The logic of computers

starting in 5:00

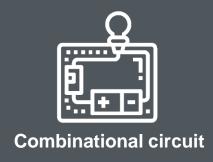
"To be, or not to be ..."

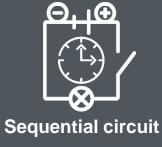
- Shakespeare



Dr. Goran Soldar Dr. Khuong An Nguyen







Boolean algebra.

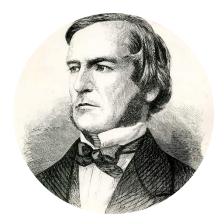
2 Combinational circuit.

3 Sequential circuit.



Boolean algebra

- Algebra on 0 and 1.
- 2 Has its root in philosophy.



1. What we start

Axioms: assuming basic objects and operations are true.

2. What we derive

Laws and theorems: manipulating Boolean expressions

3. What we build

Circuit: deriving complex digital design.

Further reading

"The Mathematical analysis of Logic". George Boole, 1847

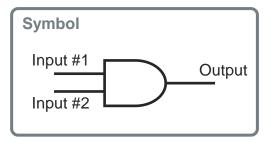


Ordinary vs Boolean algebra

	Ordinary algebra	Boolean algebra
Inputs	1	false 0
Operators		AND OR NOT XOR



AND operation



Boolean Expression

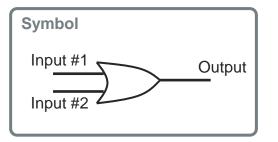
Given Input #1 = A; Input #2 = B; Output = Q

 $Q = A \bullet B$

Input #1	Input #2	Output
True (1)	True (1)	True (1)
True (1)	False (0)	False (0)
False (0)	True (1)	False (0)
False (0)	False (0)	False (0)



OR operation



Boolean Expression

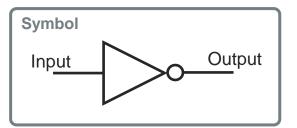
Given Input #1 = A; Input #2 = B; Output = Q

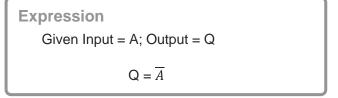
Q = A + B

Output
True (1)
True (1)
True (1)
False (0)



NOT operation

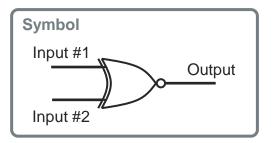




Truth table				
Input Output				
False (0)				
True (1)				



XOR (Exclusive-OR) operation



Boolean Expression

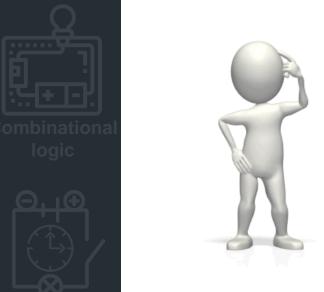
Given Input #1 = A; Input #2 = B; Output = Q

$$Q = A \oplus B = \overline{A} B + A \overline{B}$$

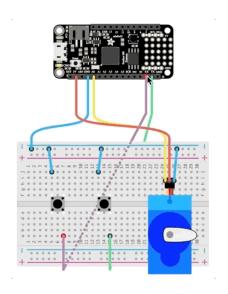
	Truth table	
Input #1	Input #2	Output
True (1)	True (1)	False (0)
True (1)	False (0)	True (1)
False (0)	True (1)	True (1)
False (0)	False (0)	False (0)



Designing a circuit



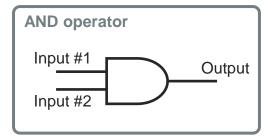






Truth table

- Lists all possible inputs and outputs.
- Two circuits are identical, if they share the same truth table.



Truth table				
Input #1	Input #2	Output		
True (1)	True (1)	True (1)		
True (1)	False (0)	False (0)		
False (0)	True (1)	False (0)		
False (0)	False (0)	False (0)		



Step 1: Writing the truth table

Specification:

A circuit that accepts 3 inputs, and returns True, only when one input is False.

Writing the truth table:

A circuit with n inputs will always have 2^n rows (possibilities).

