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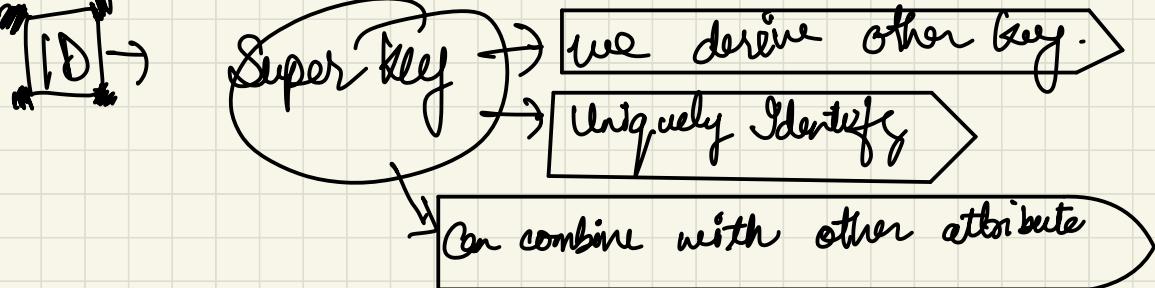
Difference between file system and dbms(DBMS vs file system)

File system	DBMS
File system is a collection of data. Any management with the file system, user has to write the procedures.	DBMS is a collection of data and user is not required to write the procedures for managing the database.
File system gives the details of the data representation and Storage of data.	DBMS provides an abstract view of data that hides the details.
In File system storing and retrieving of data cannot be done efficiently.	DBMS is efficient to use since there are wide varieties of sophisticated techniques to store and retrieve the data.
Concurrent access to the data in the file system has many problems like : Reading the file while other form of locking, deleting some information, updating some information	DBMS takes care of Concurrent access using some form of locking.
File system doesn't provide crash recovery mechanism. Eg. While we are entering some data into the file if System crashes then content of the file is lost	DBMS has crash recovery mechanism, DBMS protects user from the effects of system failures.
Protecting a file under file system is very difficult.	DBMS has a good protection mechanism.

Data
Representation

Concurrent
access

Crash
Recovery



Candidate key: Minimal chosen Super key is the Candidate key.

Candidate key

- ★ Superkeys:
 - {ID}, {SSN}, {ID, Name}, {ID, SSN}, {ID, Phone}, {Name, Phone}, {ID, Email}, {Name, SSN, Phone}, {Name, Email}, {ID, SSN, Phone}.....
- ★ Minimal super keys are called candidate keys.
- ★ Candidate Keys:
 - {ID}, {SSN}, {Name, Phone}, {Email}

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No repetition

19. Neo Academy.

Primary key: Unique but no null value

Chosen by DPA

Alternate key: Keys which are not primary

Unique key: Can also unique but also have null value

Composite key: key with more than one attribute

Foreign key:

Foreign key			
Student		Dept	
S.ID	Name	Dept_Code	Dept_Name
101	John	101	CSE
102	Robin	102	EEE
103	Alex	103	ECE
104	Vishal	104	MECH

Referential Integrity

SQL Commands

DDL - Data Definition Language

53 G Srinivas

- Create
- Alter
- Drop
- Truncate
- Rename

Statement delete
|| Data removal

Create table < table-name >

```

(
Column 1 name datatype,
Column 2 name datatype,
Column 3 name datatype( ) no constraint
);
desc table-name

```

Create table emp_detail

```

(
id int,
name varchar2(10),
Salary number(10)
);
desc emp;

```

TABLE EMP		
Column	Null?	Type
ID	-	NUMBER
NAME	-	VARCHAR2(10)
SALARY	-	NUMBER(10,0)

Download CSV
3 rows selected.

```

Autocommit Display 10
create table employee
(
id int,
name varchar(10)
);
alter table employee add address varchar(10);
desc emp11
alter table employee drop column address;
alter table employee modify id varchar(30);
alter table employee rename column id to roll_no;
alter table employee rename to emp11;
alter table emp11 add primary key (roll_no)

