

CS / EE 320 Computer Organization and Assembly Language Spring 2024 Lecture 17

Shahid Masud

Topics: From Single Cycle MIPS to Multicyle MIPS, Pipelining

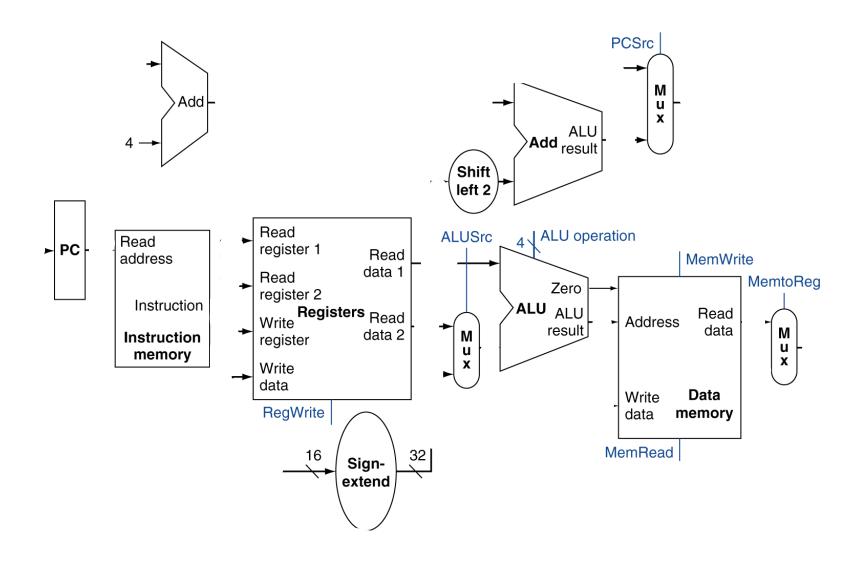
Topics

- Timing calculations in Single-Cycle Simple MIPS CPU
- How to add pipelining in simple MIPS CPU
- Detailed analysis of 5-Stage Pipelined MIPS CPU
- Calculations Throughput, Latency, Critical Path of 5-Stage MIPS CPU
- Design of individual Pipelined stages in a 5-Stage MIPS CPU
- Hazards in Pipelined MIPS CIPS
 - Structural Hazard (Hardware Related)
 - Data Hazard (Due to Assembly program)
 - Control Hazard (Architecture related)

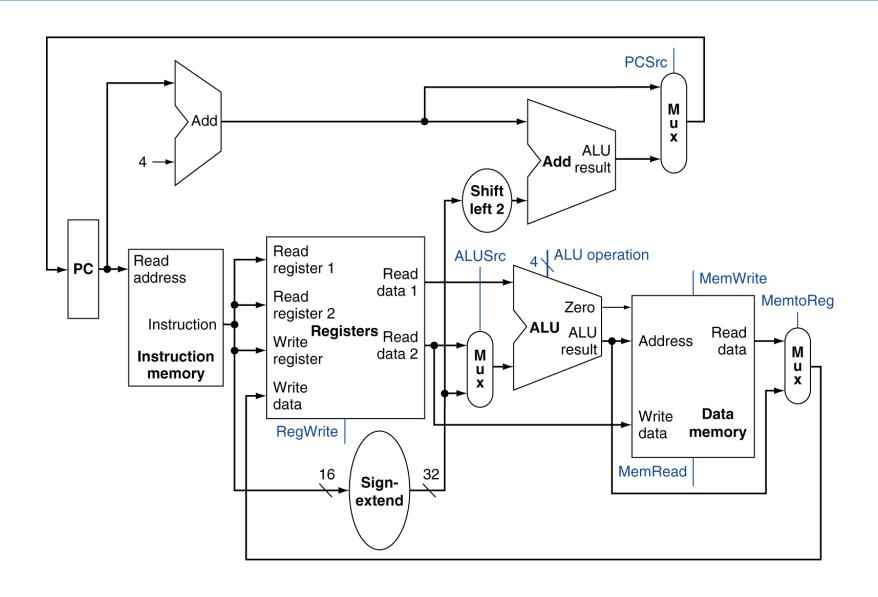


SIMPLE MIPS WITH NO PIPELINE

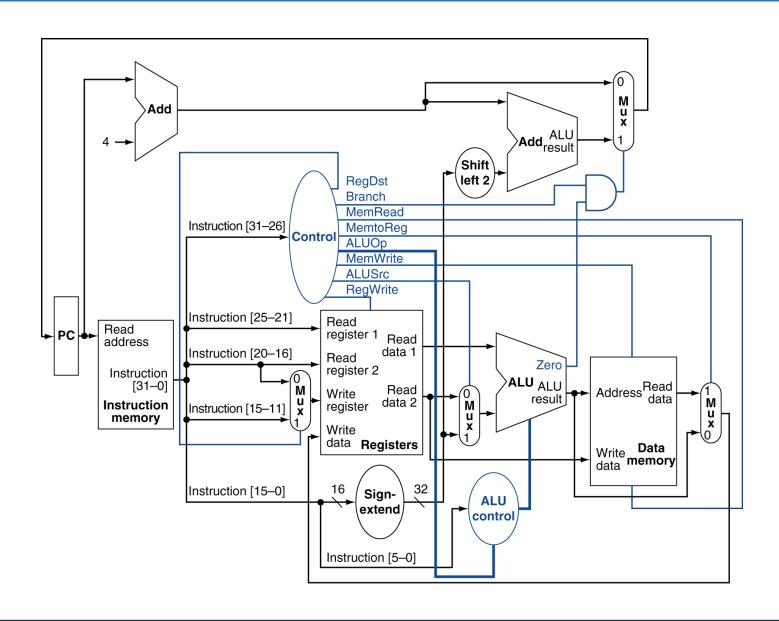
Full Datapath – for practice of connections



Full Datapath



MIPS Datapath With Control



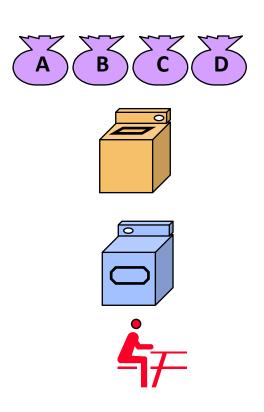


BASICS OF PIPELINING

What Is Pipelining

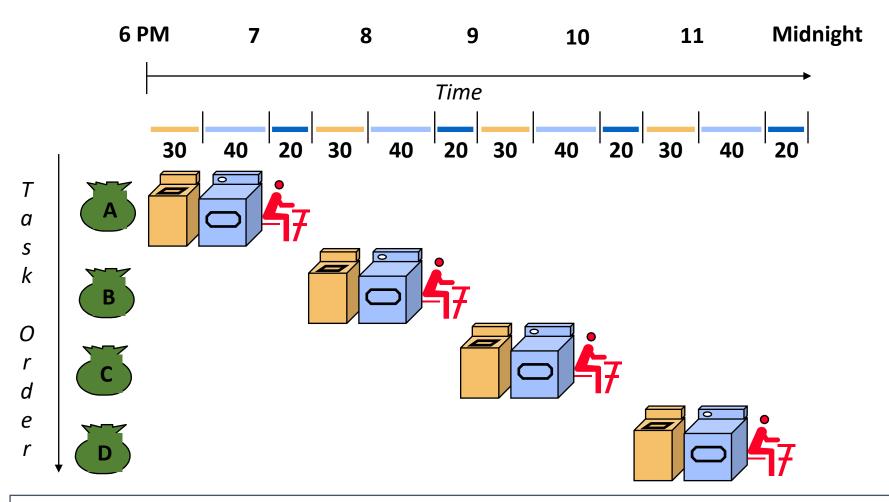


- Laundry Example
- Ann, Brian, Cathy, Dave each have one load of clothes to wash, dry, and fold
- Washer takes 30 minutes
- Dryer takes 40 minutes
- "Folding" takes 20 minutes
- One load takes 90 minutes



What Is Pipelining



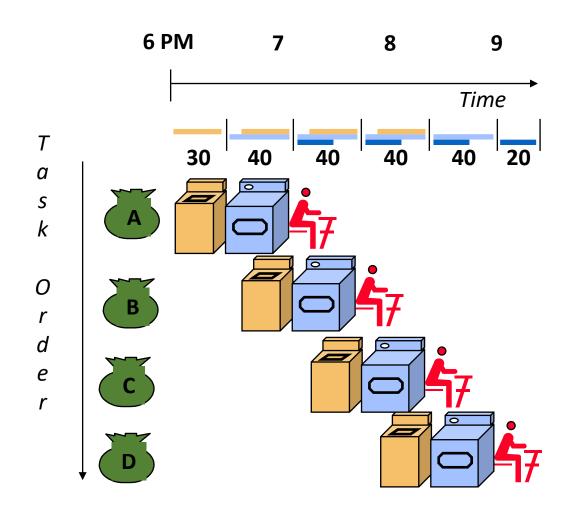


Sequential laundry takes 6 hours for 4 loads

If they learned pipelining, how long would laundry take?

