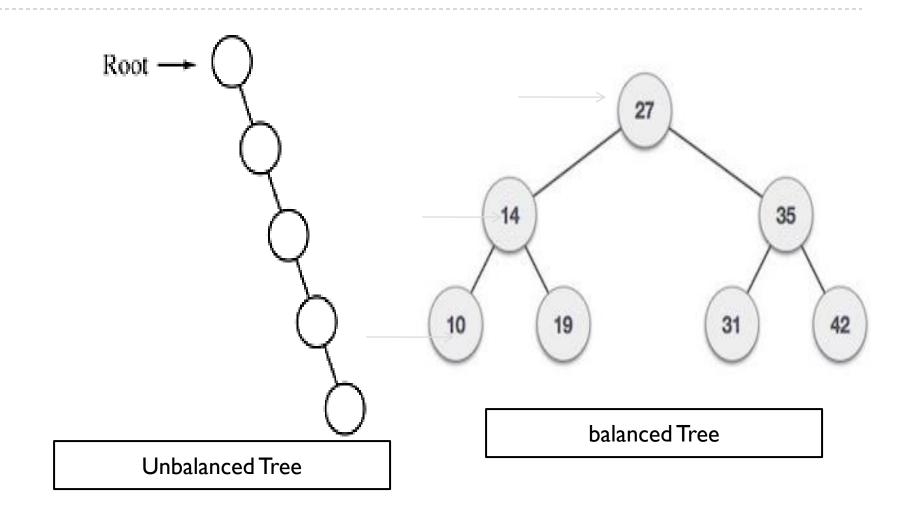
# AVL Tree

# Problem with Binary search tree

- The disadvantage of a binary search tree is that its height can be as large as N-I
- This means that the time needed to perform insertion and deletion and many other operations can be N times in the worst case.
- We want a tree with small height.
- A binary tree with N node has height at least  $log_2 (N+I)-I$
- Thus, our goal is to keep the height of a binary search tree log<sub>2</sub> (N) times
- Such trees are called balanced binary search trees. Examples are AVL tree, red-black tree.



# Binary search tree





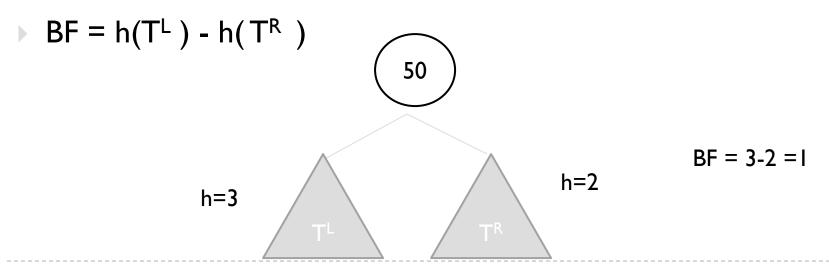
## AVL

- There is a need to maintain the binary search tree to be of balanced height, so that it is possible to obtain for the search option of log<sub>2</sub> (N) time in the worst case.
- One of the most popular balanced tree was introduced by Adelson velskii and landis (AVL)

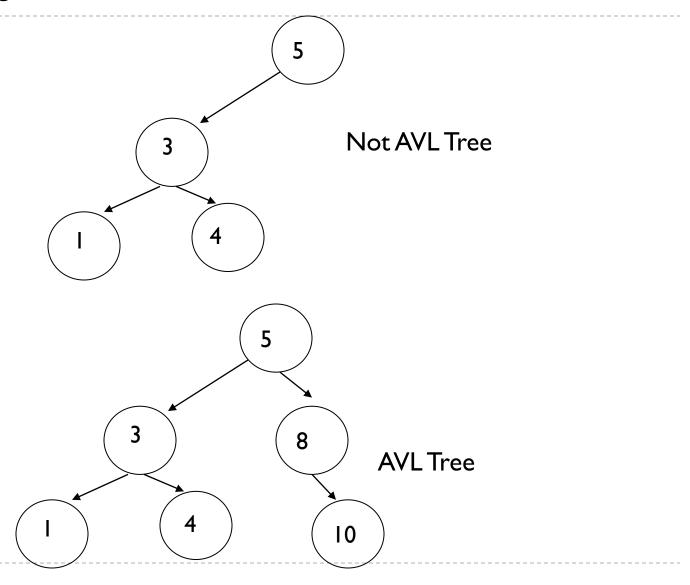


## **AVL** Tree

- An empty binary search tree is an AVL tree
- A non empty binary tree T is an AVL tree iff  $T^L$  (left subtree) and  $T^R$  (right subtree) of T and  $h(T^L)$  (height of left subtree) and  $h(T^R)$  (height of right subtree) where  $|h(T^L)| h(T^R)$  | <= 1
- Balance factor BF is difference between  $h(T^L)$  and  $h(T^R)$  and the value will be -1,0 or +1.



