

Lecture 26

EE 421 / CS 425

Digital System Design

Fall 2025

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Topics

- BRM Design and Applications
- BRM Chips
- UART Design
- UART Configuration
- UART Transmitter Operation
- UART Receiver Operation
- UART Architecture
- UART Registers
- Application of Serial Comm – Bluetooth
- Application of Serial Comm – Zigbee
- Possibly Ethernet Hardware or similar



Binary Rate Multipliers - BRM

Definition: A circuit module that transforms a stream of input clock pulses into another stream of output clock pulses

Let N_i = no. of input pulses for a particular time period

Let N_o = no. of output clock pulses for the same time period

$$\textbf{Binary Fractional Rate Multiplier} \Rightarrow N_o = \frac{B}{2^n} N_i$$

B = Rate Constant Input to module

B = $(B_{n-1}, B_{n-2}, \dots, B_2, B_1, B_0)_2$

n = no. of binary counter stages controlling the module

The clock input drives an n-bit binary counter whose outputs are labelled $(X_n, X_{n-1}, \dots, X_2, X_1)$.

Will do in next lecture due to quiz today – from previous lecture

Digital Fractional Rate Multiplier - Principle

The counter outputs (X_n to X_1 = MSB to LSB) are ANDed with the incoming Clock signals to form intermediate pulse trains P_i ($i=1,n$):

$$P_n = X_1 \cdot \text{Clock}$$

$$P_{n-1} = \overline{X_1} \cdot X_2 \cdot \text{Clock}$$

$$P_{n-2} = \overline{X_1} \cdot \overline{X_2} \cdot X_3 \cdot \text{Clock}$$

.....

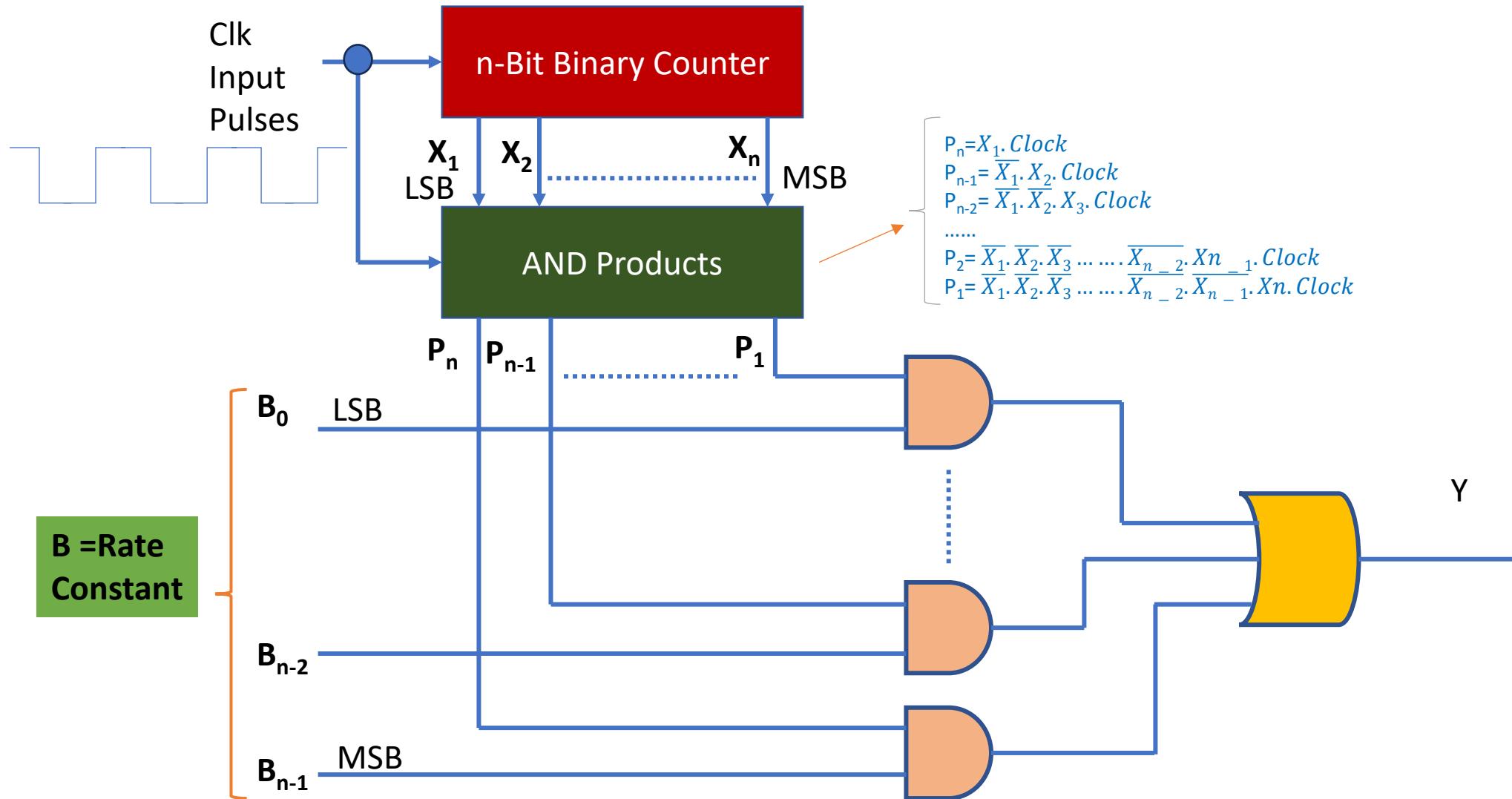
$$P_2 = \overline{X_1} \cdot \overline{X_2} \cdot \overline{X_3} \dots \dots \overline{X_{n-2}} \cdot X_{n-1} \cdot \text{Clock}$$

$$P_1 = \overline{X_1} \cdot \overline{X_2} \cdot \overline{X_3} \dots \dots \overline{X_{n-2}} \cdot \overline{X_{n-1}} \cdot X_n \cdot \text{Clock}$$

The logic output Y uses the rate constant M and the pulse train signals P_i to implement the equation $\textcolor{red}{Y} = \sum B_{i-1} \cdot P_i$

This output configuration will deliver the proper number of output pulses (N_o) as specified by rate constant B

Basic Circuit Diagram – DFRM



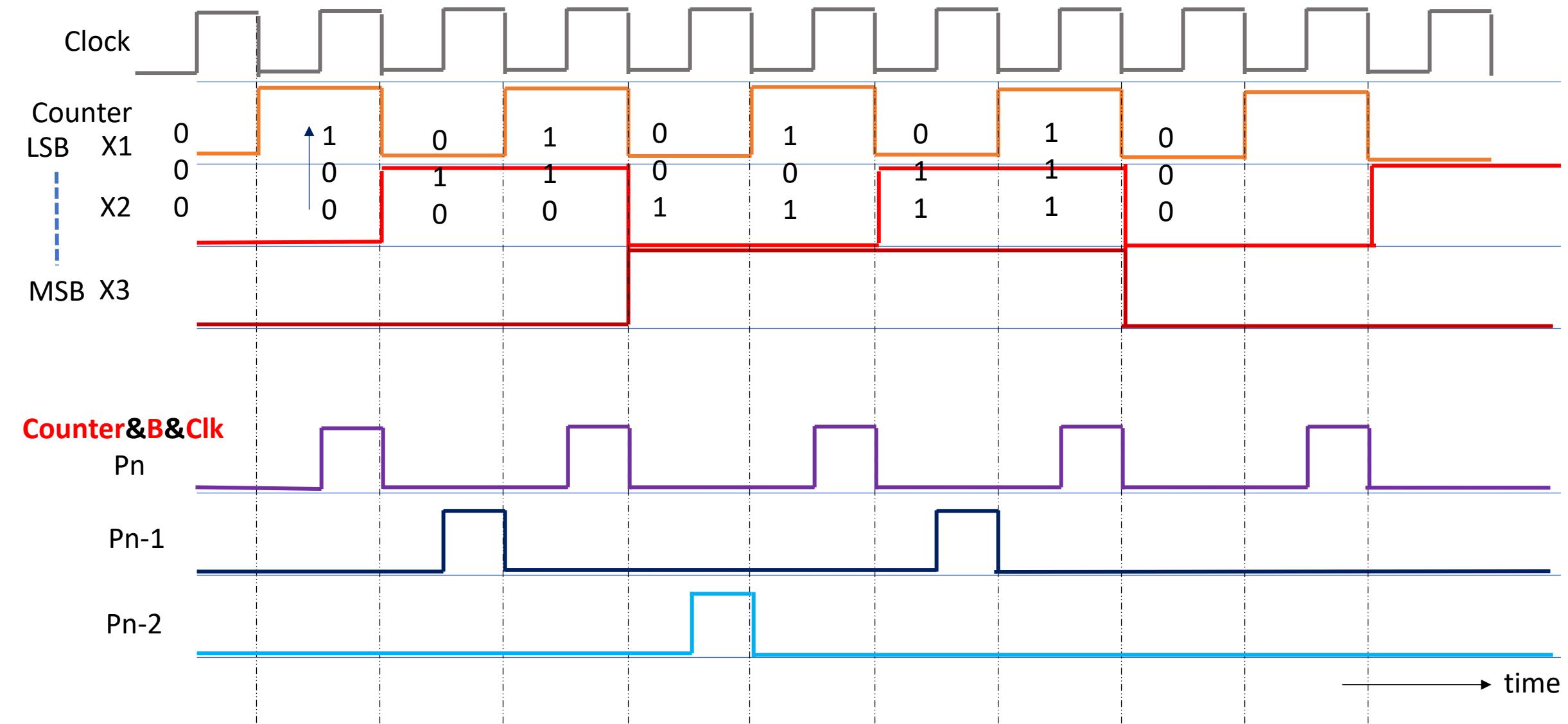
Example and timing diagram of BRM

- BRM using a **3-Bit counter**
- Rate Constant $B = 7_{10}$
- Counter generates 8 pulses per time period; Due to $B=7$, 7 pulses appear at the output and one pulse is eliminated

Analysis of timing diagram 3 bit BRM, B=7

- $Y = B_2 \cdot P_3 + B_1 \cdot P_2 + B_0 \cdot P_1$
- Where:
- $P_3 = X_1 \cdot \text{Clock}$
- $P_2 = \overline{X_1} \cdot X_2 \cdot \text{Clock}$
- $P_1 = \overline{X_1} \cdot \overline{X_2} \cdot X_3 \cdot \text{Clock}$
- And $B = (B_2 \cdot B_1 \cdot B_0) = (1.1.1)_2$
- Thus $Y = (X_1 + \overline{X_1}X_2 + \overline{X_1} \cdot \overline{X_2} \cdot X_3) \cdot \text{Clock}$
- In terms of counter output $X = (X_3 \cdot X_2 \cdot X_1)$
- Product term $X_1 \Rightarrow \sum m(1,3,5,7)$
- Product term $X_2 \cdot \overline{X_1} \Rightarrow \sum m(2,6)$
- Product term $X_3 \cdot \overline{X_2} \cdot \overline{X_1} \Rightarrow \sum m(4)$
- So that $Y(X_3 \cdot X_2 \cdot X_1) = (\sum m(1,2,3,4,5,6,7)) \cdot \text{Clock}$
- As m_0 is missing from this list so first clock from each sequence of 8 pulses will be eliminated

Timing Diagram of DFRM



How many pulses are produced?

- From the timing diagram, each of the pulse trains P_i generates 2^{i-1} pulses during one counter sequence period (2^n clock pulses)
- Thus P_1 generates 1 pulse, P_2 generates 2 pulses, P_3 generates 4 pulses, and so on
- Note that pulses do not overlap in time
- We can **OR** the respective pulses to produce desired output stream
- The output pulses can be irregularly spaced
- The output pulses are synchronized with input clock pulses

Serially Cascaded Binary Rate Multiplier BRM

Design a BRM to implement $N_{out} = \frac{63}{320} \cdot N_{in}$

Break into binary powers ($63 \times 2 = 126$; and $320 \times 2 = 640$) :

$$N_{out} = \frac{(7 \times 18)}{(10 \times 64)} \cdot N_{in}$$

$$N_{out} = \left(\frac{7}{10}\right) \cdot \left(\frac{18}{64}\right) \cdot N_{in}$$

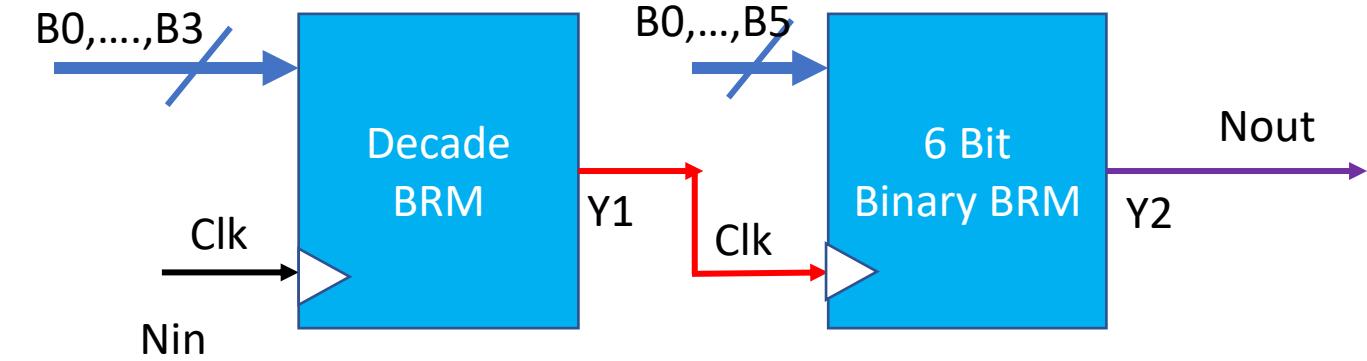
Use:

