

# Inside Buffer Overflows



James D. Murray, CISSP C|EH

@jdmurray | [www.TechExams.net/blogs/jdmurray](http://www.TechExams.net/blogs/jdmurray)

# Inside Computer Programs

1000101011

42,657,212,33

password123

Data

```
#include <stdio.h>
int main(int argc, char *argv[])
{
    printf("Hello %s!\n", argv[0]);
    return 0;
}
```

Code

# Inside Computer Programs

1000101011  
42,657,212,33  
password123

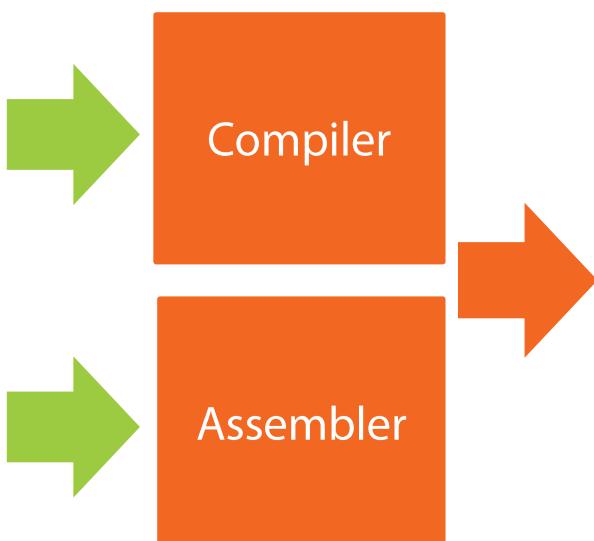
## Program Memory

```
#include <stdio.h>
int main(int argc, char *argv[])
{
    printf("Hello %s!\n", argv[0]);
    return 0;
}
```



# From Disk to RAM to CPU

```
#include <stdio.h>  
  
int main(int argc, char *argv[])  
{  
    printf("Hello %s!\n", argv[0]);  
    return 0;  
}  
  
.file    "hello.c"  
.section .rodata  
.string  "Hello %s!\n"  
.text  
.globl  main  
.type   main, @function  
main:  
.cfi_startproc
```

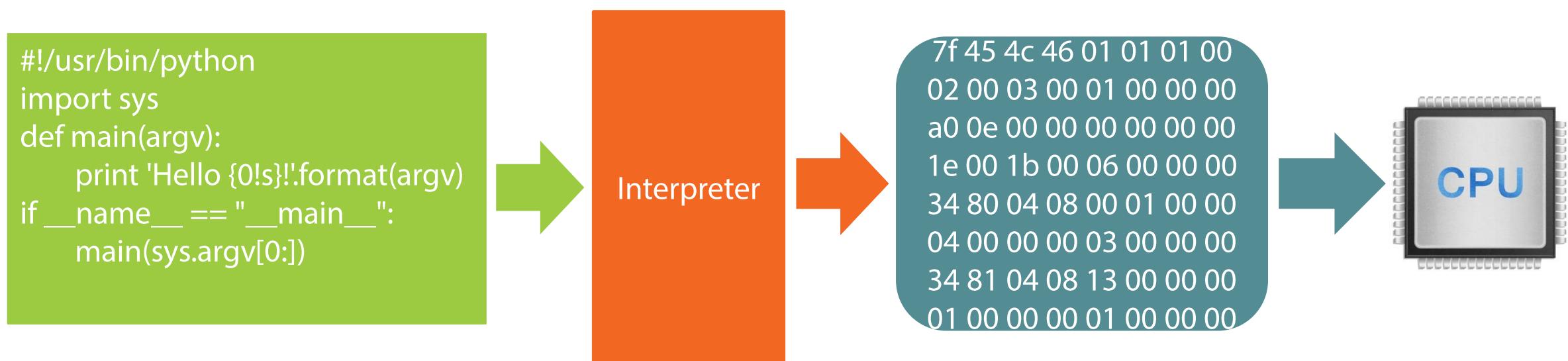


7f 45 4c 46 01 01 01 00  
02 00 03 00 01 00 00 00  
a0 0e 00 00 00 00 00 00  
1e 00 1b 00 06 00 00 00  
34 80 04 08 00 01 00 00  
04 00 00 00 03 00 00 00  
34 81 04 08 13 00 00 00  
01 00 00 00 01 00 00 00

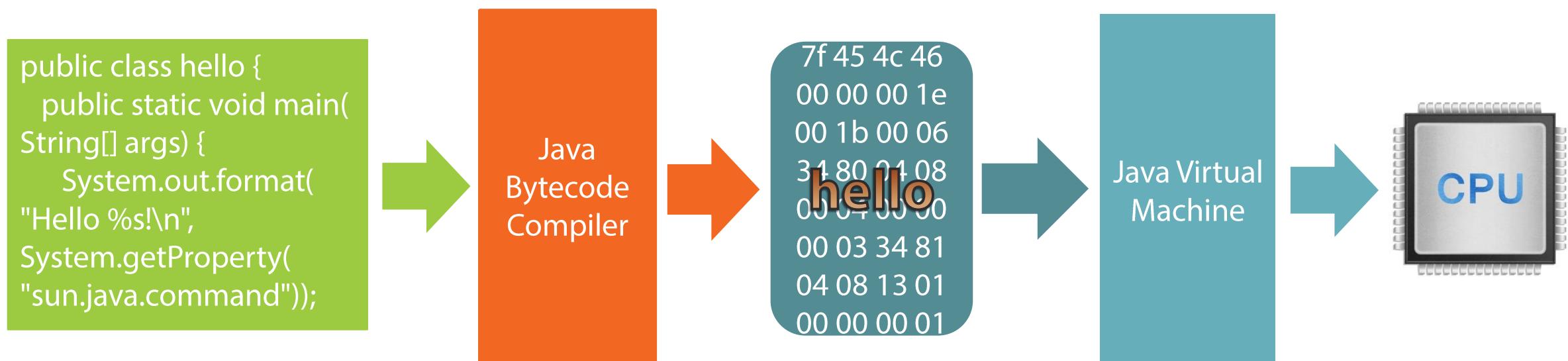
hello.exe



# From Disk to RAM to CPU

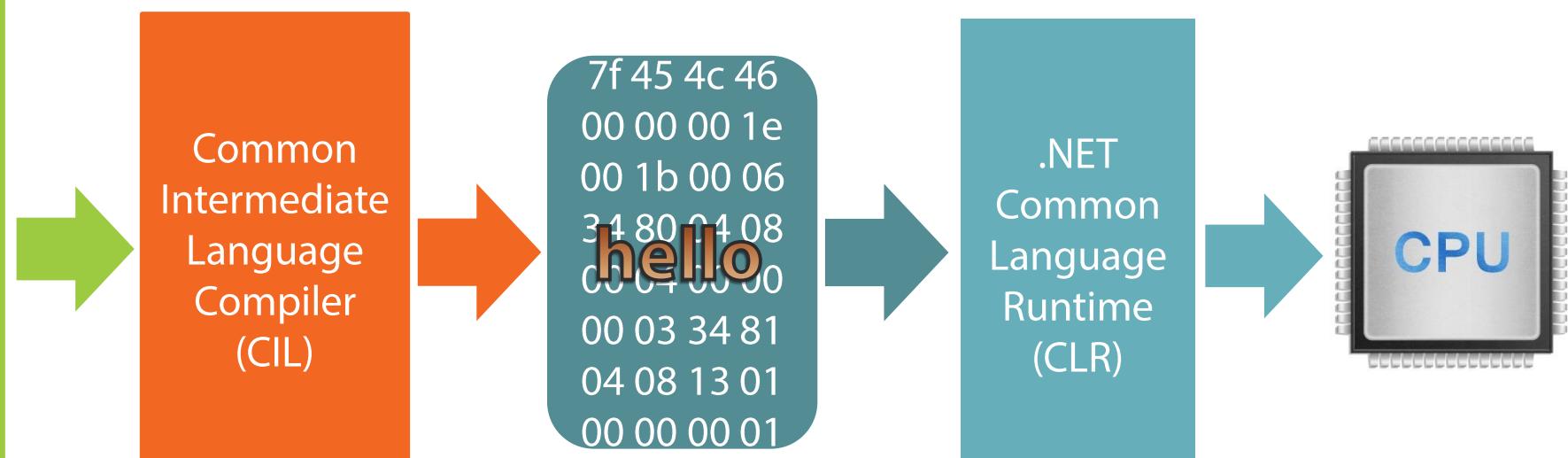


# From Disk to RAM to CPU

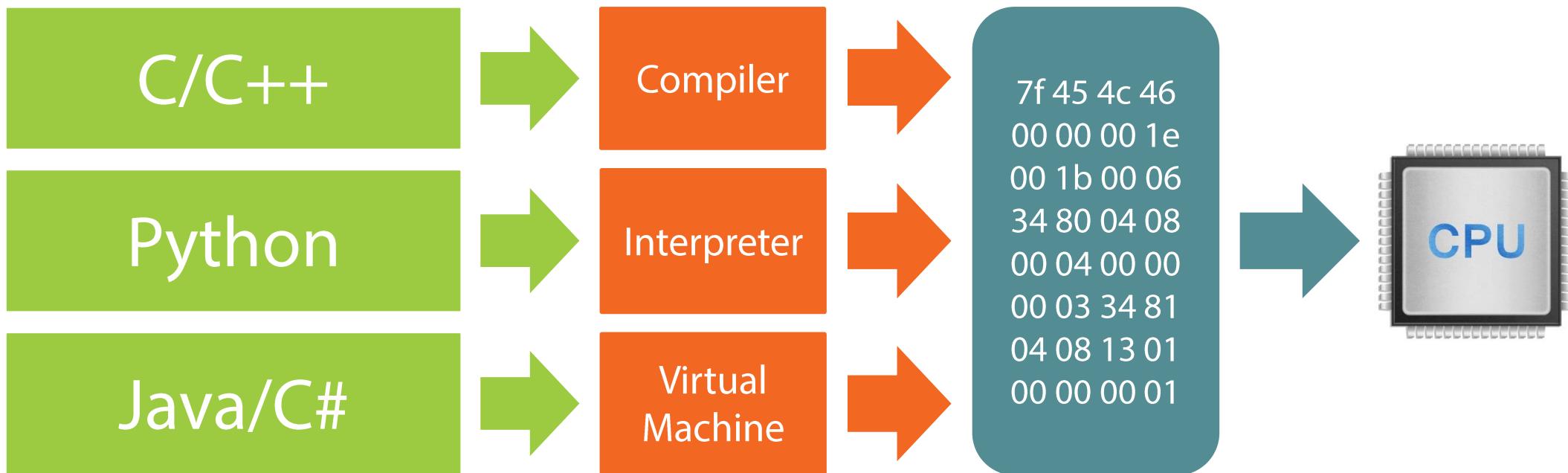


# From Disk to RAM to CPU

```
using System;
public class hello
{
    public static int
    Main(string[] args)
    {
        Console.WriteLine(
            "Hello {0}!",
            System.AppDomain.Current
            Domain.FriendlyName);
        return 0;
    }
}
```



