

Turing Algorithm for Prime Numbers and Factorization

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Abstract

This paper introduces a new algorithm to represent a whole set of numbers in one binary number representation and we use algorithm to fetch a list of Prime numbers with an execution time related to only the number of digits in the number. Some other applications for this algorithm in number theory will be factoring a number or checking if a number is prime or not.

Keywords: Prime numbers; Factorization.

1. Introduction

1.1 Binary Turing representation.

In classic computers the limitation of data types ranges makes it hard to represent decimal numbers in binary representations and doing binary operations and goes back to decimal numbers after this operations even for small numbers.

In this paper we are going to introduce new algorithm to represent numbers into binary turing layout for each number and its multipliers.

We are relying on one note that each number its multipliers are exists on the same step a way

For Example: - for number 7 and its multipliers

$$\{7, 14, 21, 28, 35, 42, 49, 56, 63, 70, 77, \dots\}$$

The difference all the time between each two consecutive numbers = 7 as those are multipliers of 7

Same will be for any other number

$$\{11, 22, 33, 44, 55, 66, 77, 88, 99, \dots\}$$

So we are going to represent each number and its multipliers in order until a specific end length or cutoff number(lentgth)

And at each step = N we are going to set this place = 1 and the rest of the digits in the number length will be 0

For example : - if we want to represent decimal system 7 number in this turing binary algorithm for length =100

[illegible]

First 1 in this binary turing number = 7 ; second 1 in this binary turning number = 14 ; third 1 in this turing number = 21 ; fourth 1 in this binary turing number = 28;

[illegible]

[illegible]