Overlapping different operations





- Pipelining doesn't help latency of single task, it helps throughput of entire workload
- Pipeline rate limited by slowest pipeline stage
- Multiple tasks operating simultaneously
- Potential speedup = Number pipe stages
- Unbalanced lengths of pipe stages reduces speedup
- Time to "fill" pipeline and time to "drain" it reduces speedup



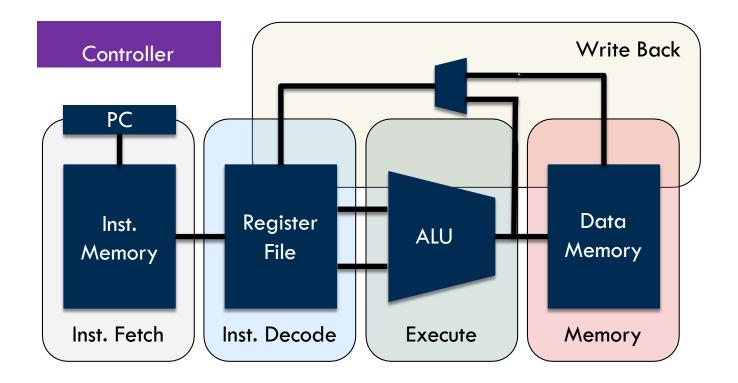
Introducing Pipelining for MIPS

Single-cycle MIPS RISC Architecture



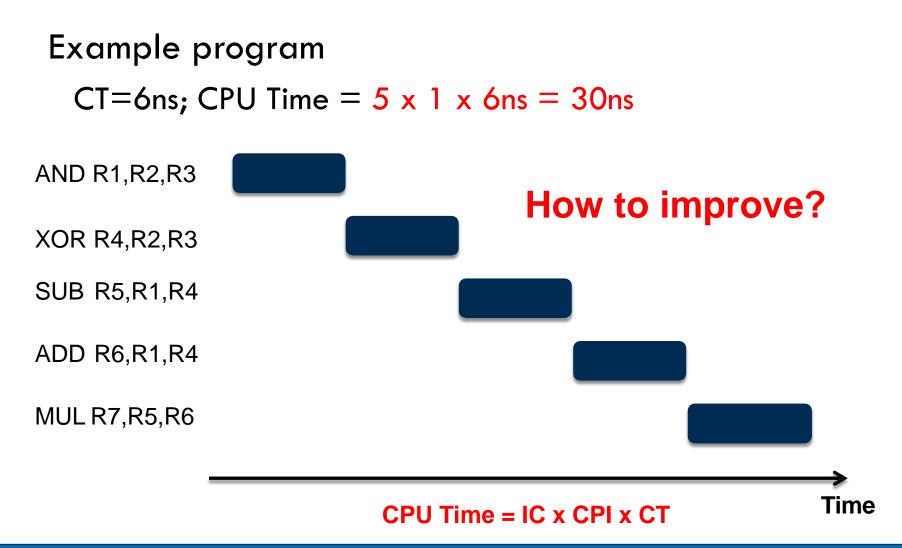
Example: simple MIPS architecture

Critical path includes all of the processing steps



Execution on Single-cycle MIPS Architecture





RISC Instruction Set

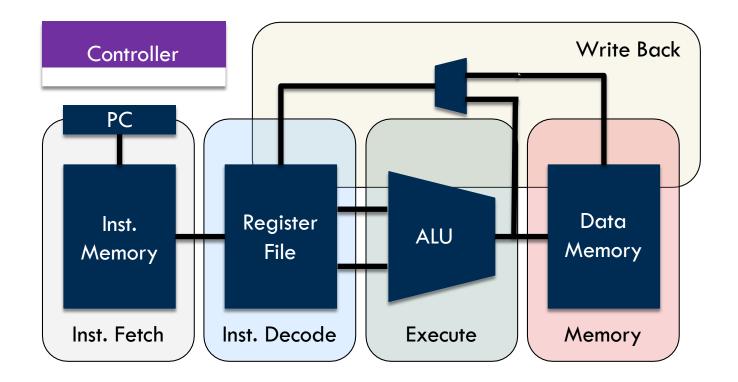


- Every instruction can be implemented in at most 5 clock cycles/ stages
 - ◆ Instruction fetch cycle (IF): send PC to memory, fetch the current instruction from memory, and update PC to the next sequential PC by adding 4 to the PC.
 - Instruction decode/register fetch cycle (ID): decode the instruction, read the registers corresponding to register source specifiers from the register file.
 - ◆ Execution/effective address cycle (EX): perform Memory address calculation for Load/Store, Register-Register ALU instruction and Register-Immediate ALU instruction.
 - Memory access (MEM): Perform memory access for load/store instructions.
 - Write-back cycle (WB): Write back results to the dest operands for Register-Register ALU instruction or Load instruction.

Re-Using Idle Resources

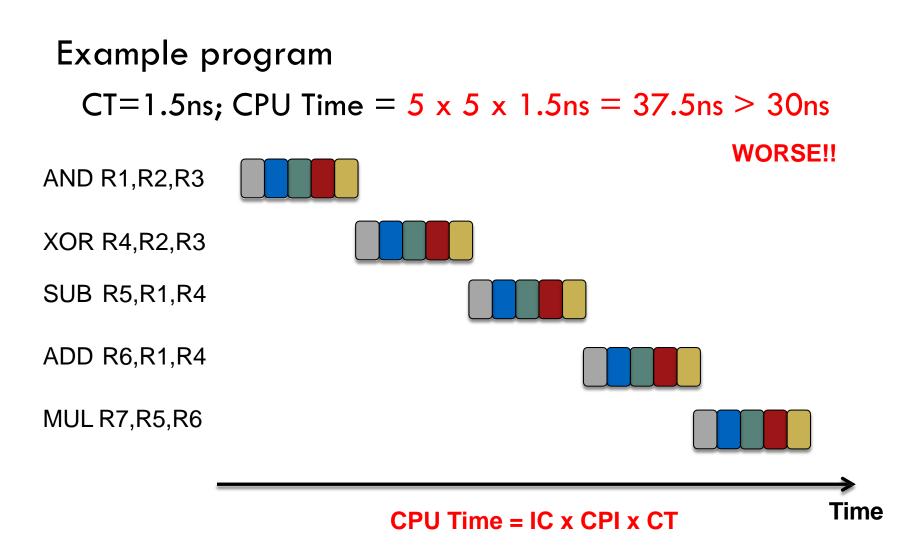


Each processing step finishes in a fraction of a cycle ldle resources can be reused for processing next instruction



Non Pipeline Architecture







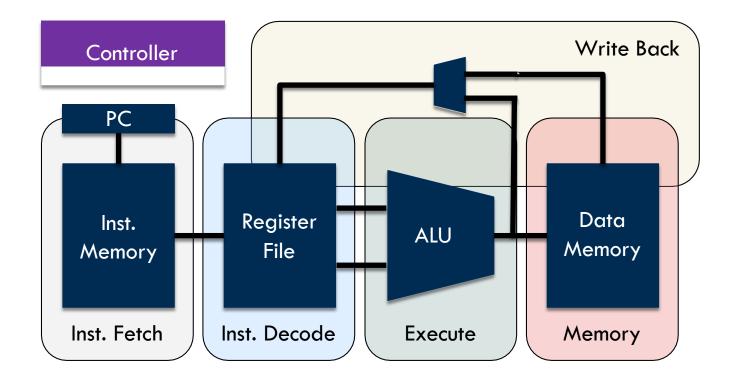
Developing a Multi-cycle Data Path

5

Develop Simple Five Stage MIPS Pipeline



A pipelined load-store architecture that processes up to one instruction per cycle