# From Program to Process



# From Program to Process

Windows Task Manager

File Options View Help

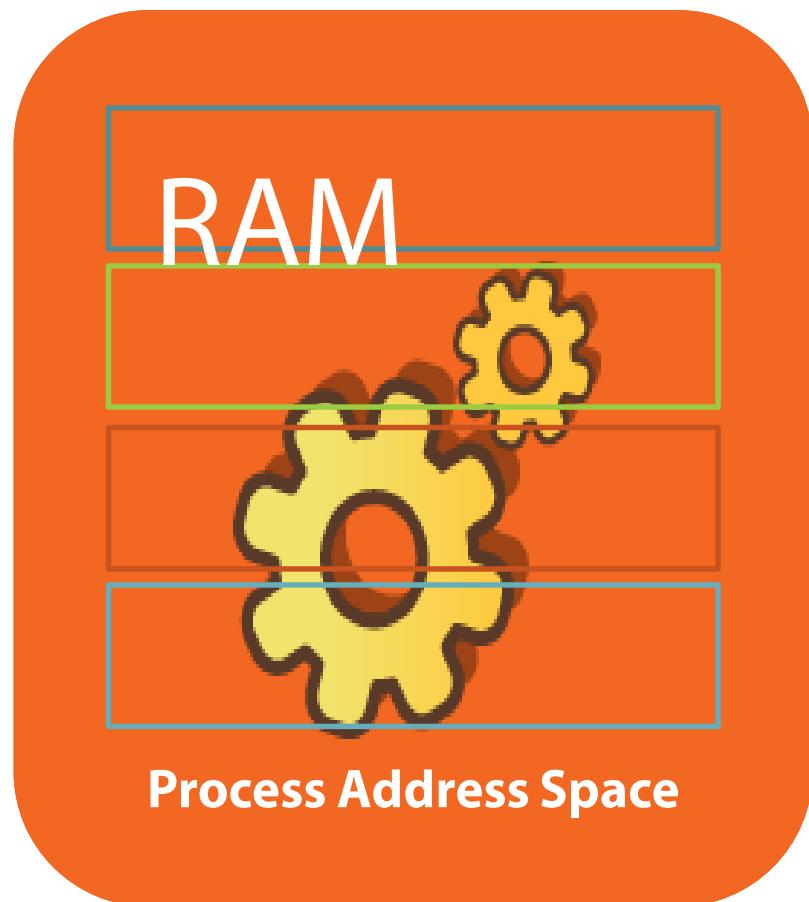
Applications Processes Services Performance Networking Users

Image Name	Description	PID	CPU	Memory...
atiedx.exe		1696	00	1,104 K
avpui.exe *32	Kaspersky Anti-Virus	3080	00	2,260 K
BrCcUxSys.exe *32	ControlCenter UX System	520	00	496 K
BrCtrlCntr.exe *32	ControlCenter Main Process	4356	00	316 K
CarboniteUI.exe *32	Carbonite User Interface	4772	00	9,348 K
CCC.exe	Catalyst Control Center: Host application	4240	00	3,788 K
chrome.exe *32	Google Chrome	7612	00	48,296 K
chrome.exe *32	Google Chrome	7764	00	19,484 K
chrome.exe *32	Google Chrome	9336	00	7,912 K
chrome.exe *32	Google Chrome	9532	00	12,572 K
chrome.exe *32	Google Chrome	10960	00	11,924 K
chrome.exe *32	Google Chrome	11660	00	18,516 K
chrome.exe *32	Google Chrome	12744	00	35,700 K
conhost.exe	Console Window Host	4520	00	344 K
conhost.exe	Console Window Host	6816	00	1,888 K
csrss.exe		876	00	19,684 K
dwm.exe	Desktop Window Manager	3172	00	1,732 K
EMET_Agent.exe	EMET_Agent	3552	00	2,904 K
explorer.exe	Windows Explorer	5360	00	107,628 K

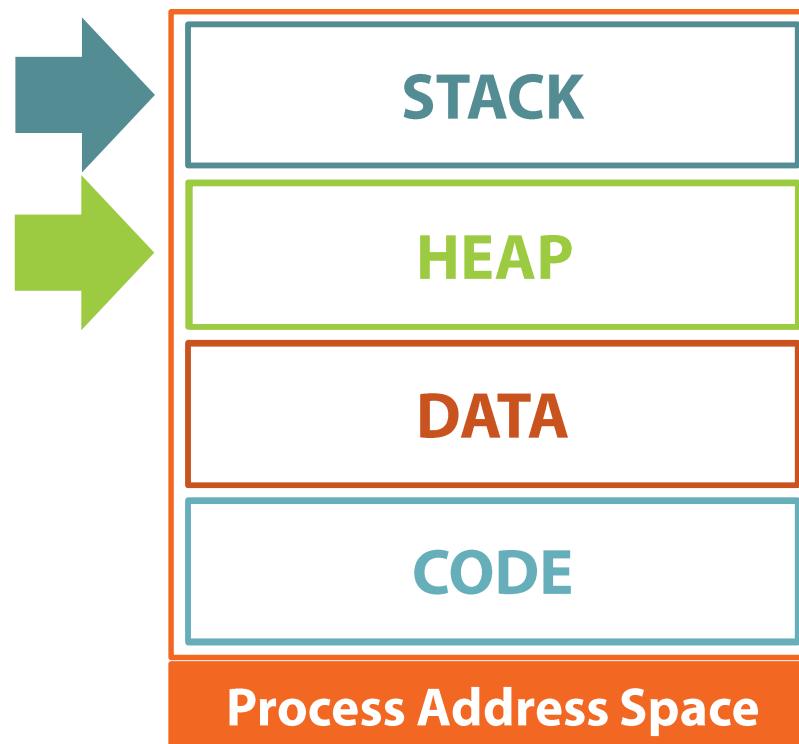
```
root@kali:~# ps -eH
```

PID	TTY	TIME	CMD
2	?	00:00:00	kthreadd
3	?	00:00:23	ksoftirqd/0
5	?	00:00:00	kworker/0:0H
7	?	00:01:40	rcu_sched
8	?	00:00:00	rcu_bh
9	?	00:00:00	migration/0
10	?	00:00:01	watchdog/0
11	?	00:00:00	khelper
12	?	00:00:00	kdevtmpfs
13	?	00:00:00	netns
14	?	00:00:00	perf
15	?	00:00:00	khungtaskd
16	?	00:00:00	writeback
17	?	00:00:00	ksmd
19	?	00:00:00	khugepaged
20	?	00:00:00	crypto

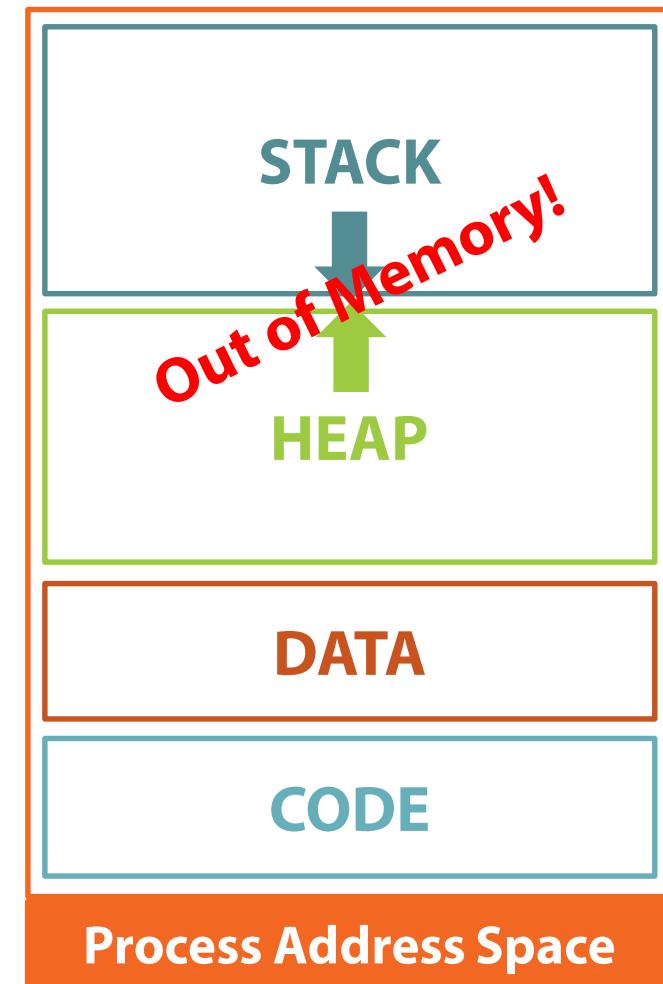
# Inside Process Memory



# Inside Process Memory

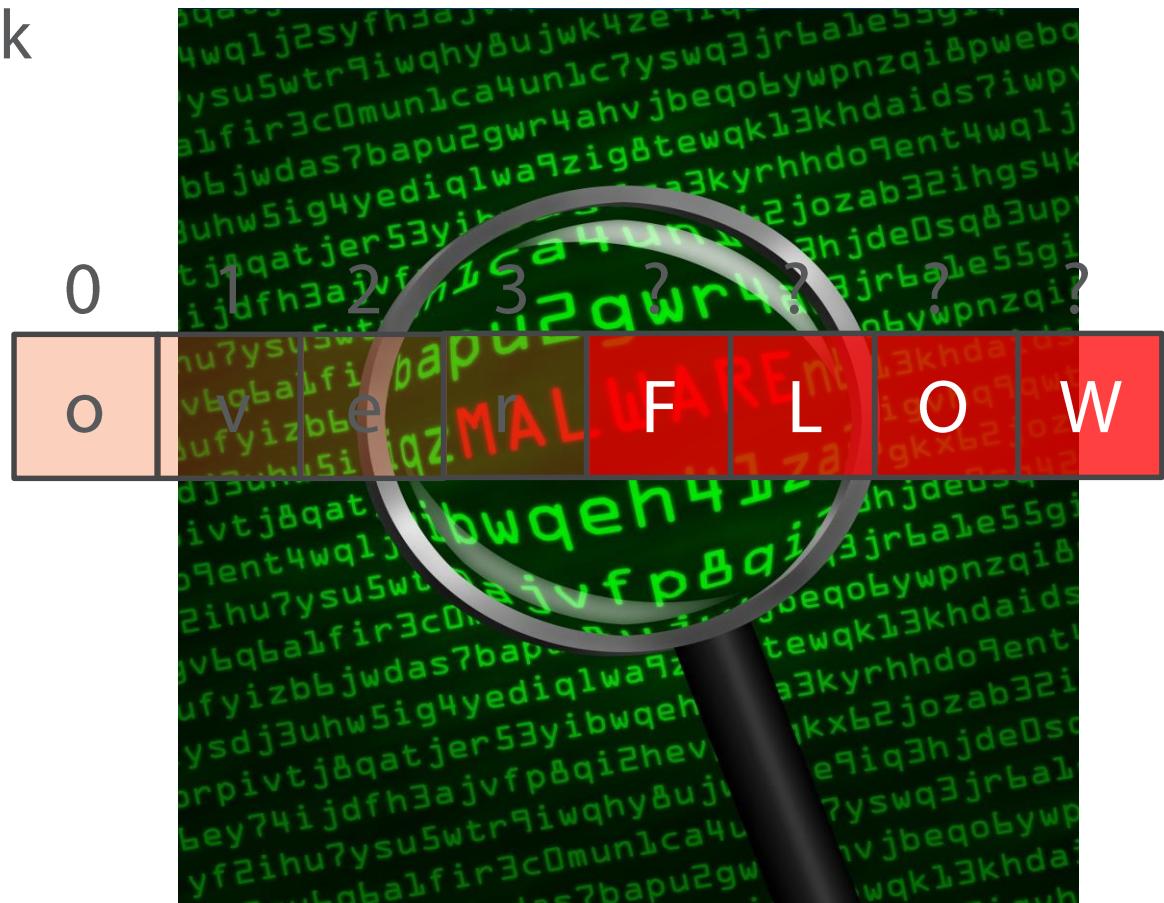


# Fixed and Dynamic Memory



# Inside the Stack

- Overflows are typically in the stack
- Stack overflow causes:
  - Process crash
  - Data leakage/exfiltration
  - Arbitrary code execution