```

Results Explain Describe Saved SQL History

Object Type TABLE Object EMP11

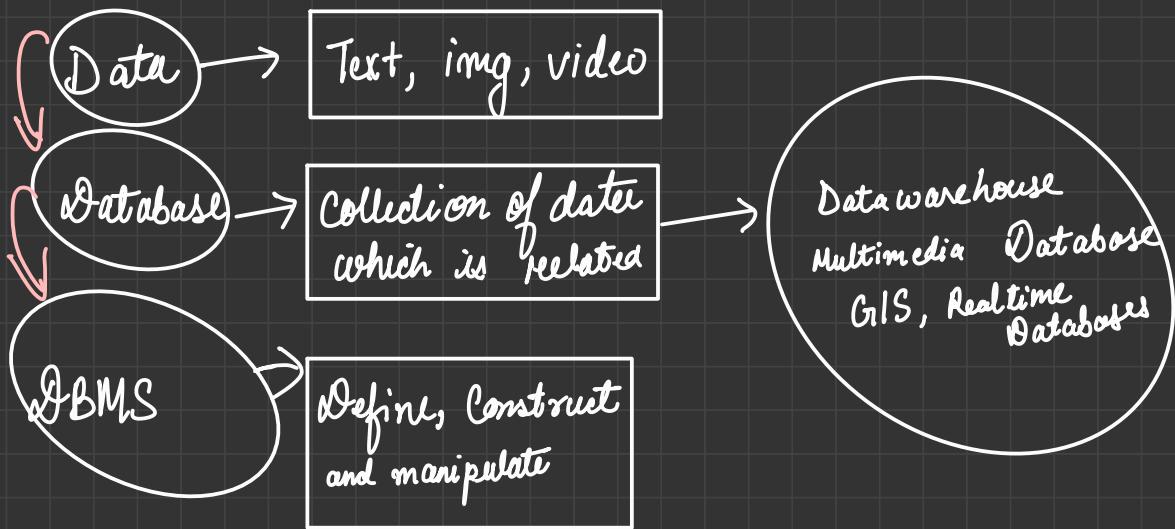
Table	Column	Data Type	Length	Precision	Scale	Primary Key	Nullable	Default	Comment
EMP11	ROLL_NO	Varchar2	30	-	-	1	-	-	
	NAME	Varchar2	10	-	-	-	✓	-	

1 - 2

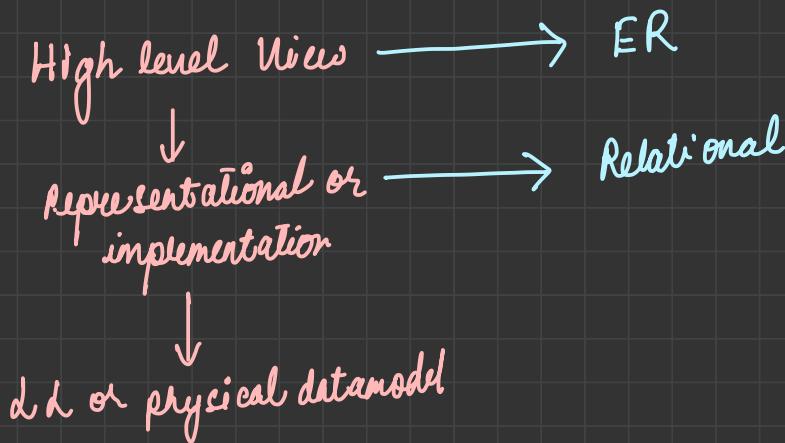
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Data Vs Database