1 x Decade Counter BRM

1 x 6-Bit Binary Counter BRM

Set B for Decade Counter = 0111

Set B for 6-Bit BRM = 10010



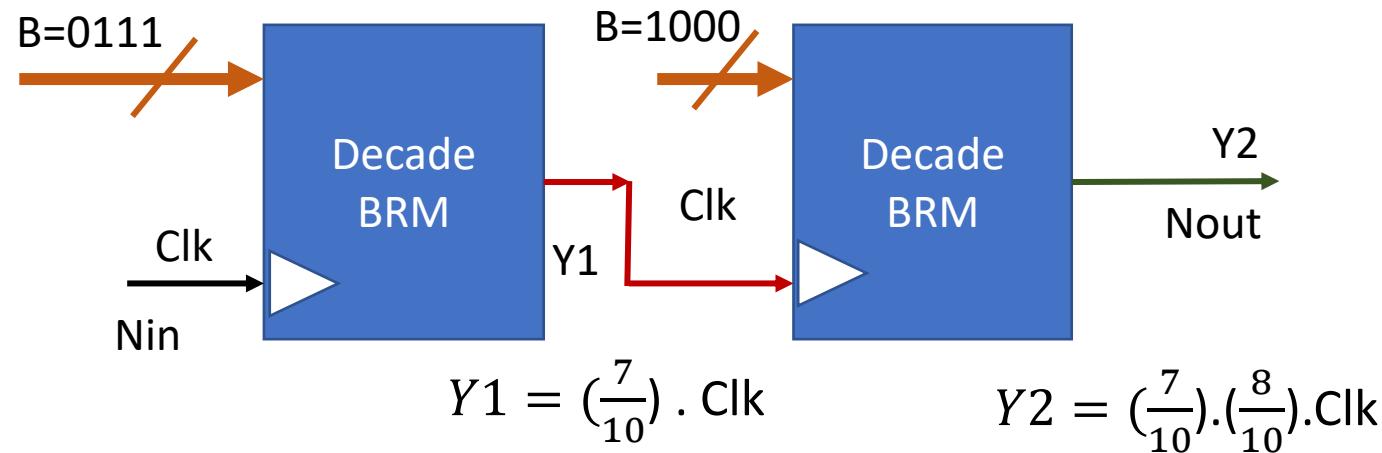
Example 2: Serially Cascaded BRM

Design a BRM circuit that produces: **0.56 of input clock rate**

We use two Decade BRMs with
Value of B is set as 7 and 8 respectively

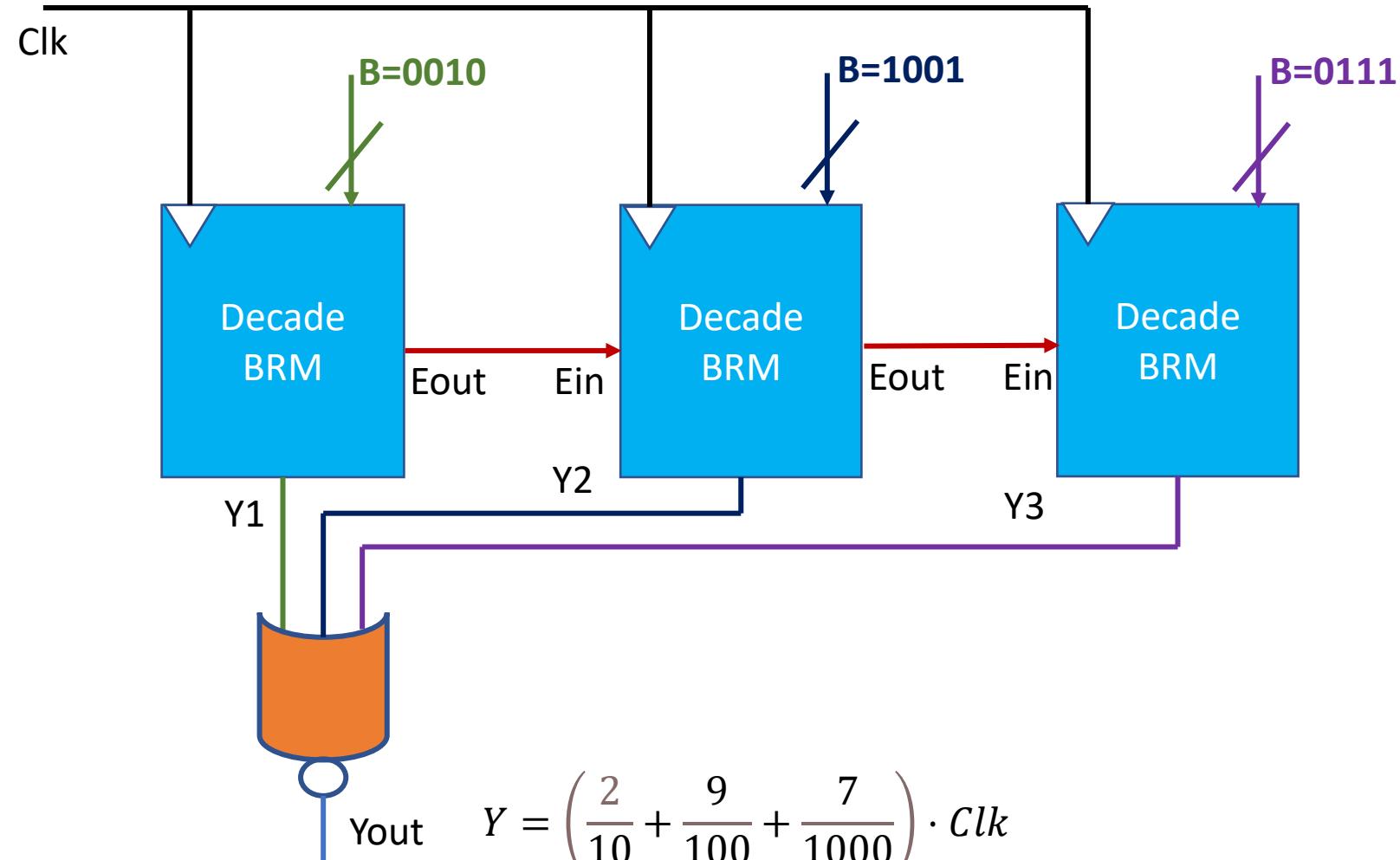
$$N_{out} = \frac{56}{100} \cdot N_{in}$$

$$N = \frac{7}{10} \cdot \frac{8}{10} \cdot N_{in}$$



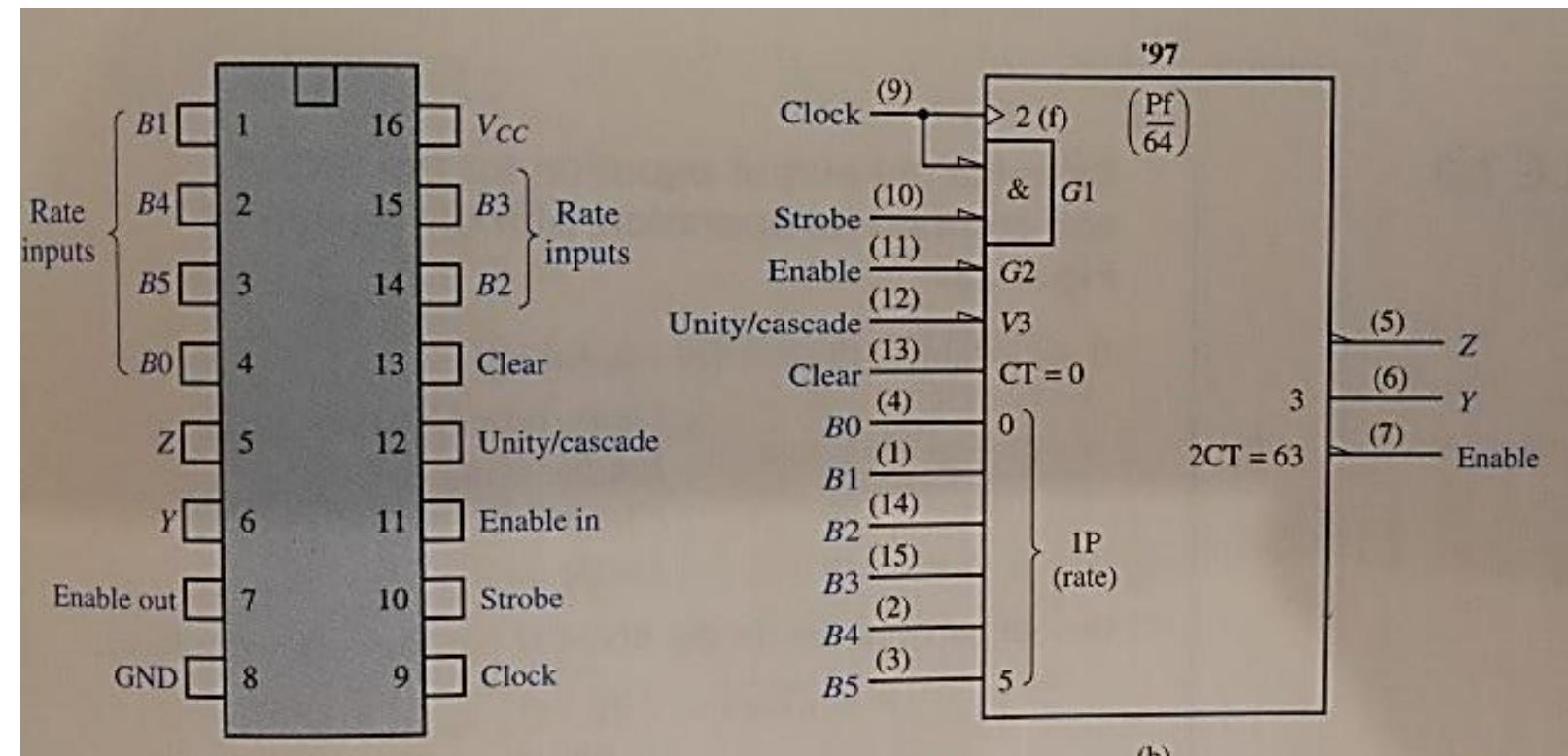
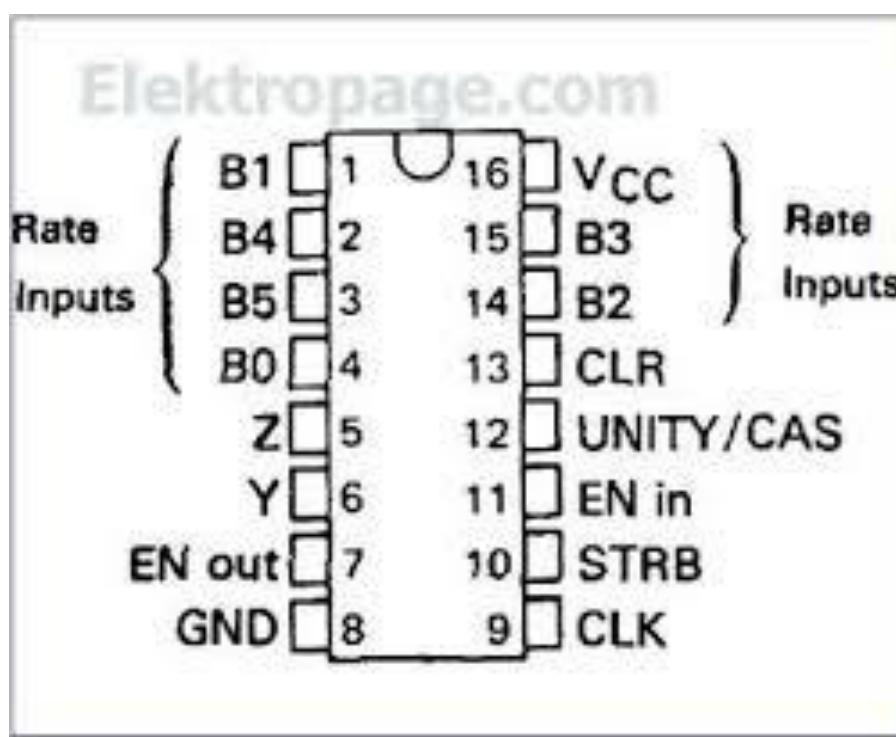
Eg 3: Parallel Cascaded Binary Rate Multiplier