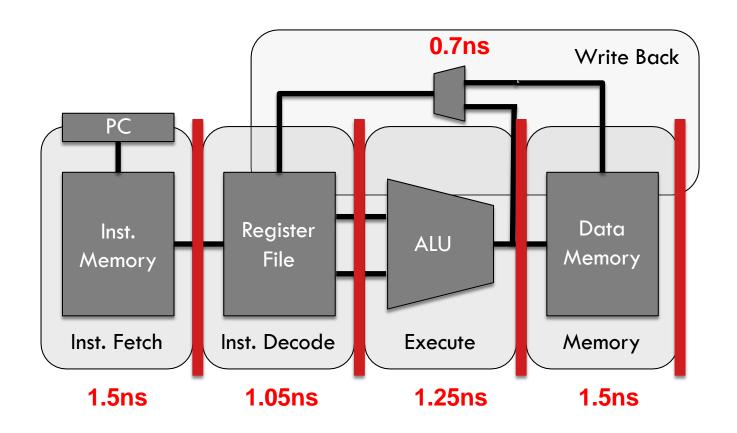


Identify Distinct Processing Stages in Architecture



Five stage pipeline

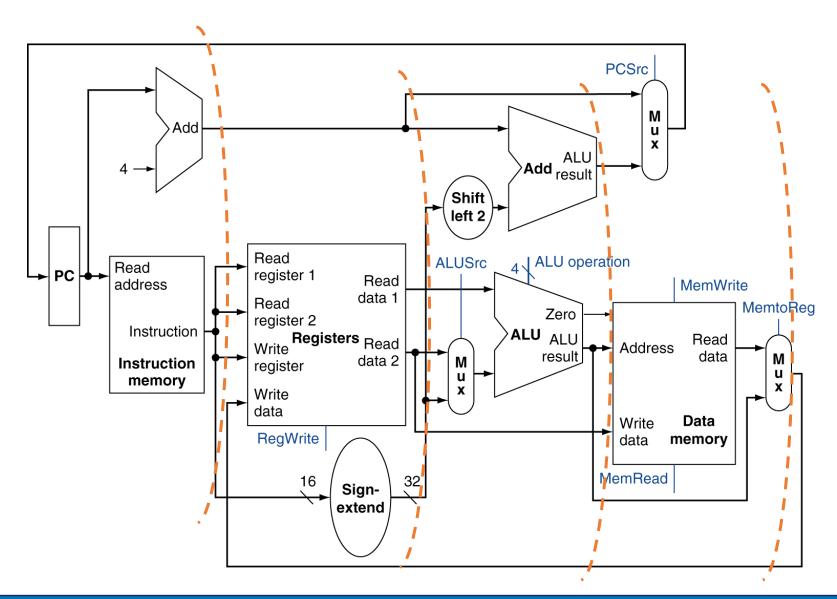
Critical path determines the cycle time





Identify CPU Operations where Pipeline Registers could be Inserted

Recap - Full Datapath



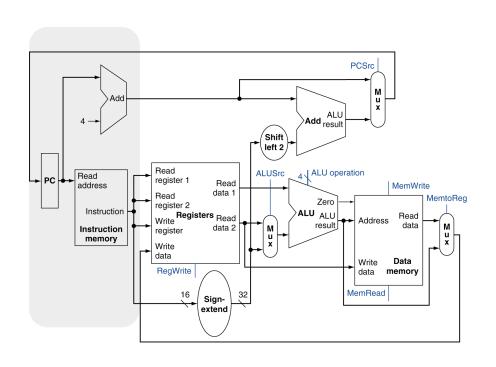
Instruction Fetch

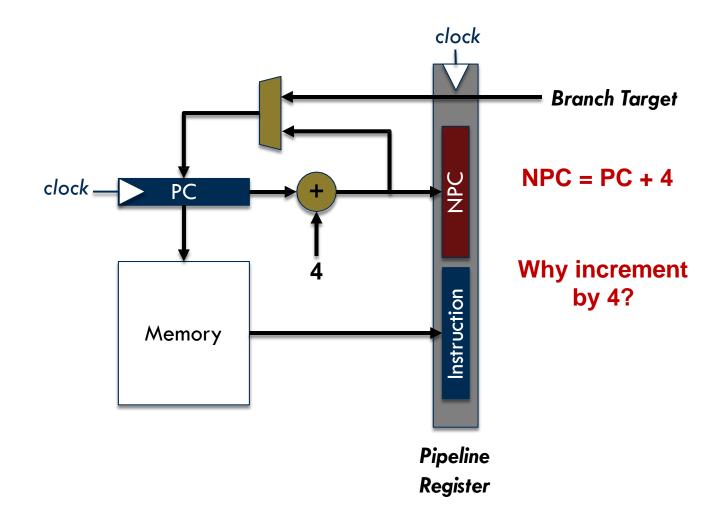


- Read an instruction from memory (I-Memory)
 Use the program counter (PC) to index into the I- Memory
 Compute New PC (NPC) by incrementing current PC
 - What about branches?
 - (later?)
- Update pipeline registersWrite the instruction into the pipeline registers

Instruction Fetch







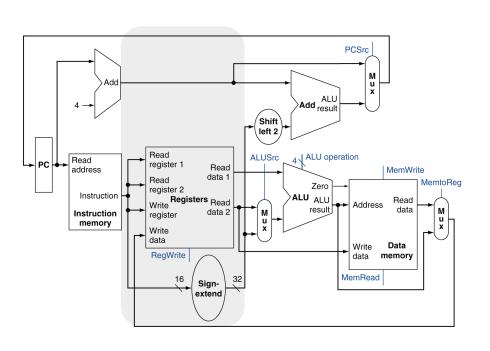
Instruction Decode

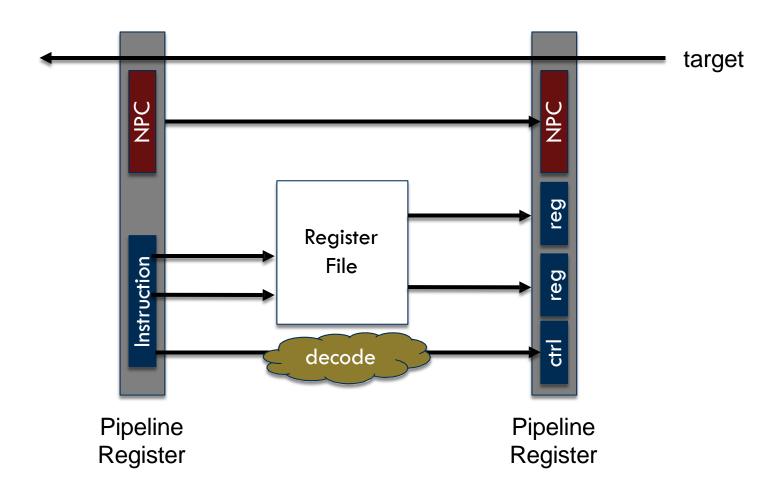


- Generate control signals for the opcode bits
- Read source operands from the register file (RF)
 Use the specifiers for indexing RF
 - How many read ports are required?
- Update pipeline registers
 Send the operand and immediate values to next stage
 Pass control signals and NPC to next stage

Instruction Decode







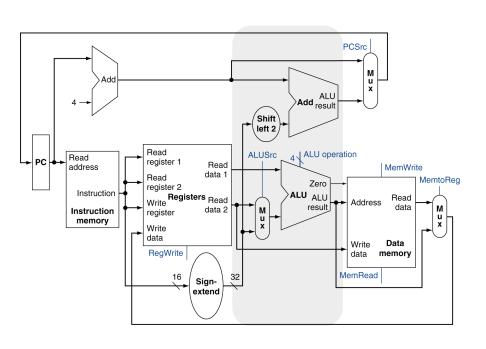
Execute Stage

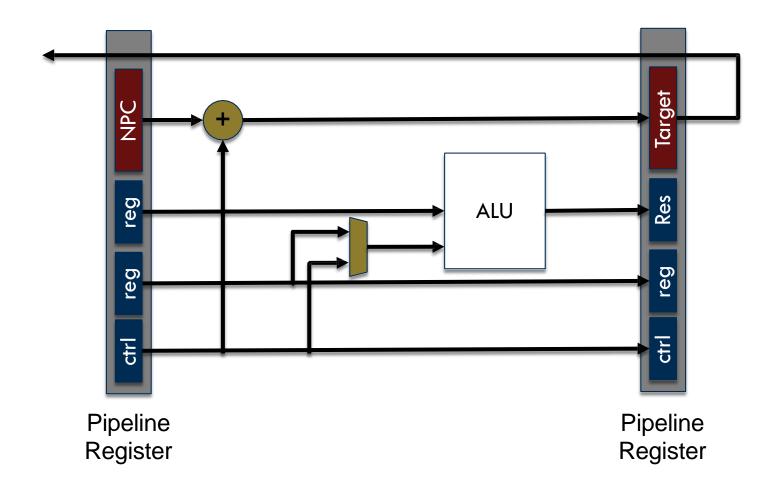


- Perform ALU operation
 - Compute the result of ALU
 - Operation type: control signals
 - First operand: contents of a register
 - Second operand: either a register or the immediate value
 - Compute branch target
 - Target = NPC + immediate
- Update pipeline registers
 - Control signals, branch target, ALU results, and destination

