CS 253: Algorithms

Chapter 22

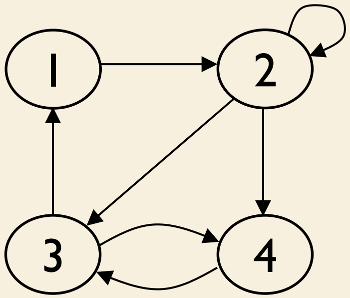
Graphs

Credit: Dr. George Bebis

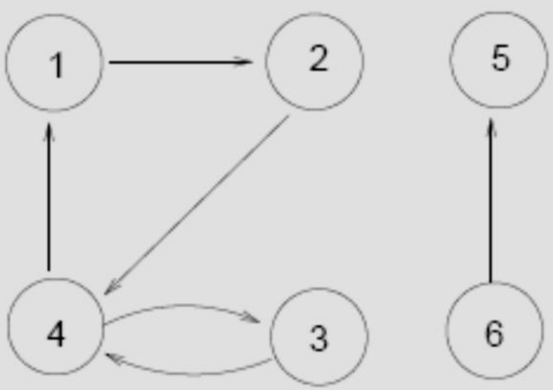
Graphs

Definition = a set of nodes (vertices) with edges (links) between them.

- G = (V, E) - graph
- V = set of vertices $|V| = n$
- E = set of edges $|E| = m$
  - Subset of V x V $=\{(u,v): u \in V, v \in V\}$

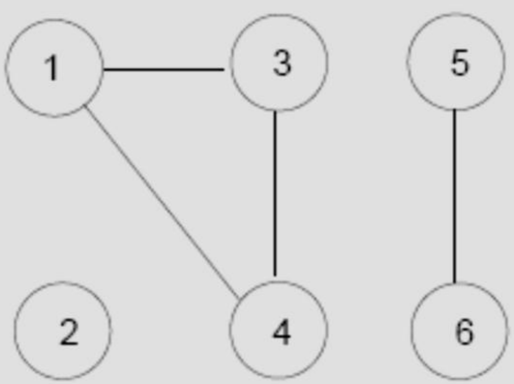


Directed graphs (digraphs)
(ordered pairs of vertices)



in-degree of v: # edges enetring v
out-degree of v: # edges leaving v

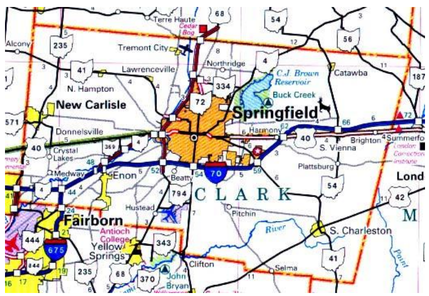
Undirected graphs
(unordered pairs of vertices)



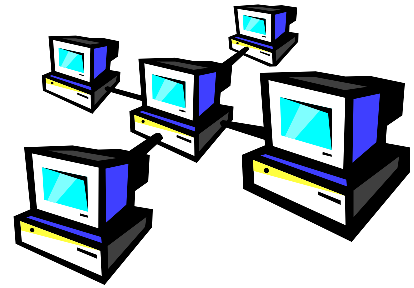
degree of v: # edges incident on v

Applications

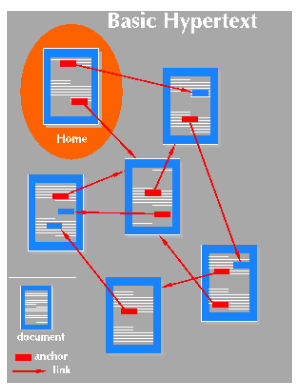
Applications that involve not only a set of items, but also the connections between them