Truth table

Input A	Input B	Input C	Output
True (1)	True (1)	True (1)	False (0)
True (1)	True (1)	False (0)	True (1)
True (1)	False (0)	True (1)	True (1)
True (1)	False (0)	False (0)	False (0)
False (0)	True (1)	True (1)	True (1)
False (0)	True (1)	False (0)	False (0)
False (0)	False (0)	True (1)	False (0)
False (0)	False (0)	False (0)	False (0)



Step 2: Deriving Boolean expressions

Find all rows with output True (1).

For each row:

if the input X is 1, write X if the input X is 0, write \overline{X}

	Truth table					
Input A	Input B	Input C	Output	Expression		
True (1)	True (1)	True (1)	False (0)			
True (1)	True (1)	False (0)	(True (1))	A.B. \overline{C}		
True (1)	False (0)	True (1)	(True (1)	$A. \overline{B}. C$		
True (1)	False (0)	False (0)	False (0)			
False (0)	True (1)	True (1)	True (1)	Ā. B. C		
False (0)	True (1)	False (0)	False (0)			
False (0)	False (0)	True (1)	False (0)			
False (0)	False (0)	False (0)	False (0)			



Step 3: Sum of the products

OR (+) all the product of the rows together.

A. B. \overline{C} + A. \overline{B} . C + \overline{A} . B. C

Truth table					
Input A	Input B	Input C	Output	Expression	
True (1)	True (1)	True (1)	False (0)		
True (1)	True (1)	False (0)	(True (1))	A.B. \overline{C}	
True (1)	False (0)	True (1)	(True (1)	$A. \overline{B}. C$	
True (1)	False (0)	False (0)	False (0)		
False (0)	True (1)	True (1)	(True (1)	Ā. B. C	
False (0)	True (1)	False (0)	False (0)		
False (0)	False (0)	True (1)	False (0)		
False (0)	False (0)	False (0)	False (0)		



Expression simplification

Axioms

○ Closure, given $a, b \in B$:

$$a+b \in B$$

 $a \bullet b \in B$

Commutative, given a, b ∈ B

$$a+b=b+a$$

 $a \bullet b = b \bullet a$

○ Identity, given $0, 1 \in B$:

$$a + 0 = a$$

 $a \cdot 1 = a$

o Distributive:

$$a + (b \bullet c) = (a + b) \bullet (a + c)$$

 $a \bullet (b + c) = a \bullet b + a \bullet c$

o Complement:

$$a + \overline{a} = 1$$

 $a \bullet \overline{a} = 0$

Laws

Operation with 0 and 1 :

$$X + 0 = X$$
 $X \bullet 1 = X$
 $X + 1 = 1$ $X \bullet 0 = 0$

o Idempotent:

$$X + X = X$$
 $X \bullet X = X$

o Involution:

$$\overline{(\bar{X})} = X$$

Complementarity:

$$X + \overline{X} = 1$$
 $X \bullet \overline{X} = 0$

o Commutative:

$$X + Y = Y + X$$
 $X \bullet Y = Y \bullet X$

Associative :

$$(X + Y) + Z = X + (Y + Z)$$

 $(X \bullet Y) \bullet Z = X \bullet (Y \bullet Z)$

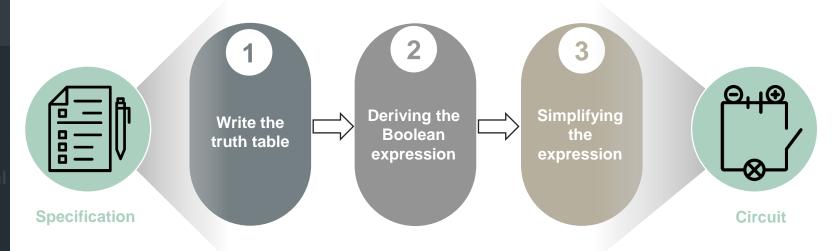
o Distributive:

$$X \bullet (Y + Z) = (X \bullet Y) + (X \bullet Z)$$

 $X + (Y \bullet Z) = (X + Y) \bullet (X + Z)$



Circuit design steps





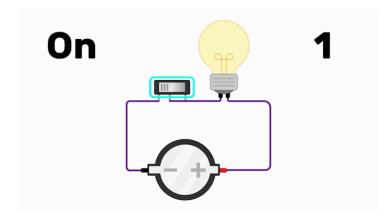






Signal delay

- We often assume the signal output happens immediately.
- CPU speed limitation: cannot perform more calculations per second than its capability.