Design a BRM circuit to produce $N_{\text{out}} = (0.297)N_{\text{in}}$



$$Y = \left(\frac{2}{10} + \frac{9}{100} + \frac{7}{1000} \right) \cdot Clk$$

7497 TTL chip 4-Bit and SN7497E chip 6-Bit BRM

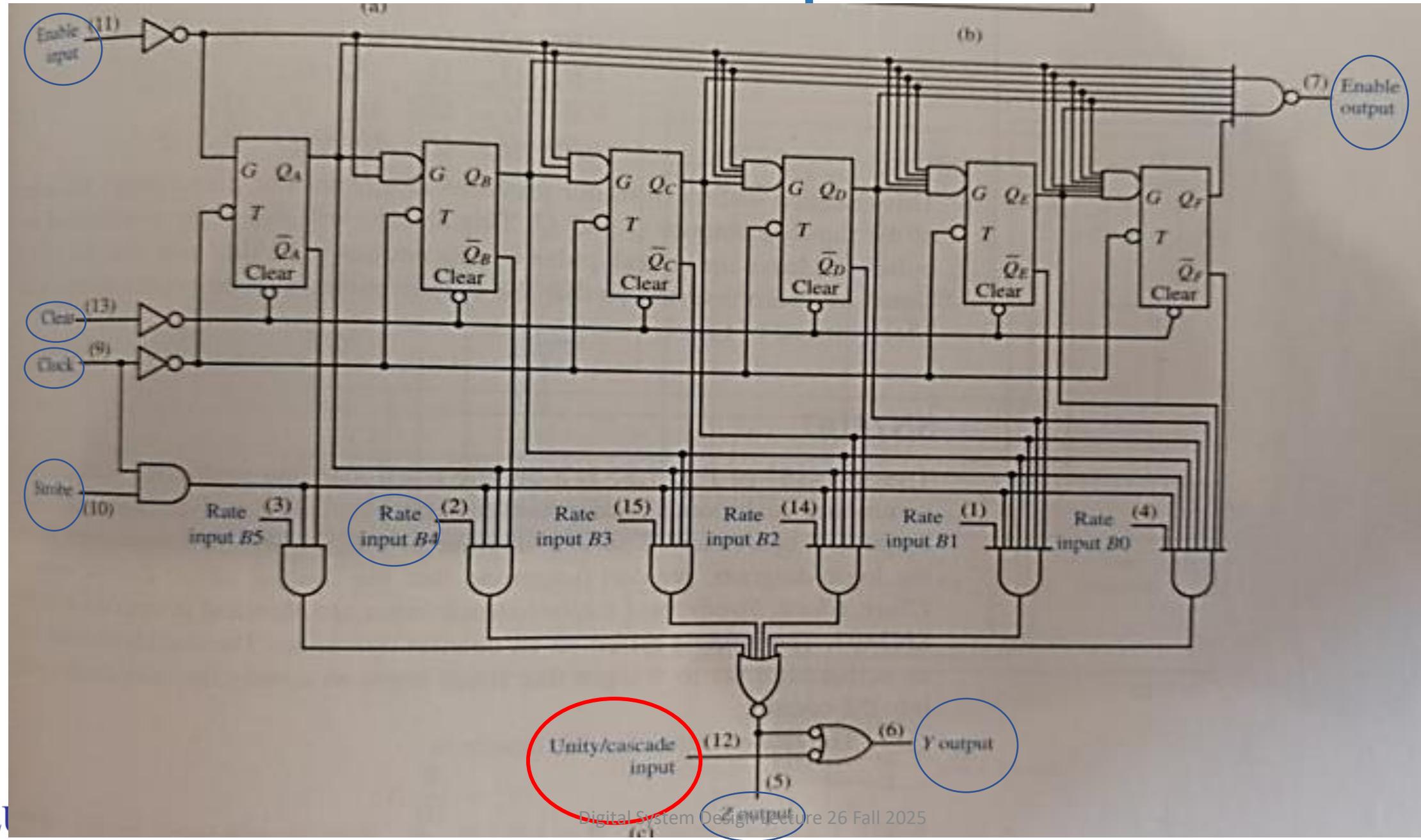


Ref: Images from internet

Pins on 7497 and SN7497E

- Clear signal is active high and asynchronous to drive all flip-flops to all-zero state
- Rising edge on Clock pulses causes counter to advance
- Strobe is an active-low Enable signal for the pulse rate output Y and Z
- Enable input and Enable output allows the modules to be cascaded
- The Unity / Cascade input may be used to combine the outputs of cascaded units, effectively ORing the pulses trains from two modules

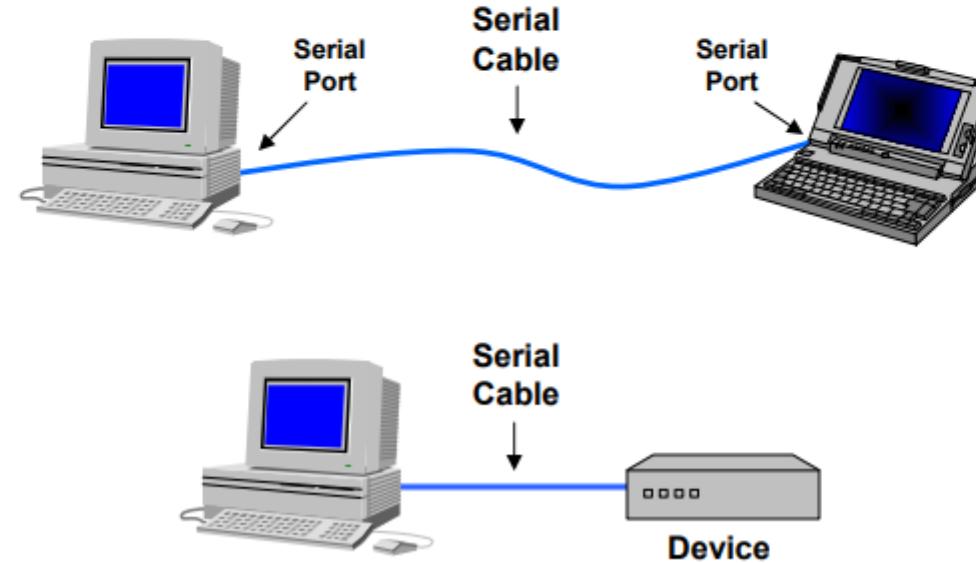
Internal view of BRM chip



UART RS232 Protocol and Hardware Implementation

UART for Serial Communication

- PC serial port is a UART!
- Serializes data to be sent over serial cable
 - De-serializes received data



UART Functions

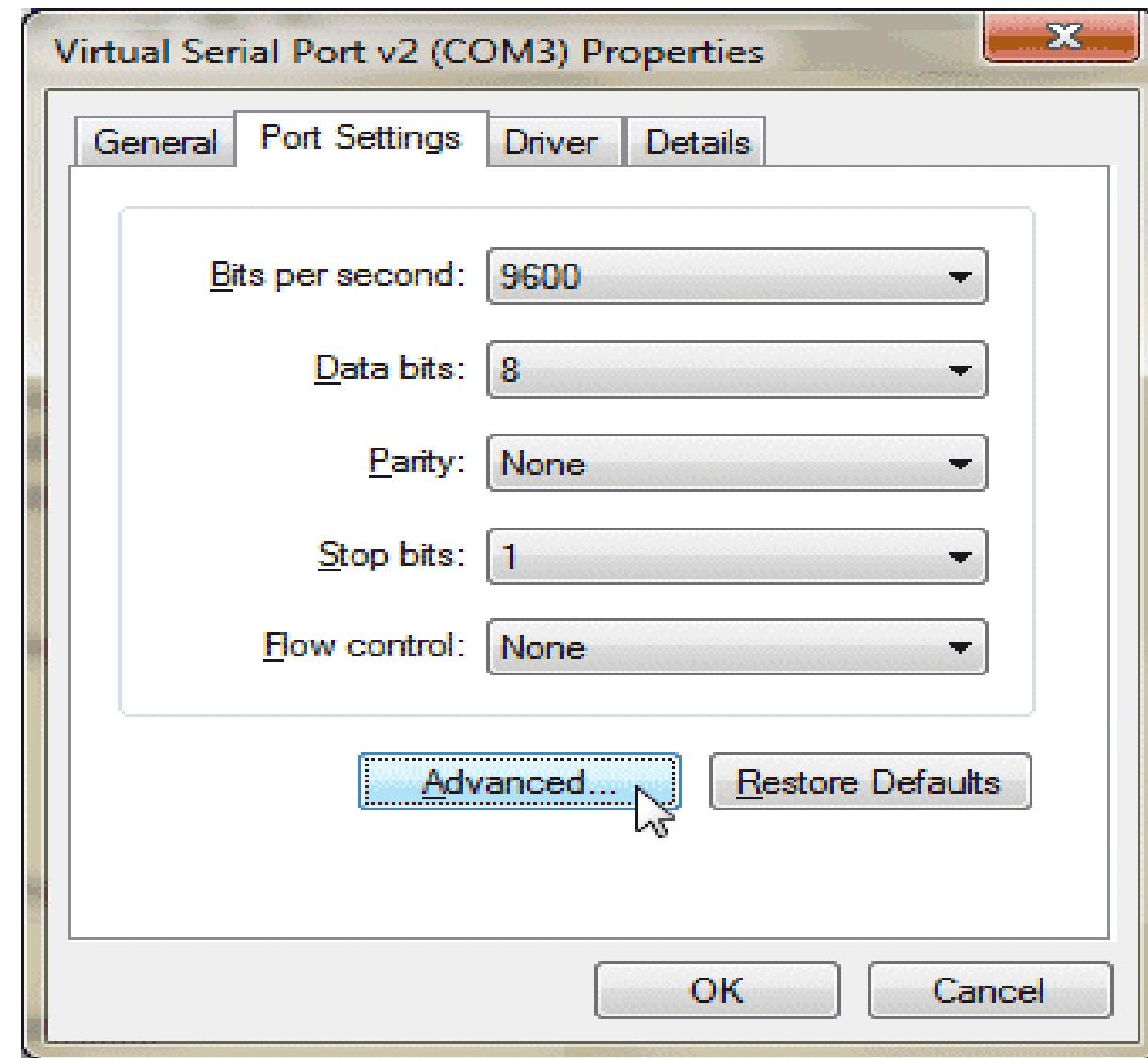
- Outbound data
 - Convert from parallel to serial
 - Add start and stop delineators (bits)
 - Add parity bit
- Inbound data
 - Convert from serial to parallel
 - Remove start and stop delineators (bits)
 - Check and remove parity bit

UART Principle

- UART translates Parallel Data (Bytes or ASCII) to Serial Data for sending and vice versa for receiving.
- Inside UART there are several registers to store status and data for sending and receiving.
- The CPU can read and write these registers through I/O ports.
- For correct communication, both UARTs must make agreement on data frame format and transmission speed (baud rate).

Serial Port Setting on PC using HyperTerminal

The transmitter and receiver have same configuration



Configurable Baud Rate

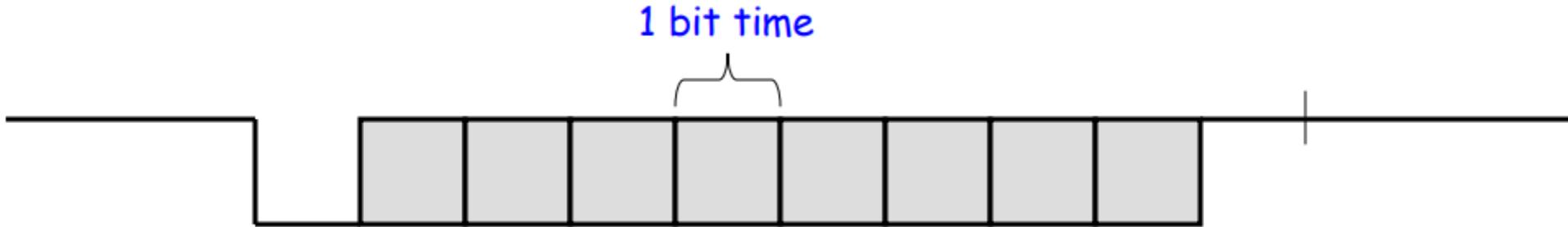
- Definition: “How many times a signal can change per second” Includes data bits plus control bits.
- Possible Baud Rates:
 - 110
 - 150
 - 300
 - 600
 - 1200
 - 2400
 - 4800
 - 9600
 - 19200
 - 38400
 - 57600, 115200, 230400, etc.

UART Throughput

- Data Throughput Example
 - Assume 19200 baud, 8 data bits, no parity, 1 stop bit
 - 19200 baud \rightarrow 19.2 kbps
 - 1 start bit + 8 data bits + 1 stop bit \rightarrow 10 bits
 - It takes 10 bits to send 8 bits (1 byte) of data
 - $19.2 \text{ kbps} \cdot 8/10 = \mathbf{15.36 \text{ kbps}}$
- How many KB (kilobytes) per second is this?
 - 1 byte = 8 bits
 - 1 KB = 1,024 bytes
 - So, $1 \text{ KB} = 1,024 \text{ bytes} \cdot 8 \text{ bits/byte} = 8,192 \text{ bits}$
 - Finally, $15,360 \text{ bps} \cdot 1 \text{ KB} / 8,192 \text{ bits} = \mathbf{1.875 \text{ KB/s}}$

UART Character Transmission

- Each bit has a fixed time duration determined by the transmission rate
- Example: a 1200 bps (bits per second) UART will have a $1/1200$ s or about $833.3 \mu\text{s}$ bit width



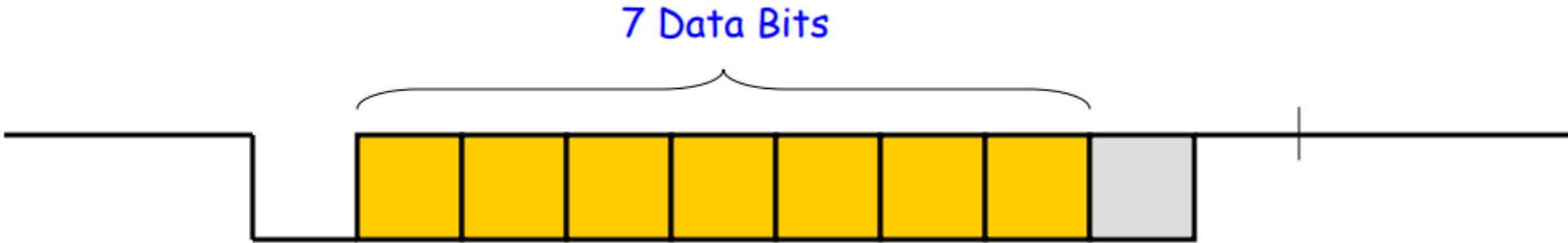
UART Start Bit in Transmission

- The **start bit** marks the beginning of a new word
- When detected, the receiver synchronizes with the new data stream



UART Data Bits in Transmission

- Next follows the **data bits** (7 or 8)
- The least significant bit is sent first



UART Parity Bit in Transmission

- The **parity bit** is added to make the number of 1's even (even parity) or odd (odd parity)
- This bit can be used by the receiver to check for transmission errors



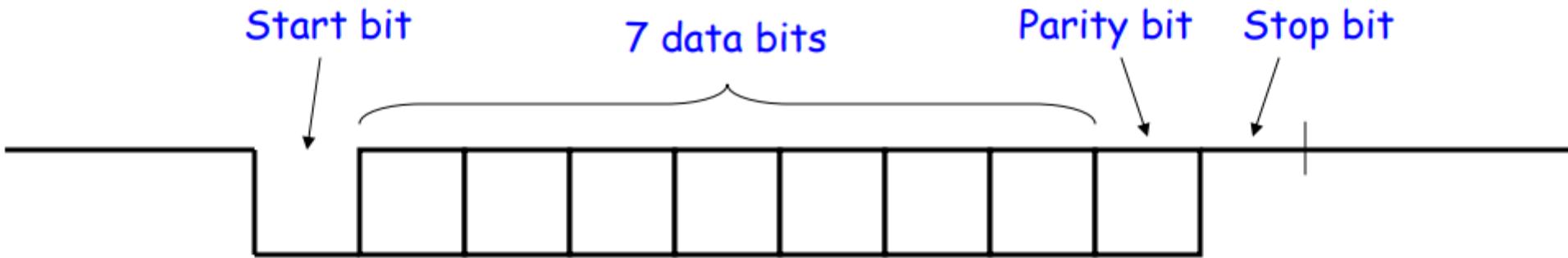
UART Stop Bit in Transmission

- The **stop bit** marks the end of transmission
- Receiver checks to make sure it is '1'
- Separates one word from the start bit of the next word



