Lecture 26
EE 421 / CS 425
Digital System Design

Fall 2024

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Topics

- BRM Design and Applications

- BRM Chips

- UART Introduction

- UART Configuration

- UART Transmitter Operation

- UART Receiver Operation

- UART Architecture

- UART Registers

- **Application of Serial Comm – Bluetooth**

- **Application of Serial Comm – Zigbee**

- **Possibly Ethernet Hardware or similar**

Binary Rate Multipliers - BRM

Definition: A circuit module that transforms a stream of input clock pulses into another stream of output clock pulses

Let Ni = no. of input pulses for a particular time period

Let No = no. of output clock pulses for the same time period

$$\textbf{\textit{Binary Fractional Rate Multiplier}} \implies \textbf{\textit{No}} = \frac{\textbf{\textit{B}}}{\textbf{\textit{2}}^{\textbf{\textit{n}}}} \textbf{\textit{Ni}}$$

B = Rate Constant Input to module

B = $(B_{n-1}, B_{n-2}, \ldots, B_2, B_1, B_0)_2$

n = no. of binary counter stages controlling the module

The clock input drives an n-bit binary counter whose outputs are labelled (Xn,Xn-1,....,X2,X1).

Will do in next lecture due to quiz today – from previous lecture

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Digital Fractional Rate Multiplier - Principle

The counter outputs (Xn to X1 = MSB to LSB) are ANDed with the incoming Clock signals to form intermediate pulse trains Pi(i=1,n):

$P_n = X_1 . Clock$

$P_{n-1} = \overline{X_1} . X_2 . Clock$

$P_{n-2} = \overline{X_1} . \overline{X_2} . X_3 . Clock$

......

$P_2 = \overline{X_1} . \overline{X_2} . \overline{X_3} ... . \overline{X_{n\_2}} . X_{n\_1} . Clock$

$P_1 = \overline{X_1} . \overline{X_2} . \overline{X_3} ... . \overline{X_{n\_2}} . \overline{X_{n\_1}} . Xn . Clock$

The logic output Y uses the rate constant M and the pulse train signals Pi to implement the equation $\boldsymbol{Y = \sum B_{i-1} . P_i}$

This output configuration will deliver the proper number of output pulses (N_o) as specified by rate constant B

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Basic Circuit Diagram – DFRM

$P_n = X_1 . Clock$
$P_{n-1} = \overline{X_1} . X_2 . Clock$
$P_{n-2} = \overline{X_1} . \overline{X_2} . X_3 . Clock$
......
$P_2 = \overline{X_1} . \overline{X_2} . \overline{X_3} ... .. \overline{X_{n\_2}} . X_{n\_1} . Clock$
$P_1 = \overline{X_1} . \overline{X_2} . \overline{X_3} ... .. \overline{X_{n\_2}} . \overline{X_{n\_1}} . Xn . Clock$
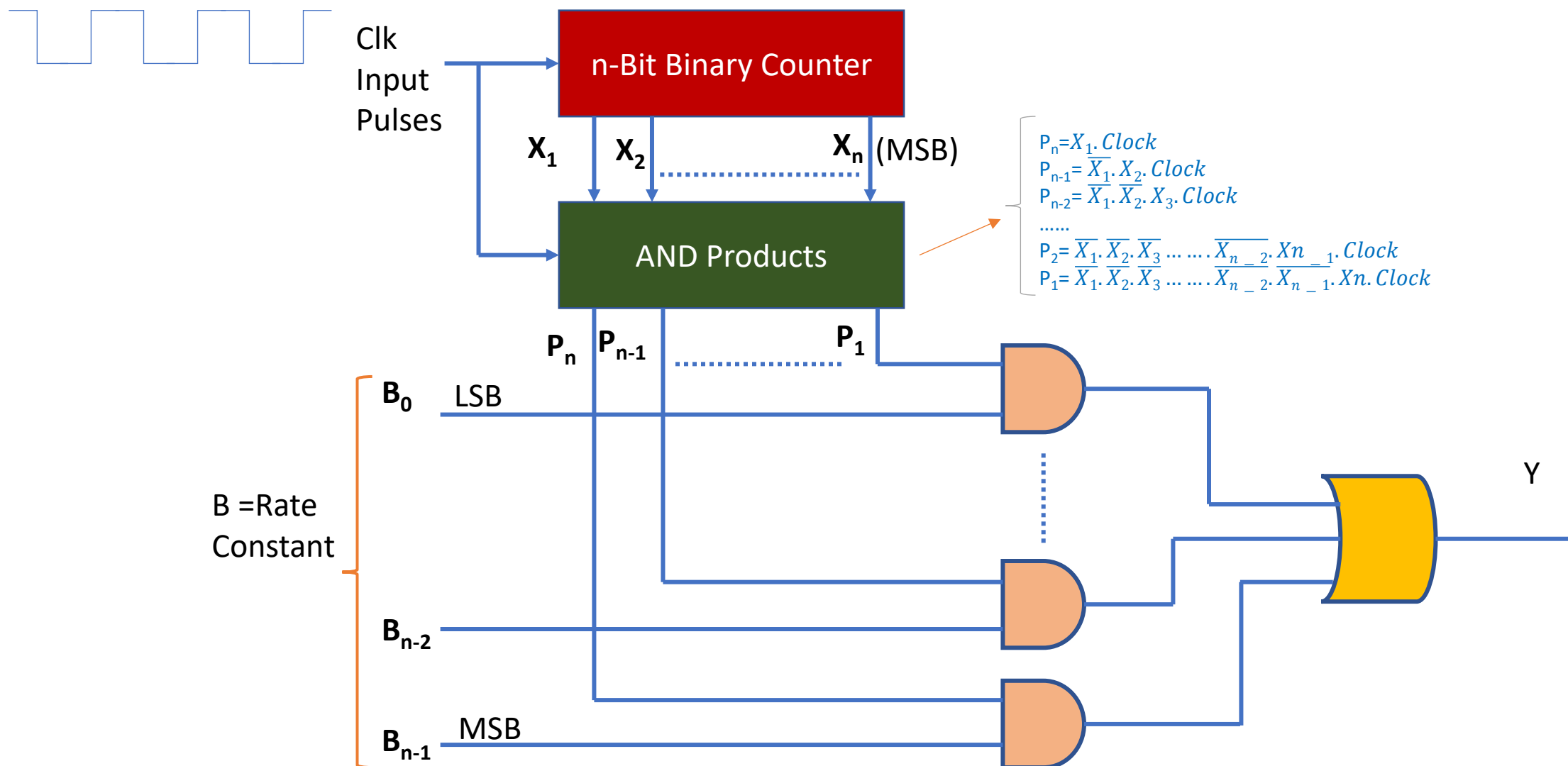
Example and timing diagram of BRM

- BRM using a 3-Bit counter

- Rate Constant B = 7_{10}

- Counter generates 8 pulses per time period; Due to B=7 , 7 pulses appear at the output and one pulse is eliminated

Analysis of timing diagram 3 bit BRM, B=7

- $Y=B_2.P_3 + B_1.P_2 + B_0.P_1$
- Where
- **$P_3=X_1.Clock$**
- **$P_2=\overline{X_1}.X_2.Clock$**
- **$P_1=\overline{X_1}.\overline{X_2}.X_3.Clock$**
- And $B=(B_2.B_1.B_0)=(1.1.1)_2$
- Thus $Y=(X_1 + \overline{X_1}X_2 + \overline{X_1}.\overline{X_2}.X_3).Clock$
- In terms of counter output $X=(X_3.X_2.X_1)$
- Product term $X_1 \Longrightarrow \sum m(1,3,5,7)$
- Product term $X_2.\overline{X_1} \Longrightarrow \sum m(2,6)$
- Product term $X_3.\overline{X_2}.\overline{X_1} \Longrightarrow \sum m(4)$
- So that $Y(X_3.X_2.X_1) = (\sum m(1,2,3,4,5,6,7)).Clock$
- As m_0 is missing from this list so first clock from each sequence of 8 pulses will be eliminated

Timing Diagram of DFRM
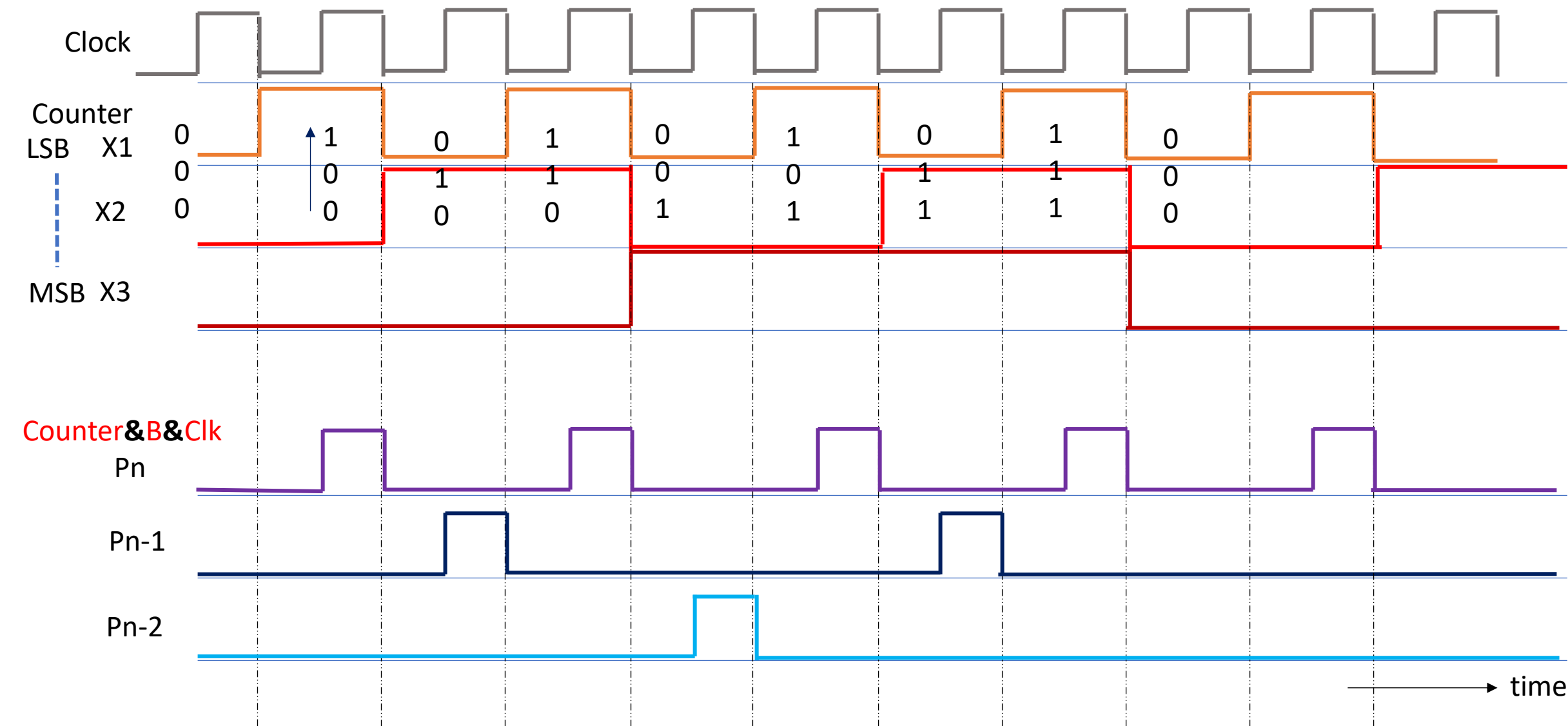
How many pulses are produced?

