

## INDEX

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### Birth Day

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## LAB - 1

## TIC-TAC-TOE GAME

Algorithm:

1) Initialize the Board:

→ Create 3x3 2D array with '-' for empty cells

2) Check if board empty function:

→ Iterate through each row &amp; check if '-' exists, if yes return true

3) Check for winner function:

→ Check each row, column &amp; diagonals for 3 matching marks &amp; return the winning mark if exists.

4) Function to display the board:

5) Start Game function:

→ Initialize the board &amp; randomly select a player &amp; take input.

→ Loop until game ends:

↳ Display the board

↳ take input &amp; validate if it is a number &amp; between 0 &amp; 2 else print invalid.

↳ Update the board with current player's mark

↳ Check if player has won &amp; display

↳ if board is full &amp; no winner print tie

↳ Switch to next player.

6) End Game.

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## Lab - 2

### Vacuum Cleaner World

Pseudo code:

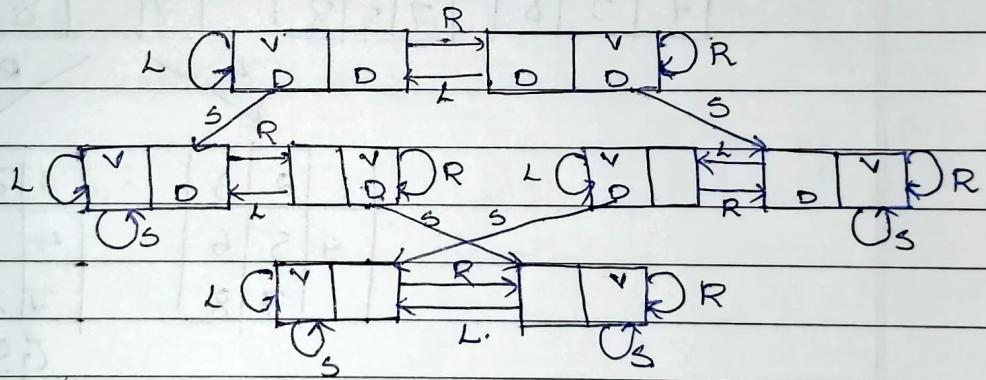
function REFLEX-VACUUM-AGENT ([location, status])  
 returns an action

if status = Dirty then return Suck

else if location = A then return Right

else if location = B then return Left

State-Space diagram:

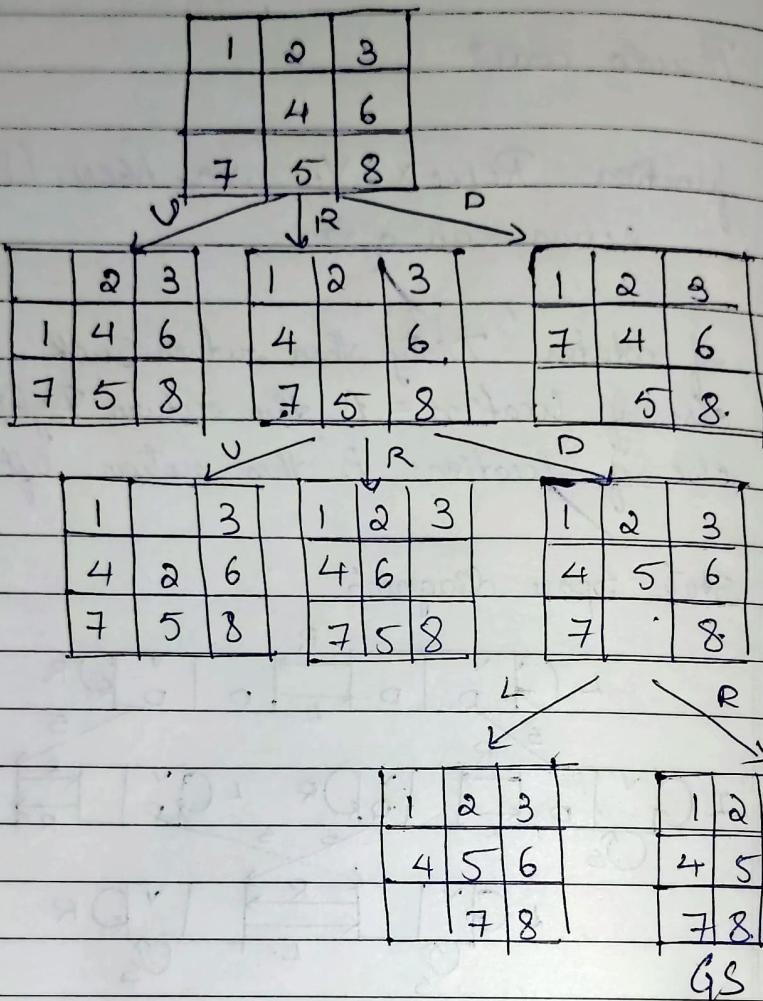


Algorithm.

- Initialize an array with status of both locations as dirty
- In a loop, take input of location & status (0/1) from user & pass it to check function
- If status is dirty print dirty suck & increment cost else print location is clean & update the status in list
- Conti Repeat until user enters status as -1

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## 8-Puzzles Problem Using BFS &amp; DFS



Algorithm: BFS

- Create a queue to store states & add the first state & mark as visited.
- In loop while queue is not empty :
  - Take first element in queue & check if it's target else try all possible moves & add them to queue & mark as visited.
- If queue is exhausted without finding solution print No solution

### DFS Algorithm:

- Create a stack to store current state  $E$ , add the first state  $E$  mark as visited.
- In loop while stack is not empty:
  - Pop the top element  $E$  & check w/ its target else mark it as visited & push it into the stack all possible states from that state.
- If stack is exhausted print No solution

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## Lab - 3

For 8-puzzle problem using A\* implementation  
to calculate  $f(n)$  using:

a)  $g(n)$  = depth of node

$h(n)$  = heuristic value

a)  $\hookrightarrow$  No. of misplaced tiles

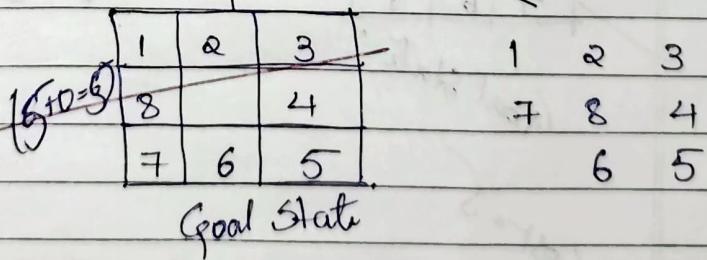
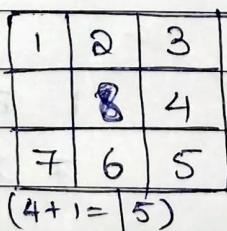
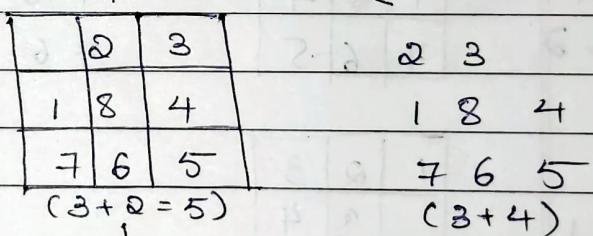
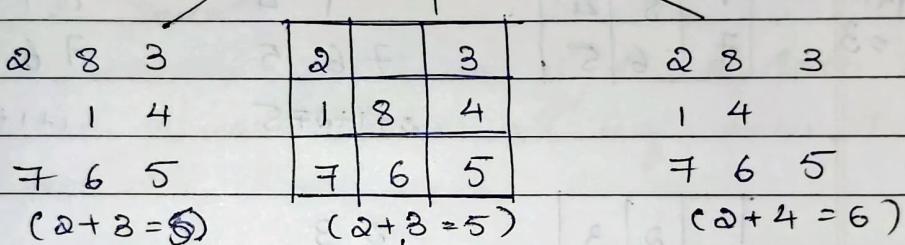
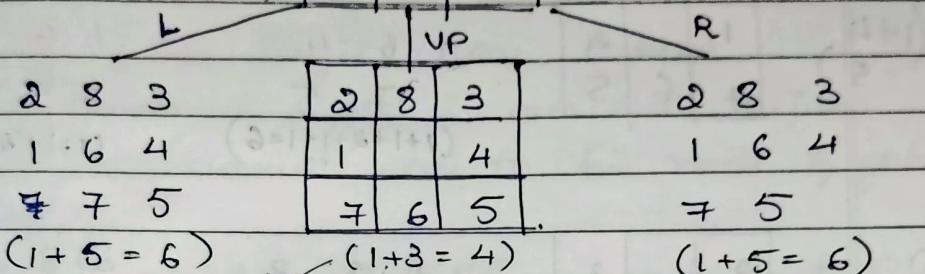
b)  $\hookrightarrow$  Manhattan distance