1 positions like a skewed diagonal lines with one origin at 1

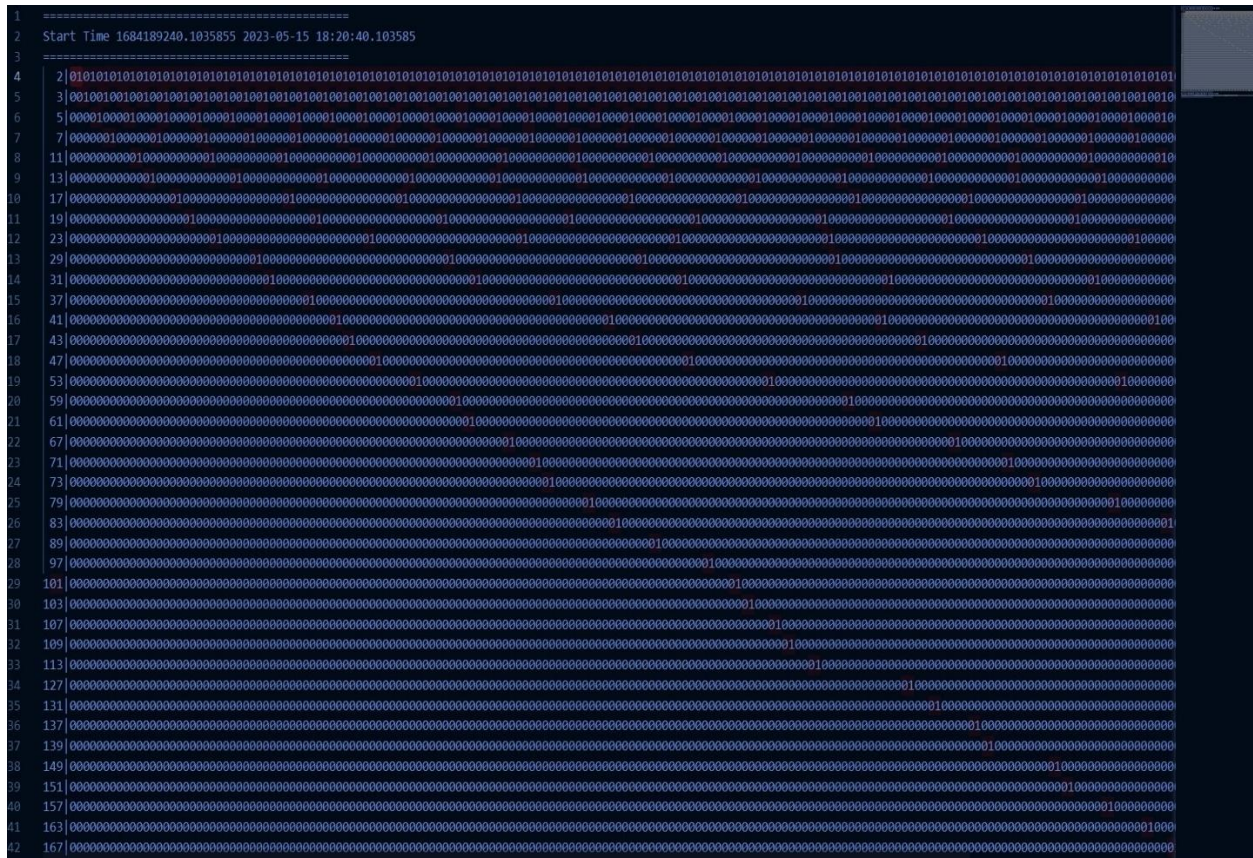


Figure 2. Turing Binary representation for Prime numbers and its multipliers

1.2 Binary Turing operations.

This representation gives us an advantage on getting the value from the index of the position of [1] in the representation. [we do not calculate value in binary system we calculate the value from the index still in base 10 system]

[illegible]

So, this representation is a set of numbers not only one number.

This representation is for SET = {7, 14, 21, 28, 35, 42, 49, 56, 63, 70, 77, 84, 91, 98} as the length we used limited to 100.

If the limit is 200

Our set will be larger until its max number ≤ 200 .

[illegible]

And this is the representation for this Set.

[7, 14, 21, 28, 35, 42, 49, 56, 63, 70, 77, 84, 91, 98, 105, 112, 119, 126, 133, 140, 147, 154, 161, 168, 175, 182, 189, 196]

[7, 14, 21, 28, 35, 42, 49, 56, 63, 70, 77, 84, 91, 98, 105, 112, 119, 126, 133, 140, 147, 154, 161, 168, 175, 182, 189, 196, 203, 210, 217, 224, 231, 238, 245, 252, 259, 266, 273, 280, 287, 294, 301, 308, 315, 322, 329, 336, 343, 350, 357, 364, 371, 378, 385, 392, 399, 406, 413, 420, 427, 434, 441, 448, 455, 462, 469, 476, 483, 490, 497, 504, 511, 518, 525, 532, 539, 546, 553, 560, 567, 574, 581, 588, 595, 602, 609, 616, 623, 630, 637, 644, 651, 658, 665, 672, 679, 686, 693, 700, 707, 714, 721, 728, 735, 742, 749, 756, 763, 770, 777, 784, 791, 798, 805, 812, 819, 826, 833, 840, 847, 854, 861, 868, 875, 882, 889, 896, 903, 910, 917, 924, 931, 938, 945, 952, 959, 966, 973, 980, 987, 994]