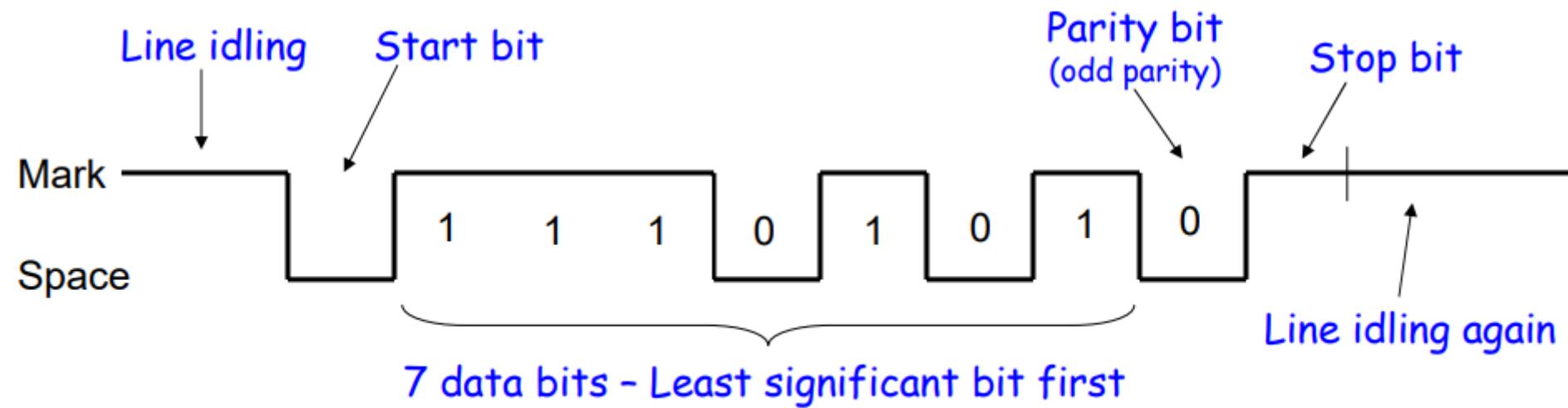
UART Transmission of 7 Bits

- In the configuration shown, it takes 10 bits to send 7 bits of data

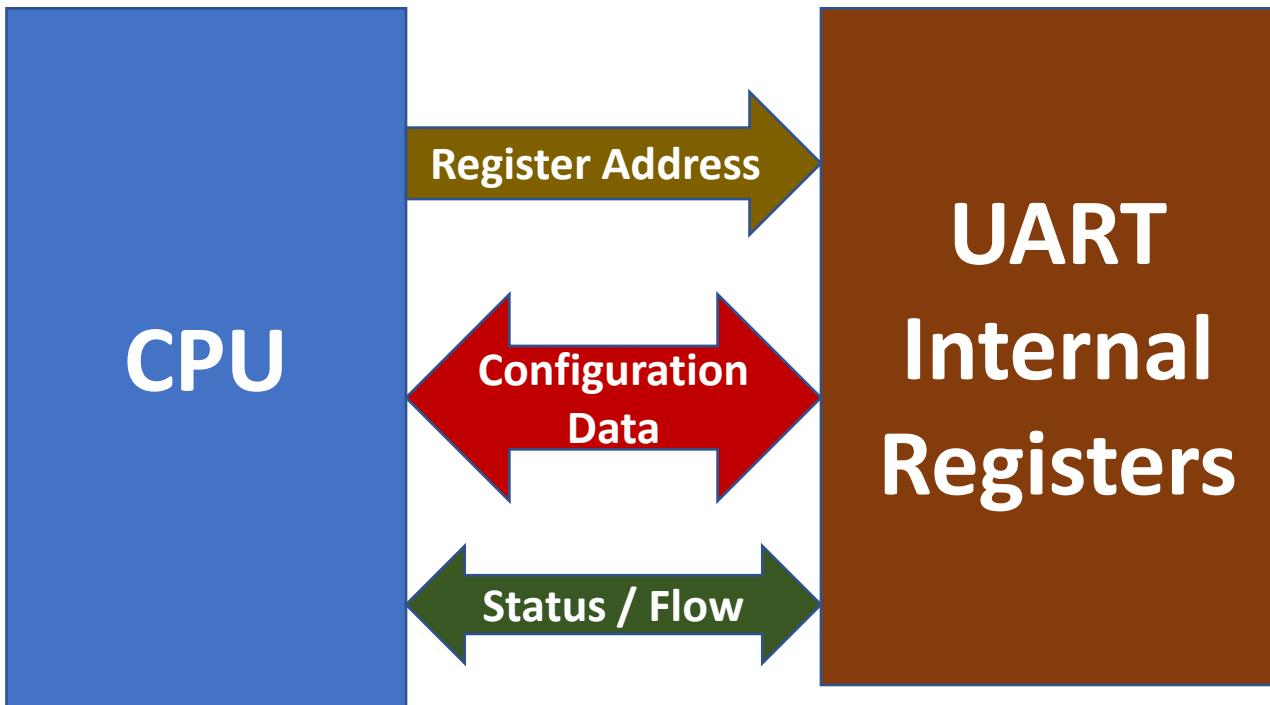


UART Transmit ASCII character 'W'

- Send the ASCII letter 'W' (1010111)

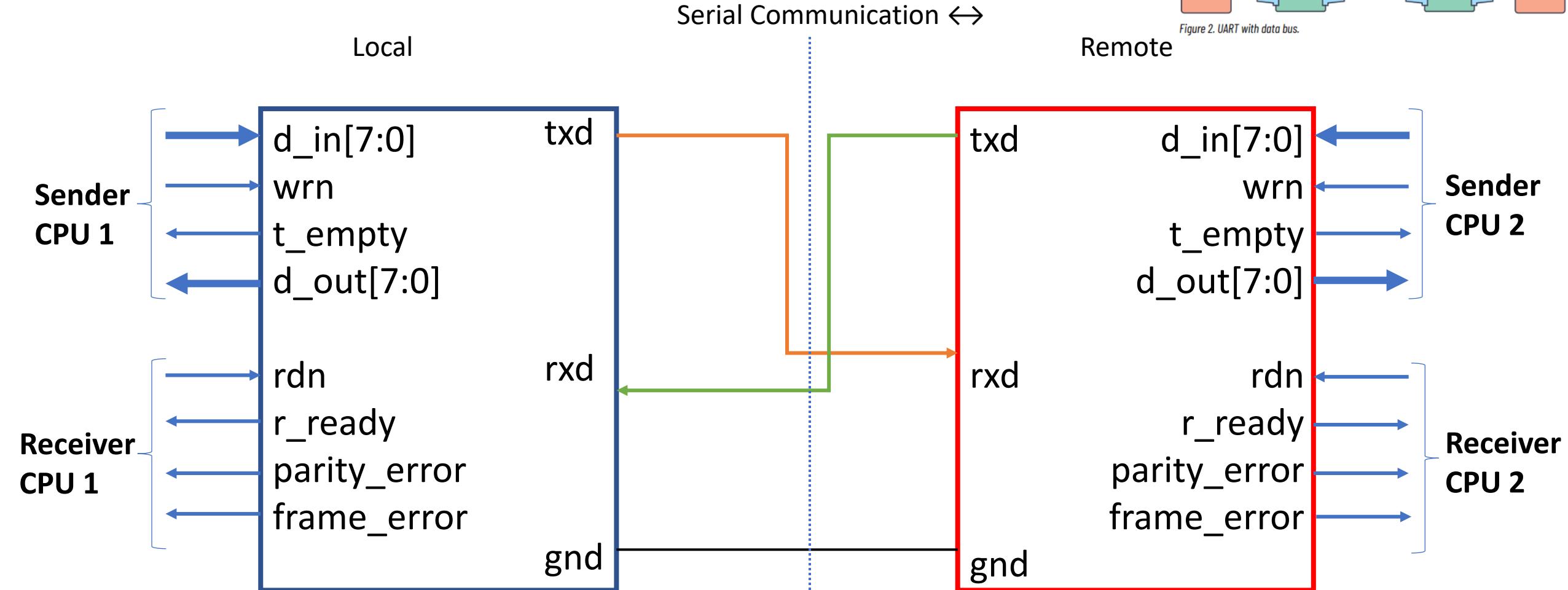
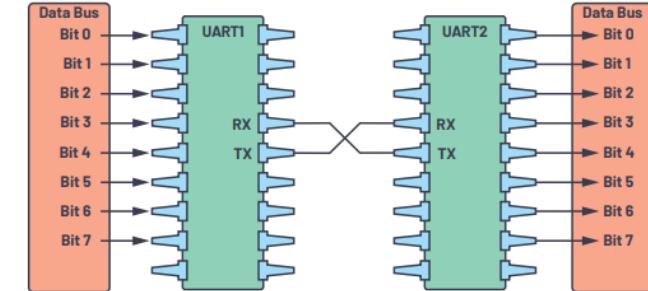


CPU and UART Configuration

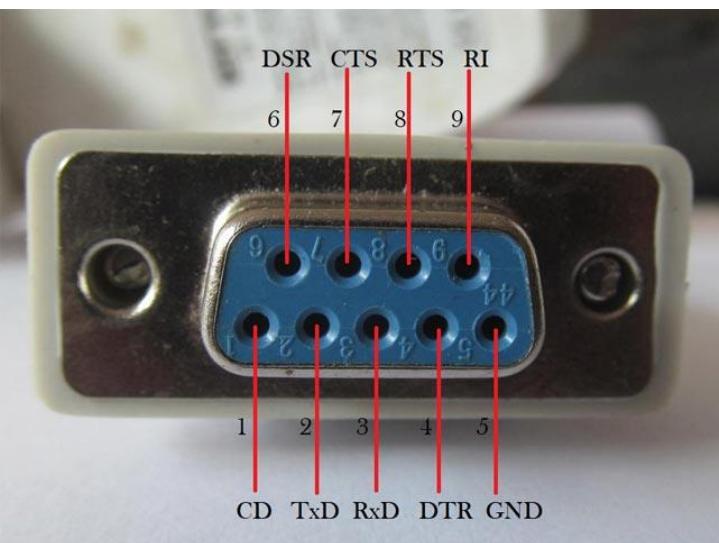
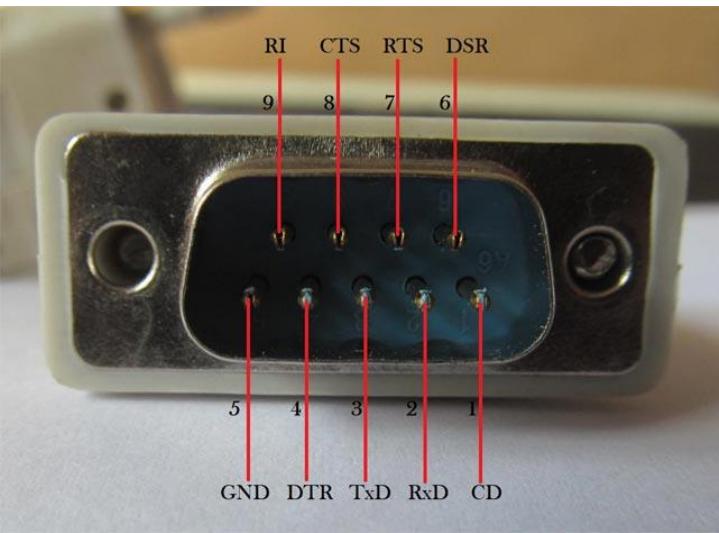


Configuration Data
Baud Rate
No. of Stop Bits (1, 2)
Parity (Yes, No, Even, Odd)
Flow Control (RTS, CTS)
Data Bits (5 to 8)

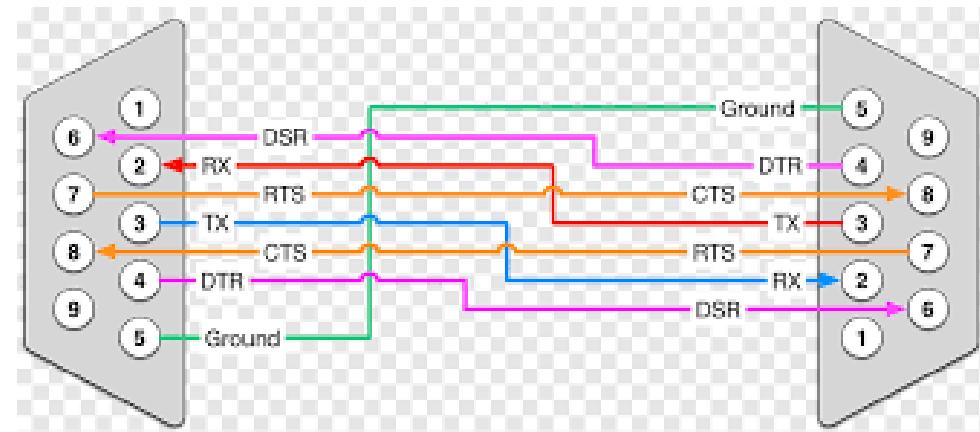
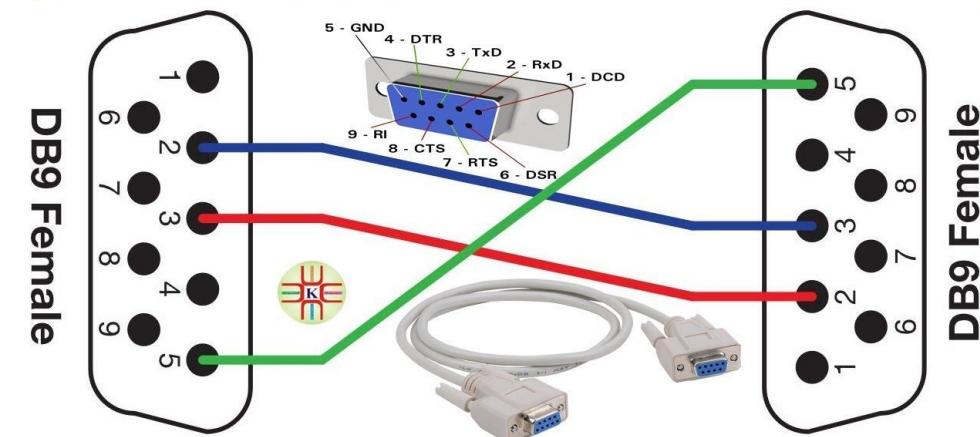
Connecting two UARTs