Execute Stage







Memory Access



Access data memory

Load/store address: ALU outcome

Control signals determine read or write access

Update pipeline registers

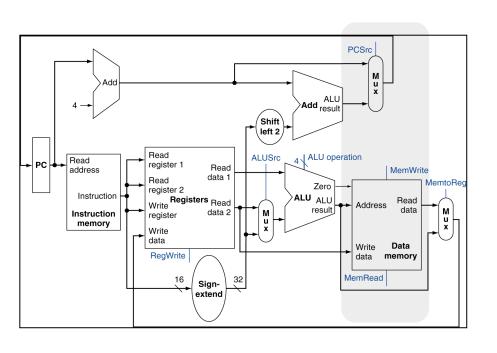
ALU results from execute

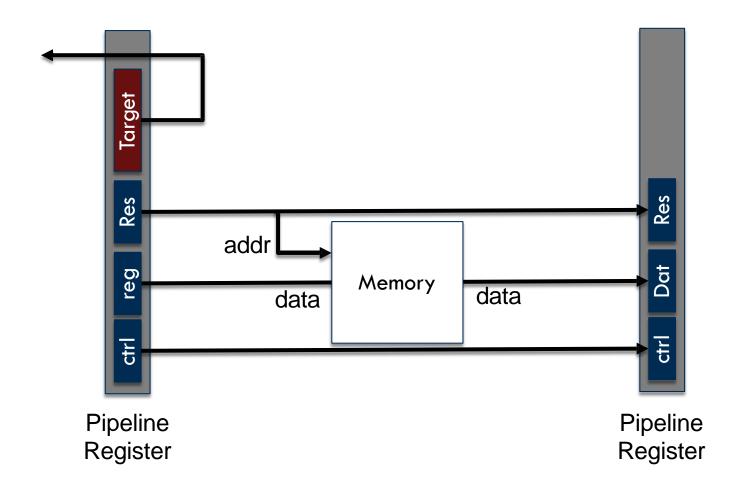
Loaded data from D-Memory

Destination register

Memory Access







Register Write Back



Update register file

Control signals determine if a register write is needed

Only one write port is required

- Write the ALU result to the destination register, or
- Write the loaded data into the register file



Improving Resource Utilization

Improving CPU Resource Utilization

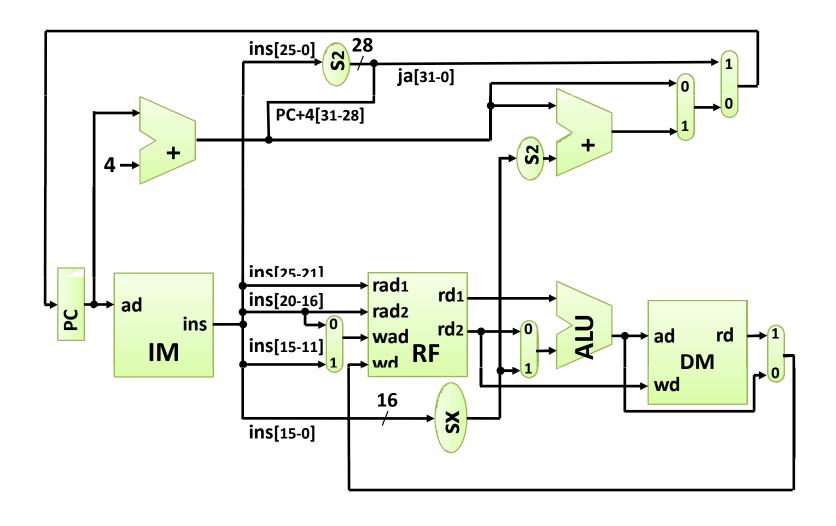


Efficient MIPS CPU Design

- Can we eliminate two adders?
- How to share (or reuse) a resource (say ALU) in different clock cycles?
- Store results in registers.
- Of course, more multiplexing may be required!
- Resources in this design: RF, ALU, MEM.