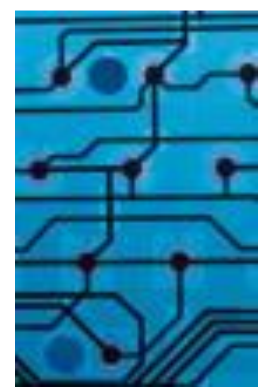
Maps



Computer networks



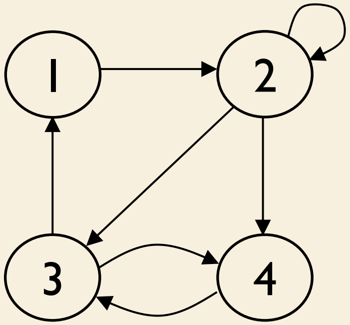
Hypertext



Circuits

Terminology

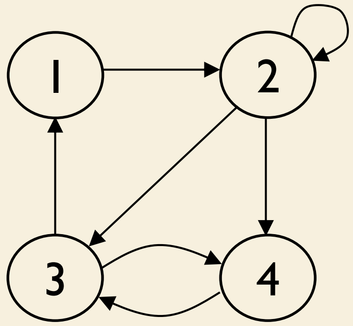
- Complete graph
  - A graph with an edge between each pair of vertices
- Subgraph
  - A graph (V', E') such that V'$\subseteq$V and E'$\subseteq$E
- Path from v to w
  - A sequence of vertices $<v_0, v_1, \ldots, v_k>$ such that $v_0=v$ and $v_k=w$
- Length of a path
  - Number of edges in the path



path from v_1 to v_4 $<v_1, v_2, v_4>$

Terminology (cont'd)

- w is **reachable** from v
  - If there is a path from v to w
- **Simple** path
  - All the vertices in the path are distinct
- Cycles
  - A path $<v_0, v_1, \ldots, v_k>$ forms a cycle if $v_0 = v_k$ and $k \geq 2$
- Acyclic graph
  - A graph without any cycles



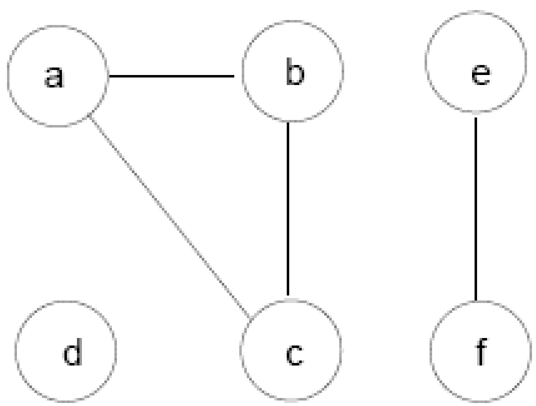
cycle from v_1 to v_1 $<v_1, v_2, v_3, v_1>$

Terminology (cont'd)

undirected graphs

connected: every pair of vertices is connected by a path

connected components: all possible connected subgraphs

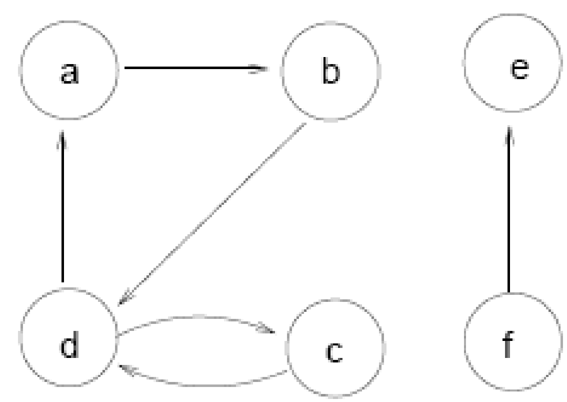


connected components: {a,b,c} {d} {e,f}

directed graphs

strongly connected: every two vertices are reachable from each other

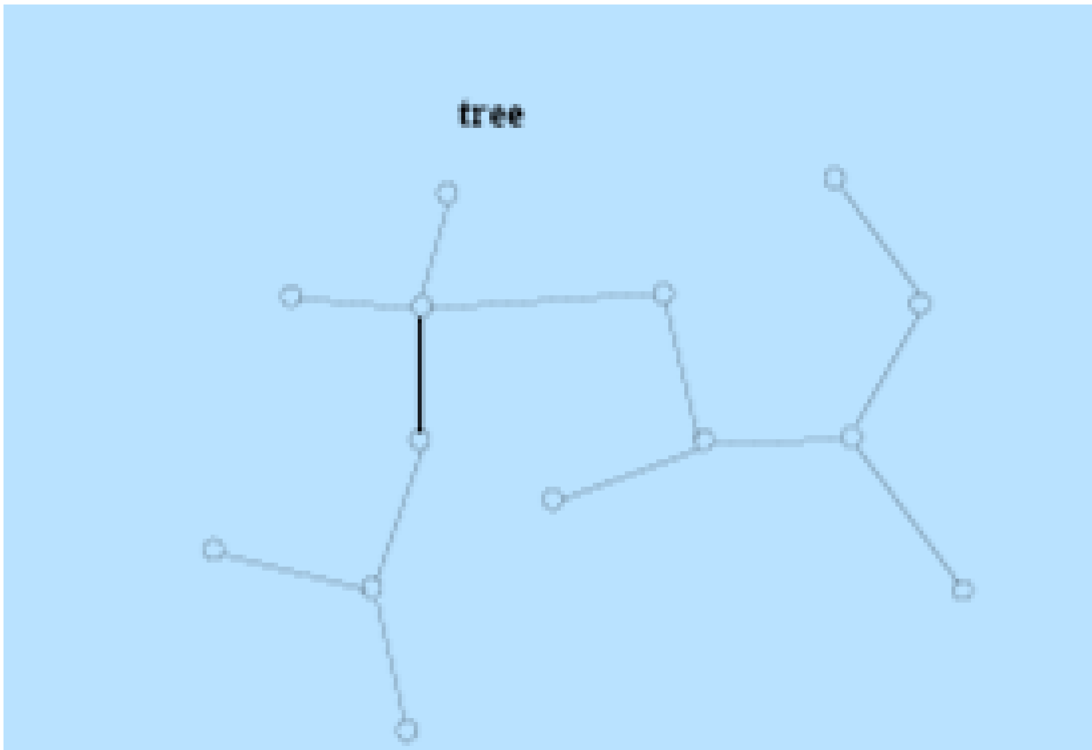
strongly connected components : all possible strongly connected subgraphs