Physical Connections through RS 232



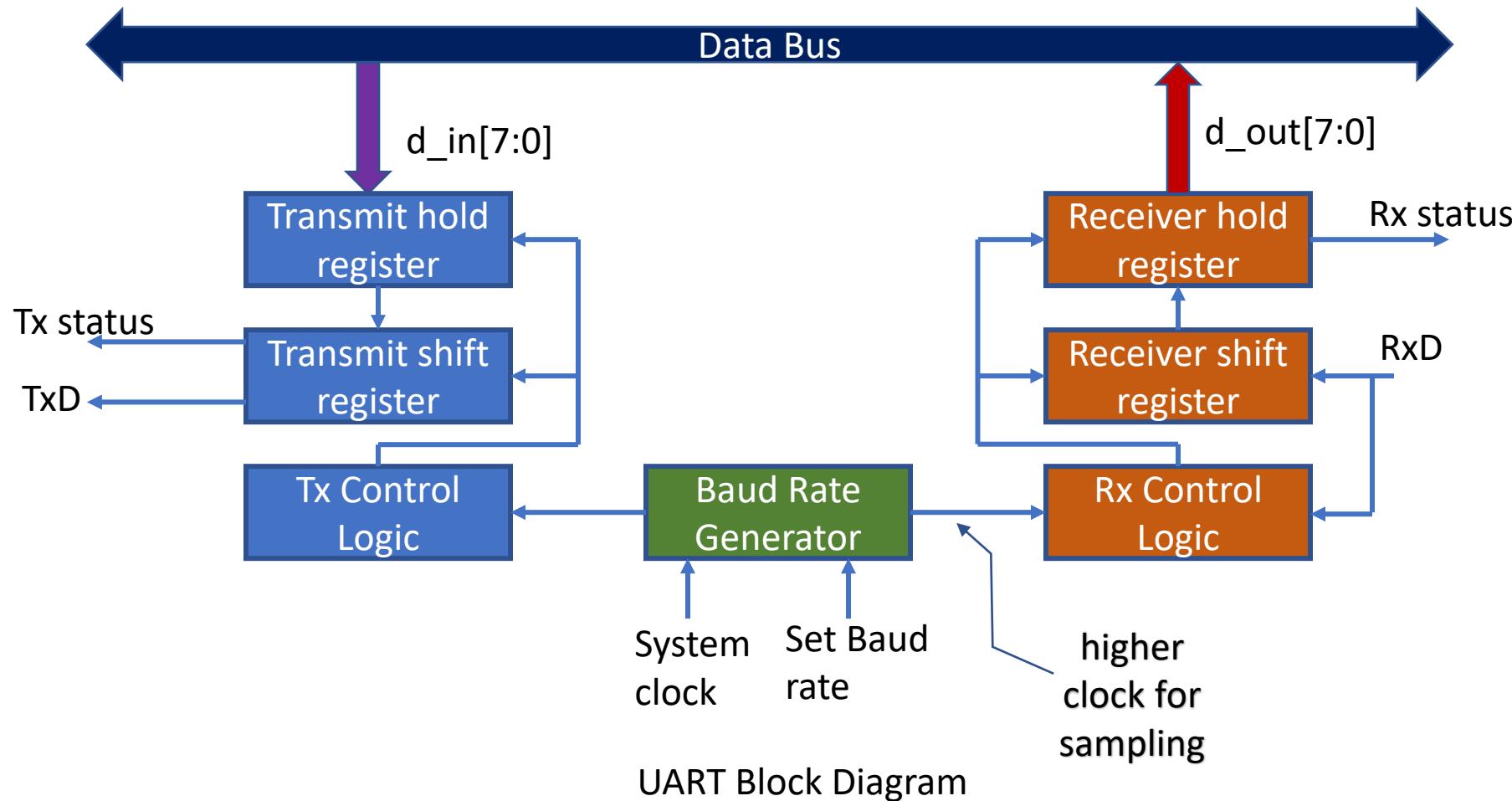
How to Make RS232 Serial Female Cable At Home



Typical UART Components

- **Baudrate clock generator:** Multiple of the bit rate to improve sampling in the middle of a bit period.
 - For generating this timing information, each UART uses an oscillator generating a frequency of about 1.8432 MHz (as an example).
 - This frequency is divided by 16 to generate the time base for communication.
 - Hence the maximum allowed communication speed is 115200 bps.
- **Input and output shift registers:** Each UART contains a [shift register](#) which is the fundamental method of conversion between serial and parallel forms. These registers shifts the data that has to be serially transmitted or serially received.
- **Transmit and receive control:** This control logic checks for the control signals from host processor to start or stop the transmission and reception of the data bits. In case of any error it also generates error signals.
- **Optional transmit and receive buffers:** Buffers can be used to hold the data temporarily.
- **Optional parallel data bus buffer:** This buffer improves the speed.
- **Optional FIFO:** The UART works by writing data from the host processor to its FIFO buffers and feeding the data from the buffer to the serial device in the format dictated by the user.

UART Architecture Block Diagram

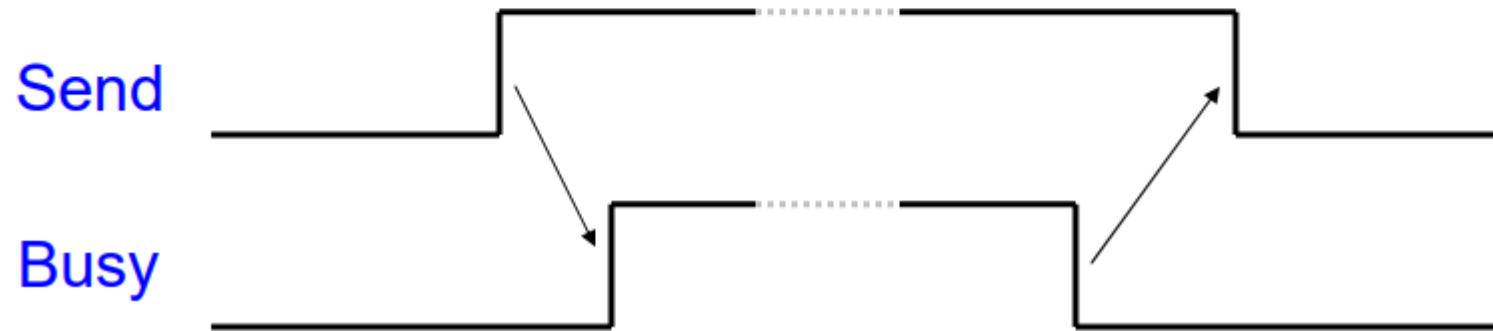


Process of Sending

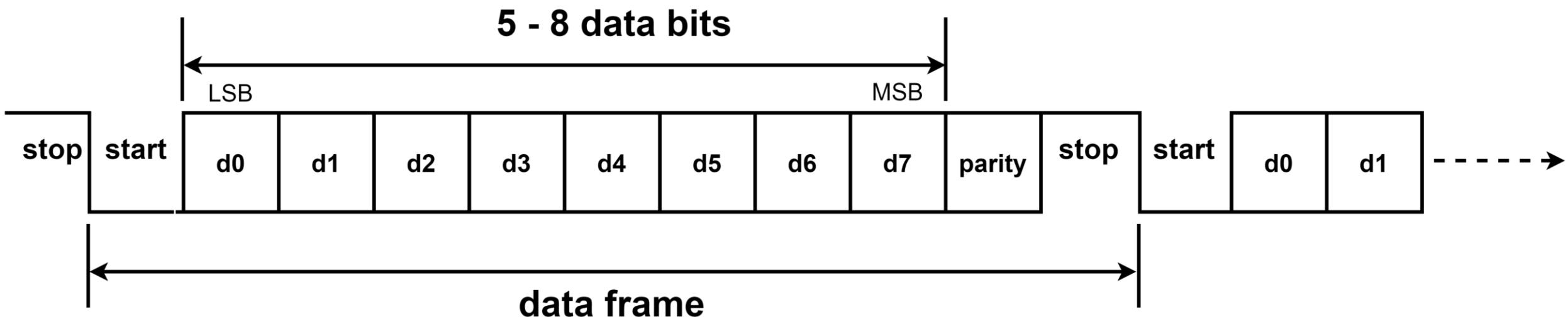
- Asynchronous transmission does not send clock signal to transmit data to the receiver. The sender and receiver must agree on timing parameters in advance and special bits such as start and stop bits are added to each word which is used to synchronize the sending and receiving units.
- A bit called the "**Start Bit**" is added to the beginning of each word that is to be transmitted. The **Start Bit** indicates the start of the data transmission, and it alerts the receiver that a word of data is about to be sent.
- Upon reception of start bit the clock in the receiver goes into synchronization with the clock in the transmitter. The accuracy of these two clocks should not deviate more than 10% during the transmission of the remaining bits in the word.
- The individual bits of the word of data are sent after the start bit. Least Significant Bit (LSB) is sent first.
- The transmitter does not know when the receiver has read at the value of the bit. The transmitter begins transmitting the next bit of the word on next clock edge.
- Parity bit is be added when the entire data word has been sent. This bit can be used to detect errors at the receiver side. Then one Stop Bit is sent by the transmitter to indicate the end of the valid data bits.

UART Transmitter / System Handshaking

- System asserts Send and holds it high when it wants to send a byte
- UART asserts Busy signal in response
- When UART has finished transfer, UART de-asserts Busy signal
- System de-asserts Send signal



Data frame format of UART



LSB of data bits is transmitted first

Start bit is logic low

Stop bit is logic high

In RESET Condition, Serial Data Transmission is ‘High’

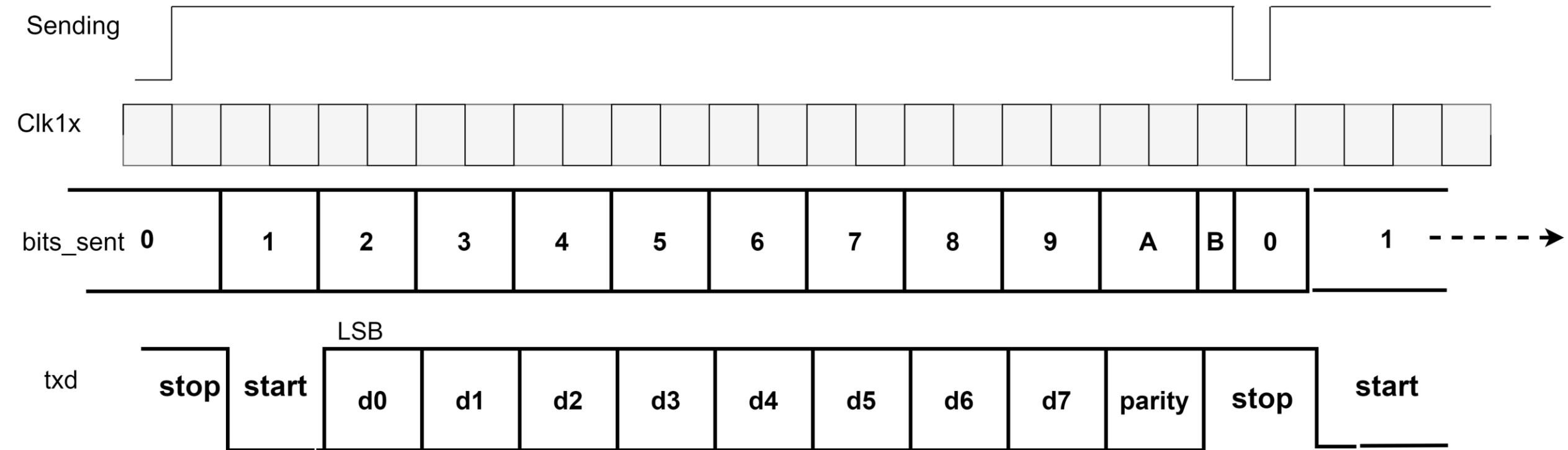
Data Frame

The data frame contains the actual data being transferred. It can be five (5) bits up to eight (8) bits long if a parity bit is used. If no parity bit is used, the data frame can be nine (9) bits long. In most cases, the data is sent with the least significant bit first.