## AVL Tree

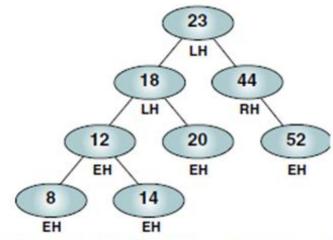


## **AVL Tree- Balance Factor**

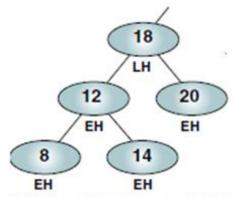
- $\blacktriangleright BF = h(T^L) h(T^R)$
- LH for left high (+1) to indicate that the left subtree is higher than the right subtree
- EH for even high (0) to indicate that the subtrees are the same height
- ▶ RH for right high (-I) to indicate that the left subtree is shorter than the right subtree.



## **AVL Tree- Balance Factor**

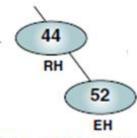


(a) Tree 23 appears balanced:  $H_L - H_R = 1$ 



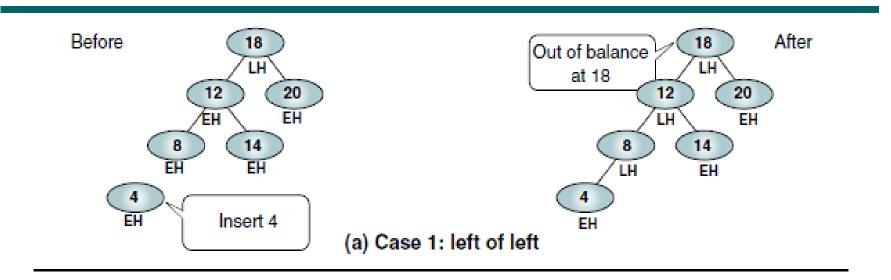
(b) Subtree 18 appears balanced:

$$H_L - H_R = 1$$



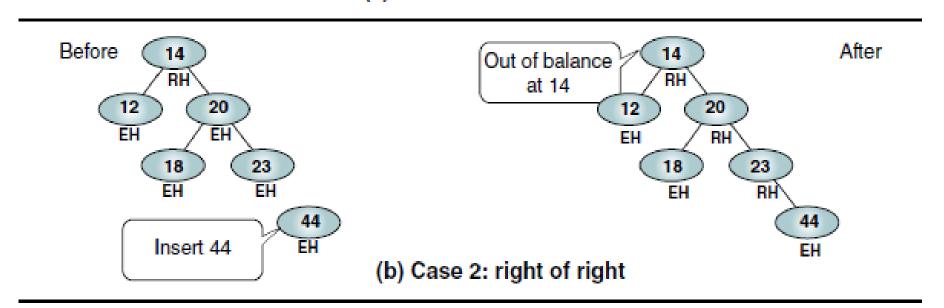
(c) Subtree 44 Is balanced:  $|H_L - H_R| = 1$ 

- Whenever we insert or delete a node from a tree, the resulting tree may be unbalanced.
  - We must rebalance it.
  - AVL trees are balanced by rotating nodes either to the left or to the right.
  - All unbalanced trees fall into one of these four cases:
  - Left of left: Inserted node is in the left subtree of left subtree of node A
  - Right of right: Inserted node is in the right subtree of right subtree of node A
  - Right of left: Inserted node is in the right subtree of left subtree of node A
  - ▶ **Left of right**: Inserted node is in the left subtree of right subtree of node A