RBR-1



RBR-2



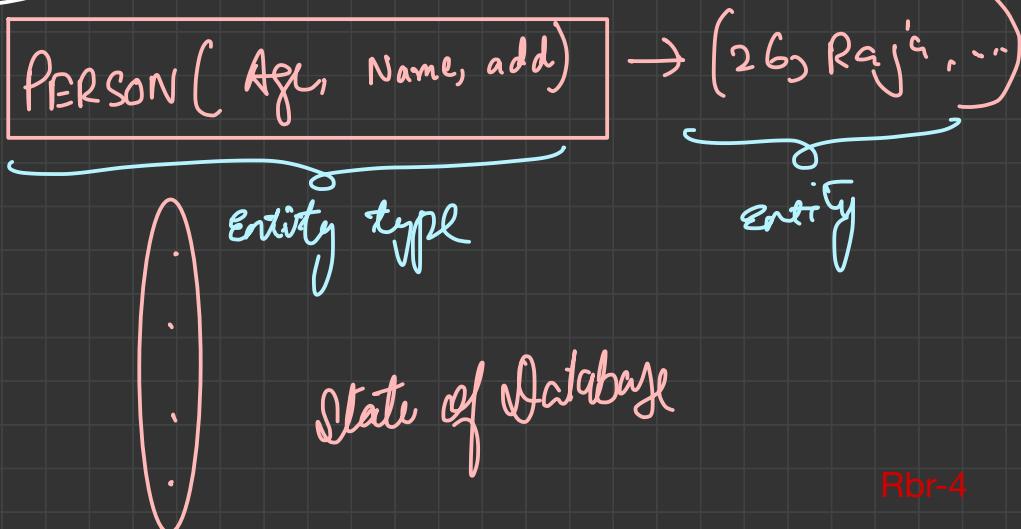
ER Model (Entity Relation Model)

RBR3

Entity → Thing → Person

Attribute → Properties → name, age

Relationship → association → works for



Rbr-4

Attributes.

collection
of simple attr
Composite vs

Name
↓
FN MN LN

Simple attribute



Single valued vs
Age

Multi valued.



Stand vs derived attribute

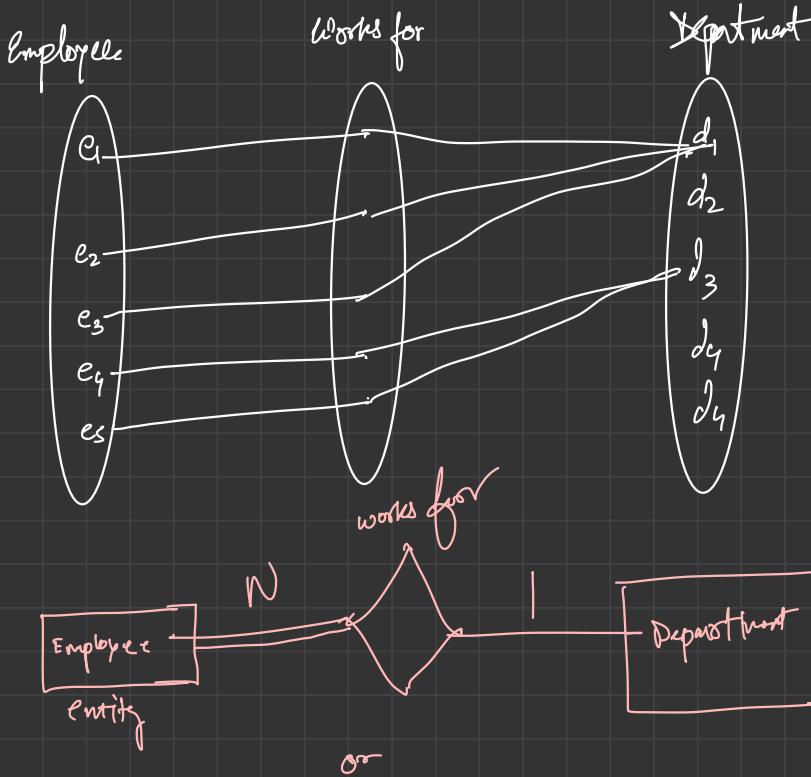
DOB

(by) can be derived
from DOB

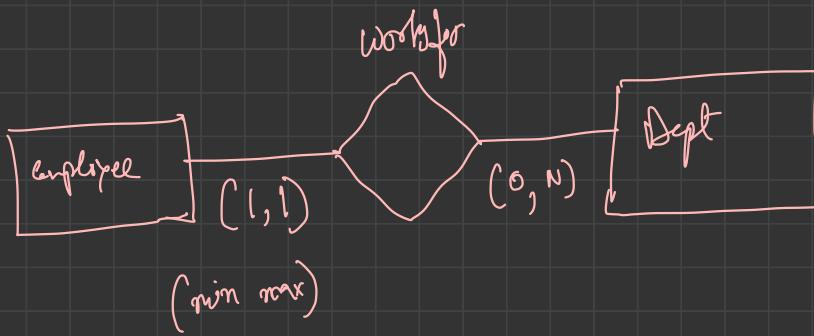
Complex attribute.

composite + multivalued.