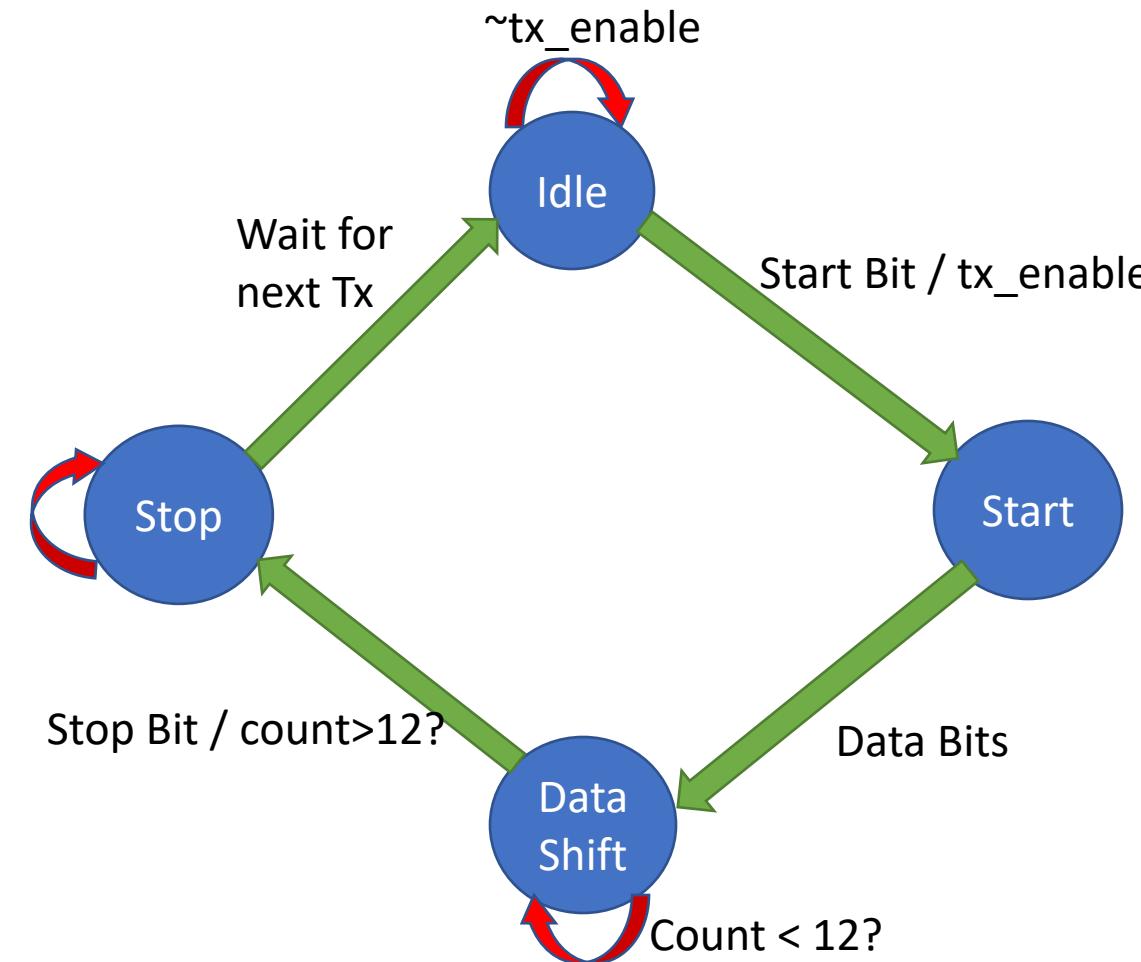


Figure 5. Data frame.

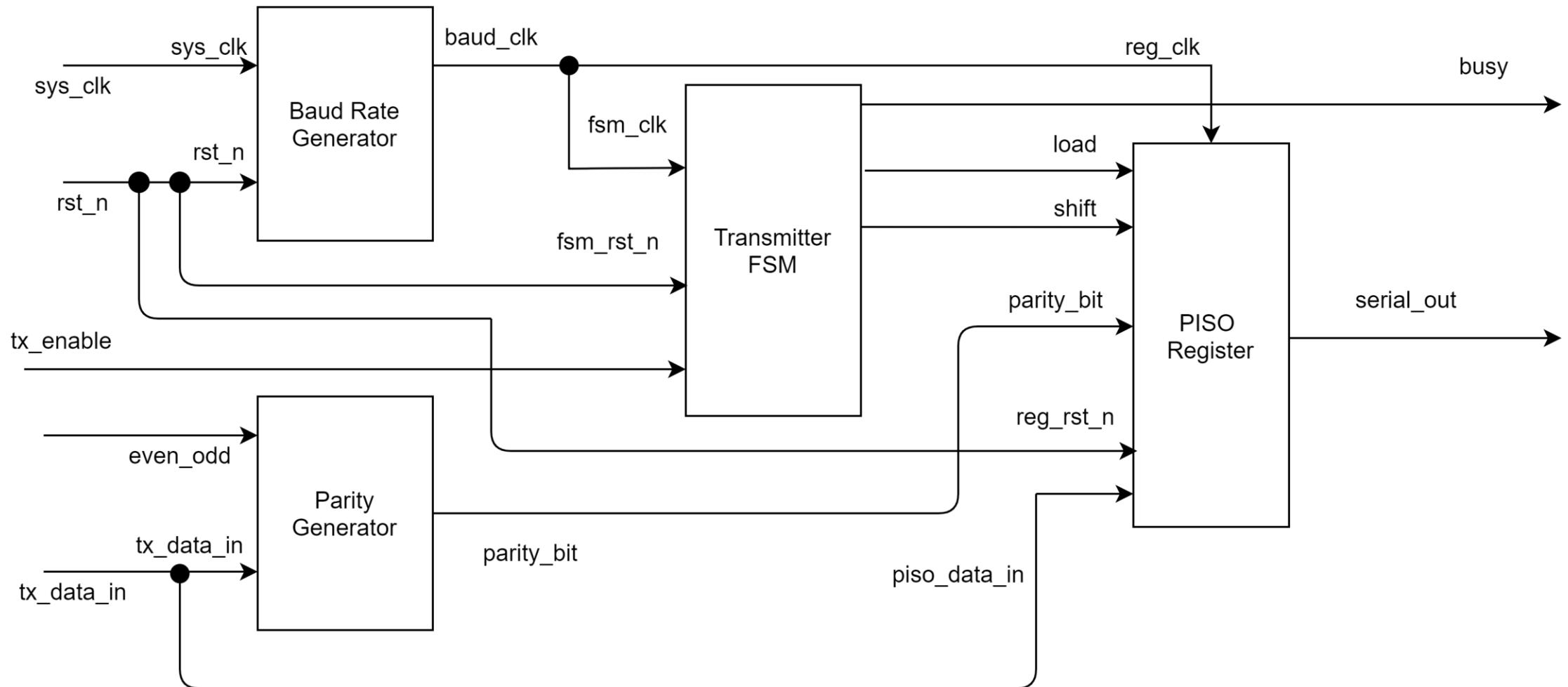
Sending clock of UART Transmitter



FSM for UART Transmitter



UART Transmitter Architecture

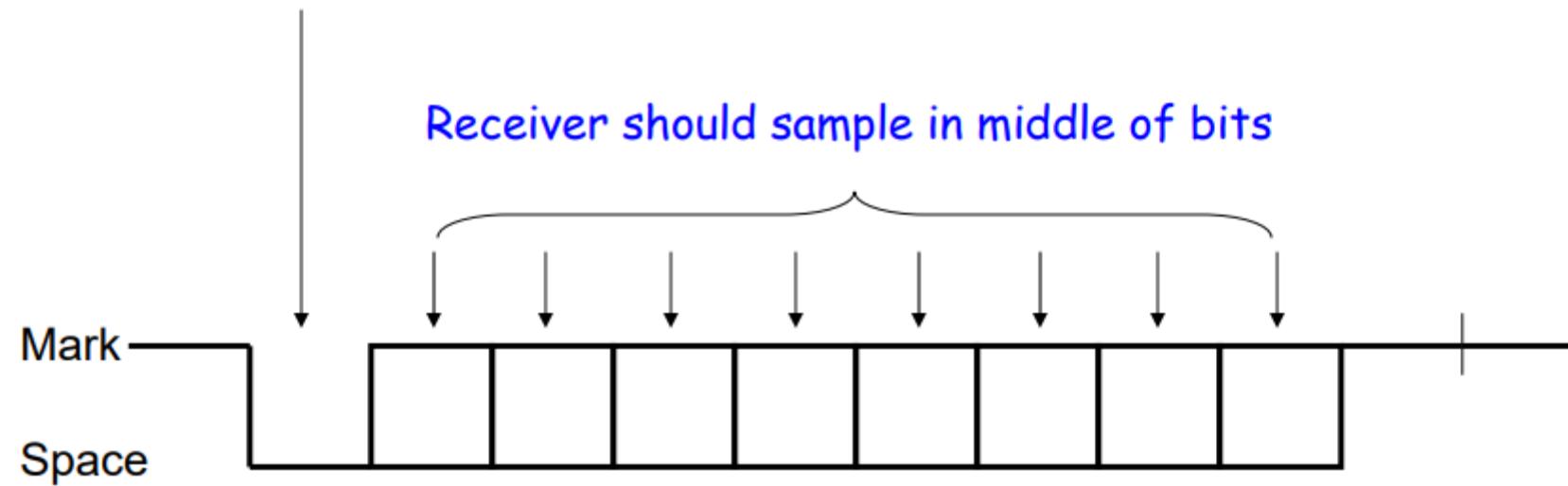


Process of Receiving

- On the receiver side once it receives all of the bits in the data word, it can check for the Parity Bits. To accomplish this task both transmitter and receiver must agree on whether a Parity Bit is to be used.
- Then Stop Bit is encountered by receiver. A missing stop bit may result entire data to be garbage. This will cause a Framing Error and will be reported to the host processor when the data word is read. Framing Error can be caused due to mismatch of transmitter and receiver clocks.
- The UART automatically discards the Start, Parity and Stop bits irrespective of whether data is received correctly or not. If the sender and receiver are configured identically, these bits are not passed to the host. To transmit new word, the Start Bit for the new word is sent as soon as the Stop Bit for the previous word has been sent.
- The transmission speed in asynchronous communication is measured by Baud Rate. A Baud Rate represents the number of bits that are actually being sent over the media. The Baud rate includes the Start, Stop and Parity bits. The Bit rate (Bits per Second-bps) represents the amount of data that is actually sent from the transmitting device to the other device. Speeds for UARTs are in bits per second (bit/s or bps), although often incorrectly called the baud rate. Standard baud rates are: 110, 300, 1200, , 115200,....., 1843200, and 2764800 bit/s.

UART Character Reception

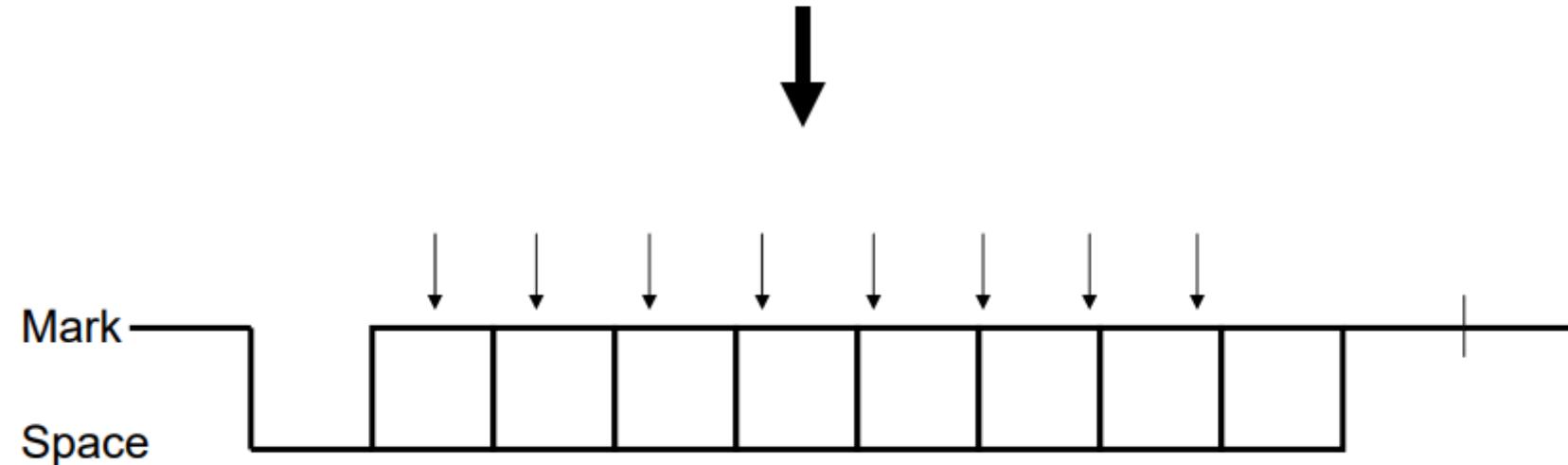
Start bit says a character is coming,
receiver resets its timers



Receiver uses a timer (counter) to time when it samples.
Transmission rate (i.e., bit width) must be known!

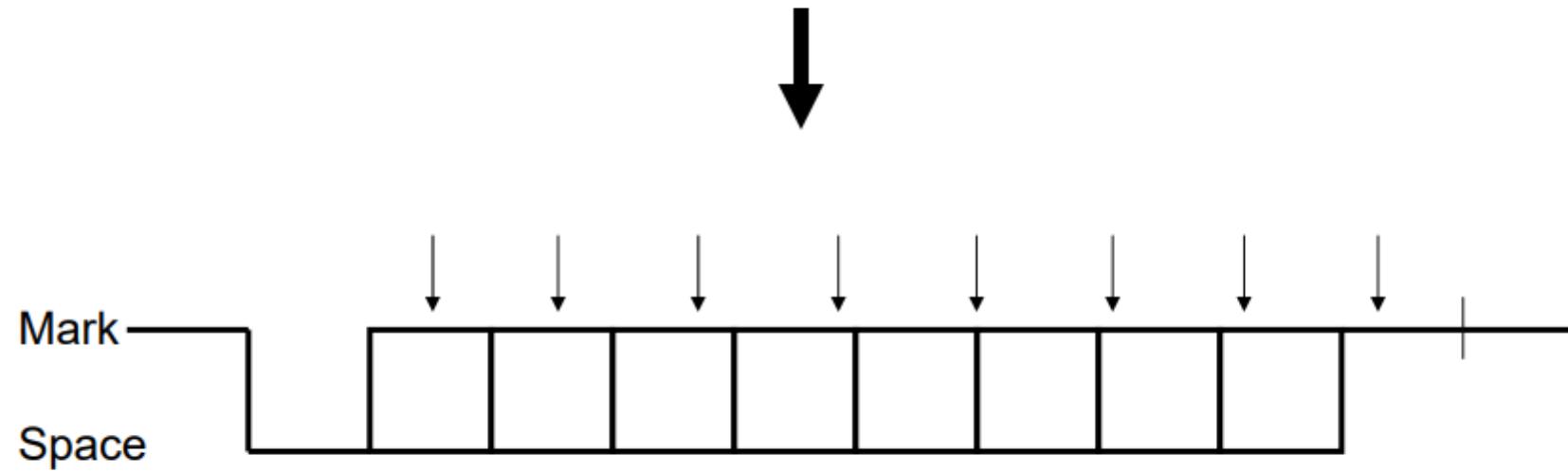
UART Receiver ‘Quick’ Sampling

If receiver samples too quickly, see what happens...



UART Receiver ‘Slow’ Sampling

If receiver samples too slowly, see what happens...



Receiver resynchronizes on every start bit.
Only has to be accurate enough to read 9 bits.