strongly connected components: {a,b,c,d} { e} {f}

Terminology (cont'd)

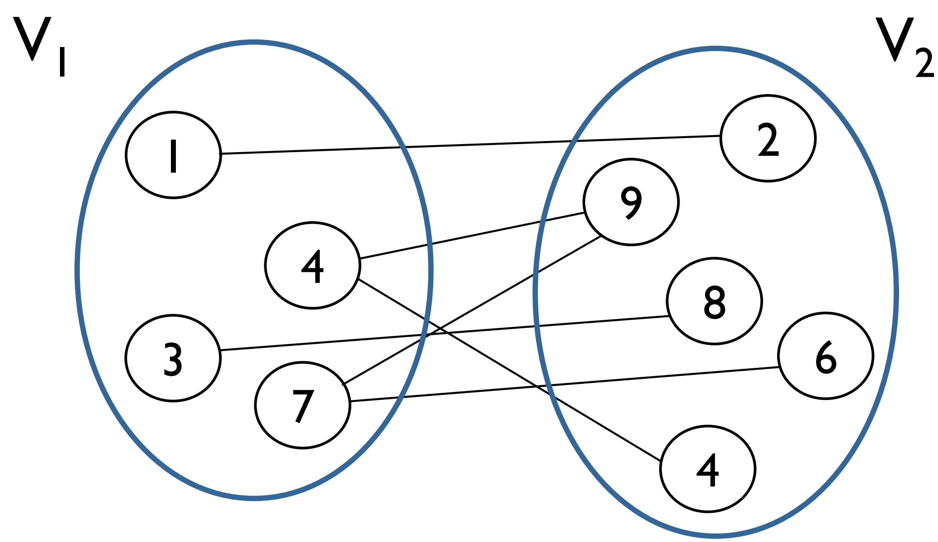
A **tree** is a connected, acyclic undirected graph



Terminology (cont'd)

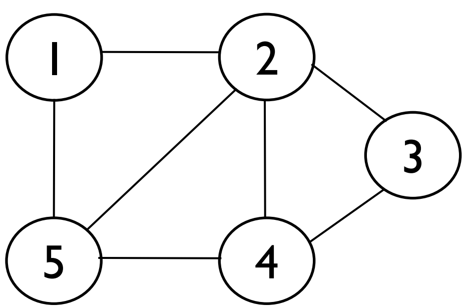
- A **bipartite graph** is an undirected graph

  $G = (V, E)$ in which $V = V_1 + V_2$ and there are edges only between vertices in V_1 and V_2