- From the timing diagram, each of the pulse trains P_i generates 2^{i-1} pulses during during one counter sequence period (2^n clock pulses)

- Thus P_1 generates 1 pulse, P_2 generates 2 pulses, P_3 generates 4 pulses, and so on

- Note that pulses do not overlap in time

- We may OR the respective pulses to produce desired output stream

- The output pulses can be irregularly spaced

- The output pulses are synchronized with input clock pulses

Serially Cascaded Binary Rate Multiplier BRM

Design a BRM to implement Nout = $\frac{63}{320}$. Nin

Break into binary powers (63 x 2 = 126; and 320 x 2 = 640) :

Nout = $\frac{(7 \times 18)}{(10 \times 64)}$. Nin

Nout = $(\frac{7}{10}).(\frac{18}{64})$. Nin

Use:

1 x Decade Counter BRM

1 x 6-Bit Binary Counter BRM

Set B for Decade Counter = 0111

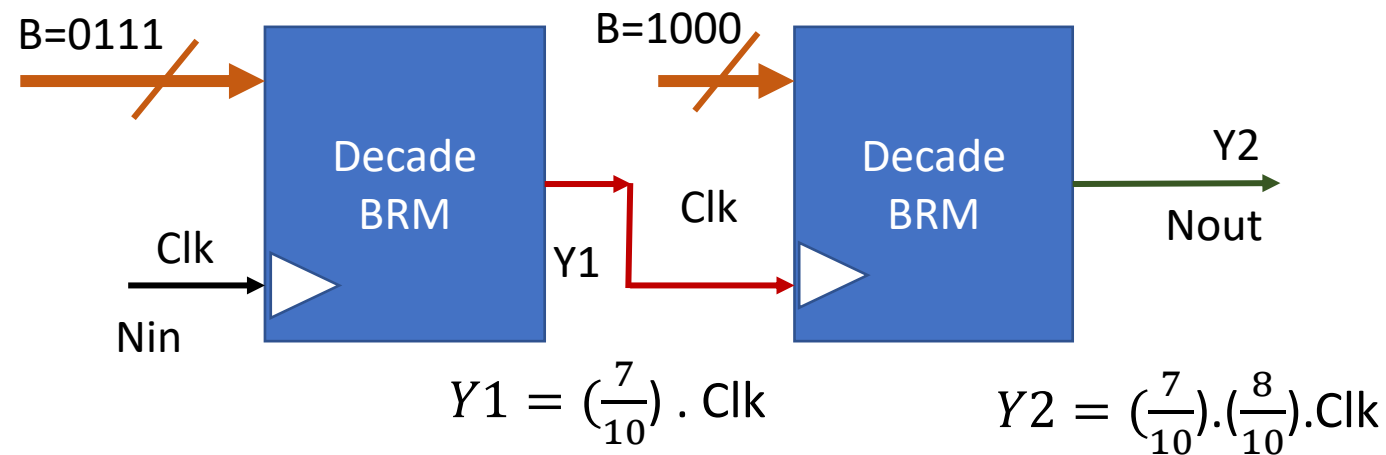Set B for 6-Bit BRM = 10010
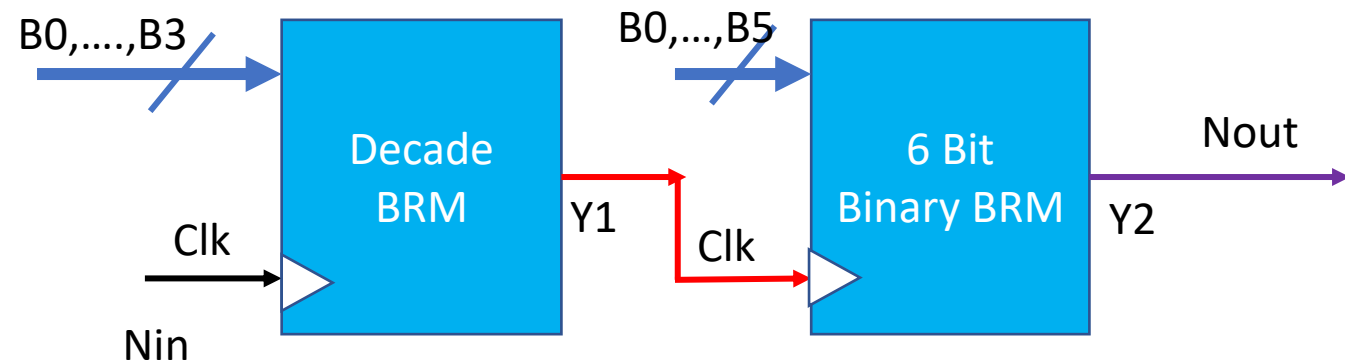
Example 2: Serially Cascaded BRM

Design a BRM circuit that produces: 0.56 of input clock rate

We use two Decade BRMs with

Value of B is set as 7 and 8 respectively

$$Nout = \frac{56}{100}. \text{Nin}$$

$$N = \frac{7}{10}.\frac{8}{10}. \text{Nin}$$

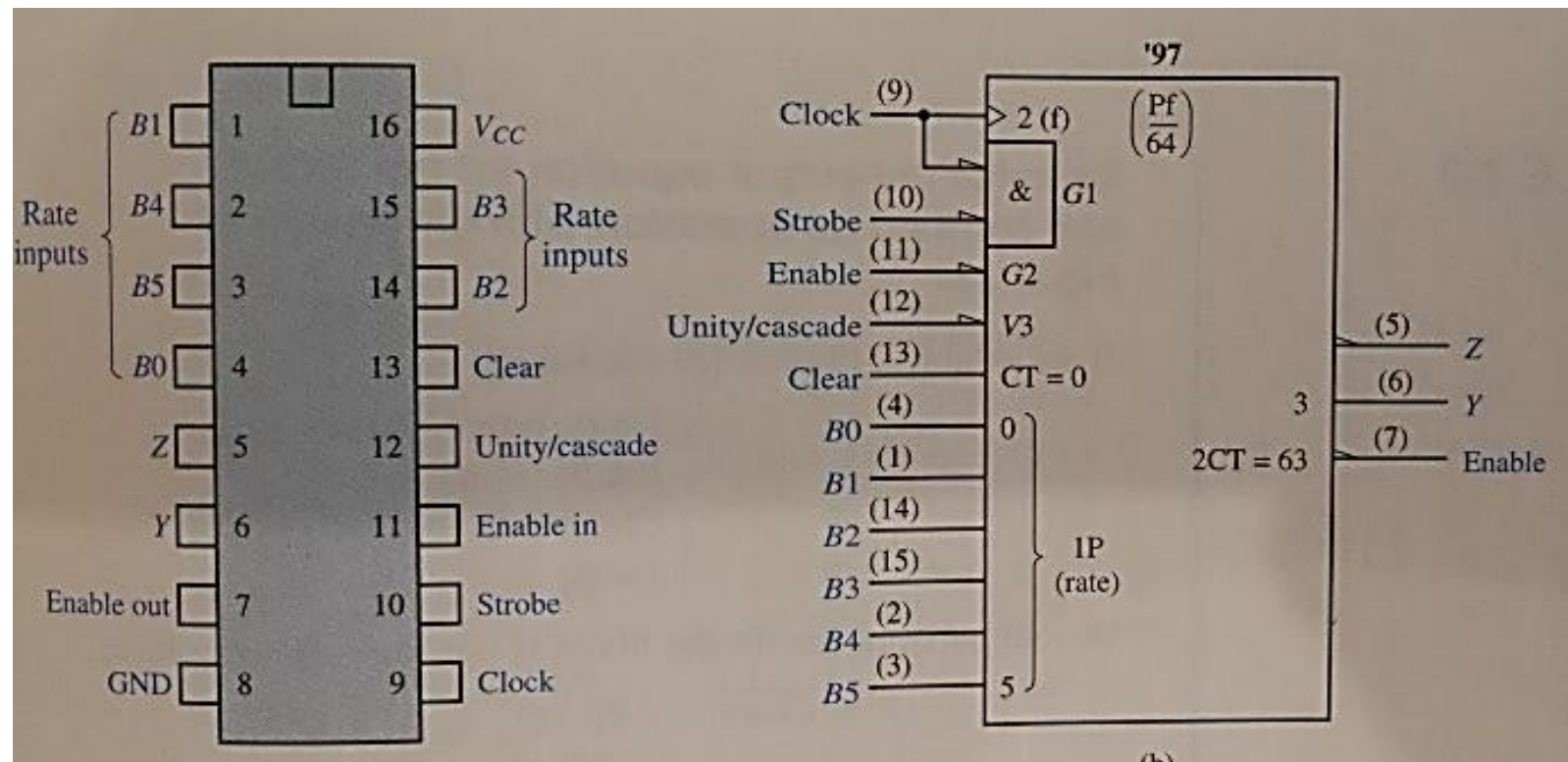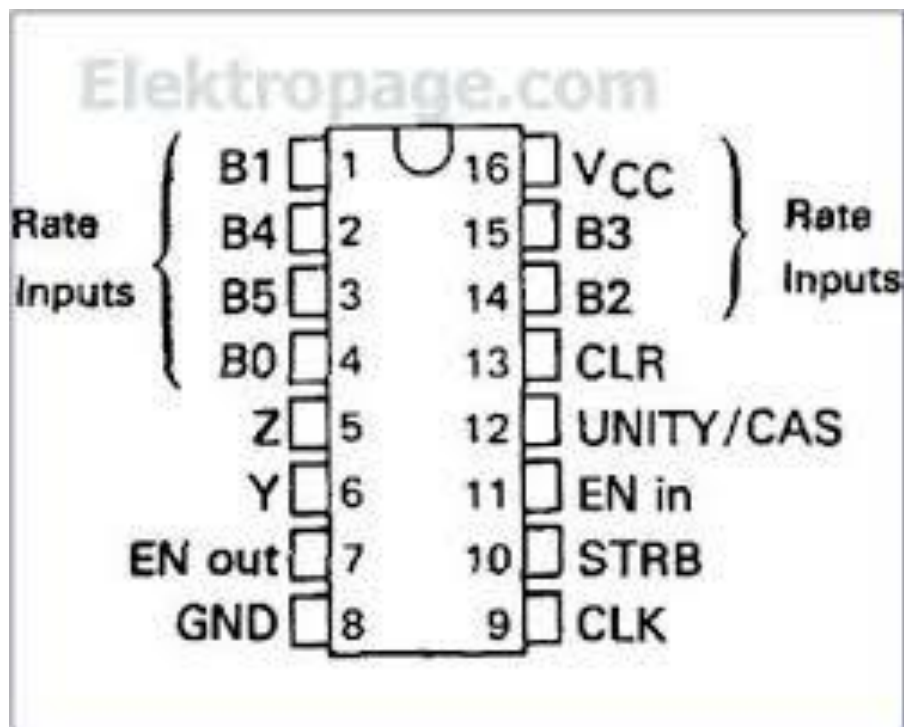$$Y1 = (\frac{7}{10}) . \text{Clk}$$