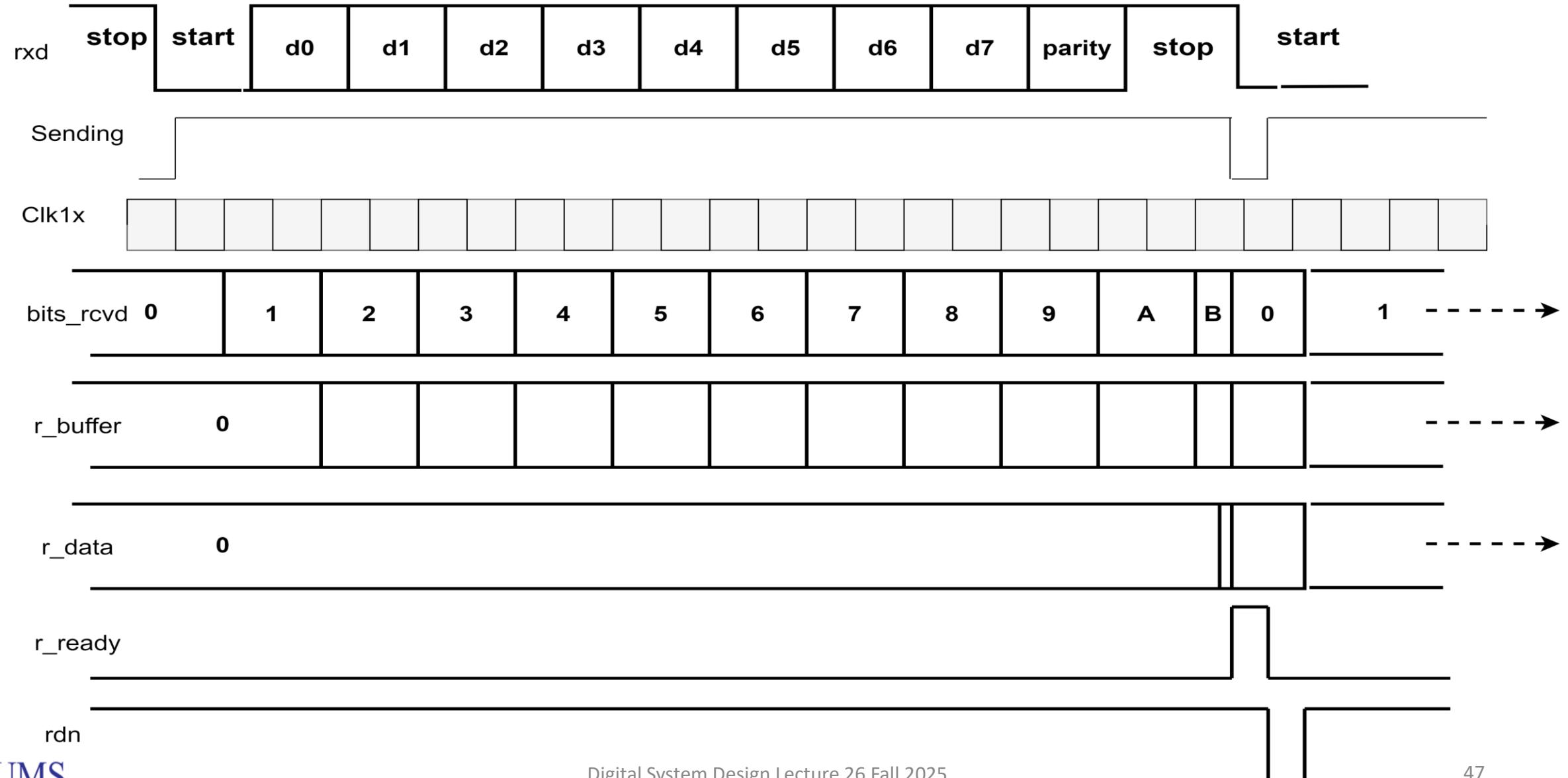
UART Receiving Stop Bit

- Receiver also verifies that stop bit is '1'
 - If not, reports "framing error" to host system
- New start bit can appear immediately after stop bit
 - Receiver will resynchronize on each start bit

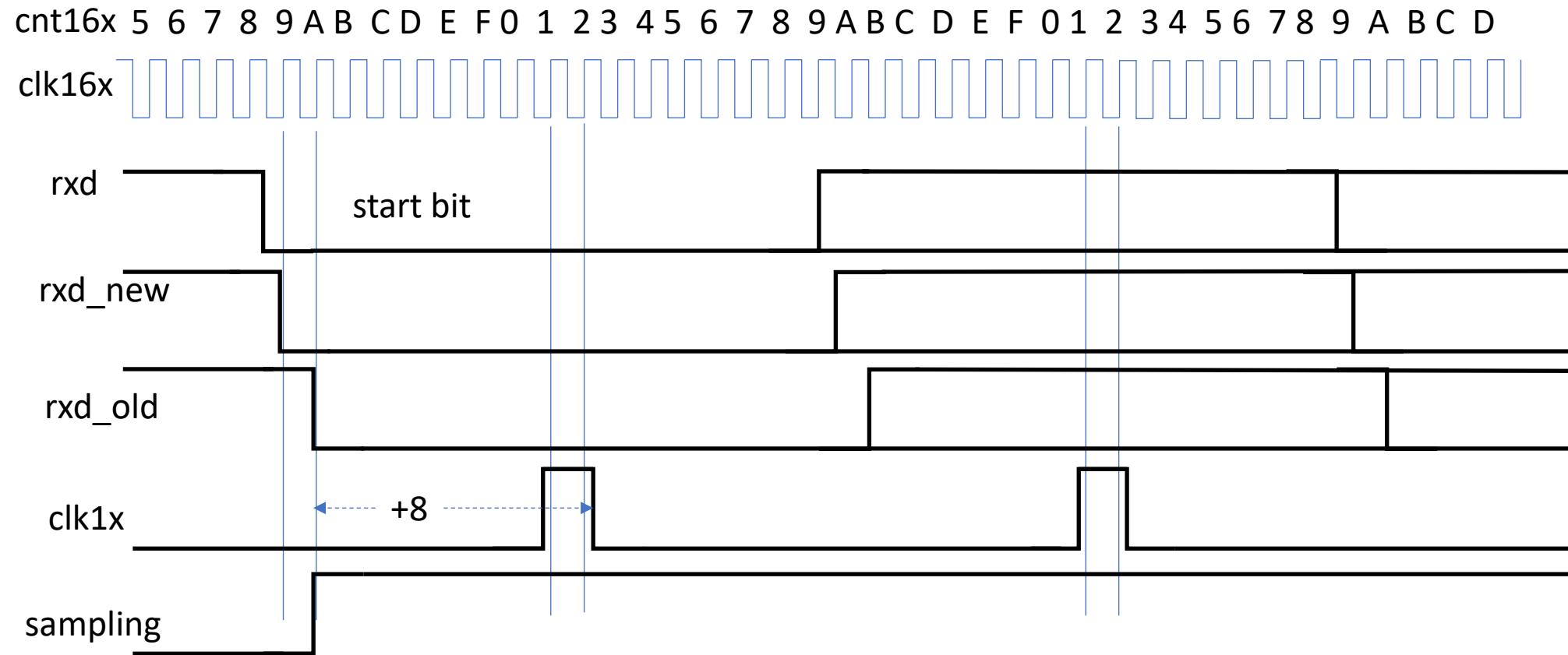
UART Receiver Configuration

- UARTs usually have programmable options:
 - **Data:** 7 or 8 bits
 - **Parity:** even, odd, none, mark, space
 - **Stop bits:** 1, 1.5, 2
 - **Baud rate:** 300, 1200, 2400, 4800, 9600, 19.2K, 38.4k, 57.6k, 115.2k...

Sampling Clock of UART Receiver



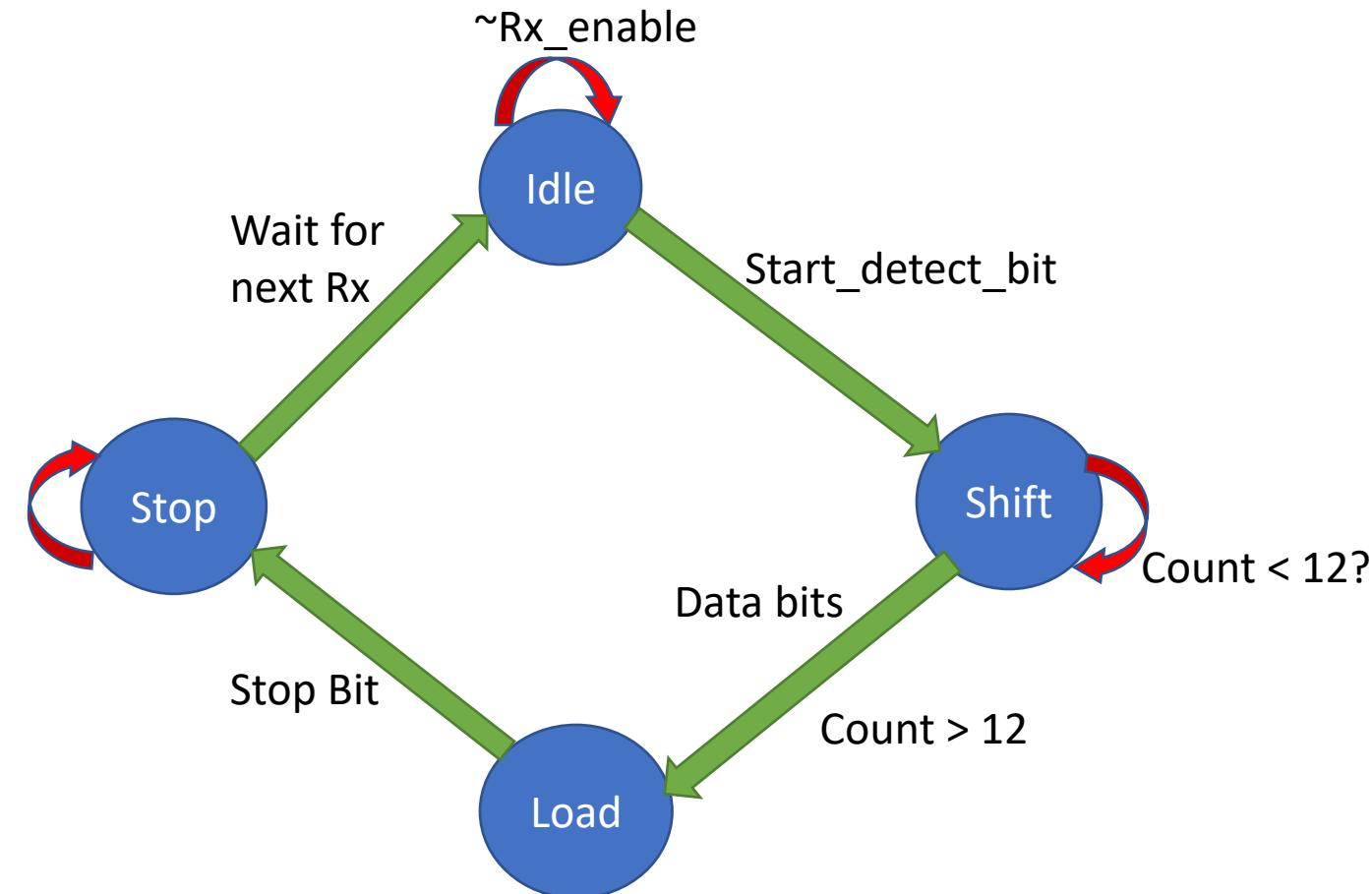
UART Receiver Timing Diagram



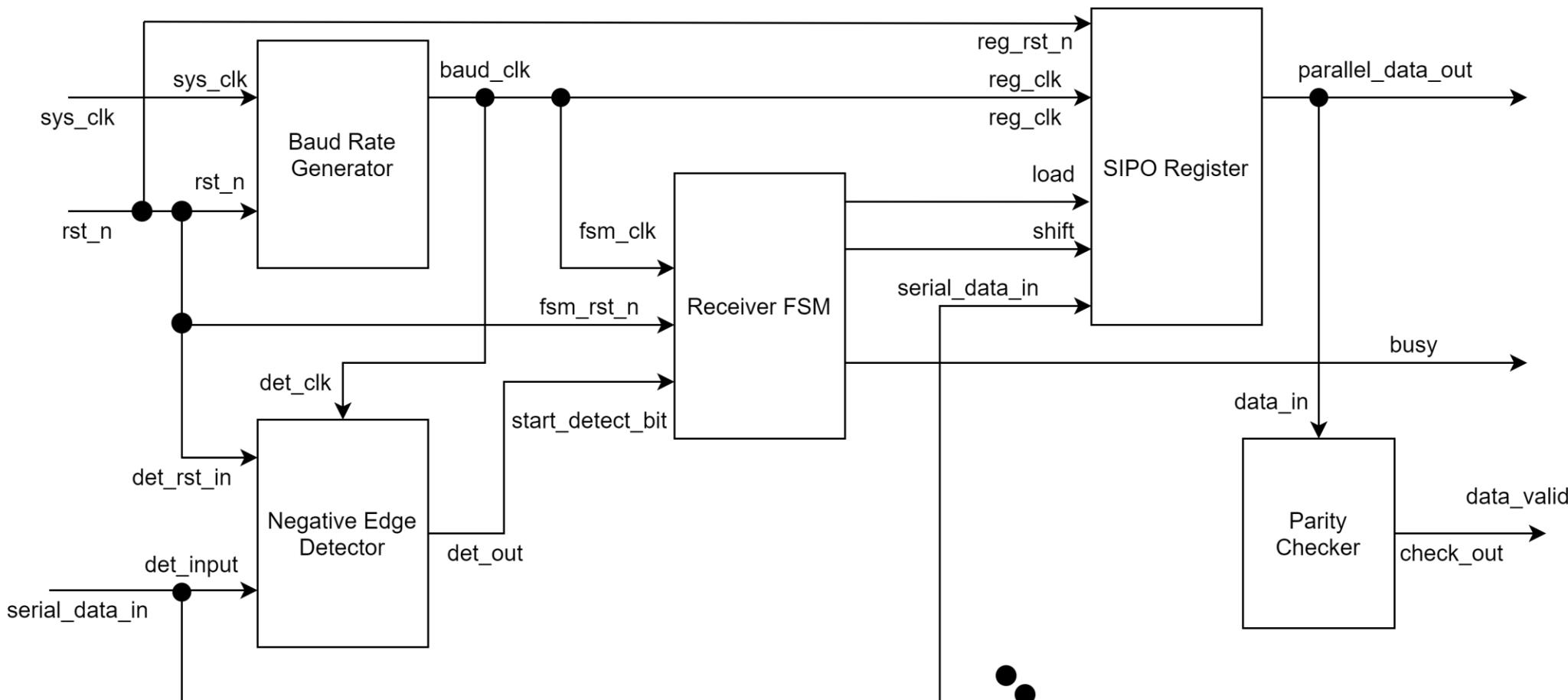
UART Receiver Starting details

- Clk16x is 16 times baud rate
- A 2-bit shift register saves input signal rxd
- When contents of this reg are 10 ($\text{rx}_d_{\text{old}}=1$, $\text{rx}_d_{\text{new}}=0$) then a falling edge is detected. We let sampling = 1 to sample rxd.
- A 4-bit counter cnt16x is used to generate a sampling pulse clk1x.
- During sampling=1, we let the rising edge of clk1x to appear at about centre place of one rxd bit; and the width of the clk1x pulse to be one clk16x clock cycle.