$$f(n) = g(n) + h(n)$$

- 1) Step 1: Place the starting node in OPEN list
- 2) Check if open list is empty or not, if empty return failure.
- 3) Select the node with smallest  $f(n)$  value  
if n is goal return success
- 4) Generate all null successors & put current node in closed list.
- 5) If it is not in open or closed compute  $f(n)$
- 6) Return to step 2.

a) Give state space diagram for:

Initial state			Goal State		
$f(n) = g(n) + h(n)$	0	3	3	1	2
$= 0 + 4$	1	6	4	8	4
$= 4$	7	5		7	6



b)

2	8	3
1	6	4
7		5

$$(1+4 = 5)$$

2	8	3
1		4
7	6	5

2, 8, 3

1, 6, 4

7, 5

$$(1+1+2+1+1=6)$$

2, 8, 3

1, 6, 4

7, 5

$$(1+1+2+1+1=6)$$

$$\begin{matrix} h(n) \\ 1+1+1 \\ =3 \end{matrix}$$

2		3
1	8	4
7	6	5

2, 8, 3

1, 4

7, 6, 5

$$2+1+2=5$$

2, 8, 3

1, 4

7, 6, 5

$$1+1+2+1=5$$

$$\begin{matrix} 1+1 \\ =2 \end{matrix}$$

	2	3
1	8	4
7	6	5

2, 3

1, 8, 4

7, 6, 5

$$1+1+1+1=4$$

1

1	2	3
	8	4
7	6	5

0

1	2	3
8		4
7	6	5

1, 2, 3

7, 8, 4

6, 5

$$\begin{matrix} 1+1 \\ 2 \end{matrix}$$

Goal state:

1	2	3
8		4
7	6	5

Cost = 5

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## LAB - 4

## Iterative Deepening Search Algorithm

- 1) For each child of the current node.  
    If it is target node, return
- 2) If <sup>current</sup> maximum depth is reached, return
- 3) Set the current node to this node & go back to 1
- 4) After exploring all children, go to next child of the parent
- 5) Increase the maximum depth. ~~by goto 1~~
- 6) If we have reached all leaf nodes, the goal node doesn't exist.

depth = 0

2	8	3
1	6	4
7		5

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depth:

$h^{(n)} = 5$

2	8	3
1	4	
7	6	5

0	8	3
1	6	4
7	5	6

2	8	3
1	6	4
7	5	

$h^{(n)} = 6$

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depth

2	3	
1	8	4
7	6	5

0	8	3
1	4	
7	6	5

2	8	3
1	4	
7	6	5

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## Lab - 4

- \* Implement Hill Climbing search algorithm to solve N-Queens Problem.

```
function HILL-CLIMBING(problem) returns a state  
    that is local maximum  
    current ← make-node(problem, INITIAL-STATE)  
    loop do  
        neighbour ← a highest-valued successor of  
        current.  
        if neighbour.VALUE < current.VALUE then  
            return current.STATE  
        current ← neighbour.
```