$$Y2 = (\frac{7}{10}).(\frac{8}{10}).\text{Clk}$$

Eg 3: Parallel Cascaded Binary Rate Multiplier

Design a BRM circuit to produce $N_{out} = (0.297)N_{in}$
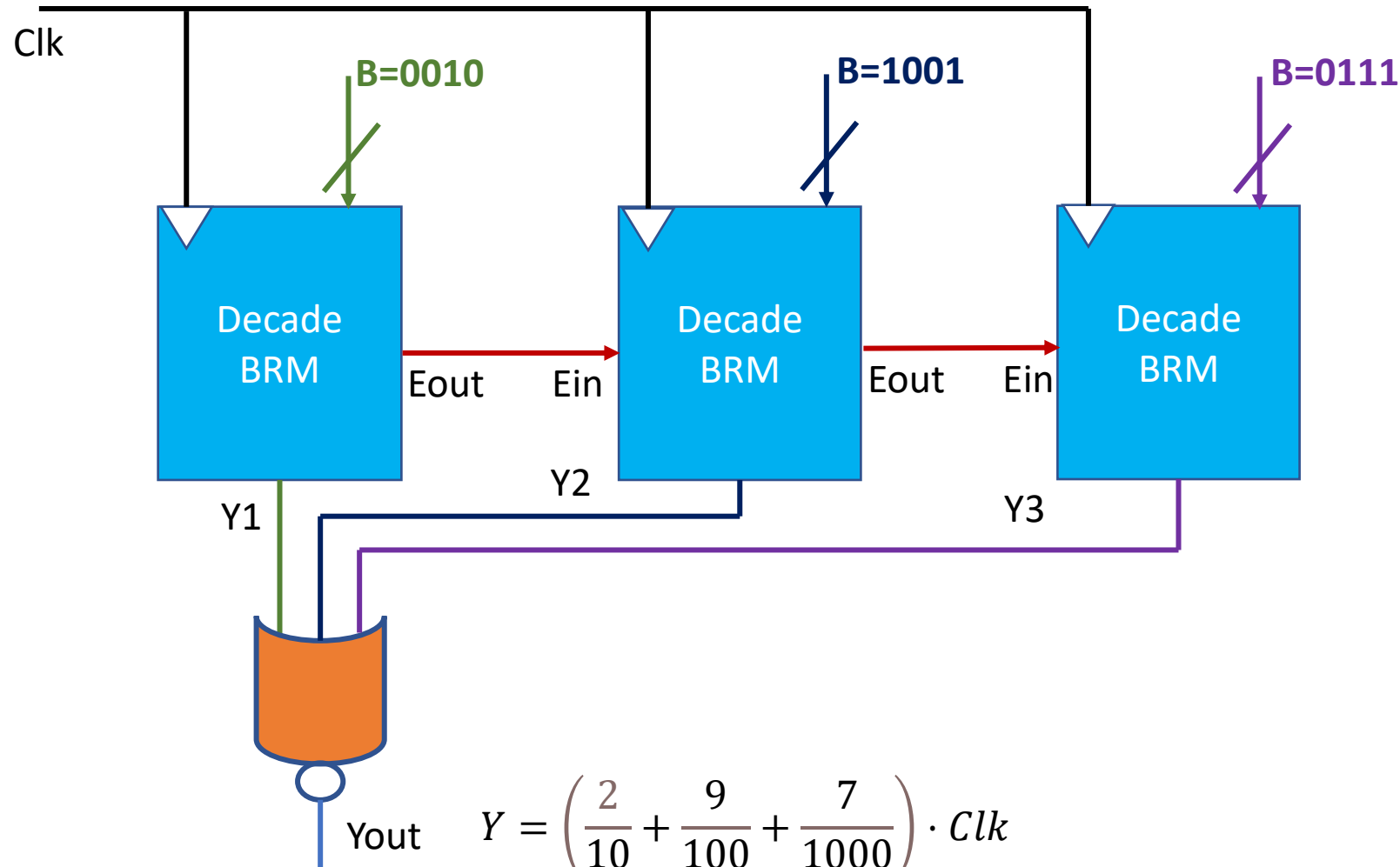
$$Y = \left(\frac{2}{10} + \frac{9}{100} + \frac{7}{1000} \right) \cdot Clk$$

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# 7497 TTL chip 4-Bit and SN7497E chip 6-Bit BRM

Ref: Images from internet

Pins on 7497 and SN7497E

- <u>Clear</u> signal is active high and asynchronous to drive all flip-flops to all-zero state

- Rising edge on <u>Clock</u> pulses causes counter to advance

- <u>Strobe</u> is an active-low Enable signal for the pulse rate output Y and Z

- <u>Enable</u> input and Enable output allows the modules to be cascaded

- The <u>Unity / Cascade</u> input may be used to combine the outputs of cascaded units, effectively ORing the pulses trains from two modules
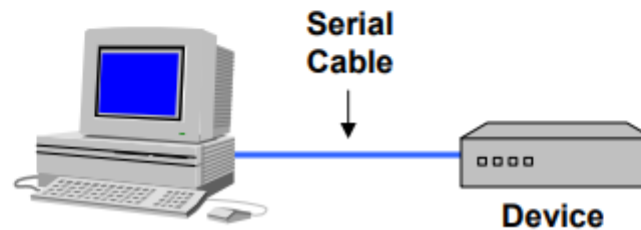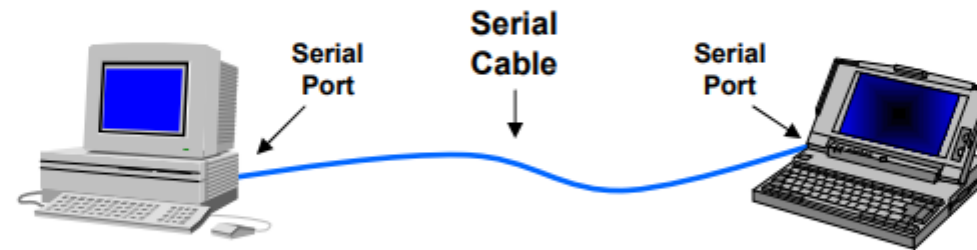
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Internal view of BRM chip
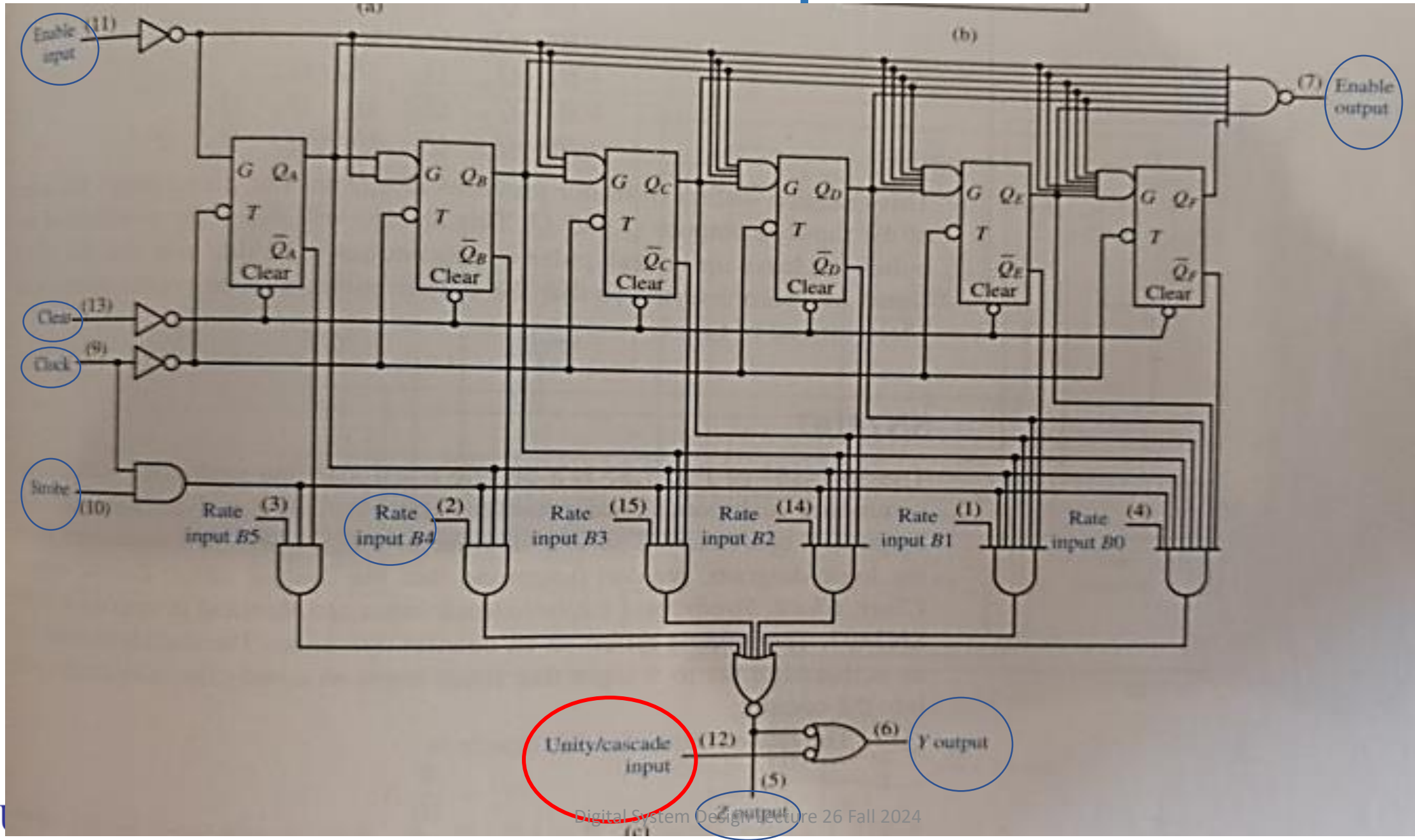
UART RS232 Protocol and Hardware Implementation

UART for Serial Communication

- PC serial port is a UART!
- Serializes data to be sent over serial cable
 - De-serializes received data

UART Functions

- Outbound data
  - Convert from parallel to serial
  - Add start and stop delineators (bits)
  - Add parity bit

- Inbound data
  - Convert from serial to parallel
  - Remove start and stop delineators (bits)
  - Check and remove parity bit

UART Principle

- UART translates Parallel Data (Bytes or ASCII) to Serial Data for sending and vice versa for receiving.

- Inside UART there are several registers to store status and data for sending and receiving.

- The CPU can read and write these registers through I/O ports.

- For correct communication, both UARTs must make agreement on data frame format and transmission speed (baud rate).

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Serial Port Setting on PC using HyperTerminal
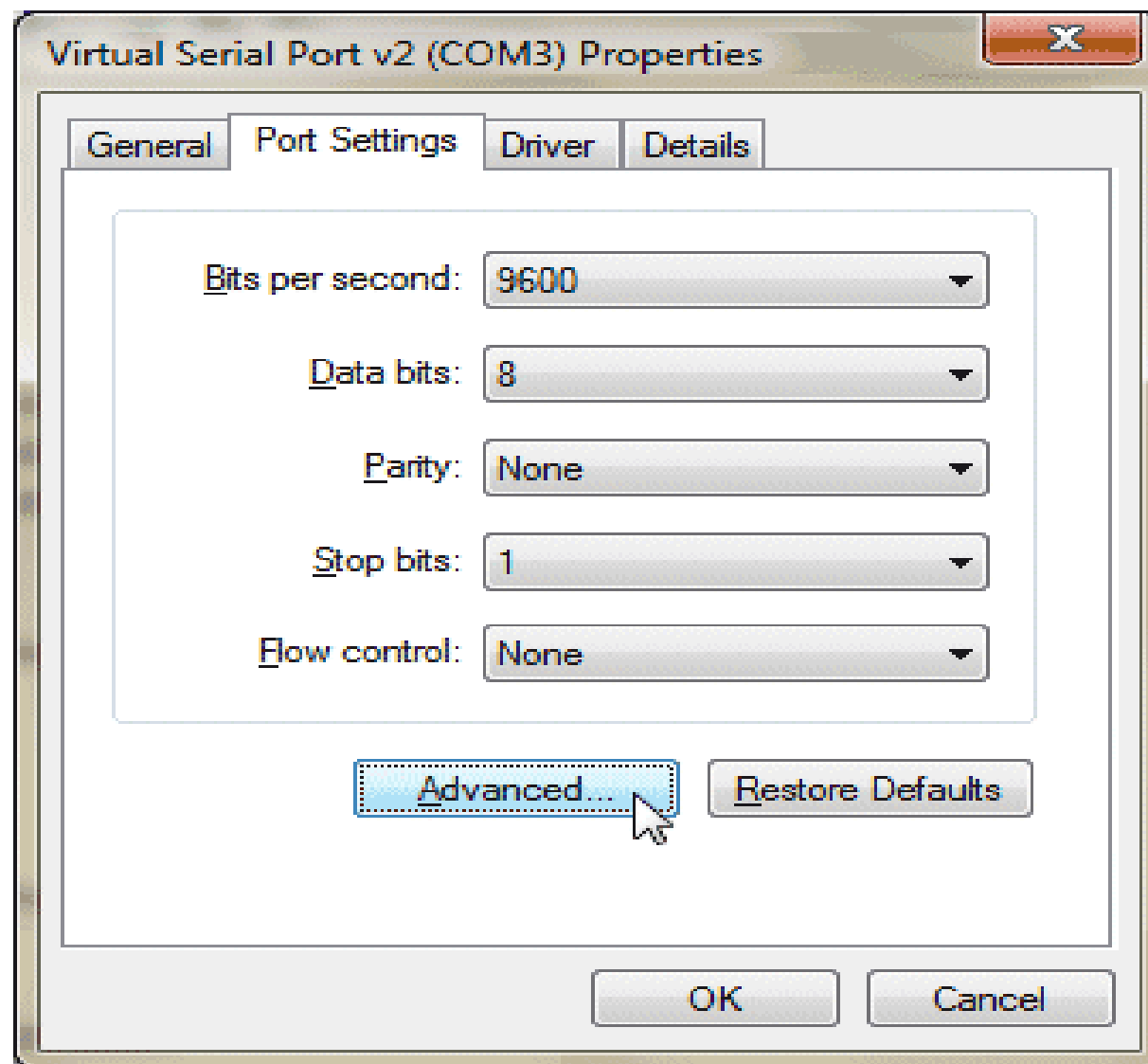
Configurable Baud Rate

- Definition: <mark>"How many times a signal can change per second" Includes data bits plus control bits.</mark>

- Possible Baud Rates:
  - 110
  - 150
  - 300
  - 600
  - 1200
  - 2400
  - 4800
  - 9600
  - 19200
  - 38400
  - 57600, 115200, 230400, etc.

