

2)

Our function uses a divide step where the array is divided into halves recursively until each small-array has just one singular element. This is of $O(\log n)$ time. Then there is a merge step, where all the small-arrays are merged. In the event of worst case, it takes $O(n)$ time to merge the small-arrays of size $n/2$ each. Since the levels of recursion are $O(\log n)$, the total complexity of this step is $O(n \log n)$.

3)

Original given array: [8, 42, 25, 3, 3, 2, 27, 3]

- Split into two small-arrays: [8, 42, 25, 3] and [3, 2, 27, 3]
- Split into two small-arrays: [8, 42] and [25, 3], [3, 2] and [27, 3]
- Merge small-arrays: [8, 42] and [3, 25], [2, 3] and [3, 27]
- Merge arrays: [3, 8, 25, 42] and [2, 3, 3, 27]
- Merge arrays: [2, 3, 3, 8, 25, 27, 42, 3]
- Final merge/sort: [2, 3, 3, 3, 8, 25, 27, 42]

As a result, the array is now sorted.

4)

The number of steps is consistent with the complexity analysis. In the worst case scenario, the number of steps is of nature $n \log n$. The merge sort algorithm divides the array into smaller halves and merges them back together, thus the expected time complexity comes to $O(n \log n)$.