State: 4 queens on the board - one queen per column

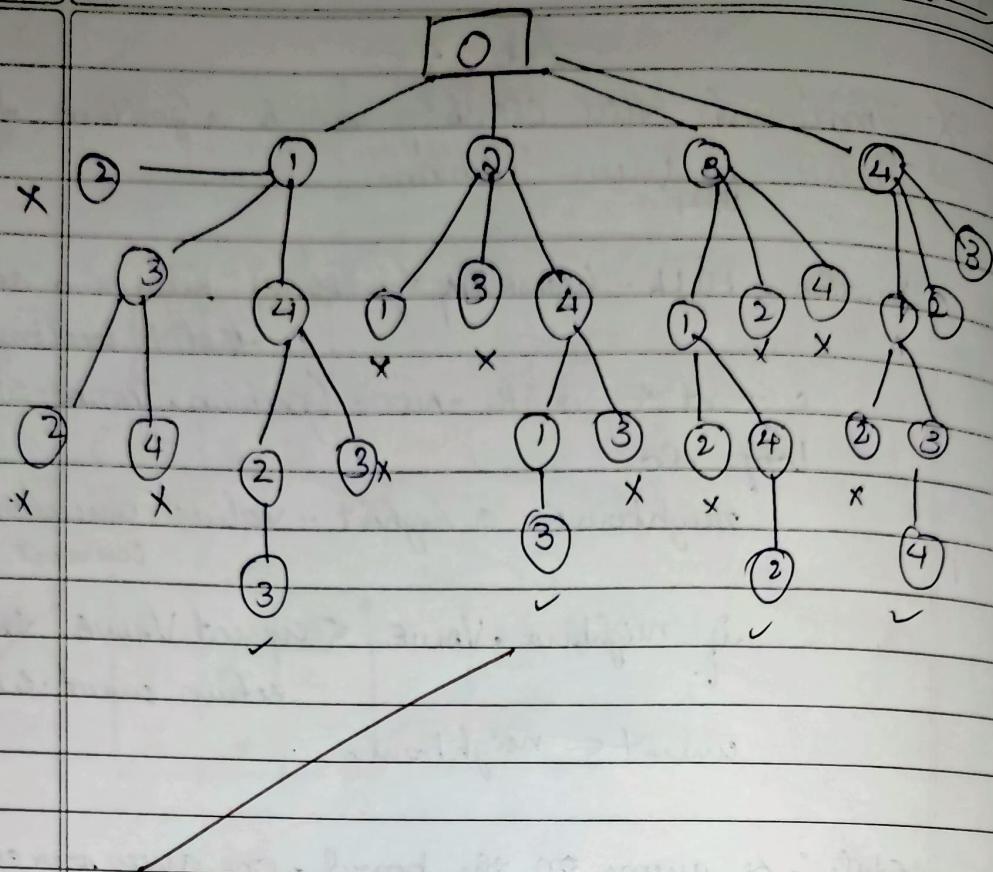
- variables:  $x_0, x_1, x_2, x_3$  where  $x_i$  is the row position of the queen in column  $i$ . Assume that there is one queen per column.
- Domain for each variable  $x_i \in \{0, 1, 2, 3\}$

Initial state: a random state

Goal state: 4 queens on the board - No pair of queens are attacking each other.

Neighbour relation: Swap row positions of 2 queens

Cost function: The number of pairs of queens attacking each other directly or indirectly.