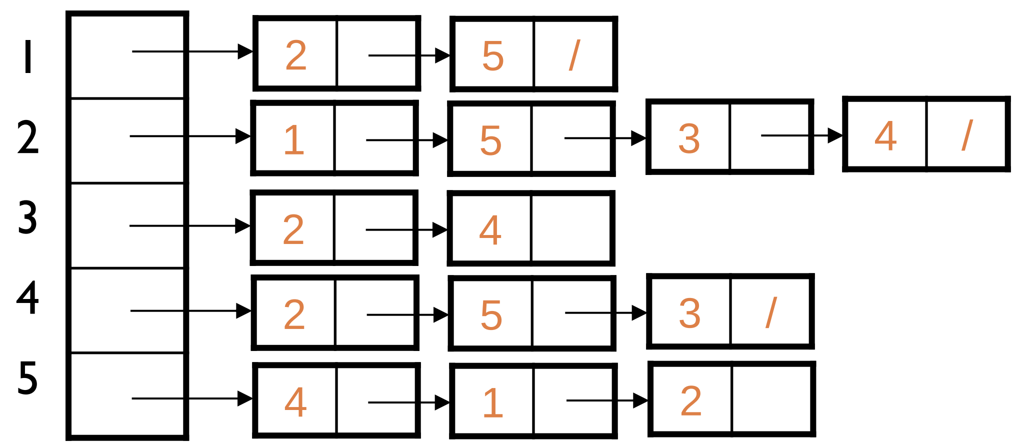


**here Graph Representation

- **Adjacency list representation** of G = (V, E)
  - An array of |V| lists, one for each vertex in V
  - Each list Adj[u] contains all the vertices v that are adjacent to u (i.e., there is an edge from u to v)
  - Can be used for both directed and undirected graphs

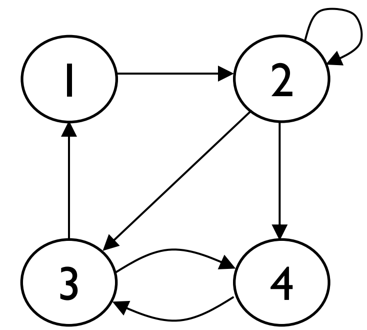


Undirected graph

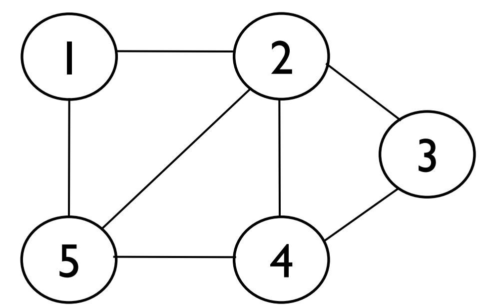


Properties of Adjacency-List Representation

- Sum of "lengths" of all adjacency lists
  - Directed graph: $|\mathbf{E}|$
    - edge (u, v) appears only once (i.e., in the list of **u**)
  - Undirected graph: $2|\mathbf{E}|$
    - edge (u, v) appears twice (i.e., in the lists of both **u** and **v**)



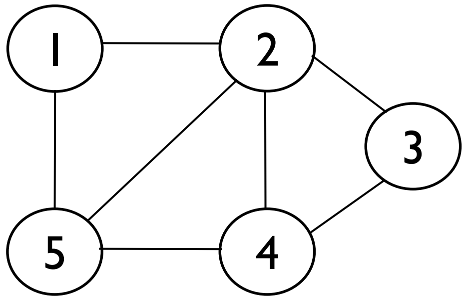
Directed graph



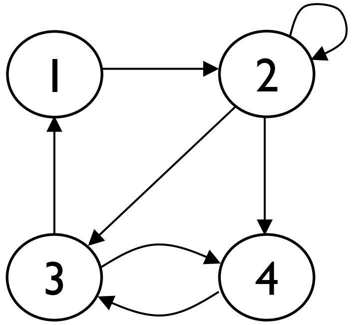
Undirected graph

Properties of Adjacency-List Representation

- Memory required = $\Theta(V + E)$
- Preferred when
  - The graph is **sparse**: $|E| << |V|^2$
  - We need to quickly determine the nodes adjacent to a given node.
- Disadvantage
  - No quick way to determine whether there is an edge between node u and v
- Time to determine if $(u, v) \in E$:
  - $O(\text{degree}(u))$
- Time to list all vertices adjacent to u:
  - $\Theta(\text{degree}(u))$



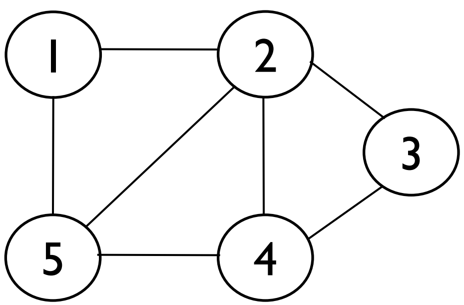
Undirected graph



Directed graph

Graph Representation

- **Adjacency matrix representation** of G = (V, E)
  - Assume vertices are numbered 1, 2, … |V|
  - The representation consists of a matrix $A_{|V| x |V|}$:
  - $a_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E \\ 0 & \text{otherwise} \end{cases}$