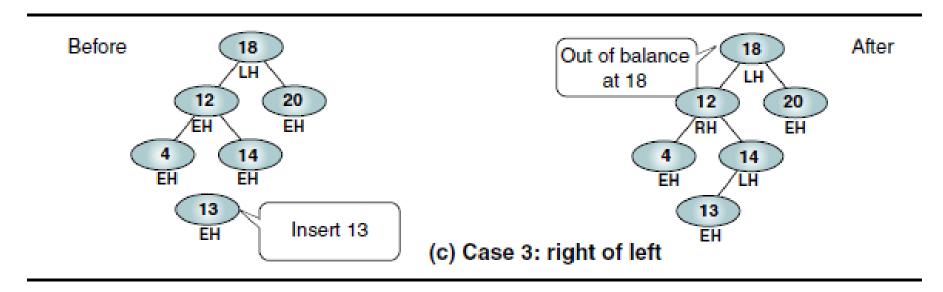
Out of balance AVL Tree





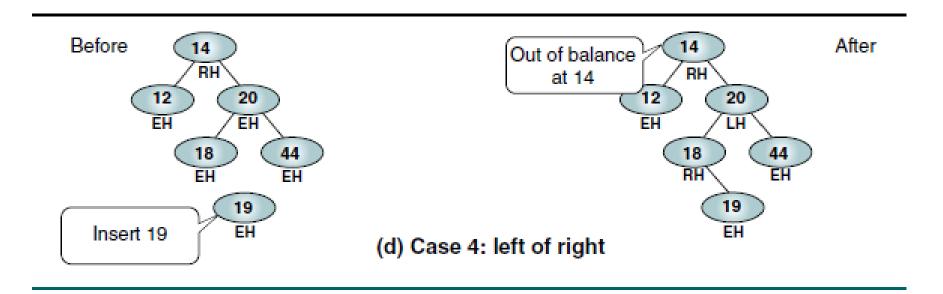
Out of balance AVL Tree





Out of balance AVL Tree



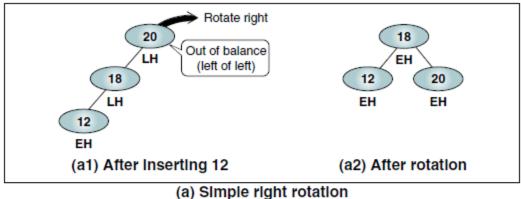


Out of balance AVL Tree



# Balancing Trees- Left of Left

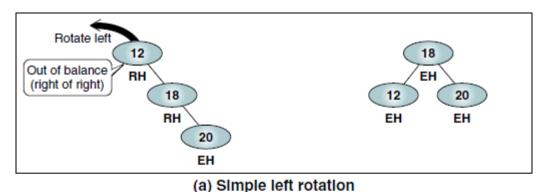
When the out-of-balance condition has been created by a left high subtree of a left high tree, we must balance the tree by rotating the out-of-balance node to the right.

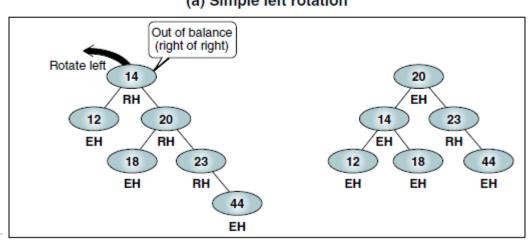


(b) Complex right rotation

# Balancing Trees- Right of Right

When the out-of-balance condition has been created by a right high subtree of a right high tree, we must balance the tree by rotating the out-of-balance node to the left.

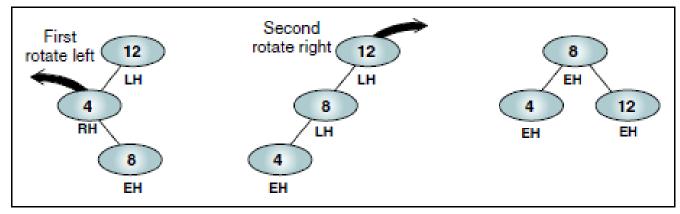




(b) Complex left rotation

# Balancing Trees- Right of Left

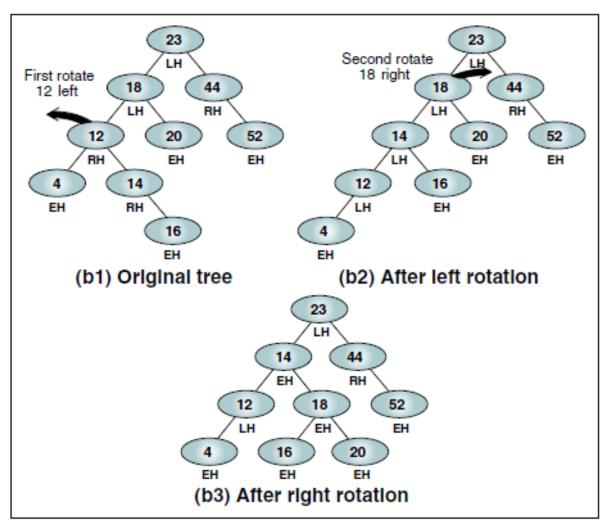
When the out-of-balance condition has been created by a right high subtree of a left high tree, we must balance the tree by rotating two nodes, one to the left and one to the right, to balance the tree.