UART Throughput

- Data Throughput Example
  - Assume 19200 baud, 8 data bits, no parity, 1 stop bit
    - 19200 baud → 19.2 kbps
    - 1 start bit + 8 data bits + 1 stop bit → 10 bits
  - It takes 10 bits to send 8 bits (1 byte) of data
  - 19.2 kbps · 8/10 = **15.36 kbps**

- How many KB (kilobytes) per second is this?
  - 1 byte = 8 bits
  - 1 KB = 1,024 bytes
  - So, 1 KB = 1,024 bytes · 8 bits/byte = 8,192 bits
  - Finally, 15,360 bps · 1 KB / 8,192 bits = **1.875 KB/s**

UART Character Transmission

- Each bit has a fixed time duration determined by the transmission rate
- Example: a 1200 bps (bits per second) UART will have a 1/1200 s or about 833.3 μs bit width
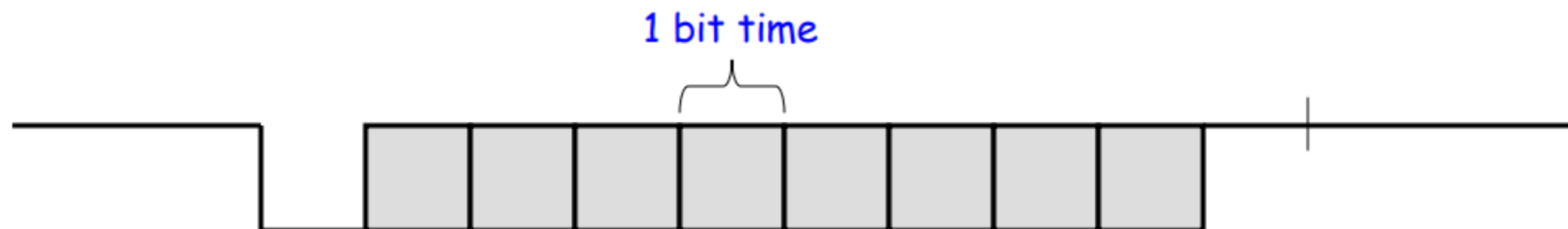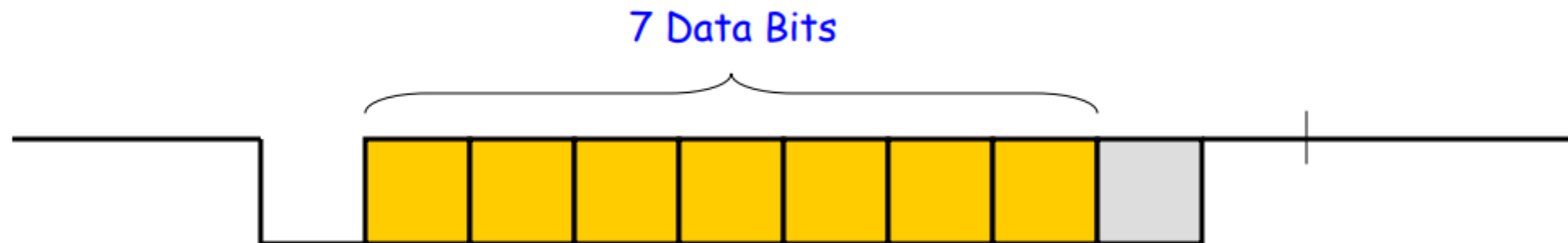
UART Start Bit in Transmission

- The **start bit** marks the beginning of a new word
- When detected, the receiver synchronizes with the new data stream

Start Bit

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UART Data Bits in Transmission

- Next follows the **data bits** (7 or 8)
- The least significant bit is sent first

7 Data Bits

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UART Parity Bit in Transmission

- The **parity bit** is added to make the number of 1's even (even parity) or odd (odd parity)
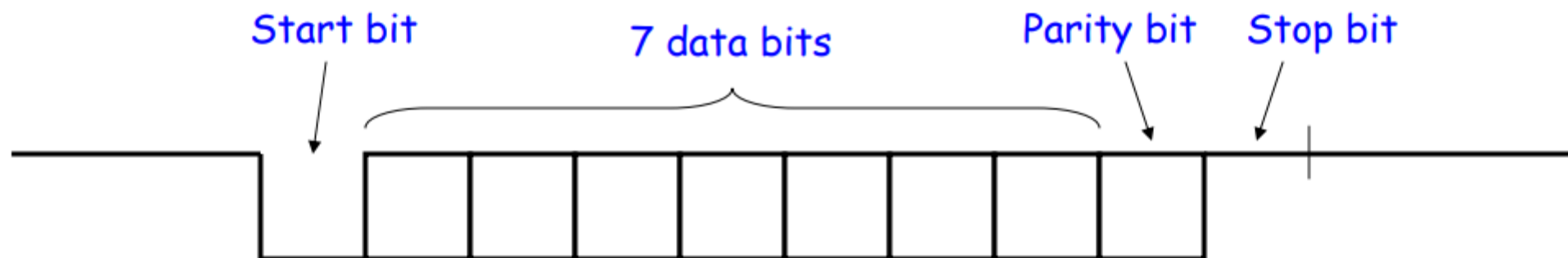- This bit can be used by the receiver to check for transmission errors

Parity Bit

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UART Stop Bit in Transmission

- The **stop bit** marks the end of transmission
- Receiver checks to make sure it is '1'
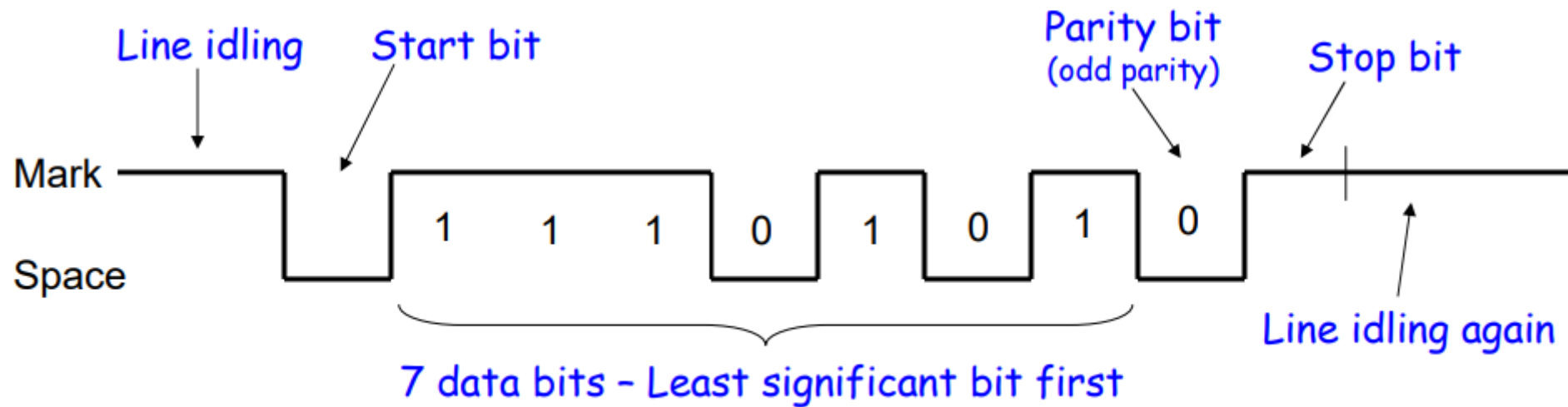- Separates one word from the start bit of the next word

UART Transmission of 7 Bits

- In the configuration shown, it takes 10 bits to send 7 bits of data
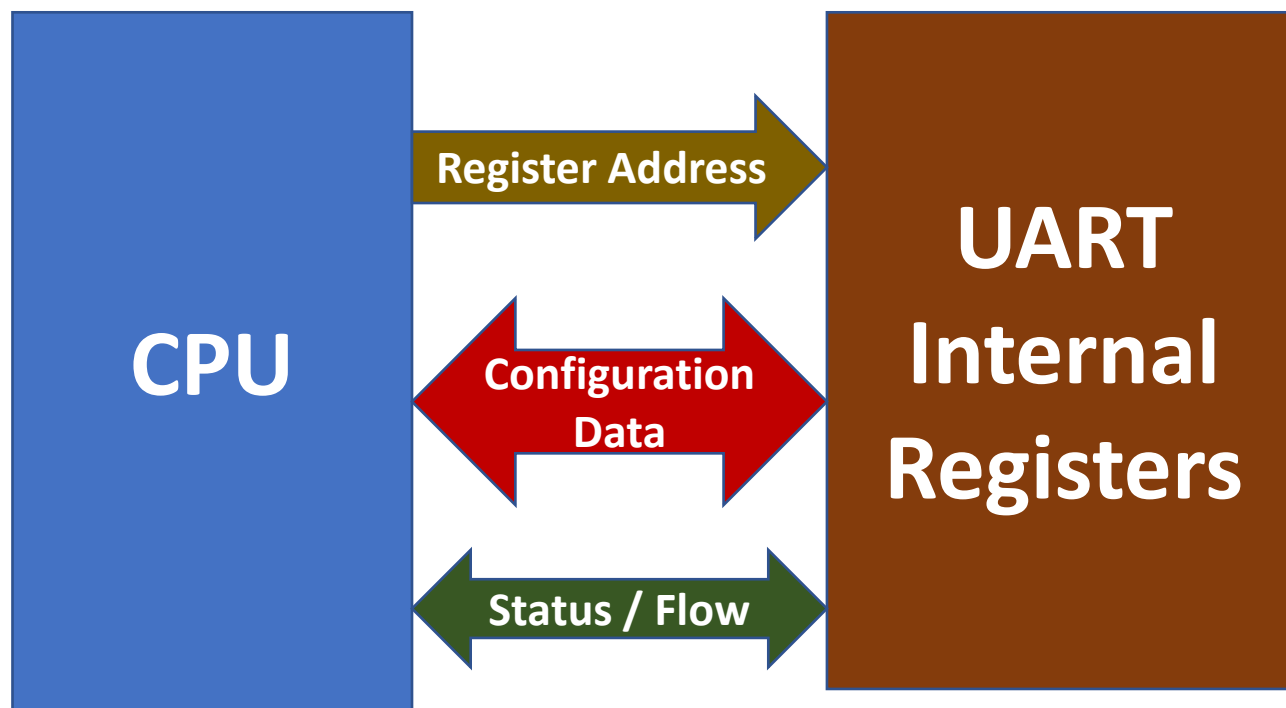
UART Transmit Ascii 'W'

- Send the ASCII letter 'W' (1010111)

Line idling   Start bit   Parity bit (odd parity)   Stop bit

Mark   1   1   1   0   1   0   1   0

Space

7 data bits – Least significant bit first

Line idling again

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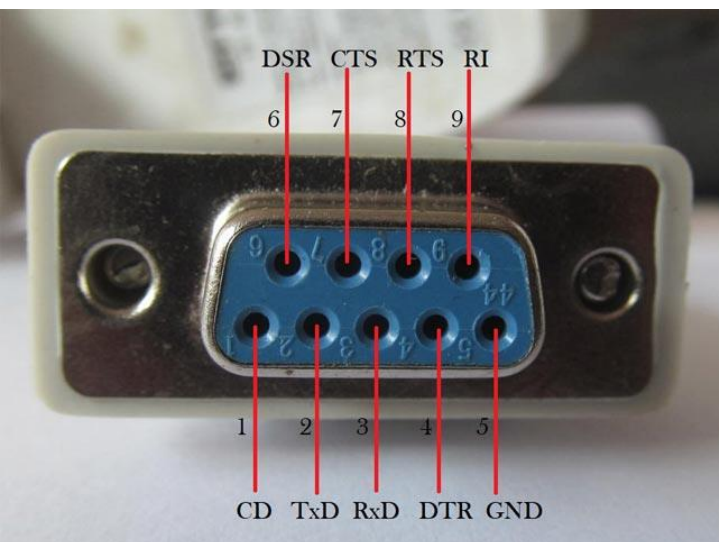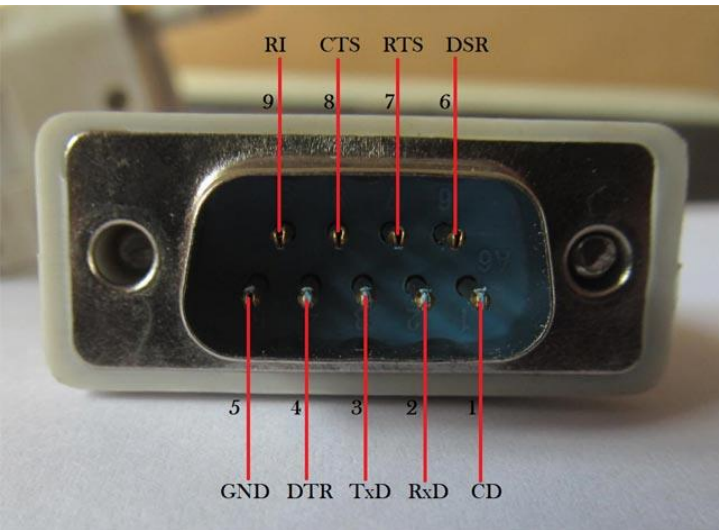# CPU and UART Configuration