Undirected graph

	1	2	3	4	5
1	0	1	0	0	1
2	1	0	1	1	1
3	0	1	0	1	0
4	0	1	1	0	1
5	1	1	0	1	0

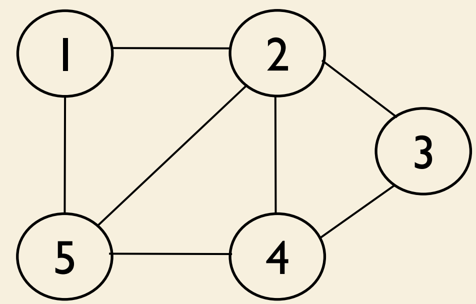
For undirected graphs, matrix A is symmetric:

$a_{ij} = a_{ji}$

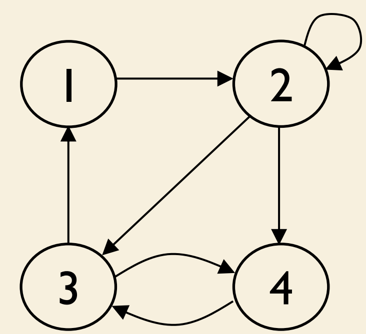
$A = A^T$

Properties of Adjacency Matrix Representation

- Memory required
  - $\Theta(V^2)$, independent on the number of edges in G
- Preferred when
  - The graph is **dense:** $|E|$ is close to $|V|^2$
  - We need to quickly determine if there is an edge between two vertices
- Time to determine if $(u, v) \in E$ ➔ $\Theta(1)$
- Disadvantage
  - No quick way to list all of the vertices adjacent to a vertex
- Time to list all vertices adjacent to u ➔ $\Theta(V)$

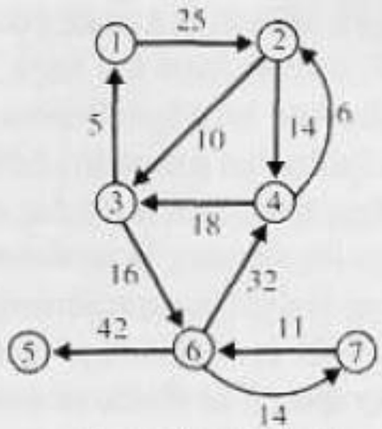


Undirected graph



Directed graph

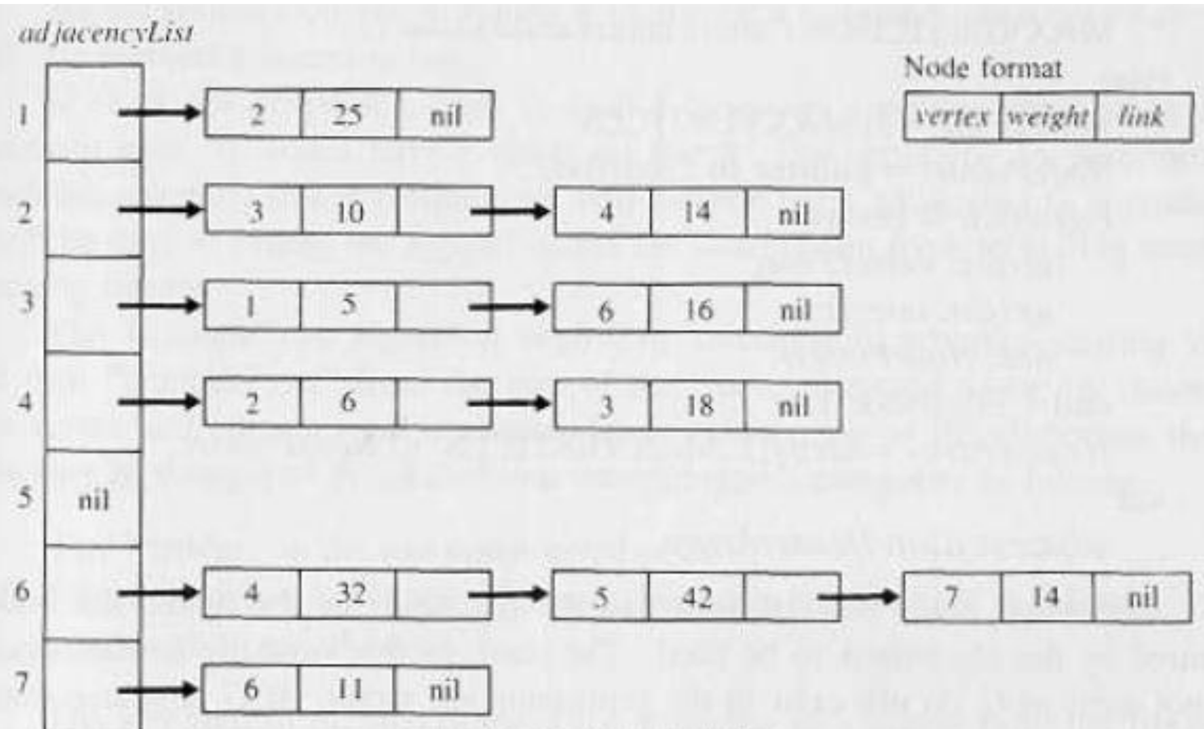
Weighted Graphs



(a) A weighted digraph.