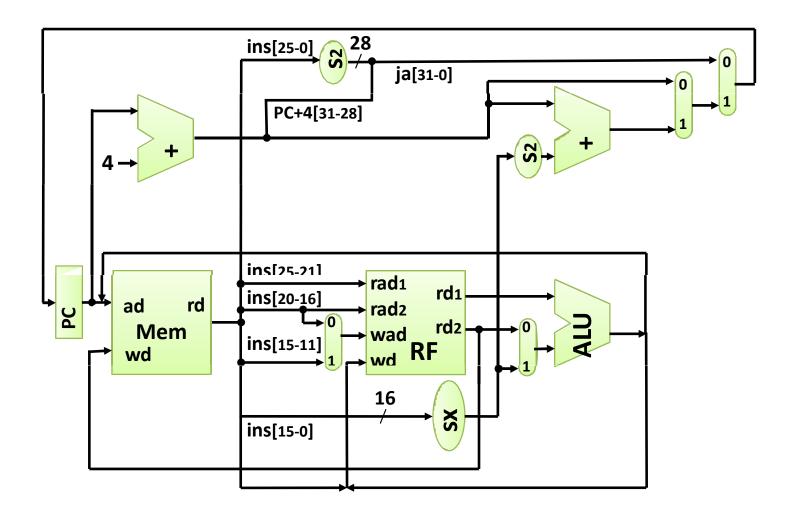
Single Cycle Datapath





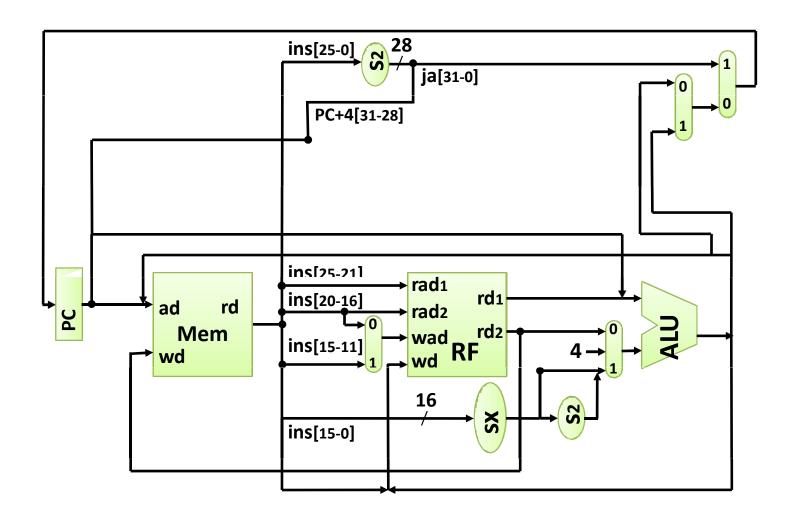
Merge IM and DM, Rearrange Diagram





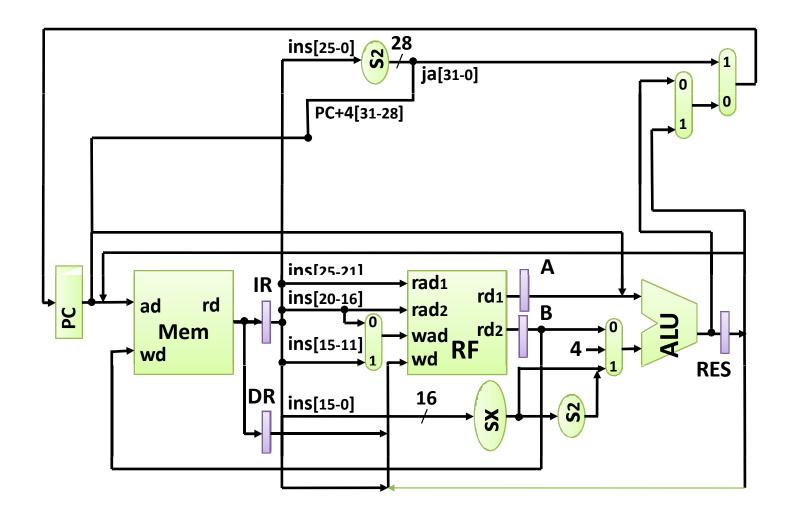
Combine Adders - Rearrange Diagram





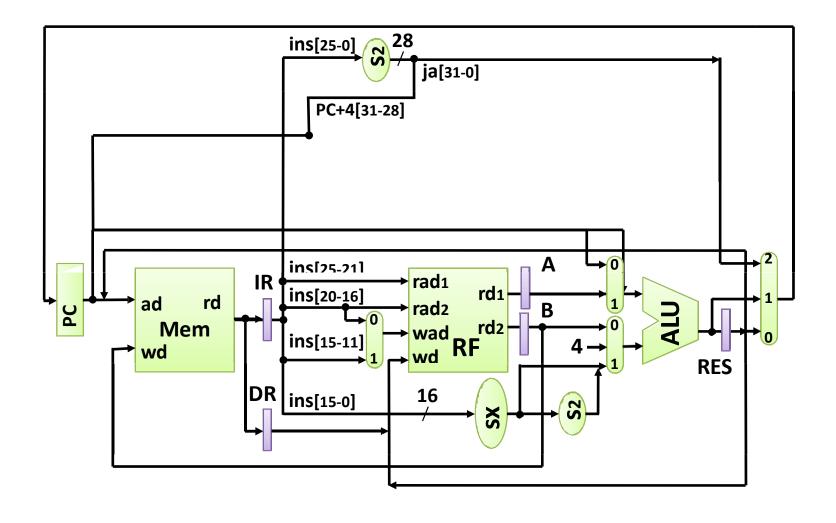
Introduce registers





Introduce ALU multiplexers at Input and Output

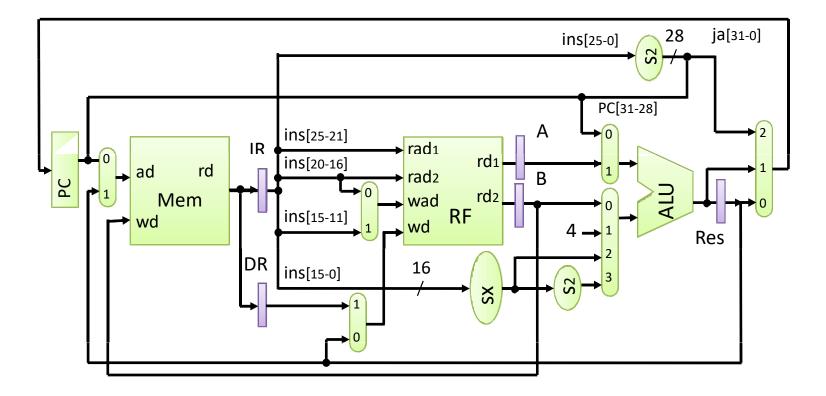




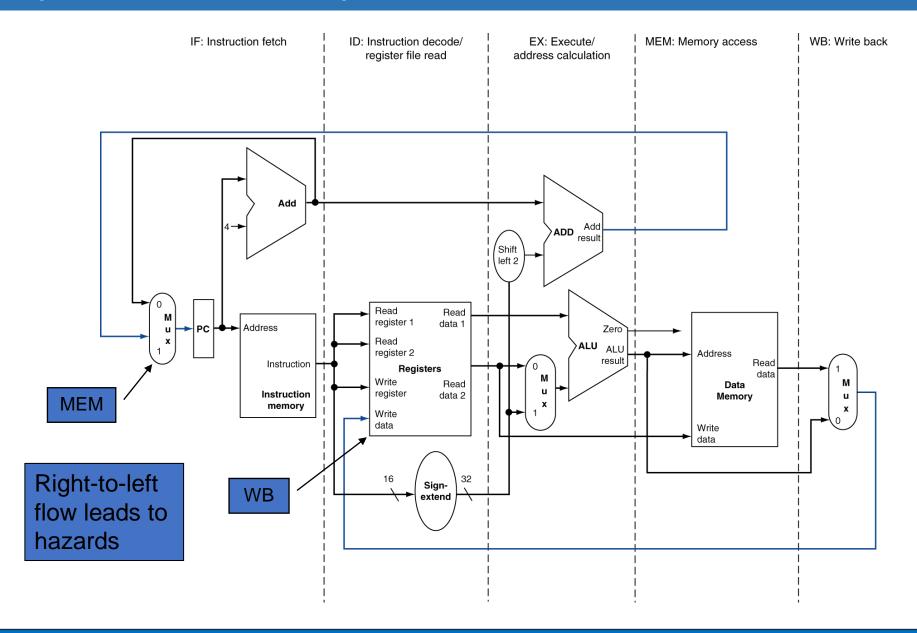
Final: Efficient Multi-Cycle Data Path



- Mux at Memory Input
- Mux at Register File Input
- One memory for lw / sw
- One ALU for EXE and Address Calculation
- Several MUXes introduced

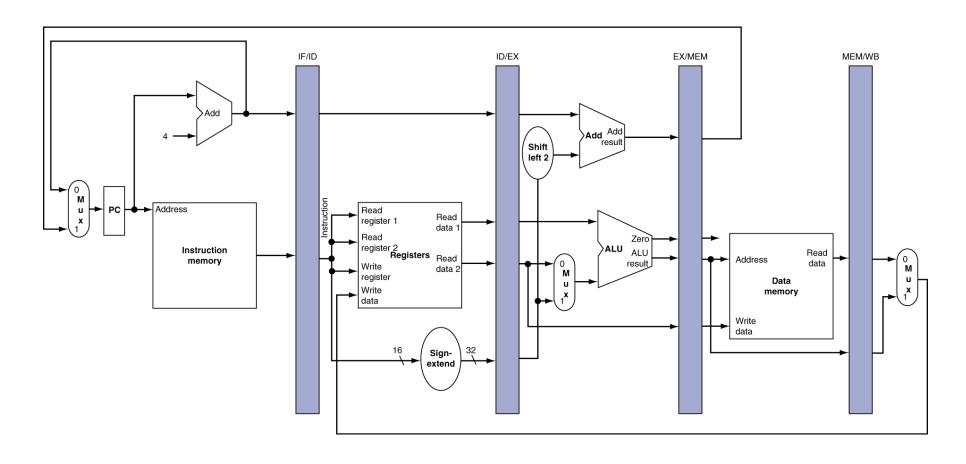


MIPS Pipelined Datapath



Pipeline registers

- Need registers between stages
 - To hold information produced in previous cycle

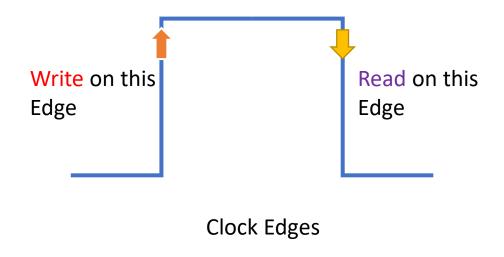


Pipeline Operation

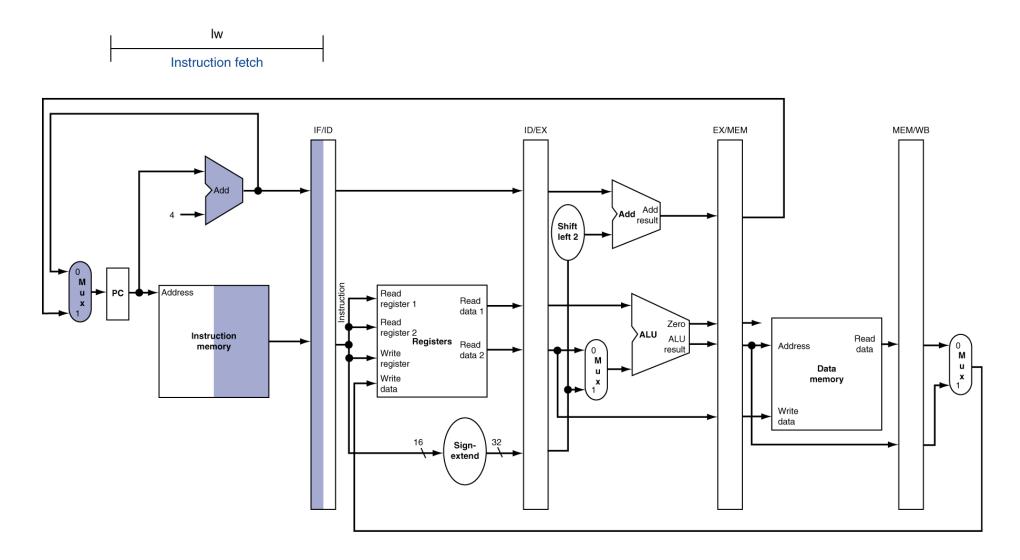
- Cycle-by-cycle flow of instructions through the pipelined datapath
 - "Single-clock-cycle" pipeline diagram
 - Shows pipeline usage in a single cycle
 - Highlight resources used
 - c.f. "multi-clock-cycle" diagram
 - Graph of operation over time
- We'll look at "single-clock-cycle" diagrams for load & store

RF is two operations in One Clock Cycle

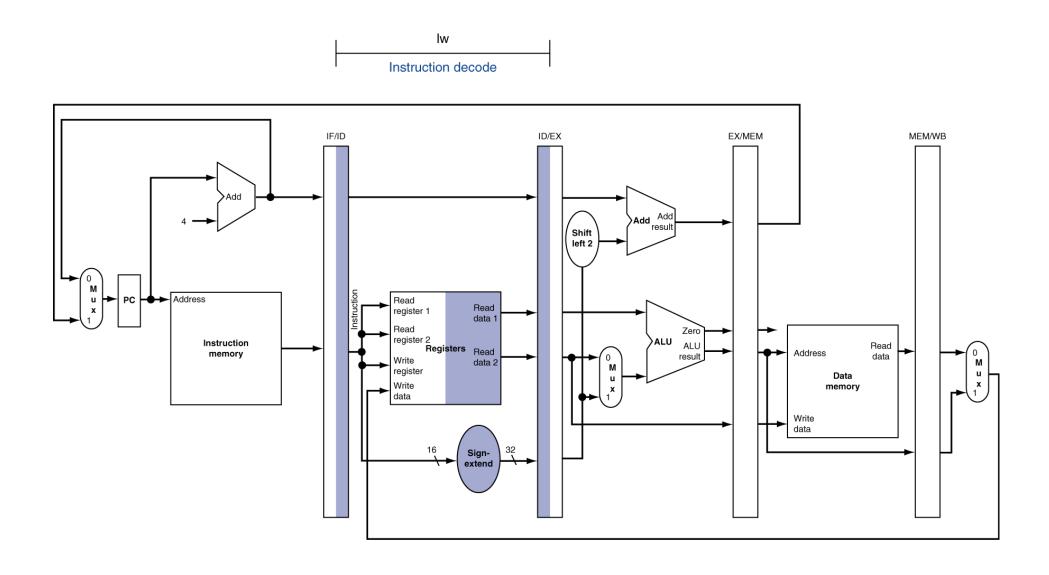
- Assume Write is when clock is going up
- Assume Read is when Clock is going down
 (you can assume otherwise too, but two operations in one clock cycle)



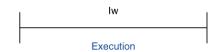
IF for Load, Store, ...