# Just Waiting in Another Queue

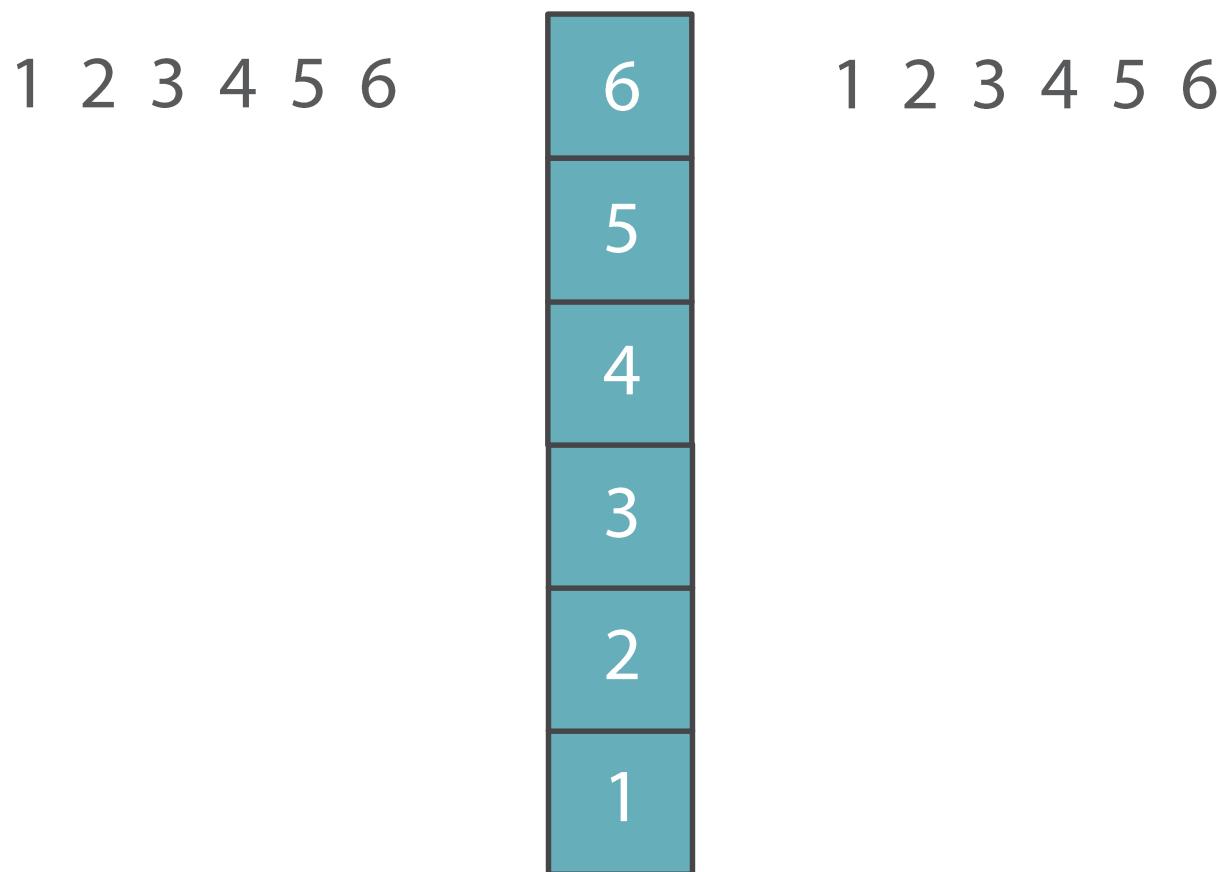
- A stack is a memory buffer
- Temporarily stores data
- Data may be in use now or used later
- Queue data structure
- Both people and data wait in queues



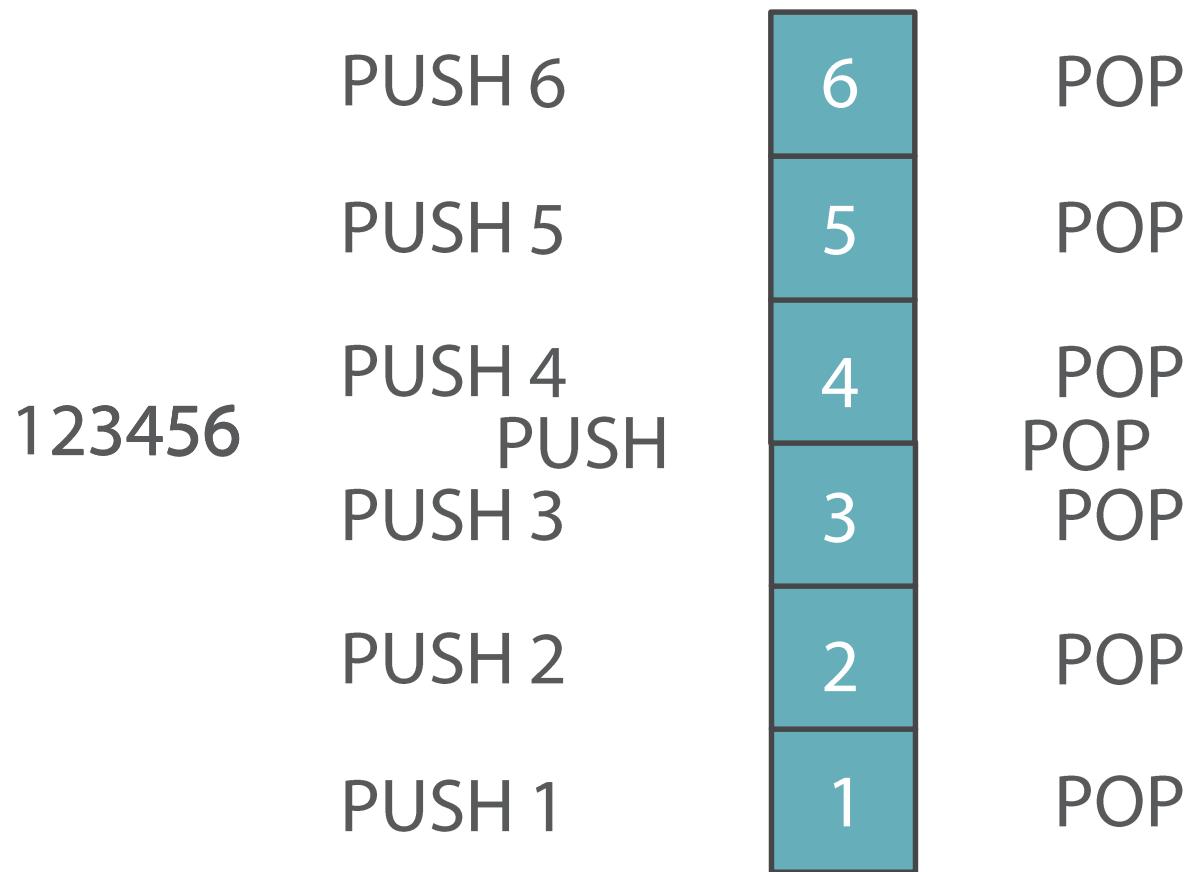
# LIFO or FILO--It's Your Choice



# Value and Order Make Unique



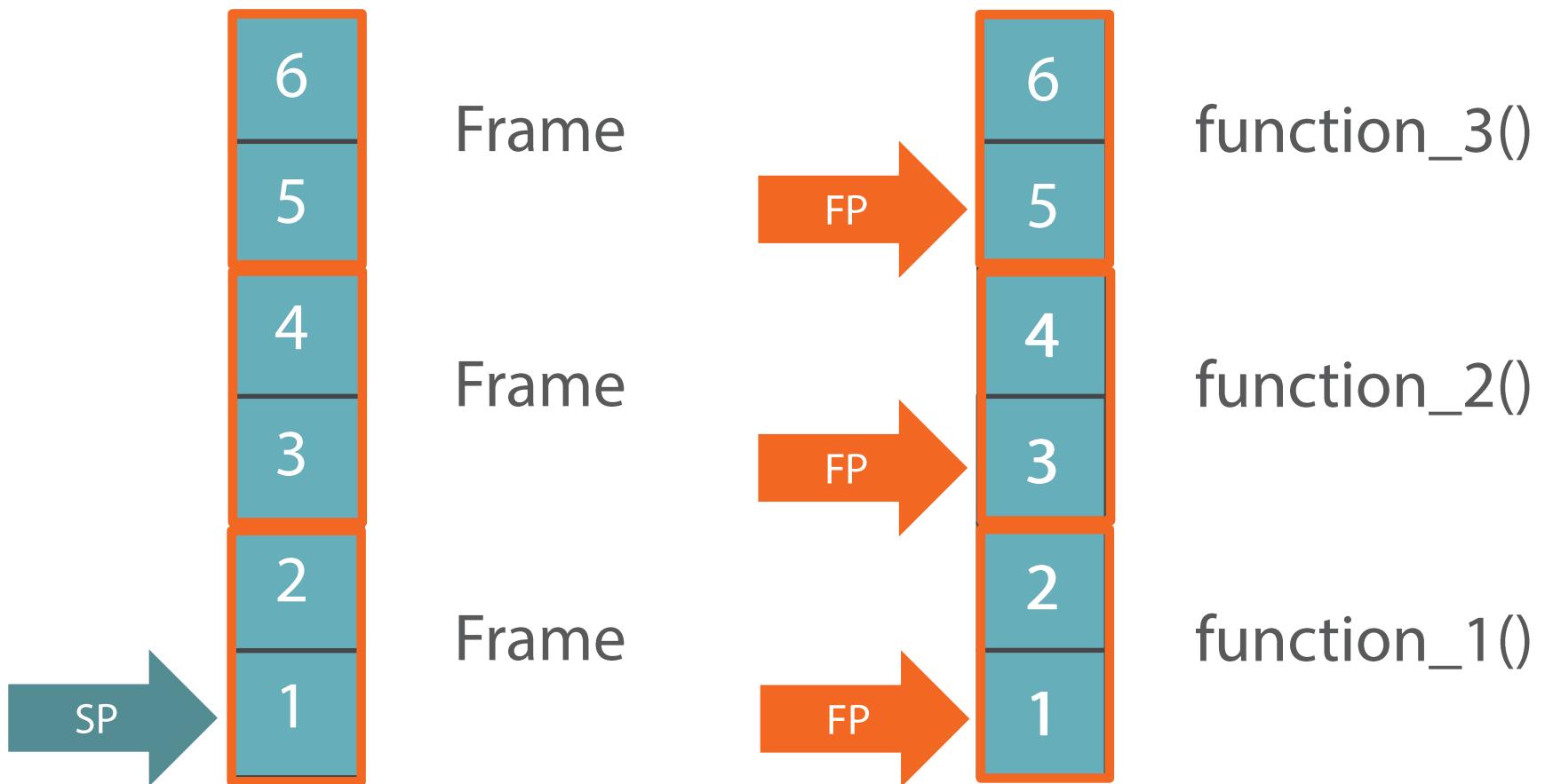
# PUSH and POP



# The Stack Pointer

- Push and Pop move data?  
 $SP = 0x1008$
  - Push copies data to the stack
  - Pop reads data from the stack
  - Data pushed and popped also stays at its source
- DATA12
- 
- POPs moving data from stack to memory:
- |   |        |
|---|--------|
| 2 | 0x1028 |
| 1 | 0x1020 |
| A | 0x1018 |
| T | 0x1010 |
| A | 0x1008 |
| D | 0x1000 |
- DATA12

# Stack Frames



# The Stack in Code

```
int main()
{
    function_1(); int function_1()
    return 0; {
        int i = 42;
        function_2(i); int function_2(int i)
        return 1; {
    }
    return 2;
}
```