(a) Simple double rotation right

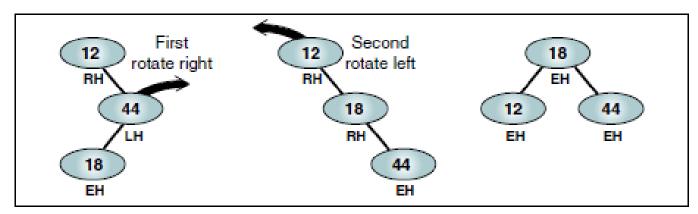


# Balancing Trees- Right of Left



# Balancing Trees- Left of Right

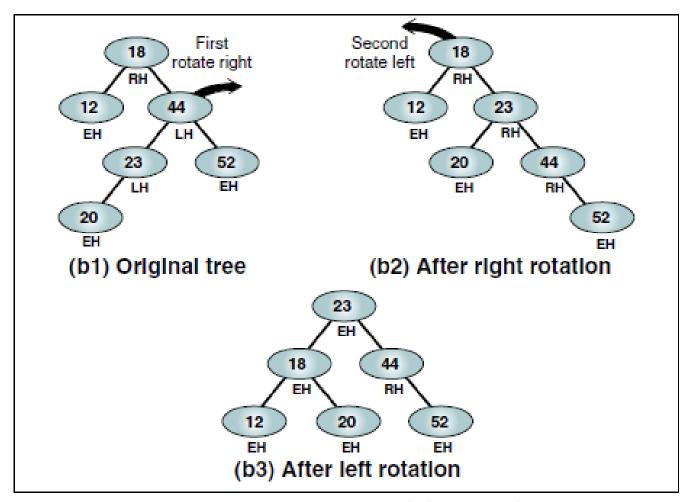
When the out-of-balance condition has been created by a lefthigh subtree of a right high tree, we must balance the tree by rotating two nodes, one to the right and one to the left.