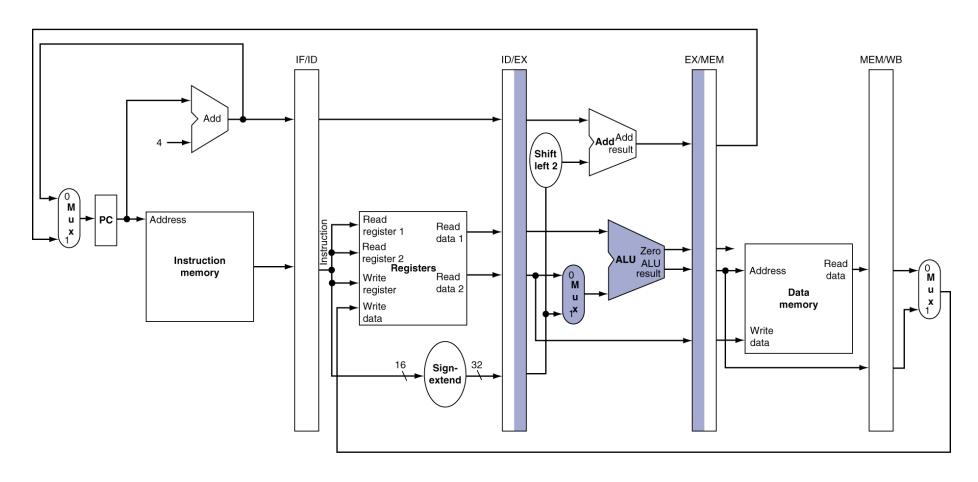


ID for Load, Store, ...