FSM for UART Receiver



UART Receiver Architecture



UART Errors

Overrun Error

An "**overrun error**" occurs when the UART cannot process the byte that just came in before the next one arrives. The host processor must service the UART in order to remove characters from the buffer. If the host processor does not service the UART and the buffer becomes full, then **Overrun Error** will occur.

Framing Error

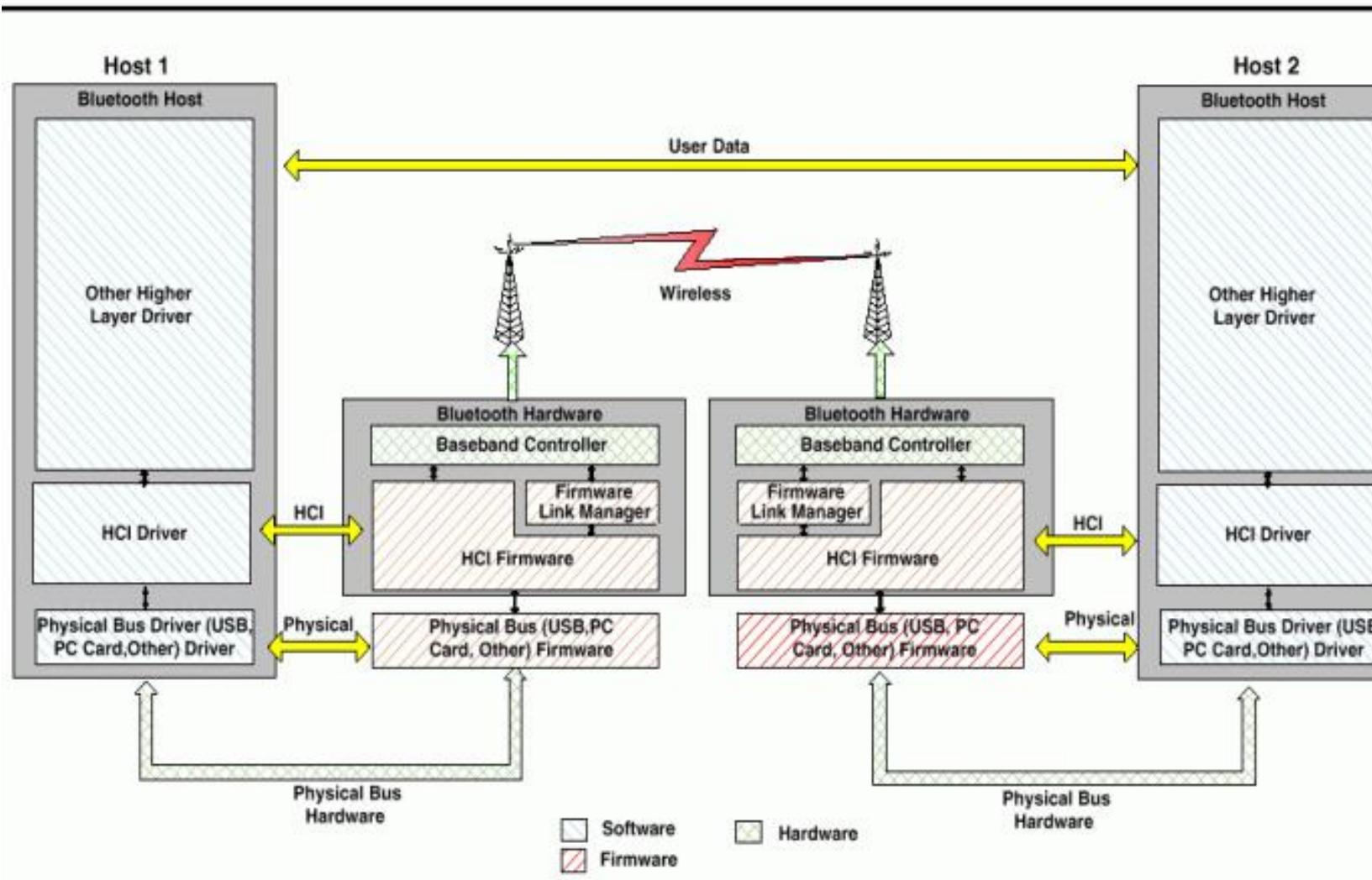
A "**Framing Error**" occurs when the designated "start" and "stop" bits are not valid. Start bit acts as a reference for the remaining bits. When the "stop" bit is expected if the data line is not in the expected idle state a **Framing Error** will occur.

Parity Error

A "**Parity Error**" occurs when the number of "active" bits does not agree with the specified parity configuration of the UART.

Bluetooth Architecture

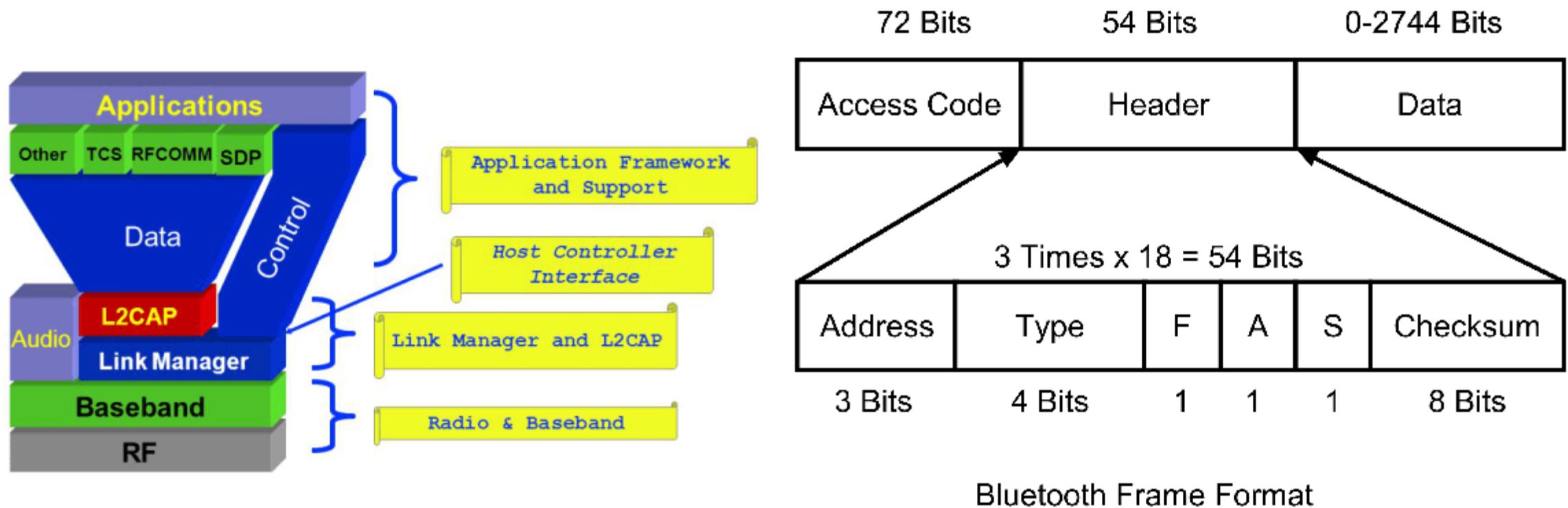
The basic structure showing how the host controller layers are fitted into the protocol stack



Comm Example - Bluetooth

- Some details of Bluetooth protocol and its hardware design

Bluetooth Architecture



Bluetooth Packet Format

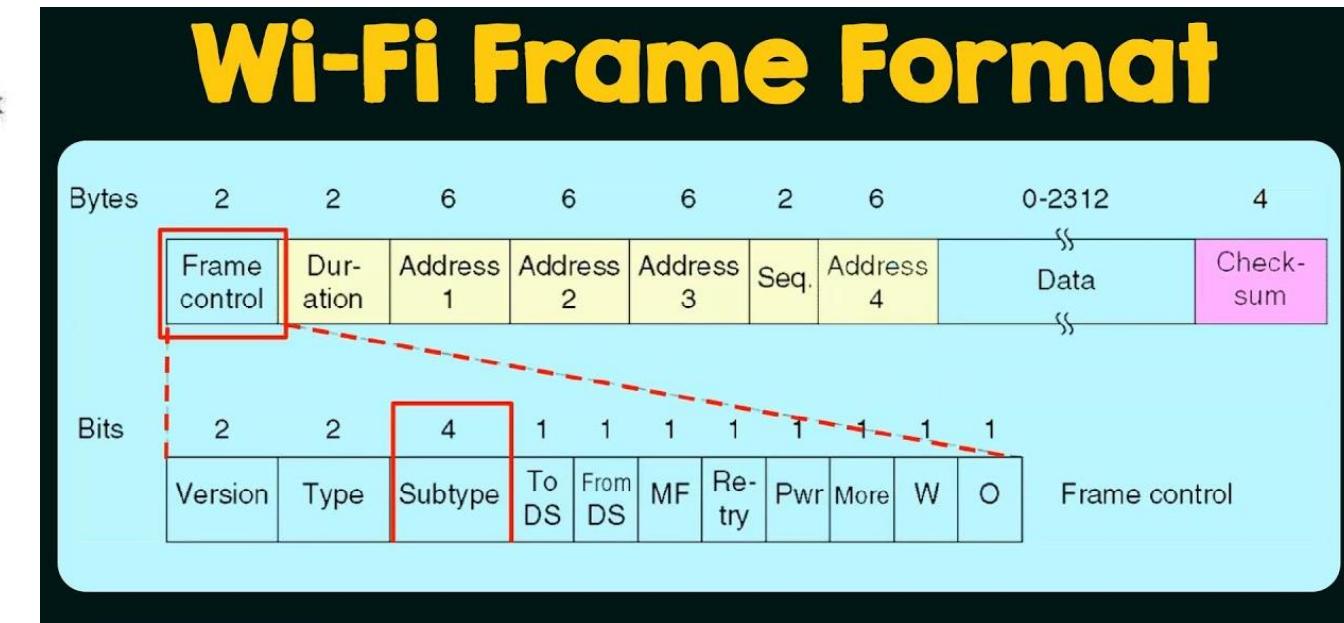
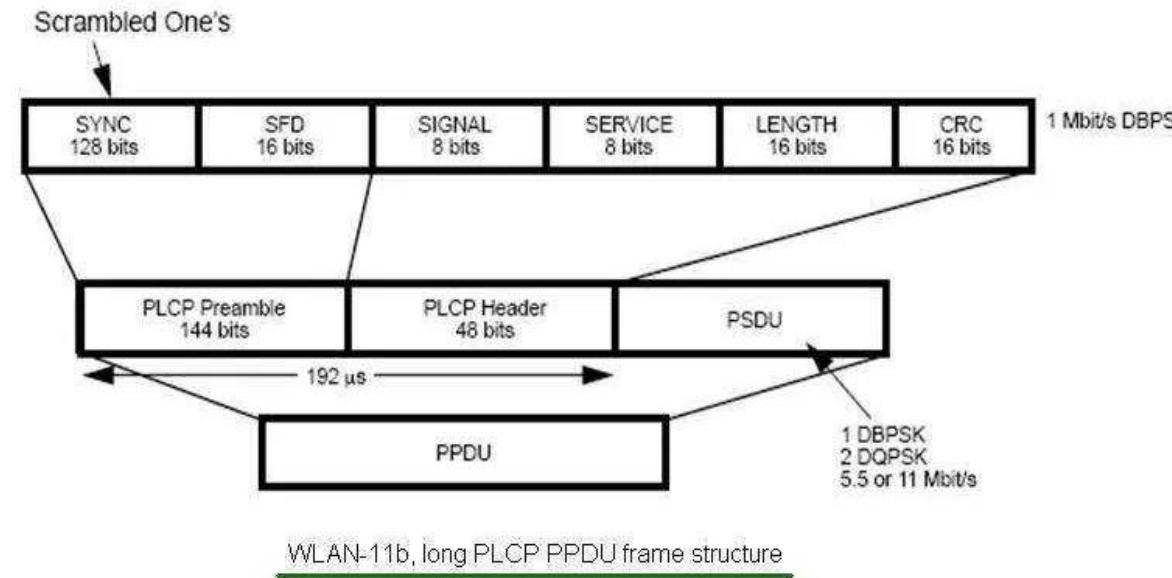
Bluetooth Packet Format

Access Code	Baseband/Link Control Header	Data Payload
72b	54b	0-2745b

- Packets can be up to five slots long. 5 slots = 3125 bits.
- Access codes:
 - Channel access code identifies the piconet
 - Device access code for paging requests and response
 - Inquiry access code to discover units
- Header: member address (3b), type code (4b), flow control, ack/nack (1b), sequence number, and header error check (8b)
18b Header is encoded using 1/3 rate FEC resulting in 54b
- Synchronous traffic has periodic reserved slots.
- Other slots can be allocated for asynchronous traffic

Comm Example – Wifi - Ethernet

- May discuss hardware of Wifi, Ethernet



Comm Example - Zibee

- Some details of Zigbee protocol and hardware design