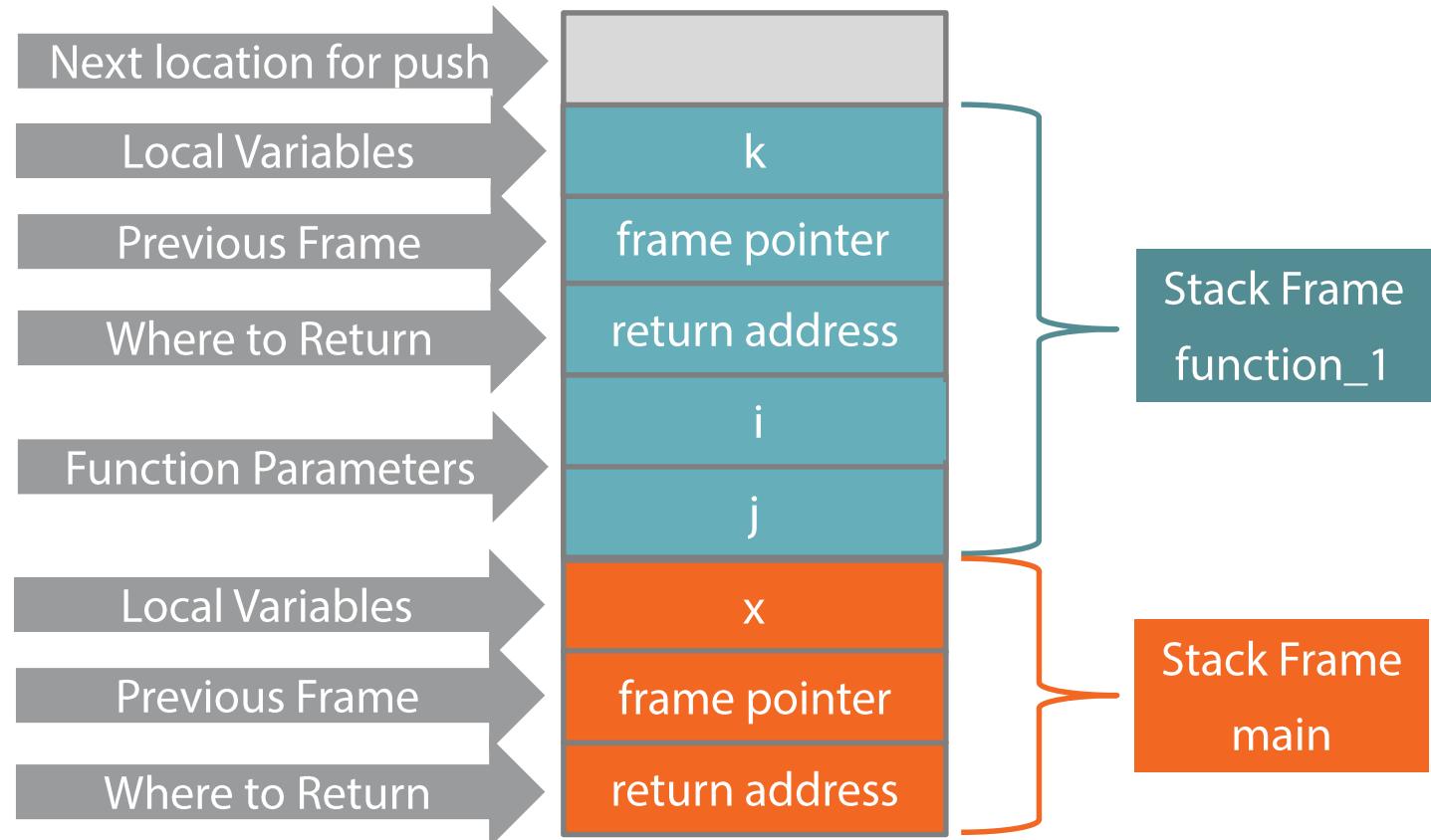
# Inside Stack Frames

- A stack frame is...
  - ...created each time a function is called.
  - ...a logical grouping of information associated with a function.
  - ...used by the program to find information about the current function.
  - ...a way to save function information that must be used later.
  - ... why functions do not overwrite the data of other functions.

# Inside Stack Frames

```
int main()
{
    int x = function_1(33, 42);
    return 0;
}

int function_1(int i, int j)
{
    int k = i + j;
    return k;
}
```



# Pointers

- A “pointer” is a data variable that stores a memory address
  - a location in the computer’s memory
  - C and C++ use pointers to locate and use objects in memory
  - A pointer contains the memory address of the object, not the object itself
  - The address of my house is its location and not the house itself
- The data type of the pointer indicates the type of data stored at that location
  - The data type determines how the memory will be read and written

# Stack Pointers

SP  
Stack Pointer

FP  
Frame Pointer

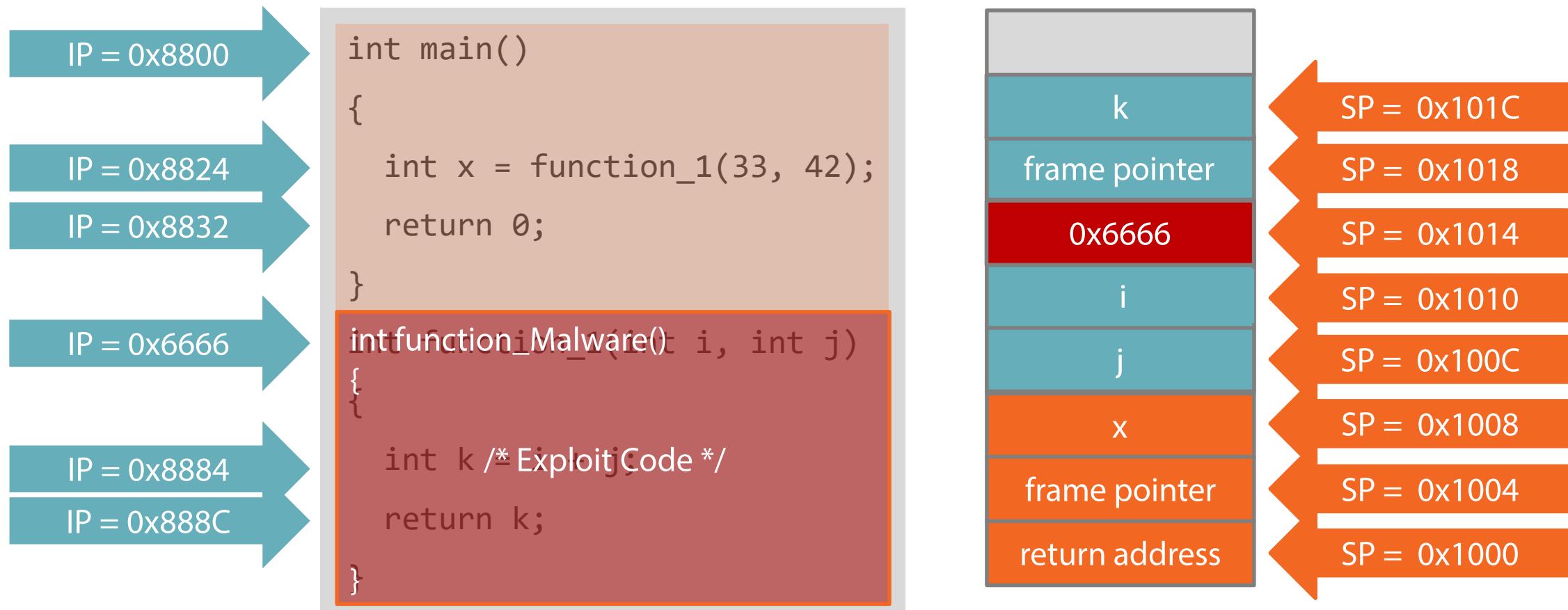
BP  
Base Pointer

IP  
Instruction Pointer

Instruction Counter  
or  
Program Counter

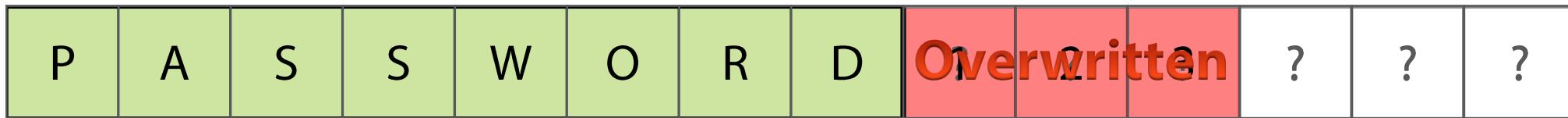
Registers are  
dedicated areas of  
CPU memory