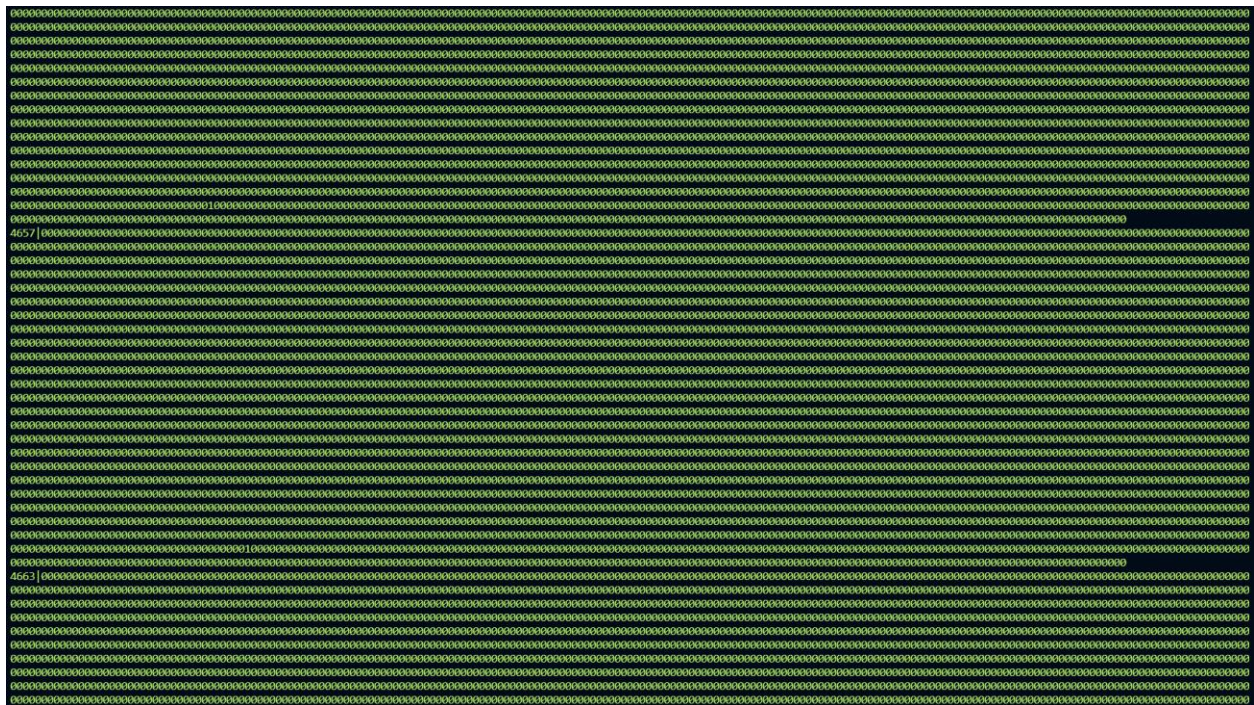


Figure 2. Turing Binary Representations for large numbers have less frequency occurrence for 1 in its representations as the first occurrence will be at position equal this huge number.



Figure 3. Turing Binary Representations for small numbers have higher frequency occurrence for 1 in each given length.

So, filter on 0 will give us nothing.

- For example: - if $ODD = 7$ the add operation results Turing binary number

+

And filter for only Zero they indexes will be set.

4- EVEN +2 = [ODD List]

+

$$=$$

Indexes of Zeros = set of ODD numbers

Now let use see if we filtered on if the value of the add = 1 what index values we are going to get

Filter result on 1:

- +

$$=$$

Indexes of Ones = All natural numbers except multipliers of the number we used in the addition.

[1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21, 23, 24, 25, 26, 27, 28, 29, 30, 31, 32, 34, 35, 36, 37, 38, 39, 40, 41, 42, 43, 45, 46, 47, 48, 49, 50, 51, 52, 53, 54, 56, 57, 58, 59, 60, 61, 62, 63, 64, 65, 67, 68, 69, 70, 71, 72, 73, 74, 75, 76, 78, 79, 80, 81, 82, 83, 84, 85, 86, 87, 89, 90, 91, 92, 93, 94, 95, 96, 97, 98, 100]

- 2- EVEN +1 = [ALL Numbers Except this EVEN Multipliers]

[illegible]

+

[illegible]
$$=$$

111111121111112111111121111111211111112111111121111111211111112111111
11211111112111111121111

[1, 2, 3, 4, 5, 6, 7, 9, 10, 11, 12, 13, 14, 15, 17, 18, 19, 20, 21, 22, 23, 25, 26, 27, 28, 29, 30, 31, 33, 34, 35, 36, 37, 38, 39, 41, 42, 43, 44, 45, 46, 47, 49, 50, 51, 52, 53, 54, 55, 57, 58, 59, 60, 61, 62, 63, 65, 66, 67, 68, 69, 70, 71, 73, 74, 75, 76, 77, 78, 79, 81, 82, 83, 84, 85, 86, 87, 89, 90, 91, 92, 93, 94, 95, 97, 98, 99, 100]

- 3- $ODD + 2 = [ALL\ EVEN\ Numbers + ALL\ this\ ODD\ Multipliers\ except\ it's\ even\ multipliers]$

As 2 has value 1 every two bits first occurrence for this odd number will be in $(\text{odd}/2 - 1/2 + 1)$

And its multipliers will be every step = ODD-1.

```
11|0000000000100000000010000000001000000000100000000010000000001000000000
001000000000001000000000010
```

+

[illegible]
$$=$$

010101010111010101010201010101011101010101020101010101110101010102010101010111010101010201010101011
1010101010201010101011

[2, 4, 6, 8, 10, 11, 12, 14, 16, 18, 20, 24, 26, 28, 30, 32, 33, 34, 36, 38, 40, 42, 46, 48, 50, 52, 54, 55, 56, 58, 60, 62, 64, 68, 70, 72, 74, 76, 77, 78, 80, 82, 84, 86, 90, 92, 94, 96, 98, 99, 100]

List of even numbers including this odd number and its odd multipliers only.