(a) Simple double rotation right



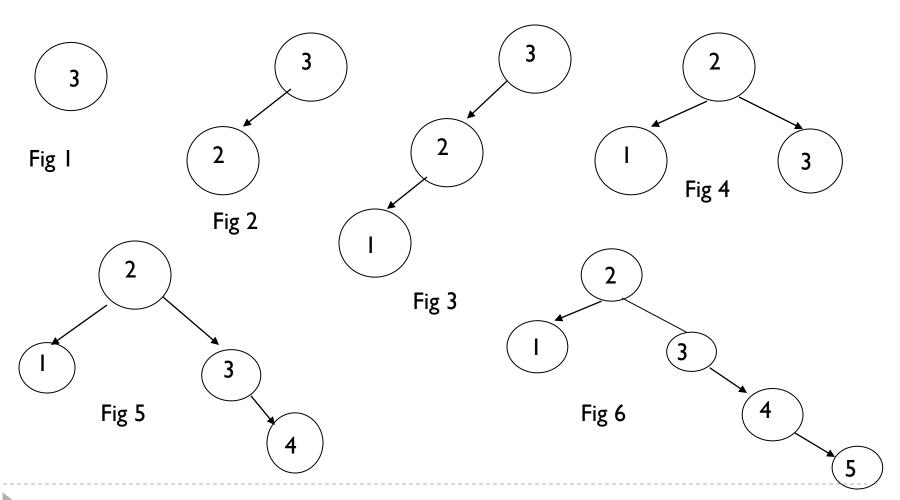
# Balancing Trees- Left of Right

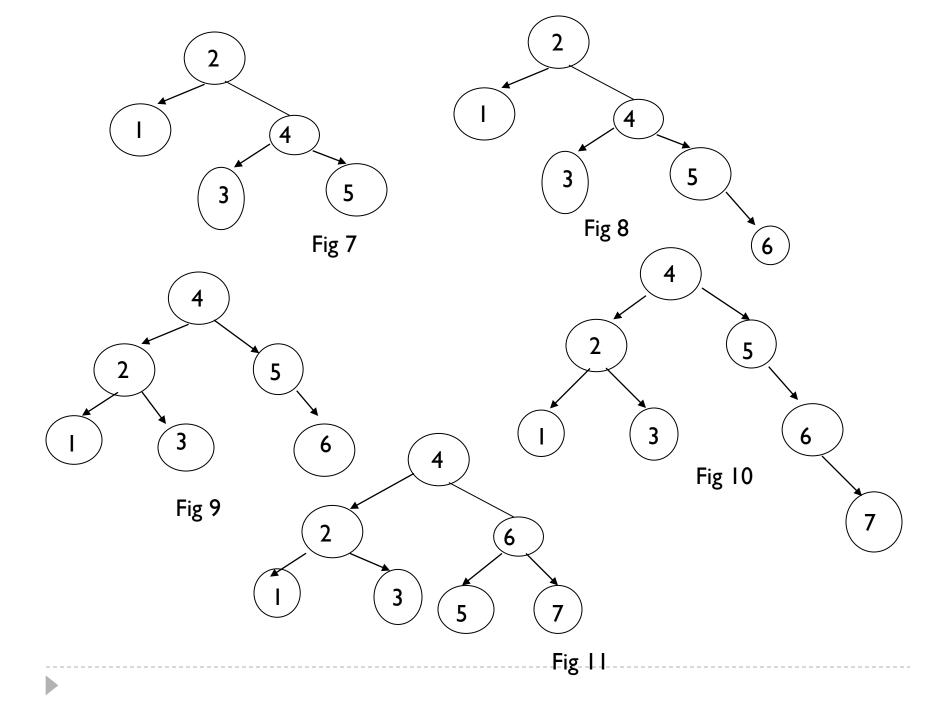


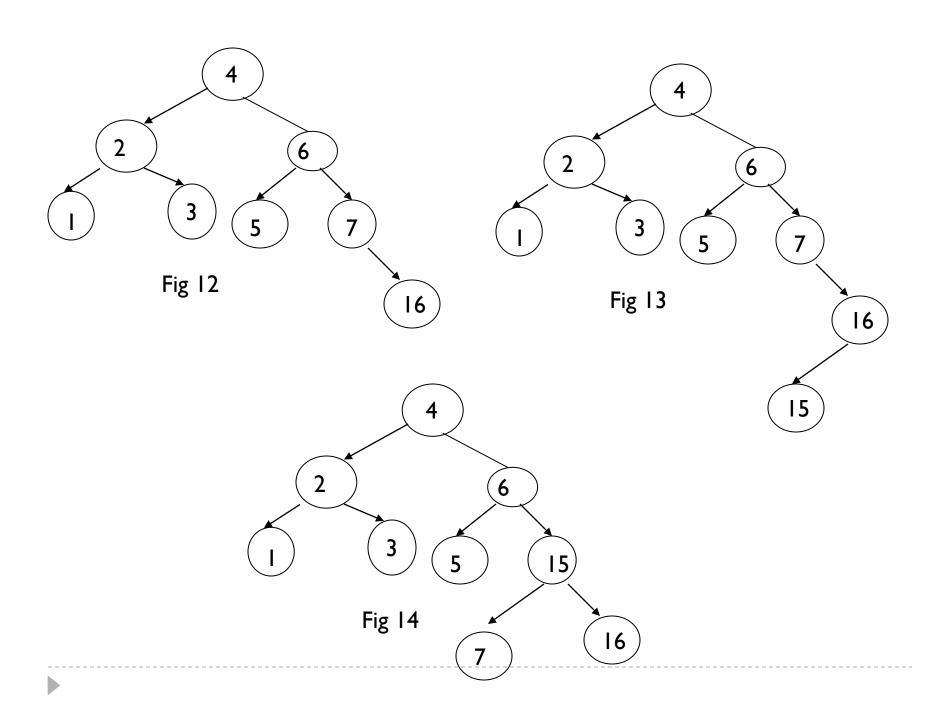
(b) Complex double rotation right

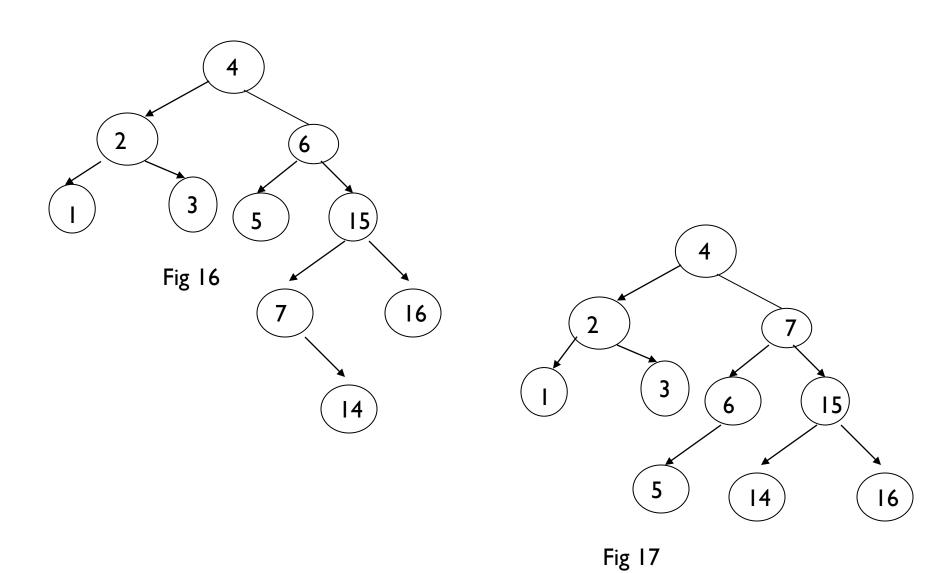
# Extended Example

Insert 3,2,1,4,5,6,7, 16,15,14









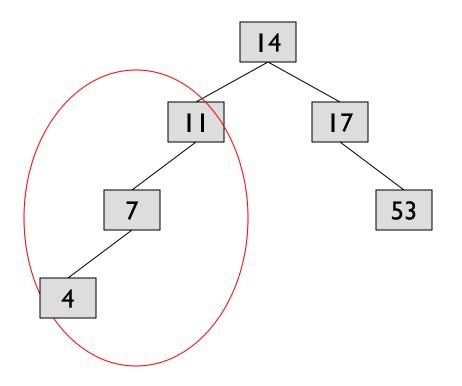
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## Deletion