# The Instruction Pointer

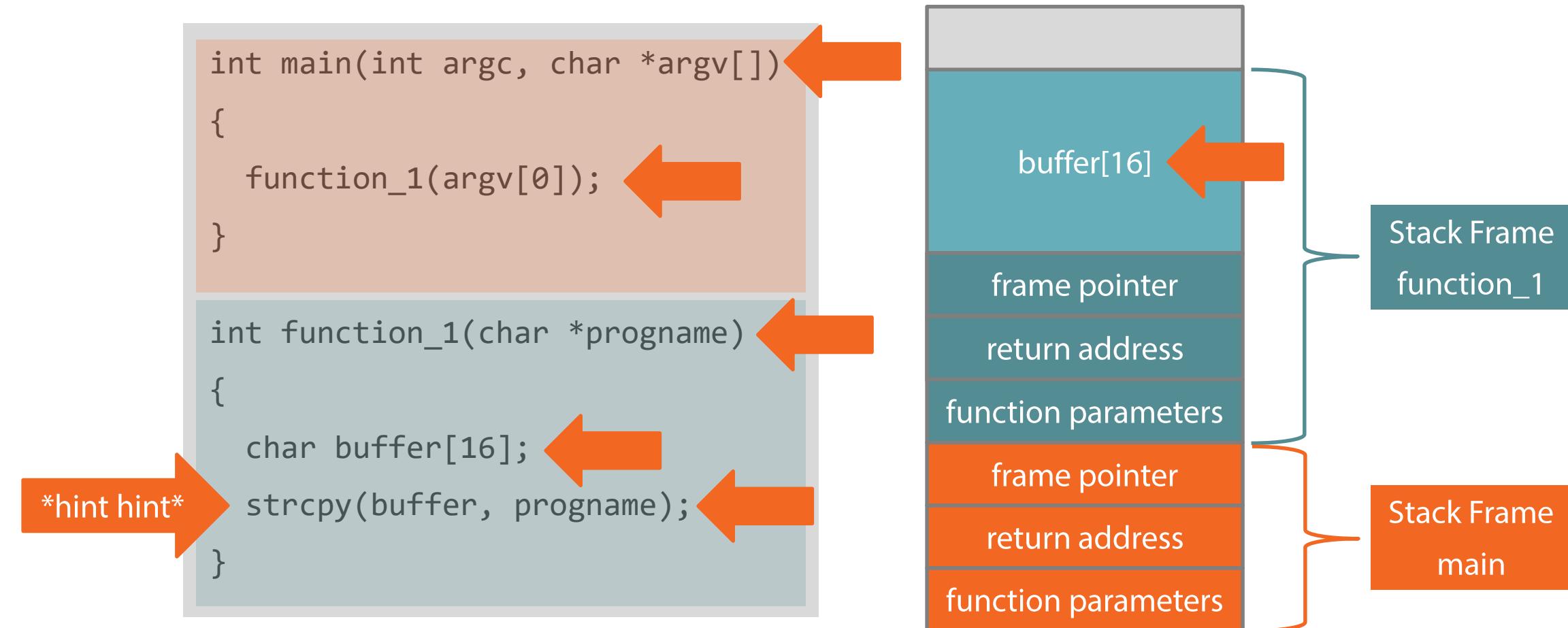


# Overflowing the Stack

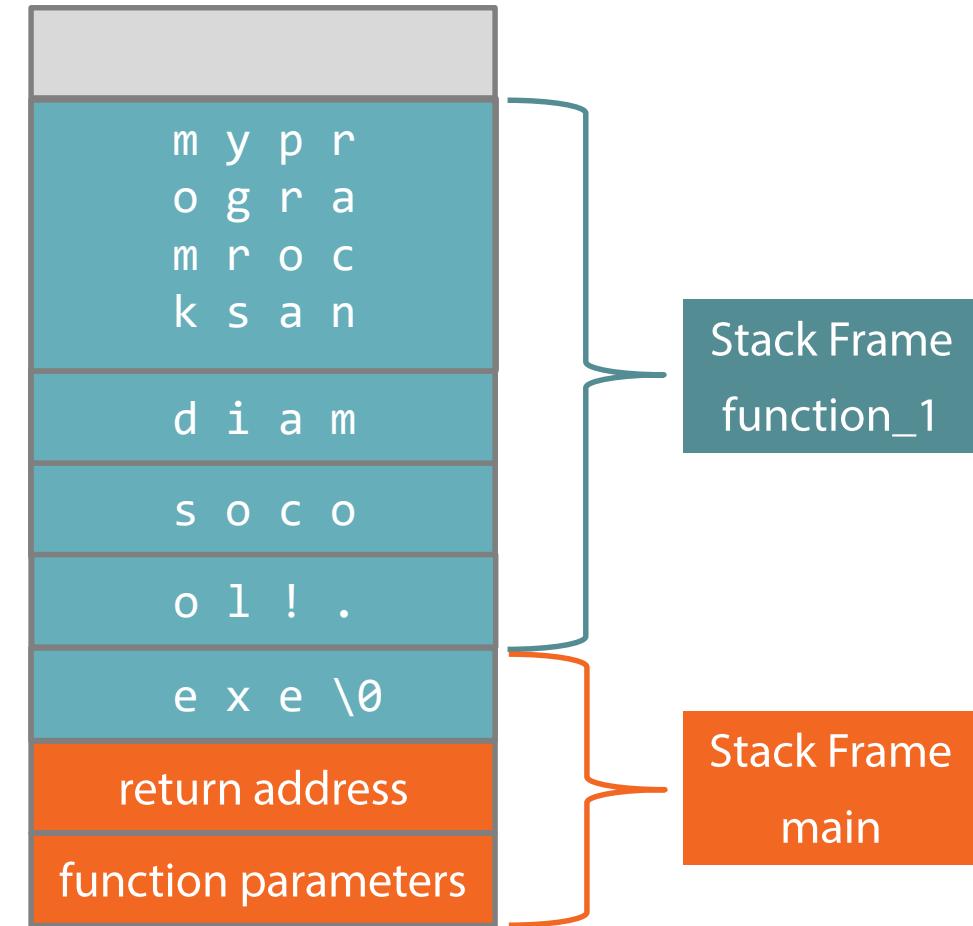
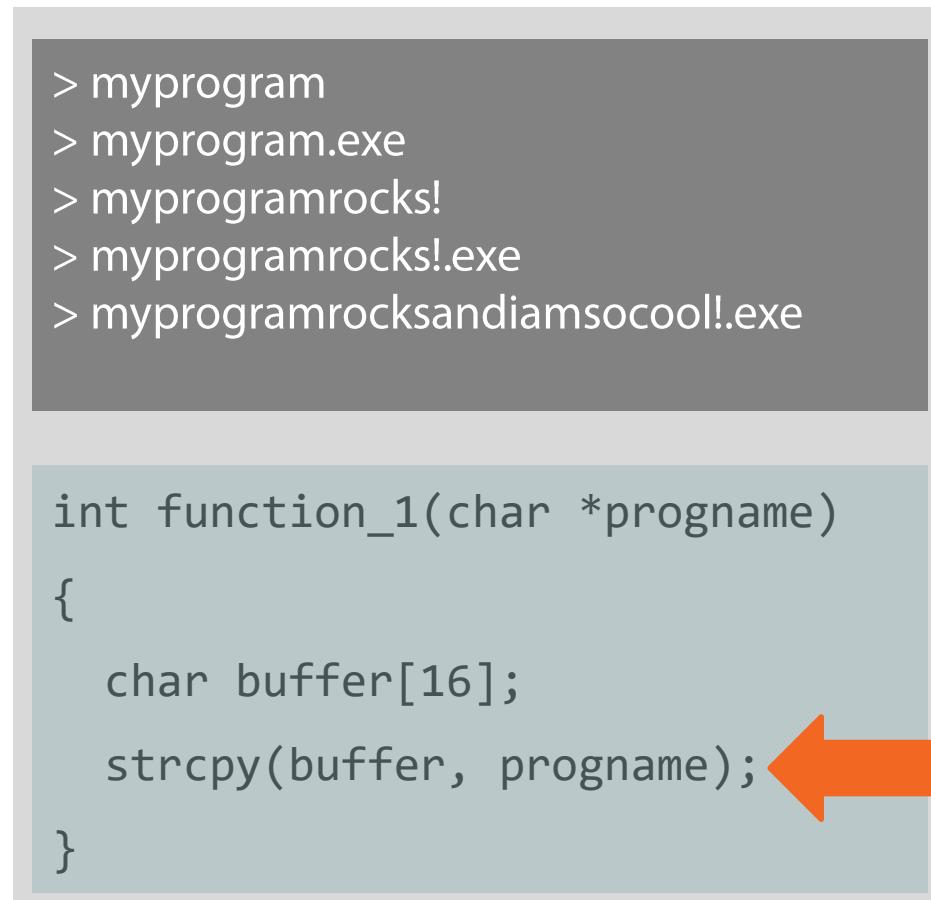
- An overflow occurs when...
  - data written to a buffer exceeds the boundaries of the buffer
  - the length of the data written is longer than the buffer
  - the extra data is written past the end of the buffer
  - anything following the buffer in memory is overwritten



# Overflow Inside the Stack



# Overflow Inside the Stack



# Exploiting Stack Overflows

What can you do  
with a buffer  
overflow  
condition?

Morris worm  
Code Red worm  
Many other Malware

Access  
memory-resident  
resources

Run installed  
programs

Inject and run  
exploit code on  
your computer

Program instability  
and abnormal  
termination

# “Exploit Code”

Any executable code that is injected into a program's buffer...

causing the buffer to overflow...

and forcing the program to run the injected code

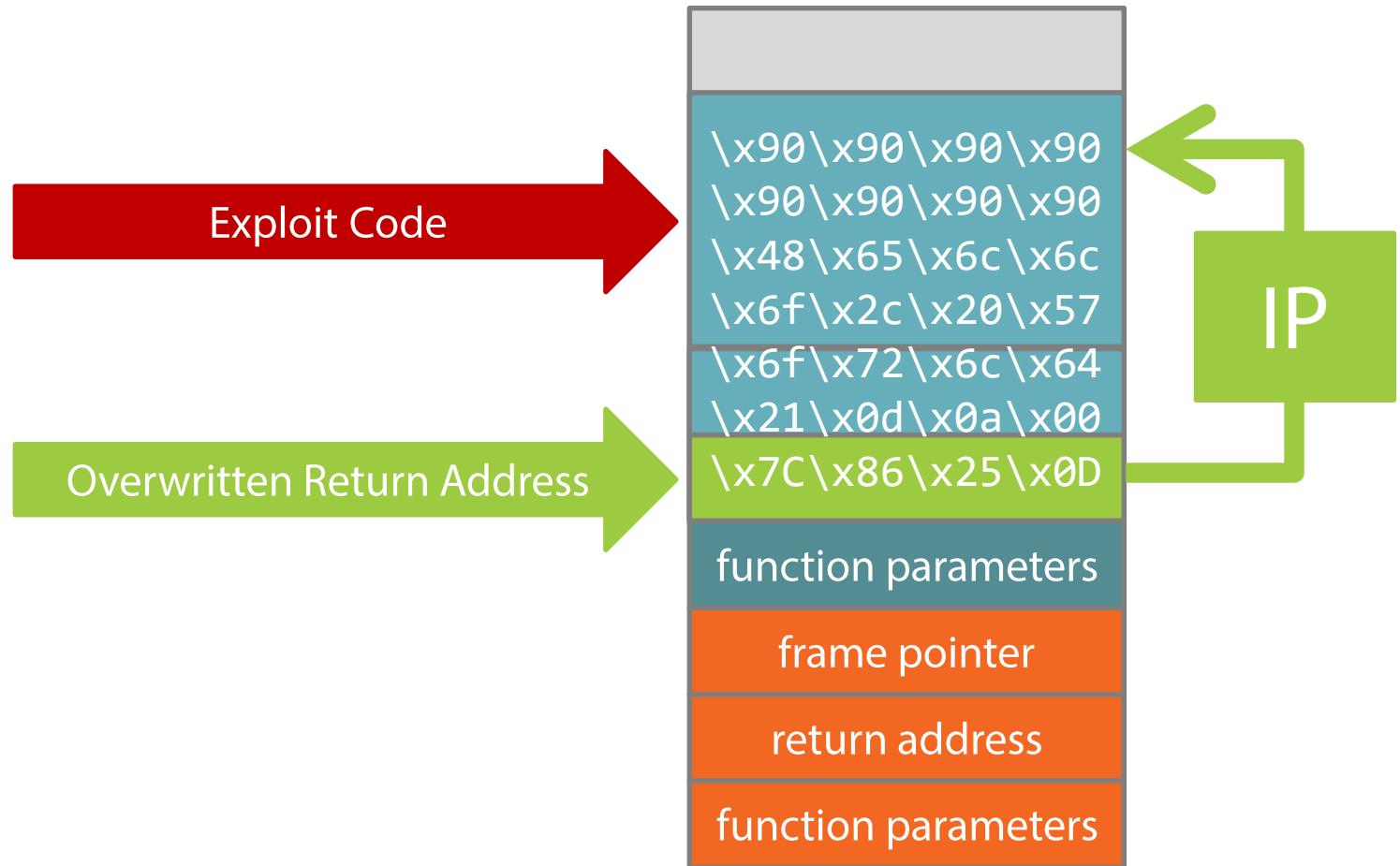
Arbitrary code  
execution

It doesn't matter what the exploit code does...

...just get it into the system and run it!