- 4- EVEN +2 = [ALL EVEN Numbers Except this Even number Multipliers]

[illegible]

+

[illegible]
$$=$$
[illegible]

[2, 6, 10, 14, 18, 22, 26, 30, 34, 38, 42, 46, 50, 54, 58, 62, 66, 70, 74, 78, 82, 86, 90, 94, 98]

Filter adds result on 2:

1- $ODD + 1 = [\text{Multipliers of this ODD Number}]$

[illegible]

+

[illegible]

=

[illegible]
$$=$$

[3, 6, 9, 12, 15, 18, 21, 24, 27, 30, 33, 36, 39, 42, 45, 48, 51, 54, 57, 60, 63, 66, 69, 72, 75, 78, 81, 84, 87, 90, 93, 96, 99]

2- EVEN +1 = [Multipliers of this EVEN Number]

```
10|0000000001000000000100000000010000000001000000000100000000010000
000001000000000010000000001
```

+

[illegible]
$$=$$
[illegible]
$$=$$

[10, 20, 30, 40, 50, 60, 70, 80, 90, 100]

3- $ODD + 2 = [EVEN \text{ Multipliers of this } ODD]$

```
13|00000000000010000000000010000000000010000000000100000000001000000000
00010000000000001000000000
```

+

[illegible]
$$=$$

0101010101011101010101010201010101010111010101010102010101010101110101010101
2010101010101011101010101

$$=$$

[26, 52, 78]

4- EVEN +2 = [Multipliers of this EVEN]

[illegible]

+

[illegible]

[illegible]

[8, 16, 24, 32, 40, 48, 56, 64, 72, 80, 88, 96]

We are going to use these operations in getting multipliers of numbers list so we can exclude it in order of getting a final list of prime numbers without doing any complex looping.

(C-1) / 6 and (C - 5) /6

[7, 13, 19, 25, 31, 37, 43,] and list [11, 17, 23, 29, 35, 41, 47, 53,]

```
def binaryI(N , L , padd):
    s=[]
    for v in range(0,L):
        s.append(0)
    for j in range(0,L,N):
        if j+N-1 > L-1 : pass
        else:
            s[j+N-1] = 1
    ''' row header'''
    s.insert(0,str(N).rjust(padd)+'|')
    return s
```



```

def PrimeList(N):
    m = []
    P = [2 , 3, 5 ]
    f = open('Prime_List.lg','w')
    padd = len(str(N))
    for j in range(0,N,6):
        a = 7 + j

        if a%5 == 0 : pass
        elif a not in m:
            Vn11 = OneBinaryI(a,N, padd)
            Vn14 = OneBinaryI(1,N, padd)
            VAn1 = [a+b for a,b in zip(Vn11[1:], Vn14[1:])]
            l1 = [i+1 for i , val in enumerate(VAn1) if val == 2 ]
            m= m + l1
            P.append(a)

        b = 11 + j

        if b > N : break
        if b%5 == 0 : pass
        elif b not in m :
            Vn21 = OneBinaryI(b,N, padd)
            Vn24 = OneBinaryI(1,N, padd)
            VAn2 = [a+b for a,b in zip(Vn21[1:], Vn24[1:])]
            l2 = [i+1 for i , val in enumerate(VAn2) if val == 2 ]
            m= m + l2
            P.append(b)
    f.write(str(P))

```

```

L = 100
lg , hnd = InitiateLogger("RunStats" , "RunStats.lg")
lg.info('=====')
start_T = time.time()
lg.info('Start Time {} {}'.format(start_T,datetime.now()))
lg.info('=====')
PrimeList(L)
lg.info('=====')
End_T = time.time()
lg.info('End Time {} {} '.format(End_T,datetime.now() ))
lg.info(" Duration in sec {} ; Duration in minutes {}".format( End_T - start_T , (End_T - start_T)/60.0 ))
lg.info('=====')