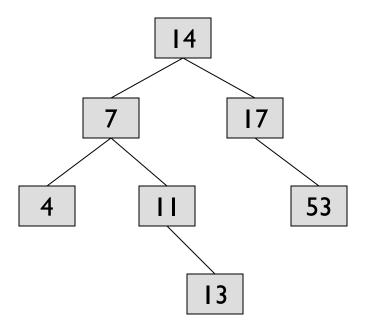
- Delete by a BST deletion by copying algorithm.
- Rebalance the tree if an imbalance occurs.
- There are three deletion cases:
  - Deletion that does not cause an imbalance.
  - 2. Deletion that requires a single rotation to rebalance.
  - 3. Deletion that requires two or more rotations to rebalance.



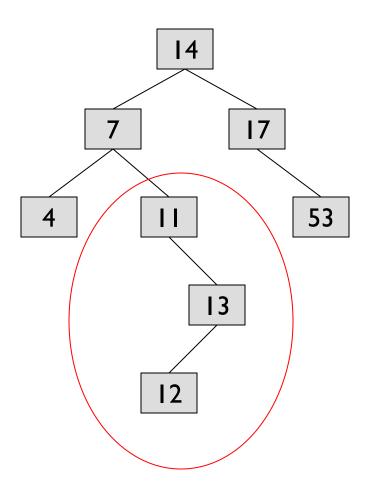
• Insert 14, 17, 11, 7, 53, 4, 13 into an empty AVL tree



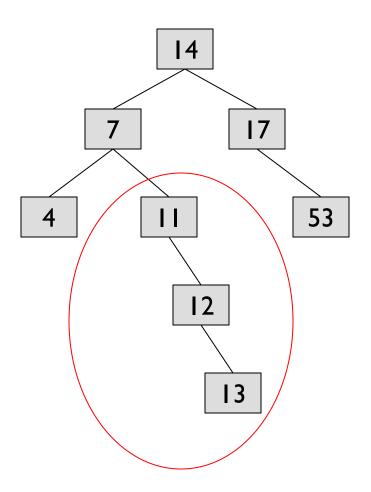
• Insert 14, 17, 11, 7, 53, 4, 13 into an empty AVL tree



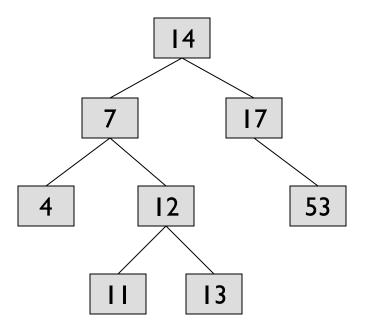
#### Now insert 12



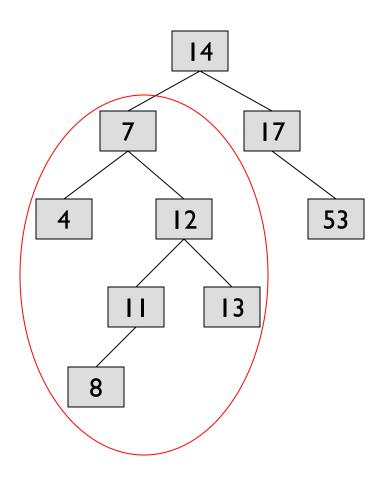
#### • Now insert 12



• Now the AVL tree is balanced.