# Stages of a BoF Exploit

The new exploit does  
overflows the buffer  
and overwrote the  
return address after  
current stack frame



# Predictability Is the Key

- Arbitrary code exploits need a predictable execution environment
- The memory location of needed objects must be known or discoverable
- System resources are mapped into the process' virtual memory space
  - 0x7C86250D : kernel32.dll.WinExec()
  - 0x7C801D7B : kernel32.dll.LoadLibraryA()
  - 0x7C802336 : kernel32.dll.CreateProcessW()
  - 0x7C80236B : kernel32.dll.CreateProcessA()
- All processes have the same resources at the same virtual memory locations
- Exploit code needs these system resources to properly run

# Where Is My Exploit Code?

The location of the exploit code in the stack or heap must be known

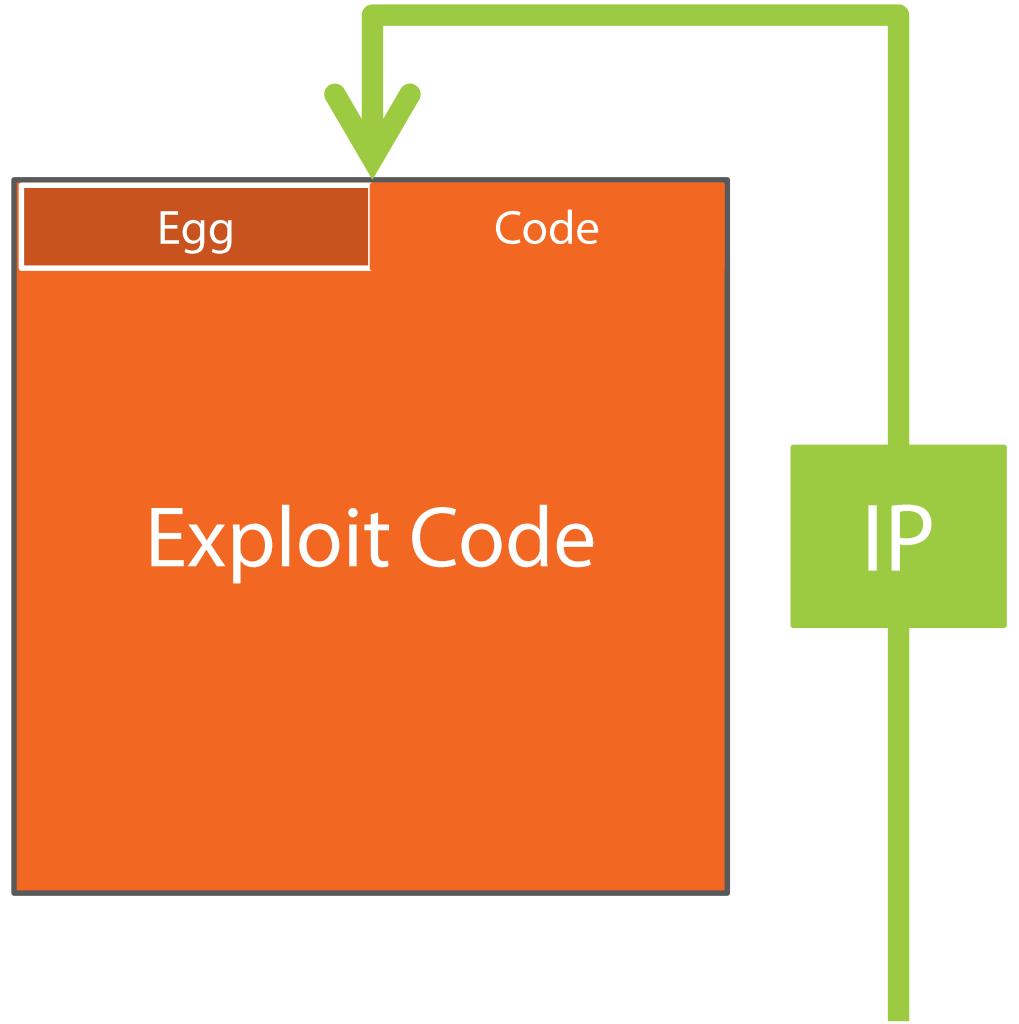
The IP must contain the memory address of the beginning of the exploit code

The exploit code may not be located at a predictable address(es)

Egghunting  
NOP Sled

# Egghunting

- The “egg” is a byte signature used to easily find exploit code in memory
- The egg is located at the beginning of the exploit code
- The egg must be unique in process memory
- The first exploit code instruction is located immediately after the egg
- Search code must be pre-loaded and memory exceptions handled



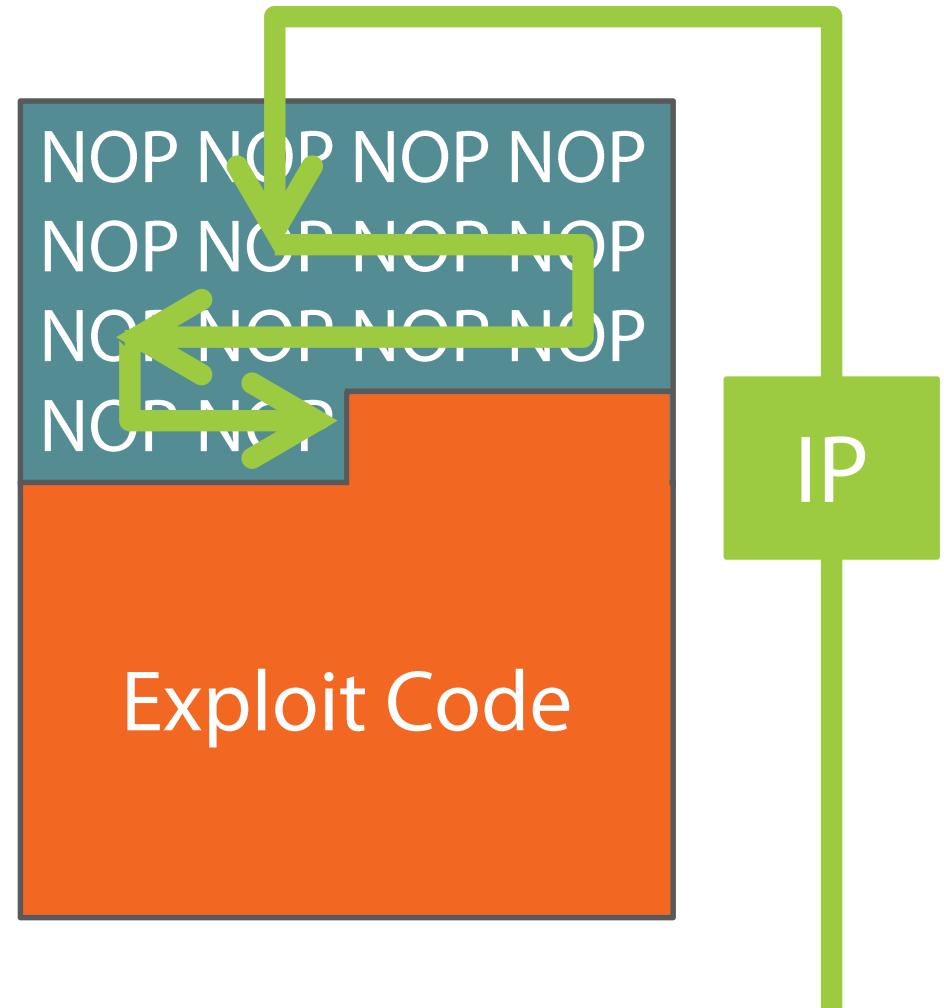
# No Operation

- Instructions command the CPU
- No Operation instruction
- NOP or NO-OP
- Do nothing and advance the IP
- NOP consumes one CPU cycle
- 1 cycle =  $10^{-9}$  seconds on 1 GHz CPU
- NOPs fix timing, synchronization, and code alignment issues

Operations	CPU Instructions
Data Transfer	MOV XCHG BSWAP
Arithmetic	ADD SUB MUL DIV CMP
Logic	AND OR XOR NOT TEST
Stack	POP PUSH ENTER LEAVE
Flow Control	JMP CALL RET INT LOOP
CPU Control	<b>NOP HLT PWRSAV RESET</b>

# NOP Sled

- Set the IP to the address of the beginning of the exploit code
  - This is not always easy or possible
- The IP only needs to point to the NOP sled
- The IP will then move down to and execute the exploit code
- A NOP sled is a sequence of NOP instructions preceding the payload, making the payload easier to find



# Pseudo-NOP Sleds

```
\0x90\0x90\0x90\0x90\0x90\0x90  
\0x90\0x90\0x90\0x90\0x90\0x90  
D\0x90\0x90\0x90\0x90\0x90\0x90  
0\0x90\0x90\0x90\0x90\0x90\0x90  
F\0x90\0x90\0x90\0x90\0x90\0x90  
6\0x90\0x90\0x90\0x90\0x90\0x90  
M\0x90\0x90\0x90\0x90\0x90\0x90  
W\0x90\0x90\0x90\0x90\0x90\0x90  
4\0x90\0x90\0x90\0x90\0x90\0x90  
W\0x90\0x90\0x90\0x90\0x90\0x90  
M\0x90\0x90\0x90\0x90\0x90\0x90  
W\0x90\0x90\0x90\0x90\0x90\0x90  
D\0x90\0x90\0x90\0x90\0x90\0x90  
24\0x90\0x90\0x90\0x90\0x90\0x90
```

Long runs of NOP instructions are easily detectable as a NOP sled

Long runs of nonsensical CPU instructions are not (easily) detectable as a NOP sled

# Shellcode

- Shellcode
  - creates an interactive OS shell
  - executes commands on the local system via command line, script, or batch file
  - /bin/sh, /bin/bash, %WINDIR%\System32\cmd.exe, explorer.exe, command.com
- Shellcode can be any type of program
  - Servers (FTP, Telnet, SSH, IRC, HTTP), monitors (sniffers, keyloggers)
- Shellcode runs with the privileges and context of the process that runs it
- Attacker: "Give me a shell with root user privileges, now!"

# Shellcode Types

- Port Binding
- Reverse
- Command Execution Code
- File Transfer
- Find Socket
- Kernel Space
- Multistage
- Process Injection
- System Call Proxy



# Creating Shellcode

- Creating shellcode
  1. Writing the shellcode in assembly language (CPU-specific)
  2. Assembling and linking the assembly code to create machine code
  3. Converting the machine code to an escaped-byte text format
- Shellcode code can be written in higher-level languages
- ...but small size is important, so stick with assembly language.

```
\xe9\x1e\x00\x00\x00\xb8
\x04\x00\x00\x00\xbb\x01
\x00\x00\x00\x59\xba\x0f
\x00\x00\x00\xcd\x80\xb8
\x01\x00\x00\x00\xbb\x00
\x00\x00\x00\xcd\x80\xe8
\xdd\xff\xff\xff\x48\x65
\x6c\x6c\x6f\x2c\x20\x57
\x6f\x72\x6c\x64\x21\x0d
```

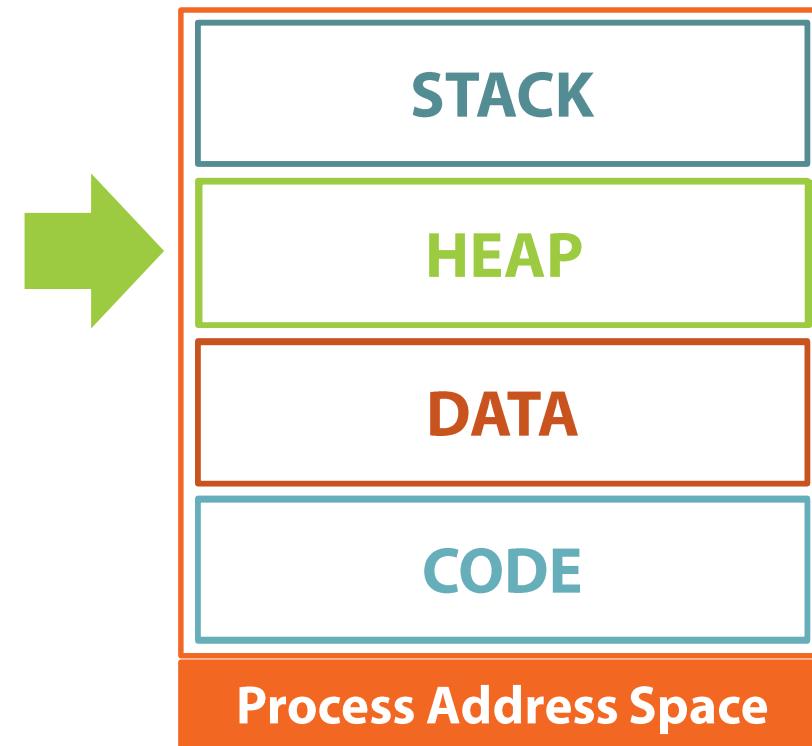
# Testing Shellcode

```
unsigned char shellcode[] = "Escaped byte codes";
int main()
{
    int func = (void (*)())shellcode;
    (*func)();
    func = (void (*)())shellcode;
    (*func)();
}
```