$$
\begin{pmatrix}
0 & 25 & \infty & \infty & \infty & \infty & \infty \\
\infty & 0 & 10 & 14 & \infty & \infty & \infty \\
5 & \infty & 0 & \infty & \infty & 16 & \infty \\
\infty & 6 & 18 & 0 & \infty & \infty & \infty \\
\infty & \infty & \infty & \infty & 0 & \infty & \infty \\
\infty & \infty & \infty & 32 & 42 & 0 & 14 \\
\infty & \infty & \infty & \infty & \infty & 11 & 0
\end{pmatrix}
$$

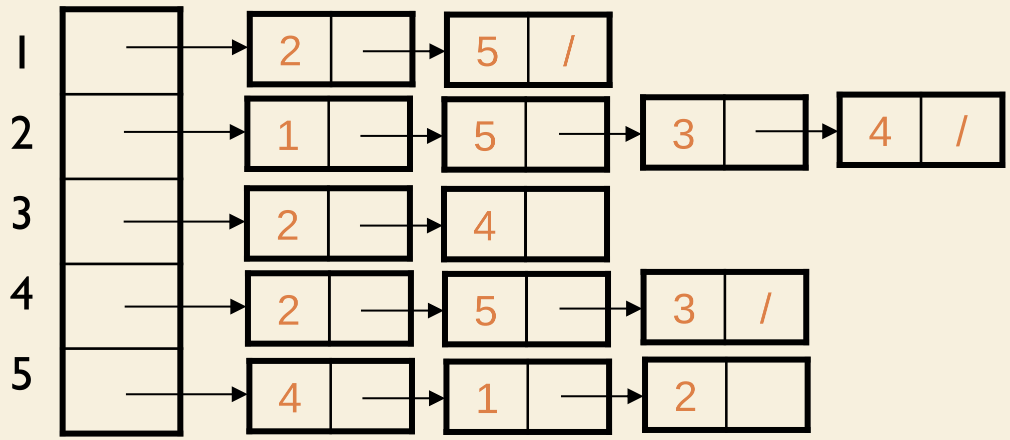
(b) Its adjacency matrix.



Problem 1

- Given an adjacency-list representation, how long does it take to compute the out-degree of every vertex?

  - For each vertex u, search Adj[u] → $\Theta(V+E)$



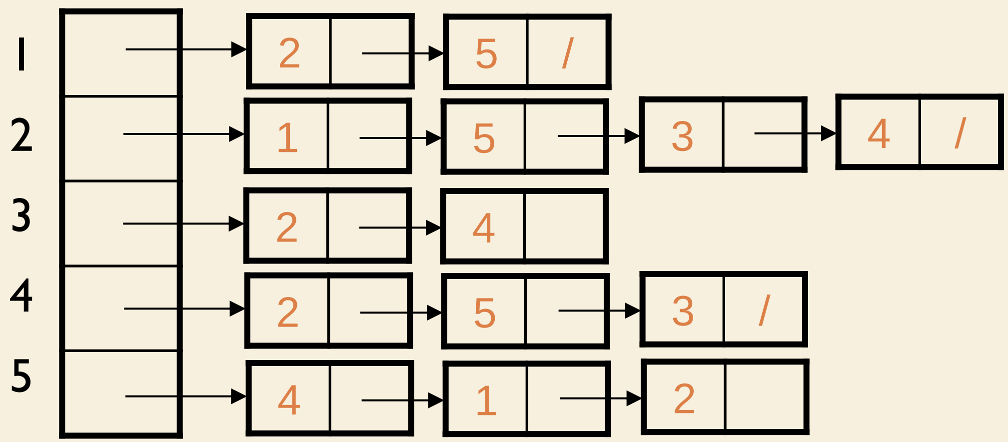
- How about using an adjacency-matrix representation?

  → $\Theta(V^2)$

Problem 2

How long does it take to compute the in-degree of every vertex?