# Connecting two UARTs
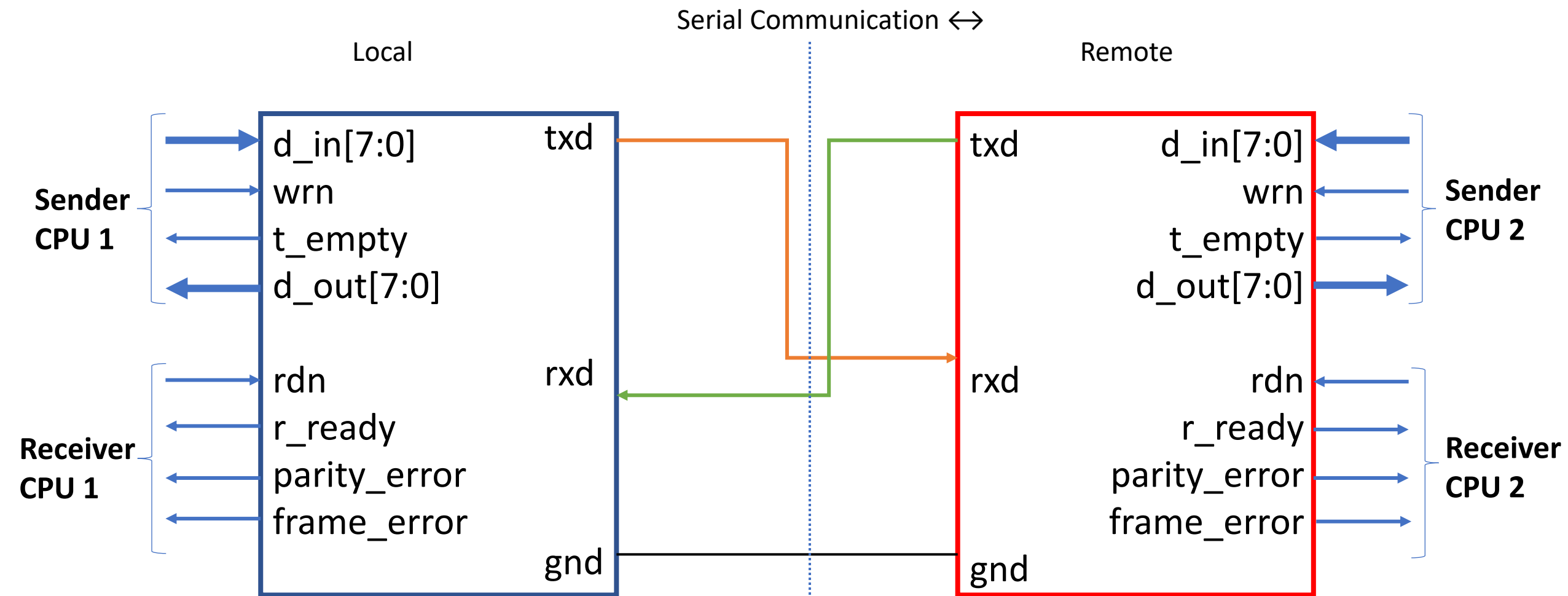
Physical Connections through RS 232
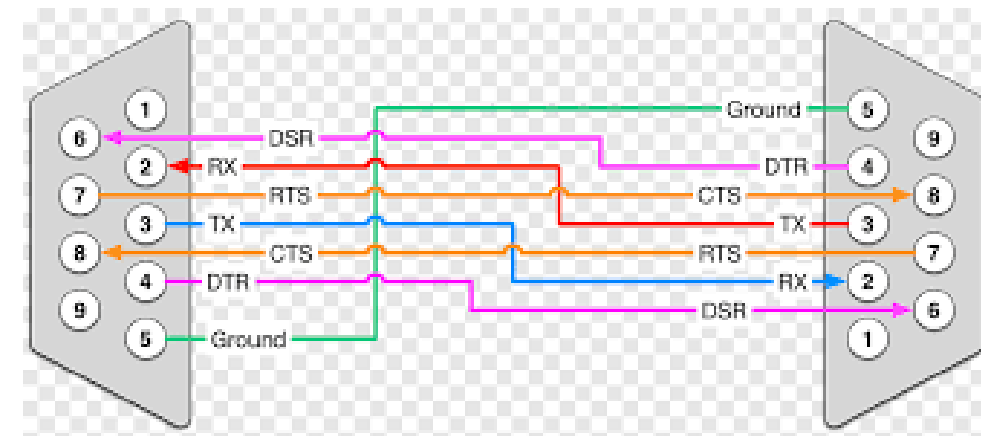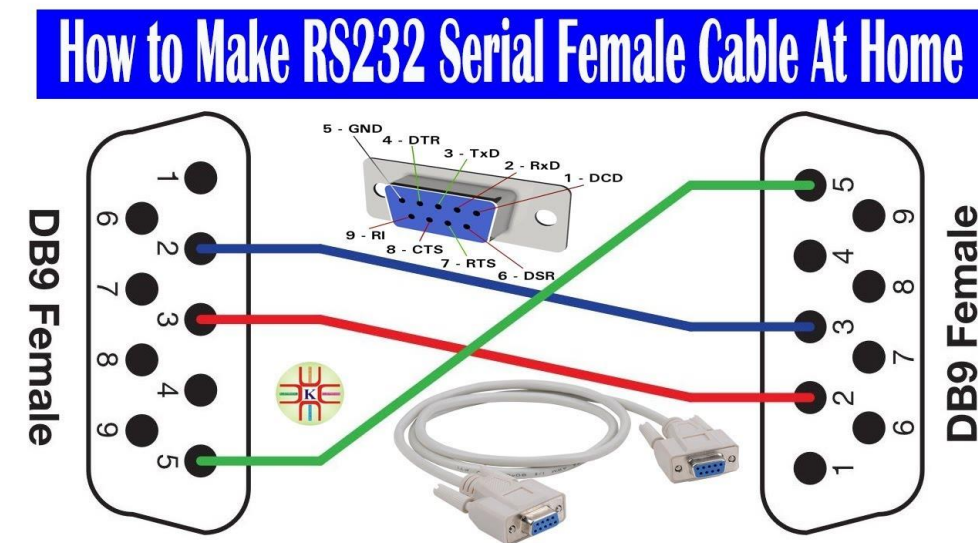
Typical UART Components

- **Baudrate clock generator**: Multiple of the bit rate to improve sampling in the middle of a bit period.
  - For generating this timing information, each UART uses an oscillator generating a frequency of about 1.8432 MHz (as an example).
  - This frequency is divided by 16 to generate the time base for communication.
  - Hence the maximum allowed communication speed is 115200 bps.
- **Input and output shift registers**: Each UART contains a shift register which is the fundamental method of conversion between serial and parallel forms. These registers shifts the data that has to be serially transmitted or serially received.
- **Transmit and receive control**: This control logic checks for the control signals from host processor to start or stop the transmission and reception of the data bits. In case of any error it also generates error signals.
- **Optional transmit and receive buffers:** Buffers can be used to hold the data temporarily.
- **Optional parallel data bus buffer**: This buffer improves the speed.
- **Optional FIFO:** The UART works by writing data from the host processor to its FIFO buffers and feeding the data from the buffer to the serial device in the format dictated by the user.

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UART Architecture Block Diagram
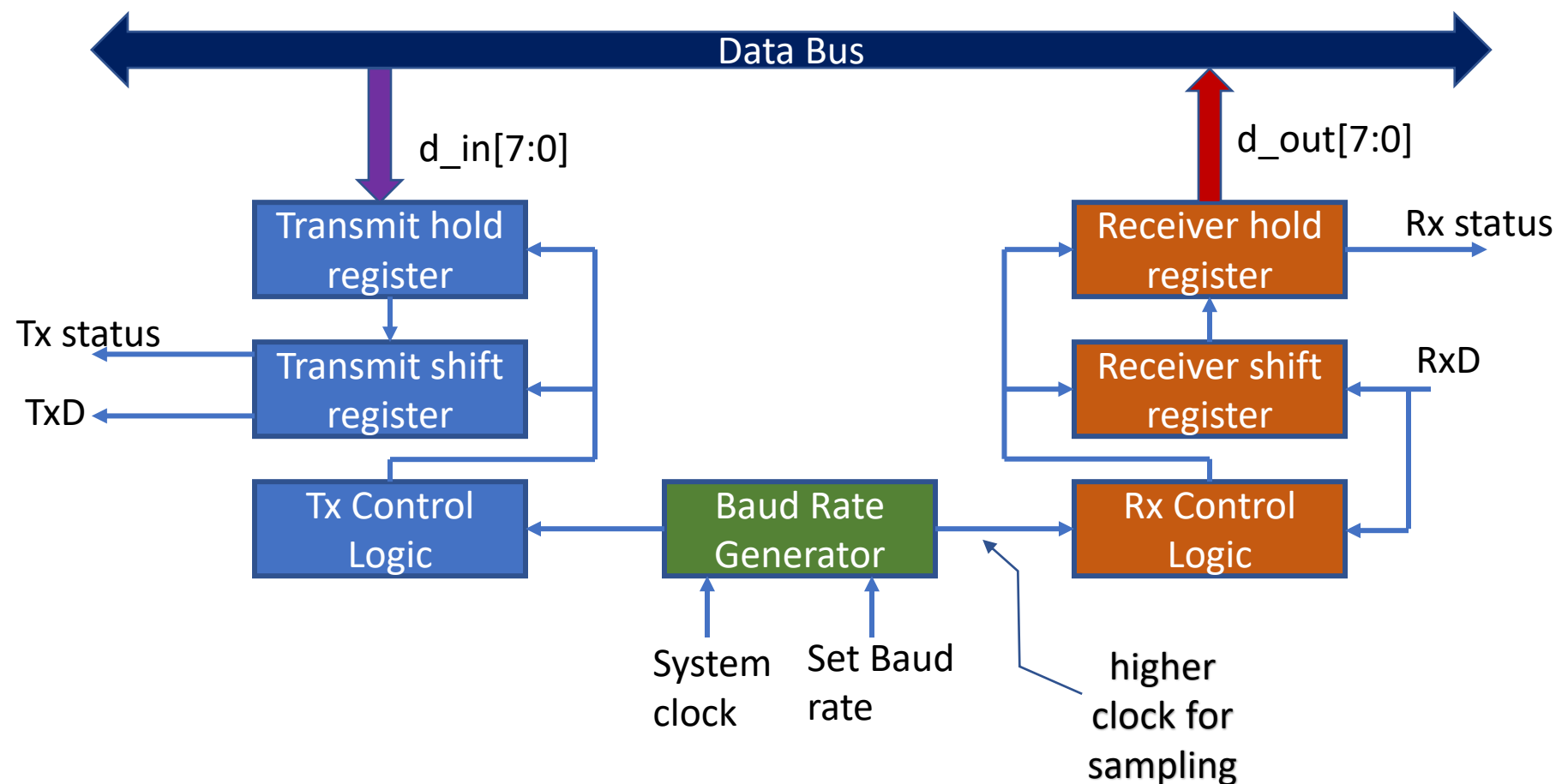
UART Block Diagram
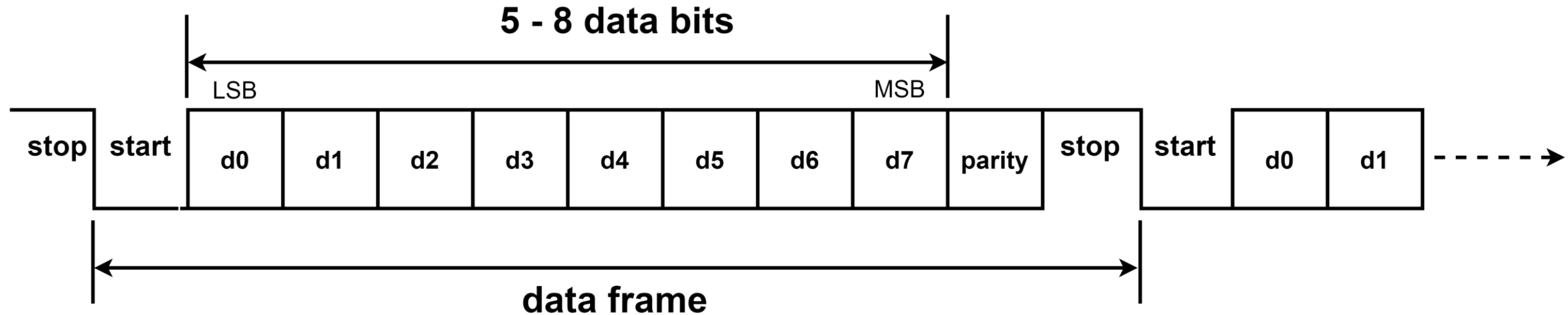
Process of Sending