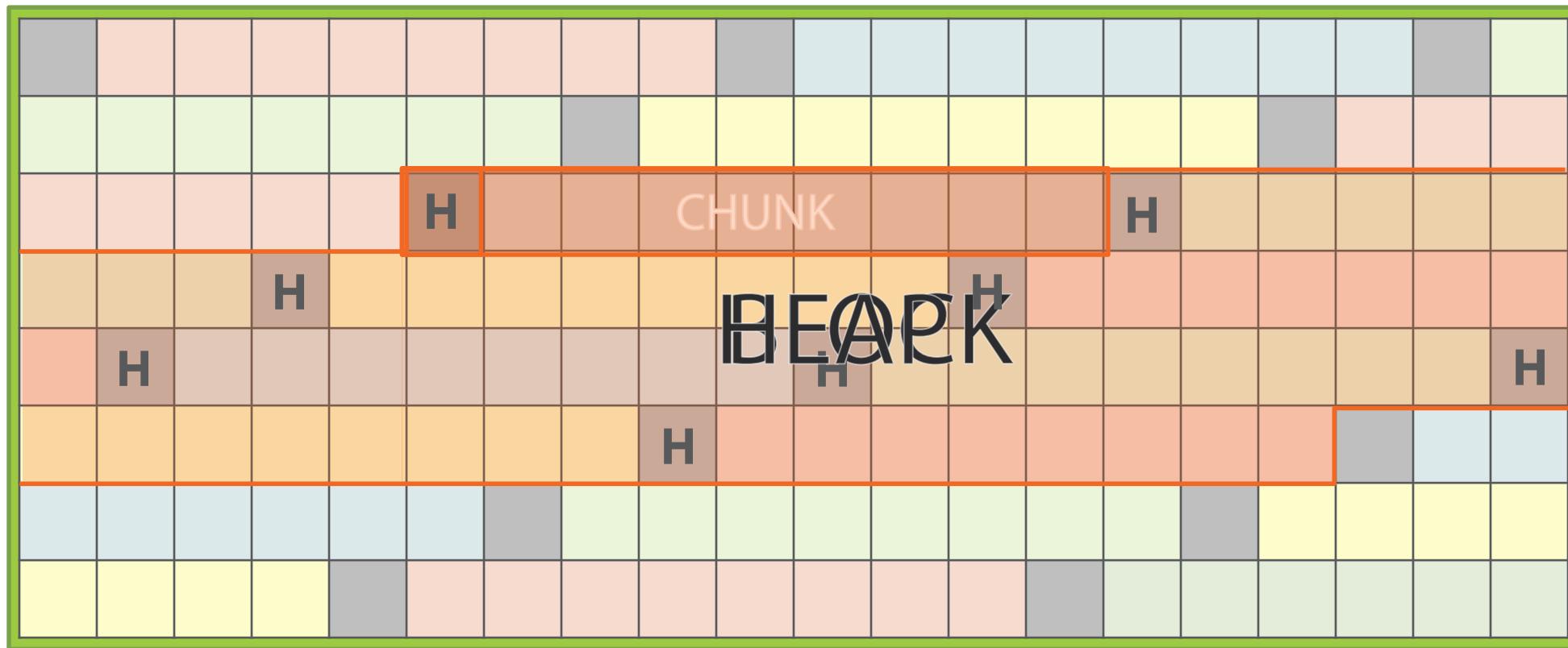
```
gcc -g -Wall -fno-stack-protector -z execstack shelltest.c -o shelltest
```

# Inside The Heap

- Heap stores program data
- The stack is fixed in size, but the heap can grow and shrink as needed
- Dynamic memory
  - allocates and releases storage space as the process needs
  - used to temporarily store data for processing
- The heap can also store and execute code, which is just a type of data



# Structure of The Heap



# Heap Chunks



# Using the Heap in Code

```
// Buffer allocated on stack only requires a variable declaration  
int bufferOnStack[20];  
  
// Buffer allocated on heap requires a function call  
int *bufferOnHeap = (int *)malloc(20);  
  
int* bufferOnHeap2 = new int[20];  
  
int *bufferOnHeap3 = (int *) HeapAlloc(hHeap, 8, 20*sizeof(int));
```

# Using the Heap in Code

```
// Releasing buffer on stack only requires exiting function block
{
    int bufferOnStack[20];
}

// Releasing a buffer allocated on heap requires a function call
free(bufferOnHeap);           bufferOnHeap = 0;
delete bufferOnHeap2;         bufferOnHeap2 = 0;
HeapFree(bufferOnHeap3);       bufferOnHeap3 = 0;
```

# Exploiting Heap Overflows

- Buffers on the heap can be overflowed just as stack buffers can
- Changes the memory addresses stored in pointer variables
  - Redirect the flow of program execution
  - Cause the process exception handler to be called
  - Crash the running process
- Code stored in the heap can be executed
  - Redirect the Instruction Pointer to the code
  - Find a heap buffer large enough for the exploit code
  - Make “fake” heap memory management structures

# Is a Buffer on the Stack or Heap?

- Important to know for writing BoF
- Look in the source code or use static/dynamic analysis
- How big is the buffer and what is it used for?
- Heap buffers are needed when:
  - very large data must be stored
  - the buffer must exist for a long time
  - the size of the buffer needs to change
  - the size of the buffer is only known at run-time

# Heap Overflow in Code

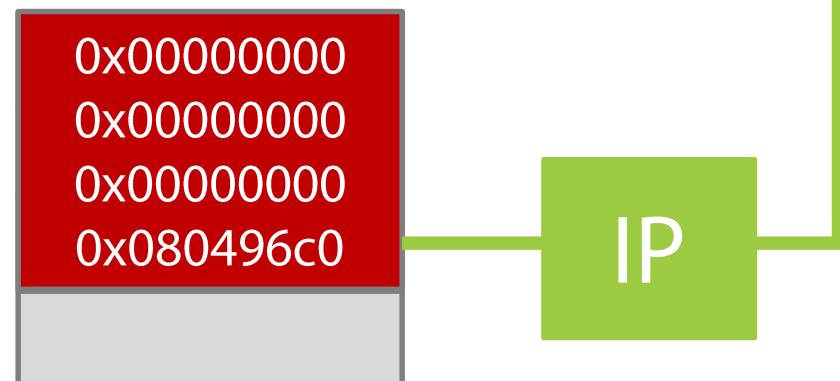
```
int main(int argc, char *argv[])
{
    char *buffer = (char*)malloc(16);
    strcpy(buffer, argv[1]);
    free(buffer);
}
```

HEAP

\x52\x4f\x21\x0a  
\x31\xc0\x50\x68  
\x78\x79\x6f\x75  
\x68\x61\x62\x72  
\x6f\x89\xe1\x6a  
\x08\x5a\x31\xdb  
\x43\x6a\x04\x58  
\xcd\x80\x6a\x17  
\x01\x01\x50\xeb  
\x12\x4c\x45\x20

# Heap Spray

- Increases the chance of finding exploit code written to a heap buffer
- Writes identical blocks of NOP sleds and exploit code to the buffer
- BoF in stack overwrites IP and redirect flow of execution to NOP sled in buffer



# Function Pointers

A variable that stores the location of a function

Functions are usually called by name in code

But which function to call in what situation?

Call a function by its memory address, not by its name

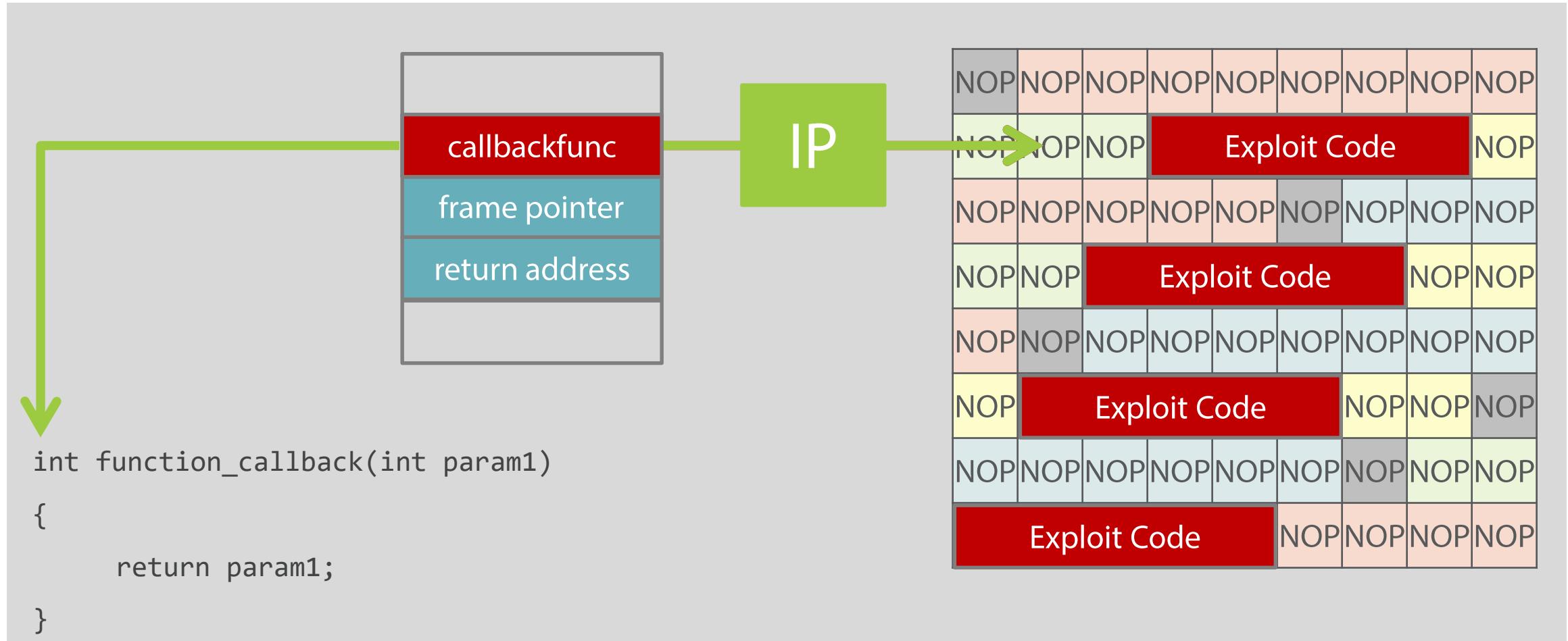
Decide which function to call while a program is running

Functions called by name must be hard-coded when the software is written

# Function Pointers in Code

```
int main()
{
    int (*callbackfunc)(int); ←
    callbackfunc = (int (*)())function_callback; ←
    int result = (*callbackfunc)(42); ←
}
int function_callback(int param1) ←
{
    return param1;
}
```