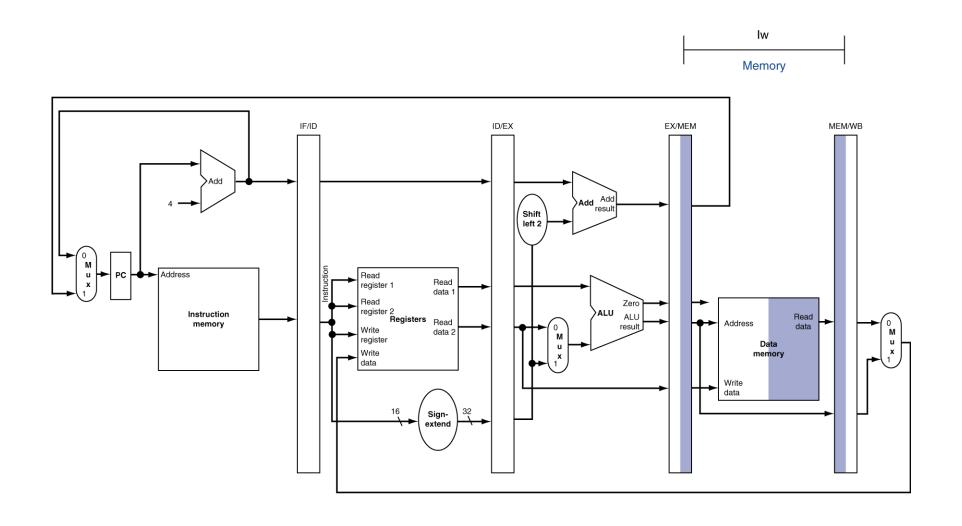


EX for Load

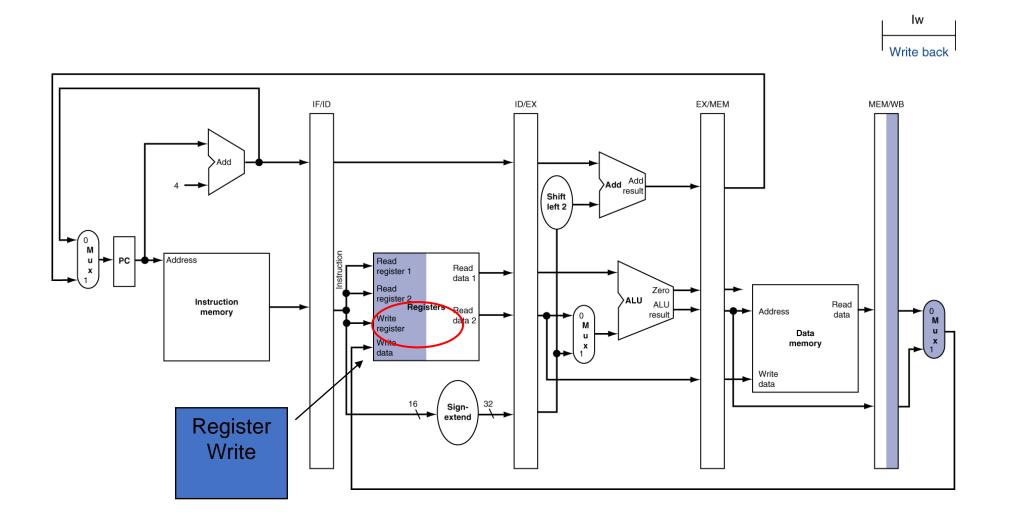




MEM for Load

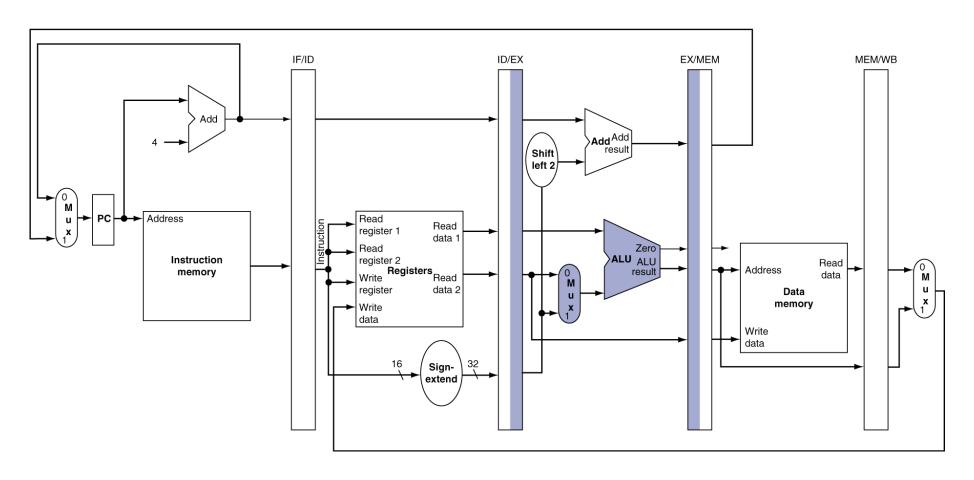


WB for Load

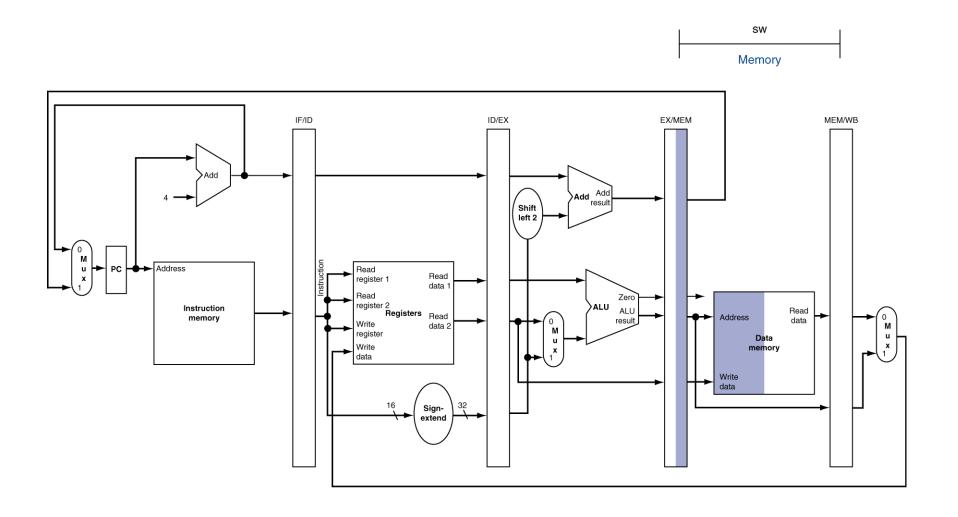


EX for Store

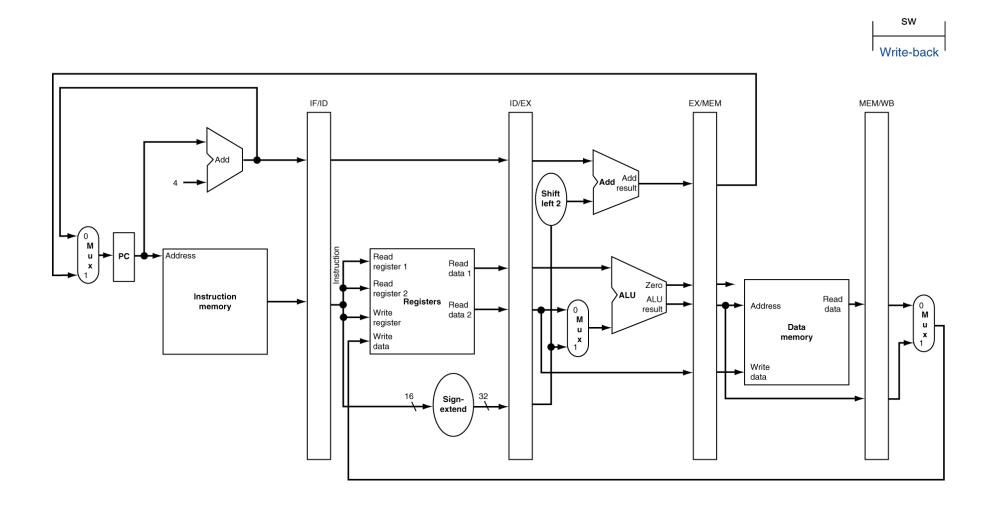




MEM for Store

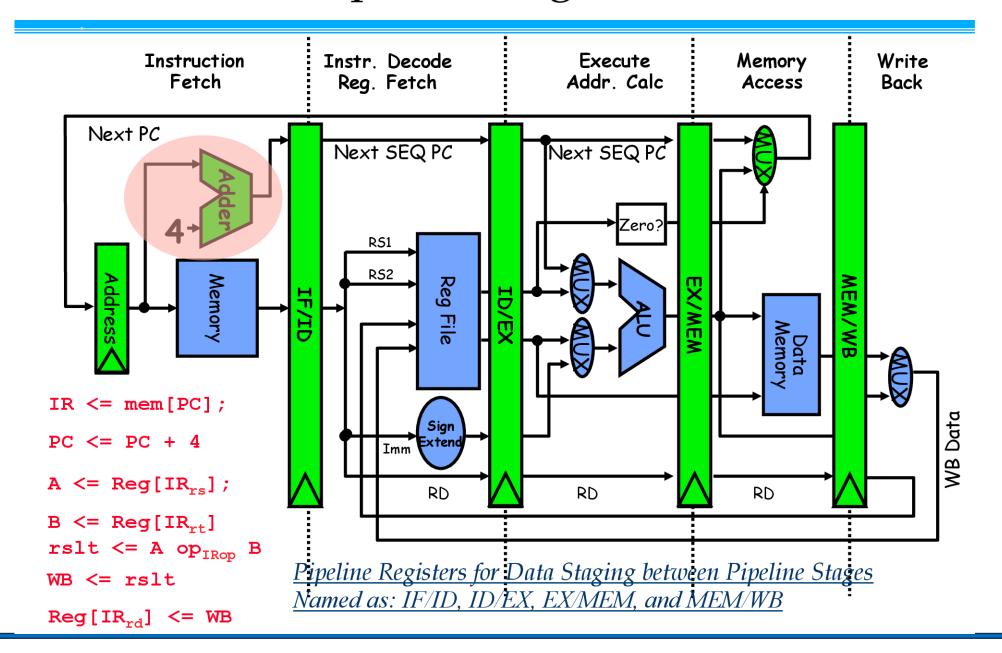


WB for Store



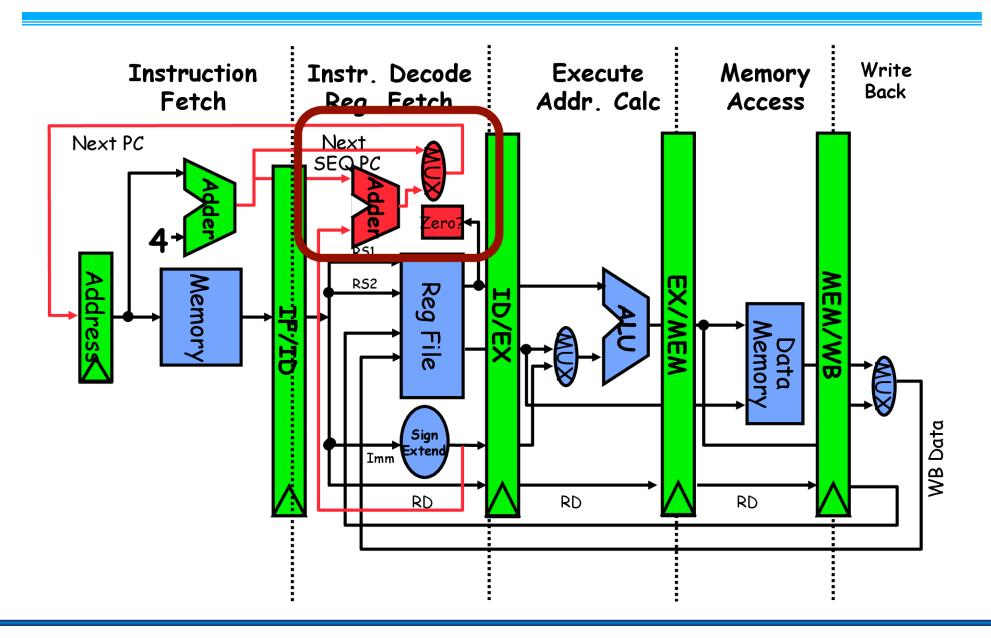
Pipeline Registers





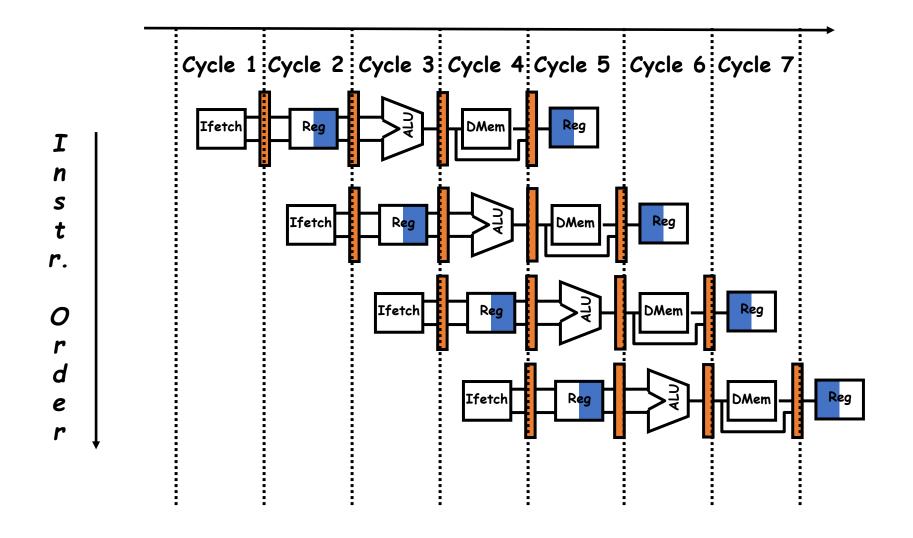
Pipelined MIPS Datapath





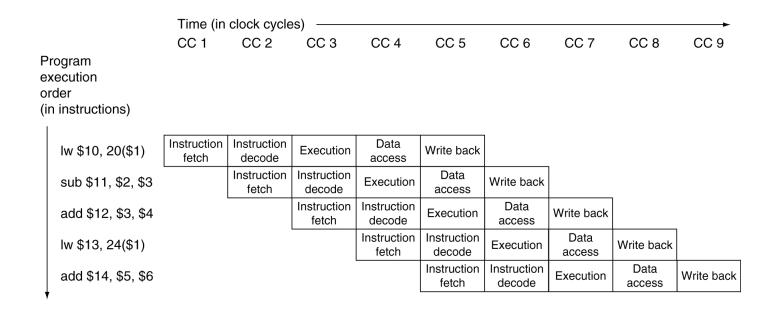
The Basic Pipeline For MIPS – Simplified Diagram





