Togannent-3 Define a macro, macro call and macro expansion Mocro - It Starts with moero pseudo - 6 P It indicates the begining of a micro definition An It is terminated with the MENS pseudo-of name as an approximate defined the use of macro Macro callname as an operator mnemenic in an arrembly program is equivalent to the corresponding more attention sequence are written with the mocro name. This is called mocrocall motro empansion - whenever there is a matro call the morro processor Substitutes the macro definition in the place of the mit morro call This is latted morrio expansion. The macro definition itself doesnot appear in the expanded source code. (2) what are the feature of a morro a Explain conditional mano expansion a Features of moero one se se se son se se AN 1 conditional macro expansion (9) (4) morro calls with morros Conditional macro expansion Classicals The sequence of macro expansion can be recorded based en some conditions mano expension Itallow condition salution of the mochine Instruction that appear in expansion of amorro call. for his we cue morno pseudo-ops AIF and A 570 AIF It is a condition at branch pseudo-cps. The format of AIF statement is as pollow AIF Lexpression > Lequence symbol)

ABGD W an unconditional tranch pseudo -op The format of A 510 Statement was follow A 670 & sequencing symbols example: consider the following program 1 Data 1 20001

2 Data 2 3 Data3

1 Data 1 2 Data 2

A 11 Data 1 doop 3 F5 DC Data 1 F 10 Data 2 DC F 15 DC Data 3

In this example the operands clabels and the number of instruction generated change in each Sequence. This program could be written as follow MADE & SHOOM

2 F140

SARGIO MOCRO VARY GOUNT Y ARGIS, WARGZ INARGY3 VAROTO 4 1, VARGI (& COUNT EOI). FINT Test if & coon AIF 2 VARUI 2 AIF (N COUNT EOZ.) FINI Test 1 & COU 7 3, &ARG 3 A 1900% FINI MEND 00003 20001 VARY 3 IDATALI DATA2 IDATA3 expanded Zence doop 2 VARY 2 DATA 3 DATA LOOPL A 1 Doctq 1 2 Data2 3 Datas F10 Loop2 A 1 Dyung A 2 DATA2 DATA1 DC F'5' Doct A2 DATAS d00p3 4 1,DATA1

8) El Statements, A 676 AIF Stadements It is a conditional branch pseudo -op The format of AIT statements has follow ATTAX expression > > Sequencing symbols wa related expression involving colore expression porameters, and sher attributes Abju Statements It is an unconditional branch pseudo -op. The format of 4670 Statements is as follows 1 670 & sequencing symbols It cencon-actionally transfers control to the statement containing sequencing symbol inits laber field. A GO statement doesn't appear in me expanded source code Explain & pay mooro -processor algorithm. two pars from - processor algorithm while implementing The following arrumption one made. (3) Mocro processor is functionally independent of the anembles Output from the moero processer will be fed into the the amombles Morro definition or call with in momo definition over not allowed because of its complication, Like of an animble, moore processed also scans and processes ene of a text. like of an amembler, micro processor also scans and processes dine of atext The lines of morro language are not so closely interrelated as in anembles i e mano definition refer to nothing outside themselves and morro call reperonly to moero definition

A moore cay substitutes 10x4 gor perameter not values example: Consider the following statements INCR YEOV TO The morro call INCRY is of ollowed by o designing X The macro processor substitues y for the macro definition It will not anigh a value 10 to Y because the EOU I tatement what processed by the moore processed It consider it as a complete line ofterd Assignment - 4 Régine Loades 9 conat are its function The leader is program which a crept. The object program Compared shere program for enecution by the Computer and initiates the enecution Function of loader Function of loader Allocation: allocate spoce in memory for the programs Linking Resolve symbolic reperence between object ducks (3) Relocation Adjust all address dependents locations such as address constants to corresponds to the allocated space Loading: physically place the machine instruction and data into

Name different types 10 ading scheme ? different types of loading schemes are 1) compile and go leaders Goneral louder scheme Relocating loaders Direct linking loaders Dynamie Joading Dynamic Linking. explain dynamic loading and dynamic binder Dynamic loading-If the total amount of core required by all there subsounting exceed the amount available, as is common with large program W5 mall compute 21 1 there is trouble There one several hordwere techniques 150ch as paging and seg mentation , short attempt to solve hus problem! In this section we will present conventional dynamic loading Schemes bared agon the use of a binder sprior to loading. Dynamic Binder and buy of 4 binder is a program that porjoins the same function as direct linexing loader in binding subscertine together but gather than process the relocated and linked text directly into memory, its output the text as a file or cord Luck. The output file is in a garmot ready to be loaded and is typically called a lead module. The loader merecy has to physically load the module into core The binder eventially performs the functions of allocation, relocation and linking: the module loader morely perform the function of landing.

There are two major classes of bindens Core image builder dinkage editos Explain Four types of cards used in direct linking Loader An External Symbol Dictionary cords. The ESD Cord Contains the information necessary to been the external symbol dicitorary or symbol Deyond the subsoutine level Degment Definition (SD) - name on START docar Definition (LD) - specified on Entry card C external regenence (ER) - specified on EXTERN card 2) TXT cards The txt Cords contain block of data and the relative address at which I the data is to be placed once the loader is divided where to lead the program relative odds the program loader addres (PLA) to the relative address and moves the data into the resulting The data on the +XT cord may be instruction, nonrelocated data or initial values of address constants. The PLD conds The RID Cords contain the following information. The location and length of each address cenetant that needs to be changed for relocation or linking

the may so bol by work on the address constant should be modified (added or subtracted) O The operation to be performed codd or subtract END cords. The end cord specifies the end of the object If the arrembles END cord has a symbol in the operand field (au subroutines) of enecution point for the entire program This on address is recorded on the END card. This is final cord required to specify the end of a cellection The 360 loadons usually either a Loader terminate (101) or END of file (FOF) cord 6) write specification of databases used in part 1 and part 2 of The database use in pass 1 and pass 2 of Direct linking loader ( pars I data pames 1 Input object ducks. A parameter the intial program Load Addrew (IPLA) supplied by the programme or the operating system that specifies the address to lead the first segment. A program Load Address (PLA) Counter, used to Keep track of each segments assigned location A table the Golobal Extuenal symbol table (GEST) that is used to store each external symbol and its corresponding awigned care address A copy of the input to be used later by pars 2