# Function Pointer Overwrite

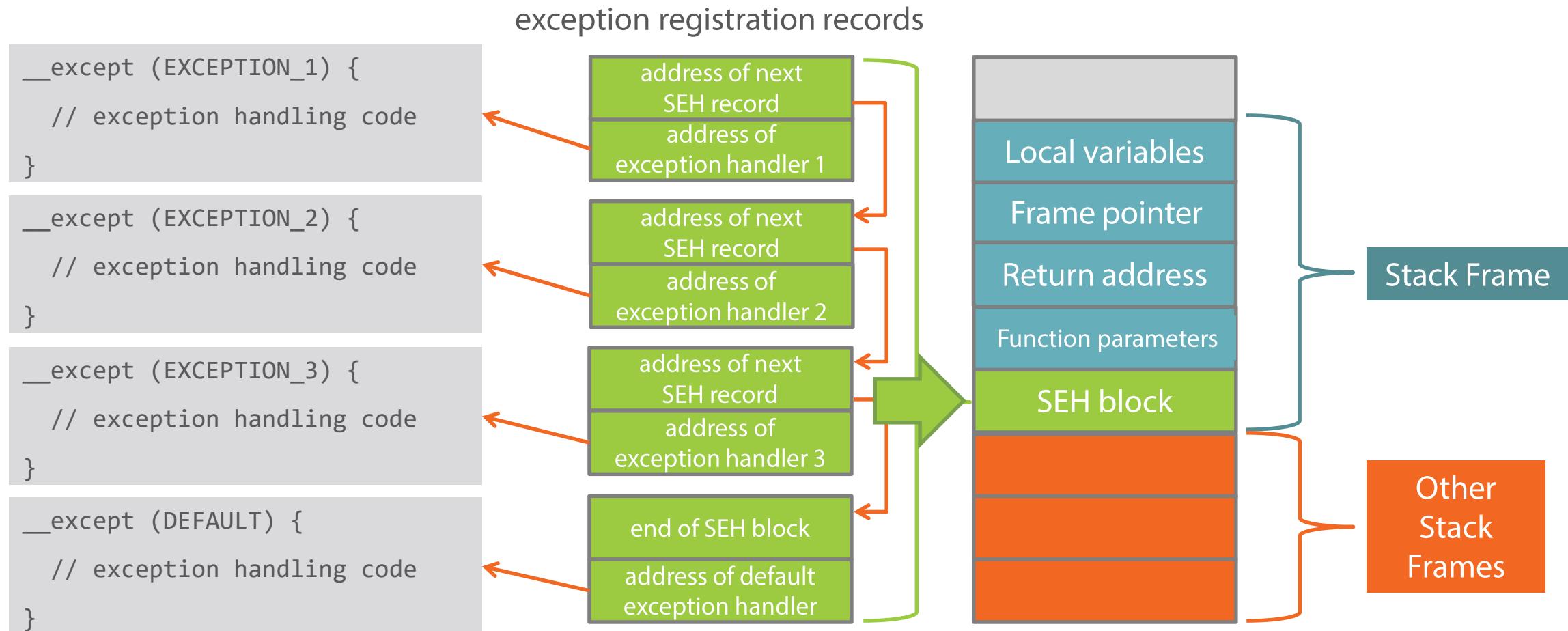


# Structured Exception Handling

- Windows OS error recovery feature
- Notification of the occurrence of error conditions
- Allows process to handle errors rather than crashing
- `__try`, `__except`, `__leave`, `__finally`, `RaiseException`
- All Window processes have SEH, not just for those written in C++

```
int main()
{
    __try
    {
        TestExceptions();
    }
    __except(EXCEPTION_EXECUTE_HANDLER)
    {
        puts("Executing SEH __except block");
    }
    return 0;
}
```

# SEH in the Stack



# Exploiting SEH

BoF authors have  
software problems  
to solve too

How to execute  
exploit code after  
it is injected?

How to keep  
exploit code  
from crashing  
the process?

Use SEH!

# CodeRed's SEH Exploit

# Inside Integer Overflows

Integer overflows  
are arithmetic  
errors that occur

when adding or  
multiplying two or  
more integers

results in a value  
too large to store

An unexpected  
value will result

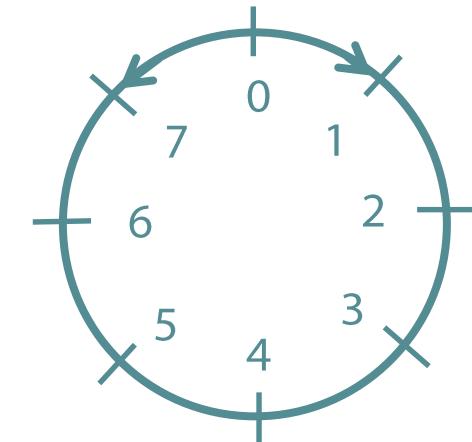
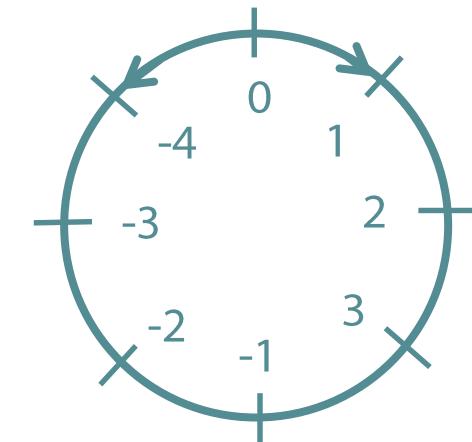
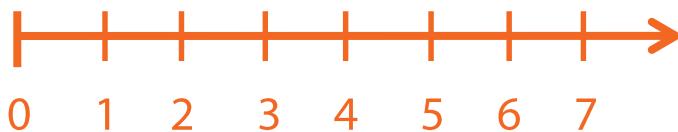
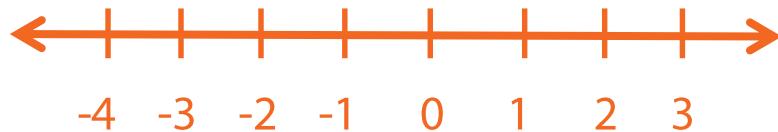
If integer overflows  
are not detected  
and corrected

problems such as  
buffer overflows  
could result

# Integer Sizes

Variable Data Type	Size	Range
signed char	8 bits	-128 to 127
unsigned char	8 bits	0 to 255
signed short int	16 bits	-32,768 to 32,767
unsigned short int	16 bits	0 to 65,535
signed long int	32 bits	-2,147,483,648 to 2,147,483,647
unsigned long int	32 bits	0 to 4,294,967,295
signed long long int	64 bits	-9,223,372,036,854,775,808 to 9,223,372,036,854,775,807
unsigned long long int	64 bits	0 to 18,446,744,073,709,551,615

# It's Loops, Not Infinity



# Overflowing an Integer

```
// unsigned char stores 0 to 255
```

```
unsigned char value = 255;  
unsigned char result = value + 1;
```

```
// result will equal zero, not 256
```

$$\boxed{255} + \boxed{1} = \boxed{0}$$

```
// signed char stores -128 to 127
```

```
signed char value = -128;  
signed char result = value - 1;
```

```
// result will equal 127, not -129
```

$$\boxed{-128} - \boxed{1} = \boxed{127}$$

# Problem or Business As Usual?

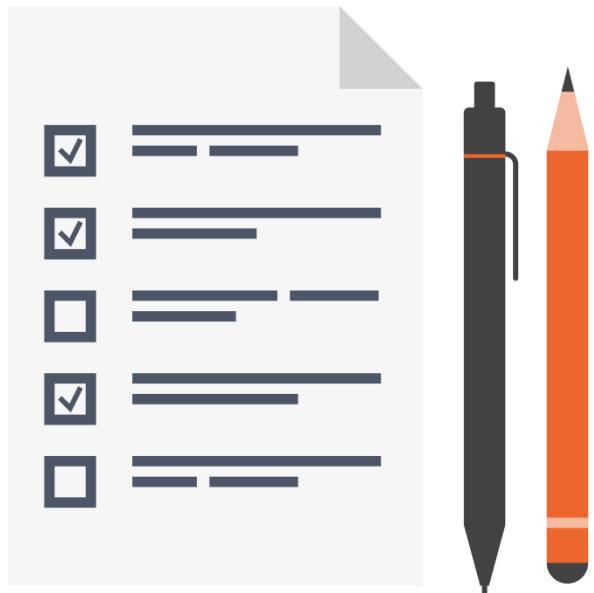
Integer overflows  
are normal and to  
be expected

Programmers must  
catch and correct in  
their code

Overflows values  
result in bad  
program decisions

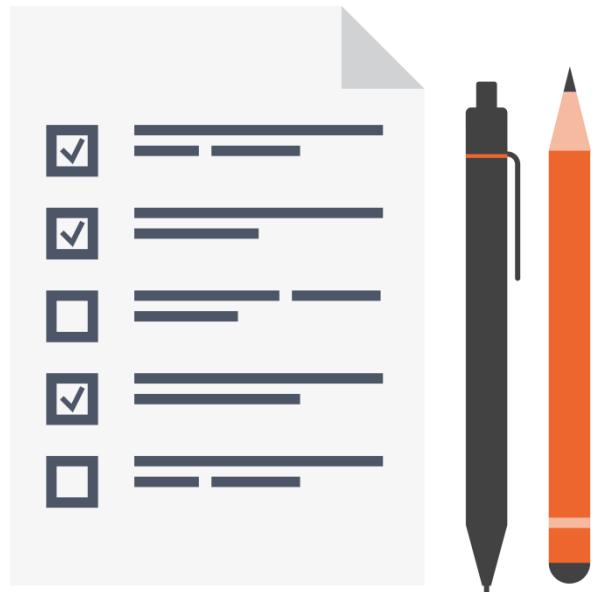
Reduce your attack  
surface and patch  
what remains

# Summary



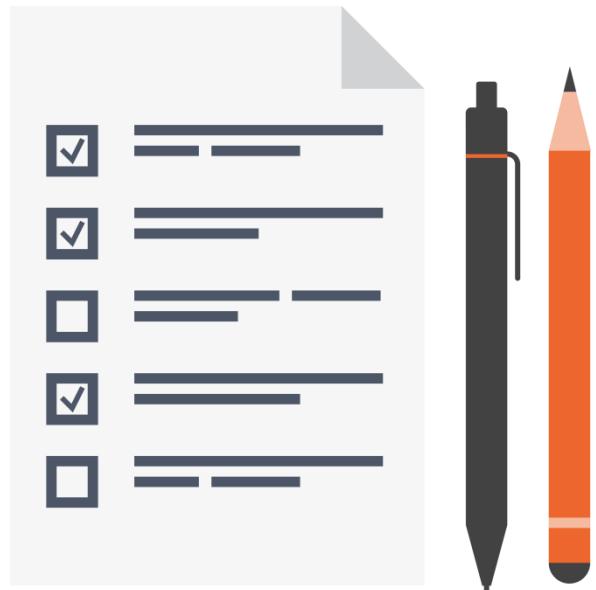
- ✓ What's inside of a computer program?
- ✓ How code become executable programs
- ✓ What happens when program is executed
- ✓ Organization of process memory

# Summary



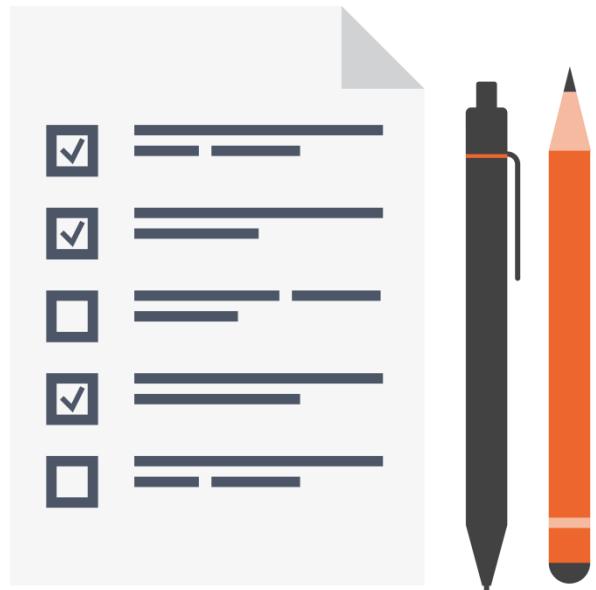
- ✓ What the stack is and how it works
- ✓ What happens when a stack overflow occurs
- ✓ How stack overflows are exploited
- ✓ What is shellcode, egghunting, NOP sleds

# Summary



- ✓ What is heap memory
- ✓ How heap buffers can be overflowed
- ✓ How heap overflows can be exploited
- ✓ Heap spray technique

# Summary



- ✓ How function pointers can be exploited
- ✓ Structured Exception Handling (SEH)
- ✓ Integer overflows



The last good thing written in C was  
Franz Schubert's Symphony number 9.

— Erwin Dieterich



Next Up:

Finding Buffer Overflows