Rbr-s

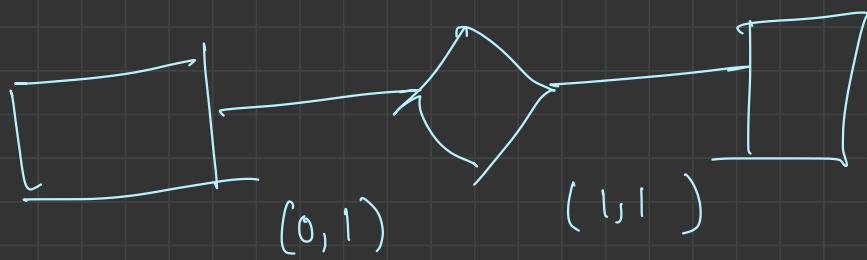
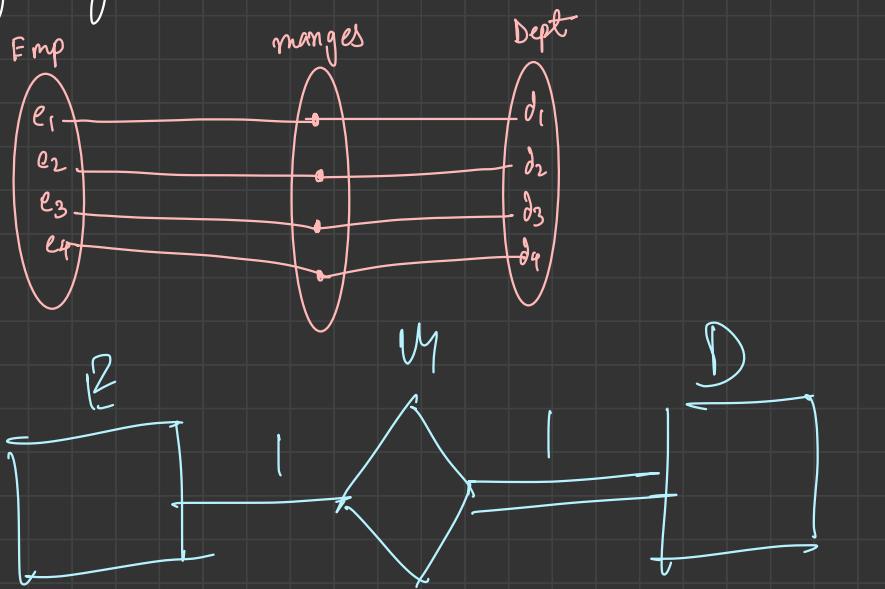


one



Every dept should have a manager and only one employee managed a dept and an employee can manage only one dept.

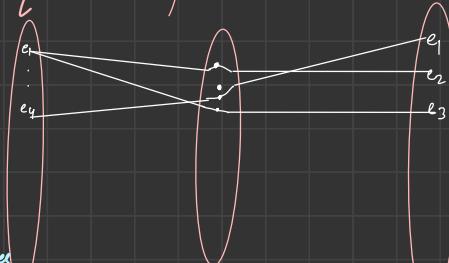
Rbr-6



There is no hard fast rule that every employee should manage a dept.

(Supervise)

Employee Supervises Employee



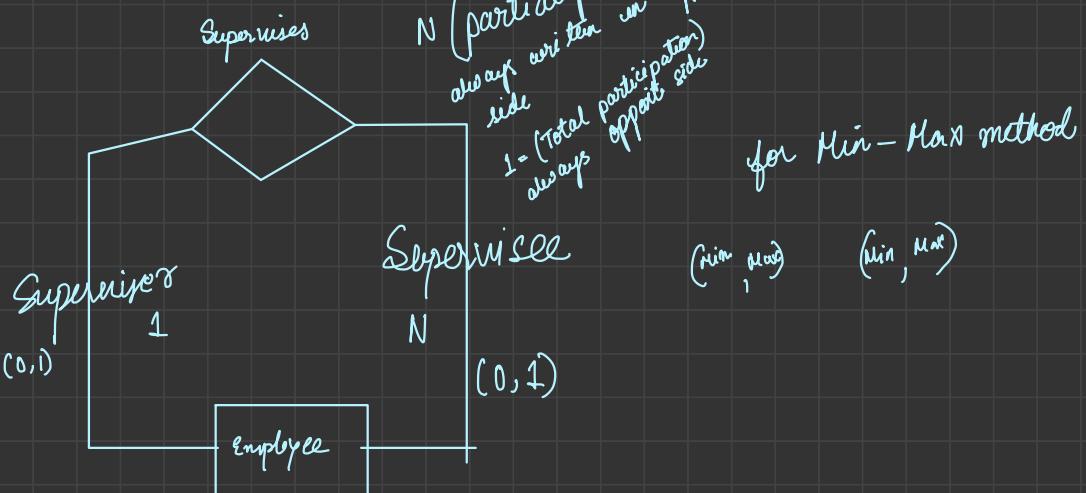
There can exist an employee
who doesn't supervise
anyone