- For each vertex u, search entire list of edges → Θ(V+E)



Problem 3

- The transpose of a graph $G=(V,E)$ is the graph $G^T=(V,E^T)$, where $E^T=\{(v,u) \epsilon V \times V: (u,v) \epsilon E\}$. Thus, G^T is G with all edges reversed.

(a) Describe an efficient algorithm for computing G^T from G, both for the adjacency-list and adjacency-matrix representations of G.

(b) Analyze the running time of each algorithm.

Problem 3 (cont'd)

Adjacency matrix

```
for (i=1; i ≤ V; i++)
    for(j=i+1; j ≤ V; j++)
    if(A[i][j] && !A[j][i]) {
      A[i][j]=0;
      A[j][i]=1;
    }
```

$O(V^2)$ complexity

	1	2	3	4	5
1	0	1	0	0	1
2	0	0	0	1	1
3	0	1	0	0	0
4	0	1	1	0	1
5	0	0	0	1	0

Problem 3 (cont'd)

Adjacency list

Allocate V list pointers for G^T (Adj'[]) ← O(V)

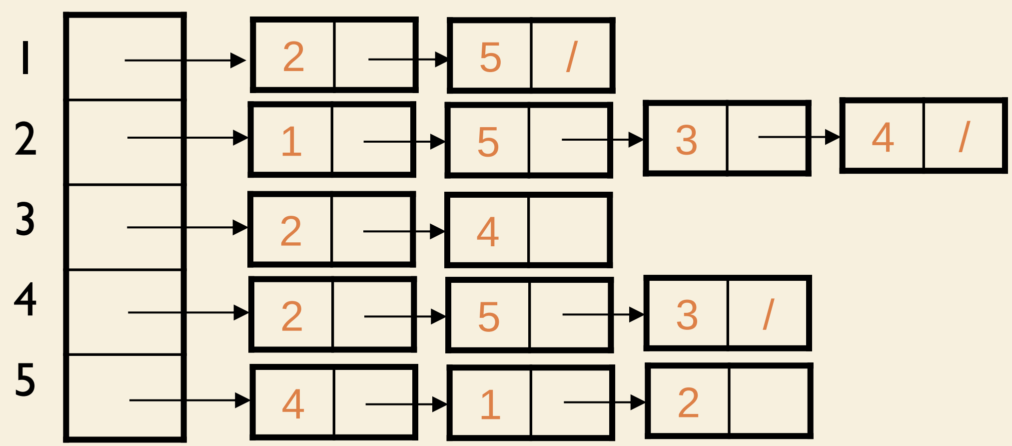
for(i=1; i ≤ V, i++)

    for every vertex v in Adj[i]

        add vertex i to Adj'[v]

← O(E)

Total time: O(V+E)

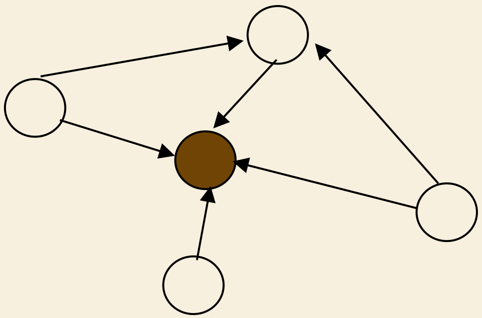


Problem 4

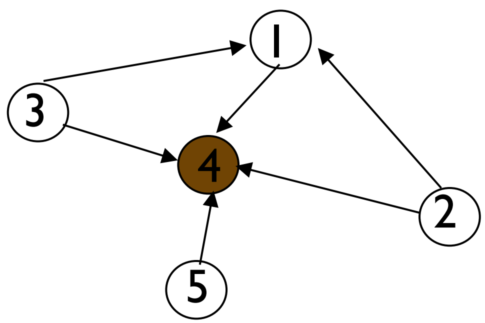
- When adjacency-matrix representation is used, most graph algorithms require time $\Omega(V^2)$, but there are some exceptions.

  Show that determining whether a directed graph G contains a **universal sink** – a vertex of in-degree |V|-1 and out-degree 0 – can be determined in time O(V).

  Example:



Problem 4 (cont.)



	1	2	3	4	5
1	0	0	0	1	0
2	1	0	0	1	0
3	1	0	0	1	0
4	0	0	0	0	0
5	0	0	0	1	0

Problem 4 (cont.)

- How many sinks could a graph have?
  - 0 or 1
- How can we determine whether a given vertex u is a universal sink?
  - The u-row must contain 0's only
  - The u-column must contain 1's only
  - A[u][u]=0
- How long would it take to determine whether a given vertex u is a universal sink?
  - O(V) time

Problem 4 (cont.)