```

Table 1. Turing Binary Execution Times		
First N Prime Numbers List	Execution Time using Turing Binary	
100	0	<pre>===== Start Time 1684218271.7750618 2023-05-16 02:24:31.775061 ===== End Time 1684218271.7750618 2023-05-16 02:24:31.775061 Duration in sec 0.0 ; Duration in minutes 0.0 =====</pre>
[2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97]		
1000	0.03 sec	<pre>===== Start Time 1684218317.0176797 2023-05-16 02:25:17.017679 ===== End Time 1684218317.0550637 2023-05-16 02:25:17.055063 Duration in sec 0.037384033203125 ; Duration in minutes 0.0006230672200520834 =====</pre>
[2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, 101, 103, 107, 109, 113, 127, 131, 137, 139, 149, 151, 157, 163, 167, 173, 179, 181, 191, 193, 197, 199, 211, 223, 227, 229, 233, 239, 241, 251, 257, 263, 269, 271, 277, 281, 283, 293, 307, 311, 313, 317, 331, 337, 347, 349, 353, 359, 367, 373, 379, 383, 389, 397, 401, 409, 419, 421, 431, 433, 439, 443, 449, 457, 461, 463, 467, 479, 487, 491, 499, 503, 509, 521, 523, 541, 547, 557, 563, 569, 571, 577, 587, 593, 599, 601, 607, 613, 617, 619, 631, 641, 643, 647, 653, 659, 661, 673, 677, 683, 691, 701, 709, 719, 727, 733, 739, 743, 751, 757, 761, 769, 773, 787, 797, 809, 811, 821, 823, 827, 829, 839, 853, 857, 859, 863, 877, 881, 883, 887, 907, 911, 919, 929, 937, 941, 947, 953, 967, 971, 977, 983, 991, 997]		
10000	2.9 sec	<pre>===== Start Time 1684218510.2120397 2023-05-16 02:28:30.212039 ===== End Time 1684218513.1682866 2023-05-16 02:28:33.168286 Duration in sec 2.956246852874756 ; Duration in minutes 0.04927078088124593 =====</pre>
[2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, 101, 103, 107, 109, 113, 127, 131, 137, 139, 149, 151, 157, 163, 167, 173, 179, 181, 191, 193, 197, 199, 211, 223, 227, 229, 233, 239, 241, 251, 257, 263, 269, 271, 277, 281, 283, 293, 307, 311, 313, 317, 331, 337, 347, 349, 353, 359, 367, 373, 379, 383, 389, 397, 401, 409, 419, 421, 431, 433, 439, 443, 449, 457, 461, 463, 467, 479, 487, 491, 499, 503, 509, 521, 523, 541, 547, 557, 563, 569, 571, 577, 587, 593, 599, 601, 607, 613, 617, 619, 631, 641, 643, 647, 653, 659, 661,..., 9511, 9521, 9533, 9539, 9547, 9551, 9587, 9601, 9613, 9619, 9623, 9629, 9631, 9643, 9649, 9661, 9677, 9679, 9689, 9697, 9719, 9721, 9733, 9739, 9743, 9749, 9767, 9769, 9781, 9787, 9791, 9803, 9811, 9817, 9829, 9833, 9839, 9851, 9857, 9859, 9871, 9883, 9887, 9901, 9907, 9923, 9929, 9931, 9941, 9949, 9967, 9973]		

100000	4.035 minutes	<pre> ===== Start Time 1684218808.7414348 2023-05-16 02:33:28.741434 ===== ===== End Time 1684219050.8653667 2023-05-16 02:37:30.865366 Duration in sec 242.12393188476562 ; Duration in minutes 4.035398864746094 ===== </pre>
500000	121.26 minutes	<pre> ===== Start Time 1684221315.3232112 2023-05-16 03:15:15.323211 ===== ===== End Time 1684228591.028291 2023-05-16 05:16:31.028291 Duration in sec 7275.70507979393 ; Duration in minutes 121.261 ===== </pre>

There will be future enhancement for this algorithm execution time on classic computers by using accumulative operations and batching to avoid hardware limitations as well.

Conclusion

Classic computers have some limitations in datatypes ranges for numbers that have big number of digits. In this paper we introduced a new algorithm that can be used to in representing set of numbers in a form of Binary Turing number controlled by the number of digits as a parameter. An enhancement for this algorithm will be to use the same algorithm but using batches instead of running the algorithm every time starting from 1 up until the max length parameter. As we showed this algorithm currently shows a promising execution time in handling large number of numbers. Also increasing the number of operations between the Turing Binary representation for each number Sets which will give better execution time and many other applications that can benefits from this algorithm. Also doing operations accumulatively, operate on more than two number at the same time (add more than two binary representations at a time accumulatively). Witch will give us an easy way for factorization without doing any division operations on classical computers.

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