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## Lab - 5

Write a program to implement Simulated Annealing

function SIMULATED ANNEALING (problem, schedule)  
 current  $\leftarrow$  MAKE-NODE(problem, INITIAL-STATE)  
 for  $t = 1$  to  $\infty$  do  
      $T \leftarrow$  schedule( $t$ )  
     if  $T = 0$  then return current  
     next  $\leftarrow$  a randomly selected Successor of current  
      $\Delta E \leftarrow$  next.Value - current.Value  
     if  $\Delta E > 0$  then current  $\leftarrow$  next  
     else current  $\leftarrow$  next only with probability  $e^{\Delta E/T}$

- 1) Start at a random point  $x$ .
  - 2) Choose a new point  $x_j$ , on a neighbourhood  $N(x)$ .
  - 3) Decide whether or not to move to new point  $x_j$
- The decision will be made based on the probability

function  $P(x, x_j, T)$

$$P(x, x_j, T) = \begin{cases} 1 \\ e^{\frac{f(x_j) - f(x)}{T}} \end{cases}$$

## Fa 8 Queens Problem :

- 1) Initialize : Start with arrangement of queens  
Set initial temp & cooling parameters
- 2) Evaluate fitness :  
Calculate no. of conflicts between queens.  
Lower conflicts = lesser fitness.
- 3) Generate Neighbour :  
Create a neighbour by swapping 2 queens.  
calculate its fitness
- 4) Accept or Reject Neighbour :
- 5) Gradually Reduce the temperature & repeat until  
the temp is very low.
- 6) Return the best arrangement found.

Output : Best State found [ 4 7 5 0 6 3 7 2 ]

Number of conflicts = 2.

~~Q8  
7 10 2  
9 1 6 3  
5 4~~

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## LAB - 6

Create a knowledge base <sup>using</sup> propositional logic & show that the given query entails the knowledge base or not

Algorithm:

- Generate all possible models: For  $n$  symbols there are  $2^n$  combinations
- For each model : evaluate each statement in  $\text{KB}$ , if all statements are true proceed to next test otherwise skip to next step.
- Check the Query : If  $\text{KB}$  is true under the model, check if query is true.
- If query is true in all cases when  $\text{KB}$  is true then  $\text{KB}$  entails the Query

Ex:-  $\text{KB} = (\text{AVC}) \wedge (\text{B} \vee \neg \text{C})$   $\alpha = \text{AVB}$

A	B	C	AVC	$\text{B} \vee \neg \text{C}$	$\text{KB}$	$\alpha$
F	F	F	F	T	F	F
F	F	T	T	F	F	F
F	T	F	F	T	F	T
F	T	T	T	T	T	T
T	F	F	T	T	T	T
T	F	T	T	F	F	T
T	T	F	T	T	T	T
T	T	T	T	T	T	T

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## LAB - 6

Create a knowledge based propositional logic & show that the given query entails knowledge base or not

Algorithm:

- Generate all possible models: For  $n$  symbols there are  $2^n$  combinations
- For each model : evaluate each statement in KB, if all statements are true proceed to next test otherwise skip to next step.
- Check the Query : If KB is true under the model, check if query is true.
- If query is true in all cases when KB is true then KB entails the Query

$$\text{Ex:- } KB = (AVC) \wedge (B \vee \neg C) \quad \alpha = AVB$$

A	B	C	AVC	$B \vee \neg C$	KB	$\alpha$
F	F	F	F	T	F	F
F	F	T	T	F	F	F
F	T	F	F	T	F	
F	T	T	T	T	T	T
T	F	F	T	T	T	T
T	F	T	T	F	F	T
T	T	F	T	T	T	T
T	T	T	T	T	T	T

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## Lab - 7

Implementation unification in first order logic

Algorithm : Unify ( $\Psi_1, \Psi_2$ )