A SIMPLE ALGORITHM TO CHECK FOR UNIVERSAL SINK:

```
IS-SINK(A, k)
let A be |V| × |V|
for j ← 1 to |V|                    ▷ Check for a 1 in row k
    do if a_kj = 1
        then return FALSE
for i ← 1 to |V|                    ▷ Check for an off-diagonal 0 in column k
    do if a_ik = 0 and i ≠ k
        then return FALSE
return TRUE
```

How long would it take to determine whether a given graph contains a universal sink if you were to check every single vertex in the graph?

$O(V^2)$

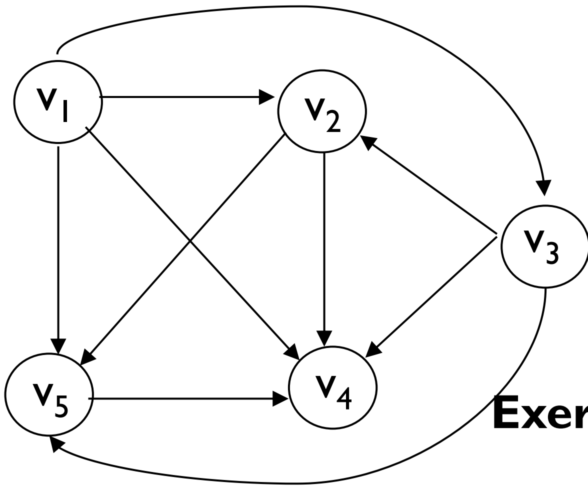
Problem 4 (cont.)

- Can you come up with a O(V) algorithm?
- Observations
  - If A[u][v]=1, then u cannot be a universal sink
  - If A[u][v]=0, then v cannot be a universal sink

```
UNIVERSAL-SINK(A)
let A be |V| × |V|
i ← j ← 1
while i ≤ |V| and j ≤ |V|
     do if a_ij = 1
            then i ← i + 1
            else  j ← j + 1
s ← 0
if i > |V|
   then return "there is no universal sink"
elseif IS-SINK(A, i) = FALSE
   then return "there is no universal sink"
else return i "is a universal sink"
```

← Why do we need this check?
ANSWER : see the last slide

Problem 4 (cont.)

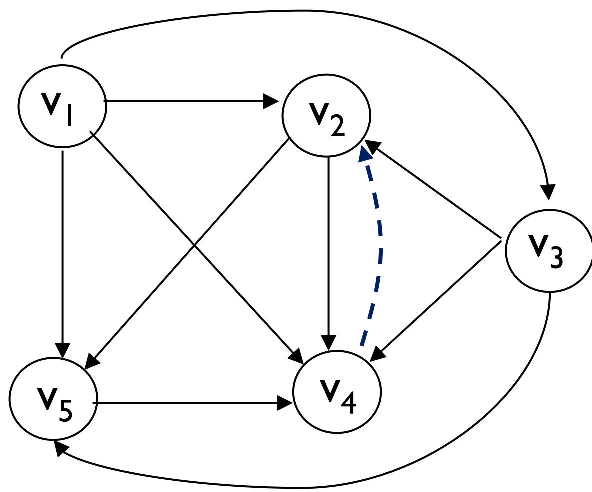


Exercise

$$
\begin{array}{c}
\begin{array}{ccccc} v_1 & v_2 & v_3 & v_4 & v_5 \end{array} \\
\left(\begin{array}{ccccc}
0 & 1 & 1 & 1 & 1 \\
0 & 0 & 0 & 1 & 1 \\
0 & 1 & 0 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0
\end{array}\right)
\end{array}
\begin{array}{c} \\ v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \end{array}
$$

- Loop terminates when $i > |V|$ or $j > |V|$
- Upon termination, the only vertex that could be a sink is i
  - If $i > |V|$, there is no sink
  - If $i < |V|$, then $j > |V|$
    - * vertices k where $1 \le k < i$ can not be sinks **Why?**
    - * vertices k where $i < k \le |V|$ can not be sinks **Why?**

Problem 4 (cont.)



$$
\begin{array}{c|c}
 & \begin{array}{ccccc} v_1 & v_2 & v_3 & v_4 & v_5 \end{array} \\
\begin{array}{c} v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \end{array} &
\begin{pmatrix}
0 & 1 & 1 & 1 & 1 \\
0 & 0 & 0 & 1 & 1 \\
0 & 1 & 0 & 1 & 1 \\
0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0
\end{pmatrix}
\end{array}
$$