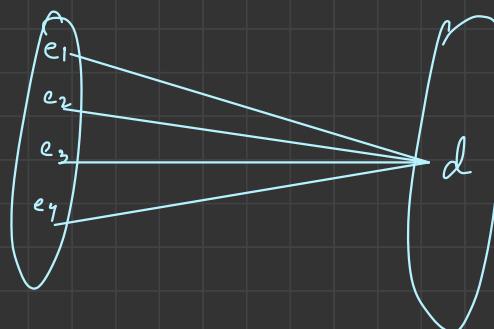
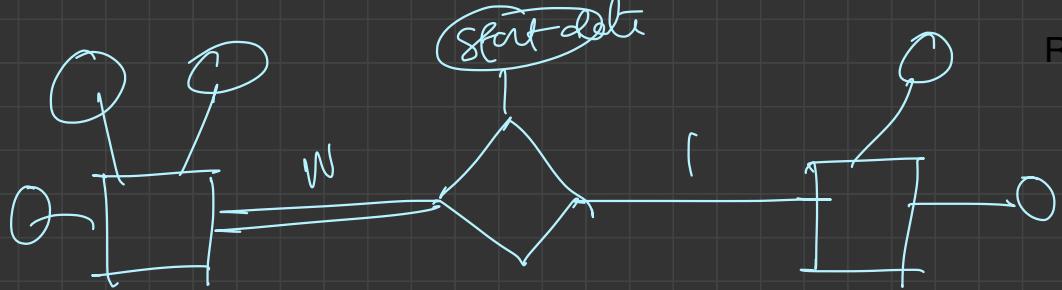
1 employee can be
managed by max 1 boss
so 1. ℓ_2

Degree of Relationship = 2

We can represent it 2 relationship which are
actually same (Recursive Relationship)

Cardinality Max^m participation = N

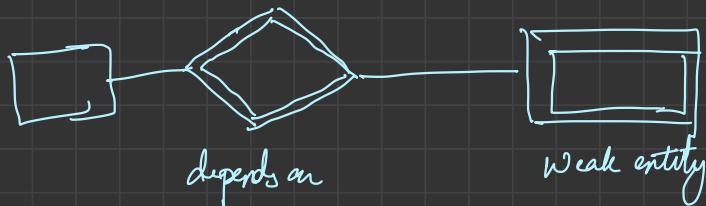




RBR-10

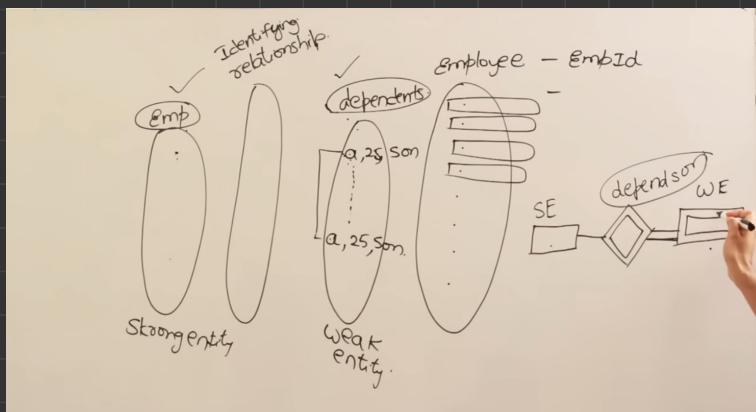
Weak entity \rightarrow No key entity

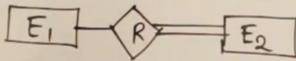
It participation Identifying relation is total