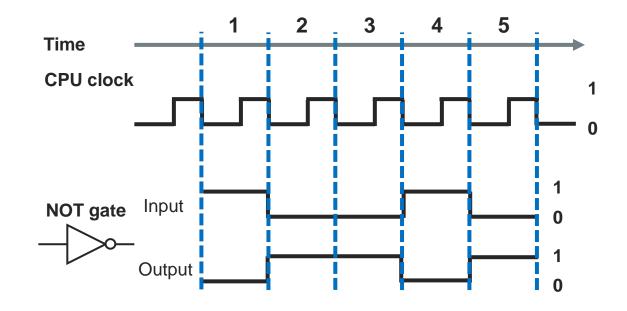








Electronic signal (in theory)



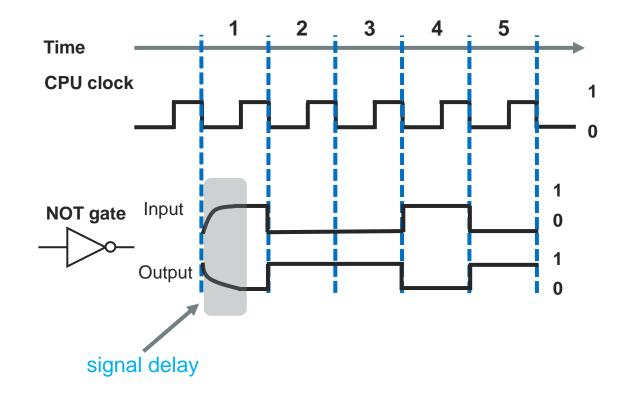








Electronic signal (in reality)



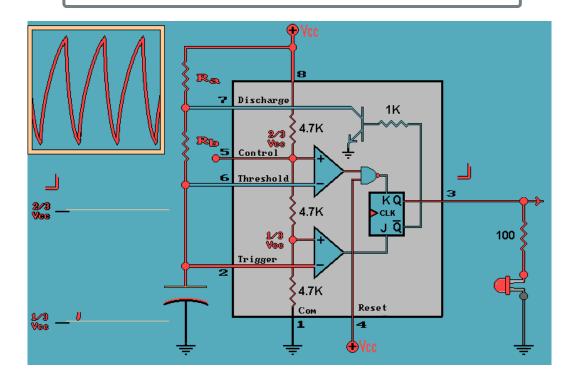




ΘΗΦ (2) Sequential logic

Processing delay

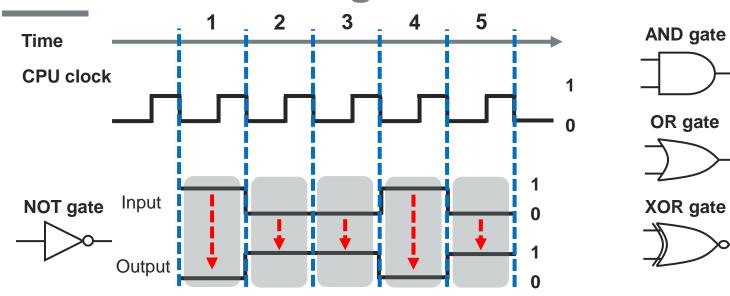
Complex circuit takes time to process the inputs.



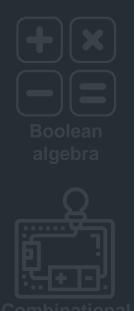
Boolean

Combinational logic

Combinational logic

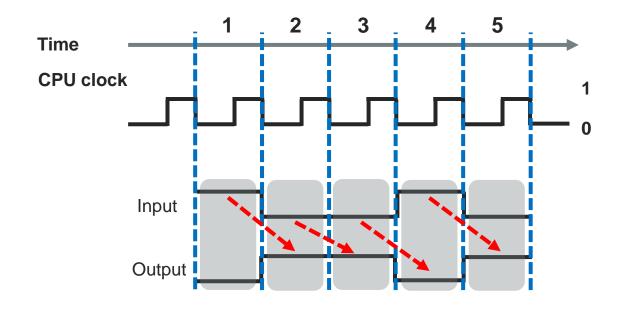


- Takes some inputs, and produces some outputs.
- All processing happens within one time unit.
- output[t] = function(input[t]).