- Step 1: If  $\Psi_1$  and  $\Psi_2$  is a variable or const:
- If  $\Psi_1$  or  $\Psi_2$  are identical, then return NIL
  - Else if  $\Psi_1$  is a variable:
    - then if  $\Psi_1$  occurs in  $\Psi_2$  then return Failure.
    - else return  $\{\Psi_1 / \Psi_2\}$
  - Elseif  $\Psi_2$  is a variable:
    - a. If  $\Psi_2$  occurs in  $\Psi_1$ , then return FAILURE
    - else return  $\{\Psi_1 / \Psi_2\}$
  - Else return FAILURE

Step 2: If the initial predicate symbol in  $\Psi_1$  and  $\Psi_2$  are not same, then return FAILURE

Step 3: If  $\Psi_1$  and  $\Psi_2$  have a different no of arguments, then return FAILURE

Step 4: Set substitution. Set (SUBST) to NIL

- Step 5: For  $i = 1$  to the no of elements in  $\Psi_1$
- Call Unify function with the  $i^{th}$  element of  $\Psi_1$  and  $i^{th}$  element of  $\Psi_2$  & put the result into S.
  - If  $S = \text{failure}$ . then return Failure
  - If  $S \neq \text{NIL}$ . then do,
    - Apply S to the remainder of both L1, L2

b.  $\text{SUBST} = \text{APPEND}(S, \text{SUBST})$

Step 6: Return  $\text{SUBST}$ .

$$\textcircled{Q} \quad \psi_1 = P(f(a), g(y))$$

$$\psi_2 = P(x, x)$$

→ predicate symbols are same

→ No of arguments are same.

→ For unification : we'll equate:

$$f(a) = x \quad g(y) = x$$

$$\Rightarrow f(a) = g(y)$$

→ but  $f$  &  $g$  are two different functions  
hence they cannot be unified.

$$\textcircled{Q} \quad \psi_1 = P(b, x, f(g(z)))$$

$$\psi_2 = P(z, f(c), f(c))$$

$$\theta_0 = b/z \quad \theta_1 = f(y)/x, \quad \theta_2 = g(c)/y$$

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## Lab - 8

Create a knowledge base consisting of first order logic statements & prove the given query using forward reasoning.

### 1) Input:

→ (condition, result)

    ↳ set of facts that must be true

    ↳ fact inferred if condition true

→ A set of initial facts known to be true

→ A query to determine if it can be inferred.

→ applied rules = True.

### 2) while applied\_rules is True:

    set applied\_rules = False

    For each (condition, result) in rules

        check if condition ⊂ facts

        if results ⊂ facts

            add result to facts

            set applied\_rules = True.

            if query ⊂ facts return True

        if no new rules applied\_rules = False. exit loop

### 3) Output:

    if query ⊂ facts return True (inferred)

    else return False (query not inferred).

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## Lab - 9

Convert a given first order logic statement into resolution.

## 1) Convert to Clause Form (CNF):

- Eliminate implication ( $A \rightarrow B$  becomes  $\neg A \vee B$ )
- Move negations inward using De Morgan's Law.
- Standardize variables
- Drop universal quantifiers
- Convert to a conjunction of disjunctions (CNF form)

## 2) Negate the statement to be Proved:

- Negate the query & add it to knowledge base.

## 3) Apply Unification

- Identify & resolve two clauses that contain complementary literals.

## 4) Derive new clauses:

- Continue resolving clauses until either:
  - A contradiction is found
  - No further resolution is possible  
(query cannot be proved)

## Lab - 10

## Implement Alpha-beta pruning

## 1) Input:

- A game tree with terminal node values
- Depth & whether maximizer or minimizer's turn

## 2) Steps:

- Initialize  $\alpha = -\infty$  &  $\beta = +\infty$
- Start at root & recursively evaluate child nodes
- For each node:
  - If ~~not~~ maximizer:
    - Initialize to  $-\infty$
    - update  $\alpha = \max$  value
    - if  $\alpha \geq \beta \rightarrow$  stop
  - If minimizer:
    - Initialize to  $+\infty$
    - update  $\beta = \min$  value
    - if  $\alpha \geq \beta \rightarrow$  stop
- Return optimal value of root