- Asynchronous transmission does not send clock signal to transmit data to the receiver. The sender and receiver must agree on timing parameters in advance and special bits such as start and stop bits are added to each word which is used to synchronize the sending and receiving units.

- A bit called the "**Start Bit**" is added to the beginning of each word that is to be transmitted. The **Start Bit** indicates the start of the data transmission, and it alerts the receiver that a word of data is about to be sent.

- Upon reception of start bit the clock in the receiver goes into synchronization with the clock in the transmitter. The accuracy of these two clocks should not deviate more than 10% during the transmission of the remaining bits in the word.

- The individual bits of the word of data are sent after the start bit. Least Significant Bit (LSB) is sent first.

- The transmitter does not know when the receiver has read at the value of the bit. The transmitter begins transmitting the next bit of the word on next clock edge.

- Parity bit is be added when the entire data word has been sent. This bit can be used to detect errors at the receiver side. Then one Stop Bit is sent by the transmitter to indicate the end of the valid data bits.

UART Transmitter / System Handshaking

- System asserts Send and holds it high when it wants to send a byte
- UART asserts Busy signal in response
- When UART has finished transfer, UART de-asserts Busy signal
- System de-asserts Send signal
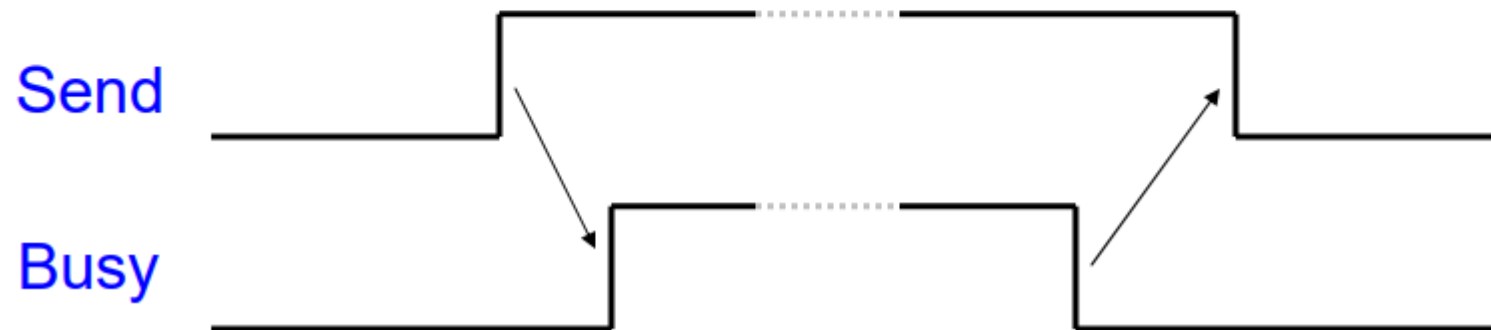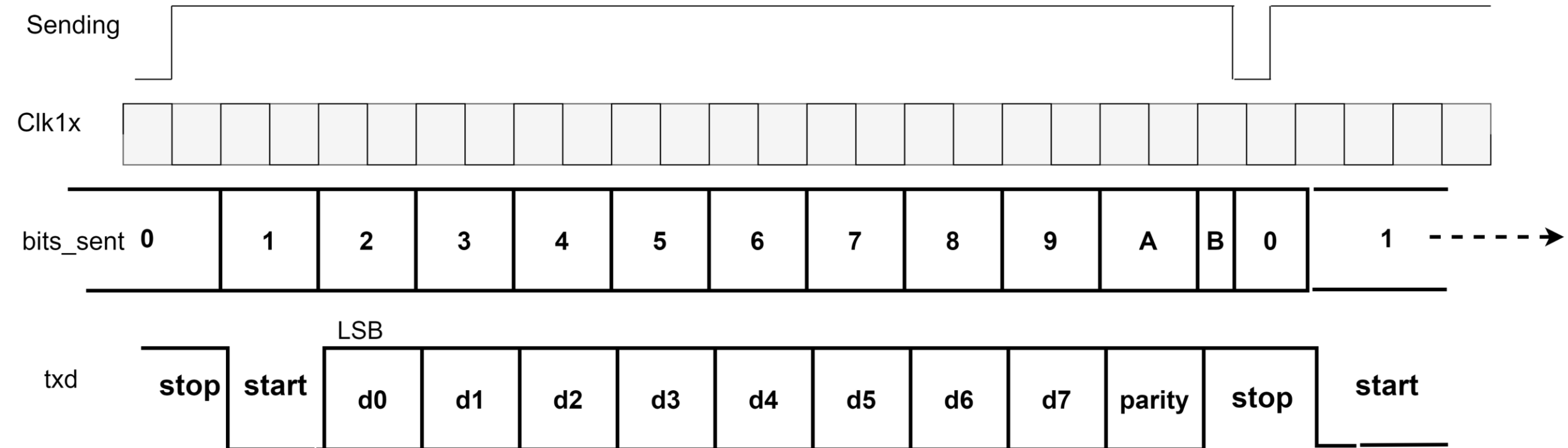
Data frame format of UART

LSB of data bits is transmitted first
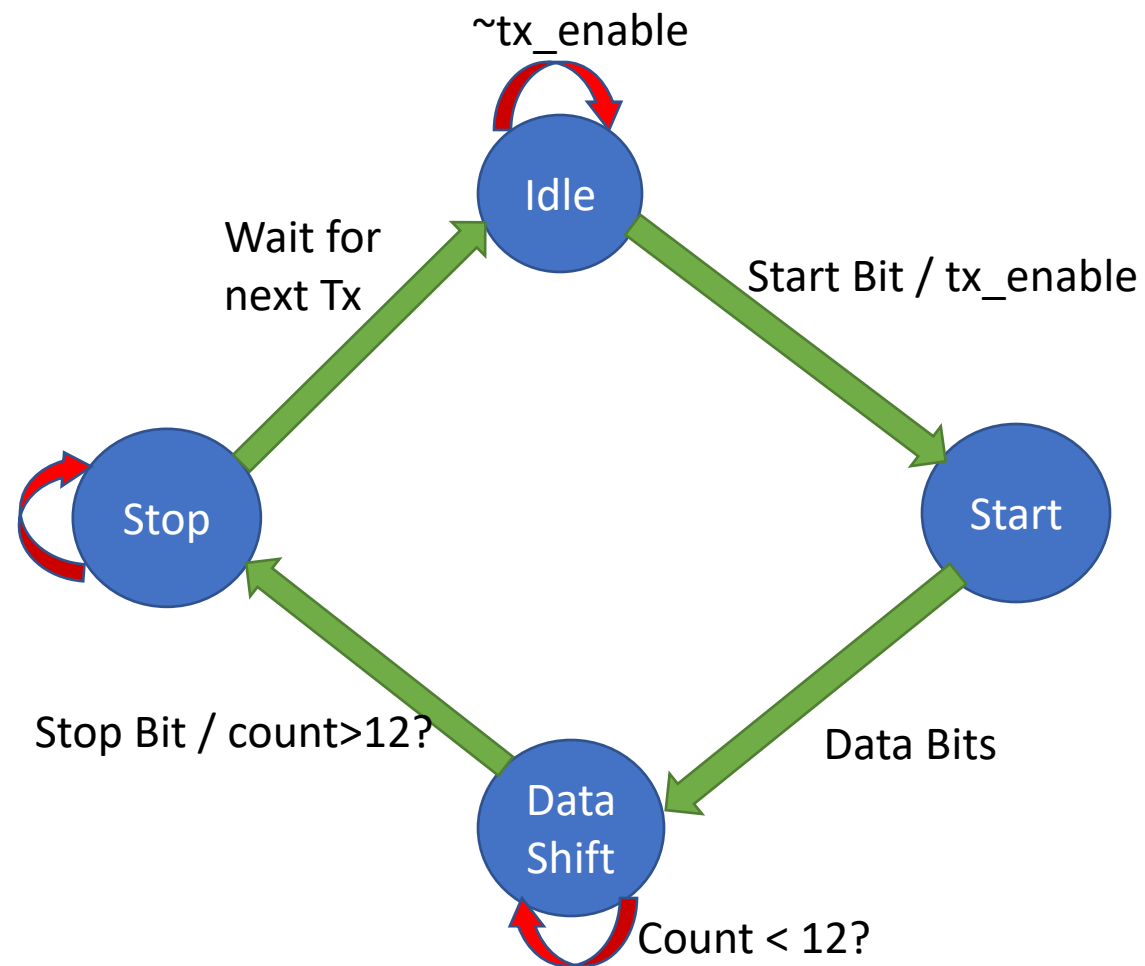
Start bit is logic low

Stop bit is logic high

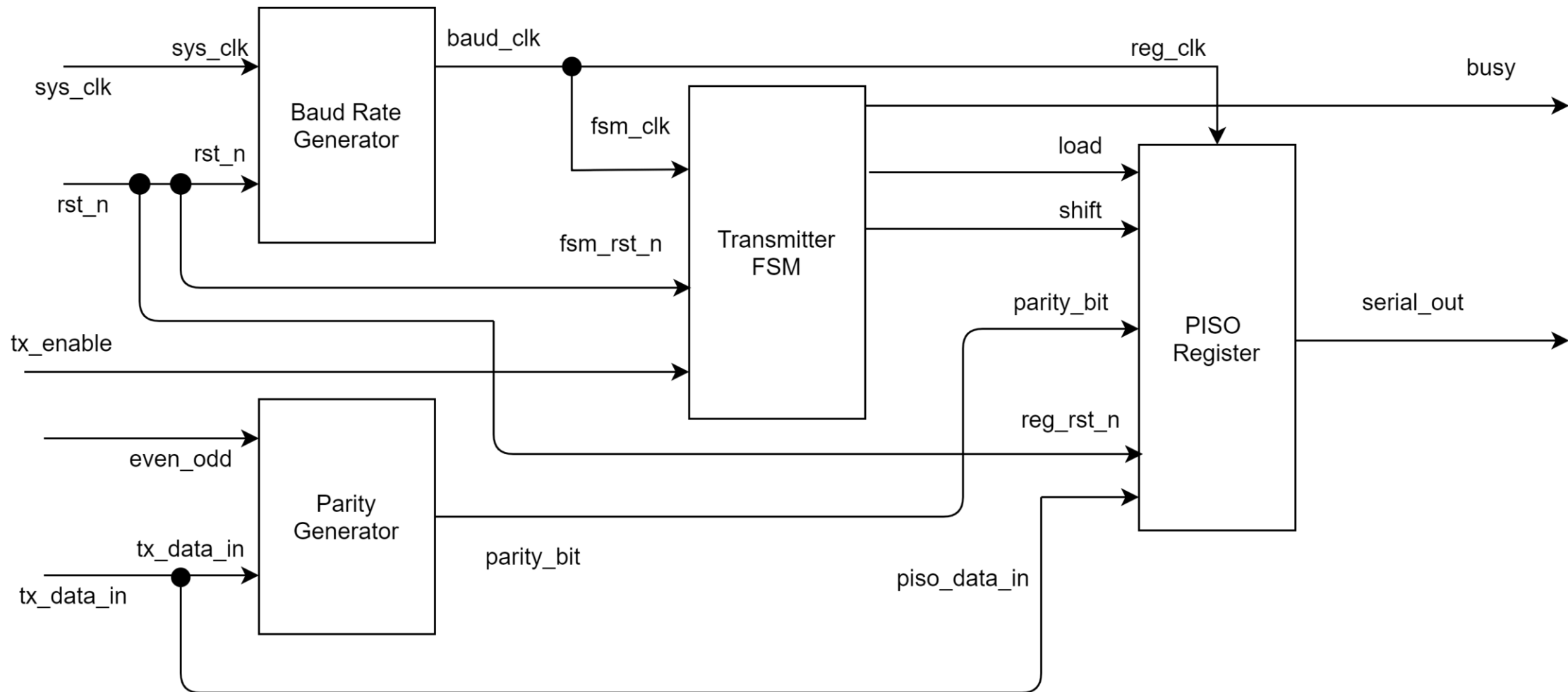In RESET Condition, Serial Data Transmission is 'High'

Sending clock of UART Transmitter

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FSM for UART Transmitter

UART Transmitter Architecture

Process of Receiving

- On the receiver side once it receives all of the bits in the data word, it can check for the Parity Bits. To accomplish this task both transmitter and receiver must agree on whether a Parity Bit is to be used.