Sequential logic

Sequential logic



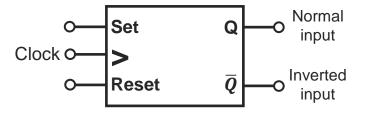
Output depends on the input, and the previous states output[t] = function(input[t - 1]).

+ X Boolean

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Sequential logic

S-R Flip-Flop

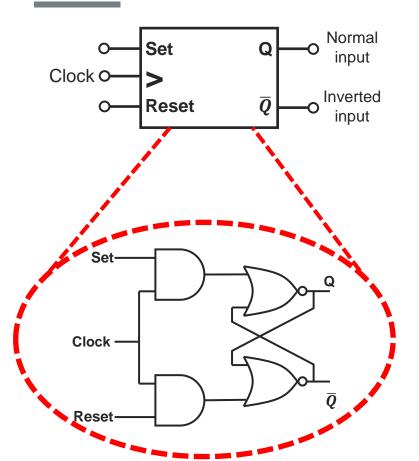


Set state	Reset state	
Q = 1	Q = 0	
$\overline{O} = 0$	$\overline{Q} = 1$	

- 1-bit storage device.
- Output can be set to store 0 or 1, depending on the inputs.
- Flip-Flop retains its outputs, even without input signals.



S-R Flip-Flop

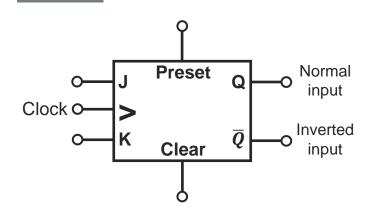


Truth table				
Set	Reset	Q	\overline{Q}	
0	0	no change		
0	1	1 0		
1	0	0	1	
1	1	invalid		



Boolean algebra

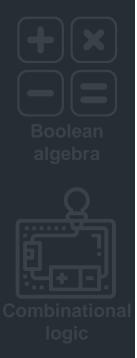
J-K Flip-Flop



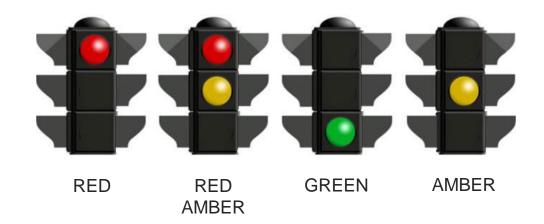
Truth table				
J	K	Q	\overline{Q}	
0	0	no change		
0	1	1 0		
1	0	0	1	
1	1	toggle		

- An improvement over S-R Flip-Flop.
 - No invalid state.
- When J = K = 1, the output toggles between 0 and 1.





Traffic light example







Questions, feedback



Cockcroft building C519 (Khuong) C537 (Goran)

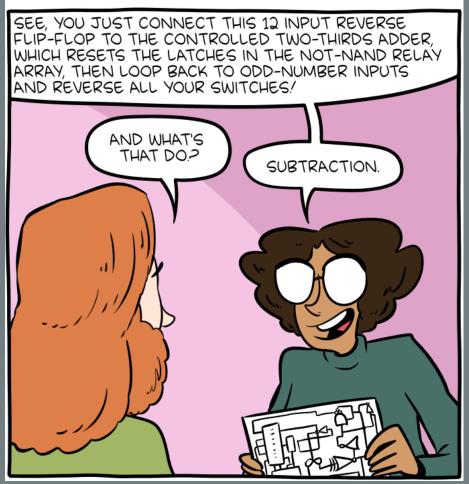


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THIS IS WHAT LEARNING LOGIC GATES FEELS LIKE



(SMBC comic)