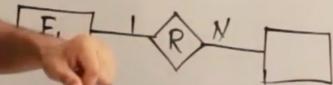
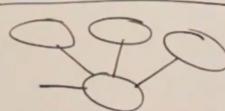
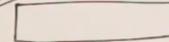
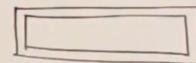
Composite attribute — 2 or more

attribute used for uniquely identify
when a single attribute can't uniquely identify



Derived attributeTotal participation of E₂ in RCardinality ratio

$$E_1 : E_2 = 1:N$$

attributeKey attributemultivalued attributeComposite attributeentityweak entityRelationshipIdentifying Relationship

NS Lectures — Unit 2

Relational model → Tables with rows & columns

Student → table

PK	PK	
Field / attribute:		
(name)	(roll no)	Scalor
Sai	121	A
Shiva		

notes available

Advantages.

- ① Simple (SQL)
- ② Accuracy (well organised & no redundancy)
- ③ Security
- ④ Flexible

DDL — Only Create Structures

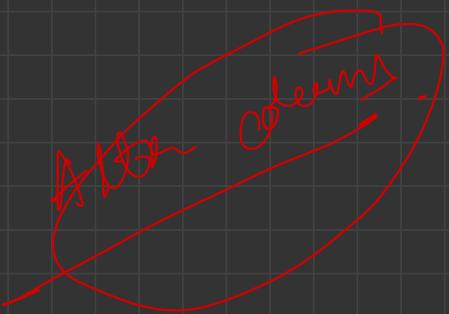
No inserting data

Create database database-name;
Create database college;
Show databases;

Alter table Student add column age int;

describle Student;

Alter ~~modify~~
alter table student



Use college.

Create table Student (

name varchar (50),
→ Text circ.

roll_no int,

Section varchar (10))

Show tables;

[Student]

describle Students;
desc

Drop: too

Drop table Student;

drop database College;

rename table Student to St;

Truncate table.

(Remove all data from
table)

only structure remain).

DML →

(int) no. ↗ ↘ needed

* Select * from student

Select name, Roll no.

Update, Delete

Select command

Select * from student, (output entire table)

age > 15 and age < 21

TCL

Commit / Rollback

DCL .

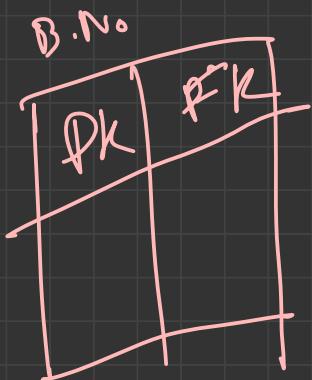
Foreign Key, Null, Duplicate
accepted

only condition should contain value only in Primary key.

View → Virtual table .

Name	Rollno	Section
Sai	1	A
Bai	2	B
Ravi	3	A

Foreign key & Primary key Connection



Relational Algebra

Unary relation Operations

- Select (σ)
- Project (π)
- Rename (δ)

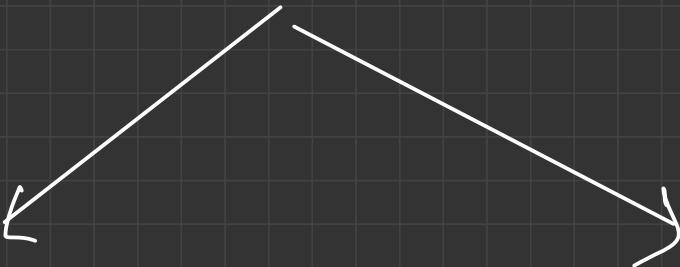
Cartesian Product

- (\times)

Set Operations

- Union (\cup)
- Intersection (\cap)
- Set difference ($-$)

Joints



Inner Joints

- Natural (*)
- Theta (θ)
- Equi (=)

Outer joints

- left outer (\cancel{X})
- Right (\cancel{X})
- Full (\cancel{X})

Selection operator duplicate possible

Projection operator duplicate not possible