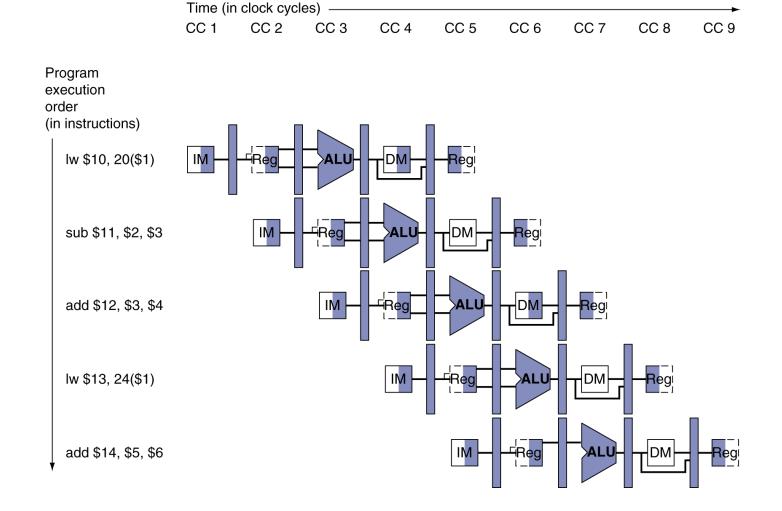
Multi-Cycle Pipeline Diagram

Traditional form



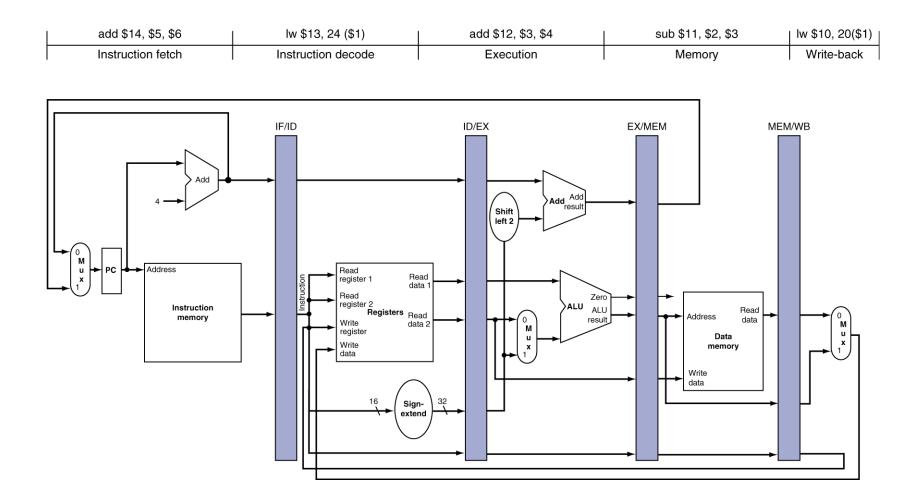
Multi-Cycle Pipeline Diagram

Form showing resource usage for different instructions



Single-Cycle Pipelined Diagram

State of pipeline in a given cycle

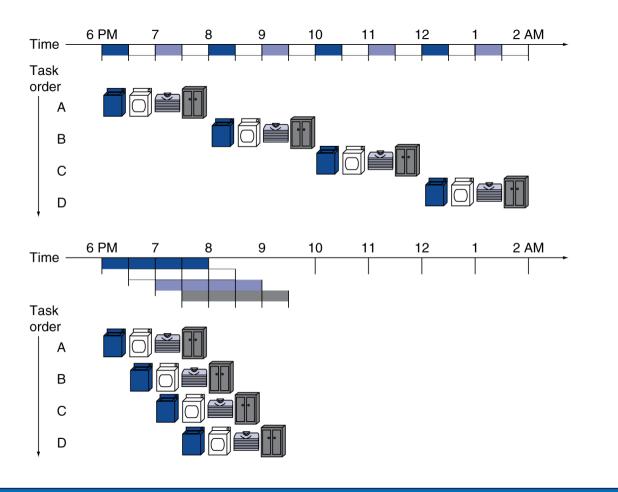




Pipeline Performance Estimation

Pipelining Analogy

- Pipelined laundry: overlapping execution
 - Parallelism improves performance



- Four loads:
 - Speedup= 8/3.5 = 2.3

Pipelined Architecture



Example program

CT=1.5ns; CPU Time =
$$9 \times 1 \times 1.5$$
ns = 13.5ns

AND R1,R2,R3

What is the cost of pipelining?

XOR R4,R2,R3

SUB R5,R1,R4

ADD R6,R1,R4

MUL R7,R5,R6

CPU Time = IC x CPI x CT

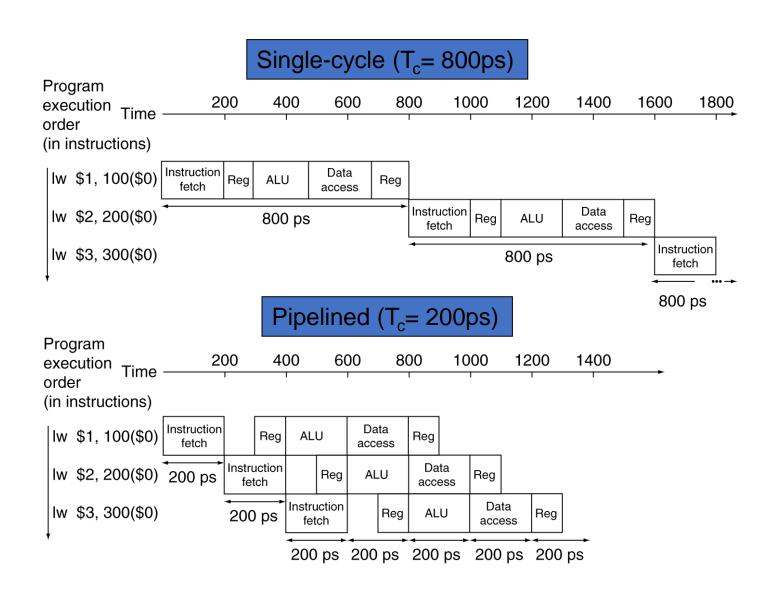
Time

Pipeline Performance

- Assume time for stages is
 - 100ps for register read or write
 - 200ps for other stages
- Compare pipelined datapath with single-cycle datapath

Instr	Instr fetch	Register read	ALU op	Memory access	Register write	Total time
lw	200ps	100 ps	200ps	200ps	100 ps	800ps
SW	200ps	100 ps	200ps	200ps		700ps
R-format	200ps	100 ps	200ps		100 ps	600ps
beq	200ps	100 ps	200ps			500ps

Pipeline Performance





Multi-cycle Control Path

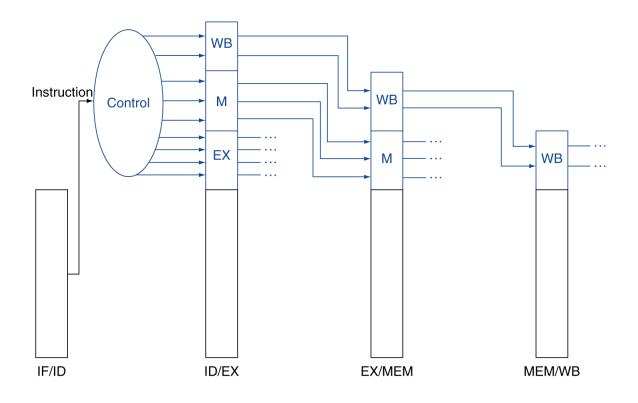
Multi-cycle Control Design



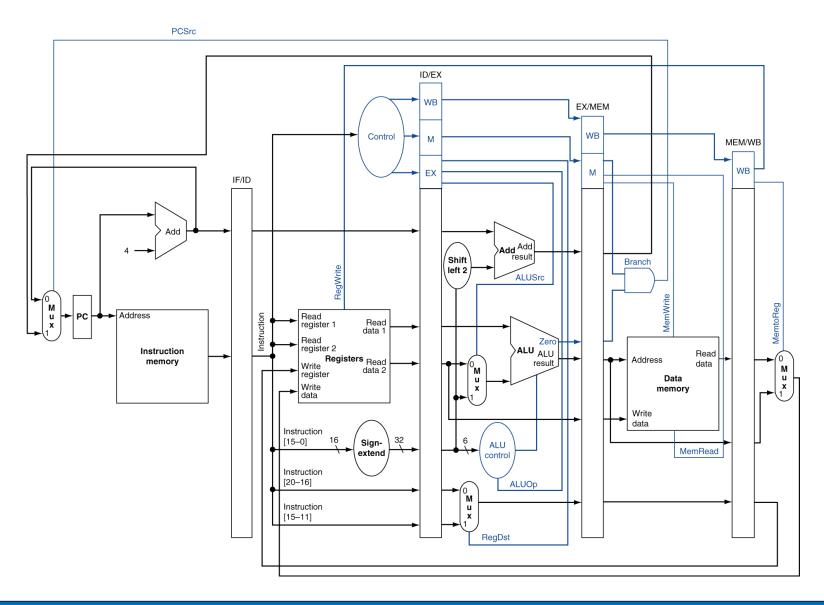
- Break instructions into cycles
- Put cycle sequences together
- Control signal groups and micro operations
- Control states and signal values
- Control state transitions

Pipelined Control

- Control signals derived from instruction
 - As in single-cycle implementation



Pipelined Control with Connections



Readings

Chap 4 of P&H Textbook