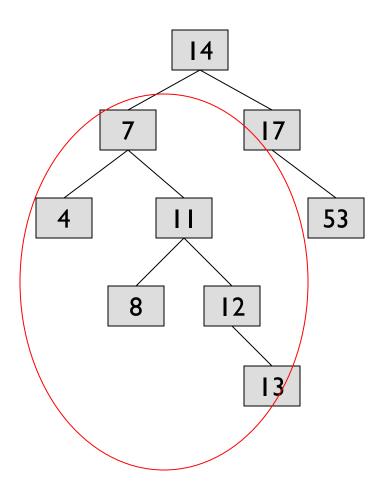


#### • Now insert 8



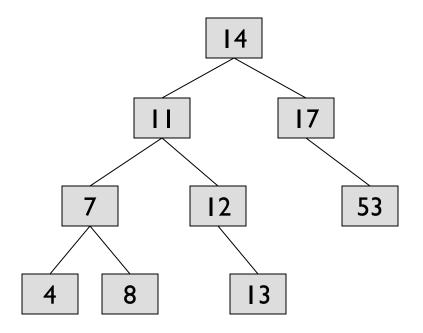


#### Now insert 8

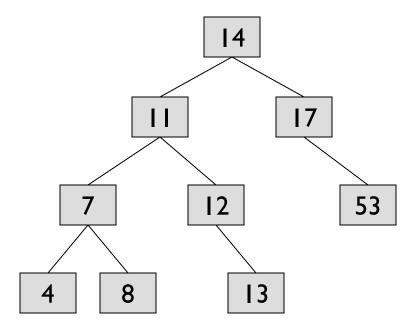




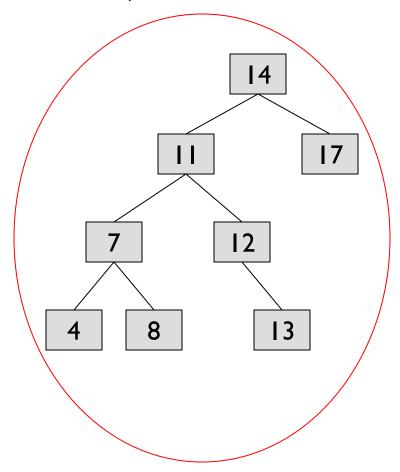
• Now the AVL tree is balanced.



