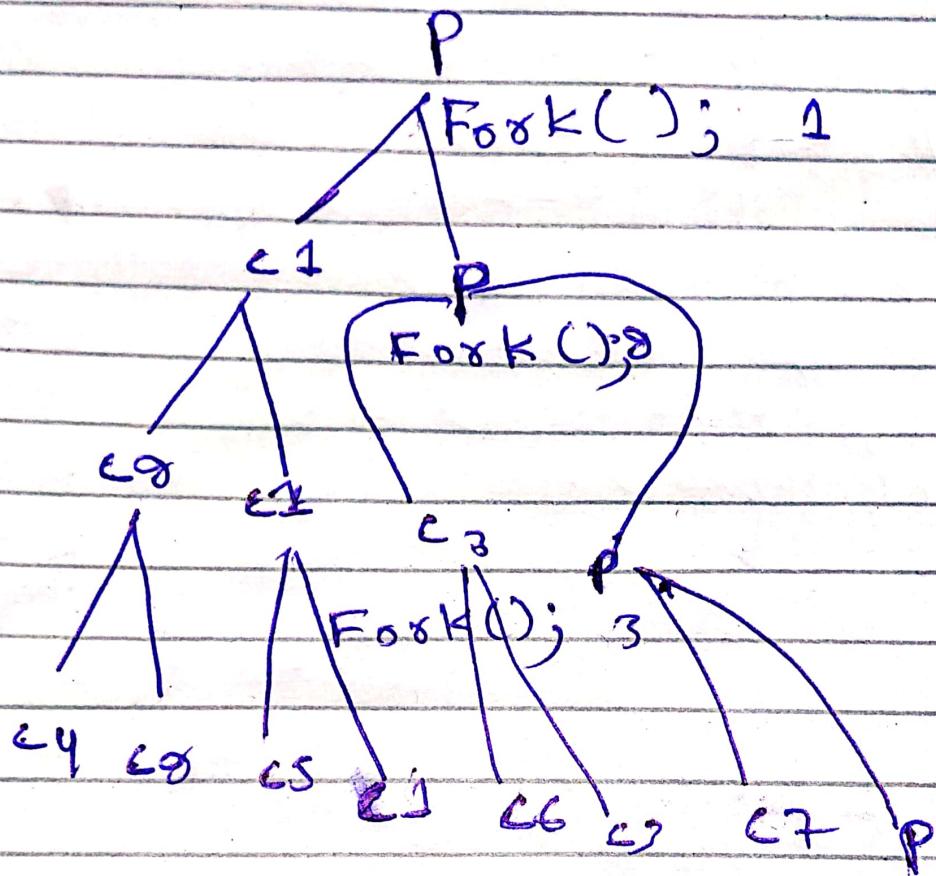


Fork();

Fork();

Fork();

printf("Hello world")



Child process, $2^n - 1$

Hello world will be
printed 2^n .