- Then Stop Bit is encountered by receiver. A missing stop bit may result entire data to be garbage. This will cause a Framing Error and will be reported to the host processor when the data word is read. Framing Error can be caused due to mismatch of transmitter and receiver clocks.

- The UART automatically discards the Start, Parity and Stop bits irrespective of whether data is received correctly or not. If the sender and receiver are configured identically, these bits are not passed to the host. To transmit new word, the Start Bit for the new word is sent as soon as the Stop Bit for the previous word has been sent.

- The transmission speed in asynchronous communication is measured by Baud Rate. A Baud Rate represents the number of bits that are actually being sent over the media. The Baud rate includes the Start, Stop and Parity bits. The Bit rate (Bits per Second-bps) represents the amount of data that is actually sent from the transmitting device to the other device. Speeds for UARTs are in bits per second (bit/s or bps), although often incorrectly called the baud rate. Standard baud rates are: 110, 300, 1200, …..,115200,…..,1843200, and 2764800 bit/s.
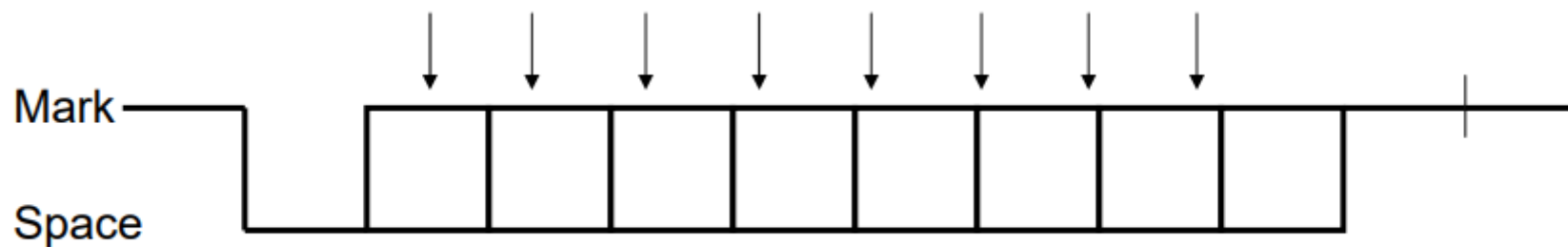
UART Character Reception

Start bit says a character is coming, receiver resets its timers

Receiver should sample in middle of bits

Mark

Space
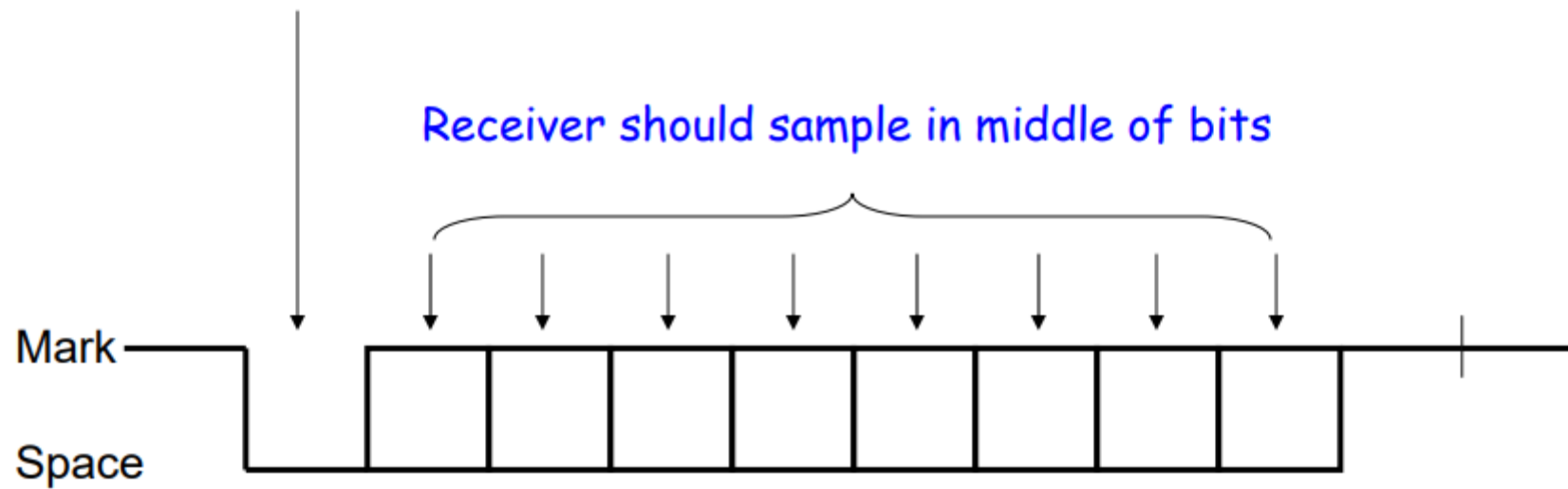
Receiver uses a timer (counter) to time when it samples. Transmission rate (i.e., bit width) must be known!
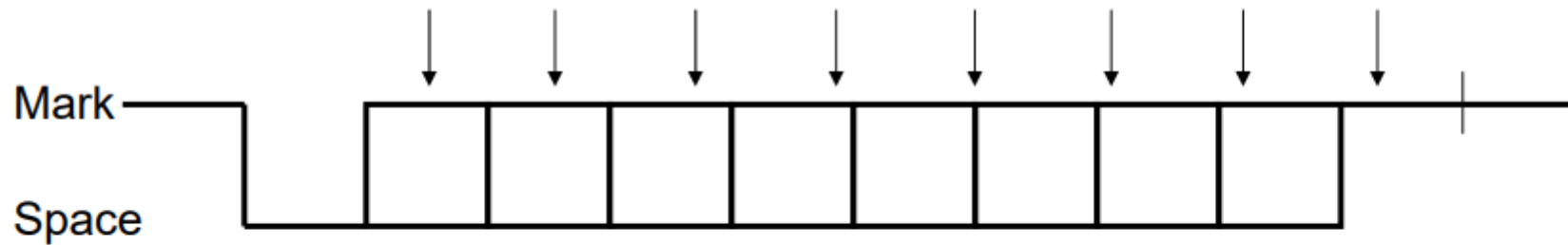
UART Receiver 'Quick' Sampling

If receiver samples too quickly, see what happens...

UART Receiver 'Slow' Sampling

If receiver samples too slowly, see what happens…

Mark
Space

Receiver resynchronizes on every start bit.
Only has to be accurate enough to read 9 bits.

UART Receiving Stop Bit

- Receiver also verifies that stop bit is '1'
  - If not, reports "framing error" to host system

- New start bit can appear immediately after stop bit
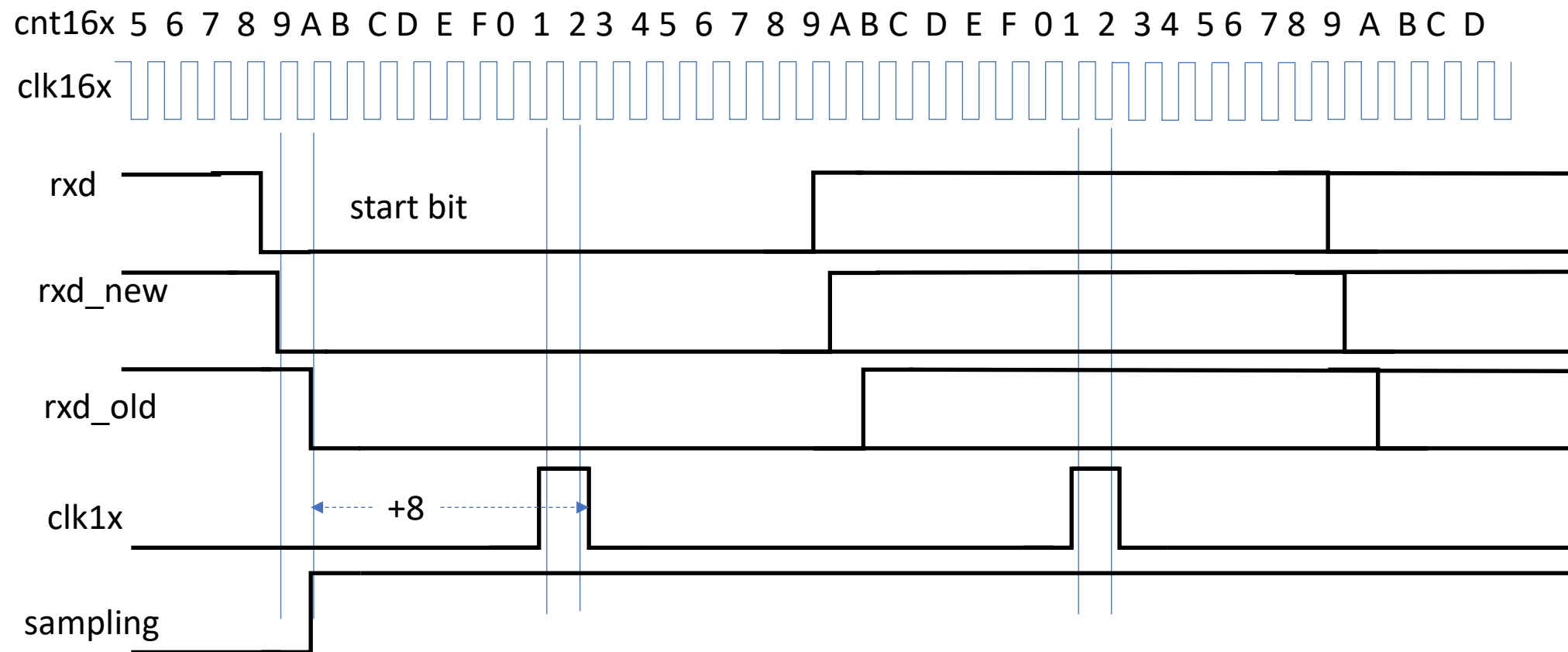 - Receiver will resynchronize on each start bit

UART Receiver Configuration

- UARTs usually have programmable options:
  - **Data:** 7 or 8 bits
  - **Parity:** even, odd, none, mark, space
  - **Stop bits:** 1, 1.5, 2
  - **Baud rate:** 300, 1200, 2400, 4800, 9600, 19.2K, 38.4k, 57.6k, 115.2k...