#### Now remove 53

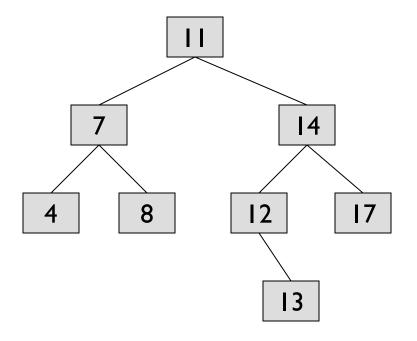


• Now remove 53, unbalanced

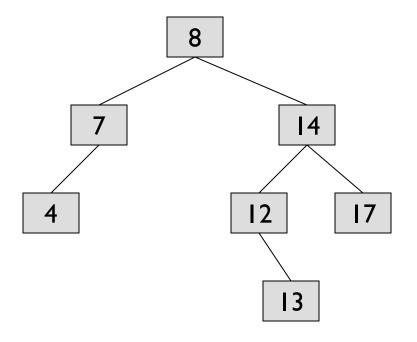




#### • Balanced! Remove II

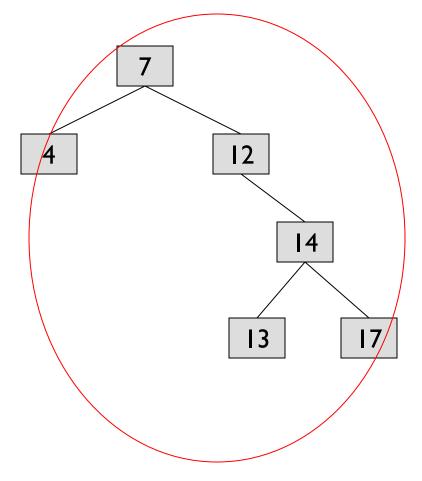


• Remove II, replace it with the largest in its left branch



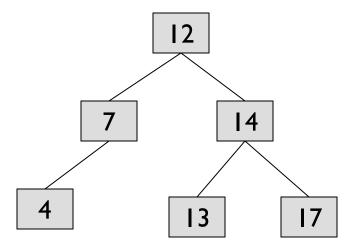
• Remove 8, unbalanced 12 17

• Remove 8, unbalanced





#### • Balanced!!

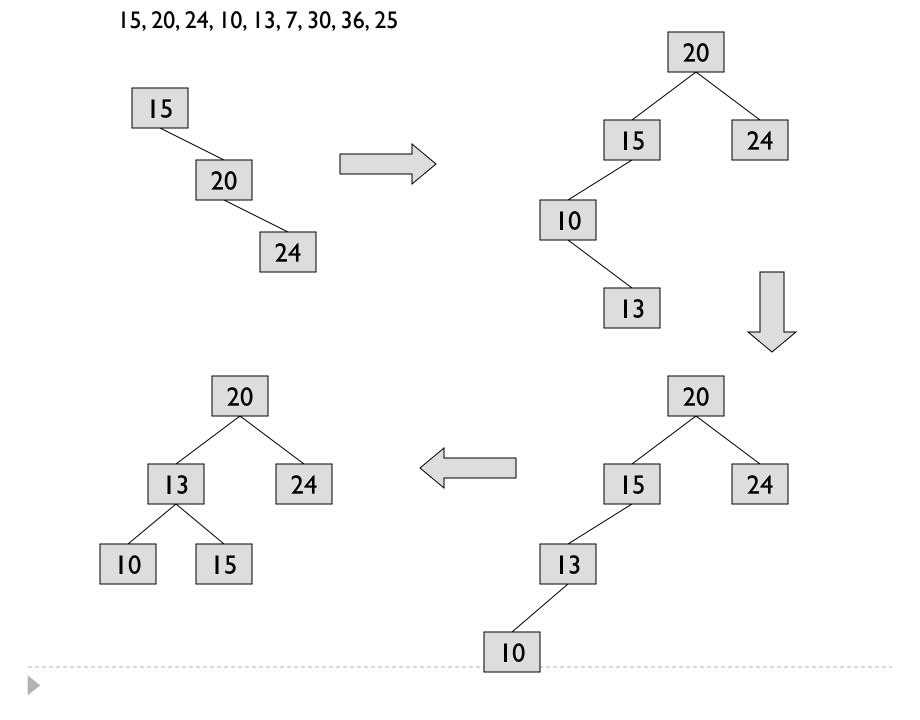




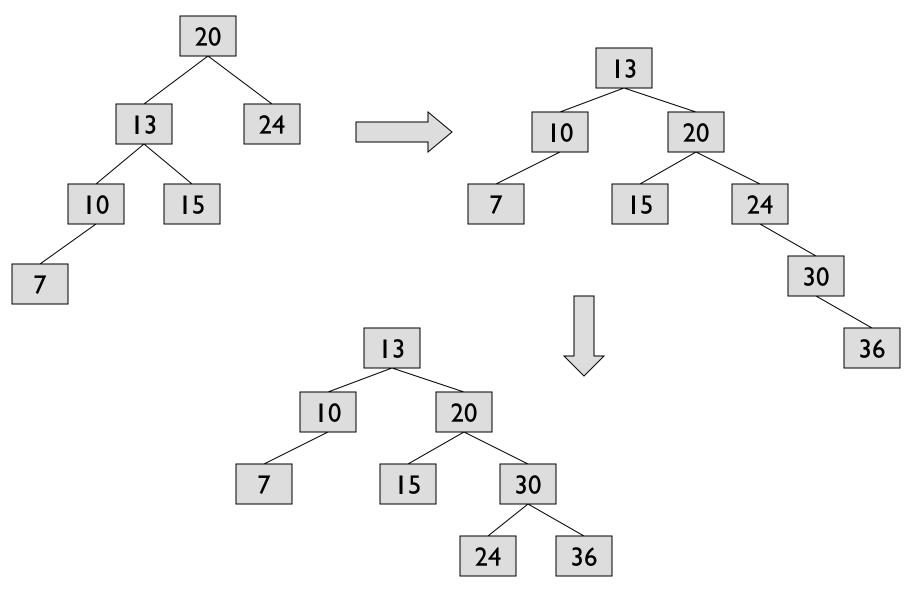
## In Class Exercises

Build an AVL tree with the following values:
15, 20, 24, 10, 13, 7, 30, 36, 25

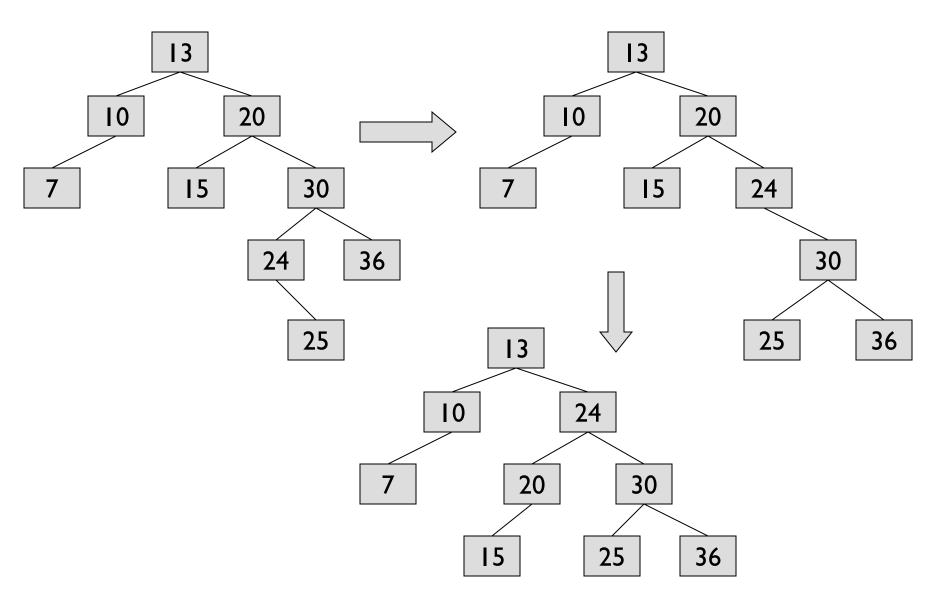




15, 20, 24, 10, 13, 7, 30, 36, 25



15, 20, 24, 10, 13, 7, 30, 36, 25



Remove 24 and 20 from the AVL tree.