Sampling Clock of UART Receiver
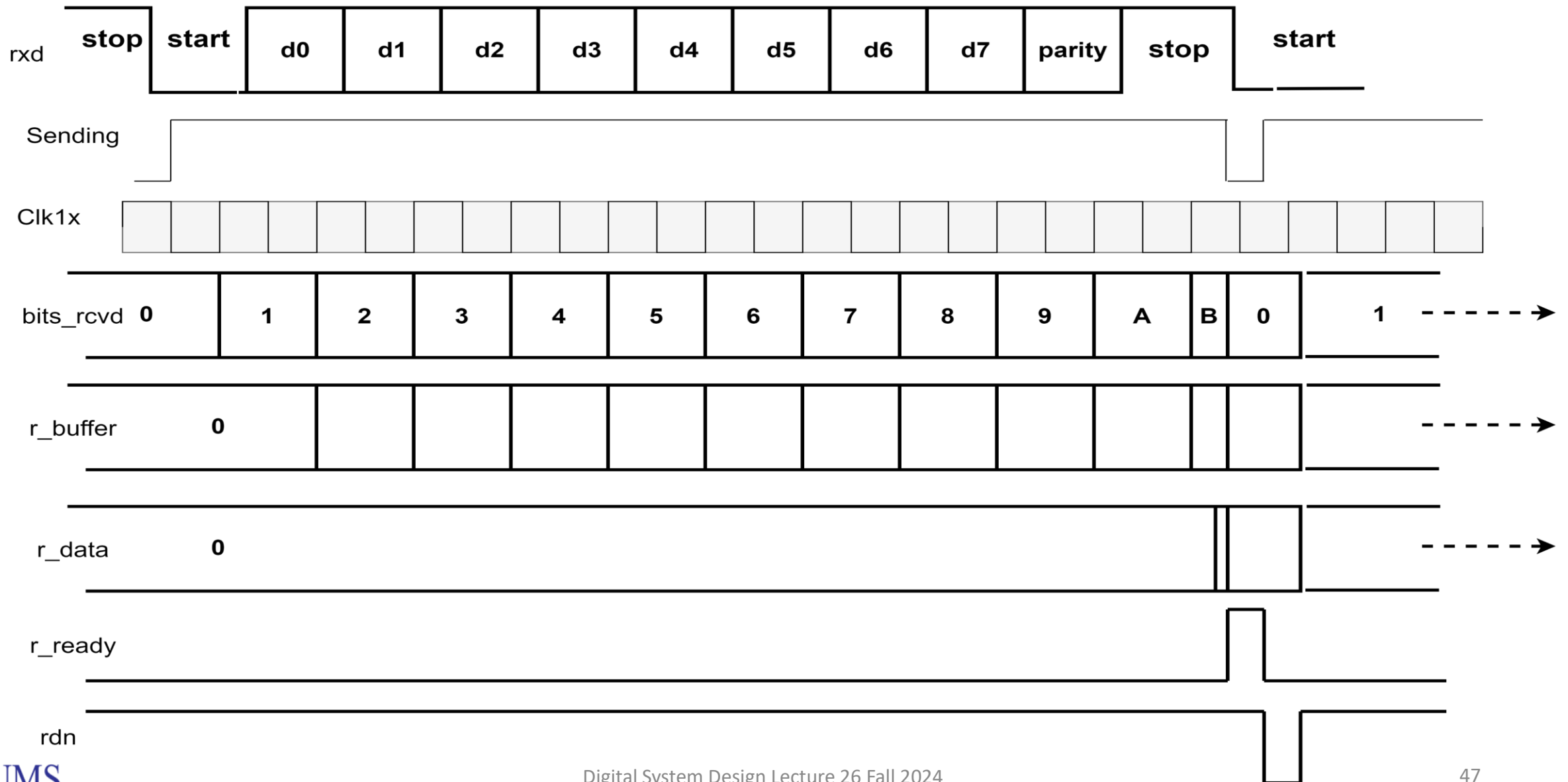
UART Receiver Timing Diagram

UART Receiver Starting details

- Clk16x is 16 times baud rate

- A 2-bit shift register saves input signal rxd

- When contents of this reg are 10 (rxd_old=1, rxd_new=0) then a falling edge is detected. We let sampling = 1 to sample rxd.

- A 4-bit counter cnt16x is used to generate a sampling pulse clk1x.

- During sampling=1, we let the rising edge of clk1x to appear at about centre place of one rxd bit; and the width of the clk1x pulse to be one clk16x clock cycle.

FSM for UART Receiver
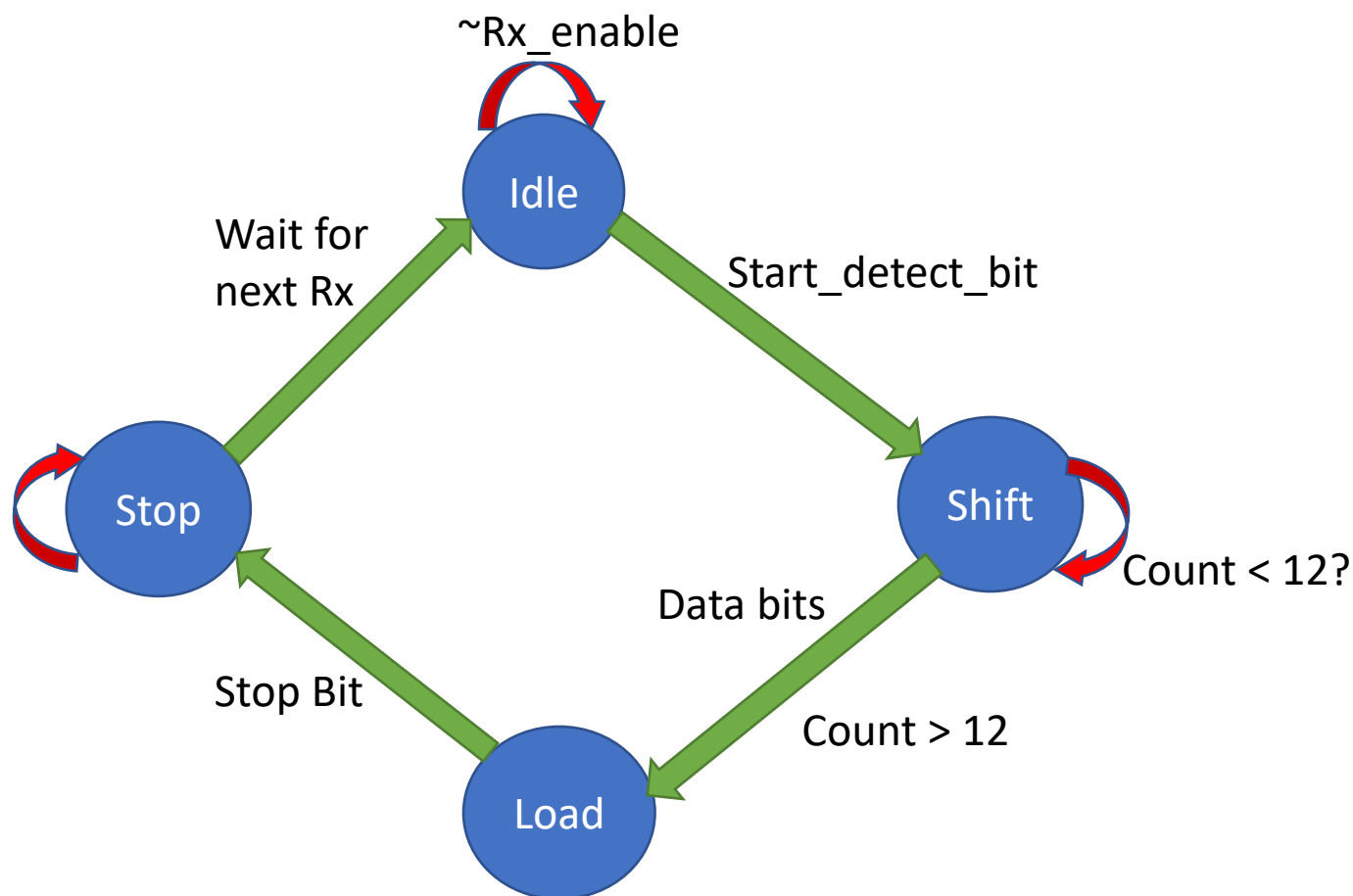
UART Receiver Architecture
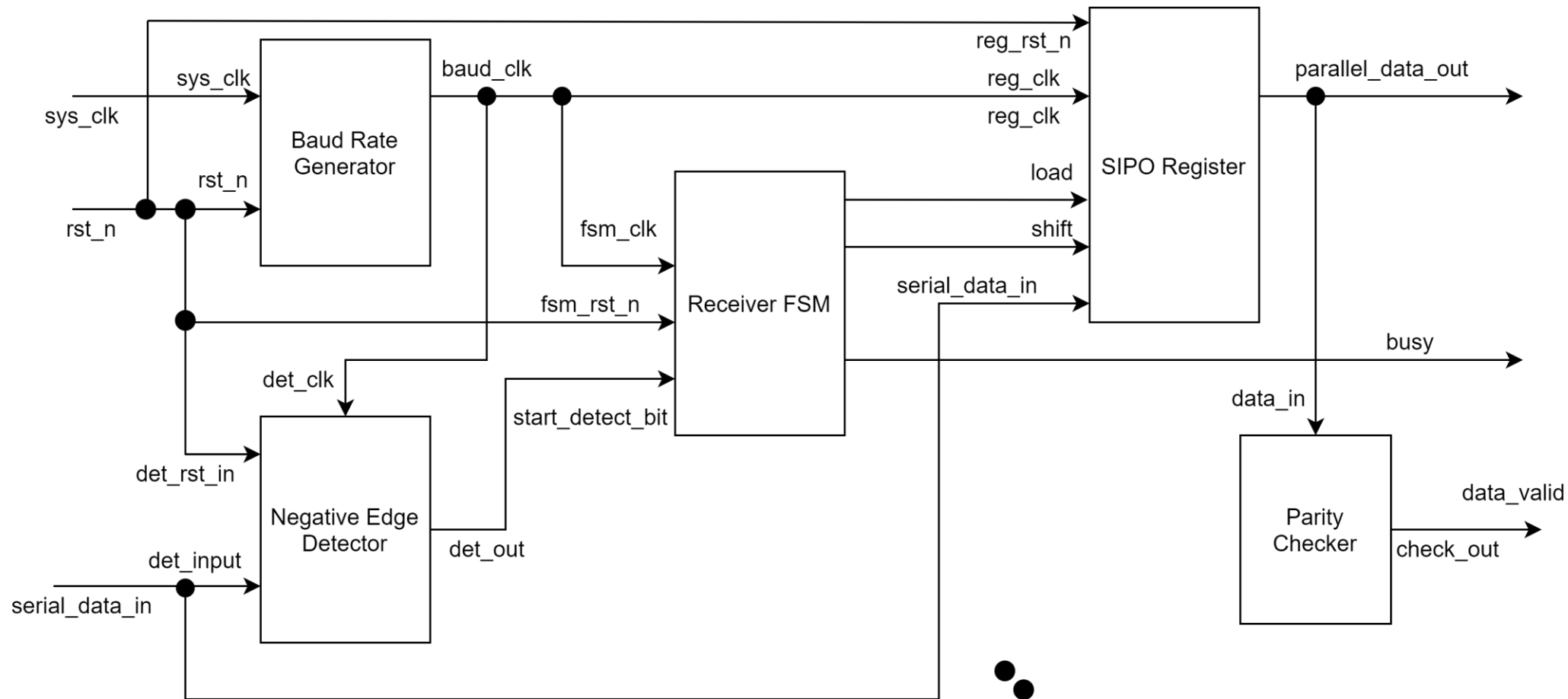
UART Errors

Overrun Error

An "**overrun error**" occurs when the UART cannot process the byte that just came in before the next one arrives. The host processor must service the UART in order to remove characters from the buffer. If the host processor does not service the UART and the buffer becomes full, then *Overrun Error* will occur.

Framing Error

A "**Framing Error**" occurs when the designated "start" and "stop" bits are not valid. Start bit acts as a reference for the remaining bits. When the "stop" bit is expected if the data line is not in the expected idle state a *Framing Error* will occur.

Parity Error

A "**Parity Error**" occurs when the number of "active" bits does not agree with the specified parity configuration of the UART.

Comm Example - Bluetooth

- Some details of Bluetooth protocol and its hardware design

Comm Example - Zibee

- Some details of Zigbee protocol and hardware design

Comm Example – Wifi - Ethernet

- May discuss hardware of Wifi, Ethernet